

# CS3EA01 Artificial Intelligence

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# Syllabus

## **Unit I:**

Introduction to artificial intelligence, various types of production systems, Characteristics of production systems, Study and comparison of breadth first search and depth first search techniques.

## **Unit II:**

Optimization Problems: Hill-climbing search Simulated annealing like hill Climbing, Best first Search. A\* algorithm, AO\* algorithms etc and various types of control strategies, Heuristic Functions, Constraint Satisfaction Problem.

## **Unit III:**

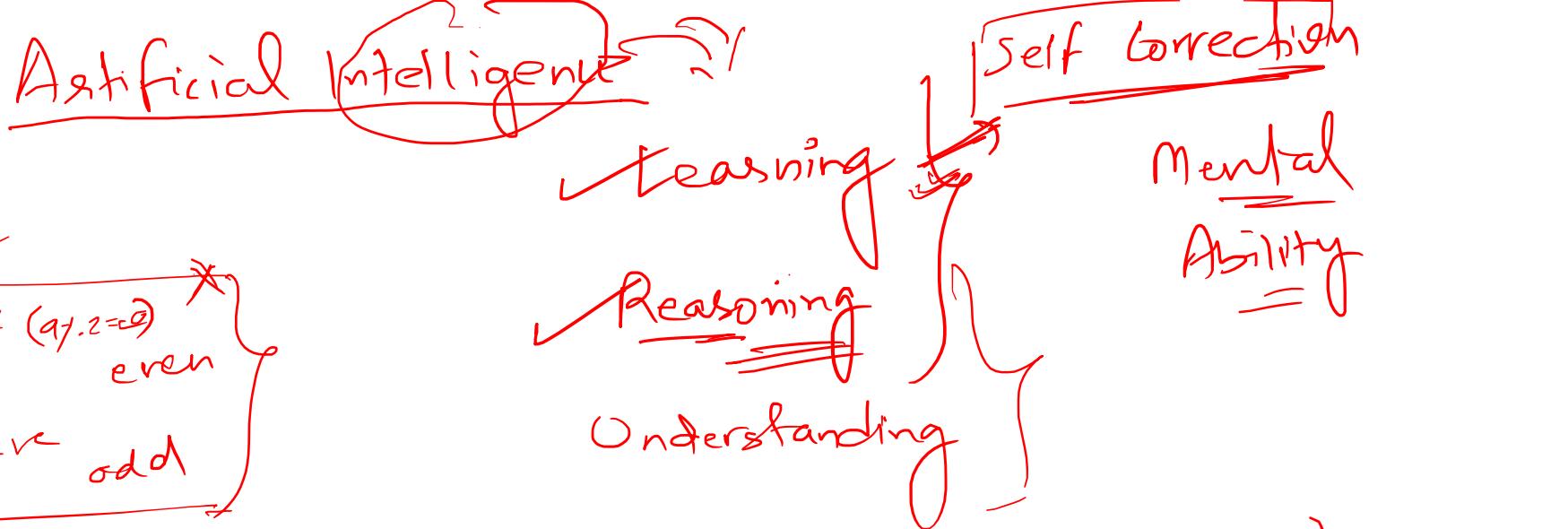
Knowledge Representation, structures, Predicate Logic, Resolution, Refutation, Deduction, Theorem proving, Inferencing, Semantic networks, Scripts, Schemas, Frames, Conceptual dependency.

## **Unit IV:**

Uncertain Knowledge and Reasoning, forward and backward reasoning, monotonic and non-monotonic reasoning, Probabilistic reasoning, Baye's theorem, Decision Tree, Understanding, Common sense, Planning

## **Unit V:**

Game playing techniques like minimax procedure, alpha-beta cut-offs etc, Study of the block world problem in robotics.



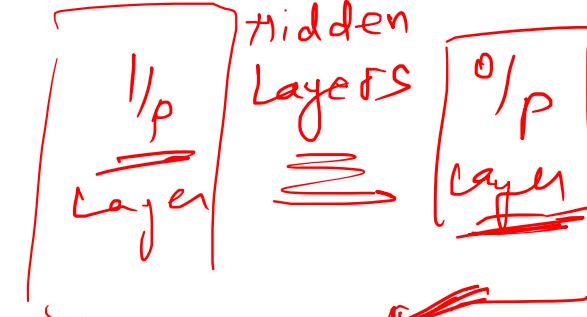
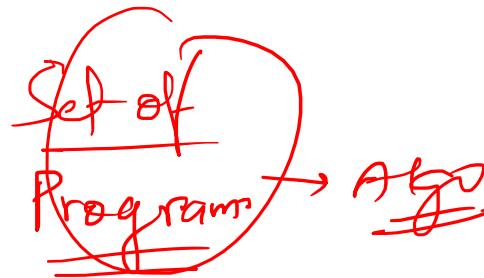
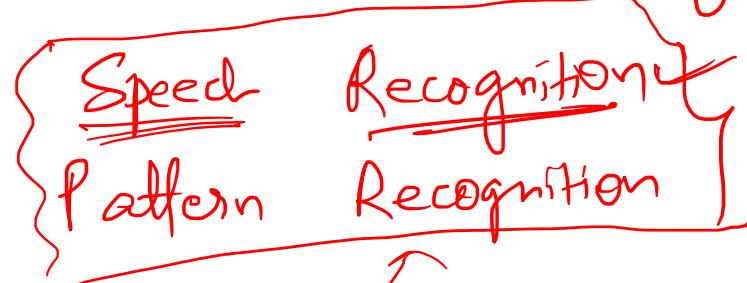
Def Intelligence exhibited by the machines (Comp / Algo / -)  
System / SW

Simulation of intelligent behaviours.

✓ Sanfoundry  
Quiz 1, 2, 3  
Asg - 1, 2

# Applications Components

✓ Machine learning



Perception

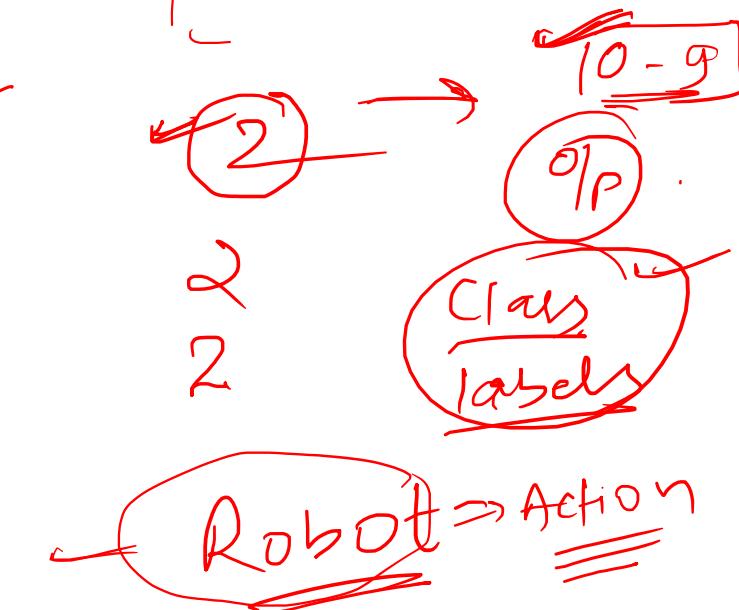
touching  
smelling  
listening

Unsupervised Learning



Natural Language Processing

Reinforcement



Pattern Recognition  $\rightarrow$  Computer Vision

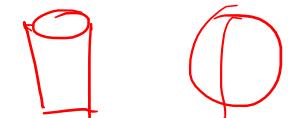
2D & 3D

Regular shapes  
Irregular shapes

Digital Images Processing

Camera Sensor

Agent



2013  
Deep Learning      Complex

facebook

~~tag~~  
~~hidden~~  
~~layer~~

Game Playing | Problem Solving  $\Rightarrow$  Searching  
Theorem Proving

DFS    BFS

Decision making  
Behavior

Expert System

Medical Consultants  
(fever)

✓ Chatbots

CADET

Applications

finance

Education / Design

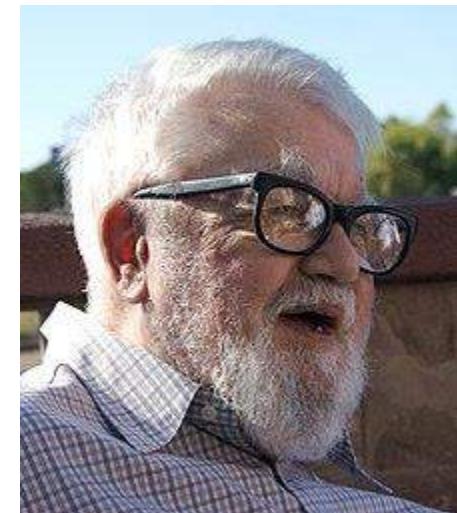
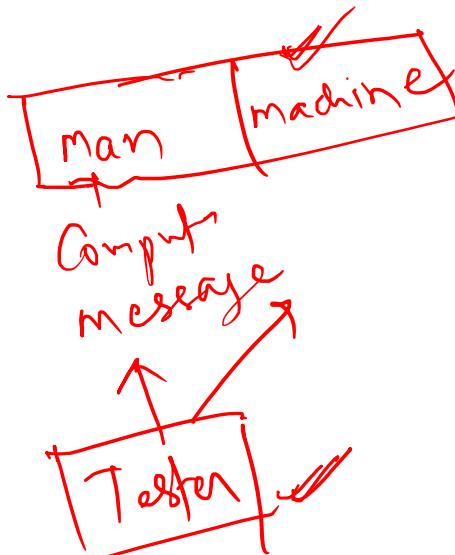
Health / Diagnosis

- ① Searching ✓
- ② Knowledge Representation
- ③ Prob. Reasoning ✓
- ④ Game Play ✓

1950 Alan Turing

Turing Test

Can m/c think?



Intelligent Behaviour  
is an ability to achieve  
human level performance  
in all cognitive tasks &  
sufficient to fool the  
interrogator

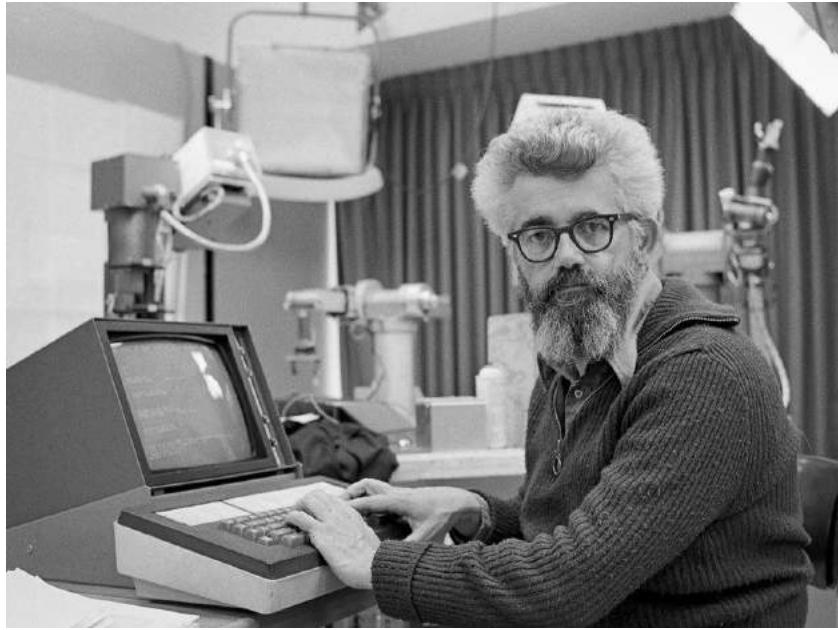
1956      father of AI

AI

2000      Chess

John McCarthy

- ✓ Algorithms
- ✓ Computing power
- ✓ Data



1956

father of AI

John McCarthy

# Turing test

During the Turing test, the human questioner asks a series of questions to both respondents. After the specified time, the questioner tries to decide which terminal is operated by the human respondent and which terminal is operated by the computer.

■ QUESTION TO RESPONDENTS ■ ANSWERS TO QUESTIONER

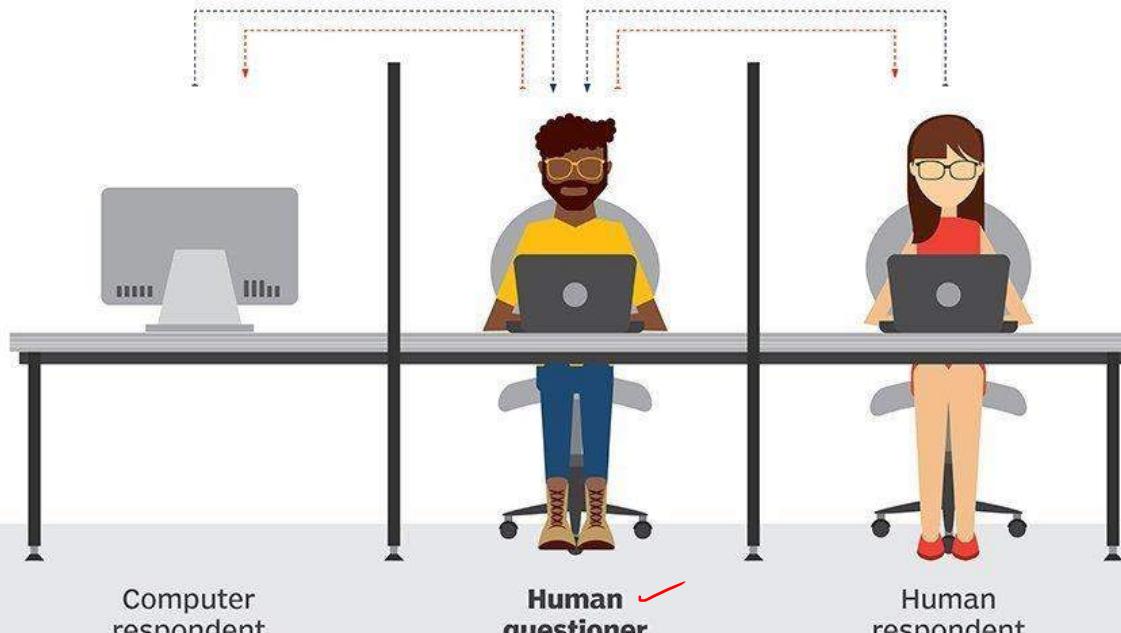


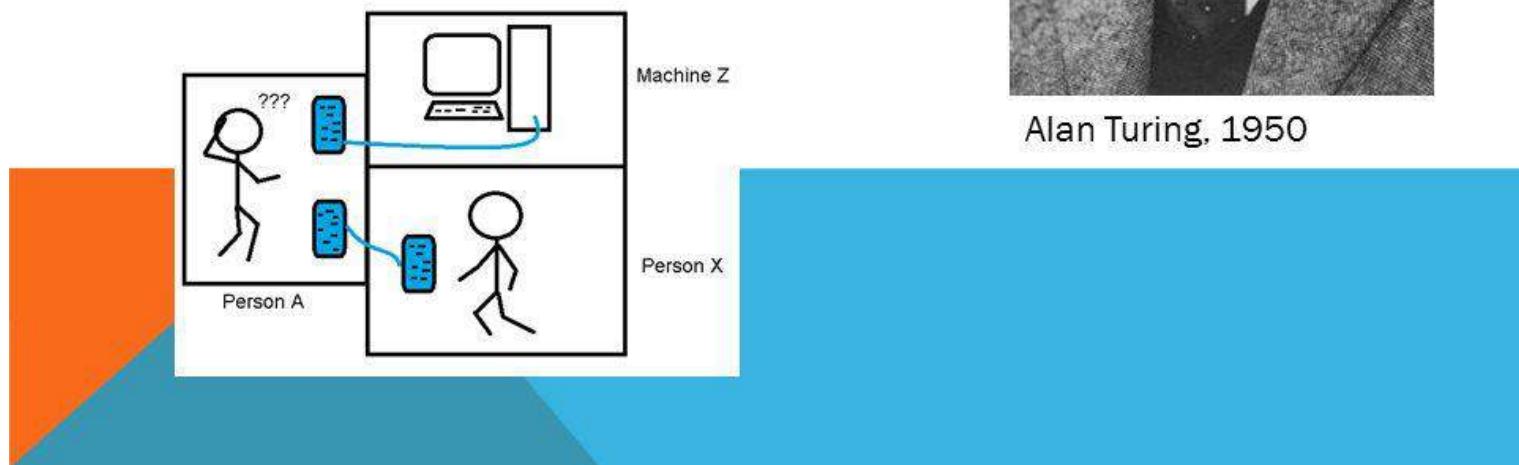
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# “CAN MACHINES THINK?”

What is the Turing Test?

- Hypothetical test called the imitation game to determine if a machine has artificial intelligence.
- Turing predicted that by 2000 it would be possible to fool an average interrogator with the probability of at least 30%



Alan Turing, 1950

# Reasoning logic 1950

A\* 1966 Alg.  
Prob. Reas AI To - 2010  $\Rightarrow$  Comp + Power  
+ Data

Google Deep mind 2015 19 x 19 Board

Formal Cognitive task

Game - chess  
checker

math.  
- logic  
- Calculus  
- T-P.

Expert fast

Engg - Design  
- fault find  
- manuf

Medical - Diagnosis  
- Stock market  
Predict

# perceptual task

C. Vision

Speech

NLP      Robot Control

Agent :- Anything that can perceive its env.  
through sensor, effectors.

Rational Agent . .

Logic

Rational thinking is the ability to consider the relevant variables of the situation & to access, organize & analyse relevant info (fact, opinion, judgement) to arrive at a sound conclusion

Rational Agent - An agent should strive to do  
right things, based on what it  
can perceive & the action it can perform

Performance

Agent type  $\Rightarrow$  Car Driver

Percepts  $\Rightarrow$  Speedom, GPS,  
microphone, camera

Action  $\Rightarrow$  Steer  
Control  
Vehicle  
Brake

Goal  
Safe

Env  
Road  
traffic

Weak AI

Machine could act as if they  
are intelligent

Strong AI

mk that act intelligent have to think  
like intelligent too

# Problem Solving

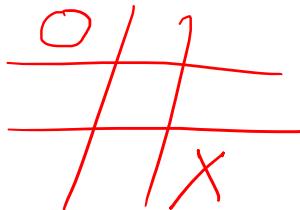
## ① Define the problem Precisely

State

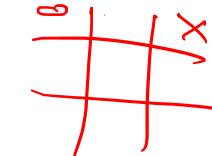
1

All info

about a particular  
env.



DFS



State space  
Repr / Tree

- (i) Initial State
- (ii) final state
- (iii) Define a state space that contains all the possible configuration of the relevant objects

Output      Path      Initial state  $\xrightarrow{\text{or start state}}$  Goal state

- ② Analyze the problem :- Various technique for solving the problem  
Knowledge DB
- ③ Isolate & Represent the task knowledge that is necessary to solve the problem
- ④ Choose the best prob-solving technique & apply it to the particular
- ⑤ Specify | Generate the Set of rules | Production Rule  
That describes the actions.  


# Production System

A P.S. is a set of programs or mechanism typically used to ~~solve~~ provide solution to particular problem which are concerned with AI

## ① Production Rules

$$C_i \rightarrow A_i$$

Conditional Part

Actions

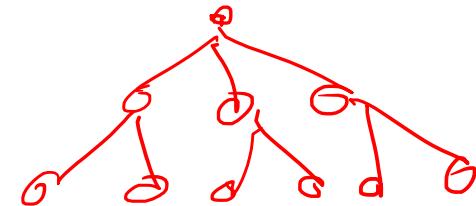
↳ functions

↳ call to another

fun

↳ call to another Production Rule

$$S \rightarrow A$$



Search is General mechanism that can be used when no more direct method is known

② Knowledge DB → Info appropriate for particular task.

③ Control Strategy

Specifies the Order in which rules will be composed to the DB & way of resolving conflict.

→ BFS, Hill Climbing.  
DFS  
A\*

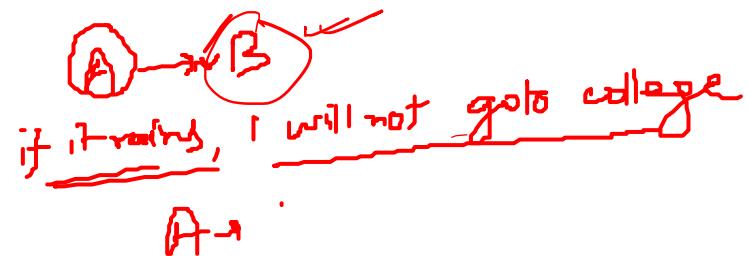
## Characteristic

### Types of Production System

- ① Monotonic Production System  
(monotonic Reasoning)

$$A, B, C, D \Rightarrow E$$

Adding more knowledge does not reduce the set of propositions that can be derived.



$$A \wedge C \rightarrow \neg B$$

Anything that could be concluded ,

clause

If a conclusion is not invalidated by adding some more knowledge.

- ② Non monotonic P.S.  $\rightarrow$  If some conclusions can be invalidated by adding some more knowledge .

③ Partially Commutative P.S. → A system in which the app of a particular sequence of rules transforms state  $x$  into state  $y$ , then any permutation of these rules that is allowable also transforms state  $x$  into state  $y$ .

Rule 1      A  
Rule 2      B  
Rule 3      C

Cond 1       $\rightarrow$  D

C      A      B      D

A      X into state Y, then any permutation of these rules that is allowable also transforms state X into state Y.

Ex math op ~~(A+B)~~

$$\begin{array}{ccc}
 X & A + B \xrightarrow{*} C \\
 & 2 + 3 \xrightarrow{*} 5 & 2 \xrightarrow{*} 5 + 2 \\
 & 5 \xrightarrow{*} 5 & 15 + 2 \\
 & 25 & \neq 17
 \end{array}$$

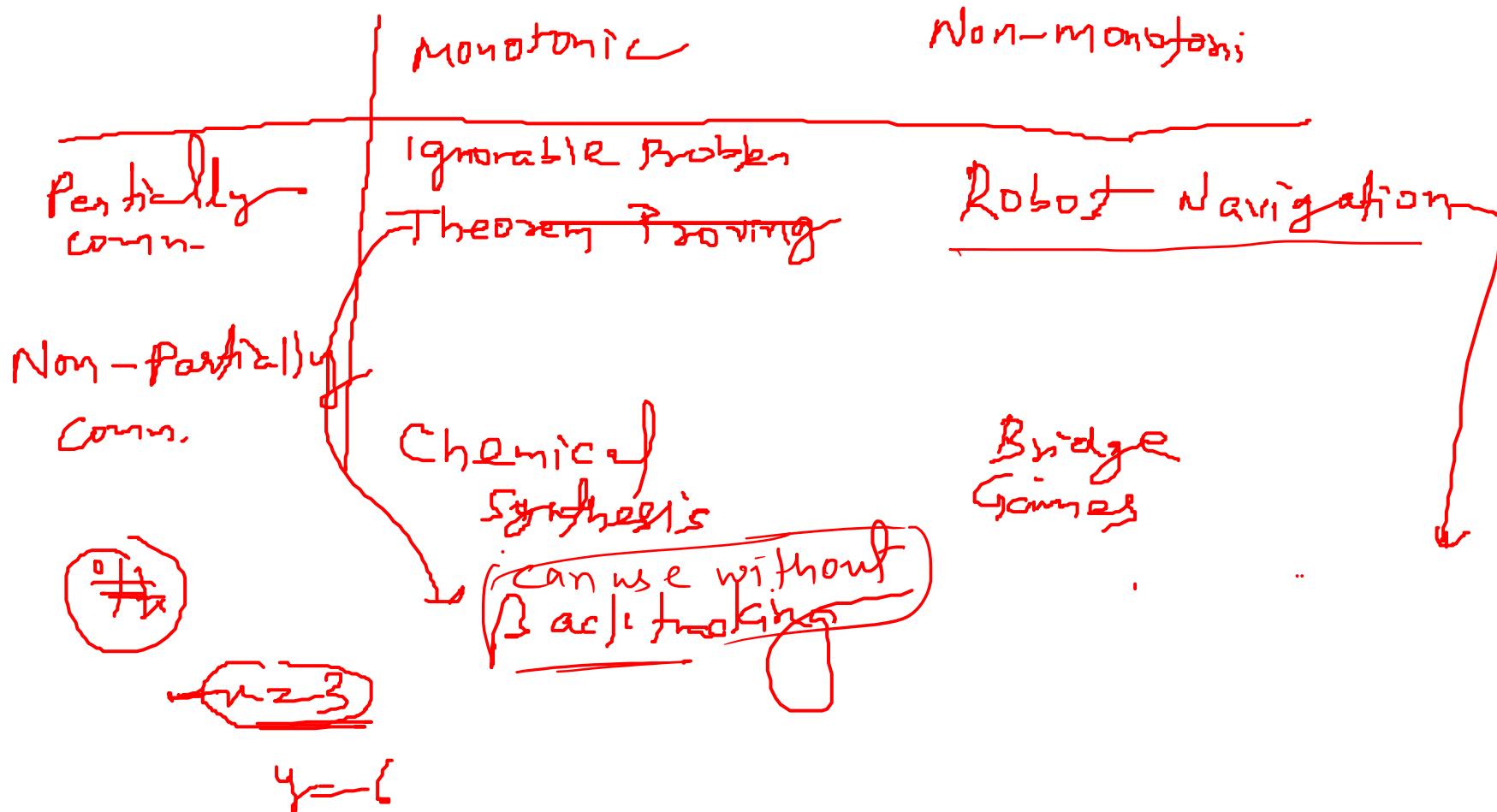
① & ③

④ Comm. Production System → Both monotonic & Partially Comm.

App

(1) Rich & Knight

## Expert System Automated Planning



Commutative  $\rightarrow$  Ignorable Problem  
(Theorem proving)

4	5	
1	2	3
6	7	8

Creating new thing rather than old things  
Can use without backtracking

~~Partially Com + Non-monotonic~~ (Robot Navigation)

AEN

Order of the app is not critical

ENH

Non-Partially + Monotonic  
(Chemical Synthesis)

8 puzzle  
problems

# Features of Production Systems

- ① Simplicity  $\Rightarrow$  if then structure  
    ↳ Knowledge Representation improves readability
- ② Modularity  $\Rightarrow$ 
  - ↳ Addn / Deletion is easy
  - ↳ discrete (knowledge)
- ③ Modifiability  $\rightarrow$  Updations of new rules
- ④ Knowledge Intensive  $\rightarrow$  stores pure knowledge

## Issues / Challenges) Drawbacks.

### ① Inefficiency $\Rightarrow$

A well devised (control) strategy to reduce this.

### ② Opacity $\Rightarrow$ 10 rules

Conflict Resolution  $\Rightarrow$  More priority to rules has less opacity

### ③ Absence of learning $\Rightarrow$ Rule-based Production

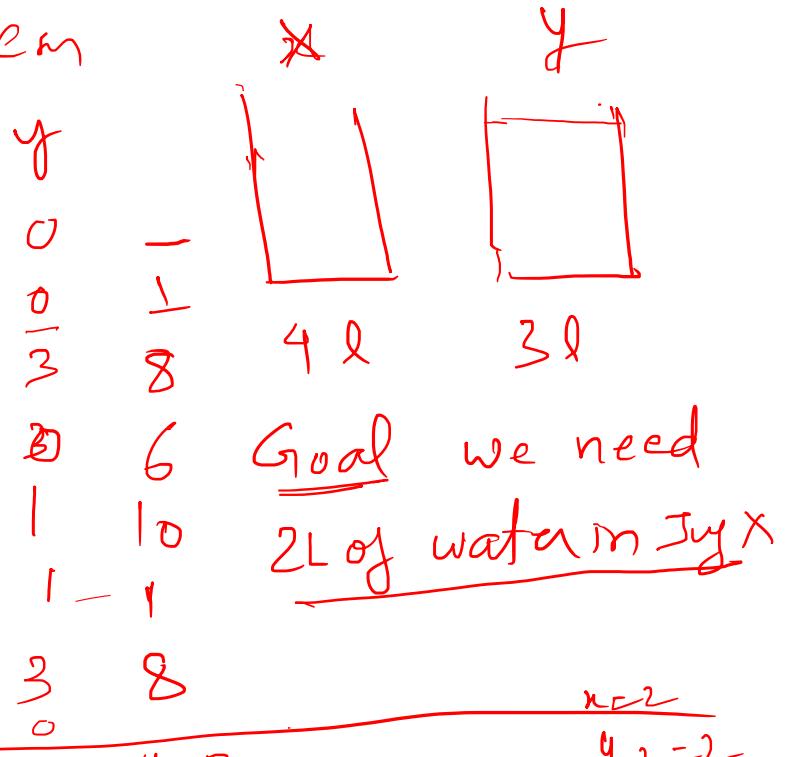
Rule 1 = Result (future use)  
Rule 2 = Result (future use)

does not use result for future use, hence it does not exhibit learning capabilities.

# Water Jug Problem

<u>Rule applied</u>	$x$	$y$	$x$	$y$
-	0	0	0	0
2	0	3	1	1
9	3	0	1	3
2	3	3	0	6
7	4	2	4	10
5	0	2	2	8
9	2	0	2	0

30 rules



- ✓ ①  $x \leq 4$   $(x, y) \rightarrow (4, y)$
- ✓ ②  $y \leq 3$   $(x, y) \rightarrow (x, 3)$
- ③  $x > 0$   $(x, y) \rightarrow (x - p, y)$
- ④  $y > 0$   $(x, y) \rightarrow (x, y - p)$
- ✓ ⑤  $x > 0$   $(x, y) \rightarrow (0, y)$
- ✓ ⑥  $y > 0$   $(x, y) \rightarrow (x, 0)$

- ⑦  $x \geq 0, x + y \geq 4$   $(x, y) \rightarrow (4, y - (4 - x))$
- ⑧  $x \geq 0, x + y \geq 3$   $(x, y) \rightarrow (x - (3 - y), 3)$
- ⑨  $x + y \leq 4, y \geq 0$   $(x, y) \rightarrow (x + y, 0)$
- ⑩  $x + y \leq 3, x \geq 0$   $(x, y) \rightarrow (0, x + y)$

## Control Strategies Requirements

To decide which rule to apply<sup>next</sup> during the process of searching for a sol<sup>n</sup> to a problem.

① It cause motion :-

② It should be systematic

It should take less time

# Search Strategies

⇒ Uninformed Search

① BFS

② DFS

Depth Limited Search

Uniform cost search

Beam Search

⇒ Informed Search

hill climbing

Best first search

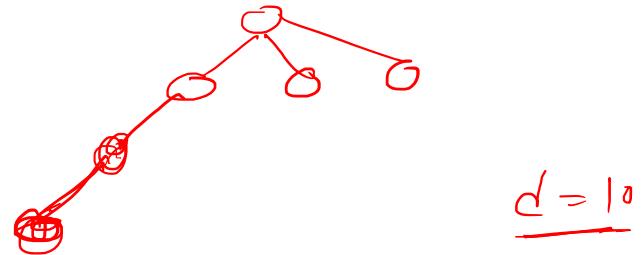
A\*

AO\*

A Search Strategy decides the order of the node expansion

✓ State Space Tree

- ↳ Completeness
- ↳ Time Complexity
- ↳ Space Complexity
- ↳ Optimality
- ↳ Systematicity

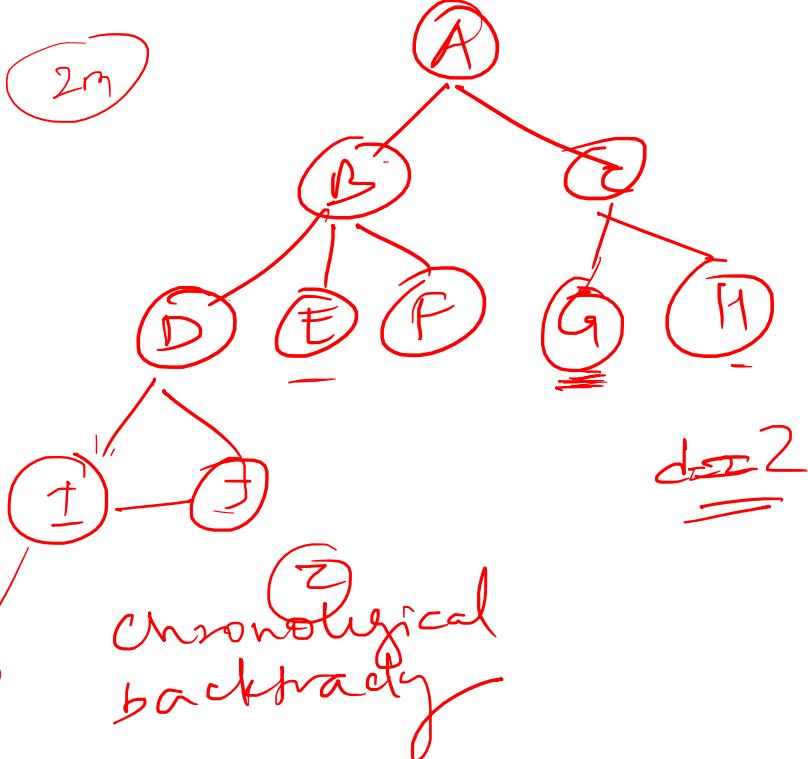


b - branching factor  
 $\stackrel{\text{max.}}{\text{branching factor}}$

d - depth (least cost)

m - max. depth of state space ( $\infty$ )

# Depth first Search



Adv. \* Requires less memory  
since only nodes on the current path are stored

- \* Find the sol<sup>n</sup> without exploring much of the state space

maintain stack of nodes to visit



ACG

depth finite

Complete? No

Optimal? No

Time Comp -  $O(b^m)$

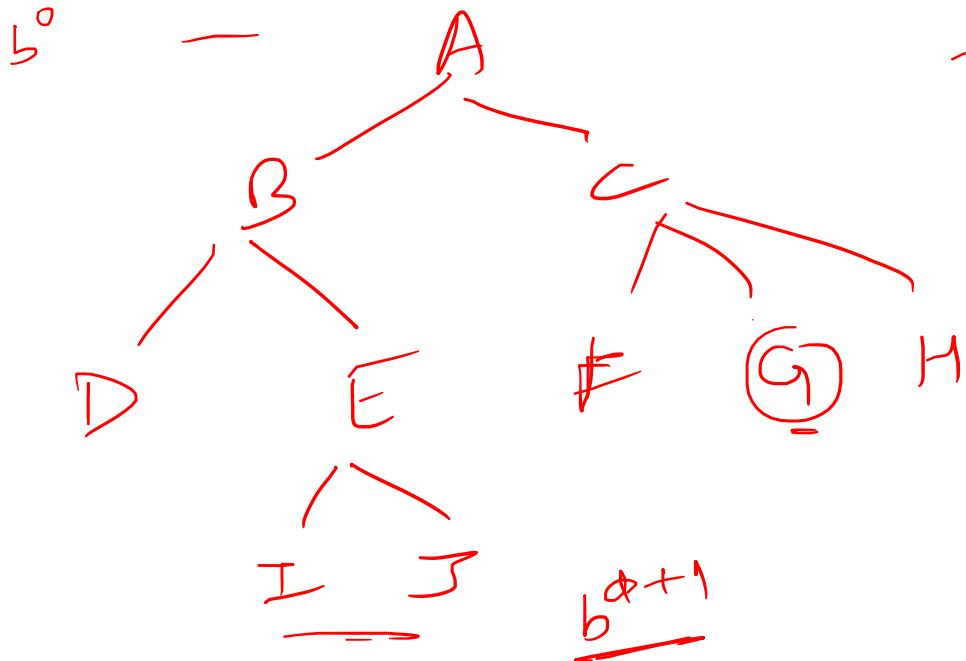
Space Comp -  $O(b^m)$

disadv. it may get trapped in a dead-loop

# Breadth first Search (BFS)

maintain Queue of nodes to visit

Layerwise



Traversal  $\overbrace{ABCDEF}^n$  nodes

draw Storage req. more

Adv = ① It doesn't get trapped in a dead loop

② It always finds the soln (complete)

Space complexity  $\rightarrow \frac{(b^d)}{(b^d)}$   
Time  $\rightarrow \frac{(b^d)}{(b^d)}$

## BFS

level by level  
level order traversal

Queue

Slower

Complete

Complexity -

→ Greedy Algo (lowest node first)

App → shortest path  
Prim | Kruskal

Web crawling

θ<sub>i</sub>

Local optimum sort

## DFS

deep

Preorder

Stack

fast

Not a complete

Complexity

Backtracking (explore one path until dead end)

App → find strongly connected component

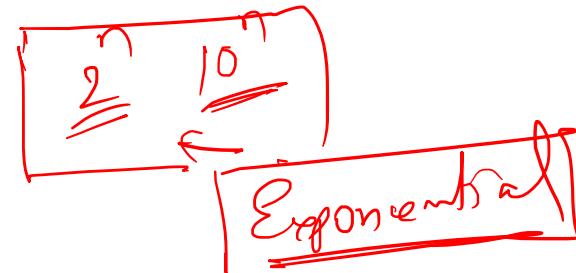
Uniform Cost Search -  $\text{edges} \infty t = 1$

Blind  
Search

Beam Search =  
(BFS)

Storage  $\propto$  nodes

Complete - NO  
Optimal - NO



Depth limited Search -  
(DFS)

d is fixed

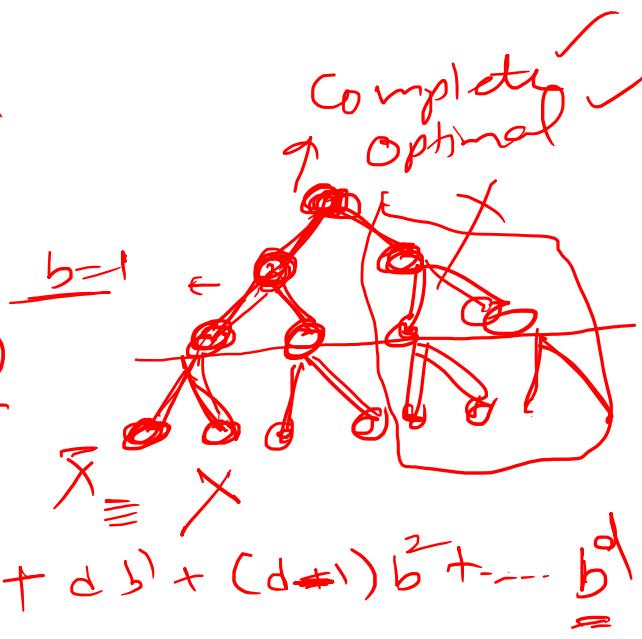
Iterative Deepening Search

$d=0$        $d=1$   
 $d=2$

Space Comp  $O(bd)$

Time Comp =

$$(d+1)b^0 + d b^1 + (d+1)b^2 + \dots b^d$$



$$\begin{array}{r}
 b \\
 | \\
 10^2 + 10^1 + 1 = 111 \\
 \hline
 \end{array}$$

## Time vs. Memory

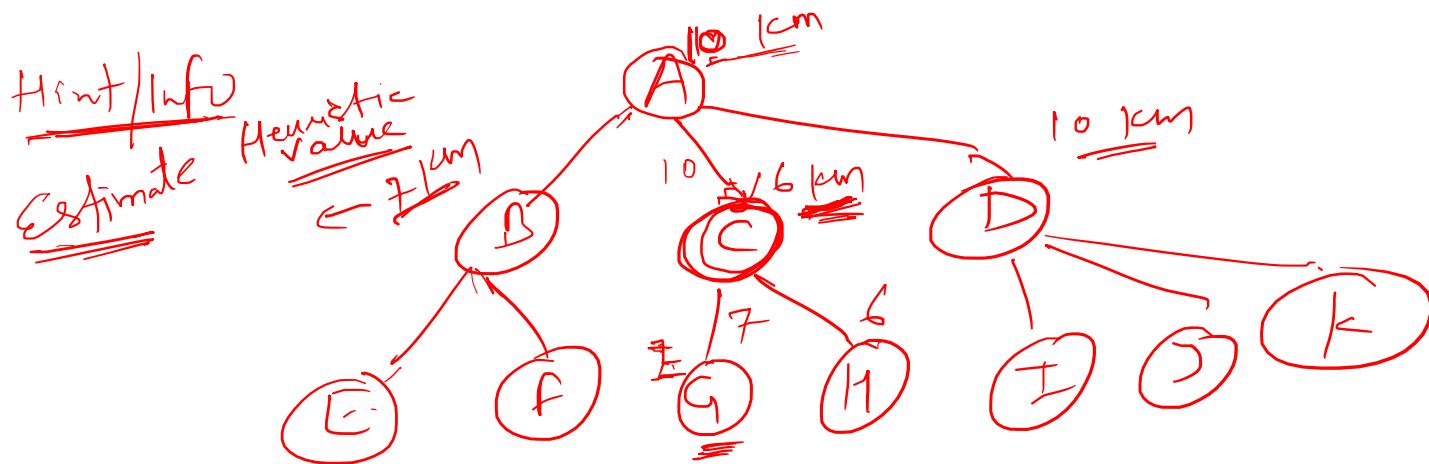
Depth	Nodes	Time	Memory
2	110	.11 milliseconds	107 kilobytes
4	11,110	11 milliseconds	10.6 megabytes
6	$10^6$	1.1 seconds	1 gigabyte
8	$10^8$	2 minutes	103 gigabytes
10	$10^{10}$	3 hours	10 terabytes
12	$10^{12}$	13 days	1 petabyte
14	$10^{14}$	3.5 years	99 petabytes
16	$10^{16}$	350 years	10 exabytes

**Figure 3.13** Time and memory requirements for breadth-first search. The numbers shown assume branching factor  $b = 10$ ; 1 million nodes/second; 1000 bytes/node.



## Unit - II

### Informed Search (Heuristic Search)



Node	$h(n)$
A	5
B	3
C	2
D	3
E	1
F	1

Evaluation function  $f(n)$

Node is selected for expansion based on an eval. function that estimates cost to goal

$g(n)$  = Cost of current node from the initial node.

$h(n)$  = Heuristic fn.: Estimated cost from current node to goal node

# (Greedy) Best first Search

It combines the advantages of BFS & DFS

It follows single path at a time, but switches path whenever some competing path looks more promising than the current one.

DS  
Priority Queue

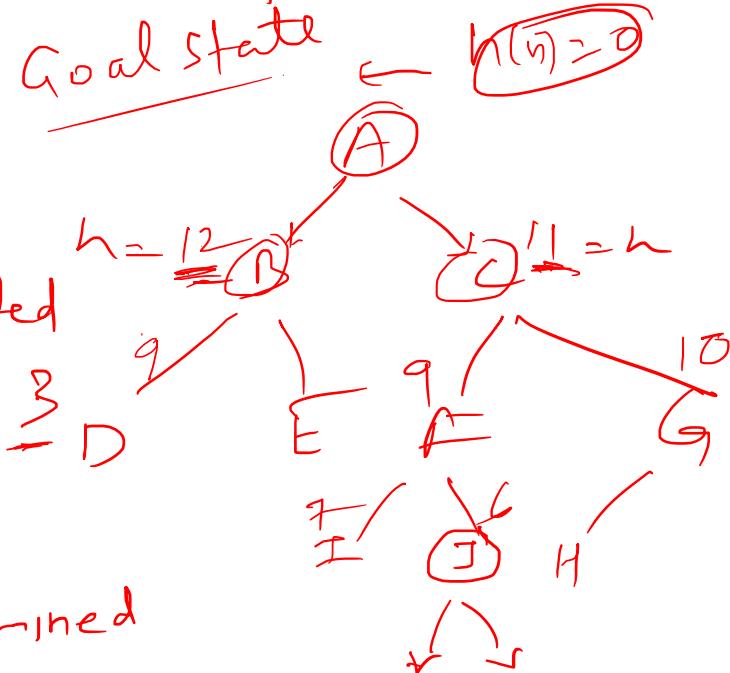
$$f(n) = h(n)$$

OPEN  $\Rightarrow$  Nodes that have been generated & have had heuristic function applied to them but which have not yet examined

CLOSED  $\Rightarrow$  Nodes that have ~~been~~ been examined

OPEN ~~A F B~~  $f(n) = h(n)$

✓ CLOSED A C F J



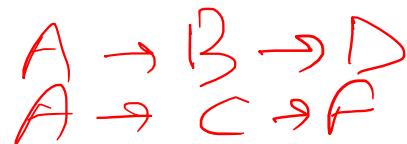
DFS

Adjacency Matrix



0	A	-	1	1	0	0
1	B	†	-	0	1	1
2	C					
3	D					
4	E					
5	F					

Adjacency list



✓ Preorder traversal

① Scanf (n)

for

for

② &G[i][j] for

③ visited[i] = 0

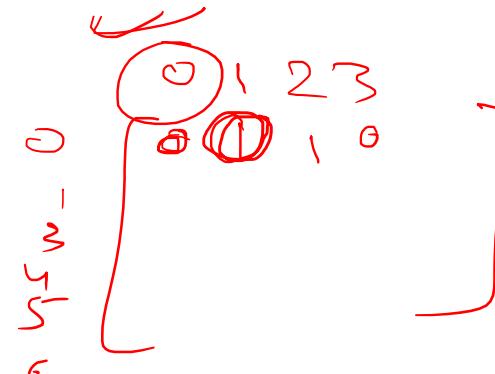
DFS(0)

DFS({ $\emptyset$ })

visited[i] = 1

for (j : visited[j])

if (!visited[j] & &G[i][j] = 1)  
DFS(j)

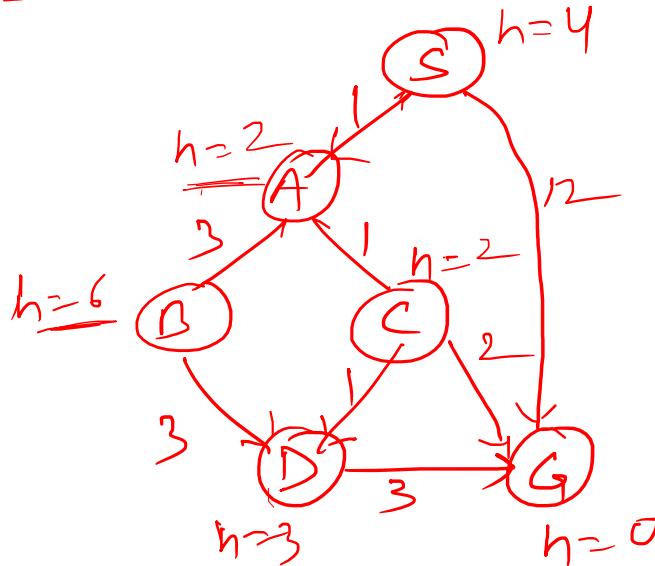


A\* Algo.

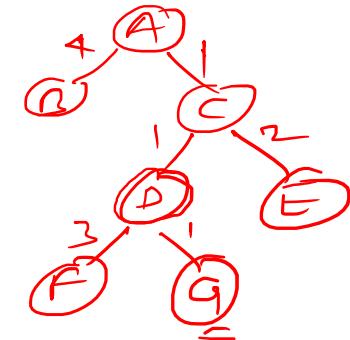
consistent

$$f(n) = g(n) + h(n)$$

Ex



$$(S \rightarrow A \rightarrow C \rightarrow G) = 4$$



$$\begin{aligned} A \text{ to } G \quad g(n) &= 1+1+1 \\ &= 3 \end{aligned}$$

$$A \text{ to } E \quad g(n) = 1+2 = 3$$

$$\begin{aligned} A \text{ to } D \quad g(n) &= 1+1 = 2 \\ h(n) &= 1 \end{aligned}$$

$$\boxed{S \rightarrow A} \quad f(n) = 1+2 = 3$$

$$\begin{aligned} S \rightarrow G & \\ f(n) &= 12+0 = 12 \end{aligned}$$

$$\begin{aligned} \overline{\text{OPEN}} \quad & \text{OPEN } CBG \\ \overline{\text{CLOSE } SA} & \\ S \rightarrow A \rightarrow C \rightarrow D & \\ f(n) &= 3+3 = 6 \end{aligned}$$

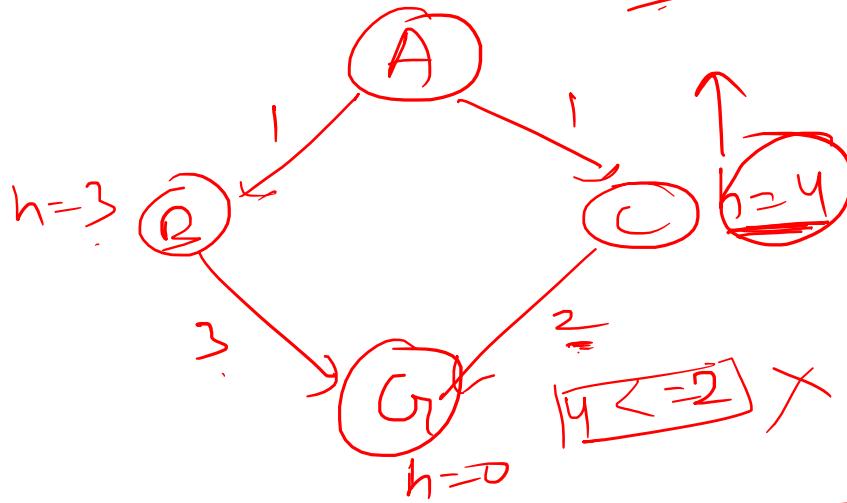
$$\begin{aligned} S \rightarrow A \rightarrow B & \quad \text{or} \\ f(n) &= 4+6 = 10 \end{aligned}$$

$$\begin{aligned} S \rightarrow A \rightarrow C & \\ f(n) &= 2+2 = 4 \end{aligned}$$

$$\begin{aligned} S \rightarrow A \rightarrow C \rightarrow G & \\ f(n) &= 4+0 = 4 \\ \underline{\text{Stop}} & \end{aligned}$$

heuristic fun<sup>n</sup>

Admissibility



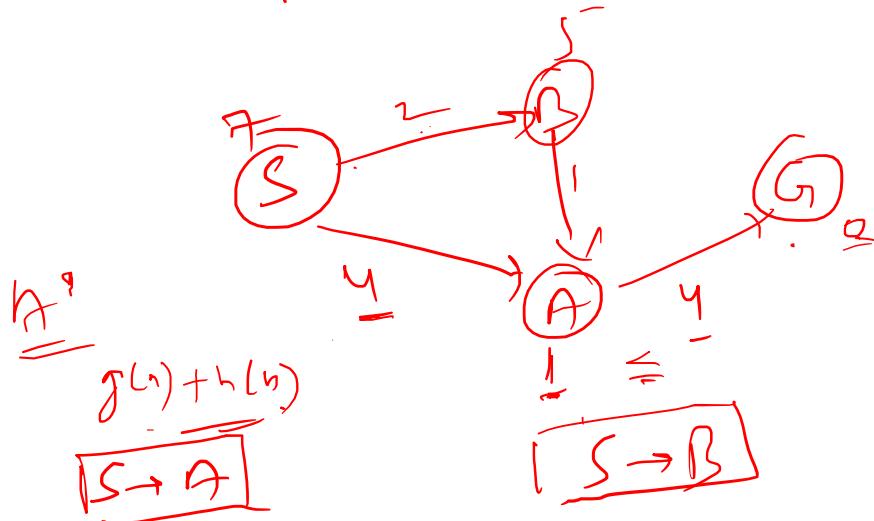
A H.F. is admissible if never overestimates the cost of reaching a goal from current node.

for every node n,

$$h(n) \leq \underline{h^*(n)} \rightarrow \begin{array}{l} \text{True value} \\ \text{cost} \\ \text{to reach the} \\ \text{goal node} \end{array}$$

If  $h(n)$  is Admissible ;  $A^*$  is optimal (tree search)

Heuristic function should be consistent (Graph searchs)



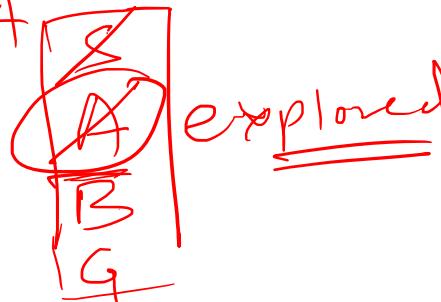
$$f(n) = 4 + 1 = 5$$

$$S \rightarrow A \rightarrow G$$

$$f(n) = 8 + 0 = 8$$

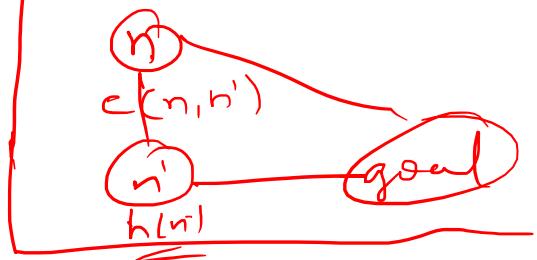
$$f(n) = 2 + 5 = 7$$

$$S \rightarrow B \rightarrow A$$



- ✓ Admissible
- ✗ Consistent

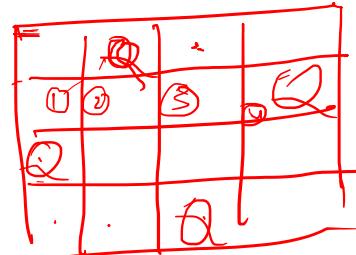
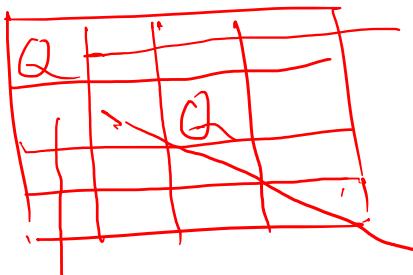
$h(n)$  is consistent if for every node  $n$  & for every successor  $n'$

$$h(n) \leq c(n, n') + h(n')$$


- # Every consistent H.F. is Admissible
  - # If  $h(n)$  is consistent,  $a^*$  is optimal
- (minimizes)
- 

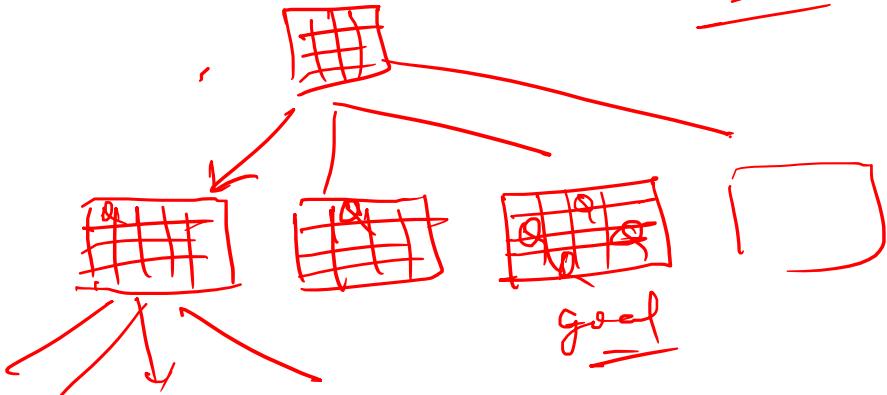
## n-Queens Problem

n=4



goal

Sampling  
⇒ Random Walk  
Generate a state  
randomly

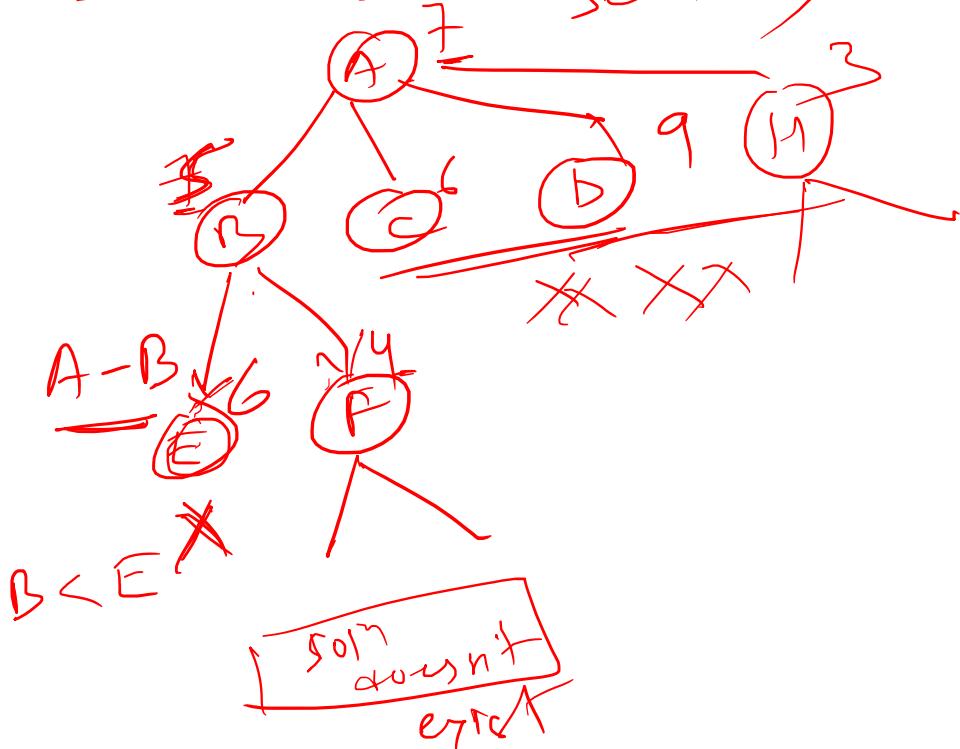


Random Walks  
Randomly picking  
a neighboring state of the current state

① Generate & Test Algo' (h-f) X  
 (Random Walk)



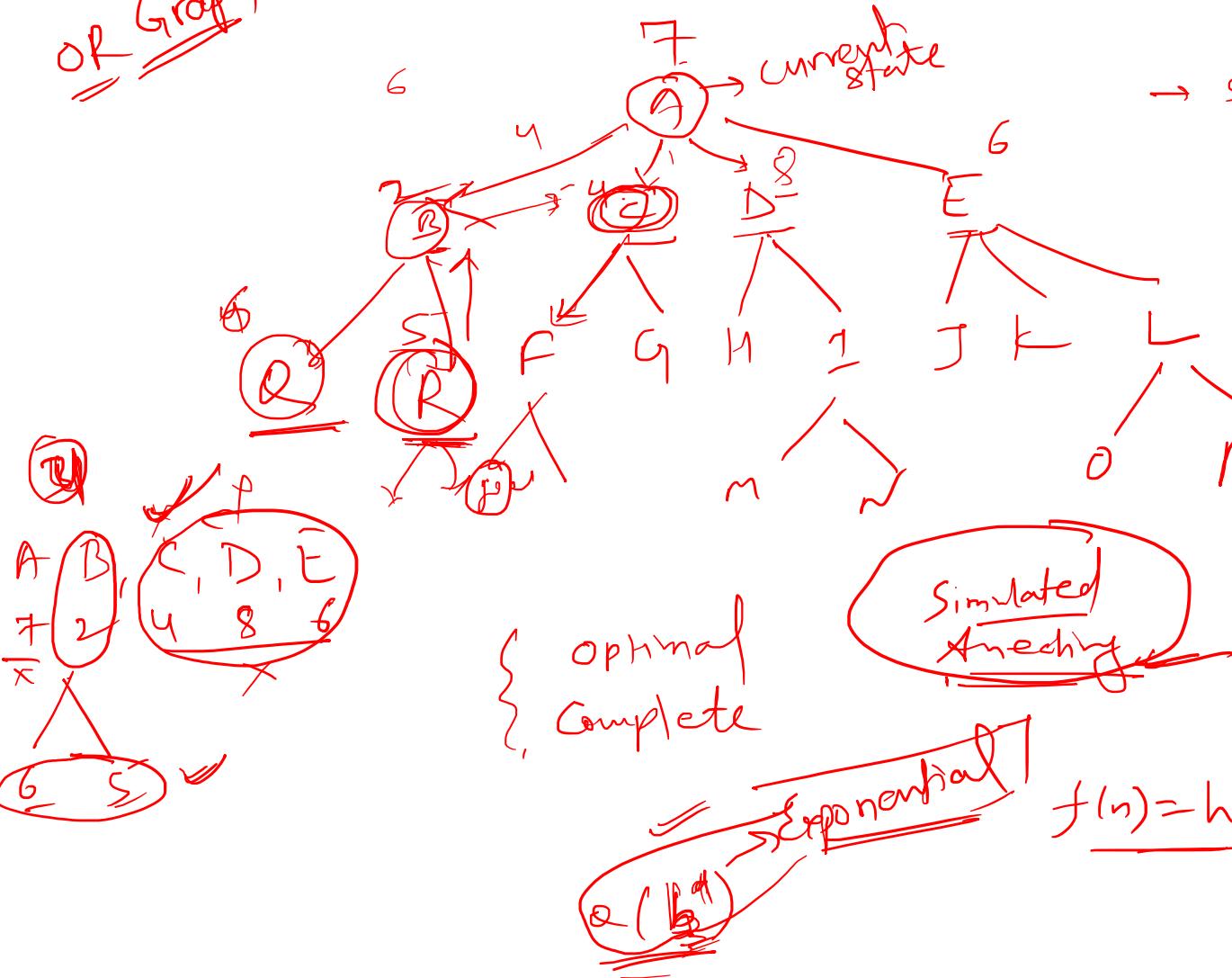
② Hill Climbing (Greedy Local Search)



X A - C - G - J X  
 X A - D - H X  
 X A - C - G - I

## Revision

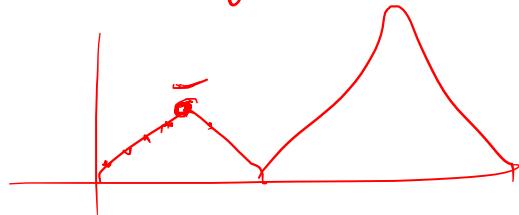
of Graph



- ① Random Sampling
  - ② Random (Get) Walk
  - ③ Hill Climbing
  - ④ Simple Steepest Ascent Hill Climbing
  - ⑤ Best first search
  - ⑥ A\* Algorithm
- $$f(n) = g(n) + h(n)$$

## Drawbacks of Hill Climbing

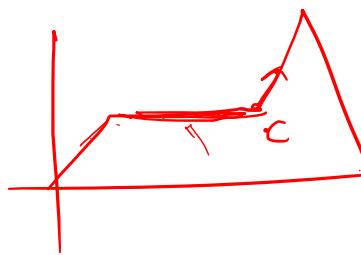
Local Maximum



State that is better than all its neighbouring state but is not better than some other state further away

Plateau

flat area of search space  
in which all the neighbours state of current state contains the same value

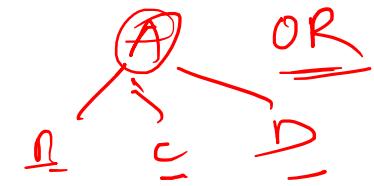
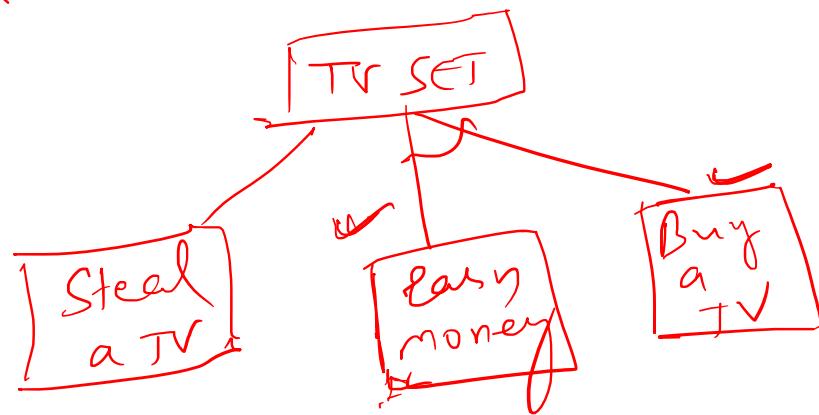


Ridge

Special form of local maxima. It has an area which is higher than surroundings but itself has a slope & cannot be reached in a single move

AO<sup>n</sup> Algo

AND-OR GRAPH

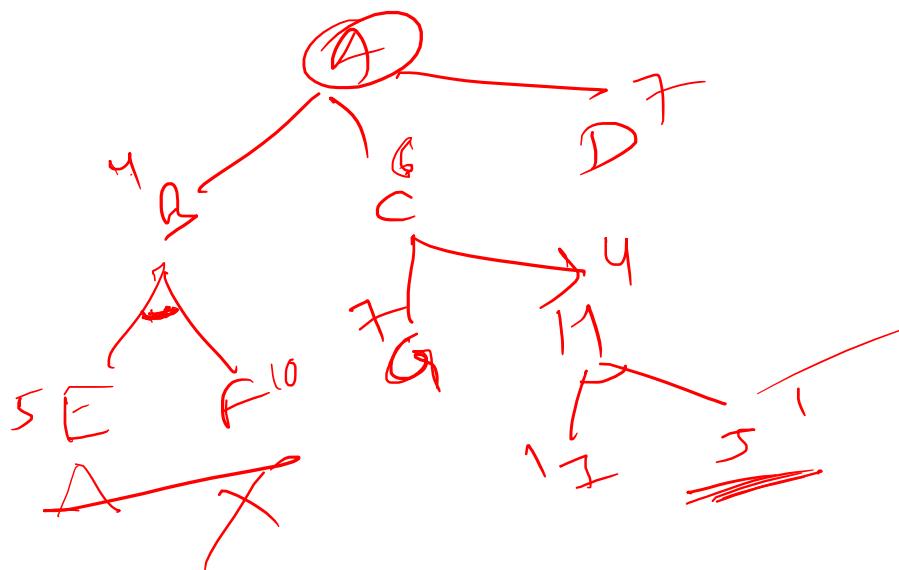


Stage 1

FUTURE

thresh<sup>1</sup>

cost < 20



# Constraint Satisfaction Prob.

## Cryptarithmetical Problem

- ① Digits  $0 \rightarrow 9$
- ② Variable  $\rightarrow$  unique & distinct
- ③ Base is 10
- ④ Number must not begin with 0

$$\begin{array}{r}
 \text{S E N D} \\
 + M O R E \\
 \hline
 \text{M O N E Y}
 \end{array}$$

$$\begin{aligned}
 E + I &= N \\
 N + R &= E + 10 \\
 E + I + R &= E + 10 \\
 R &= 10 - E - I
 \end{aligned}$$

$\times 10 \mid 8$

$$\underline{A \rightarrow 4}$$

$$\underline{B \rightarrow 6}$$

$$\times \boxed{\overline{10 \mid 23}}$$

$$\begin{array}{r}
 \text{S} \rightarrow 9 \\
 \text{M} \rightarrow 1 \\
 \text{O} \rightarrow 0 \\
 \text{R} \rightarrow 8 \\
 \text{D} \rightarrow 7 \\
 \text{E} \rightarrow 5 \\
 \text{N} \rightarrow 6 \\
 \hline
 \text{D+E = 4}
 \end{array}$$

$$\begin{array}{r}
 \cancel{H+4} \\
 \times \cancel{H+G}
 \end{array}$$

$$\begin{array}{r}
 \cancel{7+5} \\
 \times 6+5
 \end{array}$$

~~L E S~~  
~~S L S~~  
~~S S S~~

BASE  
BALL  
GAMES  
12 4/5 S E  
S E

TWO  
TWO  
~~POKE~~

A = 2  
B = 6  
G = 1  
m - 4 / 5 =  
B → 9 + 9  
8 + 8  
7 + 7  
AS → 6 + 6  
x S + S →

E  
S + 3 = S  
E + 6 = S  
S + L = E 8

L → 0  
E → 9  
S → 8

3 + 5 = 8  
S → 3