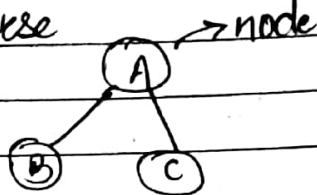


UNIT - 5 TREE & GRAPH

Root
Left
Right

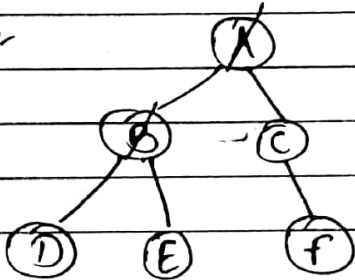
(1) Traverse



A | B | C

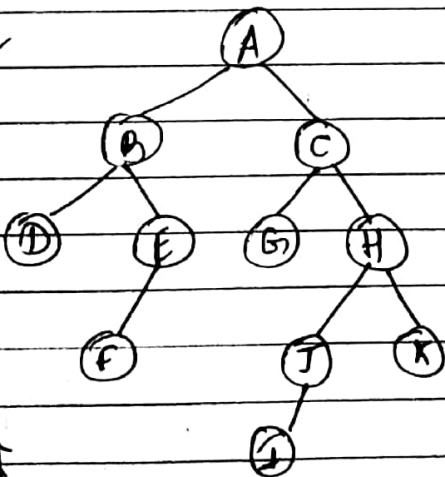
- (a) Pre → Root, left, right
- (b) Post → left, right, root
- (c) In → left, root, right

(a) Ex:-



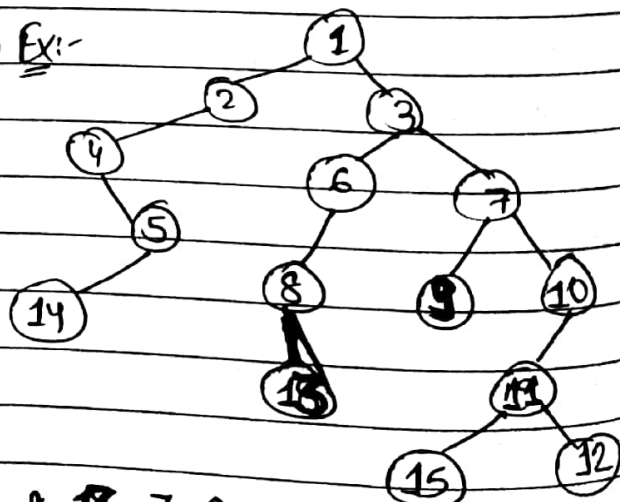
A, B, D, E, C, f (Pre)

(b) Ex:-



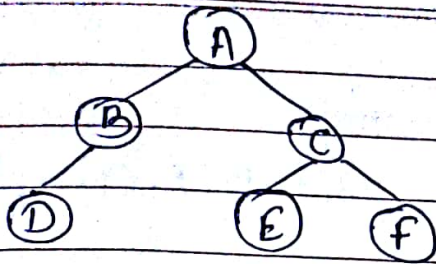
Pre,
A, B, D, E, f, C, G, H, J, I, K

(c) Ex:-



Pre, 1, 2, 4, 5, 14, 3, 6, 8, 13, 7, 9, 10, 11, 15, 12

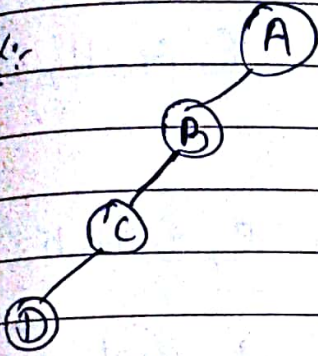
Ex:-



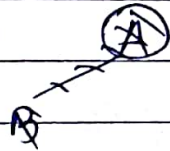
In,

D, B, A, E, C, F

Ex:-

Pre \rightarrow A, B, C, DIn \rightarrow D, C, B, APost \rightarrow D, C, B, A

Ex:-

Ex(b) \rightarrow In,

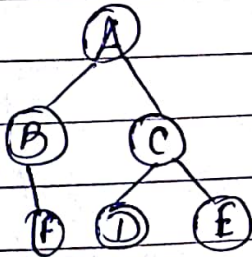
D, B, E, E, A, G, C, I, J, H, K

Ex(c) \rightarrow In,

4, 14, 5, 2, 1, 8, 6, 3, 9, 7, 15, 11

4, 14, 5, 2, 1, 8, 13, 6, 3, 9, 7, 15, 11, 12, 10

Ex:- (f)



Post,

B, F, D, E, C, A

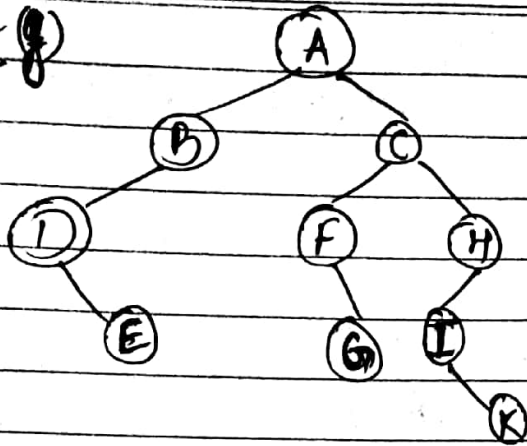
Ex:- (g) E(b) \rightarrow Post,

D, F, E, B, G, I, J, K, H, C, A

Ex(c) \rightarrow Post,

4, 14, 5, 2, 1, 8, 6, 3, 9, 15, 12, 12, 10, 7, 3, 1

Ex:- ⑧



Post,
D, E, B, F, G, I, K, H, C, A

Pre,
A, D, E, B, C, F, G, H, I, K

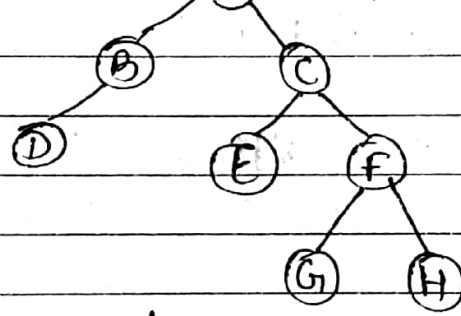
In,
D, E, B, A, F, G, C, I, H, K

Q) Inorder \rightarrow D, B, A, E, C, G, F, H

L, Root, R

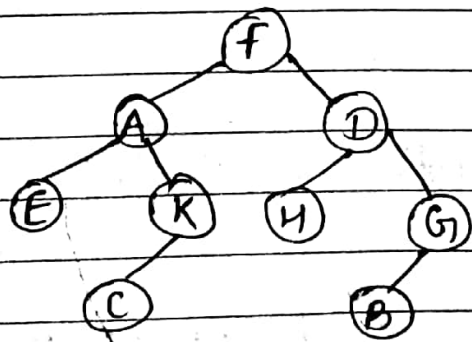
Ans) Post \rightarrow D, B, A, E, G, H, F, C, A

L, R, Root



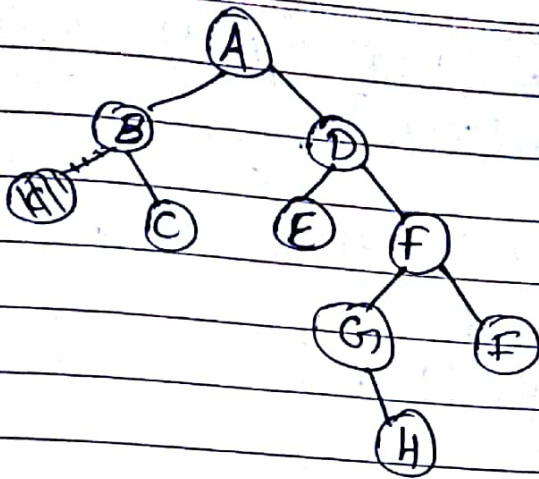
Q) In \rightarrow $\overbrace{E, A, K, K}^L, \overbrace{K, H, D, B, G}^R$
 Pre \rightarrow $\overbrace{K, A, E, K}^L, \overbrace{E, D, H, K, B}^R$

L, Root, R
 Root, L, R

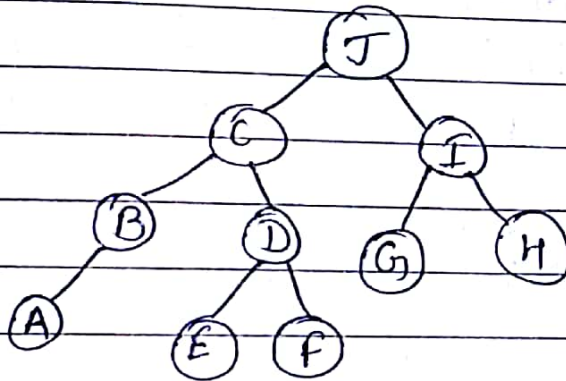


Q) Pre \rightarrow A, B, C, D, E, F, G, H, I
 In \rightarrow $\overbrace{B, C, A}^L, \overbrace{E, D, G, H, F, I}^R$

Root, L, R
 L, Root, R



Q.) Post \rightarrow A, B, ~~E~~, ~~A~~, ~~D~~, ~~E~~, G, H, I, J
 In \rightarrow A, B, ~~E~~, ~~D~~, ~~E~~, J, ~~E~~, ~~I~~, H
 L, R, Root
 L, Root, R



> Algorithm :-

- 1) Pre (node, info, left, right)
- 1) If node \neq NULL then repeat till 4).
- 2) Display info [node].
- 3) Call Pre order (left [node]).
- 4) Call Pre order (right [node]).
- 5) Exit.

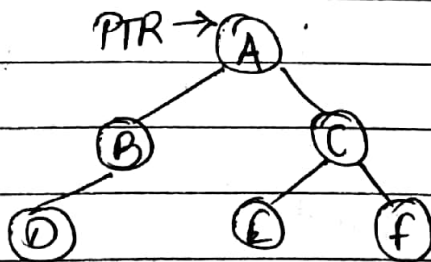
II Pre order (Info, Left, Right, Root)

- 1) Initially push NULL into the stack and initialize PTR.

Set TOP = 1 and stack [TOP] = NULL & PTR = Root

- 2) Repeat 3) to 5) while $PTR \neq NULL$.] loop
- 3) Apply PROCESS to Info[PTR].
- 4) [Right Child?]
 if $Right[PTR] \neq NULL$, then
 Set $TOP = TOP + 1$
 and $Stack[TOP] = Right[PTR]$
 (end of if)
- 5) [Left Child?]
 if $Left[PTR] \neq NULL$, then
 Set $PTR = Left[PTR]$
 else (pop from stack)
 Set $PTR = Stack[TOP]$
 and $TOP = TOP - 1$
 (end of if)
 [end of step 2) loop]
- 6) Exit.

Ex:-

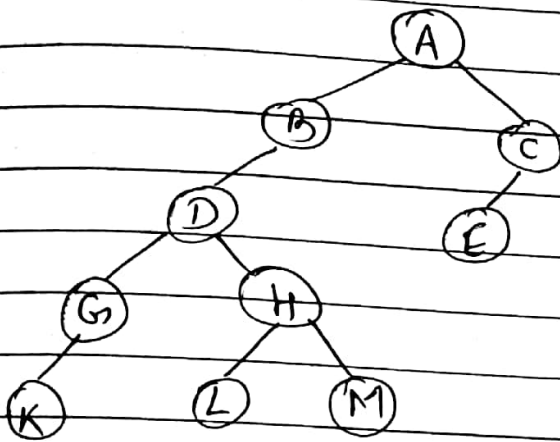


	3
	2
	1

O/P \rightarrow A, B, D, C, E, F

UNIT - 5

continued...

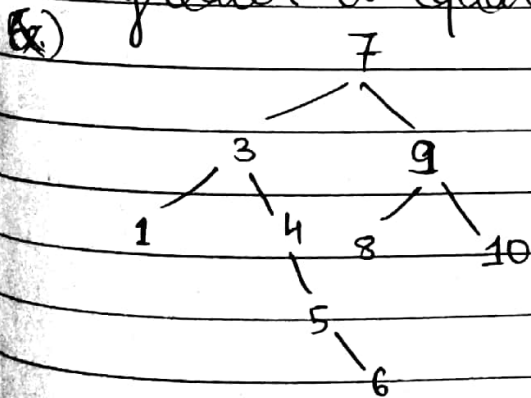


Pre \rightarrow Root, L, R
 A, B, D, G, K, H, L, M, C, E

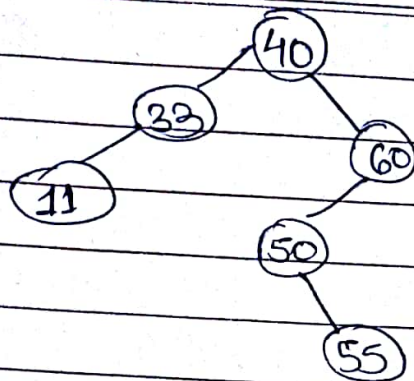
Post \rightarrow L, R, Root
 K, G

> Binary search Tree \rightarrow A binary tree is called BST if it satisfies following rules :-

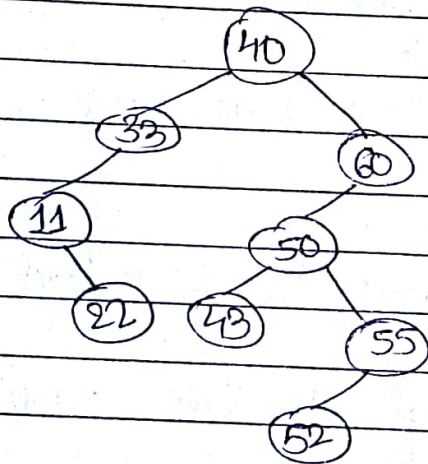
- The value of left child or left subtree should be less than the value of root.
- The value of right child or right subtree should be greater or equal to the value of root.



Q) Create binary search tree with following values :-
 40, 60, 50, 33, 55, 11

Ans.)

➤ Insert 22, 43 and 52 in existing BST.



I)

Q) Create BST → 50, 75, 25, 40, 22, 30, 33, 15, 44, 60, 90, 80

(a) Add 20 and 88.

(b) Delete 22.

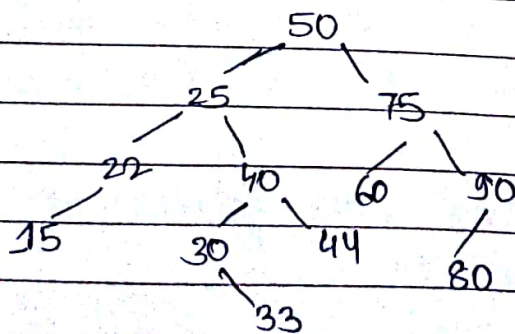
(c) Delete 25.

(d) Add 77.

(e) Delete 50 and 75

and give the output.

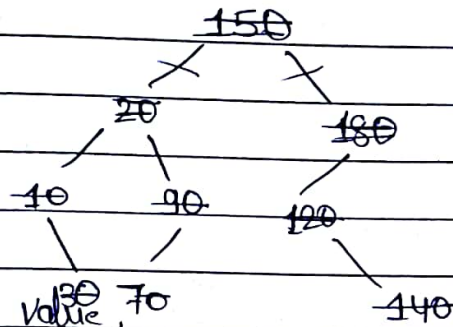
Ans.) BST →



(a)

Q) Create BST $\rightarrow 150, 180, 20, 120, 90, 140, 10, 70, 30, 50, 60, 40$

Ans.)



To delete a value, 70

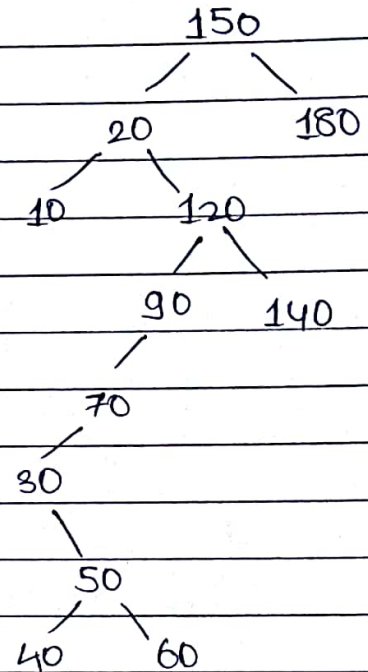
(a) find In order traversal.

In \rightarrow L, Root, R

10, 20, 30, 40, 50, 60, 70, 90,

120, 140, 150, 180

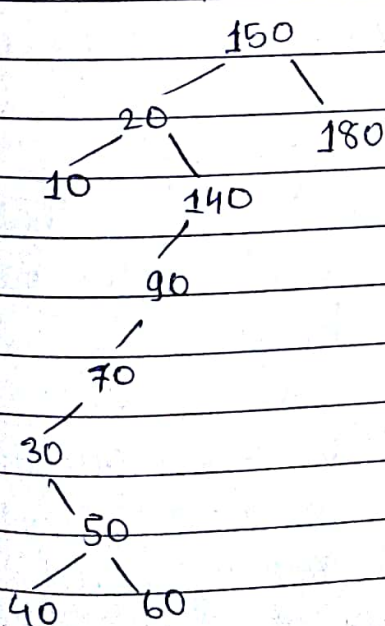
*(all the values are sorted) PC



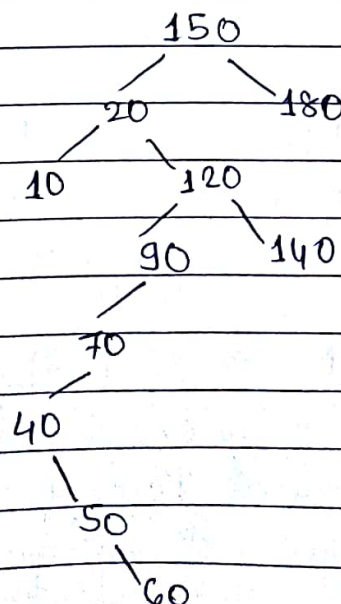
(b) Delete the element and place the In-order successor of the deleted element at its place.

4, 60,

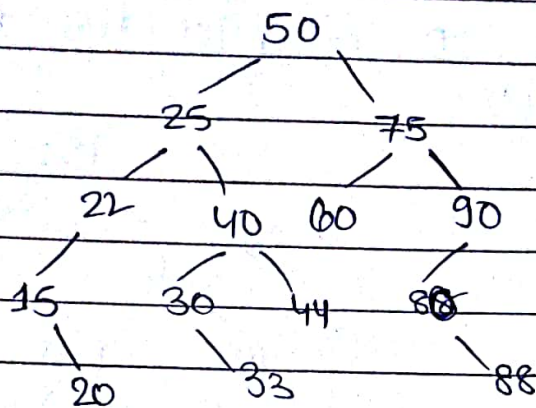
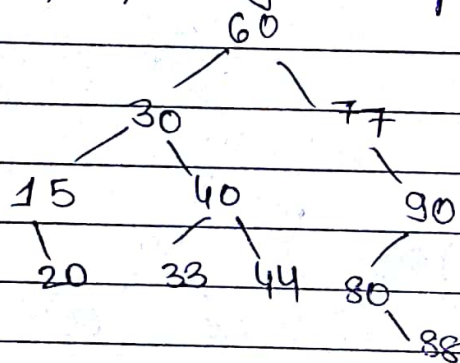
Delete 120, In-order successor = 140 (No L or R child)



Delete 30, IOS = 40 (No L or R child)



Ans. (a)

Ans. (b) In \rightarrow L, Root, R
15, 22, 30, 40After performing all operations, output \Rightarrow 

A, M, P, B, S, D, Z, C, T, N, G, Q, W, E, X

Heap sort - Heap is the DS in which values are arranged either in increasing or decreasing order. There are two types :-

(i) Max heap

they are the complete binary tree with the prop. that value at each node is either large or small to the value of its children. This prop. is called heap prop.

(ii) Min heap

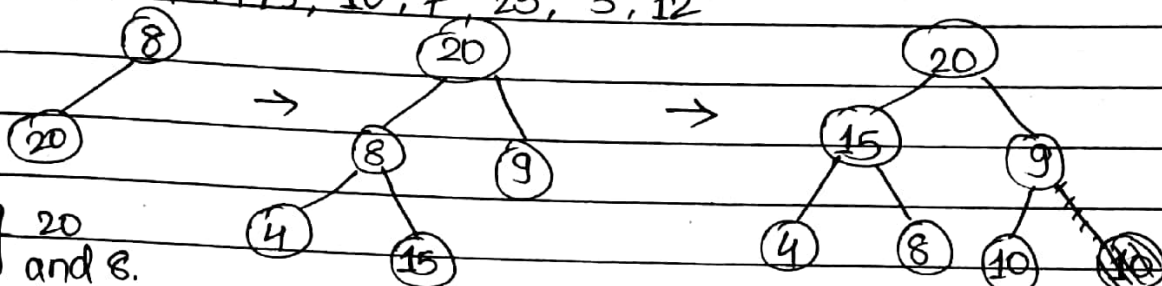
> Heapify/delete \rightarrow It is the processor for manipulating heap DS. It picks the largest child and compares it with the parent key, if parent key is larger than it quite otherwise it swap the parent key with

largest child key.

Q) Create max heap and apply heapify.

8, 20, 9, 4, 15, 10, 7, 23, 3, 12

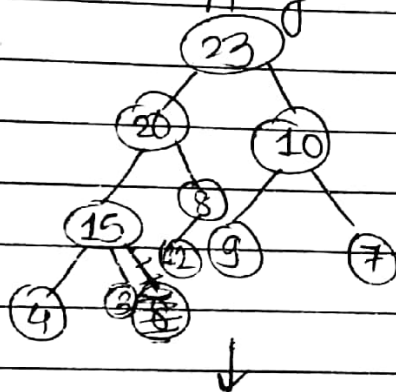
Ans)



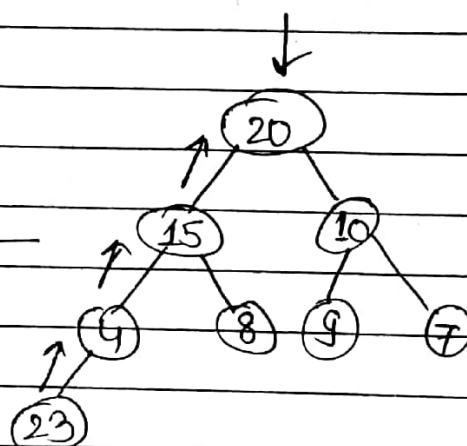
Swapping 20 and 8.

Swapping 15 and 8

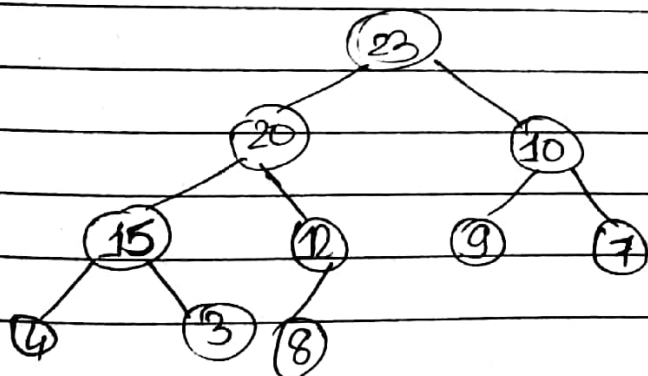
Swapping 8 and 12.



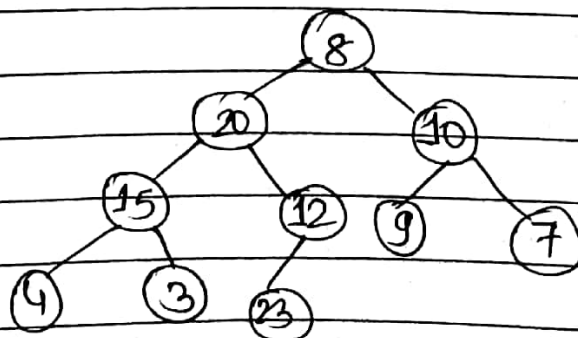
←

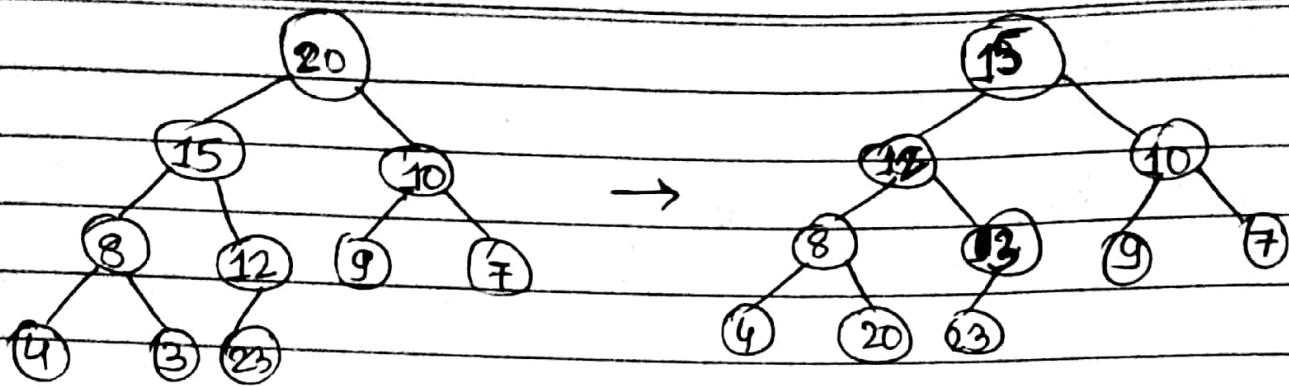


Swapping 23 and 20



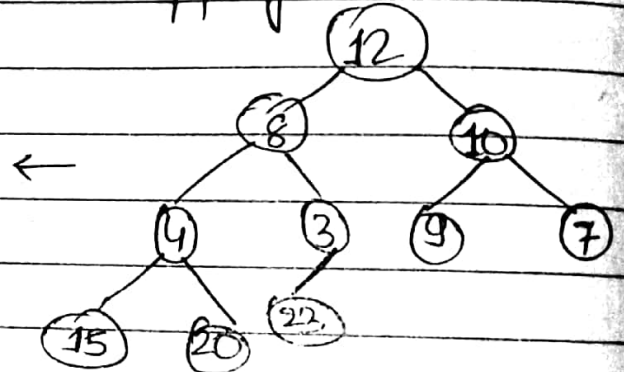
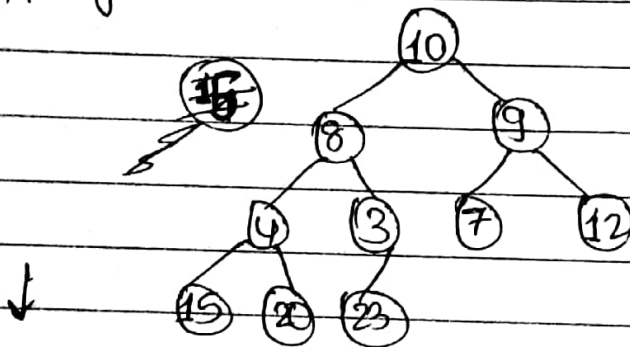
• Create minheap.





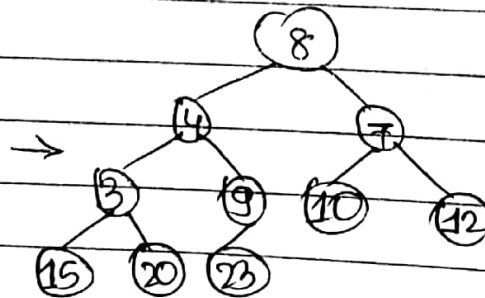
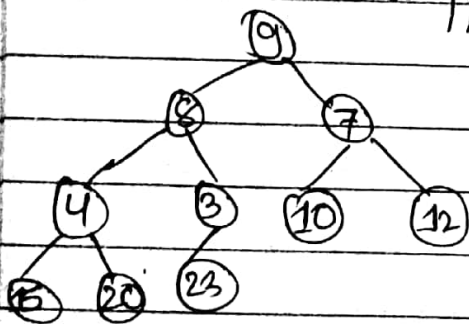
Swapping 3 and 20

Swapping 4 and 15



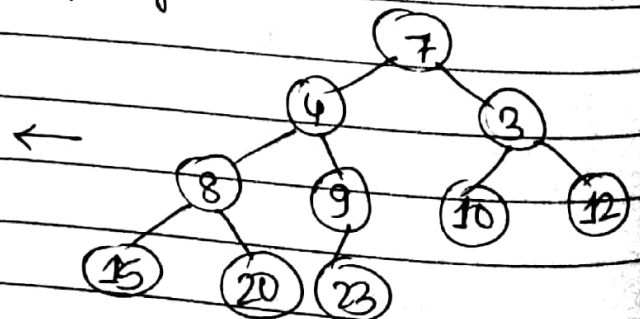
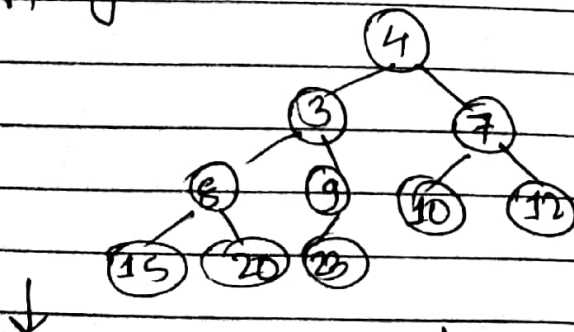
Swapping 10 and 7

Swapping 7 and 12



swapping 3 and 9

Swapping 8 and 3

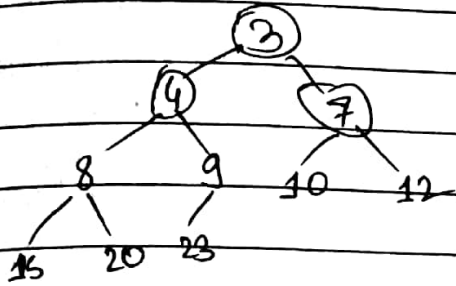


swapping 4 and 3

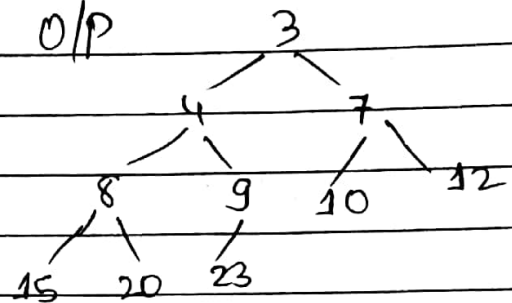
Swapping 7 and 3

25, 6, 87, 22, 47, 10, 35, 7, 63, 52, 41, 2

a) Create min heap and apply heapify.



O/P



3	4	7	8	9	10	12	15	20	23
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