

# Symbol Tables

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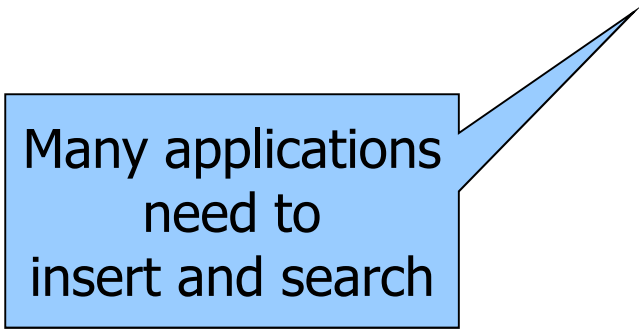
# Symbol Tables

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## ■ Definition

- A **Symbol Table** is a data structure with records including a key and allowing operations such as
  - **Insertion** of a new record
  - **Search** of a record with a given key
  - Delete, select, order, union

- Sometimes symbol tables are denoted with the term **dictionary**



Many applications  
need to  
insert and search



# Applications

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<b>Applications</b>	<b>Target, i.e., searching</b>	<b>Key</b>	<b>Return Value</b>
Dictionary	Definition	Word	Definition
Book index	Relevant pages	Word	Page list
DNS	IP address given its URL	URL	IP address
Invers DNS	URL given its IP address	IP address	URL
File system	File on disk	File name	Disk location
Web search	Web page	Keyword	Page list



# Implementations

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## ■ Linear structures

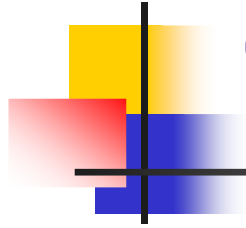
- Direct Access Tables
- Arrays
  - Unordered
  - Ordered
- Lists
  - Unordered
  - Ordered

Already studied  
To be done  
Not analysed

## ■ Tree structures

- Binary Search Trees  
(BSTs)
- Balanced Trees
  - 2-3-4
  - RB-tree
  - B-tree

## ■ Hash Tables



# Complexity: Worst Case

Data Structure	Insert	Search
Direct Access Table	1	1
Unordered Array	1	n
Ordered Array Linear Search	n	n
Ordered Array Binary Search	n	logn
Unordered List	1	n
Ordered List	n	n
BST	n	n
RB-tree	logn	logn
Hashing	1	n



## Complexity: Average Case

Data Structure	Insertion	Search	
		Hit	Miss
Direct Access Table	1	1	1
Unordered Array	1	$n/2$	$n$
Ordered Array Linear Search	$n/2$	$n/2$	$n/2$
Orderer Array Binary Search	$n/2$	$\log n$	$\log n$
Unordered List	1	$n/2$	$n$
Ordered List	$n/2$	$n/2$	$n/2$
BST	$\log n$	$\log n$	$\log n$
RB-tree	$\log n$	$\log n$	$\log n$
Hashing	1	1	1



# Direct Access Tables

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- All search algorithms analysed so far are based on comparisons
- Exception
  - Direct access tables
  - Hash tables
- In direct access tables
  - A key  $k \in U = \{0, 1, \dots, \text{card}(U)-1\}$
  - Is used as the index of the array  $st$ 
    - $st[0, 1, \dots, \text{card}(U)-1]$

The cardinality of  $U$  is small  
The keys are not necessary integer values



# Direct Access Tables

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- Given a universe  $U$  of keys, i.e., each key  $k \in U$ 
  - $\text{maxN}$  is the number of elements in  $U$ , i.e.,  $\text{maxN} = |U|$
- The array `st` is used to store the keys
- Given a key  $k$ 
  - Function `getindex(k)` returns an integer from 0 to  $\text{maxN}-1$ , acting as an array index
  - If the key  $k$  is in the table, its position is `st[getindex(k)]`
  - If the key  $k$  is not in the table `st[getindex(k)]` stores an empty element



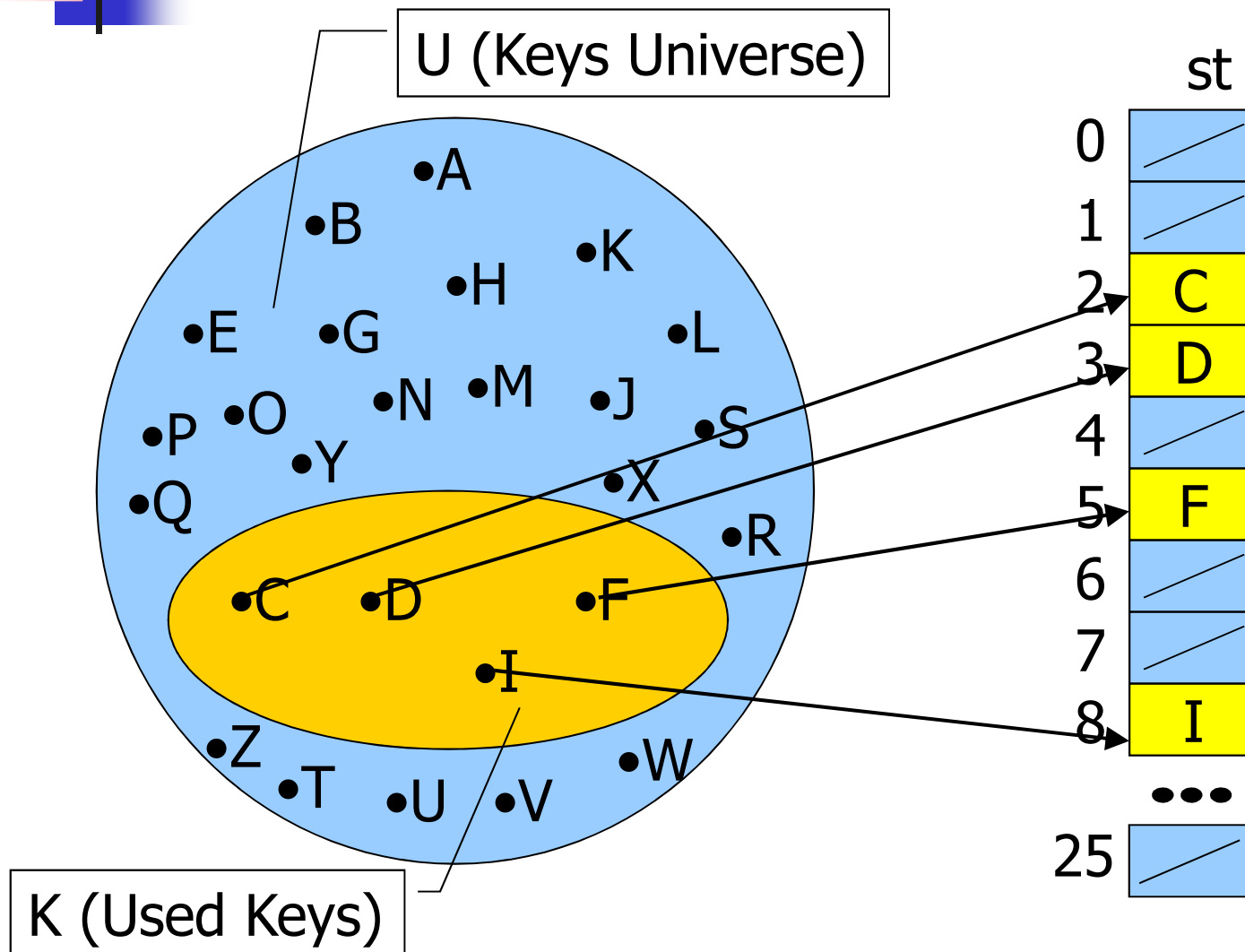


# Direct Access Tables

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- This looks simple enough, but the keys may not be integer values
  - If keys are integers from 0 to maxN-1
    - `getindex(k) = k;`
  - If keys are capital letters in the English alphabet A..Z
    - `getindex(k) = k - ((int) 'A');`
  - For generic keys `getindex` has to map those keys into integer values in the range [0, maxN-1]

# Direct Access Tables





## Pros & Cons

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- Insert, search, and delete complexity
  - $T(n) = \Theta(1)$
- Init complexity
  - $T(n) = \Theta(\max N)$
- Memory usage
  - $S(n) = \Theta(|U|) = \Theta(\max N)$



## Pros & Cons

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- Limits of direct access tables
  - For large  $|U|$  the array st cannot be allocated, i.e., they can be used only for small maxN
  - If  $|K| \ll |U|$  there is a memory loss
- In those two cases direct access tables have to be extended into **hash tables**
- Often used to convert keys into integers (and vice-versa) with a cost equal to 1