

# **Graphs**

#### **Graph Representations**

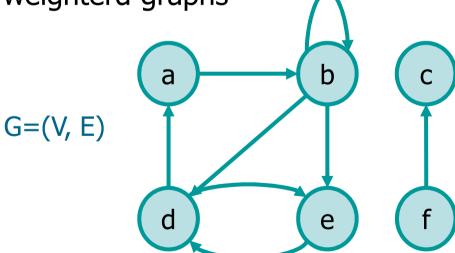
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#### Why graph?

- Many practical applications
- Hundreds of algorithms
- Interesting abstraction usable in various domains
  - Connections
  - Cycles
  - Shortest paths
  - > Etc.
- Active research area in Computer Science and Discrete Mathematics

#### **Representations of graphs**

- Representation of graphs
  - > Adjacency matrix
  - > Adjacency list
- Both of them can be applied to directed, undirected and weighterd graphs



### **Adjacency matrix**

- ❖ Given G = (V, E), its adjacency matrix is
  - > A matrix M of |V| x |V| elements

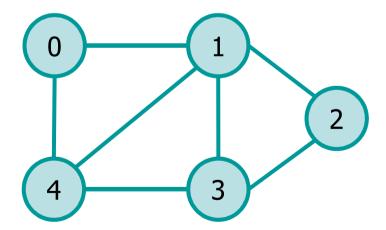
$$M[i, j] = \begin{cases} 1 & \text{if}(i, j) \in E \\ 0 & \text{if}(i, j) \notin E \end{cases}$$

For undirected graphs A is symmetric

# **Example: Undirected graph**

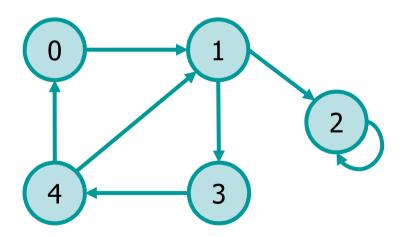
	0	1	2	3	4
0	0	1	0	0	1
1	1	0	1	1	1
2	0	1	0	1	0
3	0	1	1	0	1
4	1	1	0	1	0

Symmetric



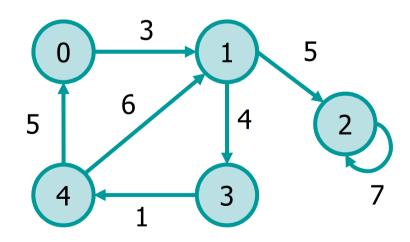
# **Example: Directed graph**

	0	1	2	3	4
0	0	1	0	0	0
1	0	0	1	1	0
2	0	0	1	0	0
3	0	0	0	0	1
4	1	1	0	0	0



# **Example: Weighted Directed graph**

	0	1	2	3	4
0	0	3	0	0	0
1	0	0	5	4	0
2	0	0	7	0	0
3	0	0	0	0	1
4	5	6	0	0	0



- Possible implementations
  - > Static 2D matrix
  - > Dynamic 2D matrix
    - Array of pointers to arrays, i.e., vertex array of structures with dynamic array of vertices
    - Use a struct when it is necessary to store edge/vertex attributes

#### Input file format

nVertex dir/undirected vertex1 vertex2 weight

. . .

#### Example

If  $0 \rightarrow$  undirected graph
If it is not present  $\rightarrow$  directed grph

2 3 4

2 4 5

6 7 1 Unweighted graph have weights equal to 1

```
/* constant declaration */
#define MAX_LINE 100
enum {WHITE, GREY, BLACK};

/* type declarations */
typedef struct vertex graph_t;

/* array (vertices) with rows of adjacency matrix */
struct vertex {
  int id, color, scc;
  int td, tq, dist, pred;
  int *rowAdj;
};

Structure declaration
  with several extra
  attributes
```

```
graph t *graph load(char *filename, int *nv) {
 graph t *g;
 char line[MAX_LINE];
 int i, j, weight, dir;
 FILE *fp;
 fp = util fopen(filename, "r");
 fgets(line, MAX_LINE, fp);
  if (sscanf(line, "%d%d", nv, &dir) != 2) {
   sscanf(line, "%d", nv);
   dir = 1;
 g = (graph t *)util calloc(*nv, sizeof(graph t));
```

```
for (i=0; i<*nv; i++) {
  q[i].id = i;
  g[i].color = WHITE;
  g[i].dist = INT_MAX;
  g[i].pred = g[i].scc = -1;
  g[i].td = g[i].tq = -1;
  g[i].rowAdj = (int *)util calloc(*nv, sizeof(int));
 while (fgets(line, MAX LINE, fp) != NULL) {
  if (sscanf(line, "%d%d%d", &i, &j, &weight) != 3) {
    sscanf(line, "%d%d", &i, &j);
    weight = 1;
  g[i].rowAdj[j] = weight;
  if (dir == 0) g[j].rowAdj[i] = weight;
fclose(fp);
return g;
```

It is often necessary to store a vertex id to matrix index correspondence

```
int graph_find(graph_t *g, int nv, int id){
  int i;
  for (i=0; i<nv; i++) {
    if (g[i].id == id) {
      return i;
                                        It is possible to use a symbol
  return -1;
                                                 table
                                                  4
           st
                idABC
                        idXYZ
                                idFOO
                                        idBAR
```

Free the graph

```
void graph_dispose(graph_t *g, int nv) {
  int i;

for (i=0; i<nv; i++) {
   free(g[i].rowAdj);
  }
  free(g);
}</pre>
```

#### **Pro's and Con's**

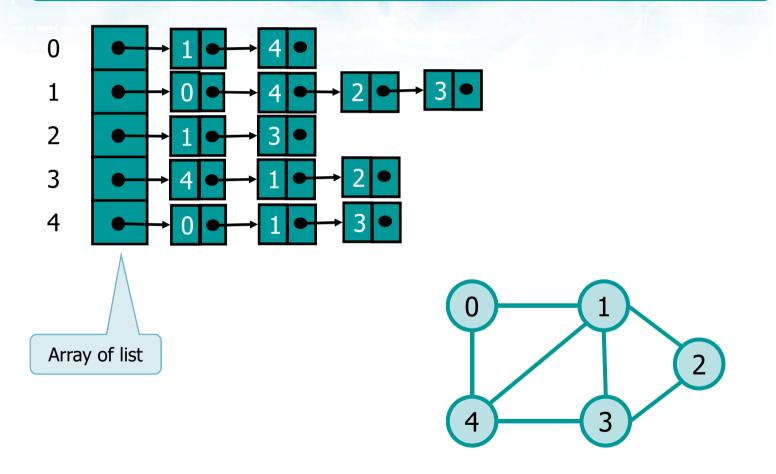
- Space complexity
  - $> S(n) = \Theta(|V|^2)$
  - > It is advantageous
    - For dense graphs, for which |E| is close to |V|<sup>2</sup>
    - When we need to be able to tell quickly if there is a connecting edge between two vertices
- No extra costs for storing the weights in a weighted graph
- Efficient access to graph topology

Boolean versus Integers

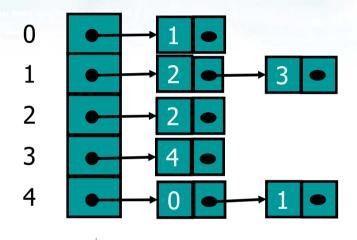
## **Adjacency list**

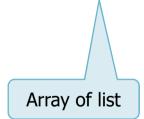
- ❖ Given G = (V, E), its adjacency list is
  - > A main list representing vertices
  - ➤ A secondary list of vertices or edges for each element of the main list
- The list of list may have different implementations
  - > A true list of lists
  - ➤ An array of lists
  - > Etc.

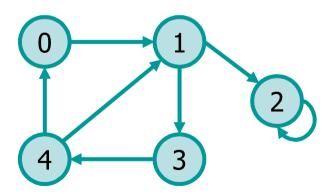
# **Example: Undirected graph**



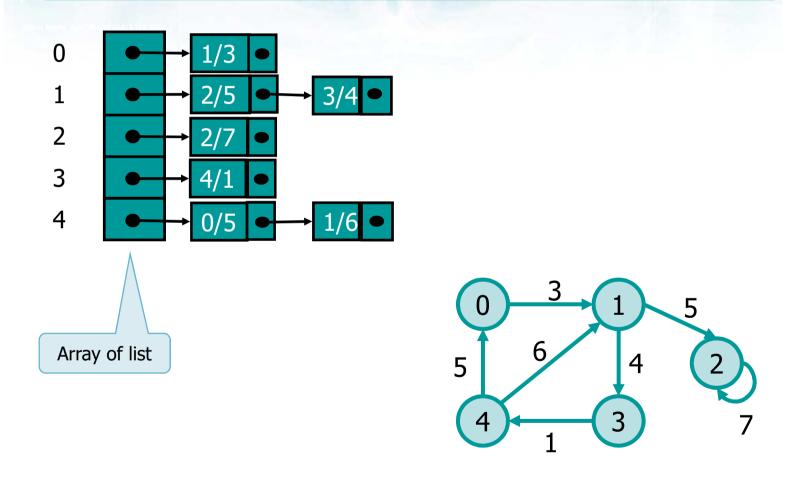
# **Example: Directed graph**







# **Example: Weighted directed graph**



```
typedef struct edge edge t;
typedef struct vertex graph_t;
struct edge {
                                        Edges of the
  int weight;
                                        vertex list of
  struct vertex *dst;
                                          edge list
  struct edge *next;
};
struct vertex {
  int id, color, scc;
  int td, tq, dist;
                                        Vertices of the
  struct vertex *pred;
                                        vertex list of
  struct edge *head;
                                          edge lilst
  struct vertex *next;
};
```

```
graph t *graph load(char *filename, int *nv) {
 graph t *g=NULL;
 char line[MAX LINE];
 int i, j, weight, dir;
 FILE *fp;
 fp = util fopen(filename, "r");
 fgets(line, MAX LINE, fp);
 if (sscanf(line, "%d%d", nv, &dir) != 2) {
   sscanf(line, "%d", nv);
   dir = 1;
 /* create initial structure for vertices */
 for (i=*nv-1; i>=0; i--) {
   g = new_node(g, i);
```

```
/* load edges */
while (fgets(line, MAX LINE, fp) != NULL) {
  if(sscanf(line, "%d%d%d", &i, &j, &weight)!= 3) {
    sscanf(line, "%d%d", &i, &j);
   weight = 1;
 new_edge(g, *nv, i, j, weight);
  if (dir == 0) {
    new_edge(g, *nv, j, i, weight);
fclose(fp);
return g;
```

```
static graph_t *new_node(graph_t *g, int id) {
   graph_t *n;

   n = (graph_t *)util_malloc(sizeof(graph_t));
   n->id = id;
   n->color = WHITE;
   n->dist = INT_MAX;
   n->pred = NULL;
   n->scc = n->td = n->tq = -1;
   n->head = NULL;
   n->next = g;
   return n;
}
```

```
static void new_edge(
  graph_t *g, int nv, int i, int j, int weight) {
  graph_t *src, *dst;
  edge_t *e;

  src = graph_find(g, nv, i);
  dst = graph_find(g, nv, j);

  e = (edge_t *)util_malloc(sizeof(edge_t));
  e->dst = dst;
  e->weight = weight;
  e->next = src->head;
  src->head = e;
}
```

It is often necessary to store a vertex id to matrix index correspondence

```
graph_t *graph_find(graph_t *g, int nv, int id) {
  graph t *n = g;
  while (n != NULL) {
    if (n->id == id) {
      return n;
    n = n->next;
                                        It is possible to use a
                                            symbol table
  return NULL;
                                         3
                                                 4
           st
               idABC
                       idXYZ
                               idFOO
                                       idBAR
```

```
void graph_dispose(graph_t *g, int nv) {
  graph t *n;
  edge_t *e;
  while (g != NULL) {
    n = g;
    while (n->head != NULL) {
      e = n->head;
      n->head = e->next;
      free(e);
    g = n->next;
    free(n);
```

#### **Pro's and con's**

- Undirected graphs
  - $\triangleright$  Total amount of elements in the lists = 2|E|
- Directed graphs
  - > Total amount of elements in the lists = |E|
- Space complexity
  - > S(n) = O(max(|V|, |E|)) = O(|V+E|)
  - ➤ It is advantageous for sparse graphs for which |E| is much less than |V|<sup>2</sup>
- Verifying the existence of edge (u,v) requires scanning the adjacency list of u
- Memory needed to represent weights in weighted graphs