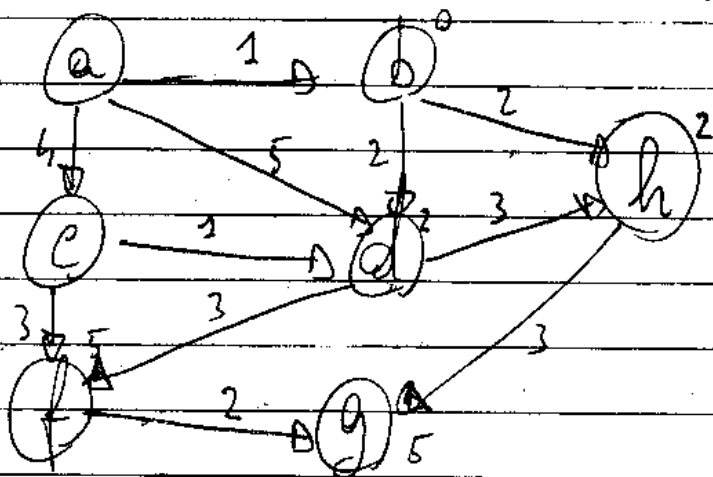
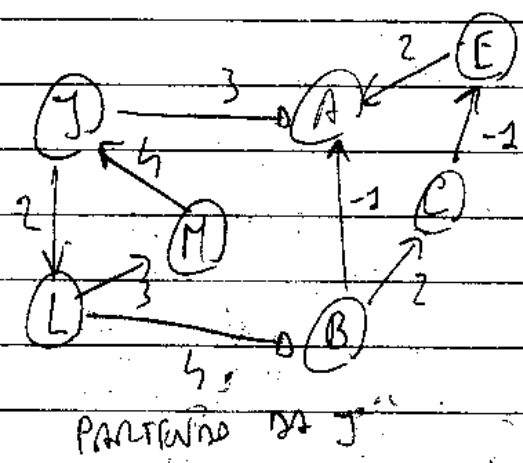


1-12 cammini minimi?

costo A (minimo)
(DISTANZA)



a	b	c	d	e	f	g	h
∞	0	∞	∞	∞	∞	∞	∞
∞	0	∞	2	2	5	5	2
∞	2	2	2	2	5	5	2
∞	2	2	2	2	5	5	2
∞	2	2	2	2	5	5	2
∞	2	2	2	2	5	5	2
∞	2	2	2	2	5	5	2

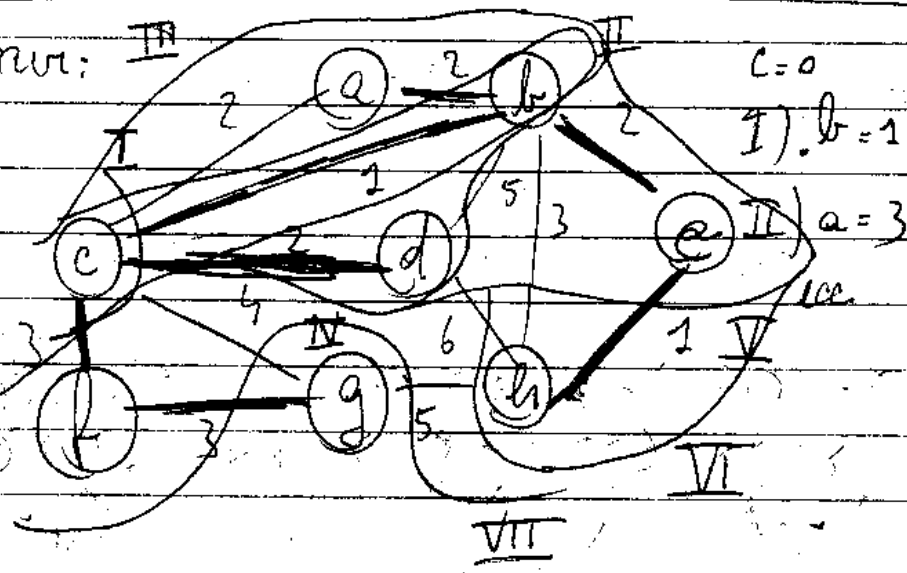


ORDINE LESSI COGRAFICO

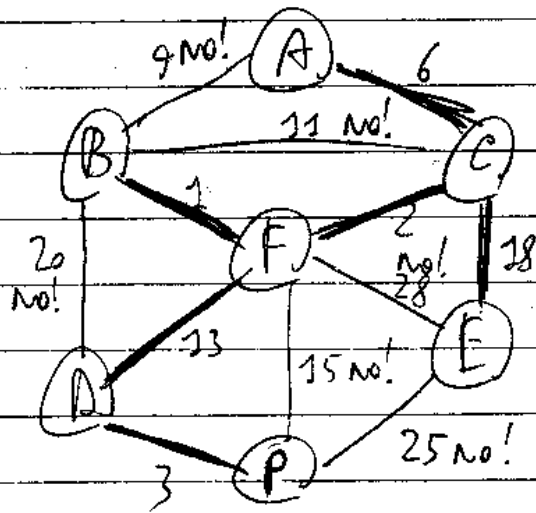
ORDINE	LESSI COGRAFICO	P ₁	P ₂	P ₃
B-A (-1)	A	∞	3	3
B-C (2)	B	∞	6	6
C-E (-1)	C	∞	8	8
E-A (2)	E	∞	7	7
J-A (3)	J	0	0	0
J-L (2)	L	∞	2	2
J-M (2)	M	∞	5	5
L-B (4)				
L-M (3)				
M-J (4)				

(BELLMAN)

ALGORITMO DI FLOYD: III



h.w.s.h.a.l



si possono anche
in peso crescente
in modo che non
introducano cicli

peso = 43

TABELLE DI HASHING: dimensioni "CANICOLA" ordine alfabetico

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16
A B C D E F G H I J K L M N O P

Q R S T U V W X Y Z

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
A	A	C	C									L	L	N	O	
1	2	3	2													

open addressing
con linear probing

$$h(h) = (k \% M + i) \% M$$

TENTATIVO

"APPELLO LAUREANDI" dim. 23 double hashing

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22
A	L				E					P		L			O	P				L		

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18
A B C D E F G H I J K L M N O P Q R
W X Y Z

19 20 21 22
S T U V

$$h(h) = [h_1(h) + i \cdot h_2(h)] \% M$$

$$h_1(h) = h \% M \quad h_2(h) = 1 + k \% C$$

$$h^1(P) = [16 \% 23 + 1(1 + 16 \% 19)] \% 23 = 10$$

$$h^1(L) = [12 \% 23 + 1(1 + 12 \% 19)] \% 23 = 2$$

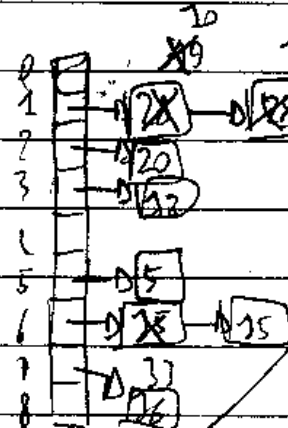
$$h^2(L) = [12 \% 23 + 2(1 + 12 \% 19)] \% 23 = [12 + 26] \% 23 = 15$$

$$h^3(L) = [12 \% 23 + 3(1 + 12 \% 19)] \% 23 = 53 \% 23 = 5 \quad h^4(L) = 20$$

19
15
23

5 28 19 15 20 33 12 16 10

slim 9 linear probing

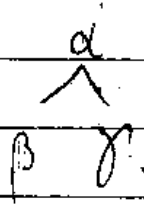


$h(h) = h \cdot 11$

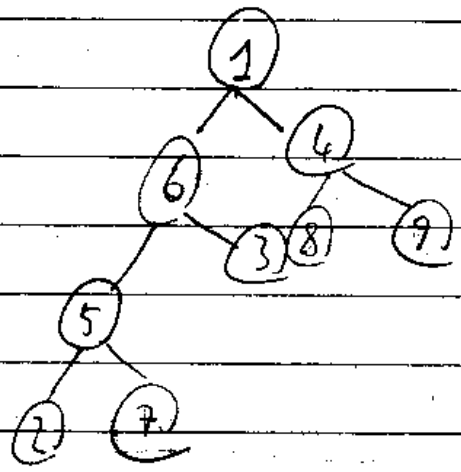
VE COLLISION SI GESTIONS
CON UNO LISTA E IN INSE-
RTION TESTA

VISITE ALBERI:

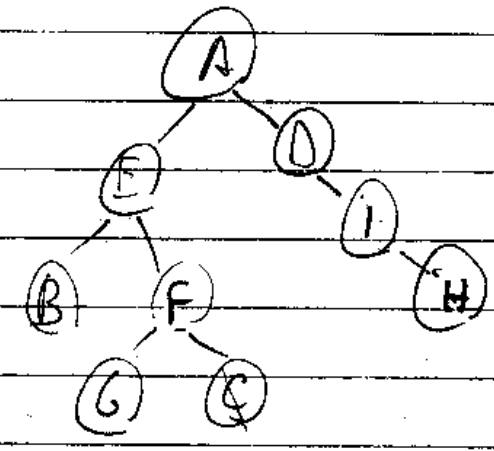
PRE-ORDER: 1 6 5 2 7 3 4 8 9
POST-ORDER: 2 7 5 3 6 8 9 4 1
IN-ORDER: 2 5 7 6 3 1 8 4 9



PRE: α β γ
IN: β α γ
POST: β γ α



PRE A E B F G C D I H
IN: B E G F C A D H I
POST B G C F E H I D A



unfolding:

$T(n) = 2T(\frac{n}{3}) + \frac{n}{3}$ $T(1) = 1$

$T(\frac{n}{3}) = 2T(\frac{n}{9}) + \frac{n}{9}$

$T(\frac{n}{9}) = 2T(\frac{n}{27}) + \frac{n}{27}$

$T(n) = 2(2T(\frac{n}{9}) + \frac{n}{9}) + \frac{n}{3} =$

$= 2[2(2T(\frac{n}{27}) + \frac{n}{27}) + \frac{n}{9}] + \frac{n}{3}$

$= 8T(\frac{n}{27}) + \frac{6}{27}n + \frac{2}{9}n + \frac{n}{3} = \sum_{i=0}^{\log_3 n} \frac{2^i}{3^i} n$

M: $\frac{n}{3^i} = 1$

$i = \log_3 n = M$

$\sum_{i=0}^M x^i < \frac{1-x^{M+1}}{1-x} < \frac{1}{1-x} = \text{Cost.}$
 $\frac{1-x^{M+1}}{1-x} < \frac{1}{1-x}$
 $\frac{x^{M+1}}{1-x} < \frac{1}{1-x}$
 $x^{M+1} < 1$
 $x < 1$

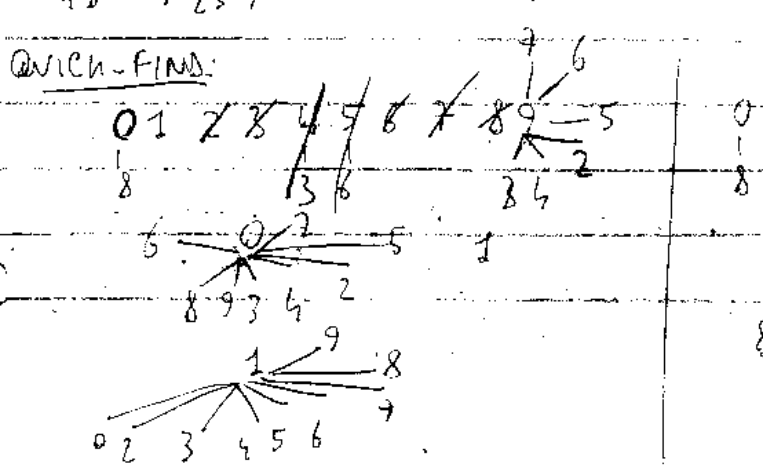
$T(n) = O(n)$
 $\sum_{i=0}^{\log_2 n} \frac{n^2}{4^i} = n^2 \sum_{i=0}^{\log_2 n} \frac{1}{4^i}$
 $\sum_{i=0}^{\log_2 n} \frac{1}{4^i} < \sum_{i=0}^{\infty} \frac{1}{4^i} = \frac{1}{1-\frac{1}{4}} = \frac{4}{3}$
 $T(n) < \frac{4}{3} n^2$

$T(n) = 4T\left(\frac{n}{2}\right) + n^2$
 $T\left(\frac{n}{2}\right) = 4T\left(\frac{n}{4}\right) + \frac{n^2}{4}$
 $T\left(\frac{n}{4}\right) = 4T\left(\frac{n}{8}\right) + \frac{n^2}{16}$
 \dots
 $T\left(\frac{n}{2^M}\right) = 4T\left(\frac{n}{2^{M+1}}\right) + \frac{n^2}{4^M}$
 $M = \log_2 n$
 $T(n) = n^2 \sum_{i=0}^{\log_2 n} 1 = O(n^2 \log_2 n)$

$T(n) = 4\left[4\left(4T\left(\frac{n}{8}\right) + \frac{n^2}{16}\right) + \frac{n^2}{4}\right] + n^2$
 $= 64T\left(\frac{n}{8}\right) + n^2 + n^2 + n^2$
 $= 3n^2 + 64T\left(\frac{n}{8}\right)$
 $M = \frac{\log_2 n}{2} = 1, M = 1 = \log_2 n$

ONLINE CONNECTIVITY: QUICK FIND

	0	1	2	3	4	5	6	7	8	9
(4,8) →	0	1	2	3	8	5	6	7	8	9
(7,3) →	0	1	2	3	8	5	6	3	8	9
(5,4) →	0	1	2	3	8	9	6	3	8	9
(2,9) →	0	1	9	3	8	9	6	3	8	9
(5,6) →	0	1	9	3	8	9	6	3	8	9
(8,0) →	0	1	9	3	8	9	6	3	8	9



QUICK-UNION

	0	1	2	3	4	5	6	7	8	9
(4,8) →	0	1	2	3	8	5	6	7	8	9
(7,3) →	0	1	2	3	8	5	6	3	8	9
(5,9) →	0	1	2	3	8	9	6	3	8	9
(9,4) →	0	1	2	3	8	9	6	3	8	9
(5,6) →	0	1	2	3	8	9	6	3	8	9
(8,0) →	0	1	2	3	8	9	6	3	8	9

