

P1. a. 通过指定-m参数调节数据包的ttl值使其能再网络中的生存周期延长来发现更多的跳转，多次实验发现centos官网的跳转数较多为23跳

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tqym@ubuntu:~$ traceroute centos.org
traceroute to centos.org (81.171.33.202), 30 hops max, 60 byte packets
 1 _gateway (192.168.1.1) 4.371 ms 4.257 ms 4.201 ms
 2 100.100.0.1 (100.100.0.1) 5.564 ms 5.495 ms 5.449 ms
 3 112.5.175.53 (112.5.175.53) 11.129 ms 5.353 ms 11.032 ms
 4 112.50.220.65 (112.50.220.65) 9.813 ms * 6.401 ms
 5 111.24.11.61 (111.24.11.61) 9.651 ms 9.606 ms 9.555 ms
 6 221.183.98.70 (221.183.98.70) 21.935 ms 111.24.5.89 (111.24.5.89) 31.047 ms 221.183.98.66 (221.183.98.66) 30.948 ms
 7 111.24.5.194 (111.24.5.194) 37.841 ms 111.24.5.182 (111.24.5.182) 27.491 ms 111.24.14.150 (111.24.14.150) 27.177 ms
 8 221.176.22.158 (221.176.22.158) 48.679 ms 221.183.68.145 (221.183.68.145) 26.705 ms 221.176.24.6 (221.176.24.6) 26.415 ms
 9 221.176.19.42 (221.176.19.42) 35.127 ms 221.176.20.2 (221.176.20.2) 25.953 ms 221.183.25.117 (221.183.25.117) 34.671 ms
10 221.183.55.81 (221.183.55.81) 47.169 ms 221.183.55.57 (221.183.55.57) 197.350 ms 221.183.55.53 (221.183.55.53) 46.692 ms
11 223.120.15.109 (223.120.15.109) 261.888 ms 223.120.15.101 (223.120.15.101) 262.394 ms 223.120.15.57 (223.120.15.57) 250.160 ms
12 223.120.10.198 (223.120.10.198) 254.830 ms 242.380 ms 248.459 ms
13 149.14.199.193 (149.14.199.193) 341.483 ms 341.410 ms *
14 be3672.ccr52.lhr01.atlas.cogentco.com (130.117.48.145) 348.646 ms * be3671.ccr51.lhr01.atlas.cogentco.com (130.117.48.137) 354.149 ms
15 be3488.ccr42.lon13.atlas.cogentco.com (154.54.60.13) 341.138 ms 341.090 ms be3487.ccr41.lon13.atlas.cogentco.com (154.54.60.5) 341.036 ms
16 be2870.ccr22.lon01.atlas.cogentco.com (154.54.58.174) 329.902 ms be2868.ccr21.lon01.atlas.cogentco.com (154.54.57.154) 325.038 ms be2871.ccr21.lon01.atlas.cogentco.com (154.54.58.186) 324.612 ms
17 151.139.40.3 (151.139.40.3) 291.300 ms 151.139.40.9 (151.139.40.9) 279.656 ms 282.181 ms
18 151.139.40.69 (151.139.40.69) 280.204 ms 151.139.40.67 (151.139.40.67) 274.183 ms 151.139.40.69 (151.139.40.69) 281.260 ms
19 * * *
20 151.139.80.3 (151.139.80.3) 279.399 ms 279.345 ms 151.139.80.6 (151.139.80.6) 279.282 ms
21 151.139.82.7 (151.139.82.7) 279.989 ms 284.387 ms 151.139.82.9 (151.139.82.9) 283.592 ms
22 * * *
23 ip-81.171.33.202.centos.org (81.171.33.202) 285.672 ms !X 280.118 ms !X 282.120 ms !X
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b.通过尝试不同的服务器，从本地到116.203.240.158主机所经过的ISP最多，为5个

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tqym@ubuntu:~$ traceroute 116.203.240.158 -A
traceroute to 116.203.240.158 (116.203.240.158), 30 hops max, 60 byte packets
 1 _gateway (192.168.0.1) [*] 37.741 ms 37.610 ms 37.545 ms
 2 1.176.30.117.broad.xm.fj.dynamic.163data.com.cn (117.30.176.1) [AS140330] 37.490 ms 37.433 ms 37.356 ms ← 1
 3 117.30.25.41 (117.30.25.41) [AS140330] 40.220 ms 117.30.25.45 (117.30.25.45) [AS140330] 37.239 ms 37.227 ms
 4 117.30.26.117 (117.30.26.117) [AS140330] 37.128 ms 61.154.236.9 (61.154.236.9) [AS140330] 39.996 ms 39.941 ms
 5 202.97.103.9 (202.97.103.9) [AS4134] 43.144 ms 202.97.40.165 (202.97.40.165) [AS4134] 43.090 ms * ← 2
 6 202.97.24.254 (202.97.24.254) [AS4134] 42.980 ms * 202.97.57.145 (202.97.57.145) [AS4134] 28.013 ms
 7 202.97.94.114 (202.97.94.114) [AS4134] 27.927 ms 202.97.94.98 (202.97.94.98) [AS4134] 27.845 ms 202.97.83.62 (202.97.83.62) [AS4134] 43.989 ms
 8 202.97.13.26 (202.97.13.26) [AS4134] 228.360 ms 202.97.89.110 (202.97.89.110) [AS4134] 228.285 ms 202.97.95.206 (202.97.95.206) [AS4134] 217.460 ms
 9 ae77.edge3.London15.Level3.net (195.50.126.217) [AS9057/AS3356] 269.608 ms 329.576 ms * ← 3
10 * ae-2-3204.edge7.Amsterdam1.Level3.net (4.69.162.181) [AS3356] 248.082 ms *
11 AS33891-NET.edge7.Amsterdam1.Level3.net (212.72.41.234) [AS9057/AS3356] 227.828 ms 227.768 ms 225.277 ms
12 ae6-2011.nbg40.core-backbone.com (80.255.14.246) [AS201011/AS33891] 215.266 ms 215.215 ms 225.111 ms ← 4
13 core-backbone.hetzner.com (5.56.20.254) [AS33891] 234.984 ms 205.714 ms core-backbone.hetzner.com (81.95.15.6) [AS33891] 244.702 ms
14 213-239-245-73.clients.your-server.de (213.239.245.73) [AS24940] 234.819 ms core12.nbg1.hetzner.com (213.239.229.165) [AS24940] 244.573 ms 233.765 ms ← 5
15 static.85-10-228-86.clients.your-server.de (85.10.228.86) [AS24940] 243.401 ms 243.145 ms spine1.cloud1.nbg1.hetzner.com (85.10.250.210) [AS24940] 253.428 ms
16 * * *
17 12546.your-cloud.host (88.99.159.151) [AS24940] 222.290 ms 222.232 ms 229.099 ms
18 static.158.240.203.116.clients.your-server.de (116.203.240.158) [AS24940] 209.025 ms 222.027 ms 221.970 ms
tqym@ubuntu:~$
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P2 XMU 平均每分钟访问速度 $\lambda = \frac{60000}{24 \times 60} = \frac{1000}{24} = \frac{125}{3} / \text{min}$

平均处理速度 $\mu = \frac{1}{3} / \text{min} \quad \therefore \frac{\lambda}{\mu} = 125 = \rho$

假设需要 k 个电话，其阻塞率

$$P_R = \frac{1}{(k+1)!} \rho^k \frac{1}{1 + \frac{1}{\rho} + \dots + \frac{1}{\rho^k}} = 1\%$$

\therefore 当 $k \geq 14$ 时, $P_R < 1\%$ \therefore 需要 14 条外部电话线

P3 单核处理速度 $v = \frac{1}{200ms} = 5/s$ \therefore 排队延迟 $w = S \frac{\rho}{1-\rho} = \frac{3}{20} s$

四核处理速度 $\mu = 4 \times v = 20/s$

web请求速度 $\lambda = 15/s$

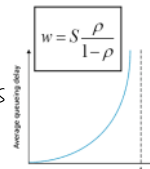
$\therefore \rho = \frac{\lambda}{\mu} = \frac{3}{4}$ $S = \frac{1}{\mu} = \frac{1}{20}$

传输延迟 $S = \frac{1}{\mu} = \frac{1}{20}$

\therefore 总服务时间为 $w + S = \frac{4}{20} = \frac{1}{5} = 0.2s$

当核心数为2时, $\mu = 2v = 10/s$

$\rho = \frac{\lambda}{\mu} = 1.5/s$, 可知平均排队延迟为无穷大



P4 a. $d_{prop} = \frac{m}{s}$

b. $d_{trans} = \frac{L}{R}$

c. end-to-end delay $= d_{prop} + d_{trans} = \frac{m}{s} + \frac{L}{R}$

d. Just left host A

e. the single link between A and B

f. Host B

g. let $\frac{m}{s} = \frac{L}{R} \Rightarrow m = \frac{L}{R} \cdot s = \frac{120}{56 \times 1000} \times 2.5 \times 10^8 = 5.357 \times 10^5 m$

P5 a. $R \cdot d_{prop} = 2 \times 1000^1 \times \frac{20000 \times 10^3}{2.5 \times 10^8} = 1.6 \times 10^5 \text{ bits}$

b. propagation delay $= \frac{20000 \times 10^3}{2.5 \times 10^8} = 0.08 s$

transmission delay $= \frac{800000}{2 \times 1000^1} = 0.4 s$

so the maximum number of bits is

$R \cdot d_{prop} = 1.6 \times 10^5 \text{ bits}$

c. the bandwidth-delay product is a measurement of how many bits can fill up a network link

d. the width is $\frac{m}{R \cdot \frac{m}{s}} = \frac{s}{R} = \frac{2.5 \times 10^8}{2 \times 1000^2} = 125 m$

it is longer than a football field (100m)

e. the width of a bit is $\frac{m}{R \cdot d_{prop}} = \frac{s}{R}$

P6 the processing delay $d_{proc} = \frac{56 \times 8}{64 \times 1000} = 7 ms$

the transmission delay $d_{trans} = \frac{L}{R} = \frac{56 \times 8}{2 \times 1000} = 0.224 ms$

the propagation delay $d_{prop} = 10 ms$

the whole time is $17.224 ms$

P7 The required transmission delay is $d_{trans} = \frac{40 \times 1000^6 \times 8}{100 \times 1000^2} = 3.2 \times 10^6 s \approx 37 \text{ days}$

(use standard SI prefixes)

\therefore using FedEx over-night delivery is faster. choose FedEx

P8 a. A circuit-switch network is more appropriate for this application.

Because it will continue running for a relatively long period of time and sending data at a steady rate. Circuit-switch can make the transmission quick, efficient and steady. Relatively speaking, packet loss and blocking will occur in packet-switch network.

b. No, the sum of the application data is less than the capacities of each and every link.

P9 a. $\frac{3 \text{ Mbps}}{150 \text{ kbps}} = 20$

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b. 10%

c. $C_{120}^n p^n (1-p)^{120-n}, p = 0.1$
 $C_{110}^n 0.1^n \times 0.9^{120-n}$

d. $1 - \sum_{n=0}^{20} C_{120}^n 0.1^n \times 0.9^{120-n}$

P10 a. The time required for the first packet switch is $\frac{L}{R} = \frac{8 \times 10^4}{2 \times 10^6} = 4 \text{ s}$
total time is $4 \times 3 = 12 \text{ s}$

b. The time required for the first packet switch is $\frac{10000}{2 \times 10^6} = 5 \text{ ms}$
After $2 \times 5 = 10 \text{ ms}$, the first switch fully receives the second packet.

c. When the first packet arrives at destination, the time spent is $5 \times 3 = 15 \text{ ms}$
After that, a packet will be received every 5 seconds. So the total time is
 $15 \text{ ms} + 799 \times 5 \text{ ms} = 4.01 \text{ s}$, which is faster than sending without message segmentation.

d. Facilitate error detection and retransmission;
Non-segmented large packets are easy to make the router cache insufficient, resulting in packet loss.

e. Sorting is required for message segmentation
The header needs to be added.

P11 The number of groups is $\frac{F}{S}$

The transmission delay is $\frac{S+80}{R}$ each link

The delay of transmitting the first packet to host B is $3 \times \frac{S+80}{R}$. After that
a packet will be received every $\frac{S+80}{R}$ seconds,
So the total delay is:

$$\text{delay} = (S+80) \times 3 + (S+80) \times \left(\frac{F}{S} - 1\right) = (S+80) \times \left(\frac{F}{S} + 2\right)$$

Derived from the above:

$$\frac{d}{dS} \text{delay} = 0$$

The solution is $S = \sqrt{40F}$

P12 电路交换电话网络和互联网两个网络在“网关”处连接在一起。当一个“Skype”用户(连接到互联网的)打一个电话给普通用户时,通过电路交换网络,在电话用户和网关之间建立了电路。由多个组构成的“Skype”的语音通过互联网发送到网关。在网关那里,语音信号被重整处理,然后发送到电路交换网络。另一方面,语音信号通过电路交换网络发送到网关,网关对语音信号进行重组处理,再发送给“Skype”用户。