Computer Networks and Network Security: Lab Assignment One

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p1.a

通过traceroute各个网站,观察跳转数量。

①www.baidu.com

```
hadoop@ubuntu:/etc$ traceroute www.baidu.com
traceroute to www.baidu.com (36.152.44.96), 30 hops max, 60 byte packets

1 * * *
2 172.31.10.33 (172.31.10.33) 56.041 ms 55.989 ms 84.587 ms

3 * * *
4 210.34.2.30 (210.34.2.30) 84.359 ms 105.453 ms 84.243 ms

5 * * *
6 * * *
7 221.183.77.169 (221.183.77.169) 21.054 ms 221.183.57.1 (221.183.57.1) 20.996 ms 221.183.77.169 (221.183.77.169) 2

0.553 ms

8 221.183.42.137 (221.183.42.137) 30.770 ms * 221.183.42.69 (221.183.42.69) 39.527 ms

9 221.183.59.54 (221.183.59.54) 33.261 ms * *
10 * * *
11 * * 182.61.216.72 (182.61.216.72) 91.631 ms
```

②m.youtube.com

③zh.m.wikipedia.org

在经过对多个网站traceroute之后,发现m.youtube.com的跃点最多,近似认为为23个

p1.b

通过尝试不同的服务器,可得到从本地到116.200.250.150经过了5个ISP

```
hadoop@ubuntu:/etc$ traceroute 116.200.250.150 -A
traceroute to 116.200.250.150 (116.200.250.150), 30 hops max, 60 byte packets

1 ***
2 172.31.10.33 (172.31.10.33) [*] 1163.381 ms 1170.041 ms 1178.602 ms

3 ***
4 210.34.2.30 (210.34.2.30) [A54538] 1248.433 ms 1252.781 ms 1267.579 ms

5 ***
6 183.250.112.21 (183.250.112 21) [A54538] 1248.433 ms 1252.781 ms 1267.579 ms

5 ***
7 112.50.220.69 (112.50.220.69) [A59808] 1260.206 ms 1088.976 ms 1084.948 ms
7 112.50.220.69 (112.50.220.69) [A59808] 196.097 ms 991.033 ms 2018.290 ms

9 111.24.11.61 (111.24.5.186) [11.24.5.186) [A59808] 296.097 ms 991.033 ms 2018.290 ms

10 11.12.4.5.186 (111.24.5.186) [A59808] 2248.665 ms 112.4.5.170 (112.45.5.170) [A59808] 2249.184 ms

11 221.176.22.158 (221.176.22.150 [A59808] 2348.665 ms 112.4.5.170 (112.45.5.170) [A59808] 2348.687 ms 111.24.5.194 (111.24.5.194) [A59808] 2249.184 ms

12 221.182.5.117 (21.183.25.117 (21.183.25.117 [A59808]) 1843.053 ms 1016.322 ms 221.183.68.145 (221.183.68.145) [A59808/A59394] 1853.203 ms

12 221.183.55.33 (221.183.55.33) [A59808/A59394] 1660.352 ms 221.176.19.210 (221.176.19.210) [A59808/A59394] 854.684 ms 221.176.19.214 (221.176.19.214) [A59808] 2249.184 ms

13 221.183.55.33 (221.183.55.33) [A59808] 393.53 ms 123.120.2.181 (223.120.2.181) [A55808/A59394] 854.684 ms 221.176.19.214 (221.176.19.214) [A59808] 393.53 ms 223.120.2.181 (223.120.2.246) [A55808] 90.673 ms 223.120.2.2177 (223.120.2.2177 [A558453] 905.676 ms

14 223.120.2.210 (223.121.2.24) [A558653] 90.404 47 ms 223.120.2.246) [A558453] 90.673 ms 223.120.2.120.2.240 [A558653] 90.404 47 ms 223.120.2.246) [A558453] 1490.193 ms

16 223.121.2.42 (223.121.2.42) [A55863] 90.404 47 ms 223.120.2.246 [A558453] 90.404 47 ms 223.120.2.246) [A558453] 1490.193 ms

17 112.174.80.98 (112.174.80.98) [A54766] 1507.152 ms 112.174.80.113 (112.174.80.113) [A54766] 1503.845 ms 112.190.29.177 (112.190.29.177) [A54766] 1020.57

18 112.174.80.93 (112.174.80.93) [A54766] 1004.385 ms 112.190.29.169 (112.190.29.169) [A54766] 1012.977 ms 112.190.29.177 (112.190.29.177)
```



M= 支/min ニP= ☆=125 個格子 Pk= 1k+1)!PK 1+ 〒P+…+ 中PK = 1%

ア $\lambda = 15/5$ $u = \frac{1}{200m5} = 5/5$ 四枚 M' = 20/5 $\rho = \frac{2}{10} = \frac{2}{10}$ $u = \frac{2}{200} = \frac{2}{10}$ $u = \frac{2}{20}$ 5 ... 排队企足 $w = 5\frac{1}{10} = \frac{2}{20}$ 5

· 总服务时间为 W+5 = 告 = 号 = 0.25 当核心为工时从"=10份户"= 六= 一片= 至7 南知平约排队延迟一口

P4. a. dprap = 15

C. end-to-end to delay = dprop + d trans = $\frac{m}{5} + \frac{L}{R}$

d it just left me host A

e. the first bit is in the single link between A and B

f. the first bit has arrived at the host 13

$$P_{5} = a. R \cdot d_{prop} = 2 \times 10^{3} \times \frac{20000 \times 10^{3}}{2.5 \times 10^{8}} = 1.6 \times 10^{5}$$

$$b. d_{trons} = \frac{1}{R} = \frac{800000}{2 \times (10^{3})^{2}} = 0.45$$

C. the bandwith-delay product is a to show how many bits the link can contain.

d. the width is
$$\frac{m}{R \cdot d_{prop}} = \frac{s}{R} = \frac{2.5 \times 10^8}{2 \times 1000^2} = 125 m$$
It's longer than a football field.

e. the width of α bit is $\frac{m}{R \cdot dprop} = \frac{S}{R}$

$$d_{proc} = \frac{56 \times 8}{64 \times 10^3} = 7m\varsigma \qquad d_{prop}^{trans} = \frac{64 \times 10^3}{2 \times (10^3)^3} = \frac{32}{10^3} = 32 \, \text{ms}$$

so the whole time is 49 ms

$$P_7$$
 if we use the link, it will take $t = \frac{40 \times (1000)^4 \times 8}{100 \times (1000)^2} = 3.2 \times 10^6 \text{s}$ ≈ 37 days.

So, chose Fed Ex my be a good choice.

- Ps. a. circuit swiched network may be more appropriate for this application because when the app is starts, it will continue running for a long period of time and the app will start in a small rate, It will cause packet hose and congestion if we use packet switch network
 - b. No, because the sum of the application data rostes is less than the capacities of each and every link.

$$P_9 \cdot a \cdot \frac{3 \times (10^3)^2}{150 \times 10^3} = 20$$

- b. 10%
- C. Go 0.1 0.9 1-n

$$d. \frac{p-1-C_{120}}{C_{120}} = \frac{1-\sum_{n=0}^{20} C_{120}}{C_{120}} = \frac{1-\sum_{n=0}^{20} C_{120}}{C_{120}}} = \frac{1-\sum_{n=0}^{20} C_{120}}{C_{120}}} = \frac{1-\sum_{n=0}^{20$$

$$P_{10}$$
 a. the total time is $3 \times \frac{8 \times 10^6}{2 \times (0^3)^3} = 125$

b. the first padet use
$$\frac{10000}{2 \times (0^3)^2} = 5 \text{ m/s}$$

the second packet use 2x5 = 10 ms

C. the total time is $5\times3+799\times5=4.015<80125$ it's faster than sending without message segmentation



d. it will also better in error detection and no much cache is needed

e. We should take another time to do message segmentation and recovery

P11.
$$n = \frac{F}{5}$$
 $d_{tilons} = \frac{5+60}{R}$ $d_{tilons} = \frac{5+60}{R}$ $d_{tilons} = \frac{5+60}{R} + \frac{(5+80)}{R} \times (\frac{F}{5}-1)$

$$= \frac{5+80}{R} (\frac{F}{5}+2)$$

$$= \frac{F}{R} + \frac{160}{R} + \frac{80F}{R5} + \frac{25}{R}$$

$$= \frac{160}{R} \times \frac{160$$

when $\frac{80F}{R5} = \frac{25}{R} = > 5 = \overline{J_{40}F}$, we can find the value

Pir. 当一个幼児用户打一个电话编新通用户时,Internet 电话网关将语言编码、压缩、打包、经路由器在印网络中传输。接收为路由器接到数据包后,根据印地收入数据传统接收为网关,该网关、将数据新包、解码、解压缩陷后传统普通电流,们电话何约即用户本传输过程亦是如此。