

Automatic Synthesis of Low Complexity Translation Operators for the Fast Multipole Method

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- Quick introduction to Taylor series based Fast Multipole Method
- Compressed Taylor Series based expansions and translations
- Results - accuracy and time complexity

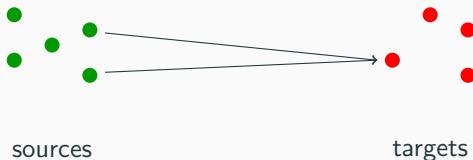
N-body problem

Let $(\mathbf{s}_j)_{j=1}^n$ be sources and $(\mathbf{t}_i)_{i=1}^n$ be targets. Potential at target \mathbf{t}_i is the sum of all potentials from the sources \mathbf{s}_j given by,

$$\sum_j \psi(\mathbf{t}_i, \mathbf{s}_j).$$

For example,

$$\psi(\mathbf{t}_i, \mathbf{s}_j) = \frac{1}{\text{dist}(\mathbf{t}_i, \mathbf{s}_j)}.$$



n sources and n targets $\implies \mathcal{O}(n^2)$ cost.

Fast Multipole Method

Algorithm by Greengard and Rokhlin (1987) to compute the potentials in $\mathcal{O}(n)$ time.

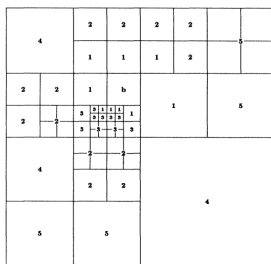


Figure 1: Carrier et al, 1988

Useful for solving PDEs with Integral equation methods.

$$\int G(x-y)\sigma_y dy.$$

Taylor Series based FMM

Local expansion:

$$\psi(\mathbf{t}, \mathbf{s}) = \sum_{|\mathbf{m}| \leq k} \underbrace{\frac{D_{\mathbf{t}}^{\mathbf{m}} \psi(\mathbf{t}, \mathbf{s}) \big|_{\mathbf{t}=\mathbf{c}}}{m!}}_{\text{depends on src/ctr}} \underbrace{(\mathbf{t} - \mathbf{c})^{\mathbf{m}}}_{\text{depends on tgt/ctr}}$$

Multipole expansion:

$$\psi(\mathbf{t}, \mathbf{s}) = \sum_{|\mathbf{m}| \leq k} \underbrace{\frac{D_{\mathbf{s}}^{\mathbf{m}} \psi(\mathbf{t}, \mathbf{s}) \big|_{\mathbf{s}=\mathbf{c}}}{m!}}_{\text{depends on tgt/ctr}} \underbrace{(\mathbf{s} - \mathbf{c})^{\mathbf{m}}}_{\text{depends on src/ctr}}$$

Taylor Series based FMM

Expansion Types:

- Special purpose expansions (Spherical harmonics, Fourier-bessel based)
- Linear Algebra (Eg: Kernel-independent FMM)
- Taylor series based expansions

Pros	Cons
<ul style="list-style-type: none">- Easily tractable symbolically for any kernel	<ul style="list-style-type: none">- Expansions $O(p^3)$ compared to $O(p^2)$- Translations $O(p^6)$ compared to $O(p^2 \log(p))$- Stability issues

Table 1: Pros and cons of Taylor series based expansions

Compressed Multipole Expansion

When ψ satisfies the Helmholtz equation,

$$\psi_{xx} + \psi_{yy} + \kappa^2 \psi = 0.$$

Recall

$$\psi(\mathbf{t}, \mathbf{s}) = \sum_{|m| \leq p} \underbrace{\frac{D_{\mathbf{s}}^m \psi(\mathbf{t}, \mathbf{s})|_{\mathbf{s}=\mathbf{c}}}{m!}}_{\text{depends on tgt/ctr}} \underbrace{(\mathbf{s} - \mathbf{c})^m}_{\text{depends on src/ctr}}$$

From the PDE we have

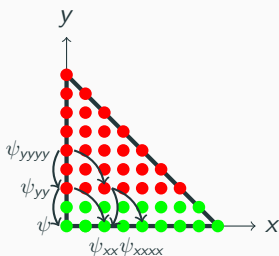
$$\begin{aligned} c_1 \psi_{xx} + c_2 \psi_{yy} + c_3 \psi &= c_1 \psi_{xx} + c_2 (-\psi_{xx} - \kappa^2 \psi) + c_3 \psi \\ &= (c_1 - c_2) \psi_{xx} + 0 \psi_{yy} + \psi (c_3 - \kappa^2 c_2). \end{aligned}$$

Compressed Multipole Expansion

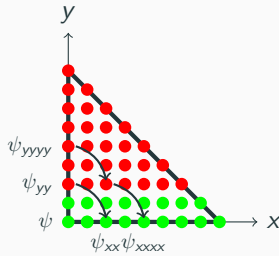
For Helmholtz equation we also have

$$\psi_{xxyy} + \psi_{yyyy} + \kappa^2 \psi_{yy} = 0,$$

$$\psi_{xxxx} + \psi_{xxyy} + \kappa^2 \psi_{xx} = 0.$$



Helmholtz 2D



Laplace 2D

All the coefficients represented by red dots get zeroed.

Count of expansion coefficients go from $\mathcal{O}(p^d)$ to $\mathcal{O}(p^{d-1})$.

Compressed Local Expansion

Recall

$$\psi(\mathbf{t}, \mathbf{s}) = \sum_{|m| \leq p} \underbrace{\frac{D_{\mathbf{t}}^m \psi(\mathbf{t}, \mathbf{s})|_{\mathbf{t}=\mathbf{c}}}{m!}}_{\text{depends on src/ctr}} \underbrace{(\mathbf{t} - \mathbf{c})^m}_{\text{depends on tgt/ctr}}$$

Out of $\mathcal{O}(p^d)$ coefficients, only $\mathcal{O}(p^{d-1})$ are independent.

This makes the number of terms of a local expansion to be $\mathcal{O}(p^{d-1})$.

Calculating derivatives for Local Expansion

Tausch (2003) proposes an algorithm which has an amortized $\mathcal{O}(p)$ time.

We found several formulae to calculate these in amortized $\mathcal{O}(1)$ time.

For Laplace 3D

$$\begin{aligned} r^2 \frac{\partial^{n+m+l}}{\partial x^n y^m z^l} \left(\frac{1}{r} \right) = & - (2n-1)x \frac{\partial^{n+m-1}}{\partial x^{n-1} y^m z^l} \left(\frac{1}{r} \right) - (n-1)^2 \frac{\partial^{n+m-2}}{\partial x^{n-2} y^m z^l} \left(\frac{1}{r} \right) - 2my \frac{\partial^{n+m-1}}{\partial x^n y^{m-1} z^l} \left(\frac{1}{r} \right) \\ & - m(m-1) \frac{\partial^{n+m-2}}{\partial x^n y^{m-2} z^l} \left(\frac{1}{r} \right) - 2lz \frac{\partial^{n+m-1}}{\partial x^n y^m z^{l-1}} \left(\frac{1}{r} \right) - l(l-1) \frac{\partial^{n+m-2}}{\partial x^n y^m z^{l-2}} \left(\frac{1}{r} \right) \end{aligned}$$

For Biharmonic 2D,

$$\begin{aligned} r^2 \frac{\partial^{n+m}}{\partial x^n y^m} \left(r^2 \log(r) \right) = & - 2(n-2)x \frac{\partial^{n+m-1}}{\partial x^{n-1} y^m} \left(r^2 \log(r) \right) - (n-1)(n-4) \frac{\partial^{n+m-2}}{\partial x^{n-2} y^m} \left(r^2 \log(r) \right) \\ & - 2my \frac{\partial^{n+m-1}}{\partial x^n y^{m-1}} \left(r^2 \log(r) \right) - m(m-1) \frac{\partial^{n+m-2}}{\partial x^n y^{m-2}} \left(r^2 \log(r) \right). \end{aligned}$$

This reduces the cost of P2L from $\mathcal{O}(p^d)$ to $\mathcal{O}(p^{d-1})$.

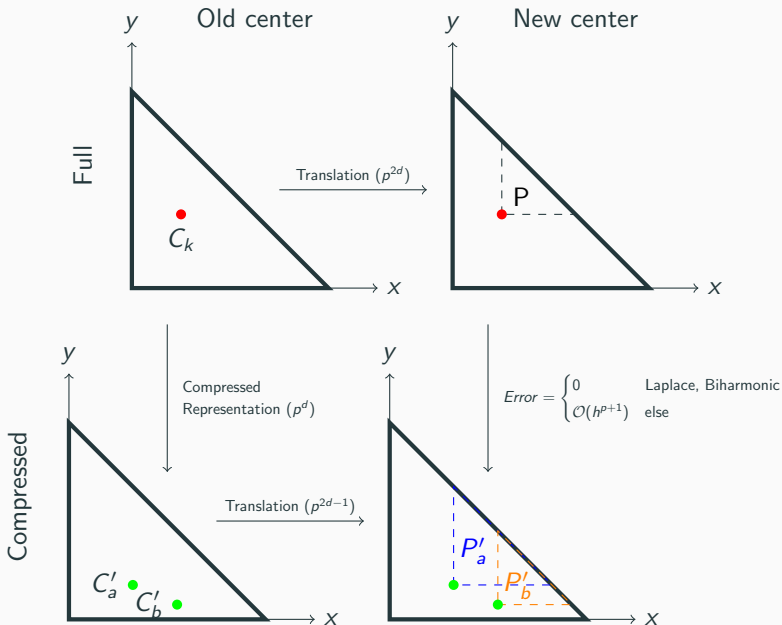
Naive Multipole Translation

Let \mathbf{c}_1 be the old center and \mathbf{c} be the new center. Then,

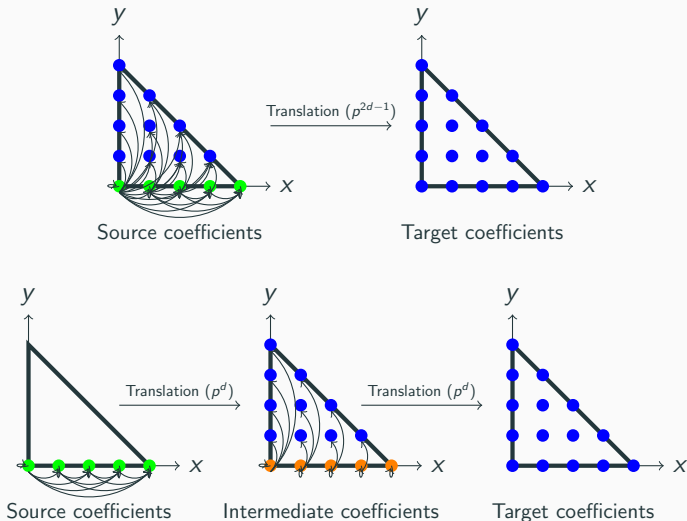
$$\begin{aligned}(\mathbf{s} - \mathbf{c})^k &= ((\mathbf{s} - \mathbf{c}_1) + (\mathbf{c}_1 - \mathbf{c}))^k \\&= \sum_{l \leq k} \binom{k}{l} (\mathbf{s} - \mathbf{c}_1)^l (\mathbf{c}_1 - \mathbf{c})^{k-l} \\&= \sum_{l \leq k} \beta_{k,l} (\mathbf{s} - \mathbf{c}_1)^l\end{aligned}$$

Cost: $\mathcal{O}(p^{2d})$.

Compressed Multipole Translation

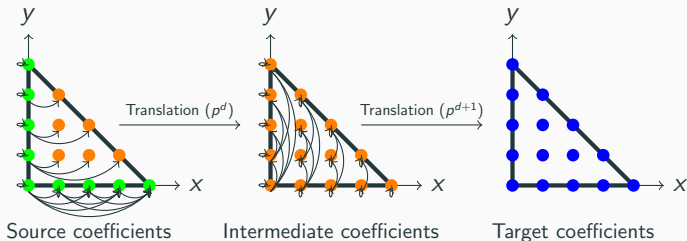


Faster Compressed Multipole Translation

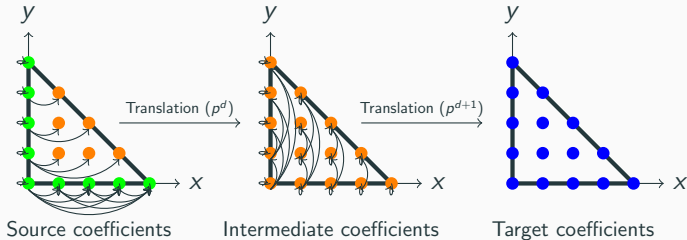


Note: For local to local translation, reverse all arrows.

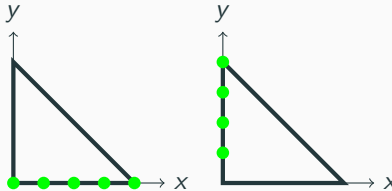
Faster Compressed Multipole Translation



Faster Compressed Multipole Translation



Divide the problem into 2 subproblems



Compressed Multipole to Local Translation

From multipole expansion, we get,

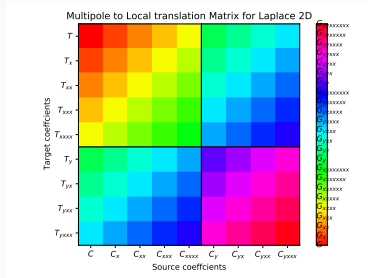
$$\psi(\mathbf{t}, \mathbf{s}) = \sum_{|m| \leq k} \underbrace{\frac{D_{\mathbf{s}}^m \psi(\mathbf{t}, \mathbf{s})|_{\mathbf{s}=\mathbf{c}}}{m!}}_{\text{depends on tgt/ctr}} \underbrace{(\mathbf{s} - \mathbf{c})^m}_{\text{depends on src/ctr}}$$

To translate this multipole expansion to a local expansion, we need to get the derivatives of the above expression and evaluate at new center.

Cost: $\mathcal{O}(p^{2d-2})$.

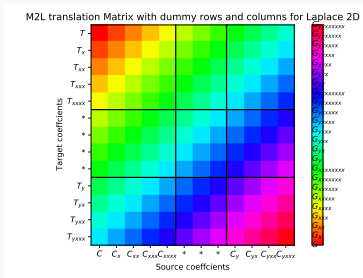
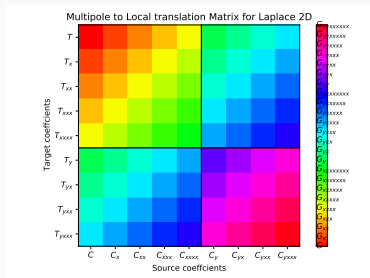
Compressed Multipole to Local Translation

Multipole to local translation matrix is a block Toeplitz matrix of smaller Toeplitz matrices.



Compressed Multipole to Local Translation

Multipole to local translation matrix is a block Toeplitz matrix of smaller Toeplitz matrices.



Use an FFT to do the translation similar to Greengard (1988).

Cost depends on number of dummy rows:

- $\mathcal{O}(p^{d-1} \log(p))$ for elliptic PDEs
- $\mathcal{O}(p^d \log(p))$ for other PDEs

Time complexities

	P2L/M2P	P2M/L2P	M2M	M2L	L2L
Taylor Series	p^3	p^3	p^6	p^6	p^6
Improved Taylor Series	p^3	p^3	p^4	$p^3 \log(p)$	p^6
Compressed Taylor Series without fast derivatives	p^3	p^3	p^3	$p^2 \log(p)$	p^3
Compressed Taylor Series with fast derivatives	p^2	p^3	p^3	$p^2 \log(p)$	p^3
Spherical Harmonic Series	p^2	p^2	$p^2 \log(p)$	$p^2 \log(p)$	$p^2 \log(p)$

Table 2: Time complexities for expansions, translations and evaluations

All operations are exact except for M2M in Compressed Taylor and M2L operations with FFT.

With Compressed Taylor generating code for Stokes

$$\begin{aligned}\mu \nabla^2 \mathbf{u} - \nabla p + \mathbf{f} &= \mathbf{0} \\ \nabla \cdot \mathbf{u} &= 0\end{aligned}$$

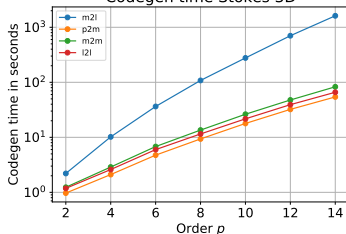
is done simply by giving the PDE as,

```
w = make_pde_syms(dim, dim+1)
mu = sym.Symbol("mu")
u = w[:dim]
p = w[-1]
pdes = PDE(mu * laplacian(u) - grad(p), div(u))
```

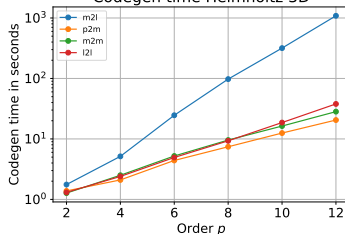
which generates code for the expansion, translations and evaluations.

Code generation

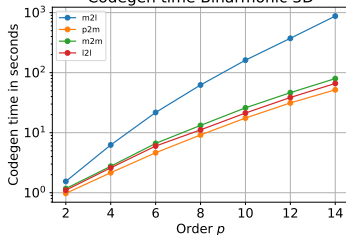
Codegen time Stokes 3D



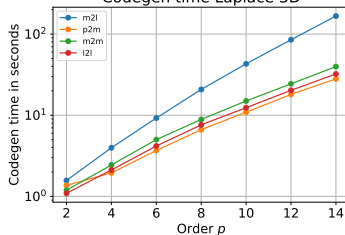
Codegen time Helmholtz 3D



Codegen time Biharmonic 3D

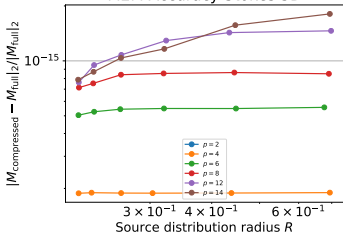


Codegen time Laplace 3D

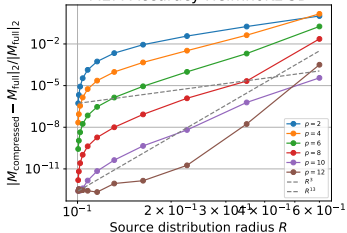


Results - Error M2M

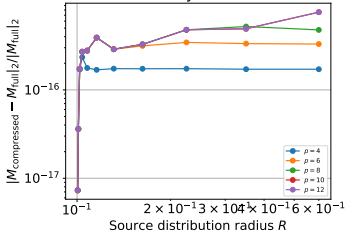
M2M Accuracy Stokes 3D



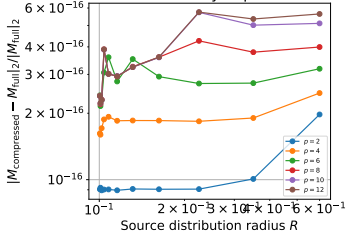
M2M Accuracy Helmholtz 3D



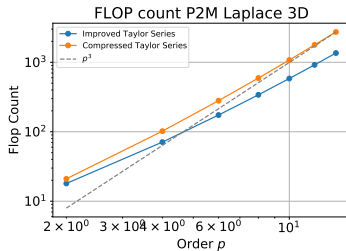
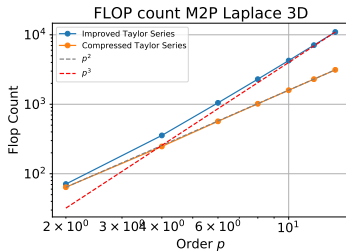
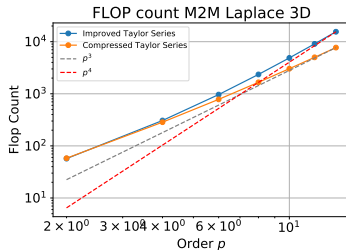
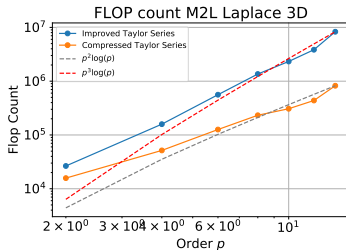
M2M Accuracy Biharmonic 3D



M2M Accuracy Laplace 3D



Results - FLOP count



Summary

- Kernel generic method for elliptic constant coefficient linear PDEs.
- Only needs the PDE and the Green's function for the PDE.
- Asymptotically better than full Taylor Series in
 - Number of FLOPs
 - Storage
- Next goal: A fast Stokes solver on a GPU.

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