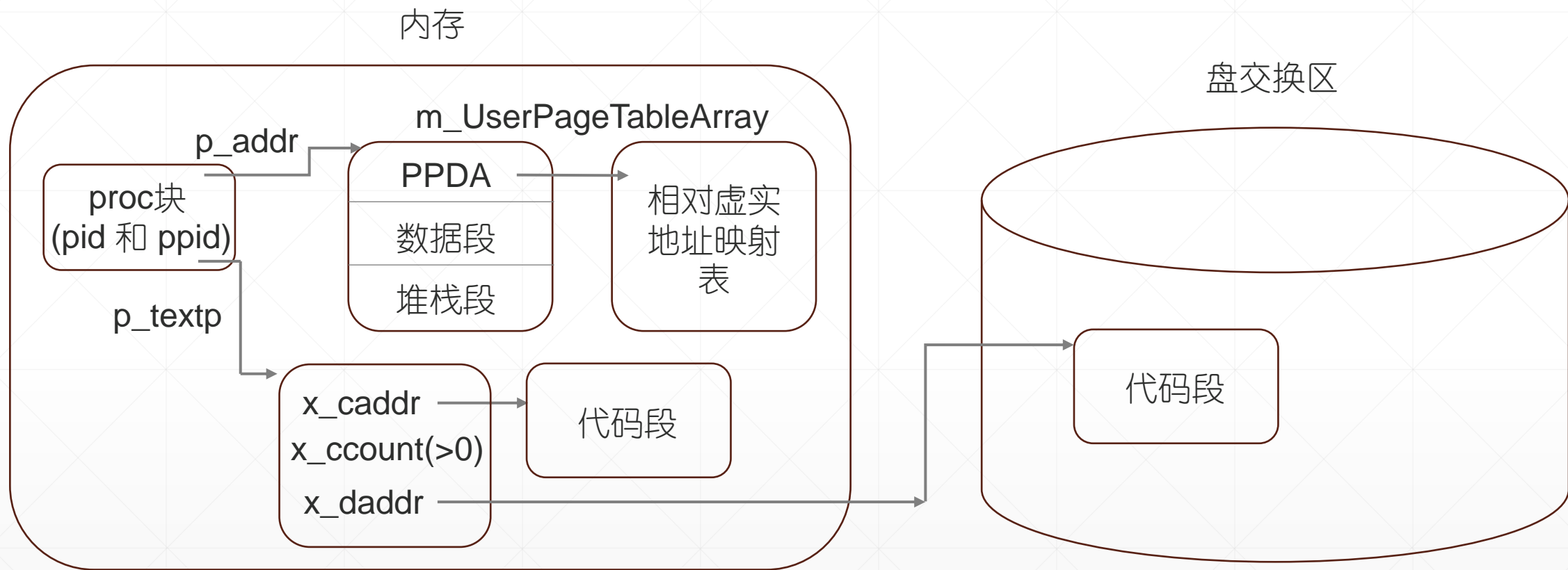
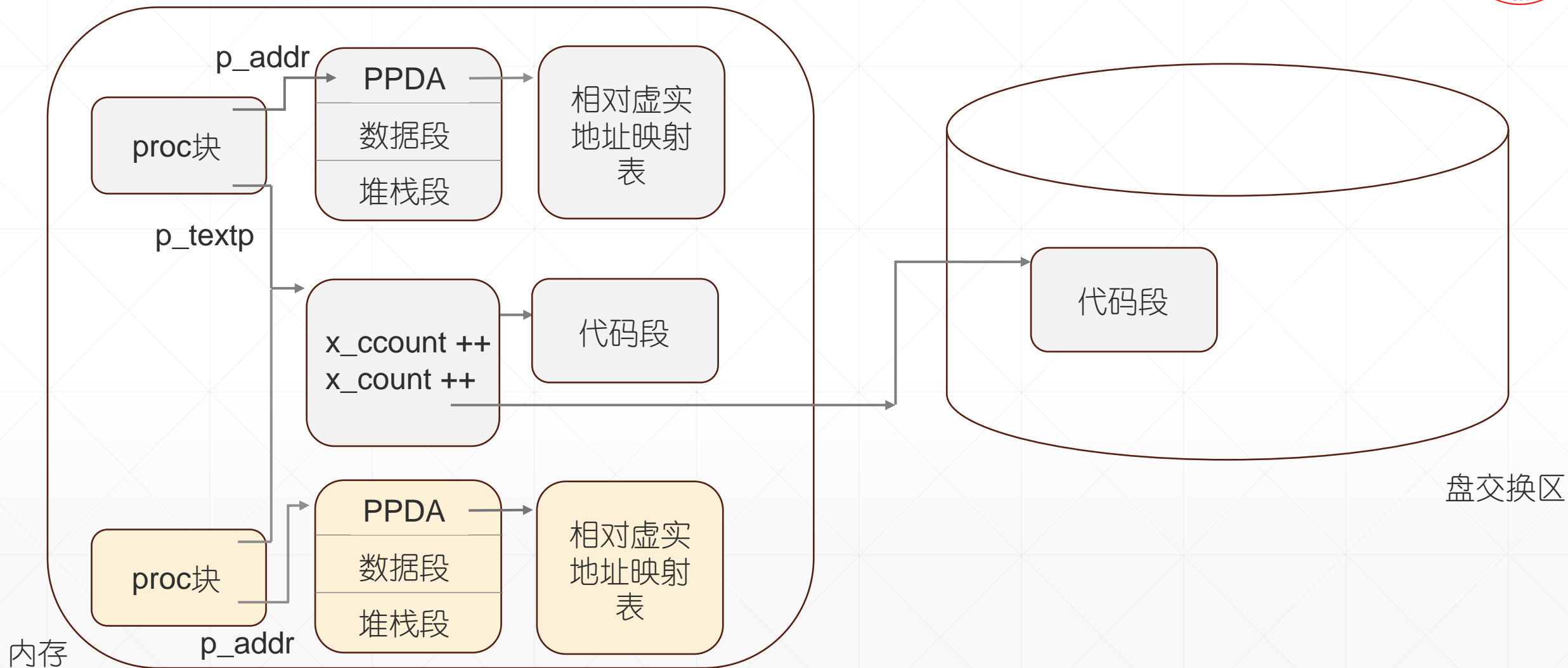


操作系统 第四章 进程管理

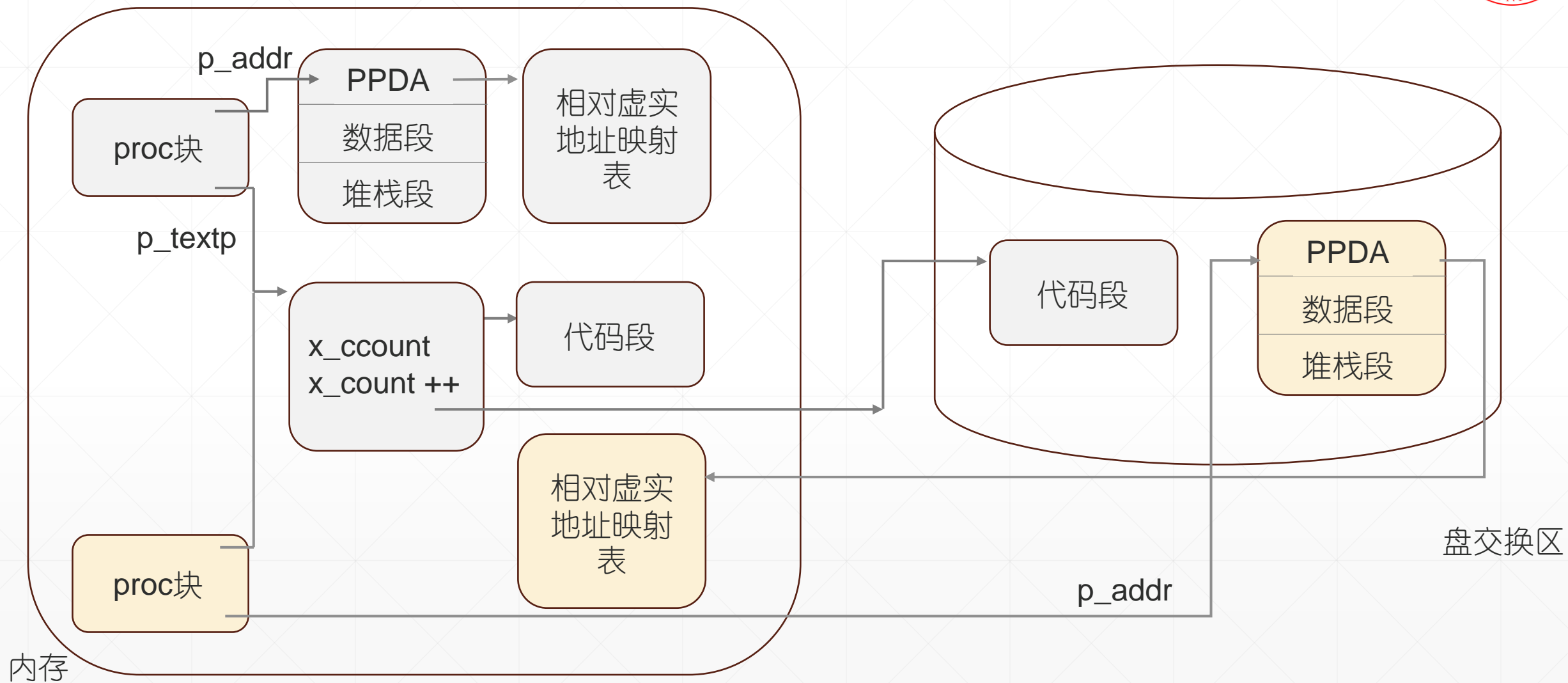
4.5 fork系统调用



创建在内存中的子进程图像



创建在盘交换区上的子进程图像



复制给子进程的 Process结构

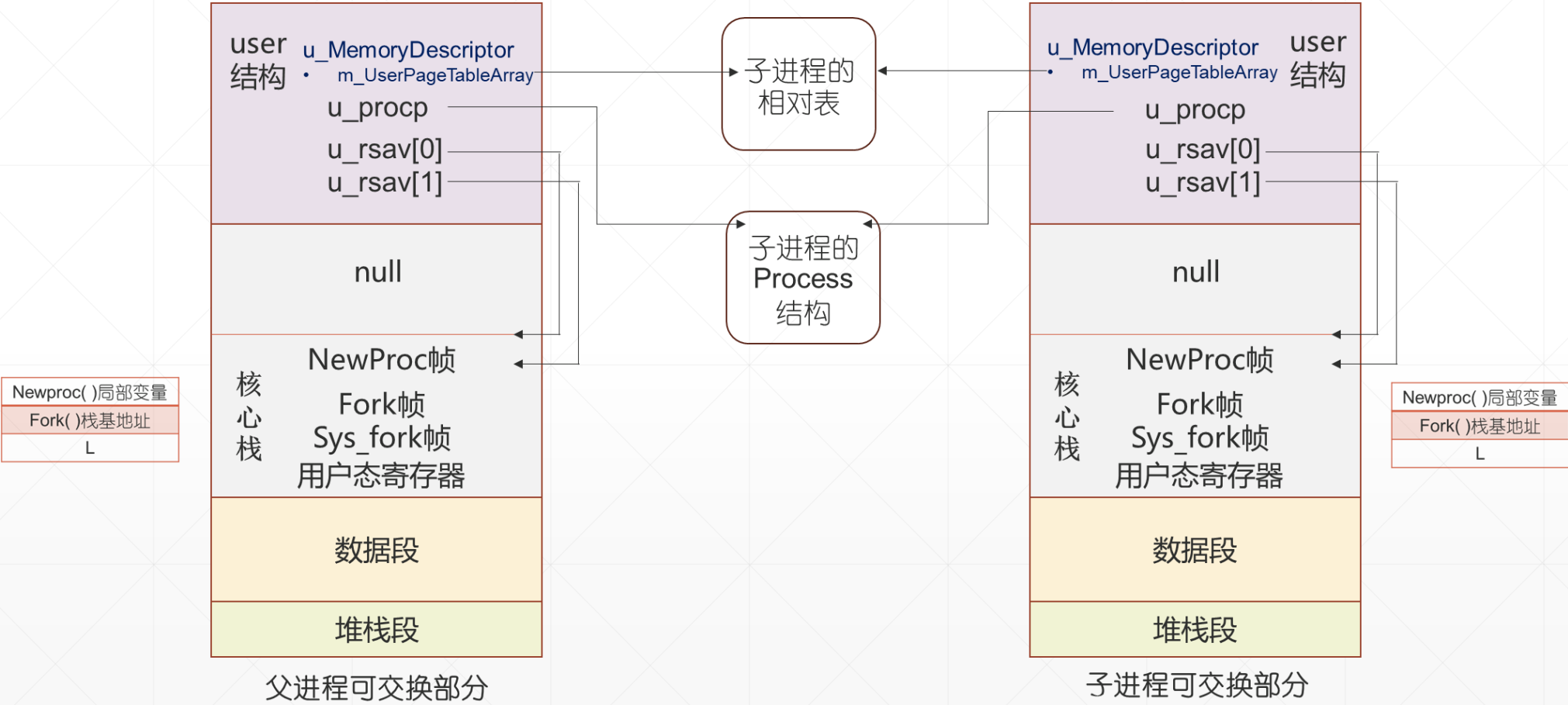
```
void Process::Clone(Process& proc) // 指针 this 和 proc 分别指向父进程和子进程的 Process 结构
{
    User& u = Kernel::Instance().GetUser();

    /* 拷贝父进程 Process 结构中的大部分数据 */
    proc.p_size = this->p_size;
    proc.p_stat = Process::SRUN;
    proc.p_flag = Process::SLOAD;
    proc.p_uid = this->p_uid;
    proc.p_ttyp = this->p_ttyp;
    proc.p_nice = this->p_nice;
    proc.p_textp = this->p_textp;

    proc.p_pid = ProcessManager::NextUniquePid(); // 为子进程分配 pid
    proc.p_ppid = this->p_pid; // 子进程登记父进程的 pid

    proc.p_pri = 0; // 确保 child 的优先数较小, 先于父进程运行 (Unix V6++ 并无此必要)
    proc.p_time = 0; // 驻留时间清 0
```

复制给子进程的可交换部分

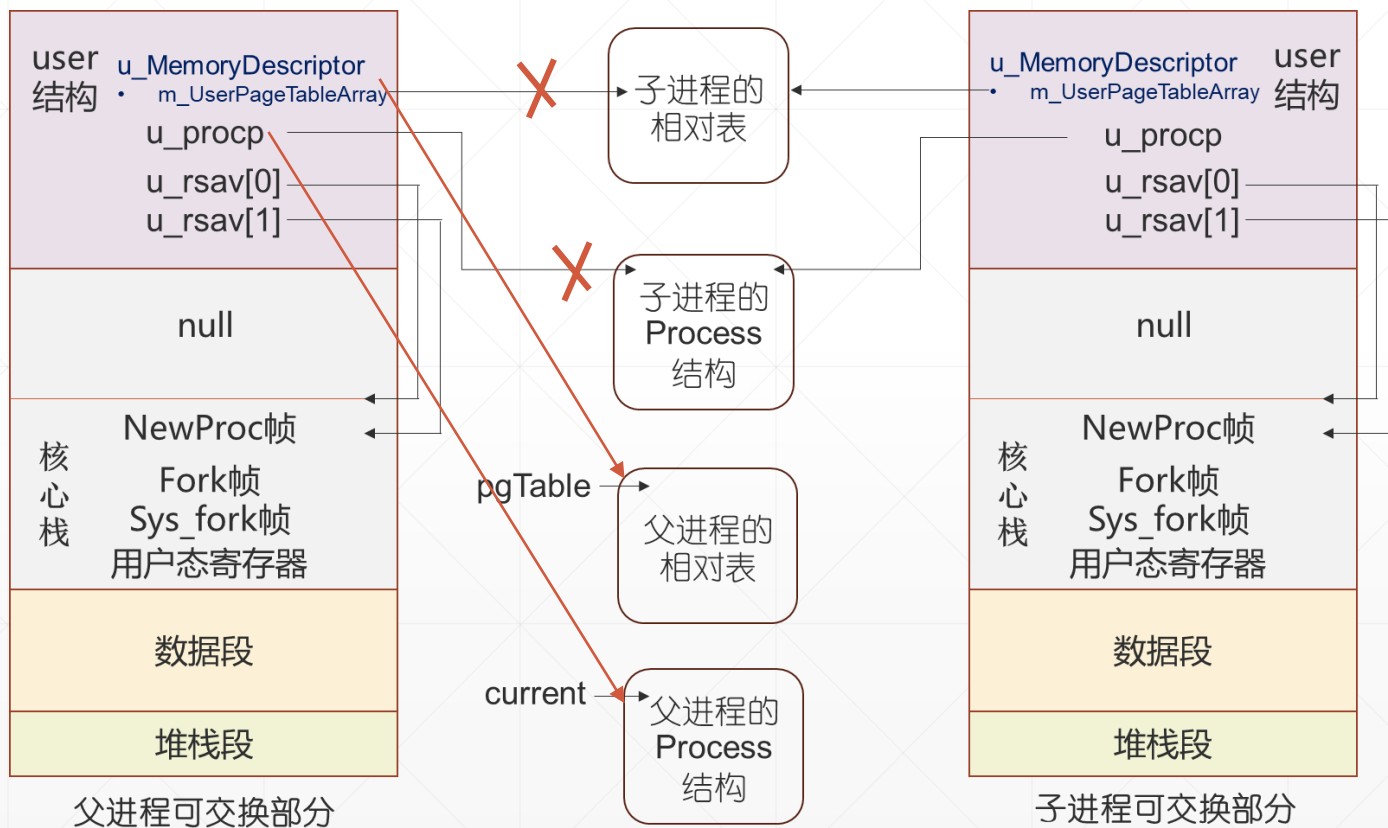
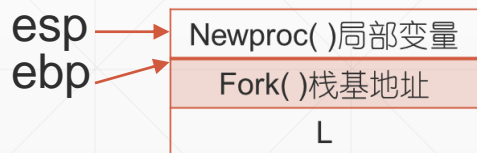


父进程修改自己的指针后，NewProc返回 0

```
u.u_procp = current;
u.u_MemoryDescriptor.m_UserPageTableArray = pgTable;
return 0;
```

```
}
```

```
mov $0, %eax
mov %ebp, %esp
pop %ebp
ret
```



NewProc() → Fork()

```
if ( this->NewProc() )
```

```
{
```

```
    u.u_ar0[User::EAX] = 0;
```

```
    u.u_cstime = 0;
```

```
    u.u_stime = 0;
```

```
    u.u_cutime = 0;
```

```
    u.u_utime = 0;
```

```
}
```

```
else
```

```
{
```

```
    u.u_ar0[User::EAX] = child->p_pid;
```

```
}
```

```
return;
```

```
}
```

```
call NewProc
```

```
L: cmp %eax, 0
```

```
    je fatherBranch
```

```
childBranch:
```

```
    .....
```

```
    jmp exit
```

```
fatherBranch:
```

```
    .....
```

```
exit:
```

父进程 fork 系统调用的返回值是 子进程pid

子进程被 Swtch() 选中

```
int ProcessManager::Swtch()
{
```

```
    User& u = Kernel::Instance().GetUser();
    SaveU(u.u_rsav);
```

```
    Process* procZero = &process[0];
    X86Assembly::CLI();
    SwtchUStruct(procZero);
    RetU();
    X86Assembly::STI();
```

```
    Process* selected = Select();
```

```
    X86Assembly::CLI();
    SwtchUStruct(selected);
    RetU();
    X86Assembly::STI();
    User& newu = Kernel::Instance().GetUser();
    newu.u_MemoryDescriptor.MapToPageTable();
```

```
    .....
```

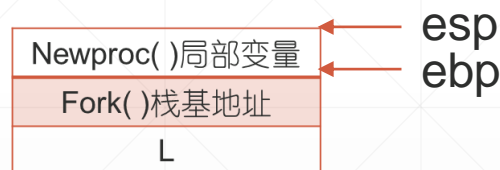
```
    return 1;
```

```
}
```

子进程被 Swtch 选中，RetU 从 u_rsav 中恢复 esp 和 ebp。建立地址映射关系后，退栈 return 1。。。实际上执行的是从NewProc，系统调用返回的工作。。。返回值是 1。



子进程可交换部分



```
mov $1, %eax
mov %ebp, %esp
pop %ebp
ret
```

Swch() → Fork()

```
if ( this->NewProc() )  
{
```

```
u.u_ar0[User::EAX] = 0;  
u.u_cstime = 0;  
u.u_stime = 0;  
u.u_cutime = 0;  
u.u_utime = 0;
```

```
}  
else
```

```
{
```

```
u.u_ar0[User::EAX] = child->p_pid;
```

```
}
```

```
return;
```

```
}
```

```
call NewProc  
L: cmp %eax, 0  
je fatherBranch  
childBranch:
```

```
.....
```

```
jmp exit
```

```
fatherBranch:
```

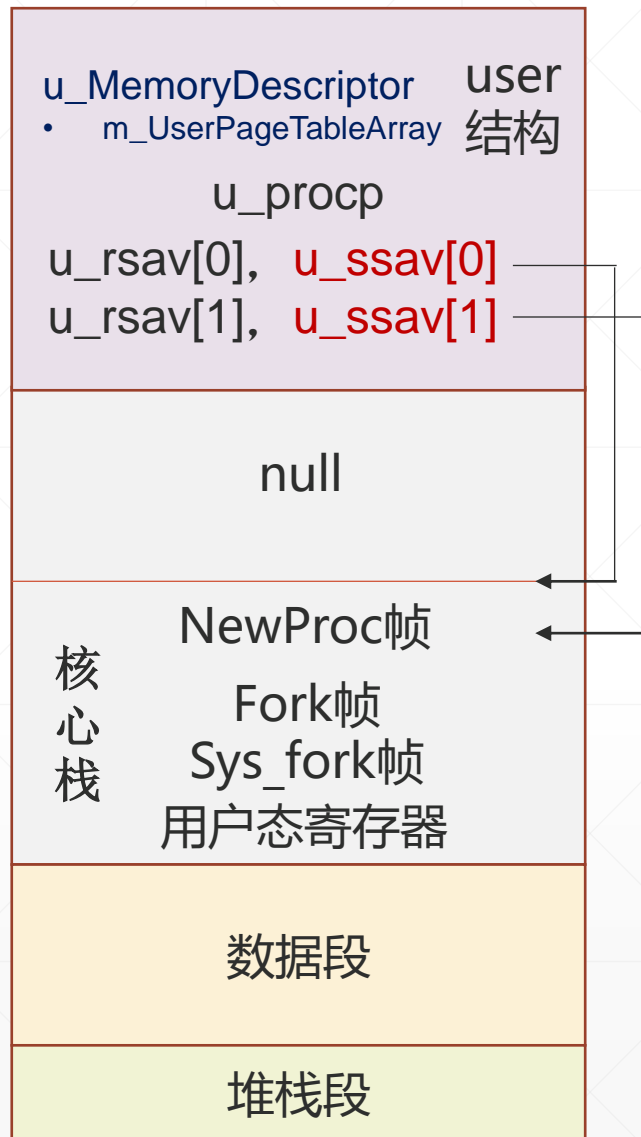
```
.....
```

```
exit:
```

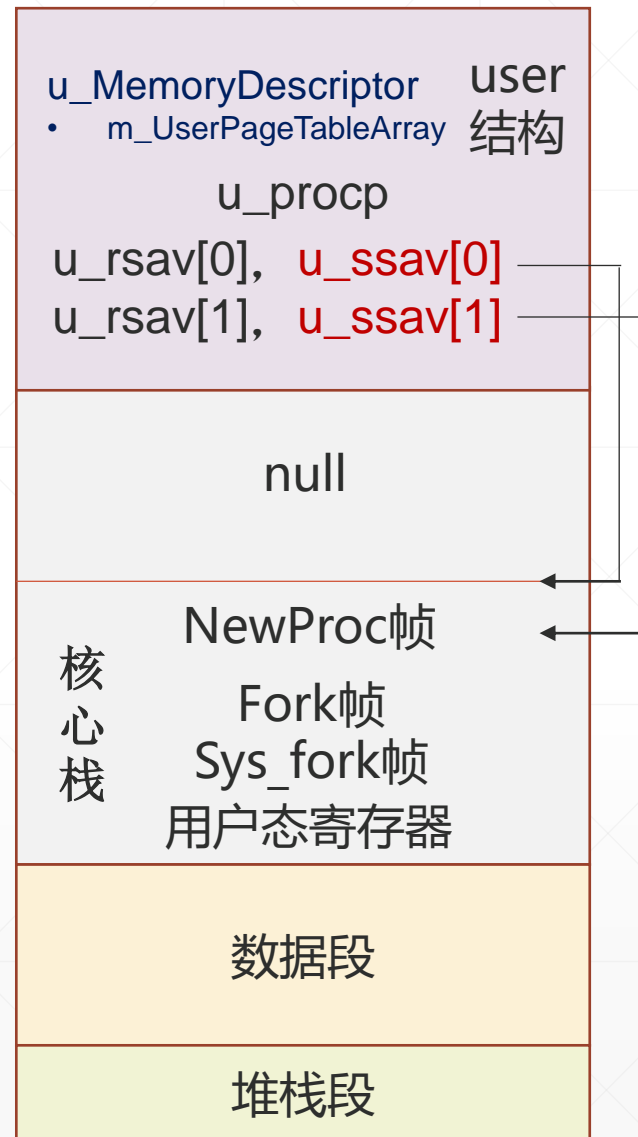
fork 系统调用的返回值是 0

XSwap 执行IO操作复制父进程图像

内存中的父进程图像



盘交换区上的子进程图像



`pProcess->p_addr:`
`blkno`
`pProcess->p_flag:`
`~SLOAD`
`pProcess->p_time`
`0`

1、NewProc帧的位置复制进u_ssav数组。

2、如果 0# 进程因 RunOut入睡，会被唤醒。

3、打 XSwap 标记



未来，0#进程会将其换入内存

被选中后

```
int ProcessManager::Switch()
{
    .....
    /* 恢复被保存进程的现场 */
    X86Assembly::CLI();
    SwtchUStruct(selected);
    RetU();
    X86Assembly::STI();

    User& newu = Kernel::Instance().GetUser();
    newu.u_MemoryDescriptor.MapToPageTable();

    if ( newu.u_procp->p_flag & Process::SSWAP )
    {
        newu.u_procp->p_flag &= ~Process::SSWAP;
        aRetU(newu.u_ssav);
    }
    return 1;
}
```



带 SSWAP 标记，从 u_ssav 中恢复 ESP 和 EBP。退栈，fork系统调用返回用户态

