

CS 1332R

WEEK 13

Introduction to Graphs

Breadth-First Search

Depth-First Search

Exam 3 Review



ANNOUNCEMENTS



Introduction to Graphs

- ❑ We define a graph (**G**) by its **vertices** and **edges**.
- ❑ We define an edge (**e**) by the two vertices it connects (**u, v**) and its weight (**w**).

V = the set of vertices

E = the set of edges

|V| = the # of vertices

|E| = the # of edges

RELEVANT TERMINOLOGY

ORDER(G) = $|V|$

SIZE(G) = $|E|$

Indegree(v) = # of edges going into a vertex

Outdegree(v) = # of edges going out of a vertex

Adjacent = describes two vertices that are connected by an edge

Connected Graph = every vertex has a path to every other vertex

Unconnected Graph = not all vertices have a path to every other vertex

Directed Graph = edges have a direction, they are distinct ordered pairs

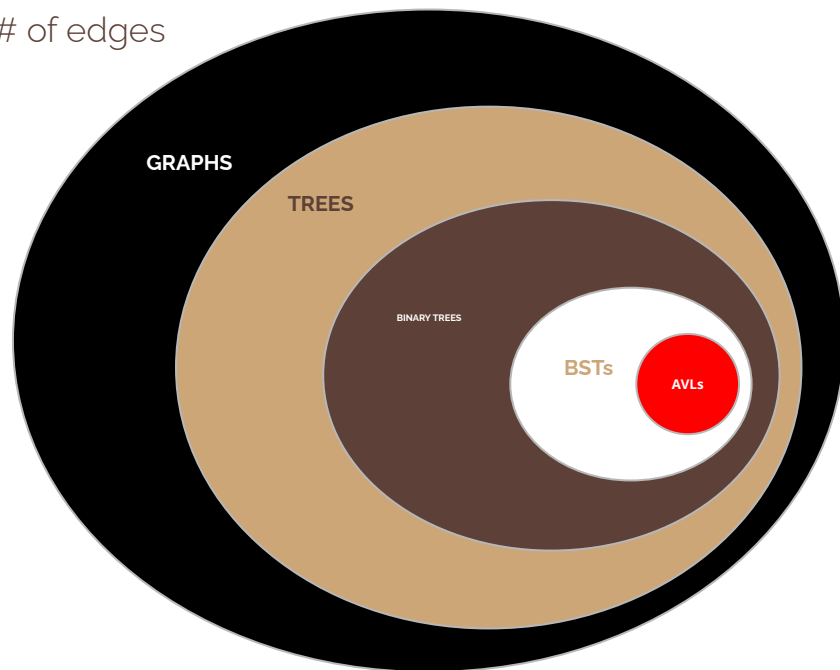
Undirected Graph = edges do not have a direction, $e(u, v)$ is the same as $e(v, u)$

Subgraph = a graph G' such that V' is a subset of V and E' is a subset of E

Cycle = a path with the same start and end vertex

Acyclic Graph = a graph containing no cycles

Tree = a minimally connected, acyclic graph



Introduction to Graphs

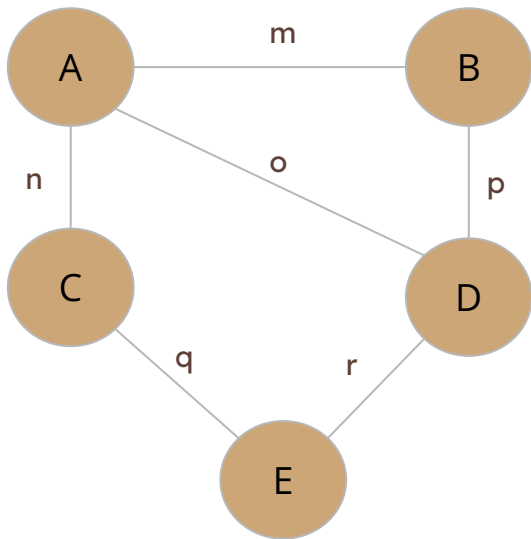
- ❑ We define a graph (**G**) by its **vertices** and **edges**.
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|V| = the # of vertices

E = the set of edges

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HOW WE REPRESENT A GRAPH

1. Adjacency Matrix
 - a. $|V| \times |V|$ 2D array
 - b. Space: $O(|V|^2)$
 - c. $\text{Matrix}[v1, v2] = \text{edge}(v1, v2, w)$
2. Adjacency List (***used in your homework***)
 - a. Map of each vertex to its incident edges
 - b. Space: $O(|V| + |E|)$
 - c. ***In your homework, we map each vertex to a list of VertexDistance objects. The VertexDistance contains an adjacent vertex and the distance to that adjacent vertex, a.k.a the weight of the edge connecting the two vertices.***
3. Edge Set (***used in your homework***)
 - a. A set of all edges in the graph
 - b. Space: $O(|E|)$

Breadth-First Search (BFS)

- ❑ **PURPOSE:** Traversing a graph from a starting vertex.
- ❑ Finds the shortest distance from the start vertex to other vertices **on an unweighted graph**.
- ❑ Use cases: level order traversal, many future algorithms (e.g. Dijkstra's)

STRUCTURES WE NEED

- **Queue<VertexDistance<T>>** : a priority queue of *VertexDistance* objects ordered by the distance, which is the cumulative distance from the start vertex to the vertex
- **Set<Vertex<T>>** : a **visited set** containing vertices we have found the shortest path to
- **List<Vertex<T>>** : a **list** of the vertices in the graph in the order they were visited

Why do we need a List and a Set to keep track of the visited vertices?

The List is ordered while a Set is unordered. The List has our actual final traversal.

BFS: Implementation

STRUCTURES

- **Queue<VertexDistance<T>>** : a priority queue of VertexDistance objects ordered by the distance, which is the cumulative distance from the start vertex to the vertex
- **Set<Vertex<T>>** : a **visited set** containing vertices we have found the shortest path to
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ALGORITHM

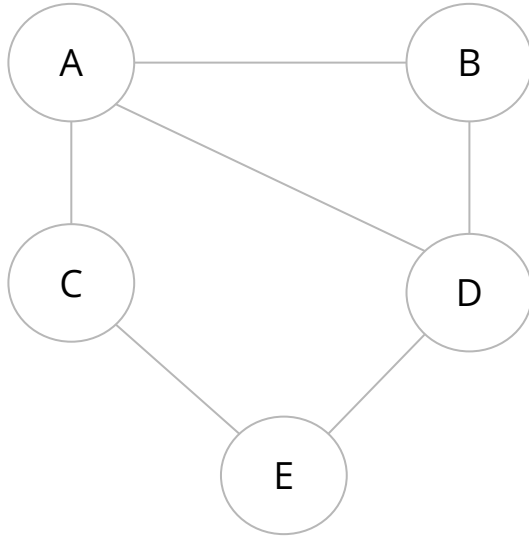
```
initialize s as start vertex
initialize queue, visitedSet, list
queue.enqueue(s)
while queue is not empty and visitedSet is not full: 2 termination conditions
    v = queue.dequeue()
    if v is not in visitedSet:
        visitedSet.add(v)
        list.add(v)
        for all u adjacent to v and u not in visitedSet:
            queue.enqueue(u)
```

Efficiency: $O(|V| + |E|)$

- Accessing every vertex and edge once
- In Java, Set is implemented as a HashSet \rightarrow .contains() is $O(1)$.

***BFS:* Practice**

Start at vertex A.



DIAGRAMMING SETUP

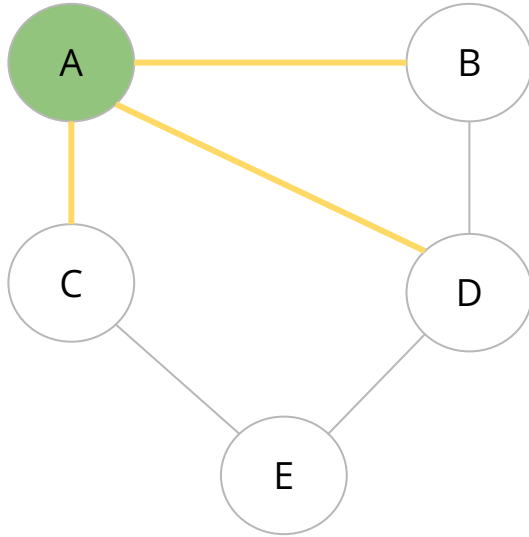
QUEUE

VISITED SET

A

BFS: Practice

Start at vertex A.



DIAGRAMMING SETUP

QUEUE

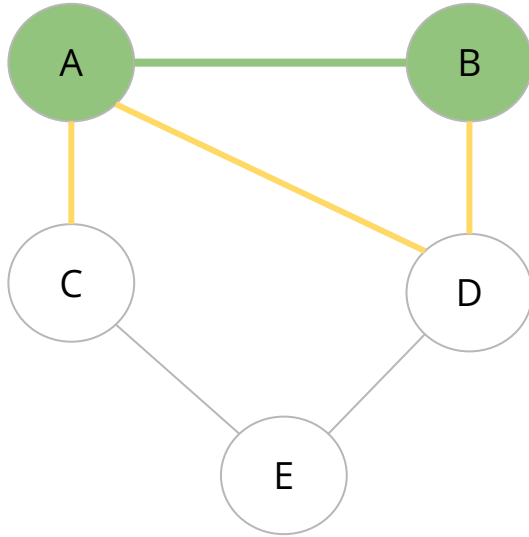
A
B
C
D

VISITED SET

A

BFS: Practice

Start at vertex A.



DIAGRAMMING SETUP

QUEUE

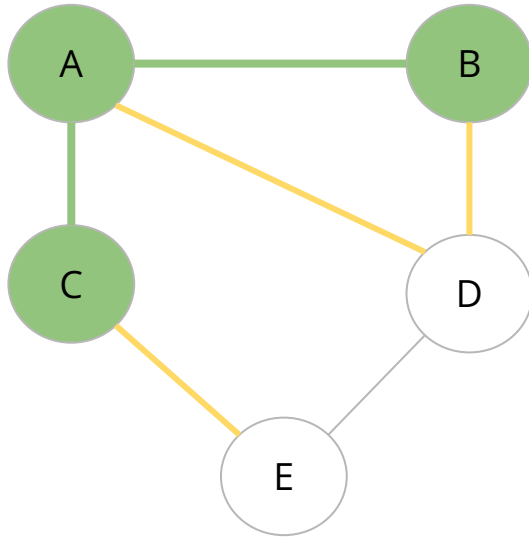
A
~~B~~
C
D
D

VISITED SET

A
B

BFS: Practice

Start at vertex A.



DIAGRAMMING SETUP

QUEUE

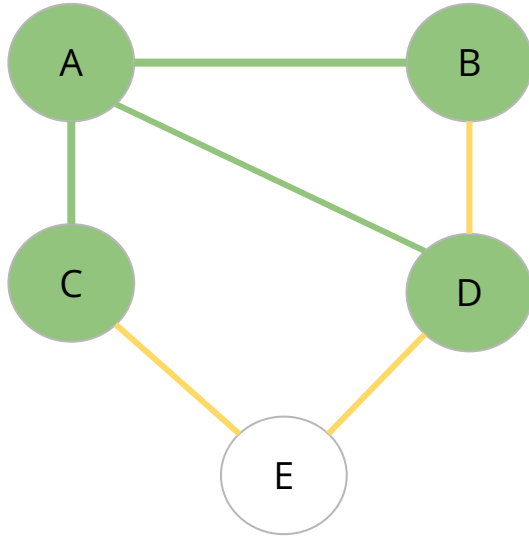
A
~~B~~
~~C~~
D
D
E

VISITED SET

A
B
C

BFS: Practice

Start at vertex A.



DIAGRAMMING SETUP

QUEUE

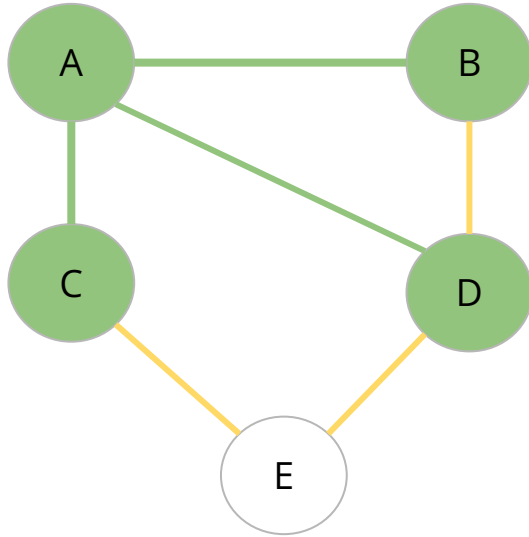
A
~~B~~
~~C~~
~~D~~
E

VISITED SET

A
B
C
D

BFS: Practice

Start at vertex A.



DIAGRAMMING SETUP

QUEUE

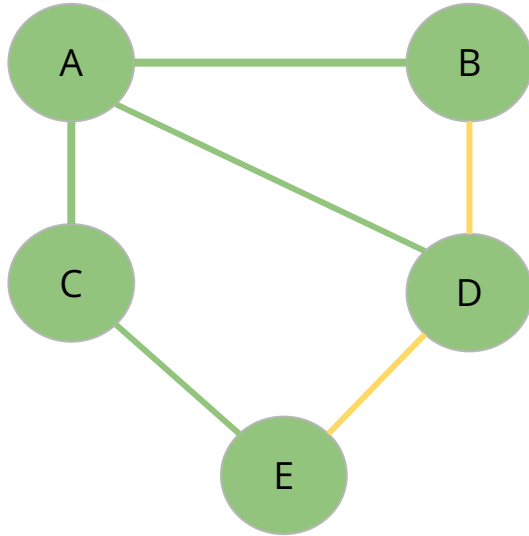
A
~~B~~
~~C~~
~~D~~
~~E~~

VISITED SET

A
B
C
D

BFS: Practice

Start at vertex A.



DIAGRAMMING SETUP

QUEUE

A
~~B~~
~~C~~
~~D~~
~~E~~

VISITED SET

A
B
C
D
E

Depth-First Search (DFS)

- ❑ **PURPOSE:** Traversing a graph from a starting vertex.
- ❑ Implemented with a stack, we will use the recursive stack
- ❑ Use Cases: pre-, post-, in-order traversals, recursion

STRUCTURES WE NEED

implicit use of a stack through recursion

- **Set<Vertex<T>>** : a **visited set** containing vertices we have found the shortest path to
- **List<Vertex<T>>** : a **list** of the vertices in the graph in the order they were visited

DFS: Implementation

STRUCTURES

- **Set<Vertex<T>>** : a **visited set** containing vertices we have found the shortest path to
- **List<Vertex<T>>** : a **list** of the vertices in the graph in the order they were visited

ALGORITHM

```
initialize s as start vertex
initialize visitedSet, list
dfs(s)
```

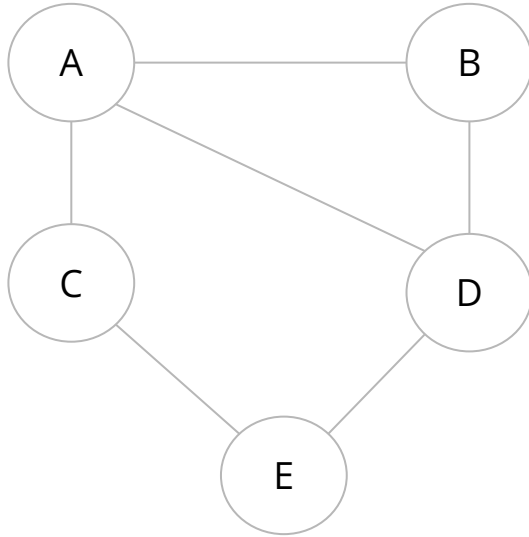
```
dfs(v):                                     base case
    if v not in visitedSet:
        visitedSet.add(v)
        list.add(v)
        for all w adjacent to v and not in visitedSet:
            dfs(w)
```

Efficiency: $O(|V| + |E|)$

- Accessing every vertex and edge once (same as BFS)
- In Java, Set is implemented as a HashSet \rightarrow .contains() is $O(1)$.

DFS: Practice

Start at vertex A.



DIAGRAMMING SETUP

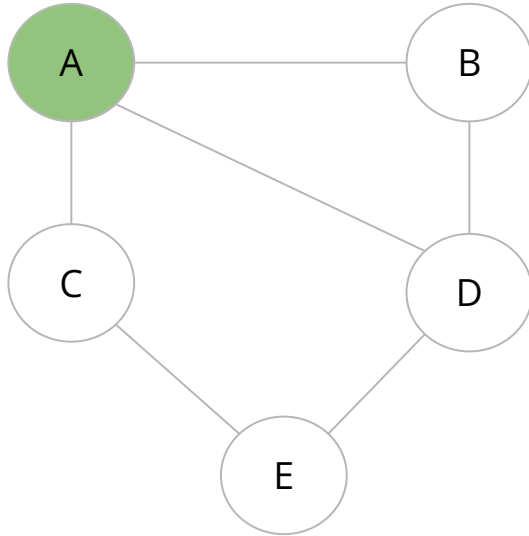
STACK

VISITED SET

dfs(A)

DFS: Practice

Start at vertex A.



DIAGRAMMING SETUP

STACK

VISITED SET

~~dfs(A)~~

A

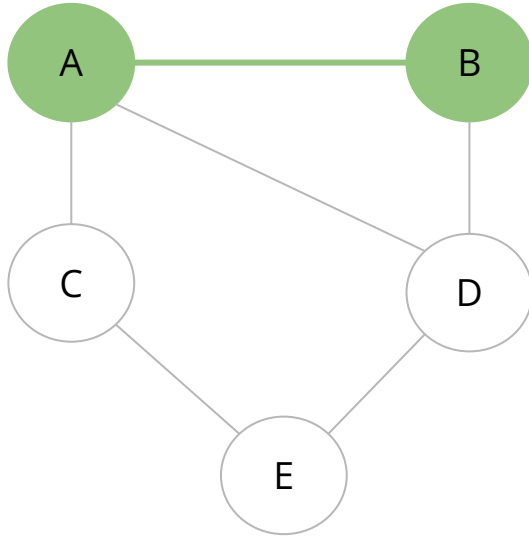
dfs(B)

dfs(C)

dfs(D)

DFS: Practice

Start at vertex A.



DIAGRAMMING SETUP

STACK

VISITED SET

~~dfs(A)~~

A

~~dfs(B)~~

B

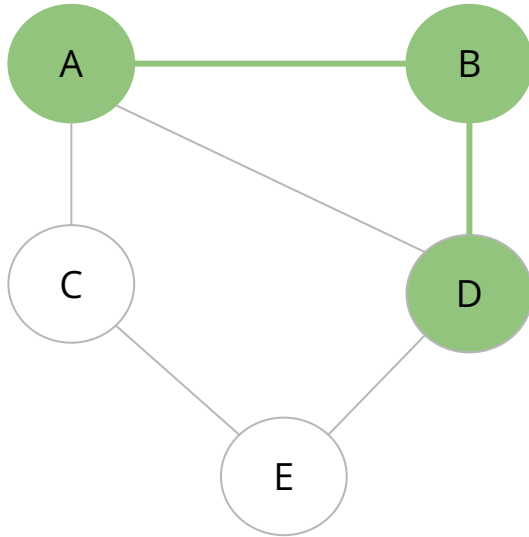
dfs(C)

dfs(D)

dfs(D)

DFS: Practice

Start at vertex A.



DIAGRAMMING SETUP

STACK

VISITED SET

~~dfs(A)~~

A

~~dfs(B)~~

B

dfs(C)

D

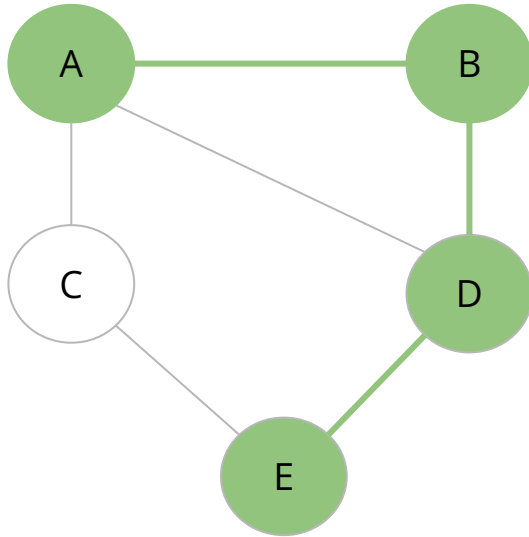
dfs(D)

~~dfs(D)~~

~~dfs(E)~~

DFS: Practice

Start at vertex A.



DIAGRAMMING SETUP

STACK

VISITED SET

~~dfs(A)~~

A

~~dfs(B)~~

B

dfs(C)

D

dfs(D)

E

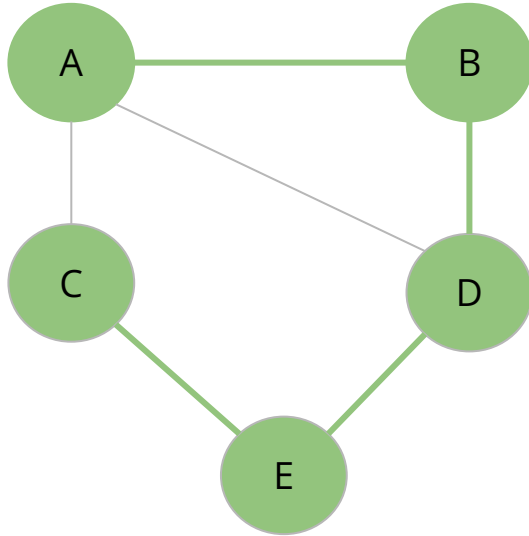
~~dfs(D)~~

~~dfs(E)~~

dfs(C)

DFS: Practice

Start at vertex A.



DIAGRAMMING SETUP

STACK

VISITED SET

~~dfs(A)~~

A

~~dfs(B)~~

B

dfs(C)

D

dfs(D)

E

~~dfs(D)~~

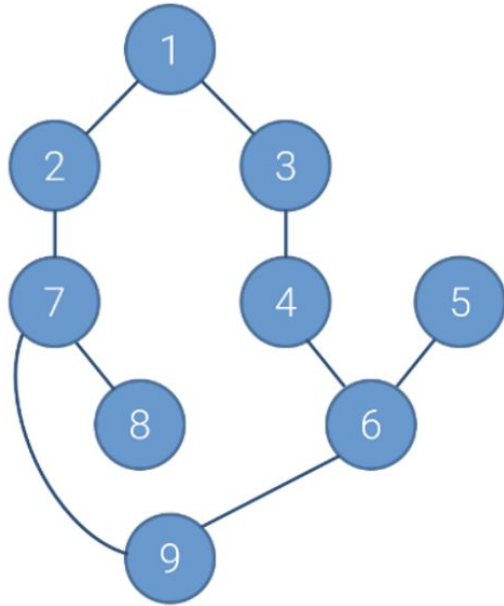
C

~~dfs(E)~~

~~dfs(C)~~

BFS/DFS: Practice

Perform BFS & DFS on this graph starting at vertex 1. Iterate through neighbors in numerical order. Return visited order for each search.



DFS: 1 2 7 8 9 6 4 3 5

BFS: 1 2 3 7 4 8 9 6 5

LEETCODE PROBLEMS

994. Rotting Oranges

785. Is Graph Bipartite?

130. Surrounded Regions

EXAM 3 REVIEW

Kahoot

Socrative: CS1332

Practice exams in Canvas: Files -> Resources -> Recitation
Materials -> Recitation Practice Exams



Any questions?

Name
Office Hours
Contact

Name
Office Hours
Contact



*Let us know if there is anything specific you want out of
recitation!*