

CS 1332R

WEEK 15

Dynamic Programming
Longest Common Subsequence
Final Exam Review



ANNOUNCEMENTS



Introduction to Dynamic Programming (DP)

- ❑ A **strategy** for solving a class of problems that can be broken down into smaller, repetitive problems **that overlap** - different than just divide & conquer recursion
- ❑ Can reduce time complexity from exponential to polynomial
- ❑ Use Cases: KMP Failure Table, shortest path algorithms, scheduling algorithms, and so many more

STEPS:

1. Break a problem down into smaller subproblems.
2. Solve these smaller problems and **save their solutions**.
3. If a subproblem is encountered again, just use the saved result ($O(1)$) instead of solving it again.

Motivation for Dynamic Programming

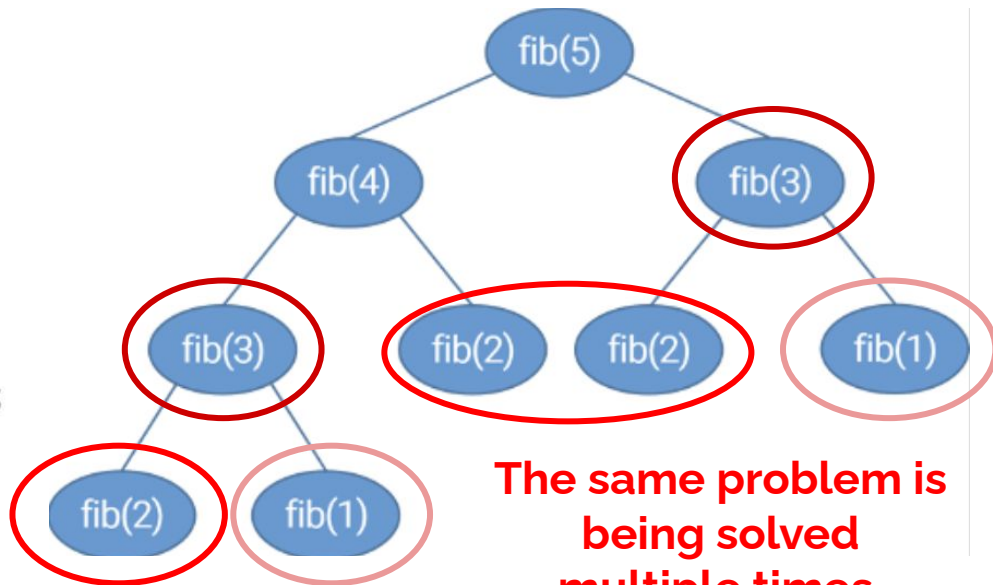
Let's explore the Fibonacci problem.

Fibonacci Sequence: 0, 1, 1, 2, 3, 5, 8, 13, ...

$$\text{fib}(n) = \text{fib}(n - 1) + \text{fib}(n - 2)$$

RECURSIVE, NON-DP SOLUTION

```
public int fib(int n) {  
    if (n == 0) return 0;  
    else if (n == 1) return 1;  
    return fib(n - 1) + fib(n - 2);  
}
```



Implementing Dynamic Programming

TWO APPROACHES

MEMOIZATION

- ❑ Still utilizes recursion, “top-down” approach
- ❑ Saves previously solved solutions using some structure (e.g. map, array, etc.)

```
Map<Integer, Integer> memo;
```

```
public int fib(int n) {  
    if (n == 0) return 0;  
    else if (n == 1) return 1;  
    if (!memo.containsKey(n)) {  
        int result = fib(n - 1) + fib(n - 2);  
        memo.put(n, result);  
    }  
    return memo.get(n);  
}
```

TABULATION

- ❑ Iterative, “bottom-up” approach
- ❑ Saves previously solved solutions in a table
- ❑ ***Faster than memoization***, no overhead space for recursion

“You’ve seen this before with KMP’s failure table!!”

```
public int fib(int n) {  
    int[] f = new int[n + 1];  
    f[0] = 0;  
    f[1] = 1;  
    for (int k = 2; k <= n; k++) {  
        f[k] = f[k - 1] + f[k - 2];  
    }  
    return f[n];  
}
```

Longest Common Subsequence (LCS)

- ❑ **SUBSEQUENCE:** A subset of a string where each character appears in the same order as in the string. The characters are not necessarily contiguous in the string.

Ex: subsequences of BROWN

BRO, BROW, BON, BOWN, BROWN

- ❑ **PURPOSE:** Given two strings, find the longest subsequence that exists in both strings.

Typically...

X = first string

Y = second string

n = length of X

m = length of Y

i = column, index of X

j = row, index of Y

Longest Common Subsequence (LCS)

STEPS

1. Create an $m+1 \times n+1$ 2D array. The value at $arr[i][j]$ is the length of the longest common subsequence of the strings $X.substring(0, i)$ and $J.substring(0, j)$.
2. Fill the first row and column with zeros.
3. Begin filling in the rows left to right, top to bottom using the following rule:

if $X[i] == Y[j]$:

$arr[i][j] = arr[i-1][j-1] + 1$

else:

$arr[i][j] = \max(arr[i-1][j], arr[i][j-1])$

choose cell to the left if equal

DIAGRAMMING SETUP

X: 0 1 2 3 4 5 6							
Y:		s	t	r	i	n	g
0	0	0	0	0	0	0	0
1	s	0					
2	t	0					
3	r	0					
4	i	0					
5	n	0					

location of final answer: $arr[m][n]$

```
else:
    arr[i][j] = max(arr[i-1][j], arr[i][j-1])
```

LCS: Practice

[illegible]


```
else:
    arr[i][j] = max(arr[i-1][j], arr[i][j-1])
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LCS: Practice

[illegible]

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LCS: Practice

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LCS: Practice

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else:
    arr[i][j] = max(arr[i-1][j], arr[i][j-1])
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LCS: Practice

[illegible]

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if X[i] == Y[j]:
    arr[i][j] = arr[i-1][j-1] + 1
else:
    arr[i][j] = max(arr[i-1][j], arr[i][j-1])
```

LCS: Practice

X: 0 1 2 3 4 5 6 7 8 9 10 11	
Y:	
0	
1	H
2	E
3	R
4	I
5	T
6	A
7	G
8	E

LCS: Practice

```
if X[i] == Y[j]:
    arr[i][j] = arr[i-1][j-1] + 1
else:
    arr[i][j] = max(arr[i-1][j],
```

LCS: Practice

[illegible]

```
if X[i] == Y[j]:
    arr[i][j] = arr[i-1][j-1] + 1
else:
    arr[i][j] = max(arr[i-1][j], arr[i][j-1])
```

LCS: Practice

		X: 0 1 2 3 4 5 6 7 8 9 10 11											
Y:			G	E	O	R	G	I	A	T	E	C	H
0		0	0	0	0	0	0	0	0	0	0	0	0
1	H	0	→ 0	→ 0	→ 0	→ 0	→ 0	→ 0	→ 0	→ 0	→ 0	→ 0	1
2	E	0	→ 0	↓ 1	→ 1	→ 1	→ 1	→ 1	→ 1	→ 1	↓ 1	→ 1	→ 1
3	R	0	0	↓ 1	→ 1	↓ 2	→ 2	↓ 2	→ 2	→ 2	→ 2	→ 2	→ 2
4	I	0	0	↓ 1	↓ 1	↓ 2	↓ 2	↓ 3	↓ 3	↓ 3	↓ 3	↓ 3	↓ 3
5	T	0	0	↓ 1	↓ 1	↓ 2	↓ 2	↓ 3	↓ 3	↓ 4	↓ 4	↓ 4	↓ 4
6	A	0	↓ 0	↓ 1	↓ 1	↓ 2	↓ 2	↓ 3	↓ 4	↓ 4	↓ 4	↓ 4	↓ 4
7	G	0	↓ 1	↓ 1	↓ 1	↓ 2	↓ 3	↓ 3	↓ 4	↓ 4	↓ 4	↓ 4	↓ 4
8	E	0	↓ 1	↓ 2	↓ 2	↓ 2	↓ 3	↓ 3	↓ 4	↓ 4	↓ 5	↓ 5	5


```
if X[i] == Y[j]:
    arr[i][j] = arr[i-1][j-1] + 1
else:
    arr[i][j] = max(arr[i-1][j], arr[i][j-1])
```

LCS: Practice

X:	0	1	2	3	4	5	6	7	8	9	10	11
Y:		G	E	O	R	G	I	A	T	E	C	H
0		0	0	0	0	0	0	0	0	0	0	0
1	H	0	0	0	0	0	0	0	0	0	0	1
2	E	0	0	1	1	1	1	1	1	1	1	1
3	R	0	0	1	1	2	2	2	2	2	2	2
4	I	0	0	1	1	2	2	3	3	3	3	3
5	T	0	0	1	1	2	2	3	3	4	4	4
6	A	0	0	1	1	2	2	3	4	4	4	4
7	G	0	1	1	1	2	3	3	4	4	4	4
8	E	0	1	2	2	2	3	3	4	4	5	5

Length of LCS: 5

To find the LCS:

1. Start at bottom right corner.
2. Go left until the value at a cell came from the upper row.
3. Go up to that cell, save the common character if it came from the upper left.
4. Repeat step 2.

Can there be multiple longest common subsequences?

YES.

LEETCODE PROBLEMS

1143. Longest Common Subsequence

70. Climbing Stairs

322. Coin Change

FINAL EXAM REVIEW

Jeopardy

Socrative: CS1332

Practice exams in Canvas: Files -> Resources -> Recitation
Materials -> Recitation Practice Exams



Any questions?

Name
Office Hours
Contact

Name
Office Hours
Contact



*Let us know if there is anything specific you want out of
recitation!*