Concurrent GCs ZGC & Shenandoah





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JVM Tuning Expert



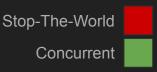
Parallel

Stop-The-World

Concurrent

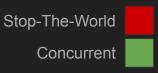
Young Generation Old Generation

Copy Mark Compact



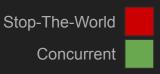
	Young Generation	Old Generation	
Parallel	Сору	Mark	Compact
CMS	Сору	Conc Mark	Conc Sweep





	Young Generation	Old Ger	Old Generation	
Parallel	Сору	Mark	Compact	
CMS	Сору	Conc Mark	Conc Sweep	
G1	Сору	Conc Mark	Compact	





	Young Generation	Old Generation	
Parallel	Сору	Mark	Compact
CMS	Сору	Conc Mark	Conc Sweep
G1	Сору	Conc Mark	Compact
ZGC		Conc Mark	Conc Compact
Shenandoah		Conc Mark	Conc Compact

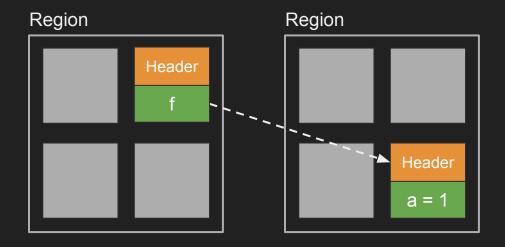


- Concurrent GCs
 - GC runs concurrently with the application

- GC races with Application
 - Marking an object "alive"
 - Compacting / Moving objects

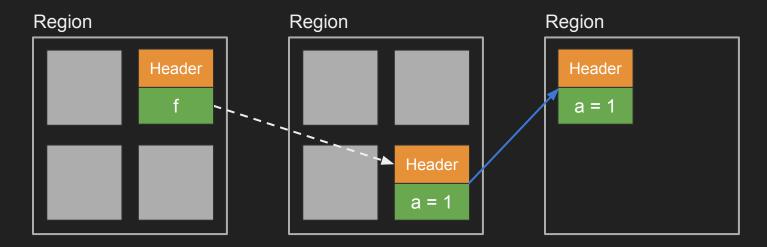


Concurrent copying example

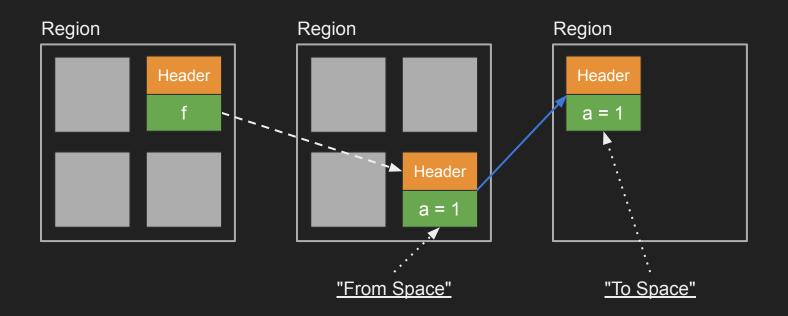




Garbage Collector compacts object







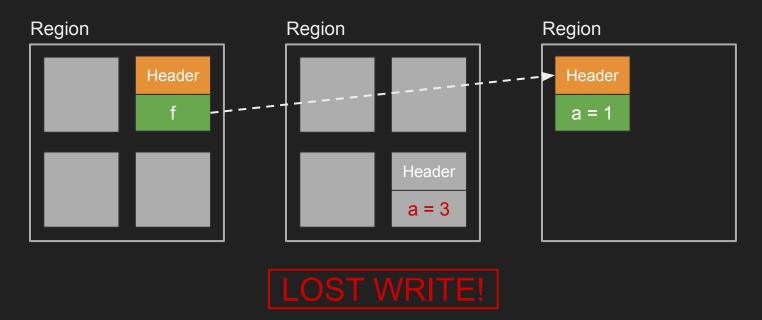


Application writes concurrently





Garbage Collector updates references





- Concurrent GCs store object metadata
 - E.g. whether an object has been marked

Concurrent GCs require JIT support

- JIT injects code that helps the GC
 - GC barriers



- Present in OpenJDK 11+
 - Available in AdoptOpenJDK builds
 - Available in Oracle OpenJDK builds
 - From https://jdk.java.net

Scalable Low Latency GC

Concurrent Compaction, Single Generation



- Only for Linux x86 64-bit
 - No compressed pointers
 - ARM port underway

- Region Based
 - ZPages similar to G1 regions
 - Small (2 MiB object size up to 256 KiB)
 - Medium (32 MiB object size up to 4 MiB)
 - Large (4+ MiB object size > 4 MiB)





STW Mark Start

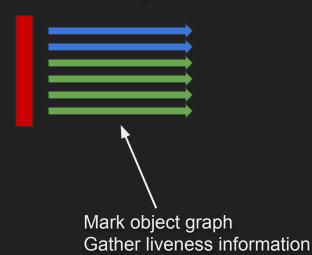


Scan threads stacks
Mark object roots



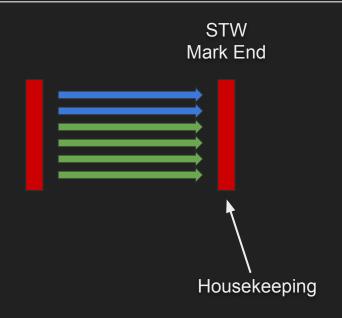


Concurrent Marking













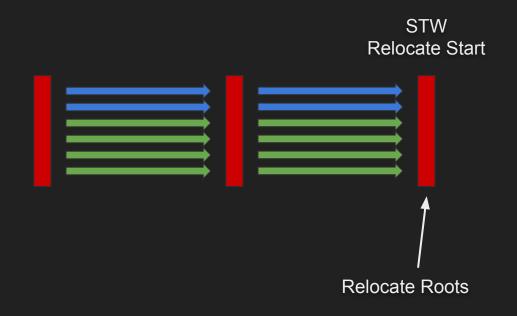
Concurrent Processing



Weak Reference Processing Free Memory Pages Unload Classes Prepare Relocation Set

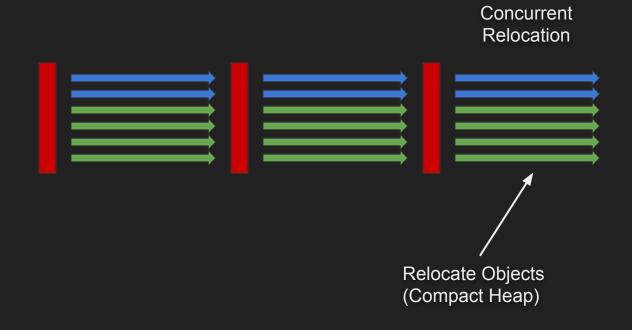








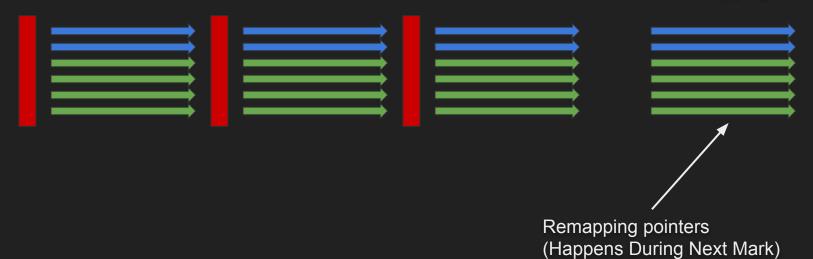








Concurrent Remapping





- ZGC solution
 - To avoid races with the application

Colored Pointers

Load GC barrier



- ZGC colored pointers
 - Color indicates GC metadata

4 color bits

Object address - 44 bits (16 TiB)

Unused - 16 bits

64 bits pointer address

- Check pointer color against GC phase
 - Wrong color => take action



ZGC load GC barrier

```
Object f = obj.field;
<load_barrier>
```



ZGC load GC barrier

```
Object f = obj.field;
if (addr_of(f) & wrong_gc_color) {
    slow_path()
}
```



ZGC load GC barrier

```
Object f = obj.field;
<load_barrier>
```

```
mov 0x10(%rdi), %rsi
test %rsi, 0x20(%r15)
jne slow_path
```



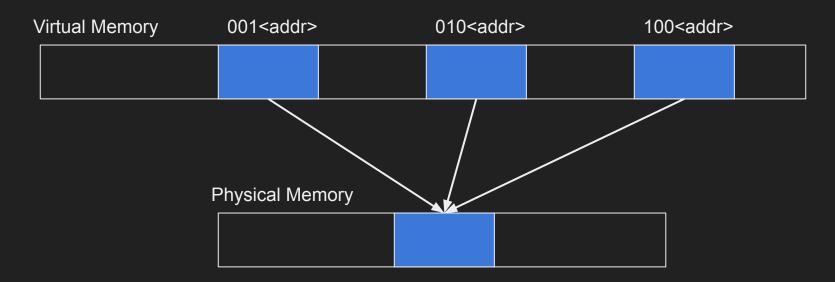
ZGC load GC barrier tests for the right color

- Wrong color -> take slow_path
 - Fix color & Run some action atomically

The action depends on the GC phase



ZGC Multi-Mapping





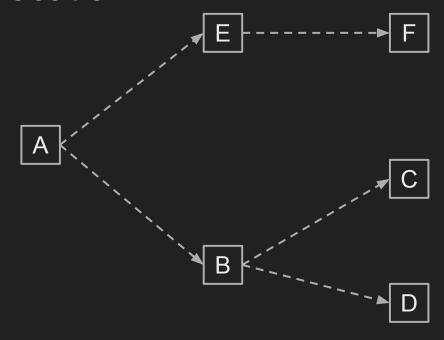
ZGC Multi-Mapping

- RSS shows 3x size
 - E.g. -Xmx=16G -> RSS~=50G
 - On Linux, use smem and track PSS

- Watch out for OOM Killers in clouds
 - o E.g. AWS

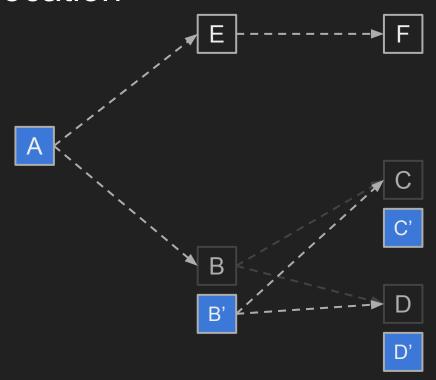


Relocation



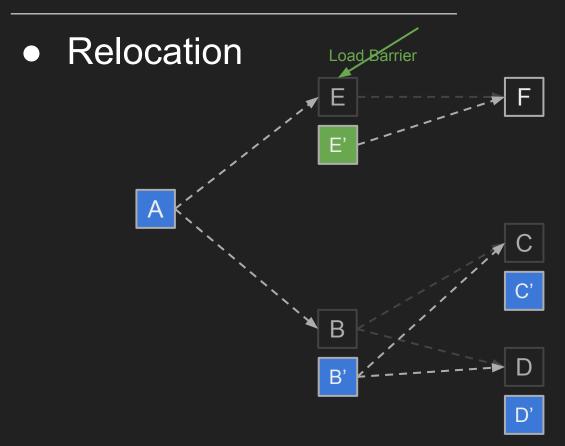


Relocation



MetaData

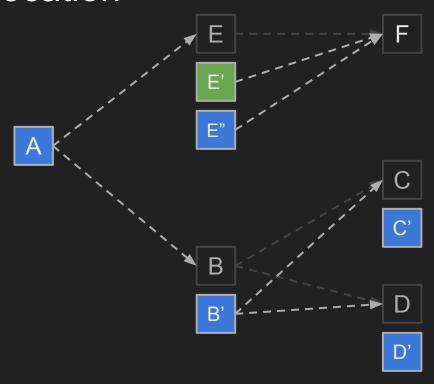




MetaData



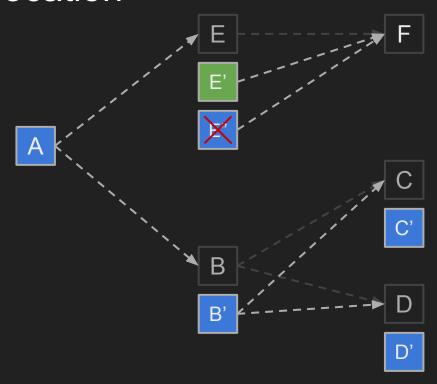
Relocation



MetaData



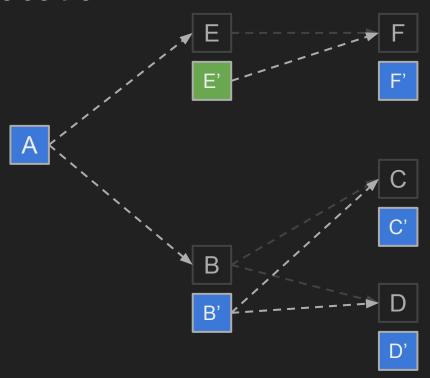
Relocation



MetaData



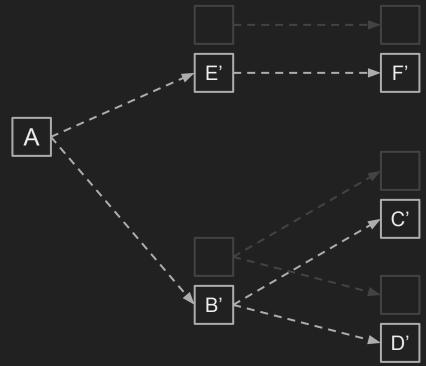
Relocation



MetaData



Remapping





- References
 - Per Lidén, Erik Österlund et al.
 - https://www.youtube.com/watch?v=7cWiwu7kYkE
 - https://www.youtube.com/watch?v=kF_r3GE3zOo
 - Project Wiki
 - https://wiki.openjdk.java.net/display/zgc/Main
 - Mailing List
 - zgc-dev@openjdk.java.net



- Present in OpenJDK 12+ repository
 - Available in AdoptOpenJDK builds
 - NOT available in Oracle's OpenJDK builds
 - Present in RedHat OpenJDK 8, 11+

Scalable low latency GC

Concurrent Compaction, Single Generation



- Available for x86 32-bit and 64-bit
 - ARM ports available

- Linux, MacOS, Windows
 - JDK 8, 11 & Latest

- Region Based
 - Derived from G1





STW Mark Start

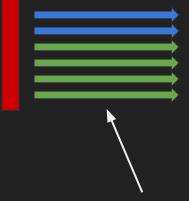


Scan threads stacks Mark object roots





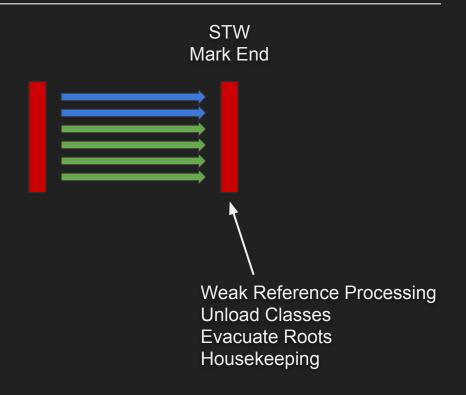




Mark object graph
Gather liveness information

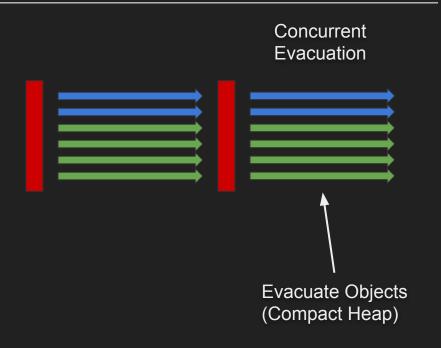




















- ShenandoahGC solution
 - To avoid races with the application

Brooks (Forward) Pointers

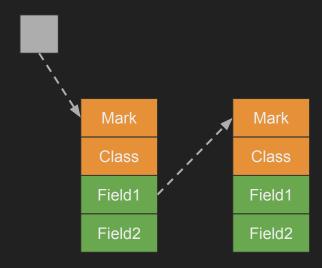
Load & Store GC Barriers



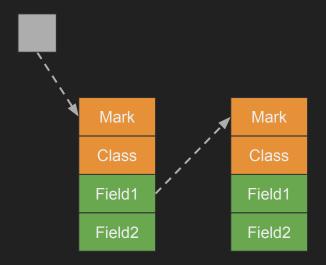
- HotSpot Object Header
 - "Word" == 32 bits on i386; 64 bits on x86-64

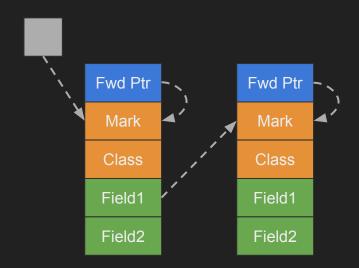
Mark Word	Class Word	Field 1	Field 2	
Mark Word	Class Word	Array Length	Slot 1	Slot 2













- Maintains weak to-space invariant
 - Reads from both from-space to-space copies
 - Writes only to to-space copy



Load GC barrier: dereference forward pointer

```
Object f = obj.field;
<load_barrier>
```



Load GC barrier: dereference forward pointer

```
Object f = obj.field;
f = deref(addr_of(f) - 8);
```



Load GC barrier: dereference forward pointer

```
Object f = obj.field;
<load_barrier>
```

```
mov 0x10(%rdi), %rsi
mov -0x08(%rsi), %rsi
```



- Store GC barrier
 - Runs some action depending on GC phase

```
obj.f = o;
<store_barrier>
```



- Store GC barrier
 - Runs some action depending on GC phase

```
obj.f = o;
if (in_evac_phase &&
    in_collection_set(obj) &&
    !is_forwarded(obj))
    slow_path();
```



- Store GC barrier
 - Runs some action depending on GC phase

```
obj.f = o;
mov 0x3d8(%r15), %r11
test %r11, %r11
...
jne <slow_path>
```

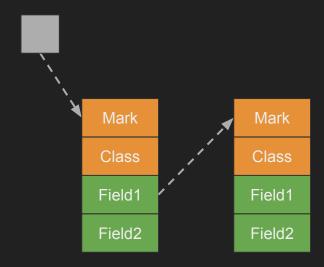


ShenandoahGC 1.0 Problems

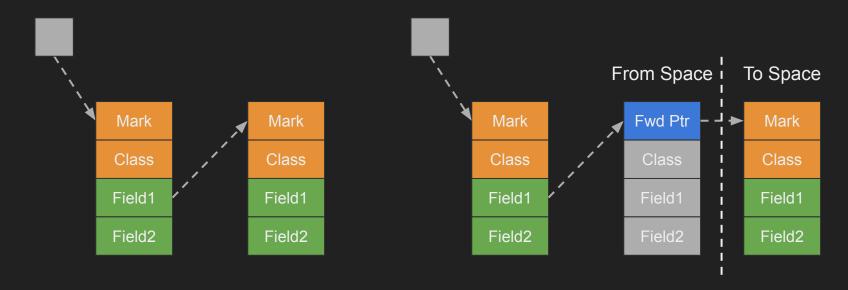
- More memory needed (due to the forward pointer)
 - Worst case 50%, common case 5-10%

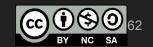
- Complicated (exotic) GC barriers
 - Load & Store GC barriers to maintain invariants











- Maintains a strong to-space invariant
 - Reads only possible from to-space copy
 - Writes only possible from to-space copy



- Load Barrier (LRB Load Reference Barrier)
 - Checks whether there may be forwarded objects
 - And whether in c-set and whether is forwarded

```
Object f = obj.f;
<load_barrier>
```



Load Barrier (LRB - Load Reference Barrier)

```
Object f = obj.f;
if (in_evac_phase &&
    in_collection_set(obj) &&
    !is_forwarded(obj))
    slow_path();
```



Load Barrier (LRB - Load Reference Barrier)

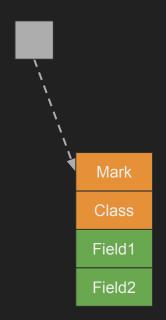
```
Object f = obj.f;
test $0x1,0x20(%r15)
...
jne <slow_path>
```



- Load and Store GC barrier are the same!
 - Big codebase simplification
 - Improved performance

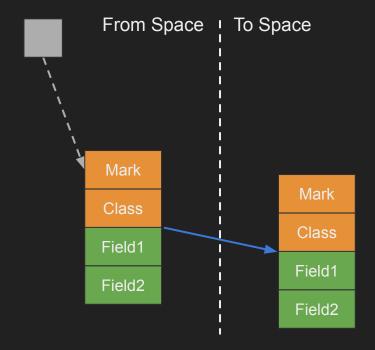


Evacuation



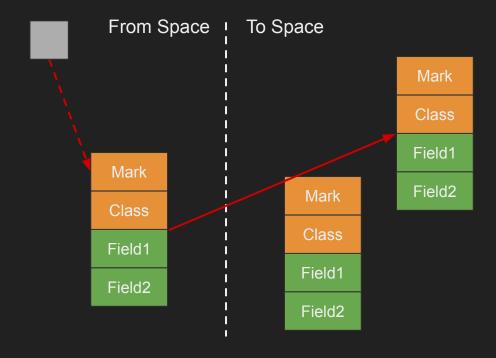


Evacuation (GC Copy)



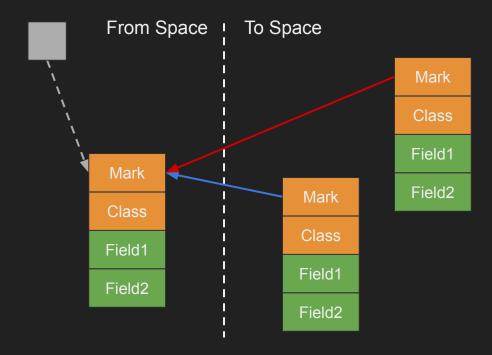


Evacuation (Application triggers LRB)



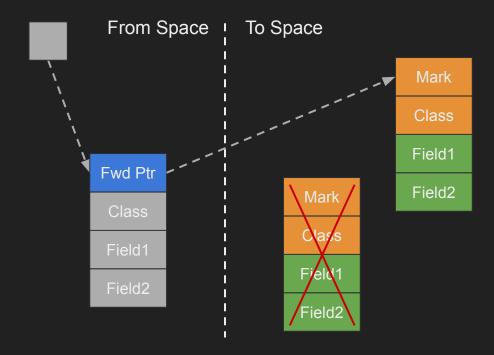


Evacuation (Atomic Fwd Ptr Update)



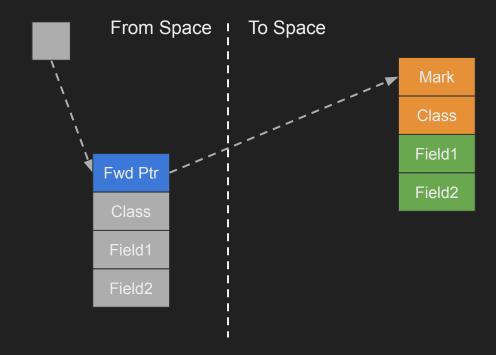


Evacuation (Application Wins)



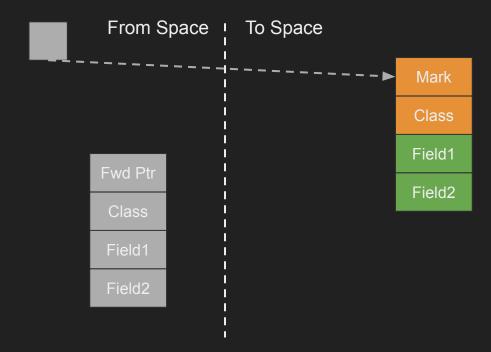


Evacuation (Application Read/Write)





Update References





Shenandoah Heuristics

- Adaptive (default)
 - Maintains free heap amount
- Static
 - Starts GC when tripping thresholds
- Compact
 - Small footprint
- Aggressive
 - Back to back GCs
- Passive
 - Non concurrent, only Full GCs
- Traversal



- References
 - Project Wiki
 - https://wiki.openjdk.java.net/display/shenandoah/Main
 - Aleksey Shipilëv, Roman Kennke et al.
 - https://www.youtube.com/watch?v=VCeHkcwfF9Q
 - https://www.youtube.com/watch?v=E1M3hNlhQCq
 - Mailing List
 - <u>shenandoah-dev@openjdk.java.net</u>



Shenandoah Questions?

Conclusions

Conclusions

Give Concurrent GCs a go

- Stuck with JDK 8?
 - Use Shenandoah

Already using JDK 11+?

ZGC or Shenandoah

Report your feedback!

