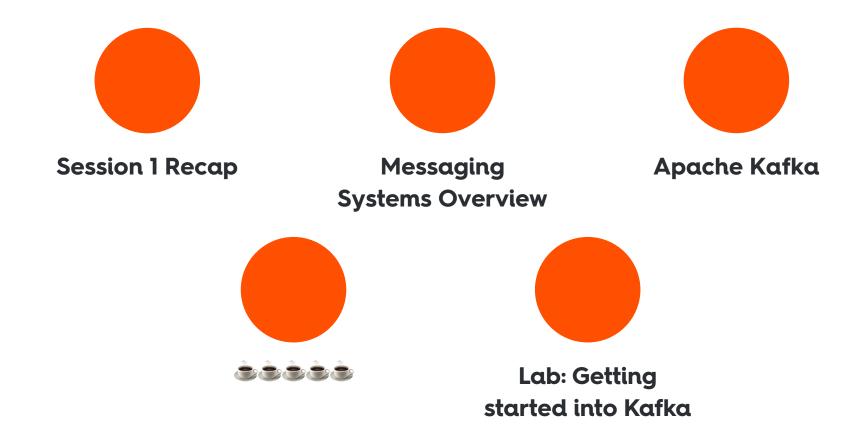
#### **CXITEO**

# Distributed Messaging Systems

Ilyas Toumlilt
i.toumlilt@criteo.com

#### Session 2 - 08/03/2023





# Session 1 Recap



#### Data-driven decision making

Information Theory: Data is a frozen information

Data drives decisions not every day, every second...

Decisions could be made by humans or by software



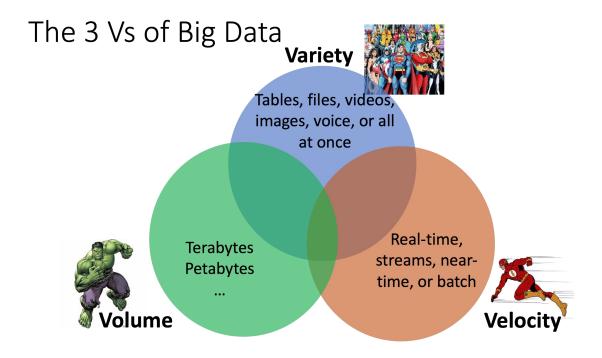
#### Problems of Data-Intensive Distributed Computing

Reliability and Availability constraints

Scalability

Data locality

The programming model





#### **Streaming 101**

Streaming vs Batch Processing

Problems arise in streaming: time, windows, state...

Common distributed system problems: *Dumb ways to die* 



#### Real classic problem

You are the owner of an application (imagine any app, remember flappy bird?) in production

It's tested, up and running with a workload of 100 users

Big buzz came and you are getting +100.000 active users

Your clients say the response time is increasing... 

Where to start looking?



#### Real classic problem

You are the owner of an application (imagine any app, remember flappy bird?) in production

It's tested, up and running with a workload of 100 users

Big buzz came and you are getting +100.000 active users

Your clients say the response time is increasing... 

Where to start looking?





#### Vertical scaling



The simplest approach to scale to higher load is to buy more powerful machines

vertical scaling or scaling up

Many CPUs, many RAM chips, and many disks can be joined together under one operating system, and a fast interconnect allows any CPU to access any part of the memory or disk.

• In this kind of shared-memory architecture, all the components together can be treated as a single machine.

Photo: Designed by pch.vector / Freepik



#### Vertical scaling



The problem with a shared-memory approach is that the cost grows faster than linearly

- A machine with twice as many CPUs, RAM, and disk capacity costs significantly more than twice the price.
- Due to bottlenecks, a machine twice the size cannot necessarily handle twice the load.

A shared-memory architecture may offer limited fault tolerance.











In a shared-nothing architecture multiple machines are used.

horizontal scaling or scaling out

Each node uses its CPUs, RAM, and disks independently.

No special hardware is required.

Any coordination between nodes is done at the software level, through the network.

While a distributed shared-nothing architecture has many advantages, it adds more complexity.

More machines may lead to more failures too.

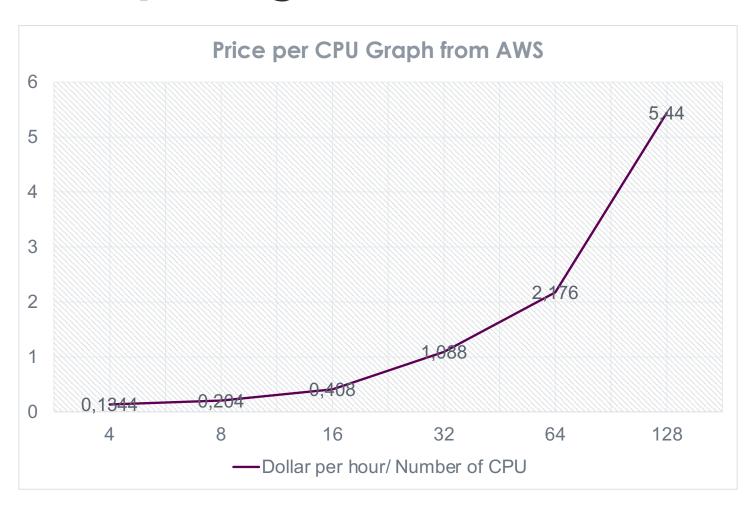


## **Distributed Systems**

- There are several reasons why you might want to distribute your data across multiple machines:
  - Scalability
  - Fault tolerance/high availability
  - Latency
  - Price



# What about pricing





There are two common ways data is distributed across multiple nodes:

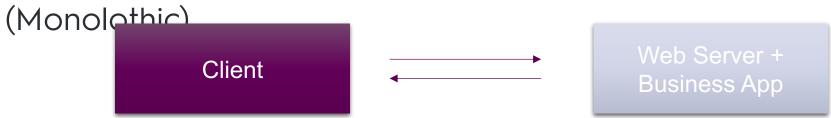
- Replication:
  - Keeping a copy of the same data on several different nodes, potentially in different locations.
  - Replication provides redundancy: if some nodes are unavailable, the data can still be served from the remaining nodes.
  - Replication can also help to improve performance(in general).
- Partitioning (sharding) e.g. partition by Country
  - Splitting a big data into smaller subsets called partitions so that different partitions can be assigned to different nodes.



# How to pass/deliver data in highly distributed environments?



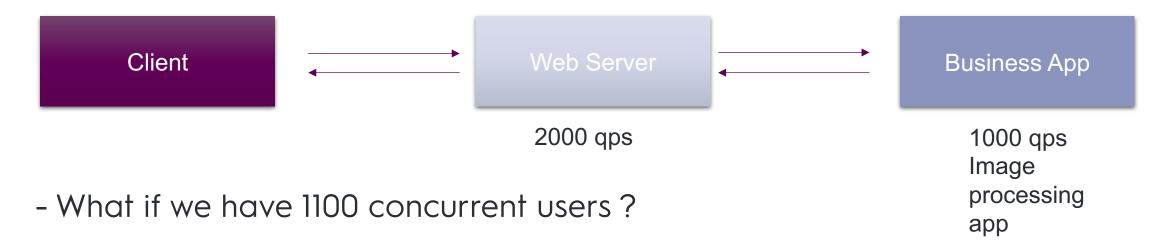
- Let's assume you are a new engineer in a start-up company, and the company provide a face detection software to their clients.
- Very very naive solution: one component does everything?



What is wrong with it?

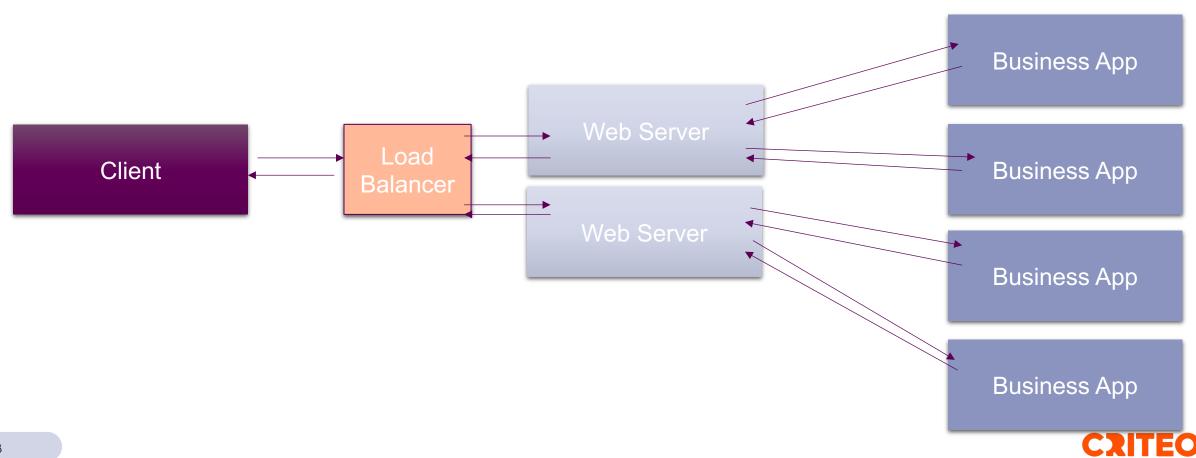


To reduce coupling, we must split core business logic from web server

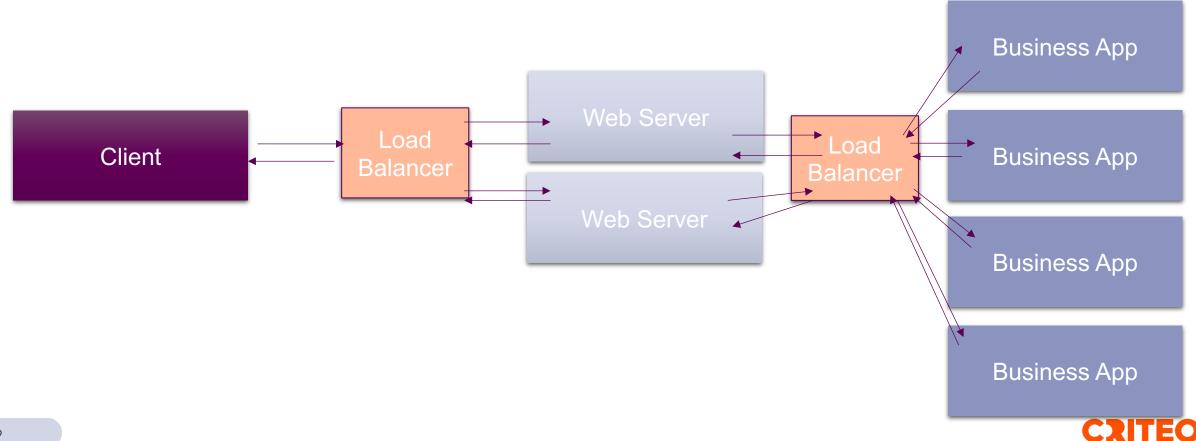


Take home note: low coupling, high cohesion

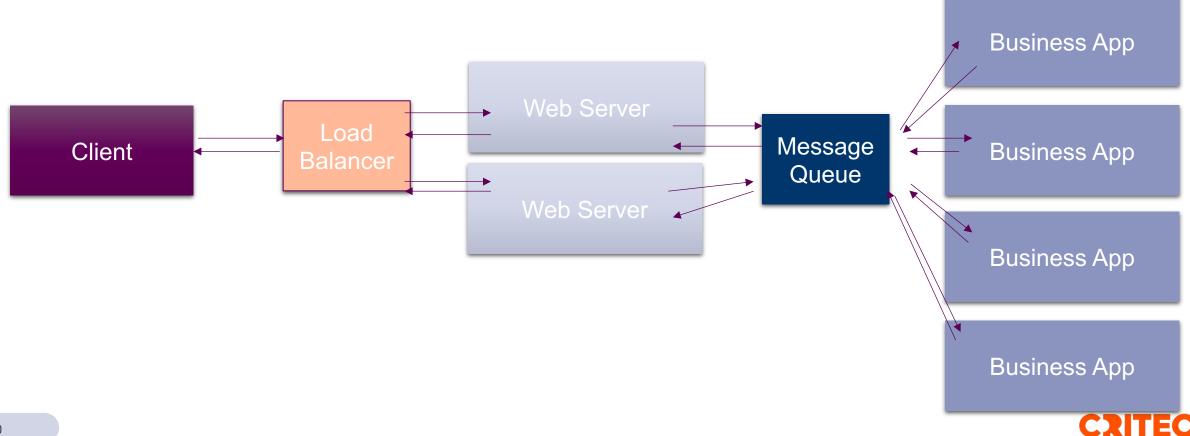




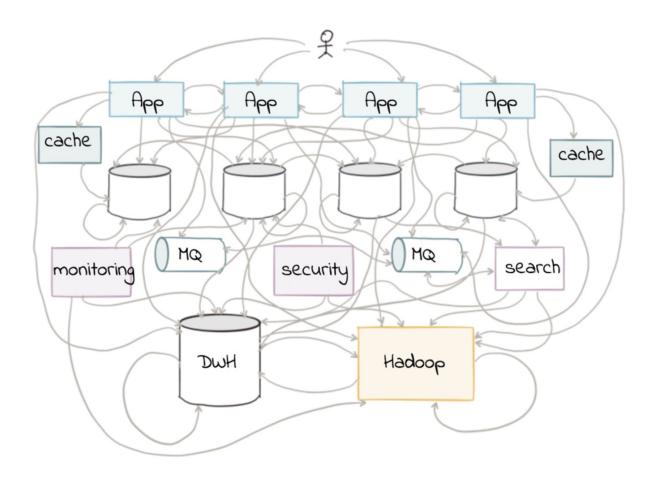
What is wrong?



What can we achieve by taking advantage of messaging queue?

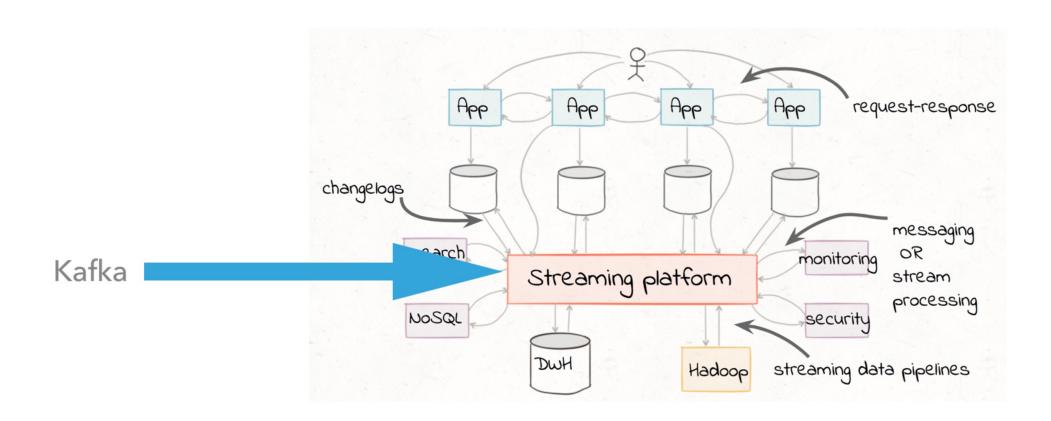


# **Pub/Sub Systems**





# **Pub/Sub Systems**

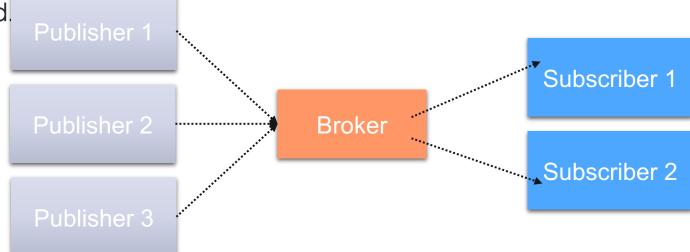




# Messaging(Queue) Systems

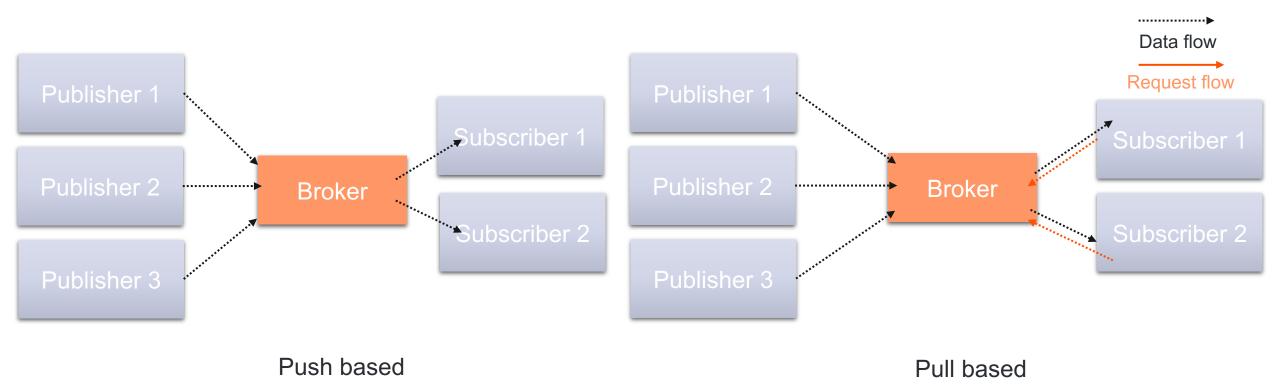
- Publisher sends a piece of data (message) not specifically directing it to a receiver.
- The publisher classifies the message somehow, and that receiver (subscriber) subscribes to receive certain classes of messages.

 Pub/sub systems often have a broker, a central point where messages are published.





## Messaging queue system (Pull vs Push based)





# Apache Kafka



#### What is Kafka?

Apache Kafka is a publish/subscribe distributed messaging system.

Also known as "distributed commit log" "distributing streaming platform".

Data stored in Kafka is durable, in order, and can be read deterministically.

Open-sourced by LinkedIn in 2011.

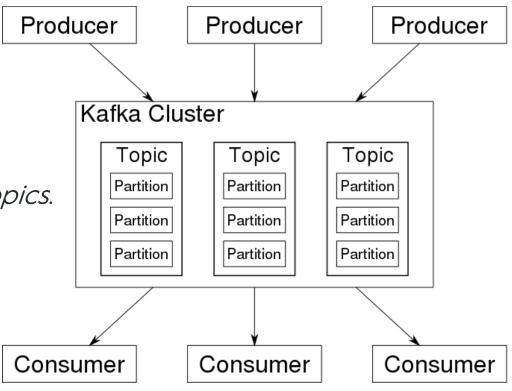
- High-throughput
- Highly distributed
- Fault-tolerant
- Low-latency





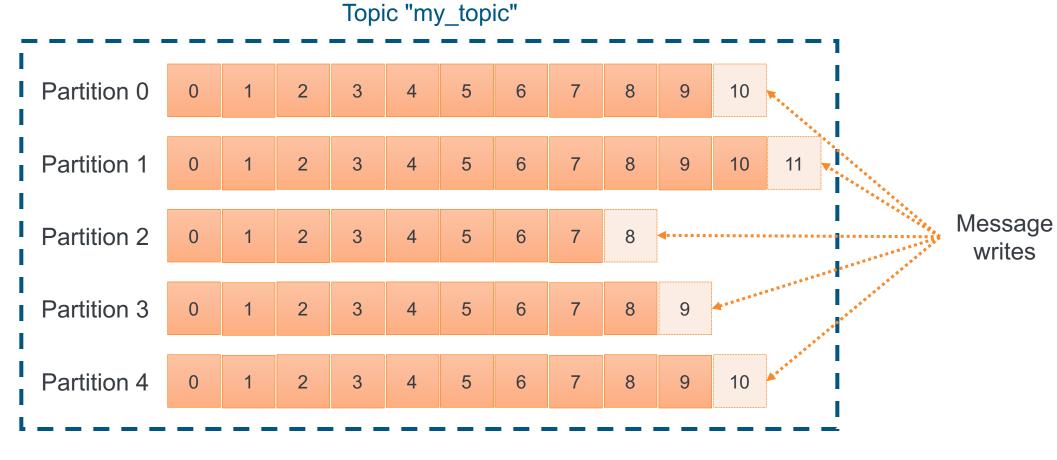
# **Anatomy of Kafka**

- Message is the unit of data.
- Producer creates a new message
- Consumer (i.e., Subscribers/Readers)
- Messages in Kafka are categorized into topics.
- Topics are broken down into a partitions.
- Messages are persistant,
- and written in an append-only fashion.





# **Anatomy of Kafka**





# Message

- The unit of data within Kafka is a message.
- A message is simply an array of bytes.
- A message can have an optional key (also an array of bytes).
- Messages contain a timestamp and optional headers.
- Messages are stored in order within a topic-partition.
- Once in Kafka topic, a message has an offset.
- For efficiency, messages are written into Kafka in batches(collection of messages).

offset key	otional y/value timestamp eaders	key	value
------------	----------------------------------------	-----	-------



#### **Topics and Partitions**

- Messages sent by a producer to a particular topic partition will be append in the order they sent.
- However, there is no guarantee of message time-ordering across the entire topic.
- Partitions are the way that Kafka provides redundancy and scalability.
- Each partition can be hosted on a different server.
  - It means that a single topic can be scaled horizontally across multiple servers.
  - Partitions also define the level of paralelism a topic can support.
- Each partition has one leader and zero or more server follower servers.



#### **Topics and Partitions**

- Partitions are replicated across the cluster
- Each partition has one leader and zero or more server follower

CON/ORC

```
bin/kafka-topics.sh --describe --bootstrap-server=localhost:9092 --topic xxx_yyy_zzz
```

Topic: xxx_yyy_zzz	Partition: 0	Leader: 19	Replicas: 19,20,17 ISR: 19
Topic: xxx_yyy_zzz	Partition: 1	Leader: 20	Replicas: 20,17,18 ISR: 18, 17, 20
Topic: xxx_yyy_zzz	Partition: 2	Leader: 17	Replicas: 17,18,16 ISR: 18, 16, 17
Topic : xxx yyy zzz	Partition: 3	Leader: 18	Replicas: 18, 16, 19 ISR: 18, 16, 19
Topic : xxx yyy zzz	Partition: 4	Leader: 16	Replicas: 16,19, 20 ISR: 16,19, 20



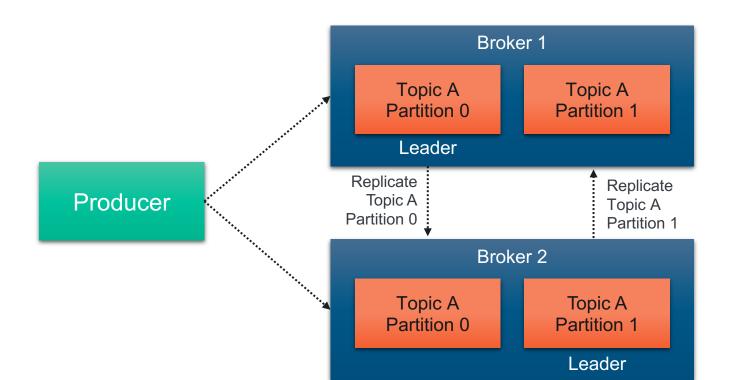
#### Kafka Broker

- A single Kafka server is called a broker.
- When interacting with producers:
  - The broker receives messages from producers;
  - Assigns offsets to them;
  - Replicates messages across other brokers;
  - Commits the messages to storage on disk.
- When serving consumers
  - It responds to fetch requests for partitions with the messages that have been committed.



#### Kafka Cluster

- Kafka brokers are designed to operate as part of a cluster.
- A partition is owned by a single broker in the cluster (leader).
- A partition can be replicated to other brokers in the cluster (replicas)



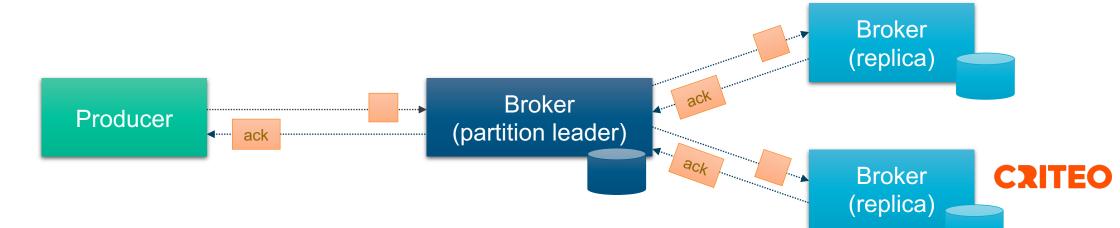


#### Kafka Producer

Producers decide to which partitions to send messages to.

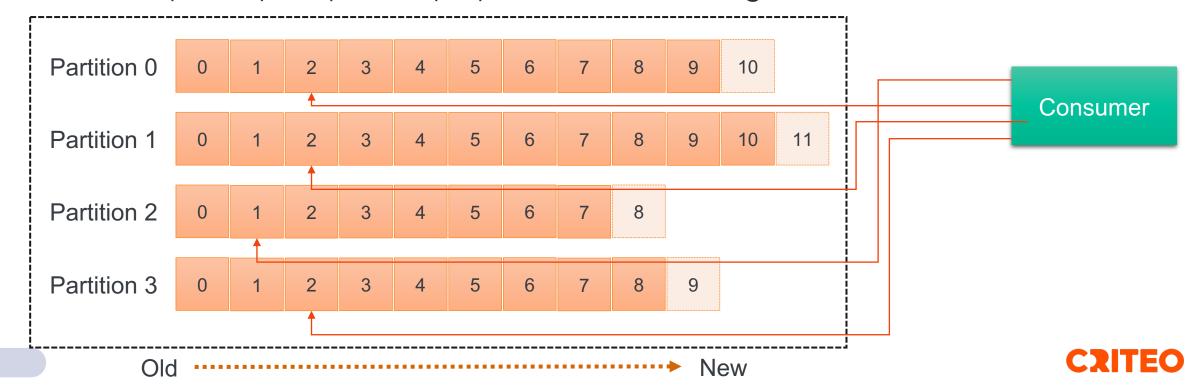
Producers wait for the acknowledgment from the broker:

- ack=0
  - Doesn't wait for confirmation from the partition leader broker.
- ack=1
  - Only waits for the acknowledgment of the leader broker.
- ack=all
  - Waits for the acknowledgments of the leader broker and all replicas.



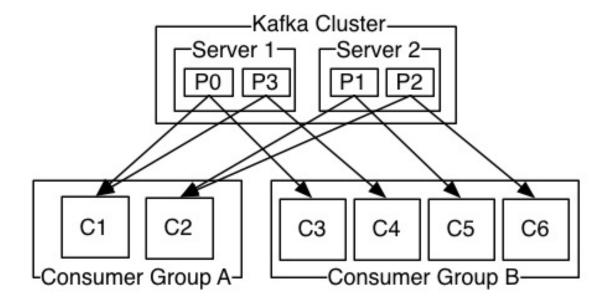
#### Kafka Consumer

- A client that consumes records from a Kafka cluster.
- Kafka consumers transparently handles the failure of Kafka brokers, and transparently adapts as topic partitions it fetches migrate within the cluster.



#### **Consumer Group**

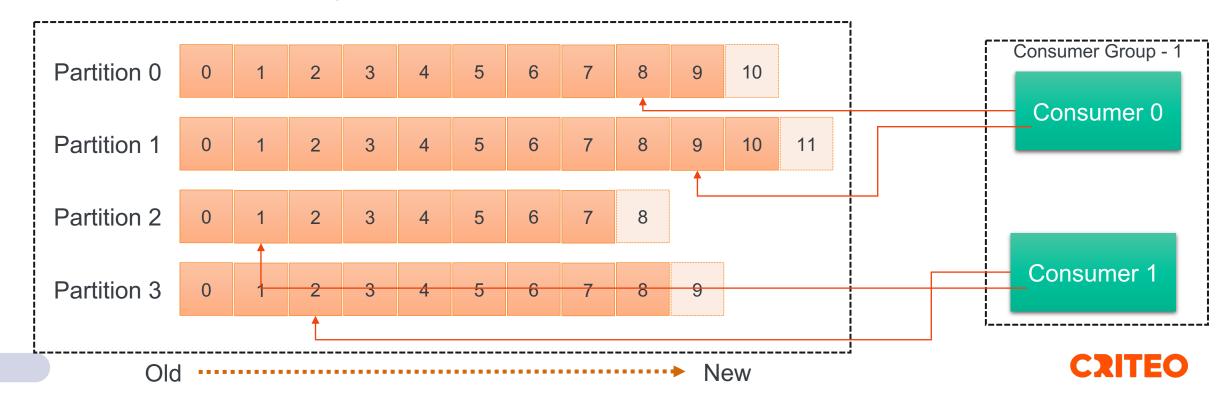
- If all the consumer instances have the same consumer group, then this works as a queue balancer load over the consumers.
- If all the consumer instances have different consumer groups, then this works like publish-subscribe and all messages are broadcast to all consumers.





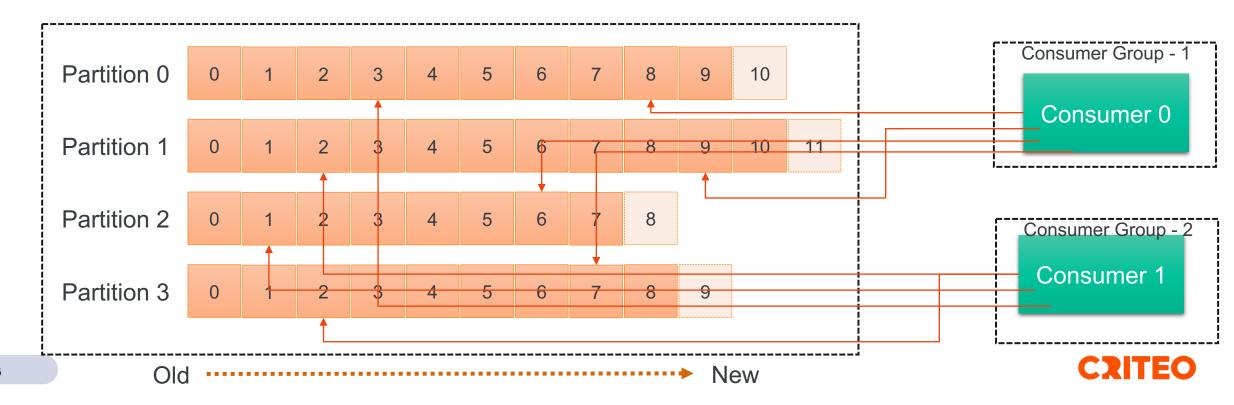
## Kafka Consumer Group

- Each partition can be consumed by only one consumer within a consumer group.
- The consumers organize themselves to balance the load.



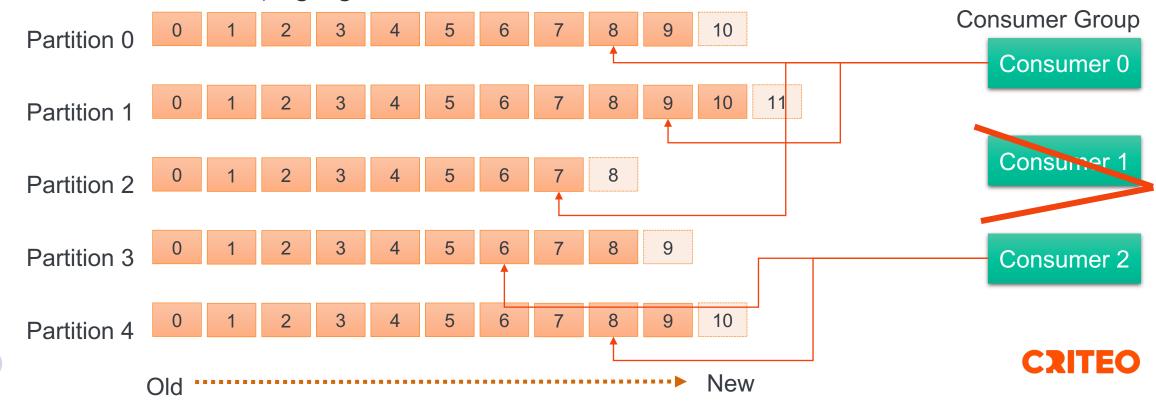
## Kafka Consumer Group

• If each consumer in different group then they don't share workload starting from scratch.



### **Consumer Rebalance**

- If a consumer is removed from the consumer group, the partitions from that consumer will be assigned to the remaining consumers.
- If a new consumer is added to the consumer group, it will receive partitions from other consumers, keeping a good balance.



# **Committing Offsets**

- If a consumer fails within a consumer group, another consumer should take place, starting from the same offsets where the failing consumer stopped.
- In order to save this state somewhere, consumers may send their offsets back to Kafka in a specific topic.
- In case of failure:
  - Committing offsets after processing it may generates duplicates.
  - Committing offsets before processing it, may end up in messages not processed.





### Kafka Guarantees

- Messages sent by a producer to a particular topic partition will be appended in the order they are sent.
  - That is, if a record M1 is sent by the same producer as a record M2, and M1 is sent first, then M1 will have a lower offset than M2 and appear earlier in the log.
- A consumer instance sees records in the order they are stored in the log.
- For a topic with replication factor N, we will tolerate up to N-1 server failures
  without losing any records committed to the log.



### Kafka Semantics

#### At most once

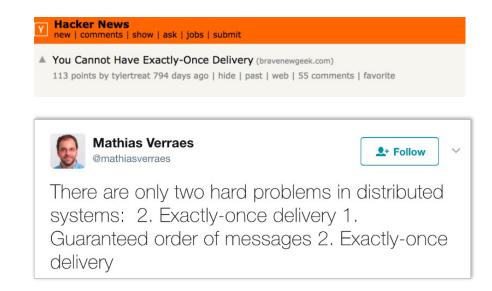
- When producers don't retry when ack times out or
- an error is returned.
- When consumers fail, they may skip some messages.

#### At least once

- When producers retry, they may send messages twice.
- When consumers fail, they may reprocess messages twice.

### Exactly once

- Kafka can deduplicate messages within the same partition (idempotent producer).
- If consumers use a transactional sink, they may commit partitions within the transaction.





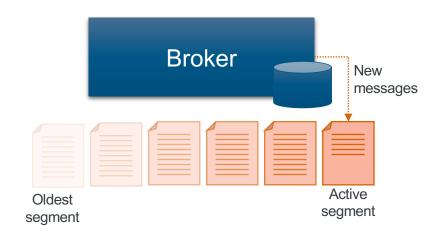
### Retention

- Kafka organizes topic-partition data based on segments.
- It stores the segments for some period of time or until it reaches a certain size.
- Messages are discarded based on retention policy if they were consumed or not.

#### •Property:

log.retention.hours (Broker-wise)
retention.ms(Topic-wise)

log.retention.bytes (Broker-wise) retention.bytes (Max size of a partition)





# Zookeeper

- Apache Zookeper is a distributed coordination service for distributed applications.
- Zookeeper provides some guarantees:
  - Sequential Consistency
  - Atomicity
  - Single System Image
  - Reliability
  - Timeliness
- Kafka uses Apache Zookeeper to keep the cluster state:
  - List of currently members of a cluster;
  - List of topics, partitions leaders and replicas;
  - Current controller; etc.
- Deprecated since Kafka 3.0.0, not required anymore





# Cap Theorem in context of Kafka

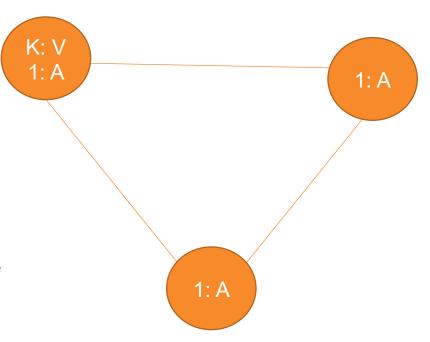
CAP Theorem(published in PODC '02) states that any distributed system can provide at most two out of the three guarantees: Consistency, Availability and Partition tolerance.

#### Partition tolerance(\*necessity)

 The system continues to operate despite an arbitrary number of messages being dropped (or delayed) by the network between nodes

#### Consistency

- Every read receives the most recent write or an error Availability
- Every request receives a (non-error) response, without the guarantee that it contains the most recent write





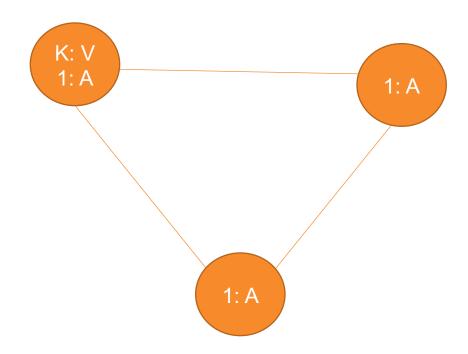
# Cap Theorem on Kafka Brokers

#### For Consistency:

- Producers acks=all
- Replicator factor >= 3
- Disable unclean leader election
- Min insync replicas=replicatorFactor -1

### For Availability:

- Enable unclean leader election
- Min insync replicas=1





# What about a break?



# **Practice time**

