```
module GaleShapley where
open import Agda.Builtin.String
open import Agda. Builtin. Equality
open import Data. Empty
open import Data. Maybe
open import Data.Nat
open import Data.Nat.Properties
open import Data.Product
open import Data.Bool
open import Data.List
open import Data.Sum
open import Relation. Nullary
open import Induction. WellFounded
open import Relation. Binary. Propositional Equality
import Data.Nat.Solver
open Data.Nat.Solver.+-*-Solver
  using (prove; solve; _:=_; con; var; _:+_; _:*_; :-_; _:-_)
-- lem2 : (4 + 6 10)
-- lem2 = solve 0 (con 4 :+ con 6 := con 10) refl
lem3: (x:) \rightarrow (2 * (x + 4) 8 + 2 * x)
lem3 = solve 1 (\lambda x' \rightarrow \text{con 2}:* (x':+ con 4) := con 8:+ con 2:* x') refl
infix 3 _>just_
infix 4 _from>_
lengthPrefs : List ( \times List ) \rightarrow
lengthPrefs | = 0
lengthPrefs ((_ , l) xs) = length l + lengthPrefs xs
```

 $compSumPrefLists\ freeMen\ engagedMen = lengthPrefs\ freeMen + lengthPrefs\ engagedMen$

```
record MatchingState: Set where
constructor mkState
field
men: List ( × List )
freeMen: List ( × List )
engagedMen: List ( × List )
women: List ( × List )
couples: List ( × List )
sumPrefLists:
sumEq: sumPrefLists lengthPrefs freeMen + lengthPrefs engagedMen
```

 $\mathsf{compSumPrefLists}: (\mathit{freeMen\ engagedMen}: \mathsf{List}\ (\ \times\ \mathsf{List}\)\) \to$

```
-- Code by @yeputons on Stack Overflow :D
is- : \rightarrow \rightarrow Bool
is- a b with a ? b
... | Dec.yes _ = true
... | Dec.no _ = false
-- Helper function to search for an element in a list of natural numbers.
positionInListHelper : \rightarrow List \rightarrow \rightarrow Maybe
positionInListHelper x \parallel \underline{\phantom{a}} = \text{nothing} --searched everywhere but couldn't find it!
positionInListHelper x (xs xss) zero with compare x xs
... | equal _ = just zero --found man at the tip of the list!
\dots \mid \underline{\quad} = \mathsf{positionInListHelper} \; x \; xss \; (\mathsf{suc} \; \mathsf{zero}) \; --\mathsf{accumulate} \; \mathsf{and} \; \mathsf{keep} \; \mathsf{searching} \dots
positionInListHelper x (xs xss) (suc n) with compare x xs
... | equal \underline{\phantom{a}} = just (suc n) --found man somewhere in the list
             = positionInListHelper x xss (suc (suc n)) --accumulate and keep searching
positionInList : \rightarrow List \rightarrow Maybe
positionInList n ns = positionInListHelper n ns zero
-- m is the proposing man
-- h is the current husband of the woman
-- prefs is the preference list of the woman
-- returns true if the woman prefers m over h
propose: (m:)(h:)(prefs: List) \rightarrow Bool
-- Woman compares
propose man h preferenceList with positionInList man preferenceList | positionInList h preferenceList
... | just p | just q = is- p q --does she prefer the new guy to the current one?
... | just p | nothing = false --shouldn't happen in an ideal world : woman received a propose
... \mid nothing \mid just q= false --shouldn't happen in an ideal world : woman married an unknown
... | nothing | nothing = false --shouldn't happen in an ideal world : who are these men??
-- We assume the "husbands" to be the proposing side,
-- therefore if women propose the pair looks like (wife, husband)
getHusband : \rightarrow List (\times) \rightarrow Maybe
getHusband woman [] = nothing --not married yet and this is the first proposal in the algo-
getHusband woman ((m, w)) []) with compare woman w
... | equal _{-} = just m --found your husband!
... | _ = nothing --not married yet
getHusband woman ((m, w) (c cs)) with compare woman w
... | equal \underline{\phantom{a}} = just m --found your husband!
\dots \mid \underline{\quad} = \operatorname{getHusband} woman (c \ cs) -- \operatorname{keep} searching
\mathsf{getWife}:\ \to \mathsf{List}\ (\ \times\ ) \to \mathsf{Maybe}
getWife man [] = nothing
getWife man((m, w)) with compare man m
... | equal _{-} = just w --found your wife!
... | _ = nothing --not married yet
```

```
getWife man((m, w) (c cs)) with compare man m
... | equal _{-} = just w --found your wife!
         = getHusband man(c cs) --keep searching
-- Simply extract a preference list from the scheme of indexed lists.
getPreferenceList : \rightarrow List ( \times List ) \rightarrow List
getPreferenceList person [] = [] --dummy case
getPreferenceList person ((p, preferences) ps) with compare person p
\dots \mid equal \_ = preferences
\dots \mid \underline{\quad} = \mathsf{getPreferenceList} \ person \ ps
-- Safely adding couples : previous marriages are unmade
safeAddNewCouple : (newCouple : \times)(previousCouples : List (\times)) \rightarrow List (\times)
safeAddNewCouple (m, w) [] = (m, w) []
safeAddNewCouple (m, w) ((a, b) []) with compare w b
\dots | equal \underline{\phantom{a}} = (m, w) []
\dots \mid \underline{} = (m, w) (a, b)
safeAddNewCouple (m, w) ((a, b) (c cs)) with compare w b
... | equal \underline{\phantom{a}} = (m, w) c cs
\dots \mid \underline{\hspace{0.5cm}} = (a \text{ , } b) \text{ safeAddNewCouple } (m \text{ , } w) \text{ } (c \text{ } cs)
\mathsf{p}:(\mathit{l}:\mathsf{List}\;(\;\times\;))(a:\;) \to a\;\; \mathsf{3} \to \mathsf{safeAddNewCouple}\;(\mathsf{2}\;,\;\mathsf{3})\;((\mathsf{1}\;,\;a)\;\;\mathit{l})\;\;((\mathsf{2}\;,\;a)\;\;\mathit{l})
p [] a e with compare 3 a
p [] zero () | w
\mathbf{p} \; [] \; (\mathsf{suc} \; \mathsf{zero}) \; () \; | \; w
p [] (suc (suc zero)) () | w
p [] (suc (suc (suc zero))) e | equal .3 = refl
p [] (suc (suc (suc (suc a)))) () | w
p(x l) zero()
p(x l) (suc .2) refl = refl
-- Safely adding new engaged men to the list : dumped man is removed
safeAddNewEngagedMan : (newEngagedMan : ( \times List ))(prevFiance : )(prevEngagedMen : List ( \times List ))
-- Dummy case: this function is only invoked if a woman is already married, so the list
safeAddNewEngagedMan (newFiance, prefs) prevFiance [] = ((newFiance, prefs), [], (0, []))
{\sf safeAddNewEngagedMan} \ (\mathit{newFiance} \ , \ \mathit{prefs}) \ \mathit{prevFiance} \ ((\mathit{m} \ , \ \mathit{prefsM}) \ \ []) \ \mathsf{with} \ \mathsf{compare} \ \mathit{prevFiance} \ \mathit{m}
... | equal \underline{\phantom{a}} = ((newFiance, prefs) [], (m, prefsM)) --kick him out!
... |  = (((newFiance, prefs) (m, prefsM) []), (0, [])) --safe to keep after all
safeAddNewEngagedMan (newFiance, prefs) prevFiance ((m, prefsM) ms engagedMen) with compare p
... | equal _ = ((newFiance, prefs) ms engagedMen, (m, prefsM)) --kick him out!
... | \_ = (m, prefsM) (proj (safeAddNewEngagedMan (newFiance, prefs) prevFiance (ms. engagedMan (newFiance))
step : MatchingState → MatchingState
-- When there are no more free men, the matching is stable and this is the last step.
step (mkState men [] engagedMen women couples <math>k p) = mkState men [] engagedMen women couples <math>k p
```

```
-- Dummy case : the function shouldn't really be invoked with a man with empty preferer
-- But otherwise Agda would question the completeness of our pattern matching.
step (mkState men((n, []) freeMen) engagedMen women couples k p) = mkState men((n, []) freeMen
-- Proposal step
step (mkState men ((n, w prefs) freeMen) engagedMen women couples k p) with getHusband w couples
... | just h with propose n h (getPreferenceList w women) --Woman has a husband, represented by hi
... | true = mkState men freeMenUpdated engagedMenUpdated women (safeAddNewCouple (n, w) couple
                      freeMenUpdated = proj_2 (safeAddNewEngagedMan (n, prefs) h engagedMen) freeMen - 0
                      engagedMenUpdated = proj (safeAddNewEngagedMan (n, prefs) h engagedMen)
... | false = mkState men((n, prefs), freeMen) engagedMen women couples (compSumPrefLists ((n, pref.
-- Woman didn't have a husband yet (represented by zero) : must accept proposal
step (mkState men ((n, w prefs) freeMen) engagedMen women couples k p) | nothing = mkState men fi
+-zero : n k \rightarrow k n + 0 \rightarrow k n
+-zero zero zero p=p
+-zero zero (suc k) p = p
+-zero (suc n) zero ()
+-zero (suc n) (suc .(n + 0)) refl = cong suc (+-right-identity n)
+-move-zero : n \ k \rightarrow k \ n+0 \rightarrow n \ k+0
+-move-zero zero zero refl = refl
+-move-zero zero (suc k) ()
+-move-zero (suc n) zero ()
+-move-zero (suc n) (suc n) (suc n) (suc n) refl = cong suc (+-move-zero n) refl
nn+m: m n \rightarrow n n + m
nn+m zero zero = -refl
nn+m zero (suc m) = -reflexive (cong suc (sym (+-right-identity m)))
nn+m (suc n) zero = zn
nn+m (suc n) (suc m) = ss (nn+m (suc n) m)
\mathsf{lengthPrefsOneSide}: \ (\mathit{freeMen\ engagedMen}: \ \mathsf{List}\ (\ \times\ \mathsf{List}\ ))(k:\ ) \to k\ \ \mathsf{compSumPrefLists}\ \mathit{freeMen\ engagedMen}
lengthPrefsOneSide [][]kp = zn, zn
[ (fst, [) engagedMen) k p = zn, proj_2 (lengthPrefsOneSide [] engagedMen k p) ]
lengthPrefsOneSide [] ((fst, x snd) engagedMen) k p = zn, -reflexive (sym p)
lengthPrefsOneSide ((fst, []) freeMen) [] k p = -reflexive (+-zero k (lengthPrefs freeMen) (+-move-zero (1
lengthPrefsOneSide ((fst, x snd) freeMen) [] k p = -reflexive (+-zero k (lengthPrefs ((fst, x snd) freeM
lengthPrefsOneSide\ ((fst\ ,\ [])\ freeMen)\ ((fst\ ,\ [])\ engagedMen)\ k\ p = lengthPrefsOneSide\ freeMen\ engagedMen
lengthPrefsOneSide\ ((fst\ ,\ [])\ freeMen)\ ((fst\ ,\ x\ snd)\ engagedMen)\ .(lengthPrefs\ freeMen\ +\ suc\ (lengthPrefs\ freeMen\ +\ suc\ freeMen\ +\ suc\ freeMen\ +\ suc\ freeMen\ +\ suc\ freeMen\ freeMen\
lengthPrefsOneSide\ ((fst\ ,\ snd)\ freeMen)\ ((fst\ ,\ [])\ engagedMen)\ .(suc\ (lengthPrefs\ ((fst\ ,\ snd)\ freeMen)\ .(fst\ ,\ snd)\ freeMen)\ .
lengthPrefsOneSide ((fst, x snd) freeMen) ((fst, x snd) engagedMen) p refl = nn+m (suc (lengthPrefs
stepsWithPrefs : m \ n \rightarrow m \ n \rightarrow m \ 1 + n
stepsWithPrefs zero zero zn = zn
```

```
stepsWithPrefs zero (suc n) zn = zn
stepsWithPrefs (suc m) zero ()
stepsWithPrefs (suc m) (suc n) (ss p) = ss (stepsWithPrefs m n p)
lengthPrefsExtLemma: (x : \times List)(xs : List (\times List))(n : ) \rightarrow lengthPrefs xs n + lengthPrefs (x : List)(xs : List (\times List))(n : ) \rightarrow lengthPrefs xs n + lengthPrefs (x : List)(xs : List)(xs : List)(xs : List)(n : ) \rightarrow lengthPrefs xs n + lengthPrefs (x : List)(xs : List)(xs : List)(xs : List)(n : ) \rightarrow lengthPrefs xs n + lengthPrefs (x : List)(xs : List)
-- Proofs! :D
 -- The type of elements that belong to a list.
data (A : Set)(a : A) : List A \rightarrow Set where
       \mathsf{now} : (as : \mathsf{List}\ A) \to a \ (a \ as)
       later : \{a': A\}\{as: \mathsf{List}\ A\} \to a \ as \to a \ (a'\ as)
 -- Bigger than comparison for Maybe -typed elements.
data \_>just\_: Maybe \rightarrow Maybe \rightarrow Set where
       \_\mathsf{from}{>}\_: \ \{m\ n:\ \} \to m > n \to \mathsf{just}\ m > \mathsf{just}\ n
- Given a man and a woman and their preferences, the condition of stability is satisfied if - another m' and v
\mathsf{leftinv} : (a:) \to \mathsf{zero} + a \ a
\mathsf{leftinv}\ a = \mathsf{refl}
\mathsf{rightinv} : (a:) \to a + \mathsf{zero} \ \ a
rightinv zero = refl
rightinv (suc a) = cong suc (rightinv a)
```