COL352: Assignment 2

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1 Question 1

We say that a context-free grammar G is self-referential if for some non-terminal symbol X we have $X \to^* \alpha X \beta$, where $\alpha, \beta \neq \varepsilon$. Show that a CFG that is not self-referential is regular.

Prove that the class of context-free languages is closed under intersection with regular languages. That is, prove that if L_1 is a context-free language and L_2 is a regular language, then $L_1 \cap L_2$ is a context-free language. Do this by starting with a DF

Given two languages L, L', denote by

$$L||L' := \{x_1y_1x_2y_2\dots x_ny_n \mid x_1x_2\dots x_n \in L, y_1y_2\dots y_n \in L'\}$$

Show that if L is a CFL and L' is regular, then L||L' is a CFL by constructing a PDA for L||L'. Is L||L' a CFL if both L and L' are CFLs? Justify your answer.

For $A \subseteq \Sigma^*$, define

$$cycle(A) = \{yx \mid xy \in A\}$$

For example if $A = \{aaabc\}$, then

$$cycle(A) = \{aaabc, aabca, abcaa, bcaaa, caaab\}$$

Show that if A is a CFL then so is cycle(A)

 \mathbf{Let}

$$A = \{wtw^R \mid w, t, \in \{0,1\}^* \ \text{ and } \ |w| = |t|\}$$

Show that A is not a CFL.

Prove the following stronger version of pumping lemma for CFLs: If A is a CFL, then there is a number k where if s is any string in A of length at least k then s may be divided into five pieces s = uvxyz, satisfying the conditions:

- 1. for each $i \geq 0$, $uv^i x y^i z \in A$
- **2.** $v \neq \varepsilon$, and $y \neq \varepsilon$, and
- **3.** $|vxy| \le k$.

Give an example of a language that is not a CFL but nevertheless acts like a CFL in the pumping lemma for CFL (Recall we saw such an example in class while studying pumping lemma for regular languages).