

CS 3205 COMPUTER NETWORKS

JAN-MAY 2019

LECTURE 10: 6TH FEB 2019

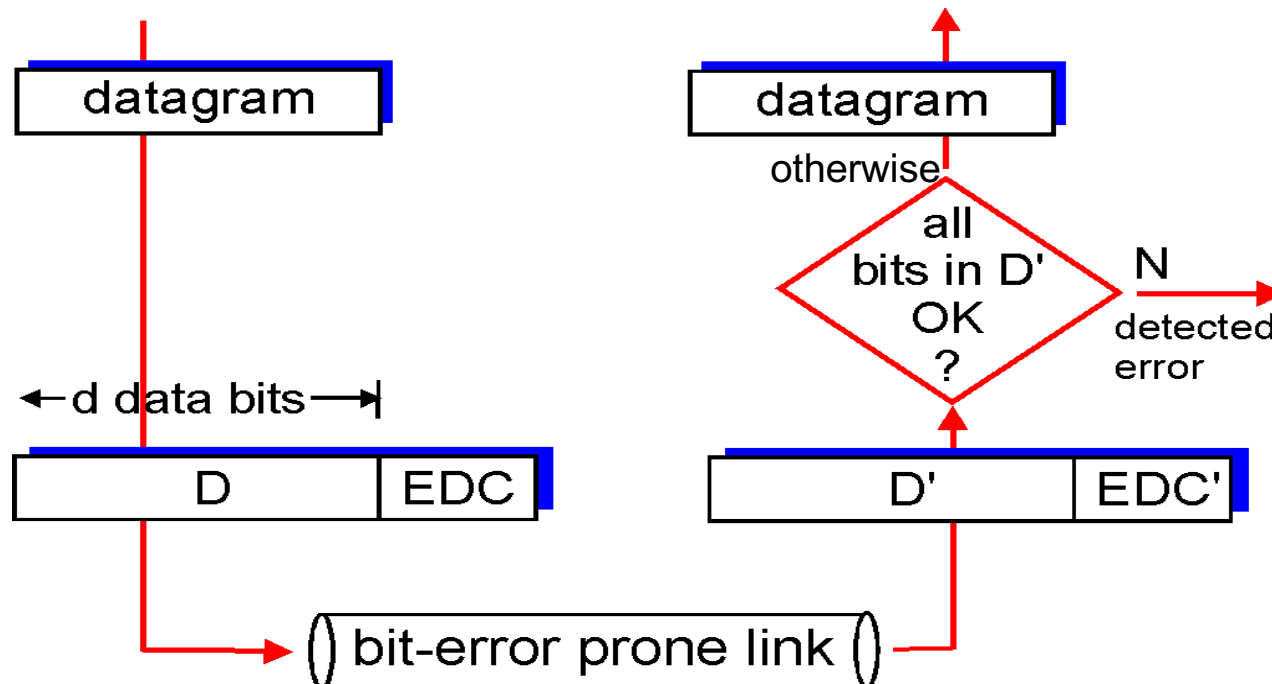
Error Detection and Correction discussed in class is covered in **Section 5.2**, Computer Networking – A top-down approach, Kurose and Ross, 6th Edition.

Error detection

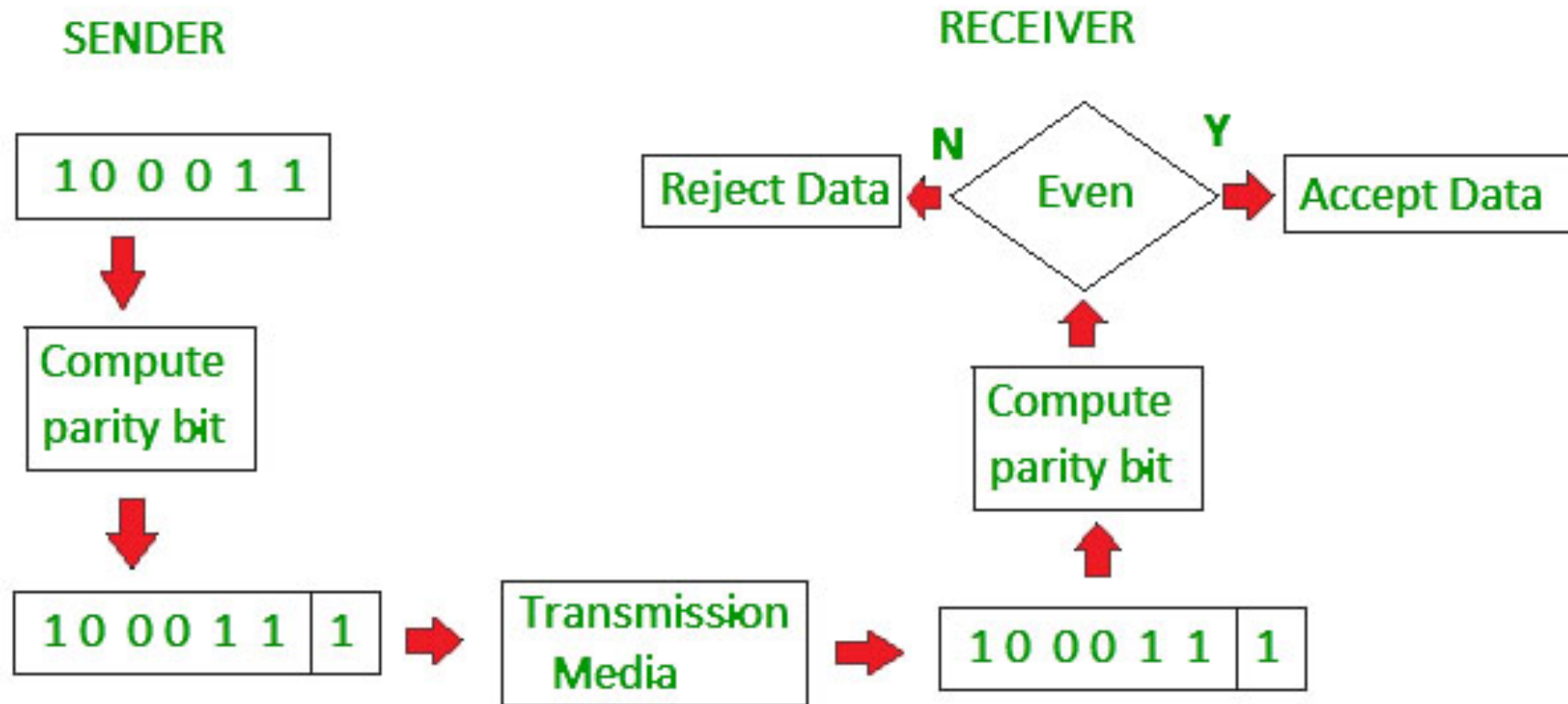
EDC= Error Detection and Correction bits (redundancy)

D = Data protected by error checking, may include header fields

- Error detection not 100% reliable!
 - protocol may miss some errors, but rarely
 - larger EDC field yields better detection and correction



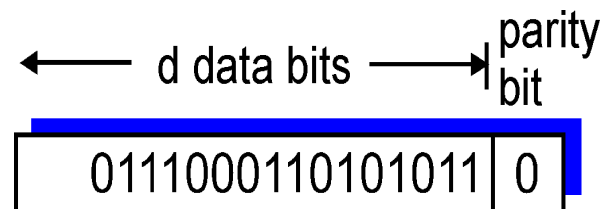
Parity checking



Parity checking

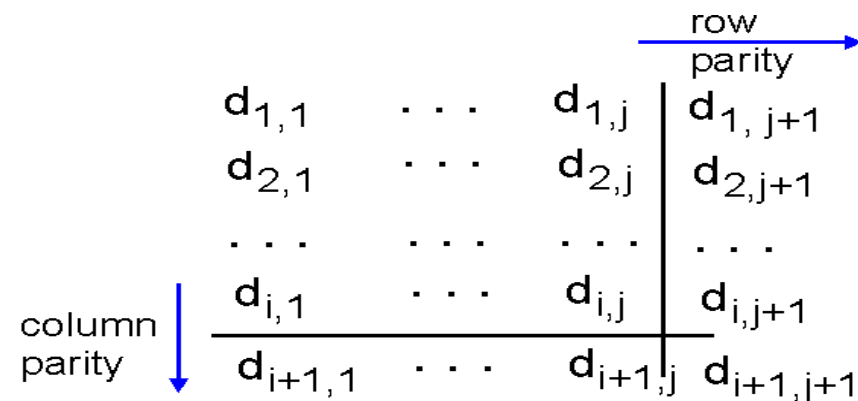
single bit parity:

- ❖ detect single bit errors



two-dimensional bit parity:

- ❖ detect and correct single bit errors



Two-dimensional bit parity helps to identify 2, 3 bit errors too, however cannot be corrected.

1	0	1	0	1	1
1	1	1	1	0	0
0	1	1	1	0	1
0	0	1	0	1	0

no errors

1	0	1	0	1	1
1	0	1	1	0	0
0	1	1	1	0	1
0	0	1	0	1	0

parity error

*correctable
single bit error*

Original Data

10011001	11100010	00100100	10000100
----------	----------	----------	----------

Row parities

10011001	0
11100010	0
00100100	0
10000100	0
11011011	0

Column parities



100110010	111000100	001001000	100001000	110110110
-----------	-----------	-----------	-----------	-----------

Data to be sent

Checksumming methods

- In checksum error detection scheme, the data is divided into k segments each of m bits.
- In the sender's end the segments are added using 1's complement arithmetic to get the sum. The sum is complemented to get the checksum.
- The checksum segment is sent along with the data segments.
- At the receiver's end, all received segments are added using 1's complement arithmetic to get the sum. The sum is complemented.
- If the result is zero, the received data is accepted; otherwise discarded.

Internet checksum (review)

goal: detect “errors” (e.g., flipped bits) in transmitted packet
(note: used at transport layer *only*)

sender:

- ❖ treat segment contents as sequence of 16-bit integers
- ❖ checksum: addition (1's complement sum) of segment contents
- ❖ sender puts checksum value into UDP checksum field

receiver:

- ❖ compute checksum of received segment
- ❖ check if computed checksum equals checksum field value:
 - NO - error detected
 - YES - no error detected.
But maybe errors nonetheless?

Original Data

10011001	11100010	00100100	10000100
----------	----------	----------	----------

1

2

3

4

k=4, m=8

Sender

```

1  1 0 0 1 1 0 0 1
2  1 1 1 0 0 0 1 0
-----
    1 0 1 1 1 1 0 1 1
      1
-----
      0 1 1 1 1 1 0 0
3  0 0 1 0 0 1 0 0
-----
      1 0 1 0 0 0 0 0
4  1 0 0 0 0 1 0 0
-----
    1 0 0 1 0 0 1 0 0
      1
-----
Sum: 0 0 1 0 0 1 0 1
Checksum: 1 1 0 1 1 0 1 0
  
```

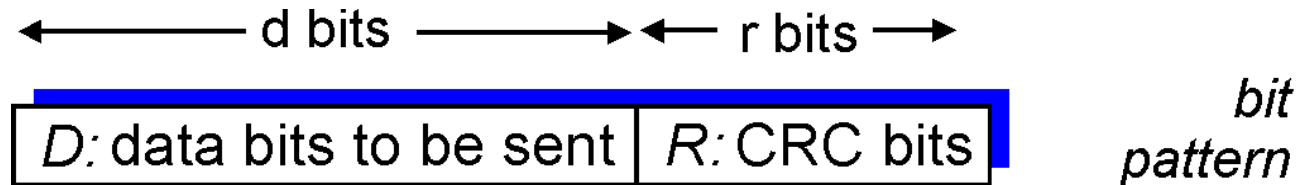
Receiver

```

1  1 0 0 1 1 0 0 1
2  1 1 1 0 0 0 1 0
-----
    1 0 1 1 1 1 0 1 1
      1
-----
      0 1 1 1 1 1 0 0
3  0 0 1 0 0 1 0 0
-----
      1 0 1 0 0 0 0 0
4  1 0 0 0 0 1 0 0
-----
    1 0 0 1 0 0 1 0 0
      1
-----
      0 0 1 0 0 1 0 1
      1 1 0 1 1 0 1 0
-----
Sum: 1 1 1 1 1 1 1 1
Complement: 0 0 0 0 0 0 0 0
Conclusion: Accept Data
  
```

Cyclic redundancy check

- ❖ more powerful error-detection coding
- ❖ view data bits, **D**, as a binary number
- ❖ choose $r+1$ bit pattern (generator), **G**
- ❖ goal: choose r CRC bits, **R**, such that
 - $\langle D, R \rangle$ exactly divisible by G (modulo 2)
 - receiver knows G , divides $\langle D, R \rangle$ by G . If non-zero remainder: error detected!
 - can detect all burst errors less than $r+1$ bits
- ❖ widely used in practice (Ethernet, 802.11 WiFi, ATM)



$$D * 2^r \text{ XOR } R$$

mathematical formula

CRC example

want:

$$D \cdot 2^r \text{ XOR } R = nG$$

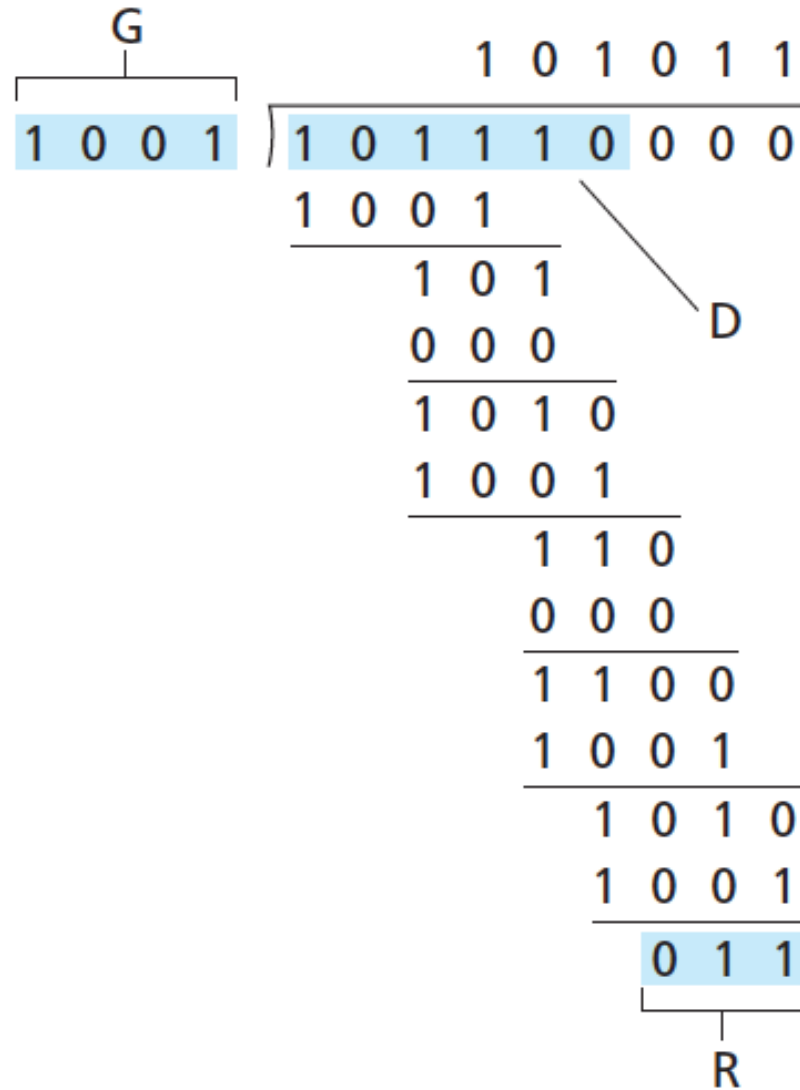
equivalently:

$$D \cdot 2^r = nG \text{ XOR } R$$

equivalently:

if we divide $D \cdot 2^r$ by G , want remainder R to satisfy:

$$R = \text{remainder}[\frac{D \cdot 2^r}{G}]$$



original message
1 0 1 0 0 0 0

@ means X-OR

Generator polynomial
 x^3+1
 $1.x^3+0.x^2+0.x^1+1.x^0$
CRC generator
1 0 0 1 4-bit

If CRC generator is of n bit then append $(n-1)$ zeros in the end of original message

Sender

```

1001 | 10100000000
@1001
-----
0011000000
@1001
-----
01010000
@1001
-----
0011000
@1001
-----
01010
@1001
-----
0011
  
```

Message to be transmitted

```

10100000000
+ 011
-----
1010000011
  
```

```

1001 | 1010000011
@1001
-----
0011000011
@1001
-----
01010011
@1001
-----
0011011
@1001
-----
01001
@1001
-----
0000
  
```

Receiver

Zero means data is accepted

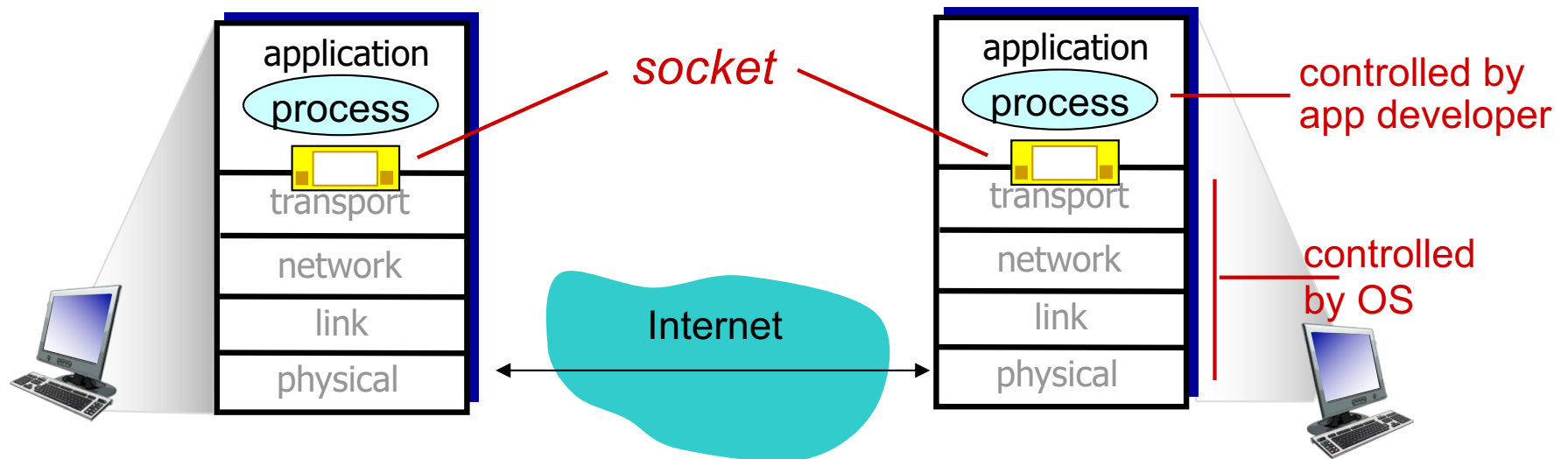
Network Application Programming

Sections 2.1.2, 2.7

Computer Networking – A top-down approach,
Kurose and Ross, 6th Edition.

Sockets

- ❖ process sends/receives messages to/from its **socket**
- ❖ socket analogous to door
 - sending process shoves message out door
 - sending process relies on transport infrastructure on other side of door to deliver message to socket at receiving process



Addressing processes

- ❖ to receive messages, process must have *identifier*
- ❖ host device has unique 32-bit IP address
- ❖ Q: does IP address of host on which process runs suffice for identifying the process?
 - A: no, *many* processes can be running on same host
- ❖ *identifier* includes both **IP address** and **port numbers** associated with process on host.
- ❖ example port numbers:
 - HTTP server: 80
 - mail server: 25
- ❖ to send HTTP message to gaia.cs.umass.edu web server:
 - **IP address**: 128.119.245.12
 - **port number**: 80
- ❖ more shortly...

Socket programming *with TCP*

client must contact server

- ❖ server process must first be running
- ❖ server must have created socket (door) that welcomes client's contact

client contacts server by:

- ❖ Creating TCP socket, specifying IP address, port number of server process
- ❖ *when client creates socket:* client TCP establishes connection to server TCP

- ❖ when contacted by client, *server TCP creates new socket* for server process to communicate with that particular client
 - allows server to talk with multiple clients
 - source port numbers used to distinguish clients (more in Chap 3)

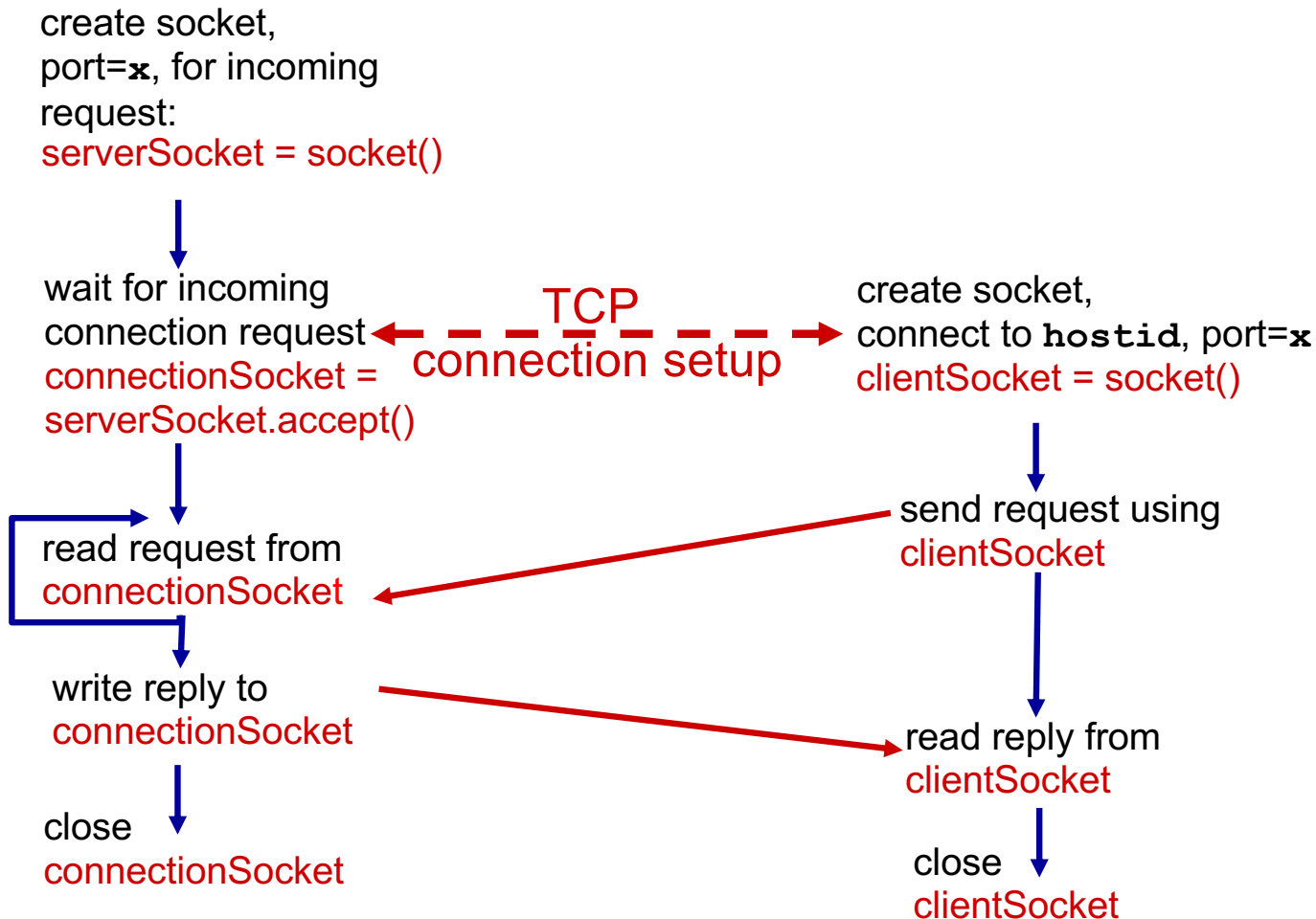
application viewpoint:

TCP provides reliable, in-order byte-stream transfer (“pipe”) between client and server

Client/server socket interaction: TCP

server (running on `hostid`)

client



Example app:TCP client

Python TCPClient

create TCP socket for
server, remote port 12000

```
from socket import *
serverName = 'servername'
serverPort = 12000
clientSocket = socket(AF_INET, SOCK_STREAM)
clientSocket.connect((serverName,serverPort))
sentence = raw_input('Input lowercase sentence:')
clientSocket.send(sentence)
modifiedSentence = clientSocket.recv(1024)
print 'From Server:', modifiedSentence
clientSocket.close()
```

No need to attach server
name, port

Example app: TCP server

Python TCPServer

create TCP welcoming socket	→	from socket import *
		serverPort = 12000
server begins listening for incoming TCP requests	→	serverSocket = socket(AF_INET, SOCK_STREAM)
		serverSocket.bind(('', serverPort))
		serverSocket.listen(1)
		print 'The server is ready to receive'
loop forever	→	while 1:
server waits on accept() for incoming requests, new socket created on return	→	connectionSocket, addr = serverSocket.accept()
read bytes from socket (but not address as in UDP)	→	sentence = connectionSocket.recv(1024)
		capitalizedSentence = sentence.upper()
		connectionSocket.send(capitalizedSentence)
close connection to this client (but <i>not</i> welcoming socket)	→	connectionSocket.close()