Information Security Technologies COMP607

Assignment 1

Student Name: Dao Minh Duc

AUT Student ID: 23199710

1.

a.

2 Rows																														
Α		Α		Ε		Н		D		S		G		N		М		В		Т		Т		Α		0		Н		Т
	0		D		Е		S		Т		R		N		0		Α		Ι		0		Ι		Ε		G		В	
3 Rows																														
Α				Α				Ε				Н				D				S				G				N		
	M		В		T		Т		Α		0		Н		T		0		D		Ε		S		Т		R		N	
		0				Α				Ι				0				Ι				Ε				G				В
4 Rows																														
Α						Α						E						Н						D						S
	G				N		M				В		T				Т		Α				0		Н				T	
		0		D				Ε		S				Т		R				N		0				Α		Ι		
			0						Ι						Ε						G						В			

Key: 4

Plaintext message: A GOOD NAME IS BETTER THAN GOOD HABITS

b.

Plaintext message: KNO WLE DGE ISN EVE RWA STE D

Key: **EQUATORIAL**

Ks = 4; 16; 20; 0; 19; 14; 17; 8; 0; 11

Р	10	13	14	22	11	4	3	6	4	8	18	13	4	21	4	17	22	0	18	19	4	3
Ks	4	16	20	0	19	14	17	8	0	11	4	16	20	0	19	14	17	8	0	11	4	16
C	14	3	8	22	4	18	20	14	4	19	22	3	24	21	23	5	13	8	18	4	8	19
	0	D	Ι	W	Е	S	U	0	Е	Т	W	D	Υ	٧	Χ	F	N	Ι	S	Е	Ι	T

The cipher text is: ODI WES UOE TWD YVX FNI SEI T

C.

i.

ii.

With 7 bits per character, there are $2^7=128$ possible values for each character in the password. With 10 such characters in a password, the key space of the password will have the size of 128^{10} .

b.

There are 10 7-bit characters in the password. Therefore, the key length in bits of it is: 10 imes 7 = 70 bits.

C.

If only 26 lowercase characters are used, we only need at most 5 bits per character to encode the password. Therefore, the key length in bits in this case will be: $10 \times 5 = 50$ bits.

d.

- (i) 128/7=18. Therefore, we need at least 18 7-bit characters.
- (ii) 128/5=25.2. Therefore, we need at least 26 characters.

The bit sequence that appears the most frequent is <code>00011111</code>, so using frequency attack, we can assume that this sequence represents letter "E", which has the ASCII encode of <code>01000101</code>.

Therefore, we can assume that the key of this cipher text is: <code>00011111</code> XOR <code>01000101 = 01011010</code>.

With that key, we can work out the rest of the message by adding it to the bits in the cipher text.

```
• 00010111 XOR 01011010 = 01001101 = m
```

- 00001110 XOR 01011010 = 01010100 = t
- 00011011 XOR 01011010 = 01011010 = a
- 00010110 XOR 01011010 = 01001100 = 1
- 00001100 XOR 01011010 = 01010110 = v
- 00010100 XOR 01011010 = 01001110 = n

The decrypted cipher text is: MEETATELEVENAM

This can be interprteted in plaintext as: Meet at eleven AM

```
~/Documents
                                                [ 12:19:10 ]
cat <u>putty.md5</u>
9047a29b7c2ed333536a7fb6d6c8bae6 putty-0.70-installed.msi
                                                [ 12:19:12 ]
~/Documents
> cp <u>putty.md5</u> putty r.md5
sending incremental file list
putty.md5
                        0.00kB/s 0:00:00 (xfr#1, to-chk=
             59 100%
0/1)
                                      [ ② 12s ] [ 12:20:15 ]
~/Documents
cat <u>putty r.md5</u>
9047a29b7c2ed333536a7fb6d6c8bae6 putty-0.70-installer.msi
~/Documents
                                                [ 12:20:19 ]
md5sum -c putty.md5 putty r.md5
putty-0.70-installed.msi: OK
putty-0.70-installer.msi: FAILED
```

Therefore, putty-0.70-installed.msi is the good copy.

md5sum: WARNING: 1 computed checksum did NOT match

```
sbt7974@Scopius:~/data$ mkdir 05
sbt7974@Scopius:~/data$ cd 05
sbt7974@Scopius:~/data/05$ cp ../notice*.* .
sbt7974@Scopius:~/data/05$ ls
notice1.hmac.txt notice2.hmac.txt notice2.txt.sig
notice1.txt notice2.txt
sbt7974@Scopius:~/data/05$ cat notice1.txt
Exam Notice

The exam is very easy and you do not need to study for it.
No need to work hard to understand the material.
Enjoy and have a good time.
Warmest regards.
Dr. Yang
```

```
sbt7974@Scopius:~/data/05$ cat notice2.txt
Exam Notice

The exam is not easy and you do need to study for it.

Do work hard to understand the material.

After the exam then you deserve to enjoy and have a good ti me.

Warmest regards.
Dr. Yang
```

```
sbt7974@Scopius:~/data/05$ openssl dgst -hmac comp607 notice1
.txt > notice1.hmac.test
sbt7974@Scopius:~/data/05$ openssl dgst -hmac comp607 notice2
.txt > notice2.hmac.test
sbt7974@Scopius:~/data/05$ cmp notice1.hmac.txt notice1.hmac.
test
notice1.hmac.txt notice1.hmac.test differ: byte 27, line 1
sbt7974@Scopius:~/data/05$ cmp notice2.hmac.txt notice2.hmac.
test
```

Therefore, notice2.hmac.txt is the authentic version.

```
~/Documents/06
> echo "This is some text file" > test1.txt

~/Documents/06
> md5sum test1.txt > test1.md5

~/Documents/06
> cat test1.md5
6b101d0e17e56ea3db991454fda2fb5f test1.txt
[ 12:36:19 ]

[ 12:36:20 ]

[ 12:36:21 ]
```

b.

```
sbt7974@Scopius:~/test$ cat test1.txt
This is a message!
sbt7974@Scopius:~/test$ sed '$ s/.$//' test1.txt > test2.txt
sbt7974@Scopius:~/test$ cat test2.txt
This is a message
sbt7974@Scopius:~/test$ md5sum test1.txt > test1.md5
sbt7974@Scopius:~/test$ md5sum test2.txt > test2.md5
sbt7974@Scopius:~/test$ cat test1.md5
f1ce17046c1dab5dba96e37d02938b96 test1.txt
sbt7974@Scopius:~/test$ cat test2.md5
9b7eb2f1b70b39d14a53846bddea2f4e test2.txt
```

Conclusion: test1.txt and test2.txt plaintext files only differs in the last byte (! character), but their MD5 hash values are completely different.

C.

$$M = 513, p = 23, q = 29$$

$$\Rightarrow N = p \times q = 23 \times 29 = 667$$

$$\Rightarrow \phi(N) = (p-1) imes (q-1) = 22 imes 28 = 616$$

Select e=5

$$\Rightarrow d = e^{-1} \mod \phi(N) = 5^{-1} \mod 616 = 493$$

Public key: (e, N) = (5, 667)

Private key: (d, N) = (493, 667)

Encrypt M=513 using public key:

$$C = M^e \mod N = 513^5 \mod 667 = 546$$

Decrypt C=546 using private key:

$$M=C^d \mod N = 546^{493} \mod 667 = 513$$

b.

$$M = 109, p = 11, q = 23$$

$$\Rightarrow N = p \times q = 11 \times 23 = 253$$

$$\Rightarrow \phi(N) = (p-1) \times (q-1) = 10 \times 22 = 220$$

Select e=3

$$\Rightarrow d = e^{-1} \mod \phi(N) = 3^{-1} \mod 220 = 147$$

Signature: $S=M^d \mod N=109^{411} \mod 253=109$

Verify signature:

$$M = S^e \mod N = 109^3 \mod 253 = 109$$

The signature is verified.

8.

a.

First, Alice and Bob agree on a prime number n=4787 and a generator g=2.

Then, they each choose a secret number, a and b, respectively.

Alice chooses a=3 and Bob chooses b=5.

They then calculate their public keys as follows:

Alice:
$$A=g^a \mod n=2^3 \mod 4787=8$$

$$\operatorname{Bob:} B = g^b \mod n = 2^5 \mod 4787 = 32$$

They then exchange their public keys.

Alice receives Bob's public key, B=32, and calculates the shared key as follows:

$$K = B^a \mod n = 32^3 \mod 4787 = 4046$$

Bob receives Alice's public key, A=8, and calculates the shared key as follows:

$$K = A^b \mod n = 8^5 \mod 4787 = 4046$$

Both Alice and Bob now have a shared key, K=4046.

b.

Both Alice and Bob can determine the value of the shared key. The shared key is calculated using the public key of the other party and the secret key of the party itself.

i.

$$M = 215, a = 198, b = 231, G = 2, p = 443$$

Alice calculates her public key:

$$A = G^a \mod p = 2^{198} \mod 443 = 144$$

Bob calculates his public key:

$$B = G^b \mod p = 2^{231} \mod 443 = 305$$

Alice sends her public key, A=144, to Bob.

Bob sends his public key, B=305, to Alice.

Alice computes the shared key:

$$K = B^a \mod p = 305^{198} \mod 443 = 321$$

Alice encrypts the message:

$$C=M\times K \mod p=215\times 321 \mod 443=350$$

Alice sends the cipher text, C=350, to Bob.

Bob derives the shared key:

$$K = A^b \mod p = 144^{231} \mod 443 = 321$$

Bob decrypts the message: $M = C imes K^{-1} \mod p = 350 imes 321^{-1} \mod 443 = 215$

ii.

$$M = 67, a = 11, b = 13, G = 2, p = 17$$

Alice calculates her public key:

$$A = G^a \mod p = 2^{11} \mod 17 = 8$$

Bob calculates his public key:

$$B = G^b \mod p = 2^{13} \mod 17 = 15$$

Alice sends her public key, A=8, to Bob.

Bob sends his public key, B=15, to Alice.

Alice computes the shared key:

$$K=B^a \mod p=15^{11} \mod 17=9$$

Alice encrypts the message:

$$C = M \times K \mod p = 67 \times 9 \mod 17 = 8$$

Alice sends the cipher text, C=8, to Bob.

Bob derives the shared key:

$$K = A^b \mod p = 8^{13} \mod 17 = 9$$

Bob decrypts the message: $M = C imes K^{-1} \mod p = 8 imes 9^{-1} \mod 17 = 67$