

# Memory Management

Johan Wahlström

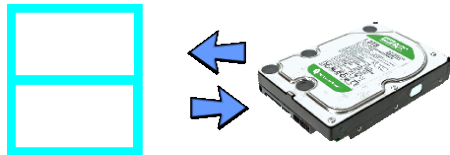
ECM1413 Computers and the Internet

# Processes

- A process is a program in execution
- Multiple processes may be active at the same time
- Each process is represented by a process control block (PCB) which is stored in memory (or on disk)
- How do we allocate space to processes in memory in an efficient way?



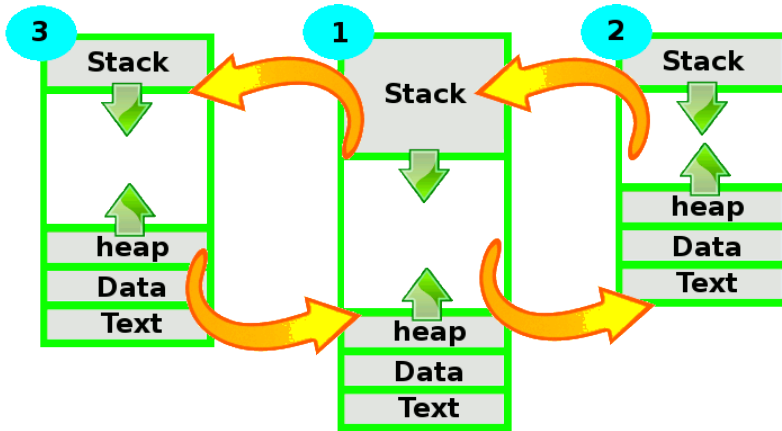
# Memory Management Tasks



Memory management tasks include

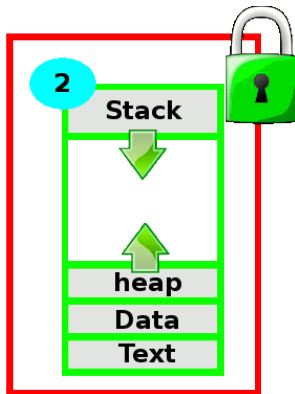
- ① relocation
- ② protection
- ③ sharing

# (1) Relocation



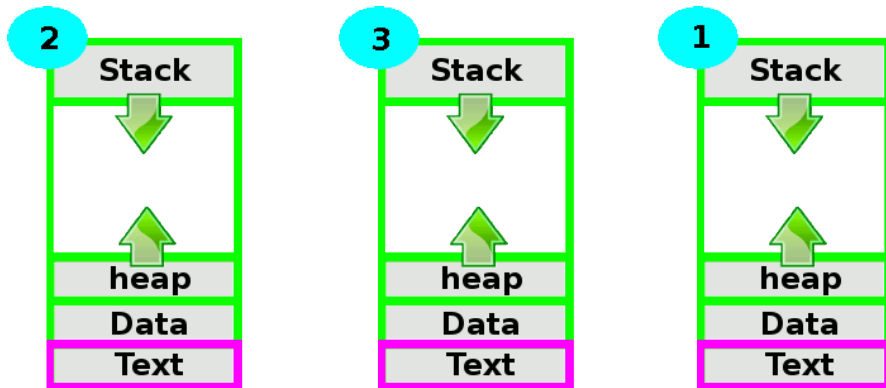
A memory management technique needs to be able to relocate a process to a new place in memory.

## (2) Protection



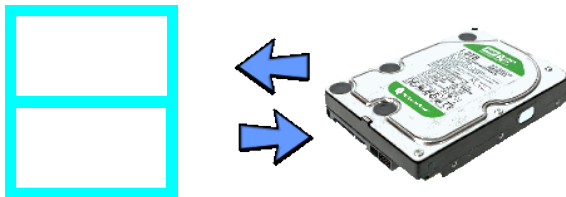
All processes need to be protected against interference from other processes.

### (3) Sharing



Processes may sometimes need to share data.

# Memory Management Techniques

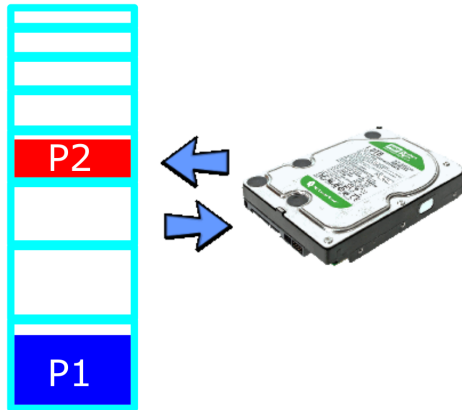


We will discuss the following memory management techniques:

- ① fixed partitioning
- ② dynamic partitioning
- ③ simple segmentation
- ④ virtual memory segmentation
- ⑤ simple paging
- ⑥ virtual memory paging

# (1) Fixed Partitioning

- Memory is divided into partitions of a fixed size; the size may vary between partitions
- A process can be loaded into any partition whose size is equal to or greater than the size of the process
- One-to-one mapping between partitions and processes



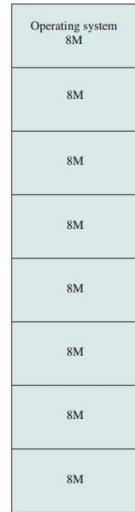


# (1) Fixed Partitioning – Unequal or Equal Partition Sizes

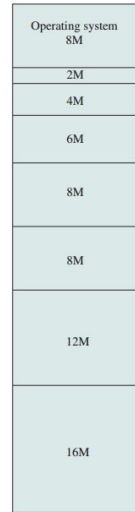
Risks with equal-sized partitions:

- Internal fragmentation: Space is wasted inside partitions when the size of the process is smaller than the size of the partition
- There may not be any partition that can fit a large process

⇒ Unequal-sized partitions are more efficient.



Equal-size partitions



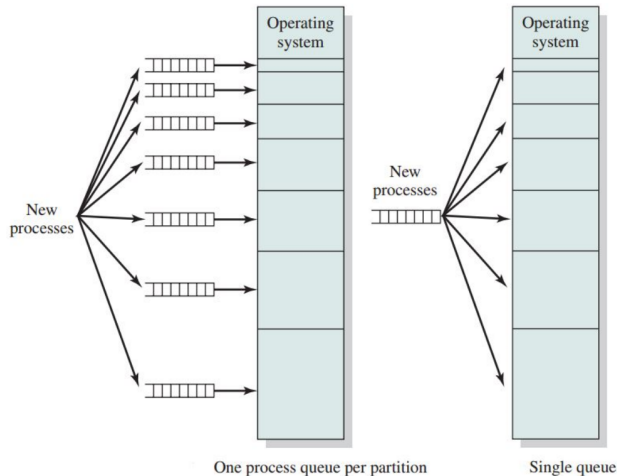
Unequal-size partitions

# (1) Fixed Partitioning – Queuing

Assume unequal-sized partitions.

How do we match processes to partitions?

- Alternative 1 (one process queue per partition): Assign each process to the smallest partition where it will fit
  - Could mean that some partitions remain empty
- Alternative 2 (single queue): Assign each process to the smallest available partition where it will fit



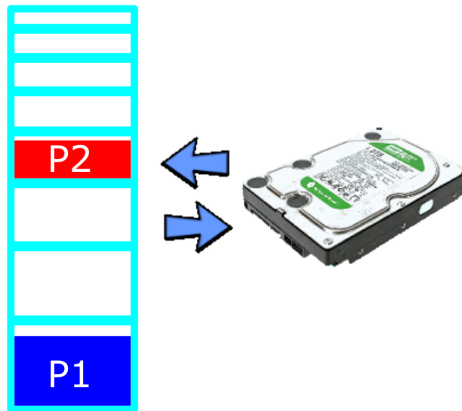
# (1) Fixed Partitioning – Pros and Cons

Pros:

- Easy to understand and implement

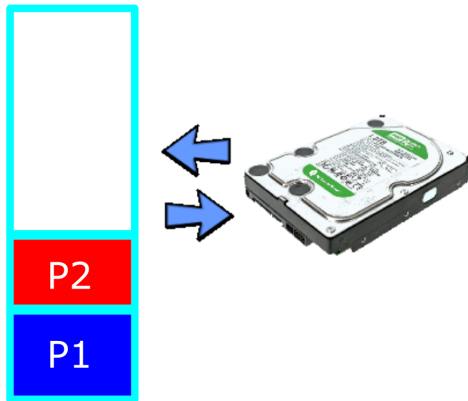
Cons:

- Internal fragmentation (wasted space inside partitions)
- A pre-specified upper limit on the number of processes
- A pre-specified upper limit on the size of the largest process

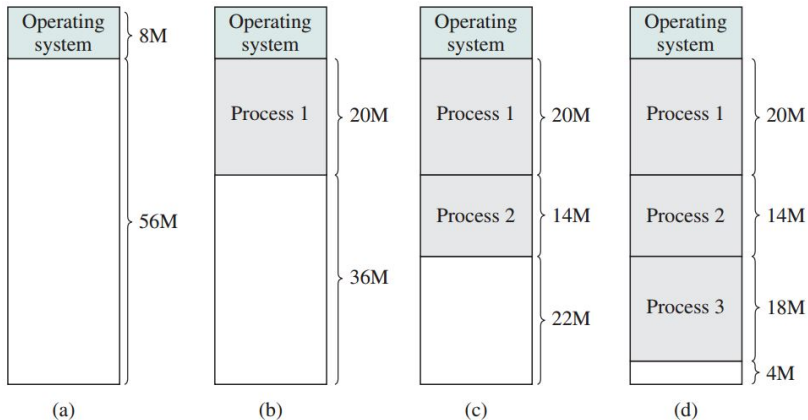


## (2) Dynamic Partitioning

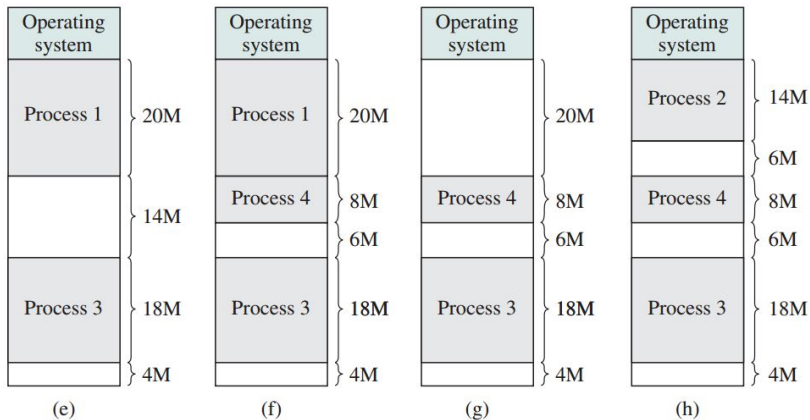
- The partition sizes adjust to the sizes of processes
- One-to-one mapping between partitions and processes



## (2) Dynamic Partitioning – Example

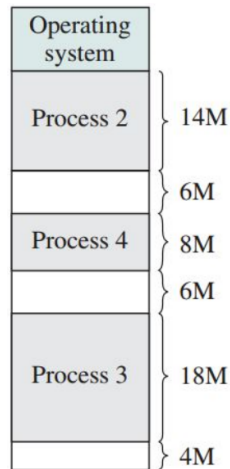


## (2) Dynamic Partitioning – Example (ctd)



## (2) Dynamic Partitioning – External Fragmentation

- External fragmentation: Space is wasted between partitions
- Compaction: Shift processes so that they are contiguous and so that all free memory is gathered in one continuous block
  - Downside: Takes a lot of time and resources



(h)

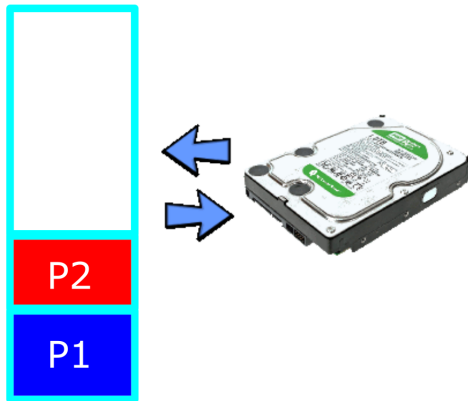
## (2) Dynamic Partitioning – Pros and Cons

Pros:

- No internal fragmentation
- No limit on the number of processes or on the size of the largest process

Cons:

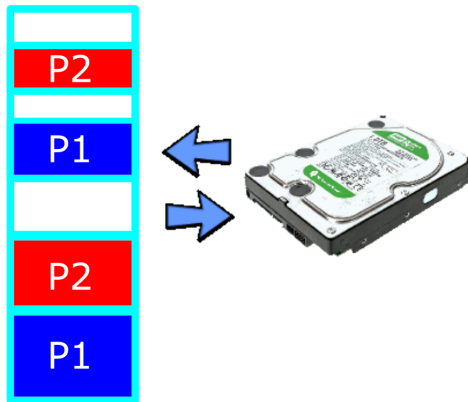
- External fragmentation (wasted space between partitions)
- Time and resources wasted on compaction





### (3) Simple Segmentation

- Each program is divided into segments (e.g., text and data segments)
- Segments are loaded into memory as in dynamic partitioning
- All segments need to be loaded into memory, but they do not need to be contiguous



### (3) Simple Segmentation – Pros and Cons

#### Pros:

- Easier to fit processes in memory than with dynamic partitioning
- No internal fragmentation
- No limit on the number of processes or on the size of the largest process

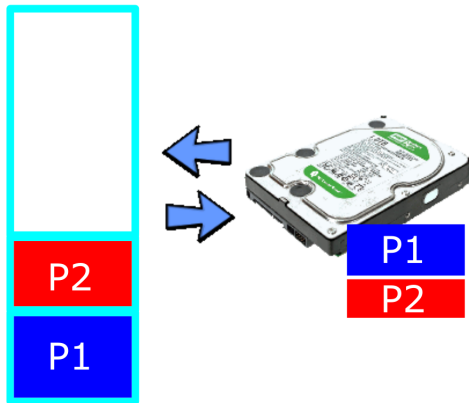
#### Cons:

- External fragmentation (less than dynamic partitioning)
- Time and resources wasted on compaction



## (4) Virtual Memory Segmentation

- Virtual memory: Use disk storage as if it was main memory
  - Makes it seem as if main memory is very big
- Segments are loaded into memory as in dynamic partitioning, except that not all segments need to be loaded at the same time



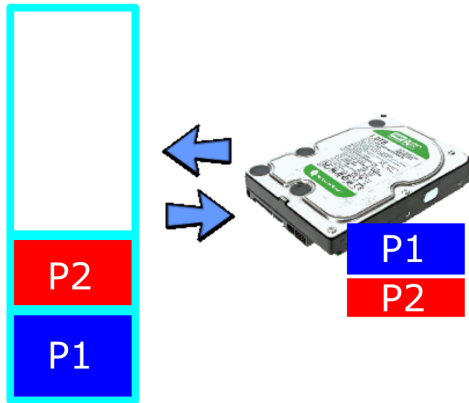
## (4) Virtual Memory Segmentation – Pros and Cons

### Pros:

- Easier to fit processes in memory than with simple segmentation
- No internal fragmentation
- No limit on the number of processes or on the size of the largest process

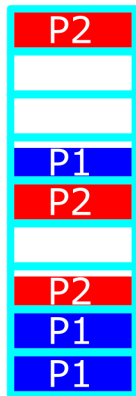
### Cons:

- More overhead required for virtual memory
- External fragmentation (less than dynamic partitioning)
- Time and resources wasted on compaction



## (5) Simple Paging

- Memory is divided into fixed-size frames of equal size
- Each process is divided into pages of the same size
- Internal fragmentation is small since
  - it only affects the last frame of each process
  - the frame/page-size is fairly small



## (5) Simple Paging – Example

Frame number	Main memory
0	
1	
2	
3	
4	
5	
6	
7	
8	
9	
10	
11	
12	
13	
14	

(a) Fifteen available frames

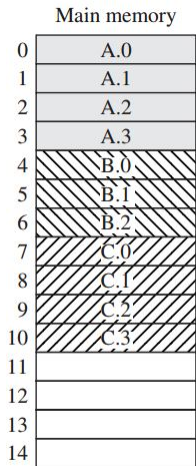
	Main memory
0	A.0
1	A.1
2	A.2
3	A.3
4	
5	
6	
7	
8	
9	
10	
11	
12	
13	
14	

(b) Load process A

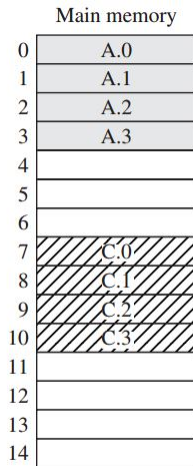
	Main memory
0	A.0
1	A.1
2	A.2
3	A.3
4	B.0
5	B.1
6	B.2
7	
8	
9	
10	
11	
12	
13	
14	

(c) Load process B

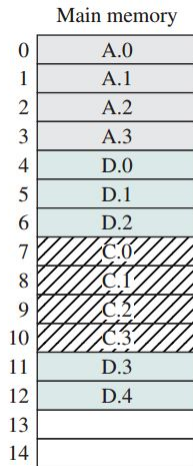
## (5) Simple Paging – Example (ctd)



(d) Load process C



(e) Swap out B



(f) Load process D

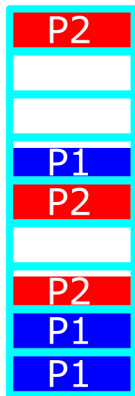
## (5) Simple Paging – Pros and Cons

Pros:

- No external fragmentation
- All pages fit perfectly into frames

Cons:

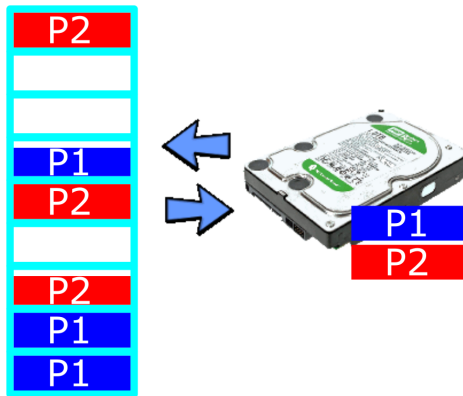
- (small) Internal fragmentation
- Each process requires a page tables (which consumes memory)





## (6) Virtual Memory Paging

- Pages are loaded into frames as in simple paging, except that not all pages need to be loaded at the same time



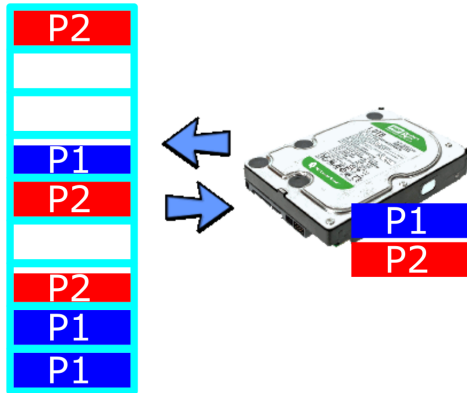
## (6) Virtual Memory Paging – Pros and Cons

### Pros:

- Easier to fit processes in memory than with simple paging
- No external fragmentation
- All pages fit perfectly into frames

### Cons:

- More overhead required for virtual memory
- (small) Internal fragmentation
- Each process requires a page tables (which consumes memory)



# Protection and Sharing within Segmentation and Paging

## Goals:

- Protection: Each process can only access its own pages/segments
- Sharing: Two processes access the same page/segment

## Implementation:

- Each frame/segment has a number of protection bits specifying
  - read-write or read-only
  - which process can access

Protection/sharing is generally easier with segmentation than with paging since

- the segment table is smaller than the page table
- segments are more natural units for protection/sharing