

# CSCI 2110 Data Structures and Algorithms

## Module 4: Ordered Lists



# Learning objectives

- *Define the ordered list data structure.*
- *Develop binary search algorithm on ordered lists.*
- *Build an OrderedList class.*
- *Understand merging operations on ordered lists.*

# What is an ordered list?

- It is a linear collection of items, in which the items are arranged in either ascending or descending order of **keys**.
- The **key** is one part of the item. It serves as the basis of ordering. **The key may vary from application to application.**
- This means that an **ordered list is sorted and maintained sorted** - even when items are added or deleted.
- **Repetition of items is normally not allowed.**
- Example of an ordered list: List of student entries in a particular course (student name, ID number, major, marks). Here the key could be the student name or the ID number.

In the first list, the key is the name;

In the second list, the key is the ID no.

item

Sorted  
↓

Name	ID	Major	Mark
Andy	9856	BCS	90.5
Boris	7859	Blnf	87.5
Dominic	3664	BA	96.6
Earl	5654	BCS	77.6
Tasha	8776	BSc.	93.4

Sorted  
↓

Name	ID	Major	Mark
Dominic	3664	BA	96.6
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Boris	7859	Blnf	87.5
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Andy	9856	BCS	90.5

# What is the advantage of having a list that is always ordered?

- The big advantage of an ordered list is that it enables fast search for an item.
- This is because ordered lists can be searched using the **Binary Search algorithm**.
- Binary search of **n items** takes only  **$O(\log_2 n)$**  in the worst case.
- We will see, however, that this advantage comes with a cost – for inserting and deleting items.

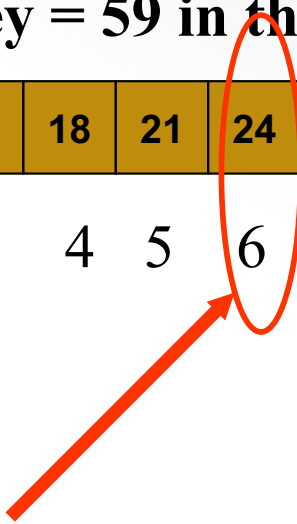
# Binary Search Algorithm

- Binary search is a powerful algorithm that can be used to search for a key in a sorted list with non-repeated items.
- *The idea in binary search is to divide the list in half, and check if the item is in the left half or the right half.*
- This procedure is repeated on smaller and smaller portions of the list.

# Binary Search Example

**Search for key = 59 in the following array:**

4	5	9	12	18	21	24	27	28	32	45	59	60
0	1	2	3	4	5	6	7	8	9	10	11	12



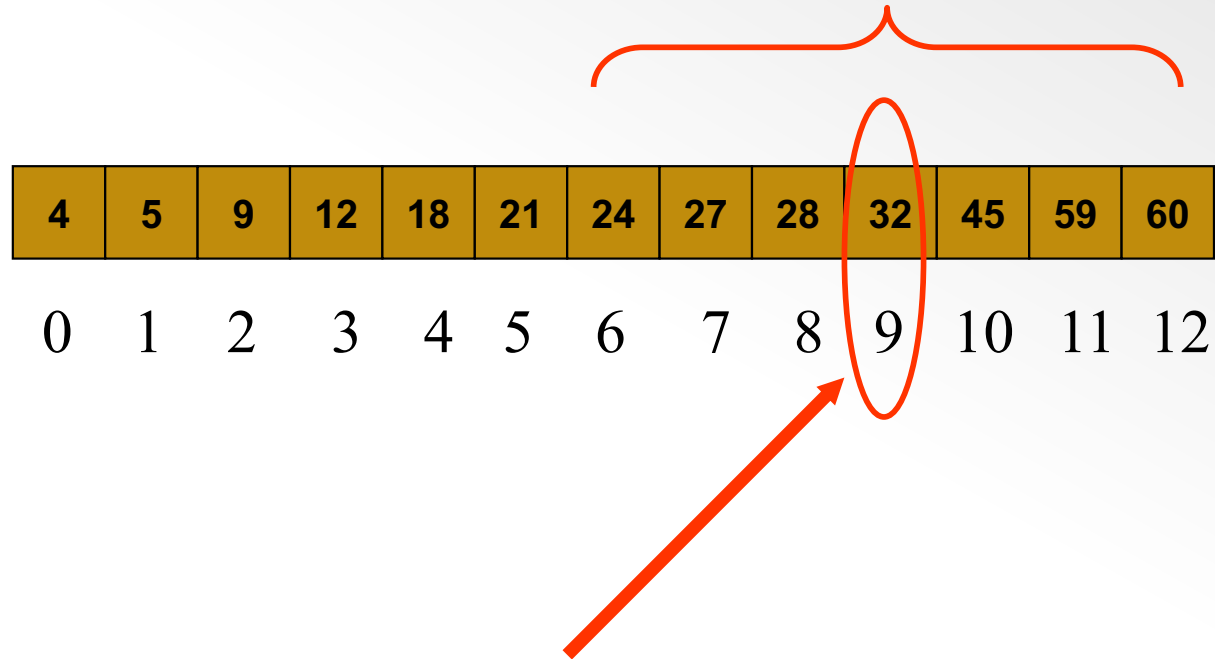
Check the element in the middle of the array.

The middle element is  $a[6] = 24$ .

The key 59 is  $> 24$

So if the key is present, it should be in the right half.

# Binary Search Algorithm



4	5	9	12	18	21	24	27	28	32	45	59	60
0	1	2	3	4	5	6	7	8	9	10	11	12

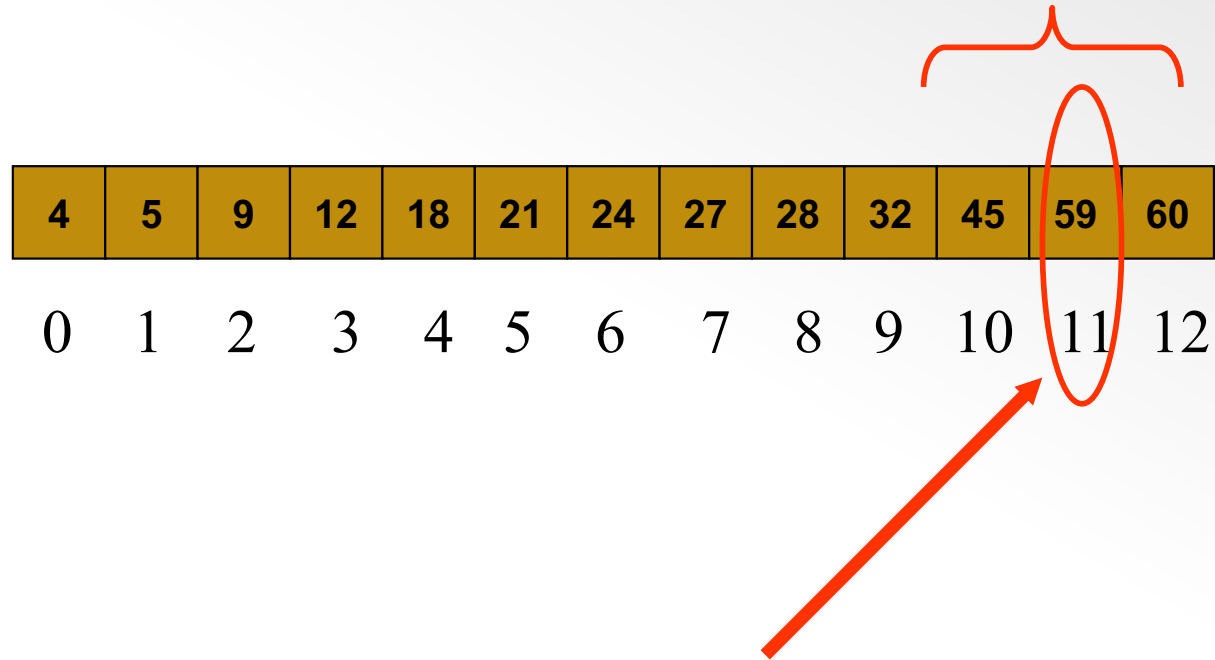
Check the element in the middle of this half.

The element is  $a[9] = 32$

The key 59 is  $> 32$

So if the key is present, it should be in the right half of this half.

# Binary Search Algorithm



The diagram illustrates a binary search step on a sorted array. The array is represented as a horizontal row of 13 yellow boxes, each containing a number. Below each box is its corresponding index, from 0 to 12. A red bracket above the array spans from index 10 to index 12, indicating the current search range. A red oval encircles the element at index 11, which is 59. A red arrow points from the bottom of the slide towards the element 59.

4	5	9	12	18	21	24	27	28	32	45	59	60
0	1	2	3	4	5	6	7	8	9	10	11	12

Check the element in the middle of this half.  
The element is  $a[11] = 59$ .  
Key found!



## BINARY SEARCH ALGORITHM IN MORE DETAIL

Let's understand binary search algorithm in more detail. Assume that the ordered list consists of Strings (names) stored in an array. The principle will be the same for binary search on any other data structure and for any other type of object.

Names are sorted in alphabetical order

Search for "Dan" → target key

Amar	Boris	Charlie	Dan	Fujian	Inder	Liz	Sam	Travis	Wendy
0	1	2	3	4	5	6	7	8	9

$lo = 0, hi = 9$   
 $mid = (lo + hi) / 2 = (0 + 9) / 2 = 4$   
 Integer Division

Check the item at index 4. Dan < Fujian (alphabetically)  
 ∴ Check the left half.

Amar	Boris	Charlie	Dan	Fujian	Inder	Liz	Sam	Travis	Wendy
0	1	2	3	4	5	6	7	8	9

$lo = 0, hi = mid - 1 = 4 - 1 = 3$   
 Find the new mid.  $mid = (lo + hi) / 2 = (0 + 3) / 2 = 1$   
 Check the item at index 1. Dan > Boris.  
 ∴ Check the right half.

Amar	Boris	Charlie	Dan	Fujian	Inder	Liz	Sam	Travis	Wendy
0	1	2	3	4	5	6	7	8	9

$lo = mid + 1 = 1 + 1 = 2, hi = 3$   
 New mid.  $mid = (lo + hi) / 2 = (2 + 3) / 2 = 2$

Search the right half  
 Dan > Charlie

Amar	Boris	Charlie	Dan	Fujian	Inder	Liz	Sam	Travis	Wendy
0	1	2	3	4	5	6	7	8	9

$lo = mid + 1 = 2 + 1 = 3$   
 $hi = 3$   
 $mid = (3 + 3) / 2 = 3$

Dan found!

Another example: Search for "Trevor"

Amar	Boris	Charlie	Dan	Fujian	Inder	Liz	Sam	Travis	Wendy
0	1	2	3	4	5	6	7	8	9

lo	hi	mid	Found?
0	9	$(0+9)/2 = 4$	No. Trevor > Fujian. Go right.
5	9	$(5+9)/2 = 7$	No. Trevor > Sam. Go right.
$7+1 = 8$	9	$(8+9)/2 = 8$	No. Trevor > Travis. Go right.
$8+1 = 9$	9	$(9+9)/2 = 9$	No. Trevor < Wendy. Go left.
9	$9-1 = 8$		

Stop.  $lo > hi$   $\therefore$  Trevor not found

### Pseudocode

Algorithm Binary Search (A, n, t)

Input: array A of length n, target t  $\leftarrow$  Key

```

lo <-- 0
hi <-- n-1
mid <-- (lo+hi)/2
while (lo <= hi)
{
    if (t == A[mid])
        key found; break out of loop

    else if (t < A[mid])
        hi <-- mid-1

    else if (t > A[mid])
        lo <-- mid + 1

    mid <-- (lo+hi)/2
}

```

$\leftarrow$  Go left

$\leftarrow$  Go right

$\leftarrow$  find the new mid

```

if (lo > hi)
    key not found
else
    key found at mid

```

# COMPLEXITY ANALYSIS OF BINARY SEARCH

Size of the list n	Maximum number of searches
n=16 $2^4$	5 { 1st search: Array of size 16 2nd search: Array of size 8 3rd search: Array of size 4 4th search: Array of size 2 5th search: Array of size 1
n=32 $2^5$	6 { 32 → 16 → 8 → 4 → 2 → 1
n=64 $2^6$	7
n=2 <sup>k</sup>	K+1

Krishna

Complexity: n is a power of two

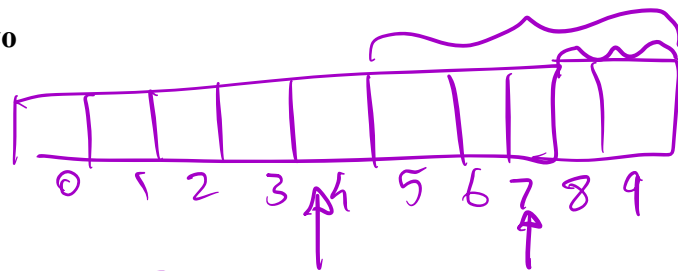
# Searches = K+1 where  $n = 2^k$   
 $= \log_2 n + 1$  that is,  $K = \log_2 n$

$\rightarrow O(\log_2 n)$

Complexity: n is not a power of two

$n = 10$

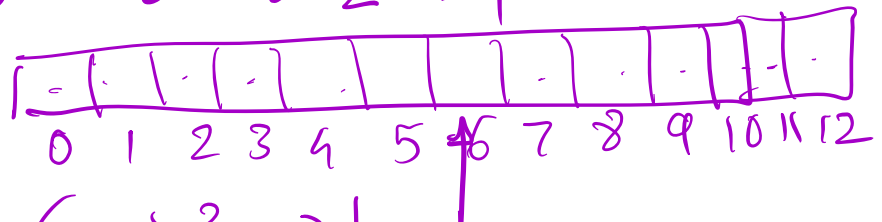
# Searches = 4



$10 \rightarrow 5 \rightarrow 2 \rightarrow 1$

$n = 13$

# Searches = 4



$13 \rightarrow 6 \rightarrow 3 \rightarrow 1$

$\lceil \log_2 n \rceil$

e.g.  $\lceil \log_2 13 \rceil \rightarrow \lceil 3.7 \rceil = 4$

$O(\log_2 n)$

## Implementation of the OrderedList class

- Our internal data structure to implement the OrderedList will be an ArrayList and not a LinkedList.
- Why? Binary search on a linked list is expensive.
- If we use the ArrayList, search is fast but it comes with a price.
  - Each time we insert an item (in the correct position) or remove an item, entries will have to be shifted.
  - This will cost  $O(n)$  for each insert or remove.
  - We will accept this cost because search is the important operation on such lists.

- We need to write a generic class.
- Recall that the entire binary search operation is based on comparisons.
- *This means we need a generic compareTo method.*
- Java has a generic interface called Comparable<T> for comparing objects of any type.
- *It has a method called compareTo with three possible return values:*
  - *0 → equal objects*
  - *Positive integer → “this” object is greater than the parameter*
  - *Negative integer → “this” object is less than the parameter.*
- Several Java classes like the String class implement this method.
- *If we need compareTo to work for any kind of object, we must define our class as follows:*

public class OrderedList<T extends Comparable<T>> {....}

*81. compareTo (82)*  
*81 = "Srinji"      82 = "Srinji"*  
*→ 0*

*81 = "Srinji"      82 = "Steve"*  
*-ve integer*

*81 = "Steve"      82 = "Srinji"*  
*+ve integer*

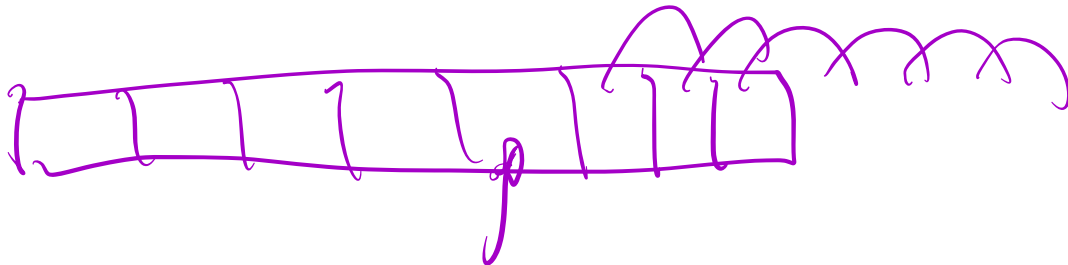
## IMPLEMENTATION OF ORDERED LIST CLASS

### Constructors

OrderedList()	Constructs an empty ordered list
---------------	----------------------------------

### Methods

Name	What it does	Header	Price tag (complexity)
size	returns size of the list	int size()	$O(1)$
isEmpty	returns true if list is empty	boolean isEmpty()	$O(1)$
clear	clears the list	void clear()	$O(1)$
get	gets the entry at the specified position	T get(int pos)	$O(1)$
first	gets the first entry	T first ()	$O(1)$
next	gets the next entry	T next()	$O(1)$
enumerate	scans the list and prints it	void enumerate()	$O(n)$
binarySearch	searches for a given item. returns position (index) if found if not found <u>returns a negative number</u>	int binarySearch(T item)	$O(\log_2 n)$
add	add a specified item at a given position	void add(int pos, T item)	$O(n)$
insert	insert a specified item at the right position	void insert(T item)	$O(n)$
remove	remove a specified item	void remove(T item)	$O(n)$
remove	remove item from a specified position	void remove(int pos)	$O(n)$





```
import java.util.ArrayList;  
public class OrderedList<T extends Comparable<T>>  
{
```

```
    //instance variables  
    private ArrayList<T> elements;  
    private int cursor;
```

```
    //create an empty OrderedList  
    public OrderedList()  
    {  
        elements = new ArrayList<T>();  
        cursor=-1;  
    }
```

] constructor

```
    //create an empty OrderedList with a given capacity  
    //another useful constructor  
    public OrderedList(int cap)  
    {  
        elements = new ArrayList<T>(cap);  
        cursor=-1;  
    }
```

] overloaded constructor

```
    //returns size of the list  
    public int size()  
    {  
        return elements.size();  
    }
```

```
    //checks if the list is empty  
    public boolean isEmpty()  
    {  
        return elements.isEmpty();  
    }
```

```
    //clears the list  
    public void clear()  
    {  
        elements.clear();  
    }
```

```
    //get the item at a given index  
    public T get(int pos)  
    {  
        if (pos<0||pos>=elements.size())  
        {  
            System.out.println("Index out of bounds");  
            return null;  
        }  
        return elements.get(pos);  
    }
```

ArrayList get method

```

//Methods first and next are useful to scan the list
//first gets the first item
//next gets the next item (wherever the cursor is)
public T first()
{
    if (elements.size()==0)
        return null;
    cursor=0;
    return elements.get(cursor);
}
public T next()
{
    if (cursor<0||cursor>=(elements.size()-1))
        return null;
    cursor++;
    return elements.get(cursor);
}

//print the list
public void enumerate()
{
    System.out.println(elements);
}

//add an item at a given position
public void add(int pos, T item)
{
    elements.add(pos, item);
}

```

will be used together to scan the ordered list.

↑ arraylist method



```
//binary search
public int binarySearch(T item)
{
```

```
    if (elements.size() == 0)
```

```
        return -1;
```

```
    int lo = 0, hi = elements.size() - 1, mid = 0;
```

```
    while (lo <= hi)
```

```
    {
        mid = (lo + hi) / 2;
```

```
        int c = item.compareTo(elements.get(mid));
```

```
        if (c == 0) return mid;
```

```
        if (c < 0) hi = mid - 1;
```

```
        if (c > 0) lo = mid + 1;
```

```
    }
```

why? ||

```
    if (item.compareTo(elements.get(mid)) < 0)
```

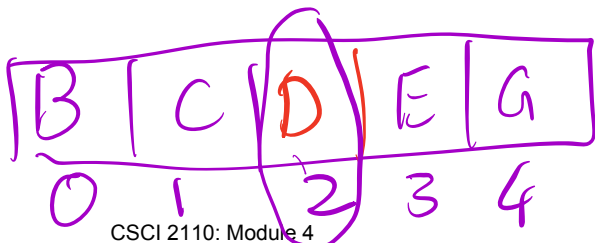
```
        return -(mid + 1);
```

else

```
        return -(mid + 2);
```

}

Example:



binarySearch ("E") → 2

binarySearch ("D") → -3

binarySearch ("F") → 4

binarySearch ("H") → -5

binarySearch ("A") → -1

insert("D") → binarySearch("D") → -3  
 Convert this to a positive integer and subtract 1.

```
//insert an item at the correct position
public void insert(T item)
{
```

```
    if (elements.size() == 0)
    {
        elements.add(item); //trivial case
        return;
    }
```

```
    int pos = binarySearch(item);
```

```
    if (pos >= 0)
    {
        System.out.println("Item already present");
        return;
    }
```

```
    else
        elements.add(-pos-1, item);
```

```
}
```

```
//removes a specified item
```

```
public void remove(T item)
{
```

```
    int pos = binarySearch(item);
```

```
    if (pos < 0)
```

```
    {
```

```
        System.out.println("No such element");
```

```
        return;
```

```
    }
```

```
    else
```

```
        elements.remove(pos);
```

```
}
```

```
}
```

```

//Simple demo to illustrate why we return -(mid+1)) and -(mid+2)) in
//binary search
public class OrderedListDemo
{
    public static void main(String[] args)
    {
        OrderedList<String> names = new OrderedList<String>();
        names.insert("B");
        names.insert("C");
        names.insert("E");
        names.insert("G");
        names.enumerate();
        System.out.println("Search E:" + names.binarySearch("E"));
        System.out.println("Search F:" + names.binarySearch("F"));
        System.out.println("Search H:" + names.binarySearch("H"));
        System.out.println("Search A:" + names.binarySearch("A"));
    }
}

```

```

//Simple orderedlist demo. Reads a text file of names and creates
//and prints the list.
import java.util.Scanner;
import java.io.*;
public class OrderedListDemo1
{
    public static void main(String[] args)throws IOException
    {
        Scanner keyboard = new Scanner(System.in);
        System.out.print("Enter the filename to read from: ");
        String filename = keyboard.nextLine();

        File file = new File(filename);
        Scanner inputFile = new Scanner(file);
        OrderedList<String> names = new OrderedList<String>();
        while(inputFile.hasNext())
        {
            String s = inputFile.nextLine();
            names.insert(s);
        }
        inputFile.close();
        names.enumerate();
    }
}

```

## MERGING ORDERED LISTS

Merging two ordered lists and related operations are important.

Suppose the first ordered list L1 is as follows:

List 1:

Amar	Boris	Charlie	Dan	Fujian	Inder	Travis
------	-------	---------	-----	--------	-------	--------

and the second ordered list L2 is as follows:

List 2:

Alex	Ben	Betty	Charlie	Dan	Pei	Travis	Zola	Zulu
------	-----	-------	---------	-----	-----	--------	------	------

The merging of L1 and L2 should produce the following:

List 3:

Alex	Amar	Ben	Betty	Boris	Charlie	Dan	Fujian	Inder	Pei	Travis	Zola	Zulu
------	------	-----	-------	-------	---------	-----	--------	-------	-----	--------	------	------

Two-finger-walking algorithm

(Pseudo code)

Result list  $L3 \leftarrow$  empty

$f1 \leftarrow 0$  & start of  $L1$

$f2 \leftarrow 0$  & start of  $L2$

while ( $f1 < \text{length of } L1$  &  $f2 < \text{length of } L2$ )

{

if (item at  $f1 < \text{item at } f2$ )

{

append item at  $f1$  to  $L3$  (add to end)

move  $f1$  ( $f1++$ )

}

else if (item at  $f2 < \text{item at } f1$ )

{

append item at  $f2$  to  $L3$

move  $f2$  ( $f2++$ )

}

else

{ append item at  $f1$  to  $L3$

$f1++$

$f2++$

}

}

if ( $f1 == \text{length of } L1$ ) append remaining items  
in  $L2$  to  $L3$

if ( $f2 == \text{length of } L2$ ) append remaining items  
in  $L1$  to  $L3$

