

Assignment 3 - Hady Ibrahim

SFWRENG 2CO3: Data Structures and Algorithms–Winter 2023

Deadline: March 5, 2023

Department of Computing and Software
McMaster University

Please read the *Course Outline* for the general policies related to assignments.

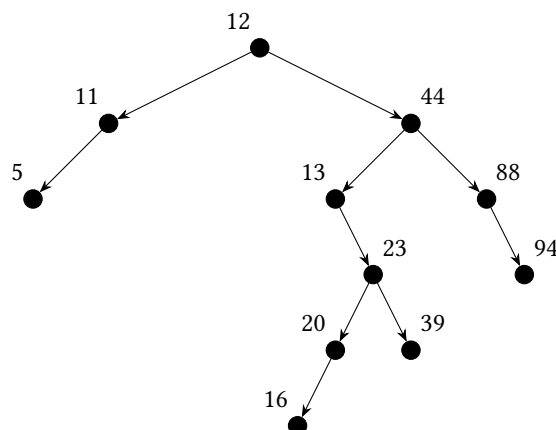
Plagiarism is a serious academic offense and will be handled accordingly.
All suspicions will be reported to the Office of Academic Integrity
(in accordance with the Academic Integrity Policy).

This assignment is an *individual* assignment: do not submit work of others. All parts of your submission *must* be your own work and be based on your own ideas and conclusions. Only *discuss or share* any parts of your submissions with your TA or instructor. You are *responsible for protecting* your work: you are strongly advised to password-protect and lock your electronic devices (e.g., laptop) and to not share your logins with partners or friends! If you *submit* work, then you are certifying that you have completed the work for this assignment by yourself. By submitting work, you agree to automated and manual plagiarism checking of all submitted work.

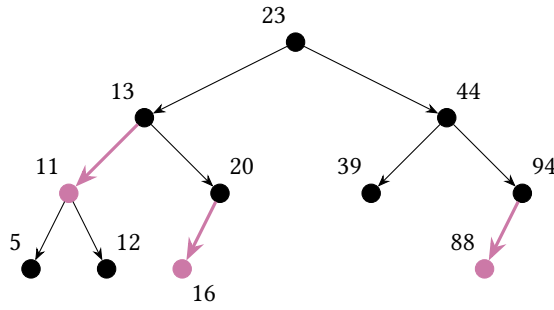
Late submission policy. Late submissions will receive a late penalty of 20% on the score per day late (with a five hour grace period on the first day, e.g., to deal with technical issues) and submissions five days (or more) past the due date are not accepted. In case of technical issues while submitting, contact the instructor *before* the deadline.

Problem 1. Consider the sequence of values $S = [12, 44, 13, 88, 23, 94, 11, 39, 20, 16, 5]$.

P1.1. Draw the binary search tree obtained by adding the values in S in sequence.



P1.2. Draw the red-black tree obtained by adding the values in S in sequence.



P1.3. Consider the hash-function $h(k) = (2k + 5) \bmod 11$ and a hash-table of 11 table entries that uses hashing with separate chaining. Draw the hash-table obtained by adding the values in S in sequence.

P1.4. Consider the hash-function $h(k) = (3k + 2) \bmod 11$ and a hash-table of 11 table entries that uses hashing with linear probing. Draw the hash-table obtained by adding the values in S in sequence.

Problem 2. We say that a hash function $h : \mathcal{U} \rightarrow \mathbb{N}$ that maps values from a set \mathcal{U} to integers in the range $[0 \dots M)$ is n -perfect if there exists at most n distinct values $u_1, \dots, u_j \in \mathcal{U}$ such that $h(u_1) = \dots = h(u_j)$.

P2.1. Consider the hash function $h(k) = (2k + 5) \bmod 11$. Is this hash function 2-perfect for the inputs $0, \dots, 21$? Explain why or why not.

Answer: This function is 2-perfect for the inputs $0, \dots, 21$ since after computation, there is only 2 values mapped to one spot in the hash table. In summary, since the function has $2k$, as you loop through the values, the new value will most often be 2 more than the previous. This is until you hit 11 as it is mod 11, then it will start back at either 0 or 1 depending on if the previous value was 9 or 10. The following is the computation.

v	$h(v)=(2v+5)\bmod 11$	v	$h(v)$	v	$h(v)$	v	$h(v)$
0	5	6	6	12	7	18	8
1	7	7	8	13	9	19	10
2	9	8	10	14	0	20	1
3	0	9	1	15	2	21	3
4	2	10	3	16	4		
5	4	11	5	17	6		

Rearrange to see that there are only two values per mapping. We can see that if there was 1 more value, it would no longer be 2-perfect:

$h(v)$	v	$h(v)$	v	$h(v)$	v	$h(v)$	v
0	3, 14	3	10, 21	6	6, 17	9	13, 2
1	9, 20	4	5, 16	7	12, 1	10	8, 19
2	4, 15	5	11, 0	8	7, 18		

P2.2. Prove that a hash function $h : \mathcal{U} \rightarrow \mathbb{N}$ can only be n -perfect if $|\mathcal{U}| \leq n \cdot M$.

Answer:

Assume $h : \mathcal{U} \rightarrow \mathbb{N}$ is n -perfect and let u_1, \dots, u_j be j distinct values in \mathcal{U} such that $h(u_1) = \dots = h(u_j)$. Then we will have that $j \leq n$.

P4.3. Professor X claims to have developed a data structure `BSETX` to which any binary string W of length $|W| = N$ can be added in $\sim N$. Furthermore, Professor X claims that `BSETX` provides *ordered iteration*: one can iterate over all $M = |S|$ strings in a set S , implemented via `BSETX`, in a lexicographical order in $\sim M + T$ in which T is the combined length of the M strings. Professor X claims that this method of sorting binary strings proves that the worst-case lower bound for sorting M binary strings is *not* $\sim M \log_2(M)$ comparisons. Argue why Professor X is wrong.

We note that strings S_1 and S_2 are *lexicographical ordered*, denoted by $S_1 < S_2$, if S_1 comes before S_2 in an alphabetical sort (e.g., as used in a dictionary). Next, we formalize $S_1 < S_2$ for binary strings: we have $S_1 < S_2$ if S_1 and S_2 are equivalent up to the $0 \leq i \leq \min(|S_1|, |S_2|)$ -th character ($S_1[0] = S_2[0]$, ..., $S_1[i-1] = S_2[i-1]$) and either $|S_1| = i < |S_2|$ or the $(i+1)$ -th character of S_1 comes before the $(i+1)$ -th character of S_2 (in which case $S_1[i] = 0$ and $S_2[i] = 1$). For example, $0 < 00$ and $00 < 01$, but not $100 < 10$.

Assignment Details

Write a report in which you solve each of the above problems. Your submission:

1. must be a PDF file;
2. must have clearly labeled solutions to each of the stated problems;
3. must be clearly presented;
4. must *not* be hand-written: prepare your report in \LaTeX or in a word processor such as Microsoft Word (that can print or exported to PDF).

Submissions that do not follow the above requirements will get a grade of zero.

Grading

Each problem counts equally toward the final grade of this assignment.