

Day - 188Heap - 4

\*  $k^{\text{th}}$  Smallest Element:

$\Rightarrow$ 

10	3	7	4	8	9	2	6
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 ,  $k=4$

$\Rightarrow$  In Brute force approach, we can sort the array & return that  $(k-1)$  element.

$\Rightarrow$  So, T.C.  $\rightarrow O(n \log n)$

S.C.  $\rightarrow O(1)$  if you use Heap sort.

$\Rightarrow$  In the next approach, we can use min heap. Then, pop the top element  $k$  times.

$\Rightarrow$  And return that popped  $k^{\text{th}}$  element.

T.C.  $\rightarrow O(n) + k \log n$   
Creation      Deletion

S.C.  $\rightarrow O(1)$ .

$\Rightarrow$  In the next approach, we will select the 4 elements & find the 4<sup>th</sup> smallest element.

$\Rightarrow$  Now, select the next element and check that element can become the part of that group of 4 by checking it with the largest element.

$\Rightarrow$  If ~~it~~ it is smaller than that, it will come into the group.



⇒ So, we will use max heap. for creating group.

⇒ So, T.C. will be as —

→  $k + (n - k) * 2 \log k$

→  $k + n \log k - k \log k$

→  $O(n \log k)$

\*  $k^{\text{th}}$  Largest Element:

⇒

6	8	2	10	5	7	4	3
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⇒ We will do same things as we do in last question.

⇒ But instead of using max heap, we will use min heap.

⇒ And if the upcoming is greater than the top element, we will add it to the heap.

⇒ We can also solve this question by using quick select method that uses quick sort. (Explain it on their own).

\*  $k^{\text{th}}$  Largest Element in a stream:

1	2	3	4	5	6
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- ⇒ We will have a stream of numbers & in that stream, we have to find  $k^{\text{th}}$  largest element.
- ⇒ Suppose, the value of  $k = 4$  then in the starting 1 will come but here there are only one element, so answer is -1.
- ⇒ In the same, we have to find all the answers & return it in vector.

\* Sum of element b/w  $k_1$  &  $k_2$  smallest element:

20	8	22	4	12	10	14
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- ⇒ First approach, we can sort the array and then find the answer.
- ⇒ We don't have to include  $k_1$  &  $k_2$ .
- ⇒ In the second approach, we will make two heap that contains  $k_1$  smallest elements &  $(k_2 - 1)$  smallest elements.
- ⇒ Then do the sum of both heaps.
- ⇒ Now, subtract both the sum & return the result.

T.C. →  $O(n \log k_2)$

S.C. →  $O(k_1 + k_2)$ , if you are using priority queue or  $O(1)$  if you are using ~~bottom~~ step-down approach.