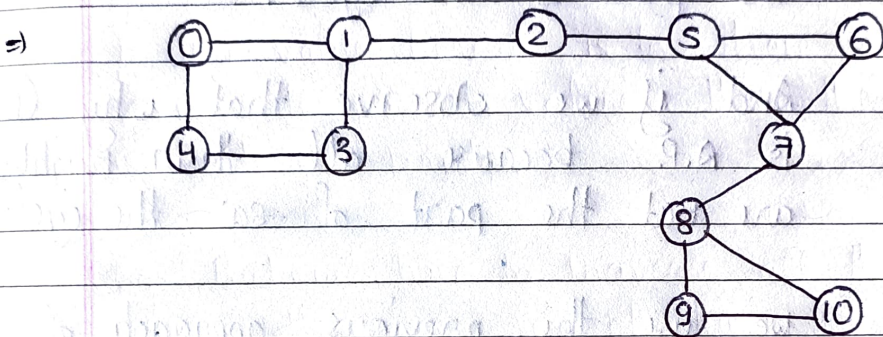


Day - 219Graph - 23\* Articulation Point?

⇒ Articulation Point is that node if we removed that node from the graph then our graph will be divided into two components.

⇒ So, here node — ①, ②, ⑤, ⑦, etc are A.P.

⇒ If we normally find every A.P.s then our T.C. will become very high.

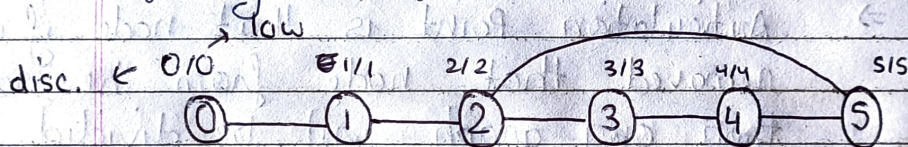
⇒ So, how to solve the question in good time complexity?



⇒ we can observe that, why ④ is not A.P. because it is the part of cycle as we can say all the neighbours of ④ are the part of the cycle.

⇒ And if we observe that why ① is A.P. because all their neighbours are not the part of the cycle.

⇒ We use the previous approach of finding bridges.



⇒  $disc[Node] \leq low[neigh]$

⇒  $\hookrightarrow$  Articulation - Point

⇒ Disc (discovery time)

⇒ At what time, the node is discovered.

⇒ Low: It store the lowest discovery time reachable from current node.



- $\Rightarrow$  so, how to start?
- $\Rightarrow$  we will apply DFS and fill the disc & low value.
- $\Rightarrow$  Then, if we get any visited node, then it will not be a AP.
- $\Rightarrow$  Also, we will compare the low of node with disc. of neigh. & select the min.
- $\Rightarrow$  When we are returning back then we check for — AP ( $\text{disc}[\text{node}] \leq \text{low}[\text{neigh}]$ )  
 $\text{low}[\text{node}] = \min(\text{low}[\text{node}], \text{low}[\text{neigh}])$
- $\Rightarrow$  We ~~if~~ will don't check for that node that parent is -1.
- $\Rightarrow$  If the root node have more than one child then the root node is A.P.
- $\Rightarrow$  Sometimes, AP can be counted twice so, that we can maintain a AP vector.