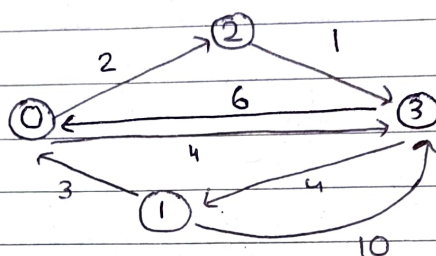
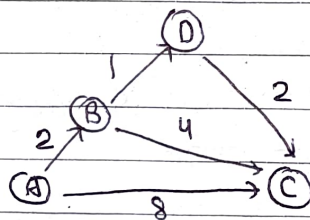


Day-211Graph-15* Floyd Warshall Algorithm:

⇒ This algorithm is used for finding multiple source shortest path.

⇒ It is only applied to directed graph.

⇒ Let's take an example,



⇒ Suppose, we want to go to C from A, then, $A \xrightarrow{8} C$

But, $[A \xrightarrow{2} B, B \xrightarrow{4} C] = A \xrightarrow{6} C$

⇒ We can also take the above path.

This is logic behind FWA

- ⇒ FWA is based on the property that if there is a path between A & B ~~and~~ and B & C.
- ⇒ Then, there must be path from A to C by going through B.
- ⇒ So, this thing or property is called Transitivity.
- ⇒ It is given by Warshall.
- ⇒ And the Floyd main task is to update the distance by choosing min distance.
- ⇒ So, for $A \rightarrow D$,
 $A \xrightarrow{2} B$, $B \xrightarrow{1} D$
 $A \xrightarrow{3} D$
- ⇒ This algo works on a very important principle that →
- ⇒ Small milate jao, large banate jao 😊 😊 (Rohit bhaiya rocks 🙌)
- ⇒ So, here comes a concept of intermediate nodes.
- ⇒ we will consider every other node as intermediate node and check for the

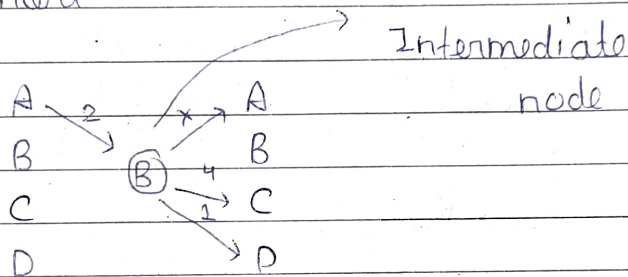
target node by indir. node.

=> Ex: $A \rightarrow C$

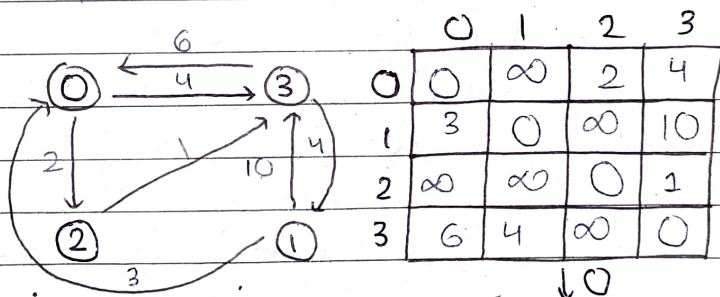
$A \xrightarrow{2} B, B \xrightarrow{4} C$ (B as IN)

$A \xrightarrow{3} B, B \xrightarrow{1} D, D \xrightarrow{2} C$ (D as IN)

=> So, in short —



=> Ex:



=>

So,

	i	k	j	0	1	2	3
0	0	1	0	0	∞	2	4
1	1	0	1	3	0	5	7
2	2	∞	2	∞	∞	0	1
3	3	6	4	8	0		

=>

$\text{path}(i, j) = \min(\text{path}(i, j), \text{path}(i, k) + \text{path}(k, j))$

	0	1	2	3	0	1	2	3		0	1	2	3		
0	0	7	2	3	0	0	∞	2	3	0	0	∞	2	4	
1	3	0	5	6	1	3	0	5	6	2	1	3	0	5	7
2	7	5	0	1	2	∞	∞	0	1	2	∞	∞	0	1	
3	6	4	8	0	3	6	4	8	0	3	6	4	8	0	

⇒ In FWA, ordering doesn't matter of taking intermediate node.

⇒ Suppose, we are taking '2' as a ~~in~~ intermediate node between i & j then it is also possible that other intermediate nodes are present.

Code

```
for(int k=0; k<n; k++){  
    for(int i=0; i<n; i++){  
        for(int j=0; j<n; j++){  
            mat[i][j] = min(mat[i][j],  
                             mat[i][k] + mat[k][j])  
        }  
    }  
}
```

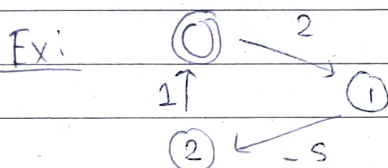
Time Complexity $\rightarrow O(V^3)$

Space Complexity $\rightarrow O(1)$

⇒ FWA can also detect negative cycle and can also work for -ve weights.

\Rightarrow Set if the diagonal element that we always have zero value but have any -ve value.

\Rightarrow That means -ve cycle present in the graph.



\Rightarrow $0 \rightarrow 1$ (in (but by 2 as $2N$))

\Rightarrow $1 \xrightarrow{-5} 2 \xrightarrow{3} 1 \rightarrow -2$

\Rightarrow $1 \xrightarrow{-2} 1$ (but self distance can't be -ve)

\Rightarrow So, -ve cycle present.

Dijkstra: Only for +ve

	<u>Sparse</u>	<u>Dense</u>
\Rightarrow	$O(E \log V) * V$	$O(V^2) * V$
\Rightarrow	$O(VE \log V)$	$O(V^3)$
\Rightarrow	$V^2 \log V$	$O(V^3)$

\Rightarrow FW: Easy to implement.