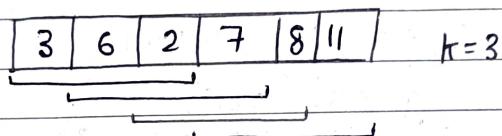


Day - 144Queue - 3

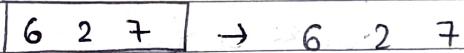
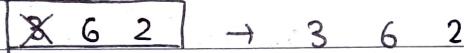
- \* Print all no. in every window of size  $k$ :



$\Rightarrow$  3 6 2  
6 2 7  
2 7 8  
7 8 11

$\Rightarrow$  we have to solve this question by using queue.

$\Rightarrow$  So,



and so on

$\Rightarrow$  So, we are popping out first element & pushing ~~the~~ the next element.

$\Rightarrow$  First, we will push  $k-1$  elements in the queue.

⇒ After that, we will push the elements & print that pop the first element & again push next element & so on.

Code

```
queue< int> q;
for( i=0; i<k-1; i++){
    q.push( arr[i] );
}
for( i=k-1; i<n; i++ ){
    q.push( arr[i] );
    display( q );
    q.pop();
}
```

}

```
void display( queue< int> q ) {
    while( !q.empty() ){
        cout << q.front() << " ";
        q.pop();
    }
    cout << endl;
}
```

}

\* First negative integer in every window of size k:

0	1	2	3	4	5	6	7
2	-3	-4	-2	7	8	9	-10
-3			8				
-3		-4		-2		0	

$k = 3$

- ⇒ Now we will do that we have done in the previous question & we will get three elements.
- ⇒ And then we will print the negative element

Code

```
queue<int> q;
for(i=0; i<k-1; i++){
    q.push(arr[i]);
}
vector<int> ans;
for(i=k-1; i<n; i++){
    q.push(A[i]);
    ans.push_back(display(q));
    q.pop();
}
return ans;
```

```
int display(queue<int>q){
    while(!q.empty()){
        if(q.front()<0){
            return q.front();
        }
    }
}
```

```
q.pop();
}
return 0;
```

T.C. → O(nk)

⇒ Now, we have to solve it in  $O(n)$ .

⇒ We will push only -ve no. into the queue.

⇒ Now, we will return the front element of the queue as first -ve no.

⇒ When, we solve this question,

we will store indexing of the elements so it will be easier to check that element will in that window or not.

⇒ For checking, we will do that —  
 $i = 3(k)$

⇒ If the index is equal to above value or less than we will pop it.

### Code

```
queue <int> q;
for( i=0; i<k-1; i++) {
    if( A[i] < 0)
        q.push(i);
}
```

```
vector <int> ans;
for( i=k-1; i<n; i++) {
    if( A[i] < 0)
        q.push(i);
```

Date \_\_\_\_\_

Page \_\_\_\_\_

```
if(q.empty())
    ans.push_back(0);
else {
    if(q.front() == i - k)
        q.pop();
    if(q.empty())
        ans.push_back(0);
    else
        ans.push_back(A[q.front()]);
}
return ans;
```

\* First non repeating char in a stream of chars:

A = "aababdc"  
⇒ If the char is not repeated yet in that stream then add it to B.  
⇒ If there is no such char then add '#'.  
B = "aab#dd"

A = "abcacdbd"  
B = "aaabbdd#"

⇒ Now, we will solve it by optimized method.

- ⇒ We will take an ~~elem~~ array that stores the repetition of the elements.
- ⇒ We will use the queue to store the non-repeated chars.
- ⇒ If queue becomes empty then add '#'.
- ⇒ Now, we will check every char that will come that it can be the answer or not.
- ⇒ We will only push that char that has not come yet.

### Code

```

string B = "";
vector<int> repeated(26, 0);
queue<int> <char> q;

for(int i = 0; i < A.size(); i++){
    if(repeated[A[i] - 'a'] >= 1){
        repeated[A[i] - 'a']++;
        while(!q.empty() && repeated[q.front() - 'a'] > 1)
            q.pop();
        if(q.empty())
            B += '#';
        else
            B += q.front();
    }
}

```

Date \_\_\_\_\_

Page \_\_\_\_\_

else{

    repeated[A[i] - 'a']++;

    q.push(A[i]);

    while(~~!q~~ repeated[q.front() - 'a'])

        q.pop();

    B += q.front();

}

}

return B;

>1)