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## 2D - Arrays

- => As we have seen excel sheets.
- => So, if we want to store the data of excel sheet then with the help of 2D Arrays, we can easily do this.
- => Let assume, we want a 4 row & 3 column table, then -
- int arr[4][3]; → 2D Array

	0	1	2
0	00	01	02
1	10	11	12
2	20	21	22
3	30	31	32

- => So, now, how 2D Array store in memory.
- => It store in two ways in memory - as row major & column major.

Row		00	01	02	10	11	12	
major	20	21	22	30	31	32		→ memory

- 2) Store row wise in contiguous manner.

0	1	2	3	4	5	6	7	8	9	10	11
00	01	02	10	11	12	20	21	22	30	31	32

Memory



⇒ Now, if we want to calculate the position of any element then —

$$\text{Index} = \text{row\_index} \times \text{col} + \text{col\_index} \quad \text{of memory} \quad \text{--- ①}$$

⇒ Now, if we have index, then we can calculate the indexing of the element.

$$\text{row\_index} \rightarrow \text{Index} / \text{col};$$

$$\text{col\_index} \rightarrow \text{Index} \% \text{col};$$

⇒ As we know,  $0 \leq \text{col\_index} < \text{col}$   
So, divide the eq<sup>n</sup> ① by col.

$$\frac{\text{Index}}{\text{col}} = \frac{\text{row\_index} \times \text{col}}{\text{col}} + \frac{\text{col\_index}}{\text{col}}$$

⇒ So, row\_index =  $\text{Index} / \text{col}$ .

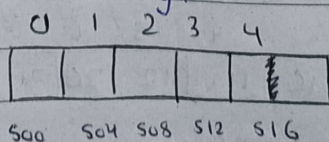
⇒ For col\_index,

$$(\text{Index}) \% \text{col} = (\text{row\_index} \times \text{col} + \text{col\_index}) \% \text{col}$$

⇒ So, col\_index =  $\text{Index} \% \text{col}$ .

⇒ `int arr[5];`

Here, arr store the starting address of the array.





=> If you want to calculate the address of particular index —  
address

$$\begin{aligned} \text{arr}[3] &= \text{base address} + \text{index} \times \text{size of element} \\ &= 500 + 3 \times 4 \\ &= 512. \end{aligned}$$

=> Same thing happens in 2D-Array —

$$\begin{aligned} \text{arr}[2][1] &= \text{base address} + \text{index} \times \text{size of element} \\ &= 500 + 7 \times 4 \\ &= 528 \end{aligned}$$

$$\text{arr}[i][j] = \text{base address} + (i \times \text{col} + j) \times \text{size of element}.$$

Initialize

$$\text{int arr}[4][3] = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12\}$$

Update or user i/p

$$\text{arr}[3][0] = 15$$

$$\text{cin} \gg \text{arr}[3][1];$$

=> Printing the elements

```
for(i = 0; i < row; i++)
    for(j = 0; j < col; j++)
        cout << arr[i][j];
```



\*

Search element in Array

```

int arr[4][3]; int x = 7;
int main() {
    int arr[4][3] = {1, 2, 3, ..., 123};
    int x = 7;
    for (int i = 0; i < 4; i++)
        for (int j = 0; j < 3; j++) {
            if (arr[i][j] == x)
                cout << "Yes";
        }
    cout << "No";
}

```

⇒

When we pass 2D Array in any function call then we have to must pass the col value.

```

void print (int arr[][4]) {
    ;
}

```

\*

Add 2 Matrix

2	3	4	5
6	7	8	9
10	11	12	13

arr1[3][4]

2	3	4	5
6	7	8	9
10	11	12	13

arr2[3][4]

⇒

$$arr[i][j] += arr1[i][j] + arr2[i][j];$$



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T.C.  $\rightarrow O(\text{row} * \text{col})$ S.C.  $\rightarrow O(\text{row} * \text{col})$ 

\*

Print row index with Max sum

	0	1	2	
0	1	2	3	$\rightarrow 6$
1	4	5	6	$\rightarrow 15$
2	7	8	9	$\rightarrow 24$ ✓

```

int sum = INT_MIN, index = 0;
for (int i = 0; i < row; i++) {
    int total = 0;
    for (int j = 0; j < col; j++)
        total += arr[i][j];
    if (sum < total) {
        sum = total;
        index = i;
    }
}

```

}

cout &lt;&lt; index;

T.C  $\rightarrow O(\text{row} * \text{col}) \xrightarrow{w} B.C.$   
 ~~$O(1) \rightarrow B.C.$~~

\*

Print sum of diagonal Element



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	0	1	2	3
0	5 <sub>00</sub>	8	3	9 <sub>03</sub>
1	6	2 <sub>11</sub>	8 <sub>12</sub>	4
2	5	3 <sub>21</sub>	2 <sub>22</sub>	2
3	2 <sub>30</sub>	8	1	9 <sub>33</sub>

⇒ Condition for this question is  $row == col$ .  
 ⇒ Here, we have two diagonal —  
 So, for calculating sum of first diagonal  
 ⇒ `int first = 0;`  
   for (int i = 0; i < row; i++) {  
     first += arr[i][i];

  int second = 0;  
   int i = 0; j = col - 1;  
   while (j >= 0) {  
     second += arr[i][j];  
     i++; j--;  
   }

T.C. →  $O(row)$ .

\* Reverse each row of matrix

	0	1	2			0	1	2
0	1	2	3		→	0	3	2
1	4	5	6			1	6	5
2	7	8	9			2	9	8



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Code

```
for( int i = 0; i < row; i++) {  
    int start = 0, end = col - 1;  
    while( start < end ) {  
        swap( arr[i][start], arr[i][end] );  
        start++; end--;  
    }  
}
```