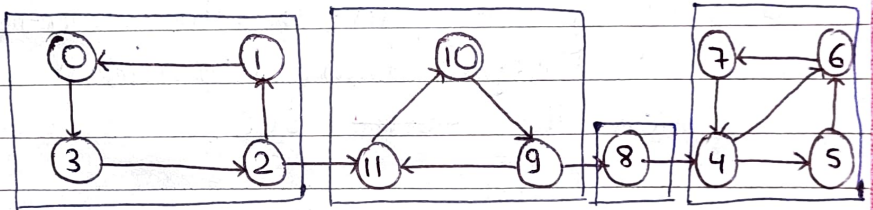


Day - 217

Graph - 21

* Strongly Connected Component:

=> In a directed graph, SCC is a subset of vertices where every vertex in the subset is reachable from every other vertex in the same subset by traversing direct edges.



=> So, here, we can see that 0 & 3 are strongly connected because we can go from 0 to 3 and from 3 to 0.

=> So, in this way, we have four strongly connected component in the group.

=> So, if we apply BFS or DFS then we ~~can~~ will traverse all the nodes in one go.

=> But we don't want this.

⇒ So, we can imagine the graph as -



⇒ So, if we change the direction ~~or~~ between two strongly component then we can solve the problem.

⇒ There will be cycle present in every SCC.

⇒ So, if we reverse the direction of the edges then also no impact on the ~~between~~ SCC.

⇒ So we will use this property.

⇒ We will follow the order of SCC to apply DFS.

⇒ So, how to achieve the order.

⇒ Order is just saying that SCC1 will be visited before SCC2.

⇒ So as we know, we have topological sort that can help in this situation.

- ⇒ So, if we use DFS method of T.S then we get ~~an~~ a stack of traversed nodes.
- =1 Also, ~~remember~~ remember that T.S. sort don't give correct answer when cycle is present in the group.
- =1 But here, it can work.
- ⇒ we use that stack of T.S for finding our answer.
- ⇒ At the time of T.S, the graph will be normal.
- =1 After that, we reverse the direction of the edges.
- ⇒ And then T.S stack will help in getting the order.
- ⇒ This whole algo is known as Kosaraju's Algorithm.
- ⇒ T.S helps in confirming that ② will always come before ①.
- ⇒ After that we start apply DFS and if we get any unvisited node from TS stack then we increase the SCC value

by 1.

=1 And apply DFS on this node.

$$T.C. \rightarrow O(V+E)$$

$$S.C. \rightarrow O(V+E)$$