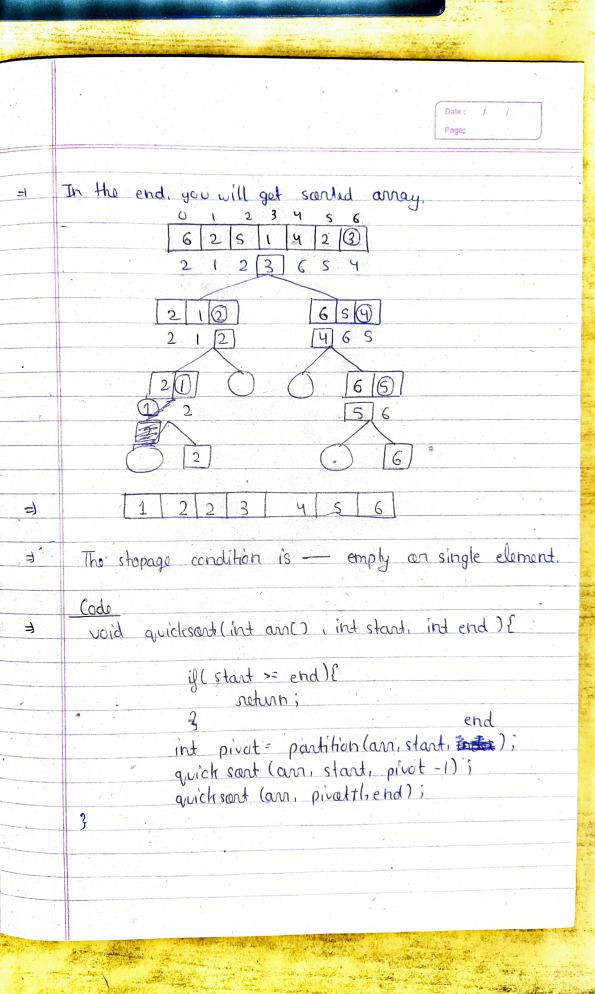
Day - 85 Quick sard It is a sarting algorithm. 6 4 2 8 13 7 11 9 3 6 First we have to select a pivot element. 4 Pivot: Any element that is on their correct position in the sorted array. For finding the cornect position of the pivoti first we  $\Rightarrow$ to find how many elements are equal on less than pivat element. suppose, in the above ex, we select 6 as pivot elem. thon & 4 elements are equal on less than 6. 23 6 8 13 7 119 = 3 6 9 6 8 1 59 3 13 11 3 46 6 7 8 9 11 137 sorted array. After putting the proof element on the correct 7 position. Then, you will have all the elements less than pivot on left & greater than pivot on right Now again do this same thing with these has コ halves.



int partition (int ara(), int start; int end) { int pos = stanti for (inti = startis i rend; i++){. if (anci) = ancend) { swap (anci), an (pos) bi pos ++i return pos-1; Time Complexity: log th 1 ← At every level, N+ N+N+N+ .-- + N (NIOgN) But this is the aug. time complexity of Q.S. =1 S.C. - log N ('in aug. case) · =)

For this ex. - 6 5 4 3 2 1 The cond. for Q.S. is worst case. N+(N-1)+(N-2)+\_+1 =  $O(N^2)$ The warst case is when our array is in descending ander order & we have to sort it in ascending order. 2 S.C. - O(N). 1 Also, when our array is already sorted then the case is worst case. 1