

Day - 196Sliding Window - 4

* Length of longest subarray with atmost k frequency:

1	2	3	1	2	3	3	4	1	2	$k=2$
---	---	---	---	---	---	---	---	---	---	-------

⇒ Ex:

1	2	3	1	2	3
---	---	---	---	---	---

 → size = 6

1	2	3	3	4	1	2
---	---	---	---	---	---	---

 → size = 7

⇒ Basic approach will be same.

⇒ Now, we will use the concept of sliding window protocol.

⇒ Start from taking as first element in the window & check the frequency.

⇒ If the frequency is less than k then increase the window size otherwise decrease the window size from starting.

Code

```
unordered_map<int, int> m count;  
int start = 0, end = 0, len = 0;  
while (end < n) {  
    count[nums[end]] ++;
```



```

len = max
while
while( count[nums[end]] > k ) {
    count[nums[start++]]--;
}
len = max(len, end - start + 1);
end++;
}
return len;

```

* Count subarray where max element appear at least k times

1	2	3	2	3	1	2	3	3	2
---	---	---	---	---	---	---	---	---	---

k = 2

⇒ Ex:

1	2	3	2	3
---	---	---	---	---

⇒ max element → 3 → 2 times

⇒ we have to count all these types of subarrays

⇒ We will SWP approach.

⇒ so, first we get a window that satisfy the conditions. Now, if we add other elements in the window the result will be same.

⇒ so, we can directly add all these subarrays. After that, we will decrease the element from starting from the window, then again check for conditions.

- ⇒ And we will repeat this process.
- ⇒ If the window is not valid then we increase the window.
- ⇒ So, first we increase the window size until $\text{count} == k$.
- ⇒ Then decrease the window size.

Code

```
int start = 0, end = 0, maxEle, count = 0,
total = 0;
while (end < n) {
    if (nums[end] == maxEle)
        count++;
    while (count == k) {
        total += n - end;
        if (nums[start] == maxEle)
            count--;
        start++;
    }
    end++;
}
return total;
```


* Subarray with k different Integer:

1	2	1	2	3
---	---	---	---	---

 $k=2$

- ⇒ We have to return good sub arrays.
 ⇒ Good subarrays are those arrays that have k different integer.

Ex:

2	1	2
---	---	---

1	2	3
---	---	---

 →

1 → 1 ✓ 1 → 1 (more
 2 → 2 2 → 2 3 → 2)

- ⇒ Here, we directly ~~don't~~ can't use the sub approach. Using this, can result in missing of some answers.

- ⇒ So, we can redefine our problem statement as Total no. of subarray, atleast k diff int.

- ⇒ Atleast (2): $\text{sub}(2) + \text{sub}(3) + \text{sub}(4) + \text{sub}(5)$
 Atleast (3): $\text{sub}(3) + \text{sub}(4) + \text{sub}(5)$
 = $\text{sub}(2)$

- ⇒ So, by using this approach, we can find out exactly ~~is~~ (k).

- ⇒ For finding atleast n we will use the previous question approach.