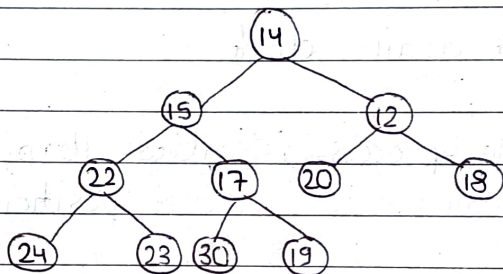


Day - 183Heap - 2* Build Heap:⇒ We have created a heap in $O(n \log n)$ T.C in previous method.

arr →

0	1	2	3	4	5	6	7	8	9	10
14	15	12	22	17	20	18	24	23	30	19

⇒

30	24	20	22	23	12	18	14	17	15	19
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⇒ We will start from index $\rightarrow 1$ then check with parent $((i-1)/2)$ and if child is greater than parent, swap them.

⇒ This method will also take $O(n \log n)$ T.C.

⇒ This is called step up.

⇒ Now, we will do step down.

⇒ Here, parent will compare with its children.

⇒ Parent $\rightarrow i$
 \downarrow

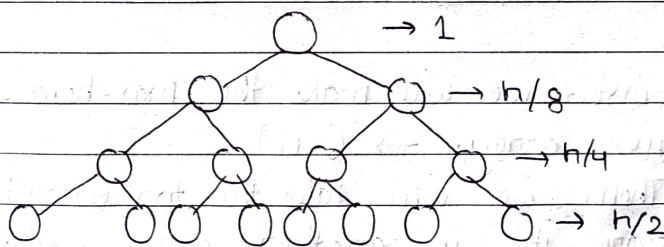
child \rightarrow left $\rightarrow 2*i+1$
 \rightarrow right $\rightarrow 2*i+2$

⇒ In this method, we will start with that node that have children, so we don't have to compare leaf nodes.

⇒ So, starting index will be —

$$\frac{(n-1)-1}{2} \Rightarrow \frac{n}{2} - 1$$

Time Complexity for step up:



⇒
$$\frac{n}{2}(\log n) + \frac{n}{4}(\log n) + \frac{n}{8}(\log n) + \dots$$

⇒
$$(n \log n) \left(\frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots \right)$$

⇒
$$n \log n (1) \Rightarrow n \log n$$

Step down:

⇒ So, at $n/4$, every node can come down one time only.

⇒ In the same way, ~~we~~ as the level decreases then the no. of down time increases.

$$\Rightarrow \frac{n}{4} \times 1 + \frac{n}{8} \times 2 + \frac{n}{16} \times 3 + \dots \log n$$

$$\Rightarrow n \left[\frac{1}{4} + \frac{2}{8} + \frac{3}{16} + \dots \right]$$

$$\Rightarrow O(n)$$

* Heap Sort:

arr \rightarrow

10	11	14	8	3	19	12
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⇒ First, we will make the max-heap of the given array. $\rightarrow O(n)$

⇒ Then, we will take the top element, store it in another array, delete it & replace it with last element.

⇒ After that, put the first element at their correct position.

⇒ So, here, we are using an extra array.

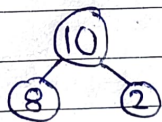
- ⇒ So, how can we solve it by using $O(1)$ S.C.
- ⇒ We know that after changing the top element with last element, there will be a space remain in the array.
- ⇒ So we will store the top element in the last element position.

* Priority Queue:

- ⇒ In PQ, every element have some priority.
- ⇒ We can also say that, there is some type of priority condition in the queue.

⇒ Ex: If we have 2, 8, 10 & the priority condition is that the biggest no. serve the first then 10 will be served first.

⇒ And this is what ... this is max heap.



- ⇒ Priority Queue is used in computer Network, OS, etc.