

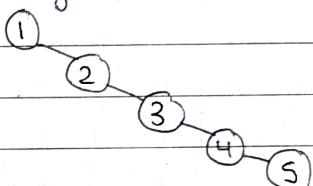
Day - 173

AVL Tree* AVL Tree:

tree.

⇒ First, we have to understand, why we need AVL tree.

⇒ In worst case, BST T.C. for various operations are —

Insertion: $O(n)$ Search: $O(n)$ Deletion: $O(n)$ 

⇒ Creation time:

Array in advance $\rightarrow n \log n$ Runtime $\rightarrow n^2$ ⇒ So, in AVL tree, we can do all three operations in $O(\log n)$ T.C.

⇒ AVL tree don't grow only on one side.

⇒ If any side is growing only then it is self balanced that side.

⇒ So, AVL tree is called as self balanced Binary Search Tree.

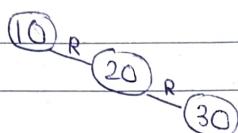
⇒ Every node should have balance between $-1 \leq \text{balance} \leq 1$

⇒ Suppose, we have three elements —

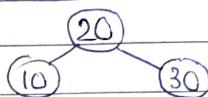
10 20 30

⇒

In BST



In AVL



⇒ So, if we have three nodes then the permutation will be —

10 20 30

30 20 10

30 10 20

10 30 20

20 30 10

20 10 30

we have to balance

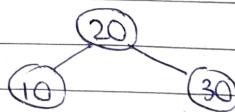
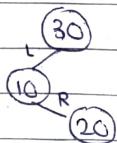
these AVL

trees.

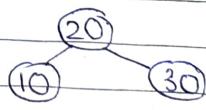
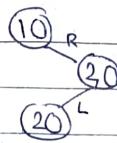
Balanced AVL

trees.

⇒



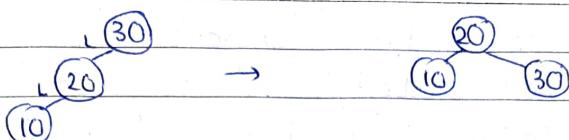
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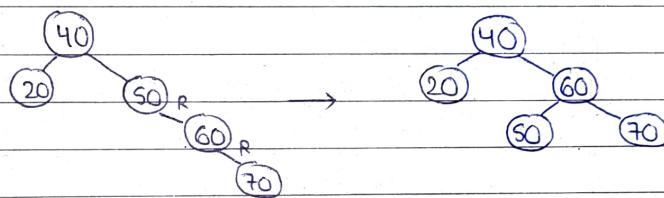
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We will only have these above cases.

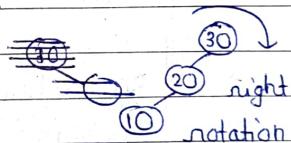
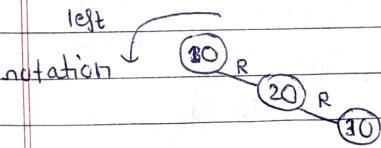
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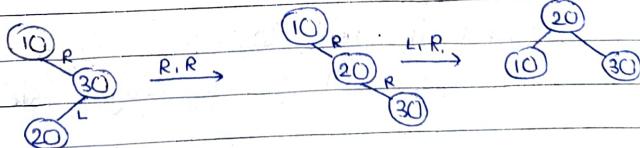
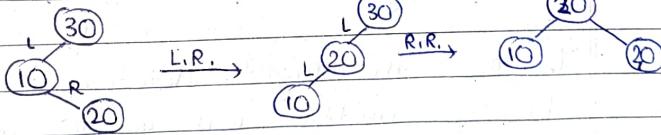
balanced

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We have two rotations —
left & right rotation.



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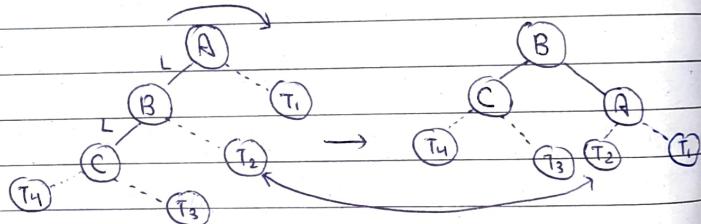


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L.R. → Left Rotation
R.R. → Right Rotation

- ① LL Case: Right Rotation (Top)
- ② RR Case: Left Rotation (Top)
- ③ LR Case: Left Rotation (Middle)
Right Rotation (Top)
- ④ RL Case: Right Rotation (Middle)
Left Rotation (Top)

⇒ There is a edge case remains here —



- ⇒ So, where will T_2 go?
- ⇒ T_2 is greater than B but less than A.
- ⇒ So, it will go to left side of A.

- ⇒ Therefore, in right rotation, right child of middle will go to left side of top.
- ⇒ And in left rotation, left child of middle will go to right side of top.
- ⇒ So, how will we know that it is which case?

- ⇒ For this, if Balance Factor (BF) =
left height - right height —
- ⇒ BF > 1 → Left case (LL or LR)
- ⇒ BF < -1 → Right Case (RR or RL)
- ⇒ For LL & LR, check where that coming node connects in the tree.
- ⇒ If on the left side then LL otherwise LR.
- ⇒ same thing for RR & RL.
- ⇒ calculating height at every time will increase the T.C.
- ⇒ So, we will create a height variable in node class.
- ⇒ So, when we return, we have to update the height.

⇒ Time Complexity
 Creation $\rightarrow O(\log n) * n$
 $\Rightarrow O(n \log n)$
 Deletion $\rightarrow O(\log n)$
 Searching $\rightarrow O(\log n)$