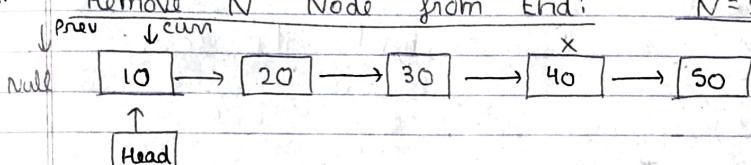


Day - 116

Linked List - 4

* Remove N^{th} Node from End: $N = 2$



- ⇒ First, we count the nodes then minus the value of k from it.
- ⇒ After that again traverse the LL and reduce the value of count by 1.
- ⇒ when the value becomes 0 then remove that current node.
- ⇒ So, we require two pointers curr & prev.

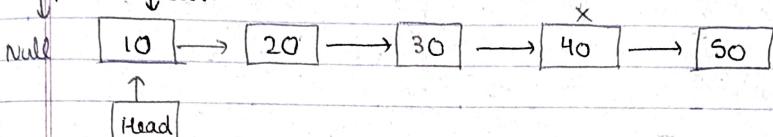
```

Code   int count = 0;
       ListNode * temp = head;
       while (temp != NULL) {
           count++;
           temp = temp → next;
       }
       count -= N;
       ListNode * curr = head, * prev = NULL;
       while (count--) {
           prev = curr;
           curr = curr → next;
       }
       prev → next = curr → next;
       delete curr;
  
```

Day - 116

Linked List - 4

* Remove N^{th} Node from End! $N = 2$



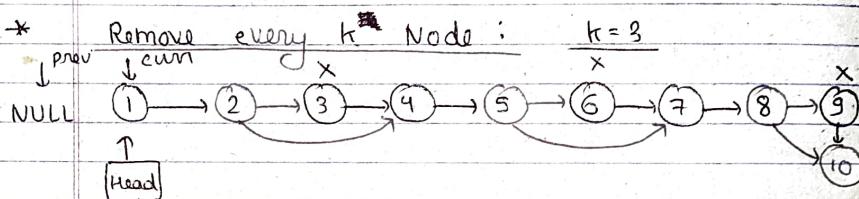
- ⇒ First, we count the nodes then minus the value of N from it.
- ⇒ After that again traverse the LL and reduce the value of count by 1.
- ⇒ When the value becomes 0 then remove that current node.
- ⇒ So, we require two pointers curr & prev.

Code

```

int count = 0;
ListNode * temp = head;
while (temp != NULL) {
    count++;
    temp = temp → next;
}
count -= N;
ListNode * curr = head, * prev = NULL;
while (count-- > 0) {
    prev = curr;
    curr = curr → next;
}
prev → next = curr → next;
delete curr;
    
```

- ⇒ There will be a edge case here when $N = 5$,
- ⇒ Because our code will not work for $N = 5$.
- ⇒ Because $\text{prev} \rightarrow \text{next}$ of prev will not exist.
- ⇒ So, we also have to handle this case.



- ⇒ So, in this question, we also require two pointers — prev & curr .

Code

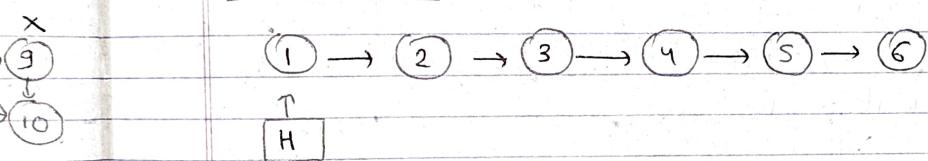
```

Node *curr = head, *prev = NULL;
int count = 1;
while(curr){
    if(k == count){
        prev->next = curr->next;
        delete curr;
        curr = prev->next;
        count = 1;
    } else {
        prev = curr;
        curr = curr->next;
        count++;
    }
}
return head;
  
```

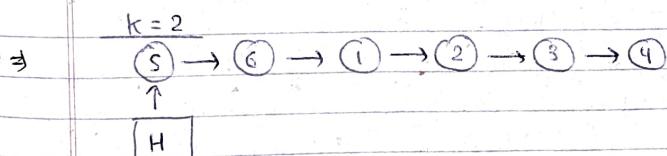
=> When the value of $k=1$, then this is a edge case —

=> if ($k==1$)
return & NULL;

* Rotate List: $k=2$



=> Rotate the List k times i.e. —



=> Our first task is to reach that k node from last.

=> For that, first we have to count the nodes then simply minus the value of k then traverse to that count value.

=> Then, we will break the bond by using curr & prev pointers.

=> We will also make prev's next to NULL.

=> After that traverse to the end of LL by tail pointer.

=> Then connect tail with first pointer & make head equal to curr.
And return head.

Code

```

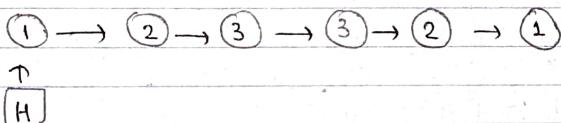
⇒ int count = 0;
ListNode *temp = head;
while( temp ){
    count++;
    temp = temp → next;
}
k = k % count; // For edge case no. 4.
count -= k;
ListNode *curr = head, *prev = NULL;
while( count-- ){
    prev = curr;
    curr = curr → next;
}
prev → next = NULL;
ListNode *tail = curr;
while( tail → next != NULL ){
    tail = tail → next;
}
tail → next = head;
head = curr;

```

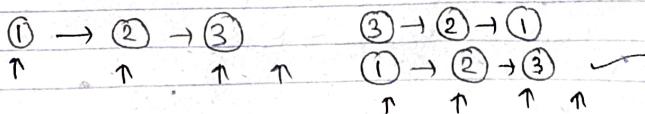
- ⇒ Edge cases for this question are →
if head is NULL then return head.
if k = 1 then return ~~NULL~~, head.
if k = 0 then return ~~NULL~~, head.
And if the value of k is more than no. of nodes then —
Rotate the list by ' $k \% \text{count}$ ' times.

- => We can also solve it by another method —
- => First we traverse the node ~~that~~ and store all the values in a array.
- => Then, the k values from end, we write in the starting. And the remaining values we write in the remaining linked list after starting.
- => But, it's also possible that this solution don't pass on the ~~Leetcode~~ Leetcode.com (NG).

* Check Palindrome:



- => First, we traverse the ~~array~~ array and store all the elements in the array.
- => After that, by using two pointers approach.
 $\downarrow \quad \uparrow$
 $1 \boxed{2} 3 3 2 1$
- => Here, we are using extra space.
- => But, we want to solve it without using extra space.
- => Now, first we break it into two pieces then reverse the second part.
 $\uparrow \quad \uparrow \quad \uparrow \quad \uparrow$
 $\uparrow \quad \uparrow \quad \uparrow \quad \uparrow$
- => And then check both the parts are same or not.



- => Then Palindrome.

⇒ List Node * temp = head;
 int count = 0;
 while(temp){
 count++;
 temp = temp → next;
 }

count /= 2;
 List Node * curr = head ; * prev = NULL;
 while(count --){
 prev = curr;
 curr = curr → next;
 }

prev → next = NULL;
 prev = NULL;
 List Node * front = NULL;
 while(curr){
 front = curr → next;
 curr → next = prev;
 prev = curr;
 curr = front;
 }

List Node * head1 = head; * head2 = prev;
 while(head1){
 if(head1 → val != head2 → val){
 return 0;
 } head1 = head1 → next;
 head2 = head2 → next;
 }
 return 1;