Page ! Date / / \* Anticulation Point Articulation Point is that node if we removed that node from the graph then our graph will be divided into two components. So, here node - D. D. S. E. etc are A.P. If we hormally find every A.P.s. then our T.E. will become very So, how to solve the question in good time complexity?

	& and a second s
=)	we can observe that, why 9 is not A.R. because it is the
	part of cycle on we can say
	part of cycle on we can say all the neighbours of 4 are
	the part of the cycle.
<u> </u>	$(\mathcal{S}) = (\mathcal{S}) = ($
4	And if we observe that why (1) is A.P. because all their heighbours are not the part of sa: the cycle.
	is A.P. because all their neighbours
	are hot the part of sa' the cycle.
<b>a</b> 0	we use the previous approach of
	finding bridges.
0	that they are brief resident and the
disc.	C 010 81/1 2/2 3/3 4/4 515
	0 0 3 4 5
	diston humanitary wills.
(=	disc Cnode] <= low [neigh]  Anticulation Point
(F. c)	Anticulation Point
	9.A Dan
3	Disc (discovery time)
=1 9 9	At what time. I the holde is discovered
3	1 Chr. 71 Chape III 10 ch
	Low: It store the lowest discovery time reachable from current hade.
6.84	reachable from current hade.
L.	Livit Commission Commi
	Company of the state of the sta

Date / / so how to start? we will apply DFS and fill the disc & low value. Then, if we got any visited node, then it will not be a AP. Also we will compare the low of hade with disc. of heigh, & select the min. when we are returning back then we check for — AP (disc Cnode) <= low (huigh))

Tow [ node ] = min [ low (node ] ; 1 low (neigh) We Will all # will don't check for that node that parent is -1. If the next node have more than one child then the nort node is A.P. Sometimes, AP can be counted twice me so, that we can maintain a AP vector.