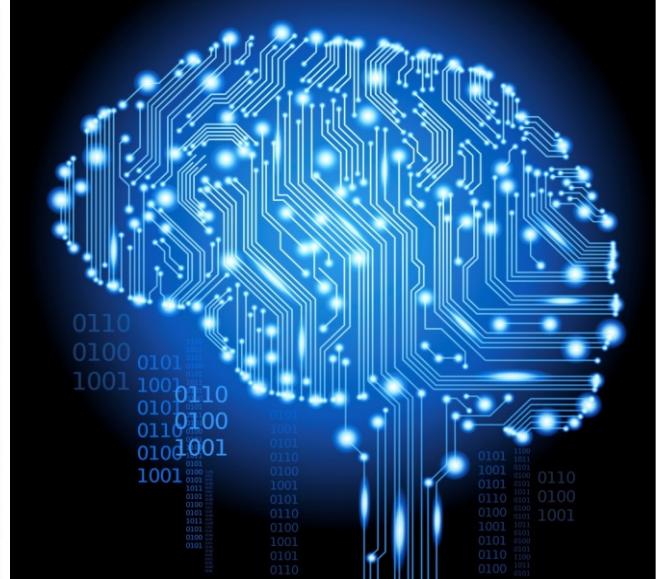
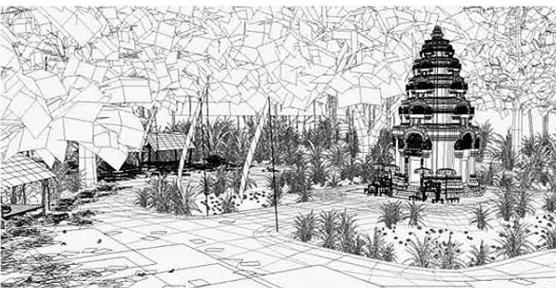


FIT1008/FIT2085 Lecture 8

Function calling

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based on D. Albrecht, J. Garcia



Where are we at:

- **Discussed MIPS architecture and main design decisions**
- **Gone through the instruction set that we will use in this unit**
- **Discussed how variables are stored and accessed**
 - Global variables in the data segment
 - Local variables in the stack segment (the system stack)
 - Memory diagrams
 - Allocation/deallocation with \$sp
 - Access by a negative offset from \$fp
- **Practiced translation of decisions**
 - selection (if) and
 - loops (while and for)
- **Discussed how to create and access arrays of integers**

Learning objectives for this lecture

- **To understand how functions are implemented in MIPS**
- **In particular:**
 - Use of the `jal` and `jr` instructions
 - Use of the `system stack` to satisfy function properties
- **To understand the reasons behind the decisions taken by the function calling convention**
- **To understand what a stack frame is, and its purpose**
- **To be able to implement a function call in MIPS**

Reminder: why using functions?

- **As encapsulation of a sequence of instructions:**

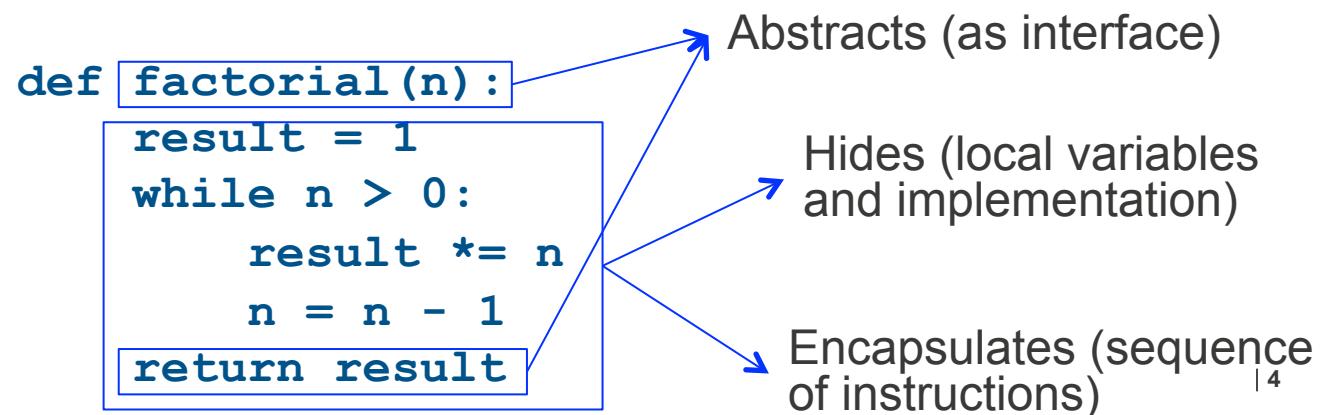
- Are self-contained
 - Can have their own private (local) variables/data
 - Can be called repeatedly (reuse)

- **As abstractions:**

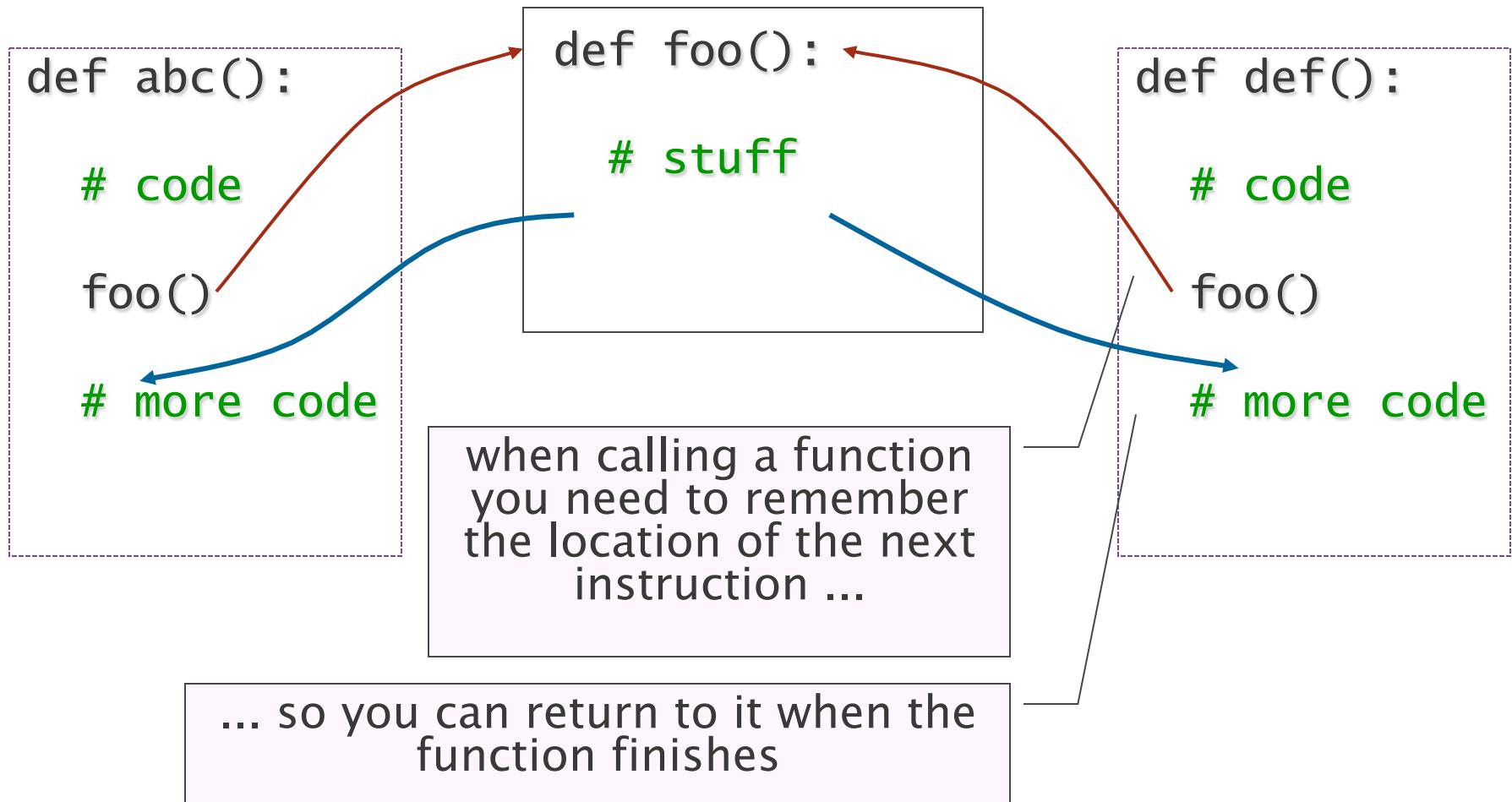
- Can be generalized by taking parameters
 - Can inform through return values

- **As hiders of information**

- Make sure caller cannot access/modify internal data



Function calling: return where?



Function calling in MIPS

- Remember: **jal** instruction for calling a function
 - **jal** is like jump (**j**), but it first **saves the address** of the instruction following the **jal** (ie. at address =PC+4) in register **\$ra**

main code

[0x004001B4]	addi \$t0, \$0, 5
[0x004001B8]	sw \$t0, 0(\$sp)
[0x004001BC]	jal foo
[0x004001C0]	sw \$v0, -8(\$fp)
[0x004001C4]	addi \$sp, \$sp, 4

function code

[0x0041F304]	foo: addi \$v0, \$0, 5
[0x0041F308]	lw \$t0, x

jal first puts address of next instruction (0x004001C0) in register \$ra then jumps to function foo at 0x0041F304

Function calling in MIPS

- Remember: **jr** instruction for returning from a function
 - **jr** jumps to the absolute address held in a specified GPR register (usually **\$ra**)

jr jumps to address
held in \$ra (was set by
jal to 0x004001C0)

```
[0x004001BC] jal foo
[0x004001C0] sw $v0, -8($fp)
[0x004001C4] addi $sp, $sp, 4
```

```
[0x0041F304] foo: addi $v0, $0, 5
[0x0041F308]           lw $t0, x
[0x0041F30C]           jr $ra
```

Function calling in MIPS

- **To write a function**
 - Put **label** at the **start** of the function
 - Write body of the function
 - End function with **jr \$ra**
- **To call a function**
 - Write **jal label**
 - When the function returns, program will continue from the next instruction
- **Is this really all it is needed to call a function?**

Passing data

- **Some functions have parameters**
 - Need a way of passing values for these parameters (i.e., arguments) from caller to function
- **Some functions return values**
 - Need a way of getting the return value safely back to caller
- **Solution: reserve some registers for these tasks**
 - Can extend the “syscall” data passing method:
 - Pass function parameters in \$a0, \$a1, \$a2, \$a3
 - Return values in \$v0, \$v1
- **This convention still has some problems**
 - Will make a better system later
 - But let's see how it works for now

Focus on the
blue parts

Example

```
# Part of program that uses
# function fact

...
# Read int from user
addi $v0, $0, 5
syscall

# Pass parameter in $a0
addi $a0, $v0, 0
# Call!
jal fact
# Put result in $t1
addi $t1, $v0, 0

# Print factorial
addi $v0, $0, 1
addi $a0, $t1, 0
syscall

# Exit
addi $v0, $0, 10
syscall
```

```
# Factorial function

fact:    addi $t0, $0, 1 # result=1
          ...
loop:    # Repeat while parameter
          # in $a0 is > 1
          addi $t1, $0, 1
          slt $t1, $t1, $a0
          beq $t1, $0, end

          # Multiply result by
          # parameter
          mult $t0, $a0
          mflo $t0

          # Decrement parameter
          addi $a0, $a0, -1

          # Loop back to test
          j loop

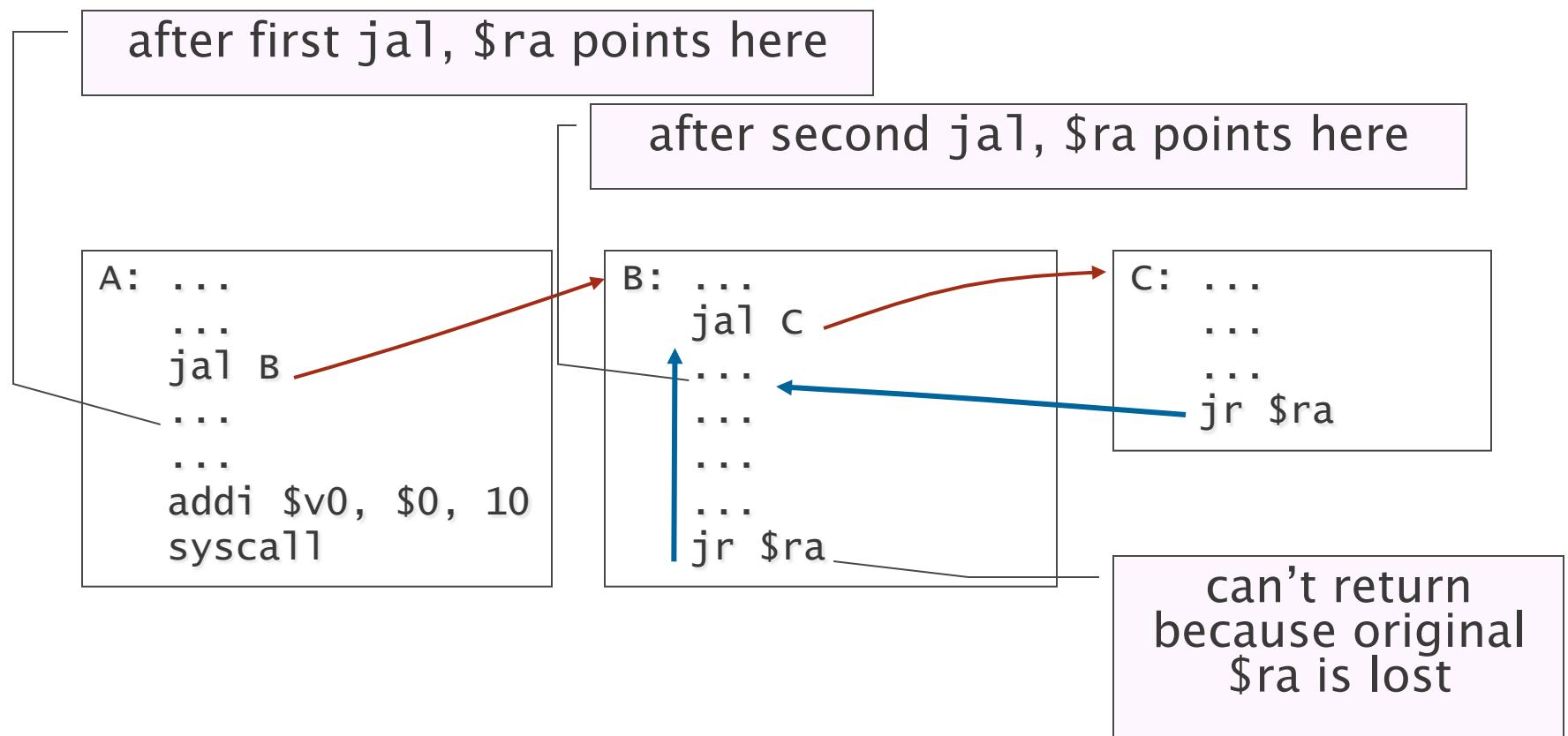
end:     # At end, $t0 contains
          # factorial of user input
          addi $v0, $t0, 0 #$v0=answer
          jr $ra               # Return
```

Limitations

- **This simple function-calling convention works, but has limits**
- **Function must:**
 - Not call other functions (i.e., it is a “leaf function”)
 - Not use more than four parameters (\$a0-\$a3)
 - Only write to “safe” registers \$v0-\$v1, \$a0-\$a3, \$t0-\$t9
 - Not have/use local variables
- **For more sophisticated function calling, one needs more sophisticated convention**

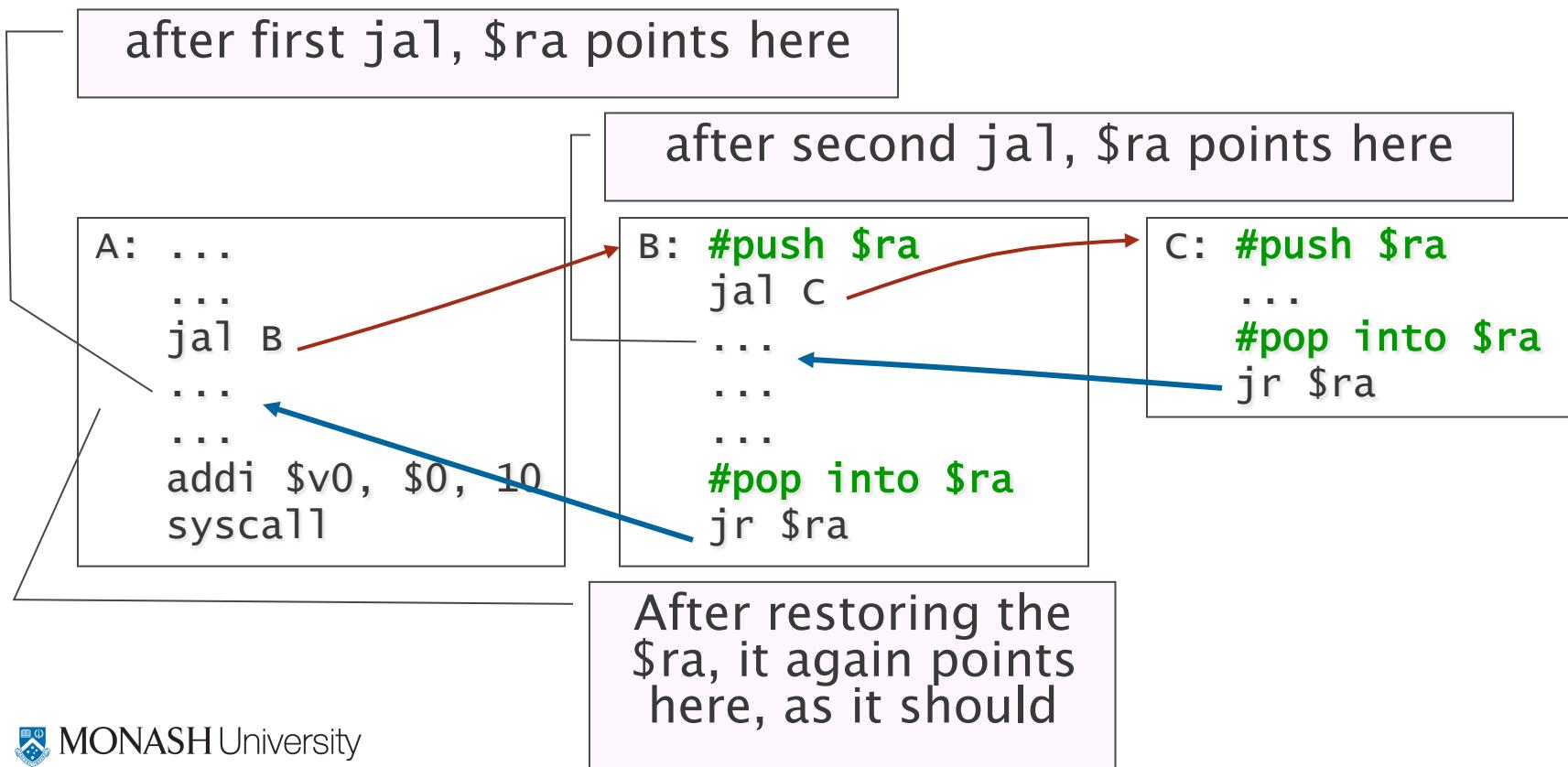
Limitation: calling other functions

- Return address register (\$ra) can hold only one value



Limitation: calling other functions

- Return address register (\$ra) can hold only one value
- Solution: save and restore \$ra register on the stack upon function entry/exit (push on entry, pop on exit)



Limitation: passing arguments

- Not enough registers (4) to pass arguments to some functions

```
addi $a0, $0, 1  
lw $a1, x  
addi $a2, $0, 0  
lw $a3, -4($fp)  
addi $??, $0, 2  
jal five  
...
```

no such register
\$a4

```
five: # takes 5  
      # parameters  
...  
# examine  
# $a0, etc  
...  
jr $ra
```

Limitation: passing arguments

- Not enough registers (4) to pass arguments to some functions
- Solution: pass arguments on stack

```
# push 2
# push global y
# push 0
# push local x
# push 1
jal five
# pop
# pop
# pop
# pop
# pop
```

1
val/addr of x
0
val/addr of y
2

```
five: # takes 5
      # parameters
      ...
      # examine
      # stack
      ...
      jr $ra
```

For simplicity, in FIT2085 we will pass **all** arguments onto the stack

Limitation: saving registers

- Function may use registers which hold important values

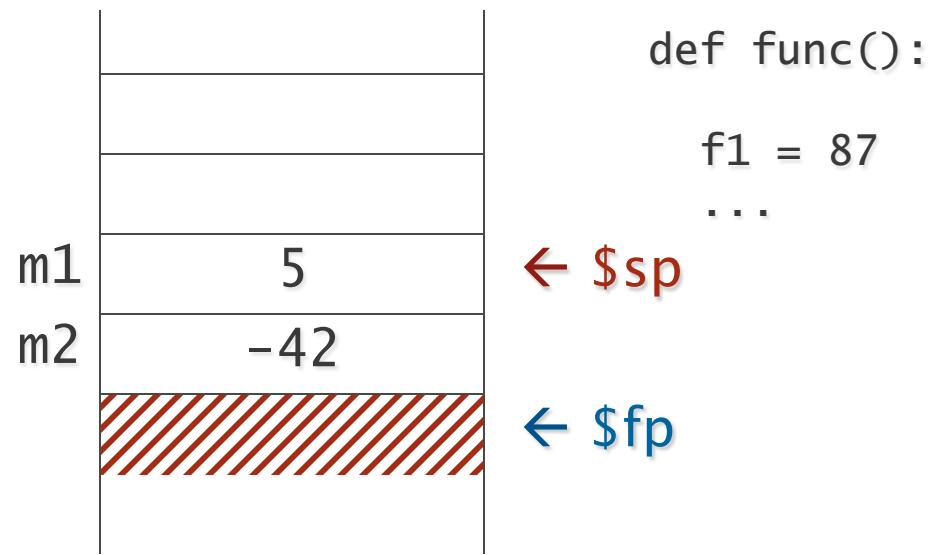
```
...
lw $t0, a
...
...
jal func
...
...
# $t0 has been
# changed!
add $t0, $t0, $v0
...
func: ...
# trashes
# $t0
lw $t0, x
...
jr $ra
```

- Solution: save/restore registers on stack (but not in FIT2085)
- In FIT2085 you are not allowed to do this; assemblers are and do!

Limitation: Local variables (reminder)

- A function needs its own local variable storage

```
def main():  
  
    m1 = 5  
    m2 = -42  
    ...  
    func()
```

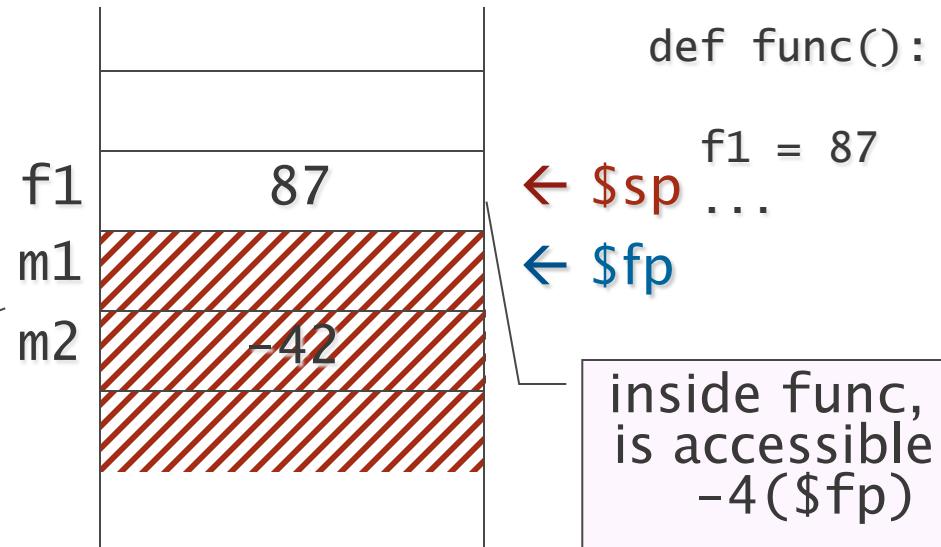


Limitation: Local variables (reminder)

- A function needs its own local variable storage
- Solution: function adjusts \$fp to top of stack, then allocates its own locals

```
def main():  
  
    m1 = 5  
    m2 = -42  
    ...  
    func()
```

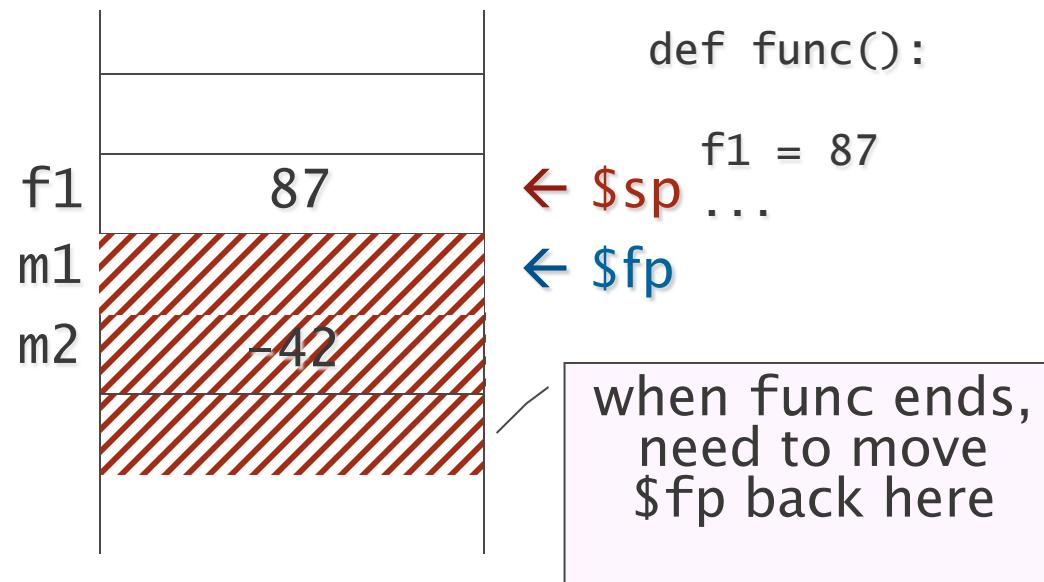
inside func, m1
and m2 are
inaccessible



Limitation: Local variables – restore \$fp

- On function return, stack state (including \$fp) must be restored

```
def main():  
  
    m1 = 5  
    m2 = -42  
    ...  
    func()
```

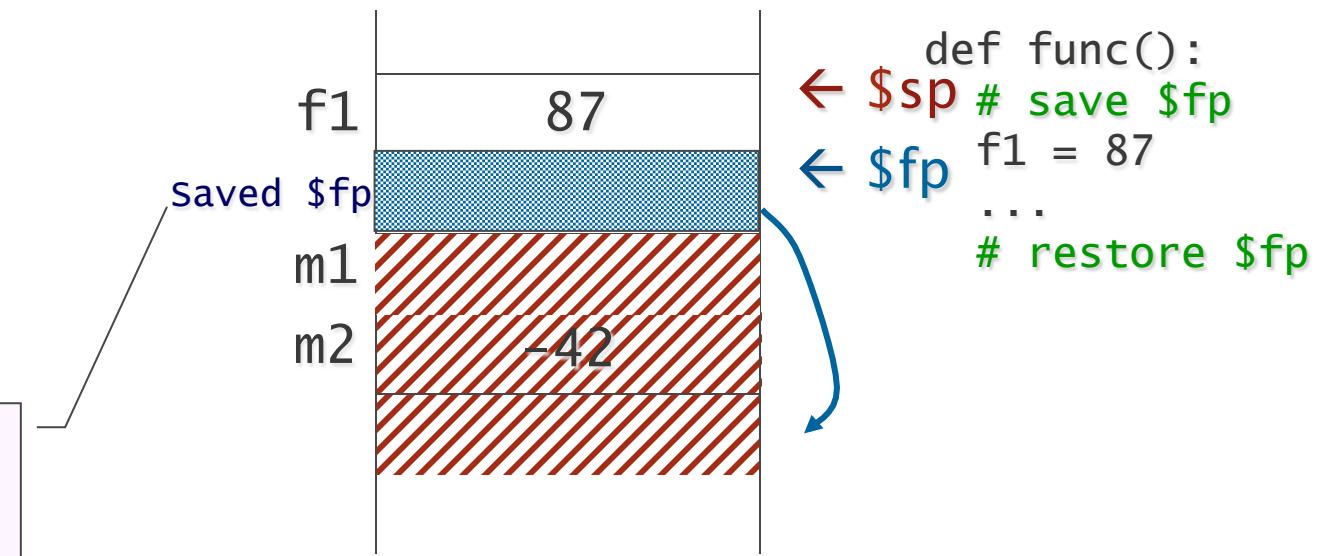


Limitation: Local variables – restore \$fp

- On function return, stack state (including \$fp) must be restored
- Solution: save/restore \$fp on stack

```
def main():  
    m1 = 5  
    m2 = -42  
    ...  
    func()
```

by saving old \$fp,
it can be restored
at function's end



Function calling convention

Task	Whose responsibility?
save / restore temporary registers	caller
pass / clear arguments	caller
save / restore \$ra	callee
save / restore \$fp	callee
allocate / deallocate local variables	callee

Function calling convention

When calling a function, **caller**:

1. Saves temporary **registers** by pushing their values on stack
2. Pushes **arguments** on stack
3. Calls the function with **jal** instruction

On function entry, **callee**:

1. Saves **\$ra** by pushing its value on stack
2. Saves **\$fp** by pushing its value on stack
3. Copies **\$sp** to **\$fp**
4. Allocates local variables



Example: caller

```
# Python program which calls a function.
```

```
def power(b, e):  
    ...  
    return result
```

```
def main():  
    base, exp, result = 0, 0, 0  
    read(base)  
    read(exp)  
  
    result = power(base, exp)  
  
    print(result)
```

This is the memory diagram at this point

base is at -12(\$fp)
exp is at -8(\$fp)
result is at -4(\$fp)

\$sp →	base	3	0x7FFF0394
	exp	4	0x7FFF0398
	result	0	0x7FFF039C
\$fp →			0x7FFF03A0

Assume user has entered
3 for base and 4 for exp



Example: caller

```
# Python program which calls a function.
```

```
def power(b, e):  
    ...  
    return result
```

```
def main():  
    base, exp, result = 0, 0 ,0  
    read(base)  
    read(exp)  
  
    result = power(base, exp)  
  
    print(result)
```

base is at -12(\$fp)
exp is at -8(\$fp)
result is at -4(\$fp)

```
main: .text  
# 3 * 4 = 12 bytes local  
addi $fp, $sp, 0  
addi $sp, $sp, -12
```

Initialize locals

```
sw $0, -12($fp)  
sw $0, -8($fp)  
sw $0, -4($fp)
```

```
addi $v0, $0, 5  
syscall  
sw $v0, -12($fp) # base
```

```
addi $v0, $0, 5  
syscall  
sw $v0, -8($fp) # exp
```

Now we are up to the
function call ...

Example: caller

caller step 1: save temporary registers by pushing their values on stack

not needed in FIT2085

\$sp →	base	3	0x7FFF038C
	exp	4	0x7FFF0390
	result	0	0x7FFF0394
\$fp →			0x7FFF0398
			0x7FFF039C
			0x7FFF03A0

Example: caller

caller step 2: push
function
arguments onto
the stack

two arguments,
called b and e

def power(b, e):

\$sp →	base	3	0x7FFF038C
	exp	4	0x7FFF0390
	result	0	0x7FFF0394
\$fp →			0x7FFF0398
			0x7FFF039C
			0x7FFF03A0

Example: caller

note offsets
of
arguments:
b at 0(\$sp)
e at 4(\$sp)

caller step 2: push
function
arguments onto
the stack

two arguments,
called b and e

def power(b, e):

\$sp →	arg 1 (b)	3	0x7FFF038C
	arg 2 (e)	4	0x7FFF0390
	base	3	0x7FFF0394
	exp	4	0x7FFF0398
	result	0	0x7FFF039C
\$fp →			0x7FFF03A0

Example: caller

caller step 3: call
function with jal
instruction

no effect visible on
the stack

\$sp →	arg 1 (b)	3
	arg 2 (e)	4
	base	3
	exp	4
	result	0
\$fp →		

Example: caller

```
# Python program which calls a function.

def power(b, e):
    ...
    return result

def main():
    base, exp, result = 0, 0 ,0
    read(base)
    read(exp)

    result = power(base, exp)

    print(result)
```

base is at -12(\$fp)
exp is at -8(\$fp)
result is at -4(\$fp)
arg1 at 0(\$Sp)
arg2 at 4(\$sp)

```
# ... continued from
# previous slide
```

```
# Call function.
# push 2 * 4 = 8 bytes
# of arguments
addi $sp, $sp, -8

# arg 1 = base
lw $t0, -12($fp) # base
sw $t0, 0($sp) # arg 1

# arg 2 = exp
lw $t0, -8($fp) # exp
sw $t0, 4($sp) # arg 2
```

```
# Call power function
jal power
```

```
# main function To Be
# Continued next lecture ...
```

Example: callee

```
# A function that gets called.

def power(b, e):

    result = 1

    # Keep going while
    # exponent is positive.

    while e > 0:

        # Multiply by base.
        result *= b

        # Less to multiply.
        e -= 1

    # Return the result to the caller.
    return result
```

Example: callee

callee steps 1 and 2: save \$ra and \$fp by pushing their values on stack

can do both these steps at once

\$sp →	arg 1 (b)	3	0x7FFF0380
	arg 2 (e)	4	0x7FFF0384
	base	3	0x7FFF0388
	exp	4	0x7FFF038C
	result	0	0x7FFF0390
\$fp →			0x7FFF0394
			0x7FFF0398
			0x7FFF039C
			0x7FFF03A0

Example: callee

saved \$fp
contains
address of
another
location on
stack

callee steps 1 and
2: save \$ra and
\$fp by pushing
their values on
stack

can do both these
steps at once

\$sp → saved \$fp
saved \$ra
arg 1 (b)
arg 2 (e)
base
exp
result
\$fp →

	0x7FFF0380
	0x7FFF0384
	0x7FFF0388
	0x7FFF038C
	0x7FFF0390
	0x7FFF0394
	0x7FFF0398
	0x7FFF039C
base	3
exp	4
result	0
\$fp →	0x7FFF03A0

Example: callee

callee step 3: copy
\$sp into \$fp

now main's local
variables are
inaccessible

\$fp = \$sp → → saved \$fp

		0x7FFF0380
		0x7FFF0384
	saved \$ra	0x7FFF0388
	arg 1 (b)	0x7FFF038C
	arg 2 (e)	0x7FFF0390
	base	0x7FFF0394
	exp	0x7FFF0398
result	0	0x7FFF039C
		0x7FFF03A0

Example: callee

callee step 4:
allocate local
variables

in this function,
one local variable
(result)

\$fp = \$sp → → saved \$fp

saved \$ra

arg 1 (b)

arg 2 (e)

base

exp

result

		0x7FFF0380
		0x7FFF0384
		0x7FFF0388
		0x7FFF038C
		0x7FFF0390
		0x7FFF0394
		0x7FFF0398
		0x7FFF039C
		0x7FFF03A0

Example: callee

callee step 4:
allocate local
variables

in this function,
one local variable
(result)

\$sp →	result	1	0x7FFF0380
\$fp →	saved \$fp	0x7FFF03A0	0x7FFF0384
	saved \$ra	0x0040005C	0x7FFF0388
	arg 1 (b)	3	0x7FFF038C
	arg 2 (e)	4	0x7FFF0390
	base	3	0x7FFF0394
	exp	4	0x7FFF0398
	result	0	0x7FFF039C
			0x7FFF03A0

Frame Pointers

\$fp always points to an old, saved, copy of \$fp	\$sp → result	1	0x7FFF0380
	\$fp → saved \$fp	0x7FFF03A0	0x7FFF0384
	saved \$ra	0x0040005C	0x7FFF0388
	arg 1 (b)	3	0x7FFF038C
	arg 2 (e)	4	0x7FFF0390
	base	3	0x7FFF0394
	exp	4	0x7FFF0398
	result	0	0x7FFF039C
			0x7FFF03A0

Stack Frames

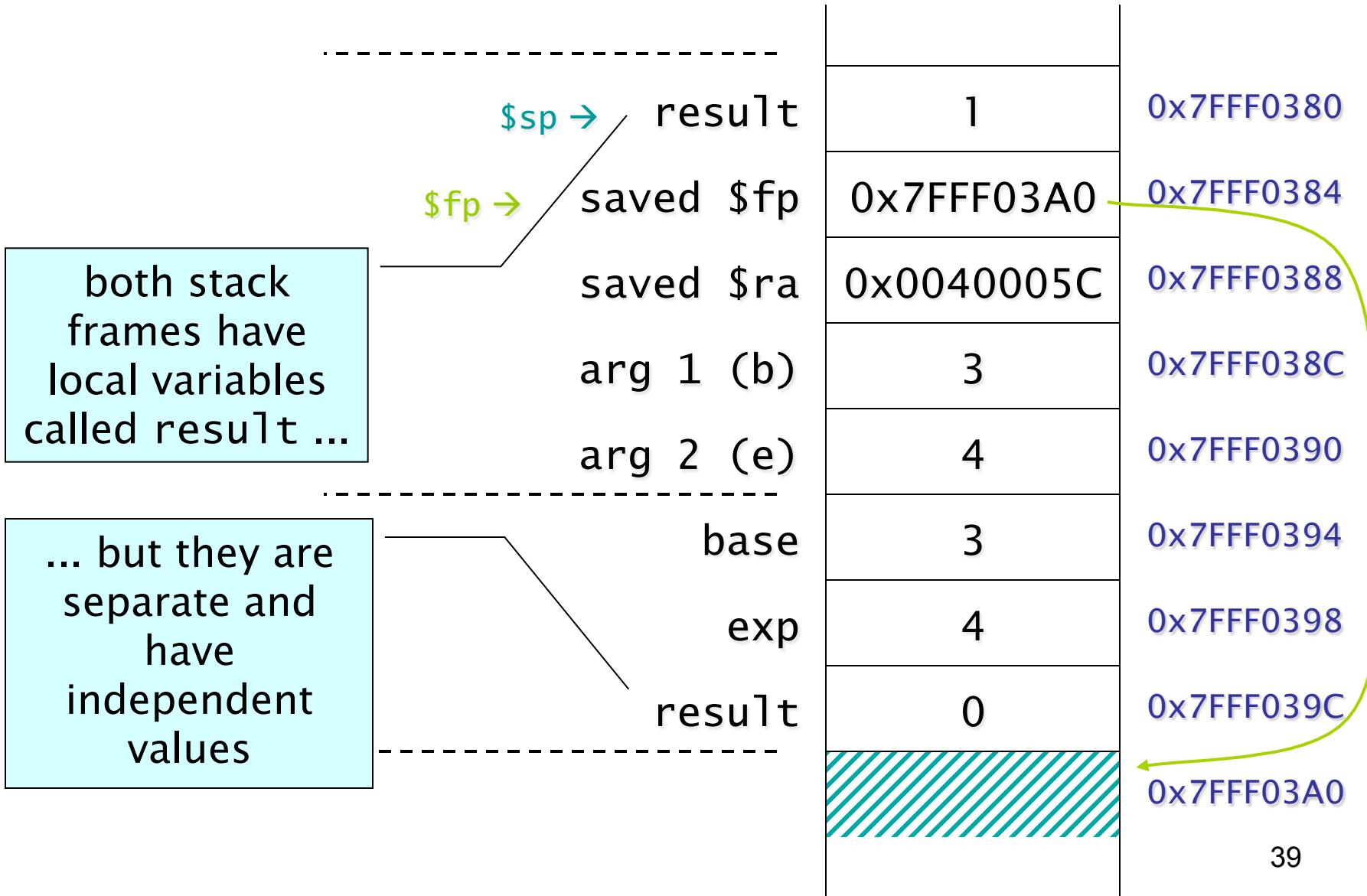
the data on
the stack
that is
associated
with a
function is
called a
stack frame

\$sp → result	1	0x7FFF0380
\$fp → saved \$fp	0x7FFF03A0	0x7FFF0384
	saved \$ra	0x7FFF0388
	arg 1 (b)	0x7FFF038C
	arg 2 (e)	0x7FFF0390
base	3	0x7FFF0394
exp	4	0x7FFF0398
result	0	0x7FFF039C

Stack Frames

power's stack frame	\$sp →	result	1	0x7FFF0380
	\$fp →	saved \$fp	0x7FFF03A0	0x7FFF0384
		saved \$ra	0x0040005C	0x7FFF0388
		arg 1 (b)	3	0x7FFF038C
		arg 2 (e)	4	0x7FFF0390
		base	3	0x7FFF0394
		exp	4	0x7FFF0398
		result	0	0x7FFF039C
				0x7FFF03A0
				0x7FFF03A0

Stack Frames



Local variables

result is at
-4(\$fp)

power's local
variables are
accessed as
main's were,
with negative
offsets from
\$fp

\$sp → result

\$fp → saved \$fp

saved \$ra

arg 1 (b)

arg 2 (e)

base

exp

result

1	0x7FFF0380
0x7FFF03A0	0x7FFF0384
0x0040005C	0x7FFF0388
3	0x7FFF038C
4	0x7FFF0390
3	0x7FFF0394
4	0x7FFF0398
0	0x7FFF039C
	0x7FFF03A0

Function parameters

b and **e** are respectively accessible at +8(\$fp) and +12(\$fp)

power's parameters (arguments) are accessible with positive offset from \$fp

\$sp → result	1	0x7FFF0380
\$fp → saved \$fp	0x7FFF03A0	0x7FFF0384
saved \$ra	0x0040005C	0x7FFF0388
arg 1 (b)	3	0x7FFF038C
arg 2 (e)	4	0x7FFF0390
base	3	0x7FFF0394
exp	4	0x7FFF0398
result	0	0x7FFF039C
		0x7FFF03A0

Example: callee

```
# A function that gets called.  
  
def power(b, e):  
    result = 1  
  
    # Keep going while  
    # exponent is positive.  
  
    while e > 0:  
  
        # Multiply by base.  
        result *= b  
  
        # Less to multiply.  
        e -= 1  
  
    # Return the result to the caller.  
    return result
```

result is at -4(\$fp)
b is at 8(\$fp)
e is at 12(\$fp)

power: # Save \$ra and \$fp
addi \$sp, \$sp, -8
sw \$ra, 4(\$sp)
sw \$fp, 0(\$sp)

Copy \$sp to \$fp
addi \$fp, \$sp, 0

Alloc local variables
1 * 4 = 4 bytes.
addi \$sp, \$sp, -4

Initialize locals.
addi \$t0, \$0, 1
sw \$t0, -4(\$fp) # result

Now we are inside
the function body.

Example: callee

```
# A function that gets called.  
  
def power(b, e):  
  
    result = 1  
  
    # Keep going while  
    # exponent is positive.  
  
    while e > 0:  
  
        # Multiply by base.  
        result *= b  
  
        # Less to multiply.  
        e -= 1  
  
    # Return the result to the caller.  
    return result
```

result is at -4(\$fp)
b is at 8(\$fp)
e is at 12(\$fp)

```
loop:      # Stop if !(e > 0)  
    lw $t0, 12($fp) # e  
    slt $t0, $0, $t0  
    beq $t0, 0, end  
  
    # result = result * b  
    lw $t0, -4($fp) # result  
    lw $t1, 8($fp) # b  
    mult $t0, $t1  
    mflo $t0  
    sw $t0, -4($fp) # result  
  
    # e = e - 1  
    lw $t0, 12($fp) # e  
    addi $t0, $t0, -1  
    sw $t0, 12($fp) # e  
  
    # Repeat loop.  
    j loop  
  
end:      # Now ready to return.  
# Continued in next lecture ...
```

Summary

- **Function calling**
 - jal and jr instructions
- **Function calling convention**
 - For making a function call
- **Structure of stack**
 - stack frames

