

## // Union Find Disjoint Set

```
vector<int>u_list, u_set;
```

// u\_list[x] contains the size of a set x    u\_set[x] contains the root of x

```
int unionRoot(int n) {
    if(u_set[n] == n)
        return n;
    else
        return u_set[n] = unionRoot(u_set[n]);
}
```

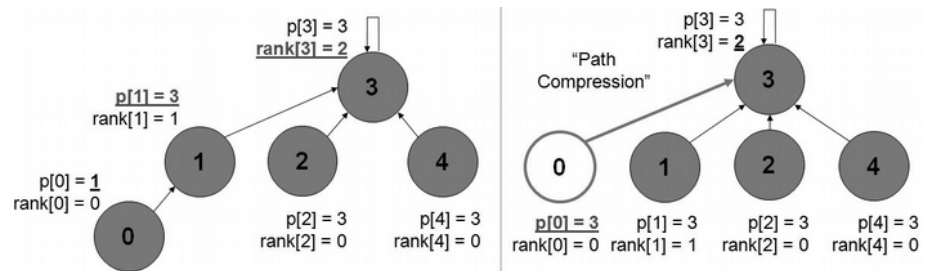
// Finds the root of a point/element  
 // If u\_set[n] == n, then n is the root of set  
 // Else keep searching for root

```
int makeUnion(int a, int b) {
    int x = unionRoot(a);
    int y = unionRoot(b);
```

// Takes two points in the same set, and return the root  
 // Find the root of x and y

```
    if(x == y)
        return x;
    else if(u_list[x] > u_list[y]) {
        u_set[y] = x;
        u_list[x] += u_list[y];
        return x;
    }
    else {
        u_set[x] = y;
        u_list[y] += u_list[x];
        return y;
    } }
```

// If root of both points are same, then nothing to do  
 // If the size of x set is larger than set y (Path Compression)  
 // Make set y the subset of x  
 // Increase the size of set x



```
void unionInit(int len) {
    u_list.resize(len+5, 1);
    u_set.resize(len+5);
    for(int i = 0; i <= len; i++)
        u_set[i] = i;
}
```

// Union Disjoint Set Initialization  
 // Space allocation  
 // At first, root of all points is the point itself

```
bool isSameSet(int a, int b) {
    if(unionRoot(a) == unionRoot(b))
        return 1;
    return 0;
}
```

// Check if two points are in same set  
 // If the root of both point are same, then they are same

## // Segment Tree

```
int arr[N], tree[4*N];
```

// Always take the tree size 4x  
 // arr[] contains values starting from index 1  
 // Builds the segment tree call : segment\_build(1, 1, len\_of\_arr)  
 // Initialization of tree position  
 // Mid point reached

```
void segment_build(int pos, int L, int R) {
    tree[pos] = 0;
    if(L==R) {
        tree[pos] = arr[L];
    }
```

```

    return; }
    int mid = (L+R)/2;
    segment_build(pos*2, L, mid);
    segment_build(pos*2+1, mid+1, R);
    tree[pos] = tree[pos*2] + tree[pos*2+1]; // Depends on usage (This is the main point to tweak)
}

void segment_update(int pos, int L, int R, int i, int val) { // Val contains the value to update SINGLE UPDATE
    if(L==R) { // Call: segment_update(1, 1, len_of_arr, pos_in_arr new_value)
        tree[pos] = val; // If L==R then this is the midpoint that contains the single value
        return; }
    int mid = (L+R)/2;
    if(i <= mid) segment_update(pos*2, L, mid, i, val); // Go for update_position
    else segment_update(pos*2+1, mid+1, R, i, val);
    tree[pos] = tree[pos*2] + tree[pos*2+1]; // Depends on usage (here summation)
}

// Query in range l-r
int segment_query(int pos, int L, int R, int l, int r) { // Finds value in l-r segment of arr[]
    if(R < l || r < L) // Out of range l-r
        return 0;
    if(l <= L && R <= r) // In range l-r
        return tree[pos];
    int mid = (L+R)/2;
    int x = segment_query(pos*2, L, mid, l, r);
    int y = segment_query(pos*2+1, mid+1, R, l, r);
    return x+y; // Return the total value
}

// Toggle bit in range of [l, r] call : update(1, 1, length_of_input, l, r) EX: Update lights in range l-r
void update(int pos, int L, int R, int l, int r) { // Lazy without propagation
    if(l <= L && R <= r) { // If segment is in range l-r
        arr[pos] ^= 1; // Tweak according to problem
        return; }
    if(r < L || R < l) // Out of range l-r
        return;
    int mid = (L+R)/2;
    update(pos*2, L, mid, l, r);
    update((pos*2) + 1, mid+1, R, l, r);
}

// Single position query
bool query(int at, int L, int R, int pos) { // Query in arr position 'pos'
    if(pos < L || R < pos) return 0; // Range is out of pos
    if(L <= pos && pos <= R) return arr[at]; // pos is in range
    int mid = (L+R)/2;
    if(pos <= mid)
        return query(at*2, L, mid, pos) ^ arr[at];
    else
        return query(at*2+1, mid+1, R, pos) ^ arr[at];
}

```

