

// Sieve, Prime, Factorization

```
bitset<10000000>isPrime;
vector<long long>primes;
```

```
void sieve(unsigned long long N) { // Only generates a number is prime or not
    isPrime.set();
    isPrime[0] = isPrime[1] = 0; // 0 and 1 are not prime
    unsigned long long lim = sqrt(N) + 5;
    for(unsigned long long i = 2; i <= lim; i++) {
        if(isPrime[i]) {
            for(unsigned long long j = i*i; j <= N; j+= i)
                isPrime[j] = 0;
        }
    }
}
```

```
void sieveGen(unsigned long long N) { // Generates a number is prime or not, and also makes an array of prime numbers
    isPrime.set();
    isPrime[0] = isPrime[1] = 0; // 0 and 1 are not prime
    for(unsigned long long i = 2; i <= N; i++) { //Note, N isn't square rooted!
        if(isPrime[i]) {
            for(unsigned long long j = i*i; j <= N; j+= i)
                isPrime[j] = 0;
            primes.push_back(i);
        }
    }
}
```

```
vector<int> primeFactor(long long n) { // Returns vector of co-efficient of prime factor
    if(isPrime[n]) { // v[x] contains the co-efficient of x
        vector<int>factor(n+1, 0);
        factor[n] = factor[1] = 1;
        return factor;
    }
    vector<int>factor(sqrt(n)+1, 0); //the size of vector must be at most sqrt(n)+1
    for(long long i = 0; i < (int)primes.size() && primes[i] <= n; i++) {
        while(n%primes[i] == 0) { //divide 1 - n with primes 1 - n
            factor[primes[i]]++; //counts how many prime in the number
            n/=primes[i]; //cuts out the prime
        }
    }
    return factor;
}
```

```
// Returns the divisors without sieve sqrt(n)
vector<unsigned long long>divisor;
void divisors(unsigned long long n) {
    unsigned long long lim = sqrt(n);
    for(unsigned long long i = 2; i <= lim; i++) {
        if(n % i == 0) {
            unsigned long long tmp = n/i;
            divisor.push_back(tmp);
            if(i != tmp) divisor.push_back(i);
        }
    }
}
```

// Prime factorization of factorials (n!)

```

vector<pair<long long, long long> > factorialFactorization(long long n) {
    vector<pair<long long, long long> > V;
    for(long long i = 0; i < (int)primes.size() && primes[i] <= n; i++) {
        long long tmp = n, power = 0;
        while(tmp/primes[i]) {
            power += tmp/primes[i];
            tmp /= primes[i];
        }
        if(power != 0)
            V.push_back(make_pair(primes[i], power));
    }
    return V;
}

long long numPF(long long n) {                //returns number of prime factors
    long long num = 0;
    for(long long i = 0; primes[i] * primes[i] <= n; i++) {
        while(n % primes[i] == 0) {
            n /= primes[i];
            num++;
        }
    }
    if(n > 1)    num++;                        //there might left a prime number which is bigger than primes[i]
    return num;
}

long long numDIFPF(long long n) {            // returns number of different prime factors
    long long diff_num = 0;
    for(long long i = 0; primes[i] * primes[i] <= n; i++) {
        bool ok = 0;
        while(n % primes[i] == 0) {
            n /= primes[i];
            ok = 1;
        }
        if(ok)    diff_num++;
    }
    if(n > 1)    diff_num++;
    return diff_num;
}

unsigned long long sumPF(long long n) {        //returns sum of prime factors
    unsigned long long sum = 0;
    for(long long i = 0; primes[i] * primes[i] <= n; i++)
        while(n % primes[i] == 0) {
            n /= primes[i];
            sum+=primes[i];
        }
    if(n > 1)    sum+= n;
    return sum;
}

```

```
vector<int>divisors[1000];
void Divisors(int n) {
    for(int i = 1; i <= n; ++i)
        for(int j = i; j <= n; j+= i)
            divisors[j].push_back(i);
}
```

// Finding Divisors without calculating prime numbers
 // We can avoid 1 and 2 if we want
 // Also, we can start from j = i+i
 // As it is known every number is divisible by itself
 // divisors[i] contains a list of numbers

// Binomial Coefficient C(n, k)
// Complexity : O(k)

```
long long binomialCoeff(long long n, long long k) {
    long long res = 1;
    if ( k > n - k )
        k = n - k;
    for (long long i = 0; i < k; ++i) {
        res *= (n - i);
        res /= (i + 1);
    }
    return res;
}
```

// Since $C(n, k) = C(n, n-k)$
 // Calculate value of $[n * (n-1) * \dots * (n-k+1)] / [k * (k-1) * \dots * 1]$
// Note: divide first then multiply to avoid overflow, decimal can be taken
// After every calculation round up the value

// Catalan Number
// Use this with Binomial Coefficient

```
long long catalan(int n) {
    long long c = binomialCoeff(2*n, n);
    return c/(n+1); }
```

//Cat(n) = $C(2*n, n)/(n+1)$;

// Euler's Totient

```
/* Euler's Totient function  $\Phi(n)$  for an input n is count of numbers in {1, 2, 3, ..., n}
* that are relatively prime to n, i.e., the numbers whose GCD (Greatest Common Divisor) with n is 1.
* Phi(4) : GCD(1, 4) = 1, GCD(3, 4)
* so, Phi(4) = 2
*/
```

```
int Phi(int n) {
    int result = n;
    for (int p=2; p*p<=n; ++p) {
        if (n % p == 0) {
            while (n % p == 0)
                n /= p;
            result -= result / p;
        }
    }
    if (n > 1)
        result -= result / n;
    return result;
}
```

// Computes phi of n
 // Initialize result as n
 // Consider all prime factors of n and subtract their multiples from result
 // p is a prime factor of n
 // Eliminate all p factors from n
 // If n is still greater than 1, then it is also a prime

```
long long phi[MAX];
void computeTotient(int n) {
    for (int i=1; i<=n; i++)
```

// Computes phi or Euler Phi 1 to n
 // Initialize

```

    phi[i] = i;
    for (int p=2; p<=n; p++) {           // Computation
        if (phi[p] == p) {               // If phi is not computed
            phi[p] = p-1;                 // p is prime and phi(prime) = prime-1;
            for (int i = 2*p; i<=n; i += p) { // Sieve like implementation
                phi[i] = (phi[i]/p) * (p-1); // Add contribution of p to its multiple i by multiplying with (1 - 1/p)
            } } }

```

// Prime Probability

// Algorithm : Miller-Rabin primality test

// This function can be used as power or mod power

```

int powMod(int x, unsigned int y, int p) {           // If pow(x, y) needed, change lines according to the comments
    int res = 1;
    x = x % p;                                     // Remove this line
    while (y > 0)
        if (y & 1)
            res = (res*x) % p;                       // res = res * x;
        y = y>>1;
        x = (x*x) % p;                               // x = x * x;
    }
    return res;
}

```

// This function is called for all k trials. It returns false if n is composite and returns true if n is probably prime.

// d is an odd number such that $d \cdot 2^r \leq n-1$ for some $r \geq 1$

```

bool miillerTest(int d, int n) {
    int a = 2 + rand() % (n - 4);           // Pick a random number in [2..n-2] . Corner cases make sure that n > 4
    int x = powMod(a, d, n);                 // Compute a^d % n
    if (x == 1 || x == n-1)
        return true;
    while (d != n-1) {                       // Keep squaring x while one of the following doesn't happen
        x = (x * x) % n;                     // (i) d does not reach n-1
        d *= 2;                             // (ii) (x^2) % n is not 1
        if (x == 1) return false;           // (iii) (x^2) % n is not n-1
        if (x == n-1) return true; }
    return false;                           // Return composite
}

```

// Note : Use k = 10 to avoid WA

```

bool isPrime(int n, int k) {                // It returns false if n is composite and returns true if n is probably prime. k is an input
    if (n <= 1 || n == 4) return false;    // parameter that determines accuracy level. Higher value of k indicates more accuracy.
    if (n <= 3) return true;               // Corner cases
    int d = n - 1;                          // Find r such that n = 2^d * r + 1 for some r >= 1
    while (d % 2 == 0)
        d /= 2;
    for (int i = 0; i < k; i++)              // Iterate given nber of 'k' times
        if (miillerTest(d, n) == false)
            return false;
    return true;
}

```

```
main() { .....
    if(isPrime(3, 10))
        cout << "This number is prime" << endl;
    .....}
```

// Pascale's Triangle

```
long long p[55][54];
void buildPascale() {                //Building Pascale of 50 rows where p[pascale_line][no_of_element] has every element values
    p[0][0] = 1;                      // Base Case
    p[1][0] = p[1][1] = 1;
    for(int i = 2; i <= 50; i++)
        for(int j = 0; j <= i; j++) {
            if(j == 0 || j == i)
                p[i][j] = 1;
            else
                p[i][j] = p[i-1][j-1] + p[i-1][j];
        }
    /* Uncomment this if you want to see the full triangle
    for(int i = 0; i <= 20; i++) {
        for(int j = 0; j <= i; j++)
            printf("%lld ", p[i][j]);
        printf("\n");
    } */
    return;
}
```

// Horner Polynomial Equation Solver $O(n \log n)$

// Naive Approach Complexity: $O(n^2)$,

// Evaluate value of $2x^3 - 6x^2 + 2x - 1 = 0$ for $x = 3$
 // Input: co_efficient[] = {2, -6, 2, -1}, $x = 3$
 // Output: 5 Algorithm Calculation : $((((2) x - 6) x + 2) x - 1)$

```
int co_efficient[1000];                // Contains the co-efficients
long long horner(long long x, long long n) { // Critical case : Check if number of co-efficient is equal to
    long long ans = co_efficient[0];        // (max power of x) + 1
    for(int i = 1; i < n; i++) {
        ans = ans*x + co_efficient[i];
    }
    return ans;
}
```

// Extended Euclid

```
int x, y, d;
void extendedEuclid(int a, int b) {
    if(b == 0) { x = 1; y = 0; d = a; return; }
    extendedEuclid(b, a%b);
    int x1 = y;
    int y1 = x - (a/b) * y;
    x = x1;
    y = y1;
}
```

// Linear Diophantine for solving equation

```

float ansX, ansY;
void linear_diophantine(int a, int b, int c) {
    extendedEuclid(a, b);
    int g = c / __gcd(a, b);
    float x0 = x*g, y0 = y*g;
    float low_n = - x0 / (b/d), hi_n = y0 / (a/d);
    low_n = ceil(low_n), hi_n = floor(hi_n);
    ansX = x0 + b * low_n;
    ansY = y0 - a * low_n;
}

```

// Contains answer of x and y respectively
 // Solving linear Diophantine equations in two variables
 // $ax + by = c$
 // $x = x_0 + b * n$ where n is an integer
 // $y = y_0 - a * n$ where n is an integer
 // If $low_n \neq hi_n$, then there exists
 // More than one solution for $low_n \leq n \leq hi_n$
 // Only getting the first solution

// Some important Functions

```

int mod(int a, int b) {
    return ((a%b) + b) % b;
}

```

// Actual mod is $(x \% m)$ biggest multiple of m which is less than x
 // $-15 \bmod 60 = 45$ (works like clock)
 // $(a + b) \% m = ((a \% m) + (b \% m)) \% m$ $(a * b) \% m = ((a \% m) * (b \% m)) \% m$

```

int gcd(int a, int b) {
    while (b) {
        int tmp = a%b;
        a = b; b = tmp; }
    return a;
}

```

```

int lcm(int a, int b) {
    return a / gcd(a, b)*b;
}

```

```

int mod_inverse(int a, int n) {
    int x, y;
    int g = extendedEuclid(a, n);
    if (g > 1) return -1;
    return mod(x, n);
}

```

// Computes b such that $ab = 1 \pmod{n}$, returns -1 on failure
 // Use extendedEuclid function

// Date & Time

```

ll age(ll y1, ll m1, ll d1, ll y2, ll m2, ll d2) {
    ll t1 = y1*10000+m1*100+d1;
    ll t2 = y2*10000+m2*100+d2;
    ll age = t1 - t2;
    if(age < 0) return -1;
    return age/10000;
}

```

// Calculates age (only year)
 // Today, Birthday Leap Years are also considered

```

bool isLeapYear(ll year) {
    if(year % 4 == 0 && year % 100 != 0) return 1;
    else if(year % 400 == 0) return 1;
    else return 0; }

```

// Returns True if leap year