```
// DSU on TREE
 1.
 2.
 3.
    int sz[maxn];
 4.
    void getsz(int v, int p){
 5.
         sz[v] = 1; // every vertex has itself in its subtree
         for(auto u : g[v])
 6.
 7.
             if(u != p){
 8.
                 getsz(u, v);
                 sz[v] += sz[u]; // add size of child u to its parent(v)
 9.
10.
    }}
11.
    // Map Style: n(logn)^2
12.
13.
14.
    map<int, int> *cnt[maxn];
15.
    void dfs(int v, int p){
        int mx = -1, bigChild = -1;
16.
        for(auto u : g[v])
17.
            if(u != p){
18.
19.
                dfs(u, v);
20.
                if(sz[u] > mx)
21.
                    mx = sz[u], bigChild = u;
            }
22.
        if(bigChild != -1)
23.
             cnt[v] = cnt[bigChild];
24.
25.
        else
26.
             cnt[v] = new map<int, int> ();
27.
         (*cnt[v])[ col[v] ] ++;
28.
29.
        for(auto u : g[v])
30.
            if(u != p && u != bigChild){
31.
                for(auto x : *cnt[u])
                    (*cnt[v])[x.first] += x.second;
32.
33.
        //now (*cnt[v])[c] is the number of vertices in subtree of vertex v that has color c.
34.
    You can answer the queries easily.
35.
36.
37.
    // Vector Style: nlogn
38.
39.
    vector<int> *vec[maxn];
40.
    int cnt[maxn];
41.
    void dfs(int v, int p, bool keep){
42.
        int mx = -1, bigChild = -1;
43.
        for(auto u : g[v])
44.
            if(u != p \&\& sz[u] > mx)
45.
                mx = sz[u], bigChild = u;
        for(auto u : g[v])
46.
47.
            if(u != p && u != bigChild)
                dfs(u, v, 0);
48.
49.
        if(bigChild != -1)
50.
             dfs(bigChild, v, 1), vec[v] = vec[bigChild];
51.
        else
52.
             vec[v] = new vector<int> ();
53.
        vec[v]->push_back(v);
        cnt[ col[v] ]++;
54.
         for(auto u : g[v])
55.
56.
            if(u != p && u != bigChild)
57.
                for(auto x : *vec[u]){
```

```
58.
                     cnt[ col[x] ]++;
 59.
                     vec[v] -> push_back(x);
 60.
         //now (*cnt[v])[c] is the number of vertices in subtree of vertex v that has color c.
 61.
     You can answer the queries easily.
         // note that in this step *vec[v] contains all of the subtree of vertex v.
 62.
         if(keep == 0)
 63.
             for(auto u : *vec[v])
 64.
                  cnt[ col[u] ]--;
 65.
 66.
     }
 67.
     // Heavy-Light-Decomposition Style: nlogn
 68.
 69.
 70.
     int cnt[maxn];
 71.
     bool big[maxn];
 72.
     void add(int v, int p, int x){
         cnt[col[v]] += x;
 73.
 74.
         for(auto u: g[v])
 75.
             if(u != p && !big[u])
 76.
                  add(u, v, x)
 77.
     void dfs(int v, int p, bool keep){
 78.
 79.
         int mx = -1, bigChild = -1;
 80.
         for(auto u : g[v])
 81.
            if(u != p \&\& sz[u] > mx)
 82.
                mx = sz[u], bigChild = u;
 83.
         for(auto u : g[v])
             if(u != p && u != bigChild)
 84.
 85.
                  dfs(u, v, 0); // run a dfs on small childs and clear them from cnt
 86.
         if(bigChild != -1)
 87.
             dfs(bigChild, v, 1), big[bigChild] = 1; // bigChild marked as big and not cleared
     from cnt
 88.
         add(v, p, 1);
 89.
         //now cnt[c] is the number of vertices in subtree of vertex v that has color c. You can
     answer the queries easily.
 90.
         if(bigChild != -1)
             big[bigChild] = 0;
 91.
 92.
         if(keep == 0)
             add(v, p, -1);
 93.
 94.
     }
 95.
 96.
     // KMP EXTRA PART
 97.
 98.
     // p is the pattern where table[] is the pre-made prefix table of pattern
 99.
     // for any index idx the nxt[idx][j] returns the new index idx where the index
100.
     // should point next, this optimizes the kmp in linear time
101.
102.
     void getState(string &p, int table[], int nxt[][27]) {
         for(int i = 0; i < p.size(); ++i) {</pre>
103
             for(int j = 0; j < 26; ++j) {
104.
105.
                  if(p[i]-'a' == j)
                      nxt[i][j] = i+1;
106.
107.
                  else
108.
                      nxt[i][j] = i == 0 ? 0 : nxt[table[i-1]][j];
109.
     }}}
110.
     // check function using nxt[idx][j]
111.
112. // idx is the index from which the string should start matching with the pattern
```

```
113. // by default idx = 0, also it refers the last index of the pattern to which
114.
     // the string matched
115.
116.
     int match(string &s, int table[], int nxt[][27], int &idx) {
117.
         int ans = 0;
         for(char c : s) {
118.
             idx = nxt[idx][c-'a'];
119.
             if(idx == p.size())
120.
121.
                  ++ans, idx = table[idx-1];
122.
123.
         return ans;
124.
     }
125.
    // LCA
126.
127.
     // Least Common Ancestor with sparse table
128.
     void dfs(int u, int p) {
129.
         in[u] = ++cnt;
130.
131.
         revIn[cnt] = u, par[u][0] = p, lvl[u] = lvl[p]+1;
132.
         for(int i = 1; i \le 20; ++i)
133.
             par[u][i] = par[par[u][i-1]][i-1];
134.
135.
136.
         for(auto v : G[u])
             if(v != p)
137.
138.
                  dfs(v, u);
139.
         out[u] = cnt;
140.
     }
141.
142.
     int LCA(int u, int v) {
143.
         if(lvl[u] < lvl[v]) swap(u, v);
         for(int p = 20; p >= 0; --p)
144.
             if(lvl[u] - (1 << p) >= lvl[v])
145.
                  u = par[u][p];
146.
147.
         if(u == v) return u;
148.
         for(int p = 20; p >= 0; --p)
             if(par[u][p] != par[v][p])
149.
150.
                  u = par[u][p], v = par[v][p];
151.
         return par[u][0];
152.
     }
153.
154.
     // LCA if the root changes, [first dfs is done with root 1 or any other fixed node]
155.
     int LCA(int u, int v, int root) {
156.
         if(isChild(u, root) and isChild(v, root))
157.
             return LCA(u, v);
158.
         if(isChild(u, root) != isChild(v, root))
             return root;
159.
160.
         int x = LCA(u, v), y = LCA(u, root), z = LCA(v, root);
         int a = lvl[root] - lvl[x], b = lvl[root] - lvl[y], c = lvl[root] - lvl[z];
161.
162.
         if(a \le b \text{ and } a \le c) \text{ return } x;
163.
         if(b <= a and b <= c) return y;</pre>
164.
         return z;
165.
166.
167.
     // ----- LCA WITH Sparse Table Vector -----
     // DFS and LCA INIT is same
168.
169.
     void MERGE(vector<int>&u, vector<int>&v) {
                                                           // Do what is to be done to merge
170.
         for(auto it : v) u.push_back(it);
                                                           // here taking lowest 10 values
```

```
171.
         sort(u.begin(), u.end());
172.
         while((int)u.size() > 10)
173.
            u.pop_back();
174. }
175.
                                       // W[u][0] will contain initial weight/weights at node u
176. vector<int> W[MAX][20];
177.
     vector<int> LCA(int u, int v) {
178.
         vector<int> T;
         if(level[u] > level[v]) swap(u, v); // v is deeper
179.
180.
         int p = ceil(log2(level[v]));
         for(int i = p ; i \ge 0; --i)
                                                 // Pull up v to same level as u
181.
             if(level[v] - (1LL<<i) >= level[u]) {
182.
183.
                 MERGE(T, W[v][i]);
184.
                v = sparse[v][i];
185.
             }
         if(u == v) {
                                                    // if u WAS the parent
186.
             MERGE(T, W[u][0]);
187.
188.
             return T;
189.
         }
190.
         for(int i = p; i >= 0; --i)
                                                                             // Pull up u and v
     together while LCA not found
             if(sparse[v][i] != -1 \&\& sparse[u][i] != sparse[v][i]) { // -1 check is if}
191.
     2^i is out of calculated range
192.
                MERGE(T, W[u][i]);
193.
                 MERGE(T, W[v][i]);
194.
                 u = sparse[u][i], v = sparse[v][i];
195.
             }
                                       // As W[x][0] denoted the x nodes weight
196.
         MERGE(T, W[u][0]);
197.
         MERGE(T, W[v][0]);
                                       // every sparse node must be calculated
         MERGE(T, W[sparse[v][0]][0]); // we can also calculate summation of distance like this
198.
199.
         return T;
200.
     }
201.
202. // ------ Overlap Path of Tree ------
203.
204. // Note: DfsTiming and isChild function required
     // a is upper node of path a-b and c is upper node of path c-d
205.
     pii overlapPath(int a, int b, int c, int d) {      // returns number of common path of c-d
206.
     and a-b
207.
         // path a-b and c-d overlaps iff b is a child of c or d or both of c&d
208.
         if(not isChild(b, c)) return {0, 0};
209.
         int u = LCA(b, d);
                                       // u is the lowest point on which c-d and a-b overlaps
210.
         if(level[a]>level[c]) {
                                            // a is below c
                                        // also u is child of a
211.
             if(isChild(u, a))
212.
                 return {a, u};
213.
         }
         else {
                                         // c is above a
214.
215.
             if(isChild(u, c))
216.
                 return {c, u};
217.
         }
218.
         return {0, 0};
                                       // no common path found
219.
     }
220.
221.
     int EdgeCount(int a, int b, int c, int d) {
                                                            // Finds number of edges if we join
     nodes a, b and
         a = Convert(a), b = Convert(b), c = Convert(c), d = Convert(d); // want to find path
222.
     from c to d
223.
       int u = LCA(a, b);
```

```
224.
         int v = LCA(c, d);
225.
         int ans = dist(c, d, v);
226.
         pii tt;
227.
         // connected paths are u->a & u->b
228.
         // query paths are v->c & v->d
229.
         // cases:
230.
         // u->a overlaps v->c
231.
         tt = overlapPath(v, c, u, a);
         ans -= tt.fi == 0? 0:dist(tt.fi, tt.se, LCA(tt.fi, tt.se));
232.
233.
         // u->a overlaps v->d
         tt = overlapPath(v, c, u, b);
234.
         ans -= tt.fi == 0? 0:dist(tt.fi, tt.se, LCA(tt.fi, tt.se));
235.
236.
         // u->b overlaps v->c
237.
         tt = overlapPath(v, d, u, a);
         ans -= tt.fi == 0? 0:dist(tt.fi, tt.se, LCA(tt.fi, tt.se));
238.
         // u->b overlaps v->d
239.
         tt = overlapPath(v, d, u, b);
240.
         ans -= tt.fi == 0? 0:dist(tt.fi, tt.se, LCA(tt.fi, tt.se));
241.
242.
         return ans;
243.
     }
244.
     // ---- return k'th node if we traverse from node u to v of a tree
245.
246.
     // NOT TESTED!!
247.
248.
     int getKthNode(int u, int v, int k, int lca) {
         int lftChain = lvl[u] - lvl[lca] + 1;
249.
         int rhtChain = lvl[v] - lvl[lca];
250.
         if(k == 1) return u;
251.
252.
         if(lca == v) {
253.
             for(int i = 20; i >= 0; --i)
254.
                  if(k - (1 << i) >= 1)
                      u = par[u][i], k -= (1 << i);
255.
256.
             return u;
257.
         }
258.
         if(lca == u) {
259.
             k = rhtChain+1-k;
             for(int i = 20; i >= 0; --i)
260.
261.
                  if(k - (1 << i) >= 0)
                      v = par[v][i], k -= (1 << i);
262.
263.
             return v;
264.
         }
265.
         if(k > lftChain) {
266.
             k -= lftChain;
267.
             k = rhtChain - k;
268.
             for(int i = 20; i >= 0; --i)
                  if(k - (1 << i) >= 0)
269.
270.
                      v = par[v][i], k -= (1 << i);
271.
             return v;
272.
         }
273.
         for(int i = 20; i >= 0; --i)
274.
             if(k - (1 << i) >= 1)
275.
                  u = par[u][i], k -= (1 << i);
276.
         return u;
277. }
278.
279.
    // ----- SUBTREE UPDATE FUNCTIONS -----
280.
     // if the root changes
281.
```

```
282. void subTreeUpdate(int u, int root, int val) {
283.
         // if u is child of root, then subtree of u
284.
         if(u == root)
285.
             DS.update(in[1], out[1], val);
286.
         else if(isChild(u, root))
             DS.update(in[u], out[u], val);
287.
         // if root is child of u
288.
         else if(isChild(root, u)) {
289.
                                           // get the first child of u
290.
             int x = getChild(root, u);
291.
             DS.update(in[1], out[1], val);
             DS.update(in[x], out[x], -val);
292.
         }
293.
294.
         else
             DS.update(in[u], out[u], val);
295.
296.
     }
297.
     11 getSubTreeSum(int u, int root) {
298.
         // if u is child of root, then subtree of u
299.
300.
         if(u == root)
301.
             return DS.query(in[1], out[1]);
302.
         if(isChild(u, root))
             return DS.query(in[u], out[u]);
303.
304.
         // if root is child of u
305.
         else if(isChild(root, u)) {
             int x = getChild(root, u);
                                             // get the first child of u
306.
307.
             return DS.query(in[1], out[1]) - DS.query(in[x], out[x]);
308.
         }
309.
         else
310.
             return DS.query(in[u], out[u]);
311.
     }
312.
     // ----- Can Give Total Spanning Tree edges for an particular set of nodes
313.
314.
315.
     set<int>nodes;
                                      // contains nodes according to dfs order
                                      // returns node Query/Insert updated distance
316.
     int nodeCost(int u) {
         auto it = nodes.insert(in[u]).first;
                                                      // inserted according to dfs in-timing
317.
         auto 1 = it, r = it;
                                                       // iterator of the inserted index
318.
319.
         if(it == nodes.begin())
             1 = --nodes.end();
320.
321.
         else --1;
322.
323.
         if(it == --nodes.end())
324.
             r = nodes.begin();
325.
         else ++r;
326.
327.
         int L = revIn[*1], R = revIn[*r];
                                             // nodes are retrieved from dfs in-timing
328.
329.
         // dst is the spanning distance if the new node is added
         int dst = lvl[u] + lvl[LCA(L, R)] - lvl[LCA(u, L)] - lvl[LCA(u, R)];
330.
331.
         return dst;
332.
     }
333.
334.
     void removeNode(int u) {
         nodes.erase(in[u]);
335.
336.
     }
337.
338.
     struct LCATree {
339.
         int tree[MAX*4], lca, n, cost;
```

```
340.
          set<int>nodes;
341.
342.
          void init(int sz) { n = sz, lca = -1, cost = 0; }
343.
          void update(int pos, int 1, int r, int idx, int v) {
344.
              if(1 == r) {
345.
                  tree[pos] += v;
346.
                  return;
347.
              }
              int mid = (1+r)>>1;
348.
349.
              if(idx <= mid) update(pos<<1, 1, mid, idx, v);</pre>
350.
                               update(pos<<1|1, mid+1, r, idx, v);
              tree[pos] = tree[pos<<1] + tree[pos<<1|1];</pre>
351.
352.
          int query(int pos, int 1, int r, int L, int R) {
353.
354.
              if(r < L \text{ or } R < 1)
                                        return 0;
355.
              if(L \le 1 \text{ and } r \le R)
                                        return tree[pos];
              int mid = (1+r)>>1;
356.
357.
              return query(pos<<1, 1, mid, L, R) + query(pos<<1|1, mid+1, r, L, R);</pre>
358.
          }
359.
          int getPar(int u, int p) {
              for(int i = 20; i >= 0; --i)
360.
                  if(p & (1 << i))
361.
362.
                       u = par[u][i];
                                                // parent sparse table
363.
              return u;
364.
          }
          int LCA() {
365.
              int u = *nodes.begin(), tot = nodes.size(), v, ret = *nodes.begin();
366.
367.
              int lo = 0, hi = lvl[u]-1;
368.
              while(lo <= hi) {</pre>
369.
                  int mid = (lo+hi)>>1;
370.
                  v = getPar(u, mid);
                  if(query(1, 1, n, in[v], out[v]) == tot)
371.
                       hi = mid-1, ret = v;
                                                          // in : dfs in time
372.
373.
                  else
                                                          // out : dfs out time
374.
                       lo = mid+1;
375.
              }
376.
              return ret;
377.
378.
          int findChainPar(int u, int t) {
                                                             // finds parent node of u having
379.
              int lo = 0, hi = lvl[u]-1, ret = u, v, mid; // active child node more than t
380.
              while(lo <= hi) {</pre>
381.
                  mid = (lo+hi) >> 1;
382.
                  v = getPar(u, mid);
383.
                  if(query(1, 1, n, in[v], out[v]) > t)
384.
                       hi = mid-1, ret = v;
385.
                  else
386.
                       lo = mid+1;
387.
              }
388.
              return ret;
389.
          }
390.
          void addNode(int u) {
391.
              int pstLca = lca;
392.
              nodes.insert(u), update(1, 1, n, in[u], 1);
393.
              if(lca == -1) {
394
                  lca = u;
395.
                  return;
396.
              }
397.
              else
```

```
398.
                  lca = LCA();
399.
              // if new LCA is same but the new node is on different chain
400.
              if(pstLca == lca and query(1, 1, n, in[u], out[u]) == 1) {
401.
                  int v = findChainPar(u, 1);
402.
                  cost += lvl[u] - lvl[v];
403.
              }
              // if new LCA changes, newLCA will always be upper from past LCA
404.
405.
              // also the node u is on different chain
              else if(lca != pstLca)
406.
407.
                  cost += lvl[u] + lvl[pstLca] - 2*lvl[lca];
408.
         }
409.
         void removeNode(int u) {
410.
              int pstLca = lca;
411.
              nodes.erase(u), update(1, 1, n, in[u], -1);
412.
              if(nodes.empty()) {
413.
                  lca = -1, cost = 0;
                  return;
414.
415.
              }
              else
416.
417.
                  lca = LCA();
418.
              if(pstLca == lca and query(1, 1, n, in[u], out[u]) == 0) {
                  int v = findChainPar(u, 0);
419.
                  cost -= lvl[u] - lvl[v];
420.
421.
              }
422.
              else if(lca != pstLca)
423.
                  cost -= lvl[lca] + lvl[u] - 2*lvl[pstLca];
424.
     }};
425.
426.
     //Trie
427.
     //Complexity : making a trie : O(S), searching : O(S)
428.
     bool found;
429.
     struct node {
430.
431.
         bool isEnd;
432.
         node *next[11];
433.
         node() {
434.
              isEnd = false;
435.
              for(int i = 0; i < 10; i++)
                  next[i] = NULL;
436.
437.
         }};
438.
439.
     //trie of a string abc, ax
     // [start] --> [a] --> [b] --> [c] --> endMark
440.
     //
441.
442.
                     [x] --> endMark
443.
     //creates trie, returns true if the trie we are creating is a segment of a string
     //to only create a trie remove lines which are comment marked
444.
445.
446.
     bool create(char str[], int len, node *current) {
447.
         for(int i = 0; i < len; i++) {
448.
              int pos = str[i] - '0';
449.
              if(current->next[pos] == NULL)
450.
                  current->next[pos] = new node();
              current = current->next[pos];
451.
452.
              if(current->isEnd) //
453.
                                   //
                  return true;
454.
         }
455.
         current->isEnd = true; //
```

```
456.
                                    //
         return false;
457.
     }
458.
459.
     void del(node *current) {
460.
         for(int i = 0; i < 10; i++)
              if(current->next[i] != NULL)
461
462.
                  del(current->next[i]);
463.
         delete current;
464.
     }
465.
     void check(node *current) {
466.
467.
         for(int i = 0; i < 10; i++) {
468.
              if(current->next[i] != NULL)
469.
                  check(current->next[i]);
470.
         if(found) return;
471.
         if(current->isEnd && !found) {
472.
              for(int i = 0; i < 10 && !found; i++)</pre>
473.
474.
                  if(current->next[i] != NULL) {
475.
                       found = 1;
476.
     }}}
477.
478.
     // NON-Dynamic implementation
     // root node is at 0 index of tree
479.
480.
     // root node counter contains total number of string insertion
     // each inserted char counter is on the child node of the edges
481.
482.
483.
     struct Trie {
484.
         struct node {
485.
              int cnt;
486.
              int nxt[60];
487.
         };
         int nodes;
488.
489.
         node tree[MAX];
490.
         void newNode() {
491.
              tree[nodes].cnt = 0;
              memset(tree[nodes].nxt, -1, sizeof tree[nodes].nxt);
492.
493.
              ++nodes;
         }
494.
495.
         void init() {
496.
              nodes = 0;
497.
              newNode();
498.
         }
499.
         int getId(char x) {
500.
              if(x \ge 'A') and x \le 'Z')
                  return (x - 'A' + 27);
501.
502.
              return (x - 'a' + 1);
503.
         void insert(string &str, int len = 0, int idx = 0) {
504.
505.
              tree[idx].cnt++;
506.
              if(len == str.size()) return;
507.
              int id = getId(str[len]);
508.
              if(tree[idx].nxt[id] == -1) {
                  tree[idx].nxt[id] = nodes;
509.
510.
                  newNode();
511.
512.
              insert(str, len+1, tree[idx].nxt[id]);
513.
         }
```

```
514.
         int search(string &str, int len = 0, int idx = 0) {
515.
             if(len == str.size())
516.
                 return -2;
517.
             int id = getId(str[len]);
             if(tree[idx].nxt[id] == -1)
518.
519.
                 return -1;
             if(tree[idx].cnt == 1)
520.
521.
                 return len;
             return search(str, len+1, tree[idx].nxt[id]);
522.
523. }};
```