

- Case - $\left\{ \begin{array}{l} \text{Best: } O(1) \\ \text{Average: } O(\log n) \\ \text{Worst: } O(\log n) \end{array} \right.$

Binary Search

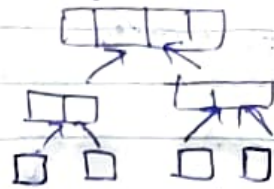
1. Def. binary Search (A, x):
2. $n = \text{len}(A)$
3. $\text{beg} = 0$
4. $\text{end} = n - 1$
5. $\text{result} = -1$
6. While ($\text{beg} \leq \text{end}$):
7. $\text{mid} = (\text{beg} + \text{end}) / 2$
8. If ($A[\text{mid}] \leq x$):
9. $\text{beg} = \text{mid} + 1$
10. $\text{result} = \text{mid}$
11. Else:
12. $\text{end} = \text{mid} - 1$
13. Return result

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- Case - $\begin{cases} \text{Best} = O(n \log n) \\ \text{Average} = O(n \log n) \\ \text{Worse} = O(n \log n) \end{cases}$

Merge Sort

1. Start
2. Declare array & left, right, mid, variable
3. Perform merge function
if left > right
return



- $mid = (left + right) / 2$
 merge sort (array, left, mid)
 merge sort (array, mid+1, right)
 merge (array, left, mid, right)
4. Stop

Merge Sort (array, l, r)

5. If $l > r$
 - find middle point to divide array into two halves:
 - middle $m = l + (r - l) / 2$
 - call merge sort for first half:
 - call merge sort (array, l, m)
 - call merge sort for second half:
 - call merge sort (array, m+1, r)
 - Merge the two halves sorted in step 2 & 3.
 - call merge (array, l, m, r)

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Case $\left\{ \begin{array}{l} \text{Best: } O(n \log n) \\ \text{Average: } O(n \log n) \\ \text{Worst: } O(n^2) \end{array} \right.$

Quick Sort

QUICKSORT (array A, start, end)

1. {
2. if (start < end)
3. {
4. p = partition (A, start, end)
5. QUICKSORT (A, start, p-1)
6. QUICKSORT (A, p+1, end)
7. }
8. }

Partition Algorithm :

Partition (array A, start, end)

1. { pivot = A[end]
2. i = start - 1
3. for j = start to end - 1 {
4. do if (A[j] < pivot) {
5. then i = i + 1
6. swap A[i] with A[j]
7. }
8. swap A[i+1] with A[end]
9. return i + 1
10. }

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