

Lecture 2

Propositional Logic

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FIT2014 Theory of Computation

Overview

- Propositions
- Connectives
- Tautologies
- Arguments
- Representing Knowledge

Definition

A proposition is

- A statement

which is

- Either true or false.

Examples

- **$2 + 2 = 4$**
 - is a proposition which is **true**.
- **The moon is made of cheese**
 - is a proposition which is **false**.
- **It will rain tomorrow**
 - is a proposition.
- **This statement is false**
 - is **not** a proposition.
- **Vote for Mickey Mouse**
 - is **not** a proposition.

Negation

P : I have three children.

$\neg P$: I do **not have three children.**

Other Notation: $\sim P$, \overline{P}

Truth table:

P	$\neg P$
F	T
T	F

Connectives

- And \wedge (&)
- Or \vee
- Implies \Rightarrow (\rightarrow)
- Equivalence \Leftrightarrow (\leftrightarrow)

Conjunction

P : This subject is interesting.

Q : I am tired.

$P \wedge Q$:

- This subject is interesting **and** I am tired.
- This subject is interesting **although** I am tired.
- This subject is interesting **but** I am tired.

P	Q	$P \wedge Q$
F	F	F
F	T	F
T	F	F
T	T	T

Disjunction

P : Sue is a football player.

Q: Bob is lazy.

$P \vee Q$: Sue is a football player **or Bob is lazy.**

Note: Disjunction is inclusive.

P	Q	$P \vee Q$
F	F	F
F	T	T
T	F	T
T	T	T

De Morgan's Laws

$$\neg(P \vee Q) = \neg P \wedge \neg Q$$

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Can prove using truth tables.

E.g.:



Augustus De Morgan (1806-1871)

http://www-history.mcs.st-andrews.ac.uk/Biographies/De_Morgan.html

P	Q	$P \vee Q$	$\neg(P \vee Q)$
F	F	F	T
F	T	T	F
T	F	T	F
T	T	T	F

$\neg P$	$\neg Q$	$\neg P \wedge \neg Q$
T	T	T
T	F	F
F	T	F
F	F	F

Conditional

P: It is Tuesday.

Q: We are in Belgium.

$P \Rightarrow Q$:

- **If** it is Tuesday **then** we are in Belgium.
- It's being Tuesday **implies** we are in Belgium.
- It's Tuesday **only if** we are in Belgium.
- It's being Tuesday is **sufficient** for us to be in Belgium.

Conditional Table

P	Q	$P \Rightarrow Q$
F	F	T
F	T	T
T	F	F
T	T	T

P: Grace is a COBOL expert.

Q: Grace can program.

$P \Rightarrow Q$:

If Grace is a COBOL expert then she can program.



Grace Hopper (1906-1992)

[http://www.cs.uoregon.edu/
groups/wics/projects.php](http://www.cs.uoregon.edu/groups/wics/projects.php)

Biconditional

P : The world will blow up.

Q : KAOS rules the world.

$P \Leftrightarrow Q$:

- The world will blow up **if and only** if KAOS rules the world.
- KAOS ruling the world is a **necessary and sufficient** condition that the world will blow up.

Biconditional Table

P	Q	$P \Leftrightarrow Q$
F	F	T
F	T	F
T	F	F
T	T	T

P: It will rain tomorrow.

Q: I will wear a raincoat tomorrow.

$P \Leftrightarrow Q$:

I will wear a raincoat tomorrow if and only if it rains.

Interpretation

The truth table for a statement provides a record for all possible interpretations of that statement.

Example:

S: If the price is less than \$30 and I have at least \$50, then I will buy that CD.

- P: The price of the CD is less than \$30.
- L: I have at least \$50.
- B: I will buy that CD.
- S: $(P \wedge L) \Rightarrow B$

P	L	B	S
F	F	F	T
T	F	F	T
F	T	F	T
T	T	F	F
F	F	T	T
T	F	T	T
F	T	T	T
T	T	T	T

Definitions

A **tautology** is a statement whose interpretation is always true.

We say two statements P and Q are **logically equivalent** if they always have the same interpretation.

- **Their truth tables are the same.**
- **i.e., $P \Leftrightarrow Q$ is a**

Examples of logical equivalence

$P \Rightarrow Q$ is logically equivalent to
 $\neg P \vee Q$

$P \Leftrightarrow Q$ is logically equivalent to
 $(\neg P \vee Q) \wedge (P \vee \neg Q)$

These can be proved using truth tables.

Fitness Example

S: If I exercise then I will get fit, and I do exercise, so I will get fit.

E: I do exercise.

F: I will get fit.

$$\mathbf{S: (E \Rightarrow F) \wedge E \Rightarrow F}$$

Definition

An **argument** consists of:

- A set of propositions, $\mathbf{P}_1, \dots, \mathbf{P}_n$, called the **premises**.
- Another proposition, \mathbf{C} , called the **conclusion**.

An argument is called **valid** if the statement

$$\mathbf{P}_1 \wedge \dots \wedge \mathbf{P}_n \Rightarrow \mathbf{C}$$

is a tautology.

Mary's Exam Example

Today Mary has a Law exam or a Computer Science exam or both. She doesn't have a Law exam. Therefore she must have a Computer Science exam.

L: Mary has a Law exam today.

C: Mary has a Computer Science exam today.

Premises: **$L \vee C, \neg L$**

Conclusion: **C**

Wumpus World

[based on a video game by Gregory Yob, c1972]

- The idea of the game is to find the gold.
- The cave has rooms that lie in a grid.
- Dangers
 - **The Wumpus** (You can smell a Wumpus in the next room or the same room)
 - **Pits** (You can feel a breeze in the next room)
 - **Bats** (You can hear the bats in the next room)
- You can use to kill a Wumpus with an arrow.

Example

			P
W	G	P	
			B
S		P	

- W represents the Wumpus
- P represents a Pit
- B represents Bats
- S is the Starting Position
- G represents where the gold is.

A Game

- No stench or breeze in square 1,1
 - **Therefore Wumpus is not in 1,2 or 2,1**
- Suppose you move to square 2,1 and detect a breeze and no stench there.
 - **Therefore there is a Pit in either 2,2 or 3,1**
- So you go back and up to square 1,2 and detect a stench and no breeze.
 - **Therefore the Wumpus is in square 1,3.**

Notation

- $W_{1,1}$
 - **The Wumpus is in square 1,1.**
- $S_{1,2}$
 - **A stench was detected in square 1,2.**
- $B_{2,1}$
 - **A breeze was detected in square 2,1.**

Etc.

Knowledge

- $P_1: \neg \mathbf{S}_{1,1} \wedge \neg \mathbf{B}_{1,1}$
- $P_2: \neg \mathbf{S}_{2,1} \wedge \mathbf{B}_{2,1}$
- $P_3: \mathbf{S}_{1,2} \wedge \neg \mathbf{B}_{1,2}$
- $P_4: \neg \mathbf{S}_{1,1} \Rightarrow \neg \mathbf{W}_{1,1} \wedge \neg \mathbf{W}_{1,2} \wedge \neg \mathbf{W}_{2,1}$
- $P_5: \neg \mathbf{S}_{2,1} \Rightarrow \neg \mathbf{W}_{1,1} \wedge \neg \mathbf{W}_{2,1} \wedge \neg \mathbf{W}_{2,2} \wedge \neg \mathbf{W}_{3,1}$
- $P_6: \neg \mathbf{S}_{1,2} \Rightarrow \neg \mathbf{W}_{1,1} \wedge \neg \mathbf{W}_{1,2} \wedge \neg \mathbf{W}_{2,2} \wedge \neg \mathbf{W}_{1,3}$
- $P_7: \mathbf{S}_{1,2} \Rightarrow \mathbf{W}_{1,3} \vee \mathbf{W}_{1,2} \vee \mathbf{W}_{2,2} \vee \mathbf{W}_{1,1}$

Wumpus Argument

- Want to obtain the conclusion: $W_{1,3}$
- Need to show:

$$\mathbf{P_1 \wedge P_2 \wedge P_3 \wedge P_4 \wedge P_5 \wedge P_6 \wedge P_7 \Rightarrow W_{1,3}}$$

is a tautology.

- Truth Table has $2^{12} = 4096$ rows.

Problems

- Need to introduce a lot of propositions to represent any useful knowledge.
- Using truth tables to show validity requires:
 - **Exponential Space**
 - **Exponential Time**

History

- George Boole
1815-1864



<http://www-history.mcs.st-andrews.ac.uk/Biographies/Boole.html>

Revision

- Propositions
- Connectives
- Tautologies & Logical Equivalence
- Arguments

Reading

- Sipser, pp. 14-15, 21-25.