

# Lecture 28

## Iteration vs Recursion

FIT 1008  
Introduction to Computer Science



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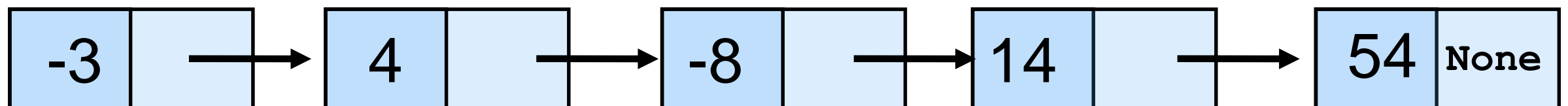
Do not remove this notice.

Operation	Class Method
str(obj)	__str__(self)
len(obj)	__len__(self)
item in obj	__contains__(self,item)
y = obj[ndx]	__getitem__(self,ndx)
obj[ndx] = value	__setitem__(self,ndx,value)
obj == rhs	__eq__(self,rhs)
obj < rhs	__lt__(self,rhs)
...	
obj + rhs	__add__(self,rhs)
...	

```
class List:
    def __init__(self):
        self.head = None
        self.count = 0
```

```
def __len__(self):
    return self.count
```

?



count from head to access elements

```
class List:  
    def __init__(self):  
        self.head = None  
  
    def __len__(self):
```

```
class List:
    def __init__(self):
        self.head = None

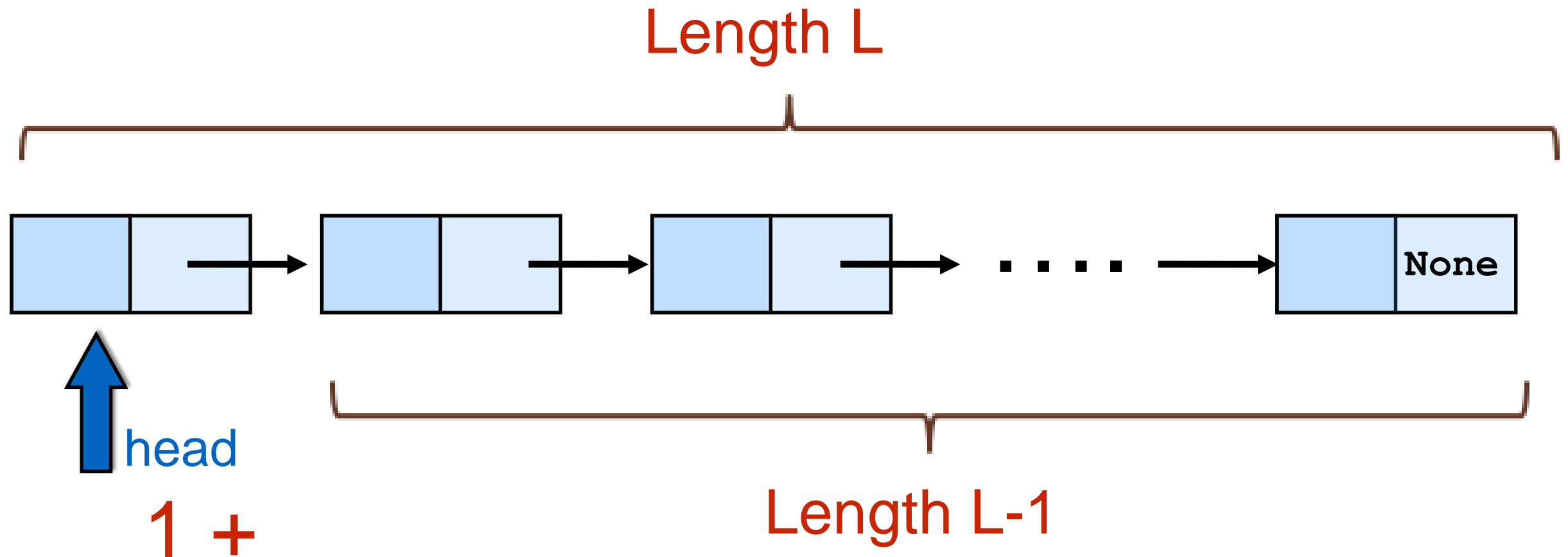
    def __len__(self):
        current = self.head
        count = 0
        while current is not None:
            current = current.next
            count += 1
        return count
```

```
class List:
    def __init__(self):
        self.head = None

    def __len__(self):
        current = self.head
        count = 0
        while current is not None:
            current = current.next
            count += 1
        return count
```

**Complexity:**  $O(n)$  where  $n$  is the size of the list.

What about recursively?



**Convergence:** Call recursion with  $L-1$ . Use variable *current*.

**Base case:** Empty? Size of empty list is 0.

**Combining solutions:** Add up result of recursive call +



```
def __len__(self):
```

```
def __len__(self):  
    if self.head is None:  
        return 0  
    else:  
        return 1 + self.__len__(self.head.next)
```



```
def __len__(self):  
    if self.head is None:  
        return 0  
    else:  
        return 1 + self.__len__(self.head.next)
```

TypeError: \_\_len\_\_() takes 1 positional argument but 2 were given

```
def length(self, current):
```

```
def _length(self, current):  
    if self.current is None: # base case  
        return 0  
    else:  
        return 1 + self._length(current.next)
```

Base case

Convergence

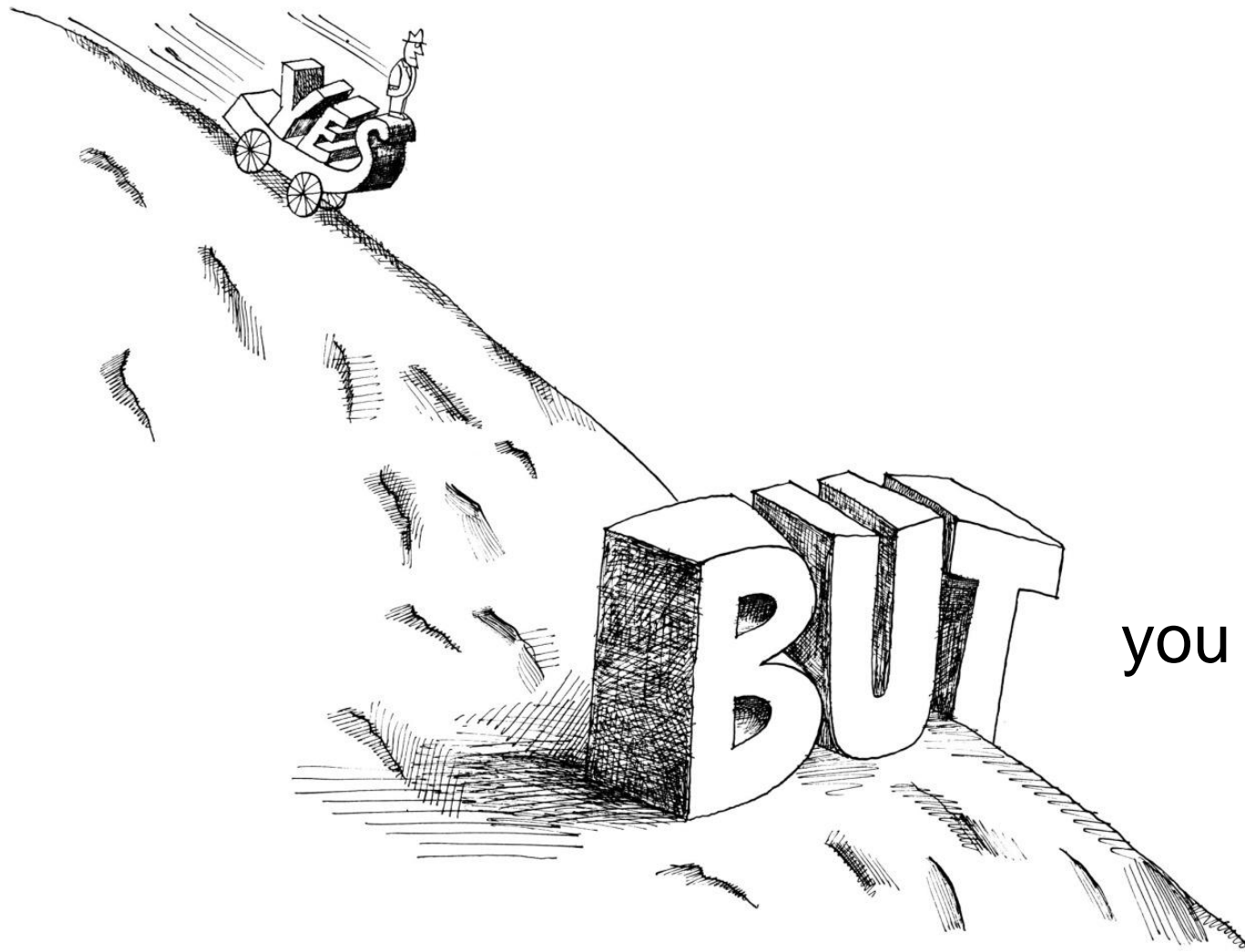
Combination

```
def __len__(self):  
    return self._length(self.head)
```

Auxiliary method sets up the initial  
parameters

# Recursion vs Iteration

- Can every **iterative function** be implemented using **recursion**?  
Yes, it is **straightforward**.
- Can every **recursive function** be implemented using **iteration**?



you might also need to **store** past results.

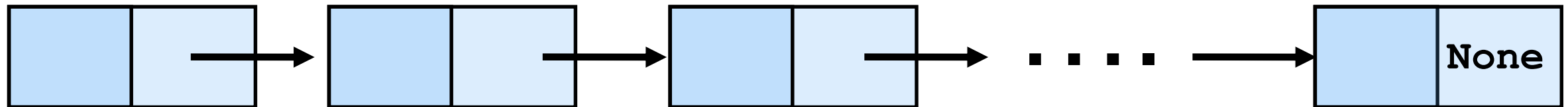
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[https://66.media.tumblr.com/3f1bb6a13c96f90cb332b385a6363b2e/tumblr\\_nkaavk9owg1qz6f4bo2\\_1280.jpg](https://66.media.tumblr.com/3f1bb6a13c96f90cb332b385a6363b2e/tumblr_nkaavk9owg1qz6f4bo2_1280.jpg)

Operation	Class Method
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obj < rhs	__lt__(self,rhs)
...	
obj + rhs	__add__(self,rhs)
...	

\_\_\_contains\_\_\_

current



head



# \_\_contains\_\_

```
def __contains__(self, item):
```

# \_\_contains\_\_

```
def __contains__(self, item):  
    current = self.head  
    while current is not None:  
        if current.item == item:  
            return True  
        current = current.next  
    return False
```

**Complexity:** Worst case -  $O(n)$  where  $n$  is the size of the list.

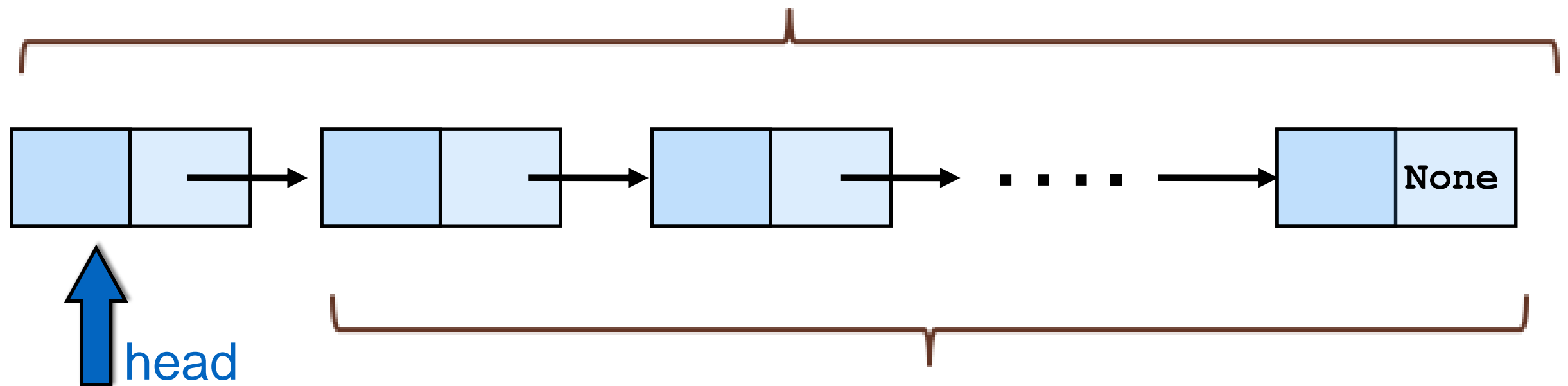
# \_\_contains\_\_

```
def __contains__(self, item):  
    current = self.head  
    while current is not None:  
        if current.item == item:  
            return True  
        current = current.next  
    return False
```

- **Best case** complexity when found first:  $O(1) * \text{CompEq}$  where  $\text{CompEq}$  is the complexity of  $==$  (or  $__eq__$ )
- **Worst case** when not found:  $O(n) * \text{CompEq}$  where  $n$  is the length of the list.

What about recursively?

\_\_contains\_\_(6) in L



$6 == \text{head.item}$  **or** \_\_contains\_\_(6) in L-1

**Convergence:** Call recursion with L-1. Use variable *current*.

**Base case:** Empty or Element Found. We need both.

**Combining solutions:** it's in the head **or** in the remaining list.

```
def __contains__(self, item):
```

Auxiliary method sets up the initial parameters

```
def __contains__(self, item):  
    return self._contains_aux(self.head, item)
```

```
def _contains_aux(self, current, item):  
    if current is None:  
        return False  
    return current.head == item or self._contains_aux(current.next, item)
```

Base case

Convergence

Combination

# Alternative coding

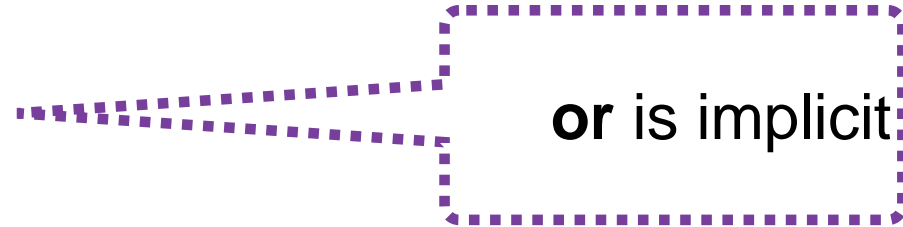
```
def __contains__(self, item):  
    return self._contains_aux(self.head, item)
```





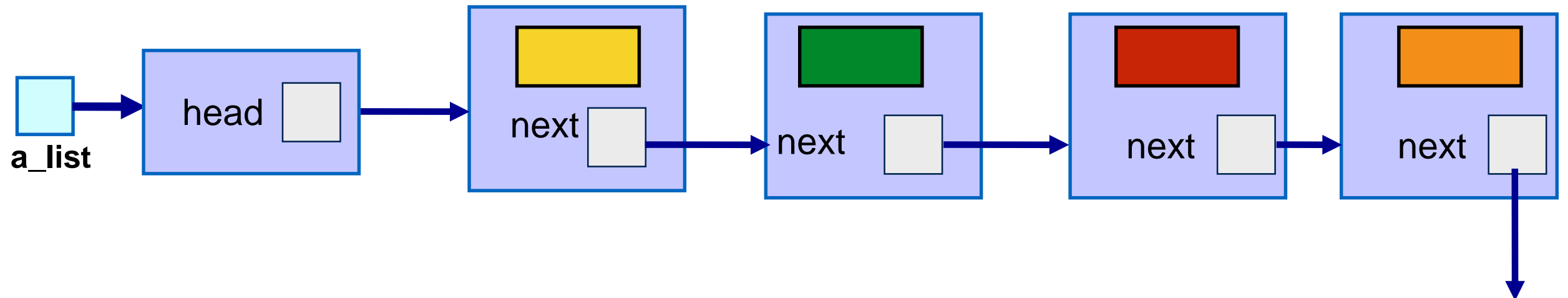
# Alternative coding

```
def __contains__(self, item):  
    return self._contains_aux(self.head, item)
```

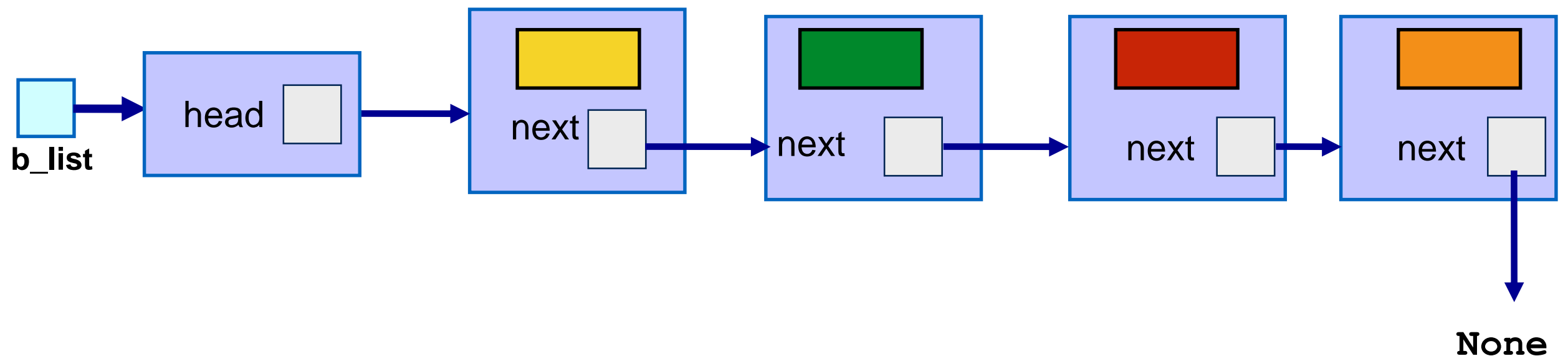
```
def _contains_aux(self, current, item):  
    if current is None:  
        return False  
    elif current.head == item:   
        return True  
    else:  
        return self._contains_aux(current.next, item)
```

If complexity is the same, why bother with recursive implementations?

# Copy Linked Lists



```
b_list = a_list.copy()
```

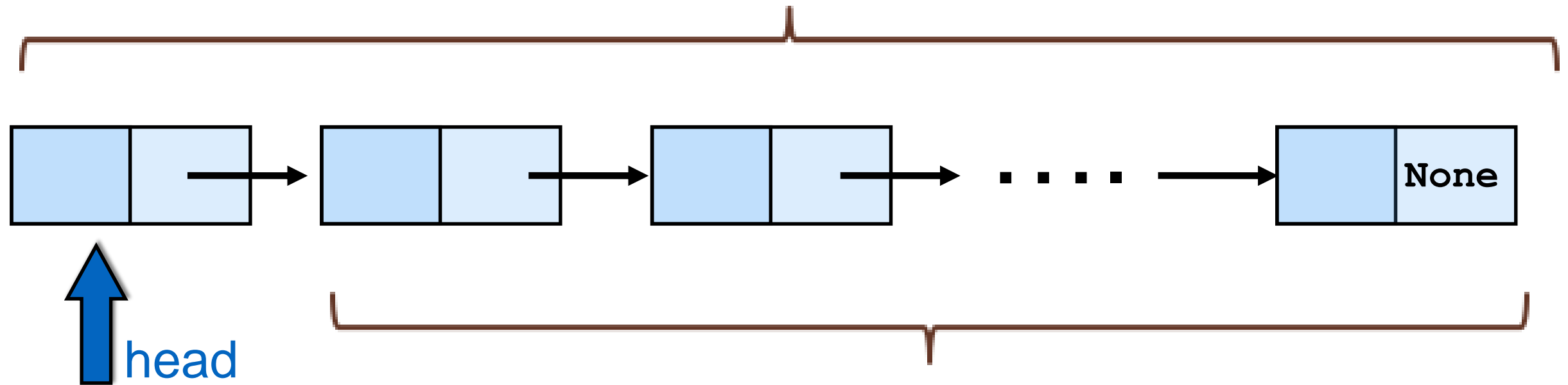


```
def copy(self):
```

```
def copy(self):  
    new_list = List()  
    for item in self:  
        new_list.insert(len(new_list), item)  
    return new_list
```

What about recursively?

`_copy(self.head, new_list)`



`_copy(self.head.next, new_list)`

`new_list.insert(0, head.item)`

# copy

```
def copy(self):
```



# copy

Auxiliary method sets up the initial parameters

```
def copy(self):  
    new_list = List()  
    self._copy_aux(self.head, new_list)  
    return new_list
```

Convergence

```
def _copy_aux(self, node, new_list):  
    if node is not None:  
        self._copy_aux(node.next, new_list)  
        new_list.insert(0, node.item)
```

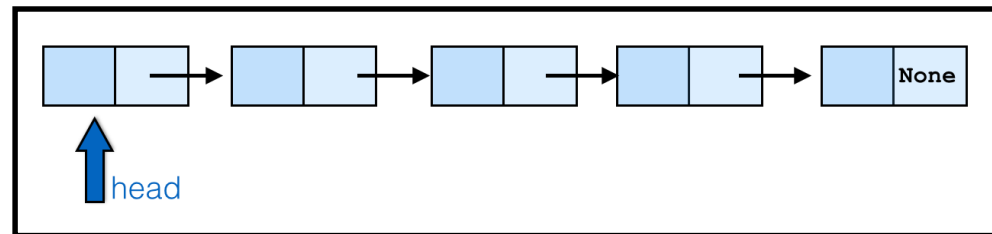
Base case

Combination

```
def copy(self):
    new_list = List()
    self._copy_aux(self.head, new_list)
    return new_list

def _copy_aux(self, node, new_list):
    if node is not None:
        self._copy_aux(node.next, new_list)
        new_list.insert(0, node.item)
```

$O(n)$



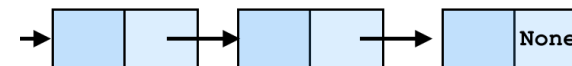
`_copy_aux`  $O(1)$



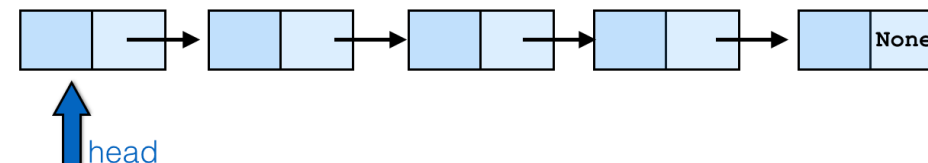
`_copy_aux`  $O(1)$



`_copy_aux`  $O(1)$



`_copy_aux`  $O(1)$



$n$  times

# Using iterators...

```
def copy(a_list):  
    new_list = List()  
    copy_aux(iter(a_list), new_list)  
    return new_list
```

```
def copy_aux(iter, new_list):  
    try:  
        item = next(iter)  
        copy_aux(iter, new_list)  
        new_list.insert(0, item)  
    except StopIteration:  
        pass
```

# copy

```
def copy(self):  
    new_list = List()  
    for item in self:  
        new_list.insert(len(new_list), item)  
    return new_list
```

$O(n^2)$

```
def copy(self):  
    new_list = List()  
    self._copy_aux(self.head, new_list)  
    return new_list  
  
def _copy_aux(self, node, new_list):  
    if node is not None:  
        self._copy_aux(node.next, new_list)  
        new_list.insert(0, node.item)
```

$O(n)$

# Or....

Have a tail attribute and you can have  
 $O(1)$  append and hence  $O(n)$  copy



# Advantages/Disadvantages of Recursion

- Advantages:
  - More natural
  - Easier to prove correct
  - Easier to analyse
- Disadvantages:
  - Run-time overhead depending on the quality of the compiler
  - Memory overhead (fewer local variables versus stack space for function call)





what's easy to prove  
Just keep doing 'you', and being who  
~~you are~~ and doing what feels natural  
to you.  
balancing your time and memory use  
— Taylor Swift —

AZ QUOTES

# Summary

- Recursive algorithms are characterised by:
  1. Existence of base cases
  2. Decomposition into simpler subproblems
  3. Combination of solutions to subproblems
- Recursive methods require:
  1. One or more base cases
  2. One or more recursive calls
  3. Convergence in the recursive calls
  4. Combination of sub-solutions