# Lecture 28 Iteration vs Recursion

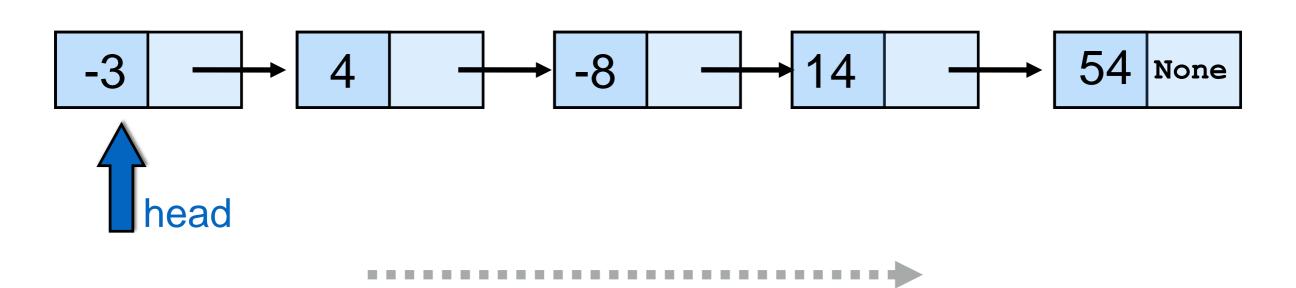
FIT 1008
Introduction to Computer Science



Operation	
str(obj)	str(self)
len(obj)	len(self)
item in obj	
y = obj[ndx]	
obj[ndx] = value	
obj == rhs	
obj < rhs	
obj + rhs	

```
class List:
    def __init__(self):
        self.head = None
    self.count = 0

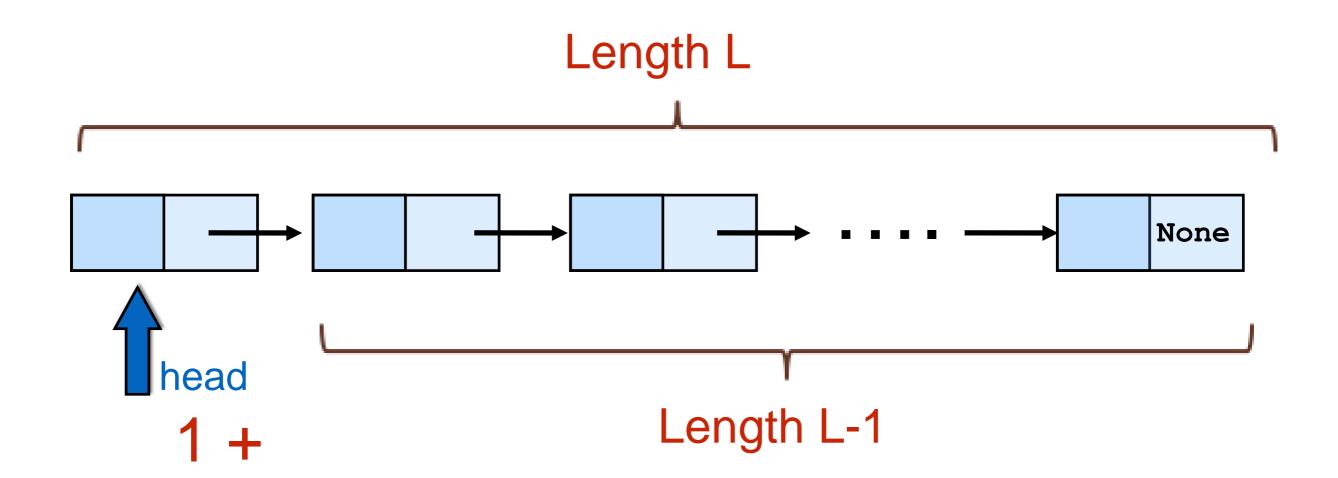
def __len__(self):
    return self.count
?
```



count from head to access elements

```
class List:
    def ___init___(self):
        self.head = None
    def ___len__(self):
        current = self.head
        count = 0
        while current is not None:
            current = current.next
            count += 1
        return count
```

## What about recursively?



Convergence: Call recursion with L-1. Use variable current.

Base case: Empty? Size of empty list is 0.

Combining solutions: Add up result of recursive call +

```
def __len__(self):
    if self.head is None:
        return 0
    else:
        return 1 + self.__len__(self.head.next)
```



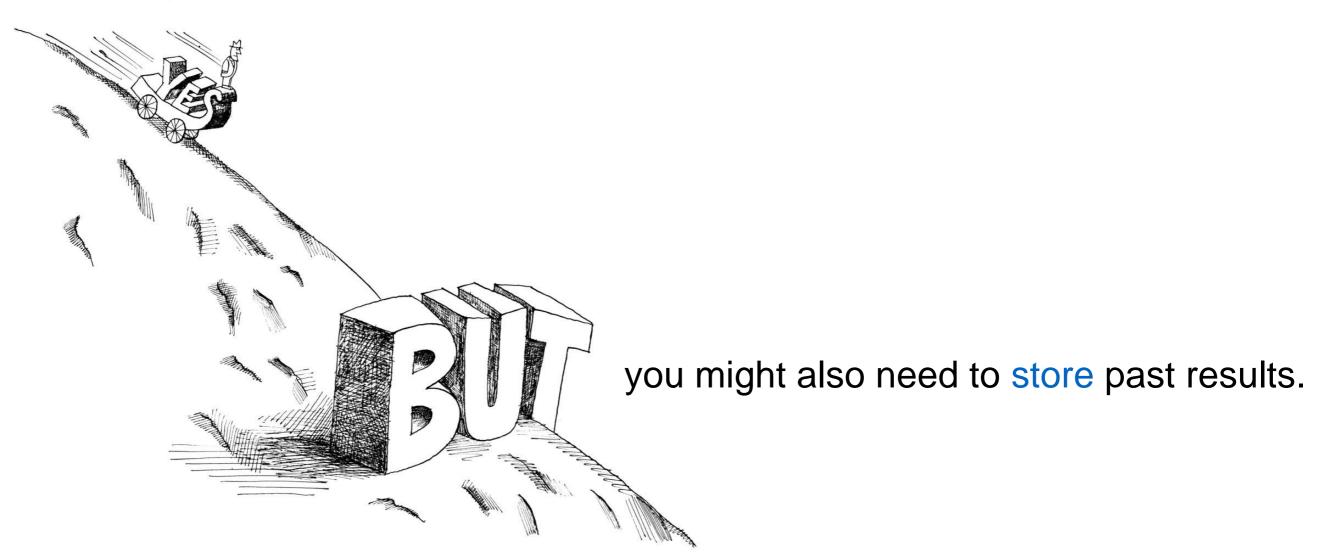
```
def __len__(self):
    if self.head is None:
        return 0
    else:
        return 1 + self.__len__(self.head.next)
```

TypeError: \_\_len\_\_() takes 1 positional argument but 2 were given

```
Base case
def _length(self, current):
    if self.current is None: # base case
         return 0
                                                 Convergence
    else:
         return 1 + self._length(current.next)
                          Combination
def len (self):
    return self._length(self.head)
                             Auxiliary method sets up the initial
                                    parameters
```

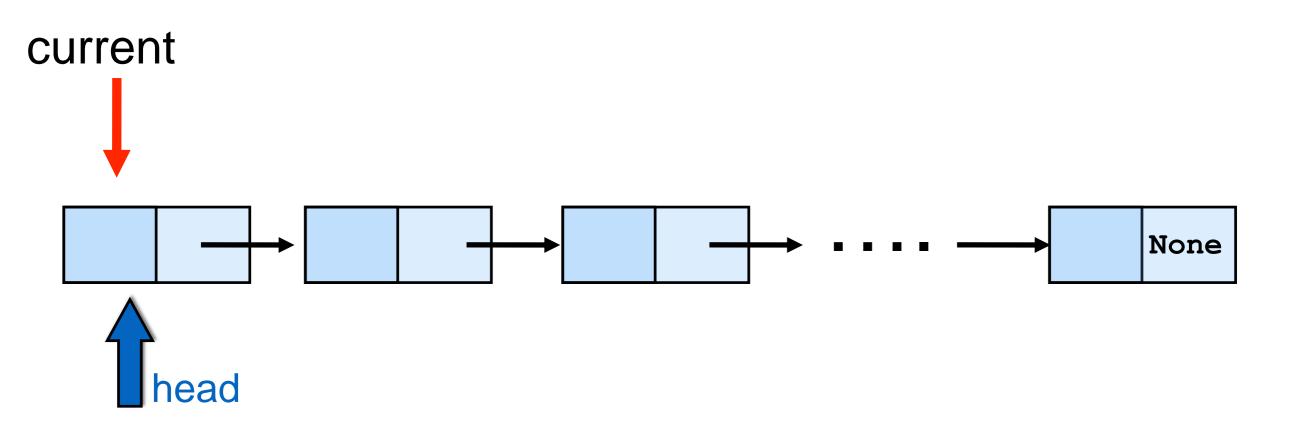
#### Recursion vs Iteration

- Can every iterative function be implemented using recursion?
   Yes, it is straightforward.
- Can every recursive function be implemented using iteration?



Operation	
str(obj)	
len(obj)	
item in obj	contains(self,item)
y = obj[ndx]	
obj[ndx] = value	
obj == rhs	
obj < rhs	
obj + rhs	

## \_\_contains\_\_\_



### \_\_contains\_\_\_

```
def __contains__(self, item):
    current = self.head
    while current is not None:
        if current.item == item:
            return True
        current = current.next
    return False
```

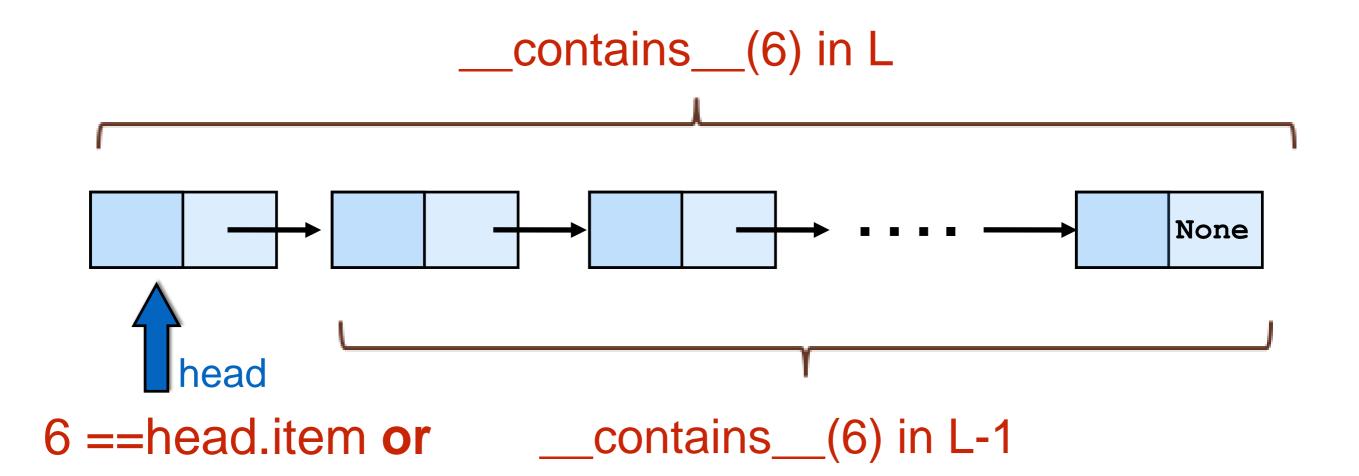
**Complexity:** Worst case - O(n) where n is the size of the list.

#### contains

```
def __contains__(self, item):
    current = self.head
    while current is not None:
        if current.item == item:
            return True
        current = current.next
        return False
```

- Best case complexity when found first: O(1)\*CompEq
  where CompEq is the complexity of == (or \_\_eq\_\_)
- Worst case when not found: O(n)\*CompEq where n is the length of the list.

## What about recursively?



Convergence: Call recursion with L-1. Use variable *current*.

Base case: Empty or Element Found. We need both.

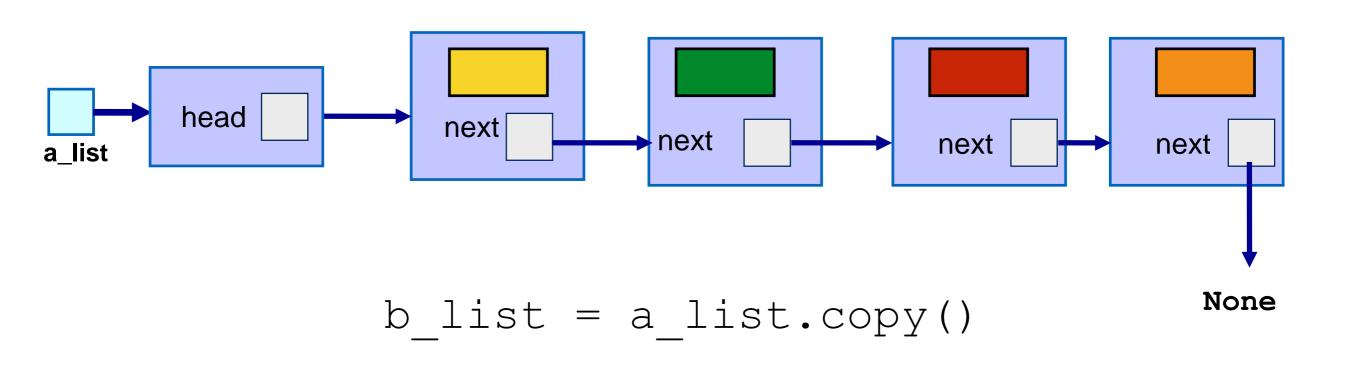
Combining solutions: it's in the head or in the remaining list.

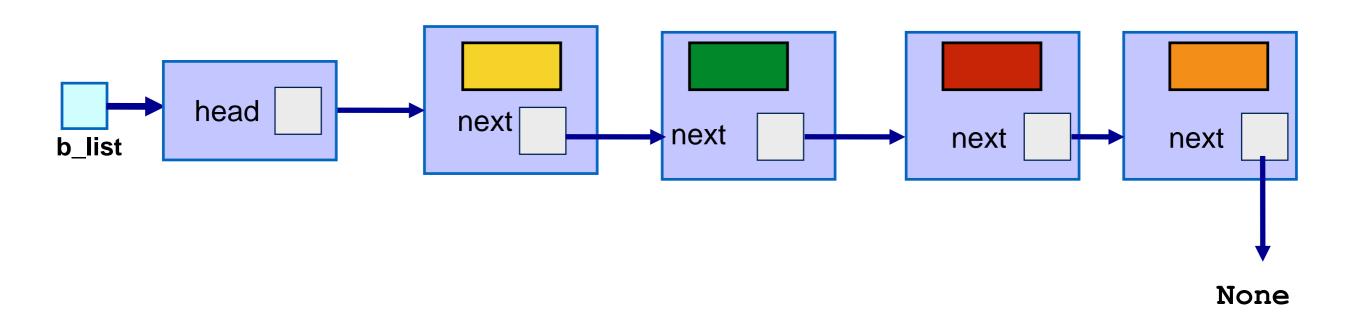
```
Auxiliary method sets up the initial
                      parameters
def __contains__(self, item):
    return self._contains_aux(self.head, item)
                                                                  Convergence
def _contains_aux(self, current, item):
                                                   Base case
    if current is None:
        return False
    return current.head == item or self._contains_aux(current.hext, item)
                                         Combination
```

#### Alternative coding

## If complexity is the same, why bother with recursive implementations?

#### Copy Linked Lists

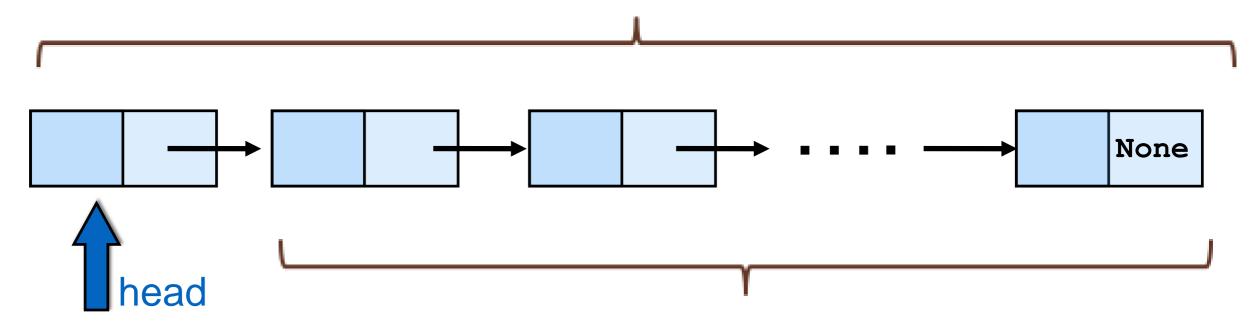




```
def copy(self):
    new_list = List()
    for item in self:
        new_list.insert(len(new_list), item)
    return new_list
```

## What about recursively?

#### \_copy(self.head, new\_list)



\_copy(self.head.next, new\_list)

new\_list.insert(0, head.item)

#### copy

```
Auxiliary method so
                                          up the initial parameter
def copy(self):
    new list = List()
    self._copy_aux(self.head, new_list)
     return new_list
                                            Convergence
def _copy_aux(self, node, new_tist):
    if node is not None:
         self._copy_aux(node.next, new_list)
         new_list.insert(0, node.item)
                                          Combination
    Base case
```

```
def copy(self):
                                 new_list = List()
                                 self._copy_aux(self.head, new_list)
                                 return new_list
                              def _copy_aux(self, node, new_list):
                                 if node is not None:
                                     self._copy_aux(node.next, new_list)
                                     new_list.insert(0, node.item)
                                head
 copy_aux
                  0(1)
                  0(1)
 copy_aux
                                                                None
_copy_aux O(1)
                                                                None
                                                                             n times
                                                                None
copy_aux
```

### Using iterators...

```
def copy(a_list):
        new_list = List()
        copy_aux(iter(a_list), new_list)
        return new_list
def copy_aux(iter, new_list):
    try:
        item = next(iter)
        copy_aux(iter, new_list)
        new_list.insert(0, item)
    except StopIteration:
        pass
```

#### copy

```
def copy(self):
    new_list = List()
    for item in self:
        new_list.insert(len(new_list), item)
    return new_list
```

```
O(n^2)
```

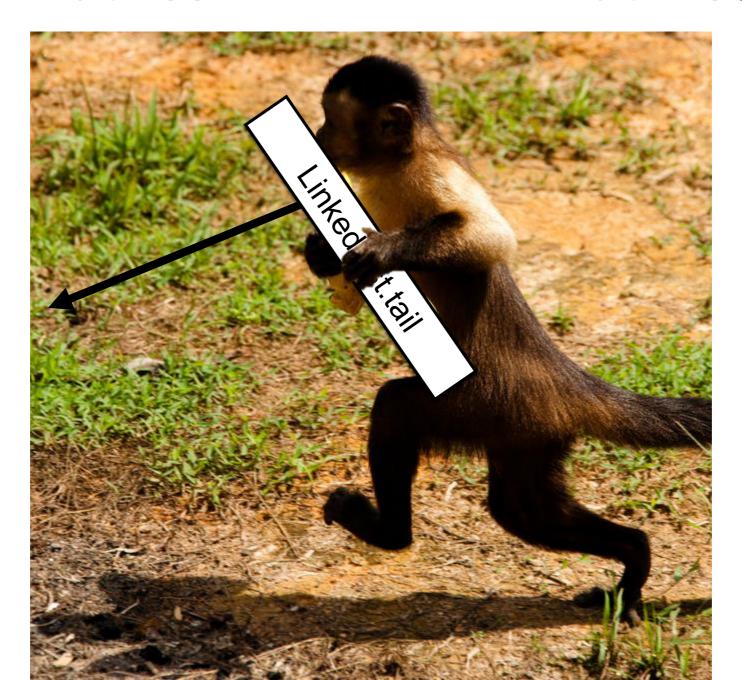
```
def copy(self):
    new_list = List()
    self._copy_aux(self.head, new_list)
    return new_list

def _copy_aux(self, node, new_list):
    if node is not None:
        self._copy_aux(node.next, new_list)
        new_list.insert(0, node.item)
```

**O(n)** 

### Or...

Have a tail attribute and you can have O(1) append and hence O(n) copy



## Advantages/Disadvantages of Recursion

- Advantages:
  - More natural
  - Easier to prove correct
  - Easier to analyse
- Disadvantages:
  - Run-time overhead depending on the quality of the compiler
  - Memory overhead (fewer local variables versus stack space for function call)



what's easy to prove

Just keep doing 'you', and being who
you are and doing what feels natural
to you.

balancing your time and memory use

— Taylor Swift —

AZ QUOTES

## Summary

- Recursive algorithms are characterised by:
  - 1. Existence of base cases
  - 2. Decomposition into simpler subproblems
  - 3. Combination of solutions to subproblems
- Recursive methods require:
  - One or more base cases
  - One or more recursive calls
  - 3. Convergence in the recursive calls
  - 4. Combination of sub-solutions