

Lecture 36

Heaps and Heap sort

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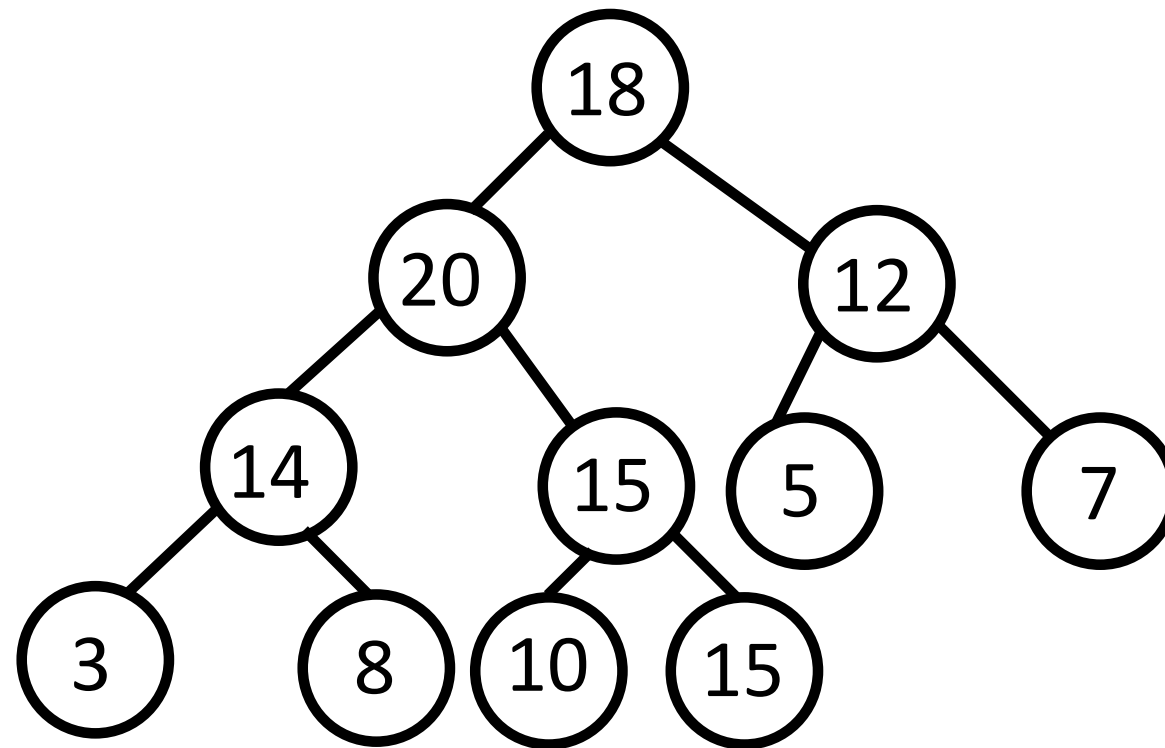
Operations

add:

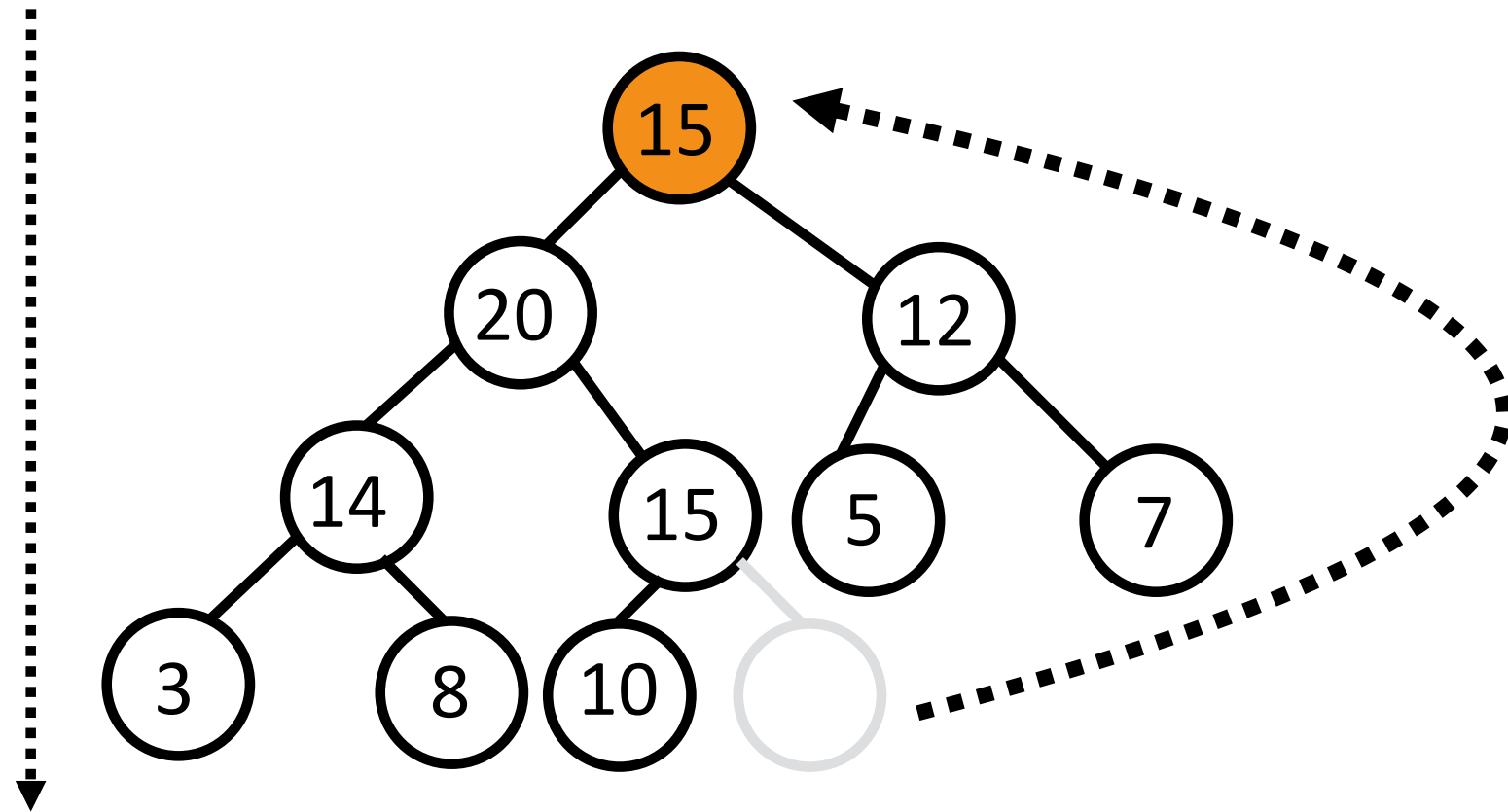
- put at the bottom
- while order is broken, rise.

get_max:

- swap root with last item
- remove last item
- while order is broken, sink.

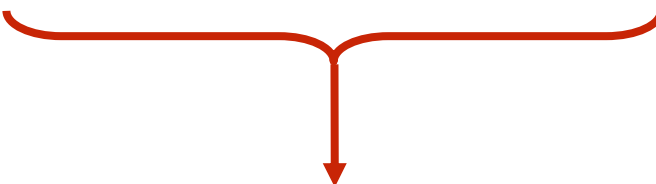


sink
(swapping with
largest child)



A concrete implementation

```
def largest_child(self, k):  
    if self.array[2*k] > self.array[2*k+1]:  
        return 2*k  
    else:  
        return 2*k+1
```



what if k only has one child

Make the item at index k sink to the correct position.

```
def sink(self, k):  
    while 2*k <= self.count:  
        child = self.largest_child(k)  
        if self.array[k] >= self.array[child]:  
            break  
        self.swap(child, k)  
        k = child
```

left child in last position means
k has only one child



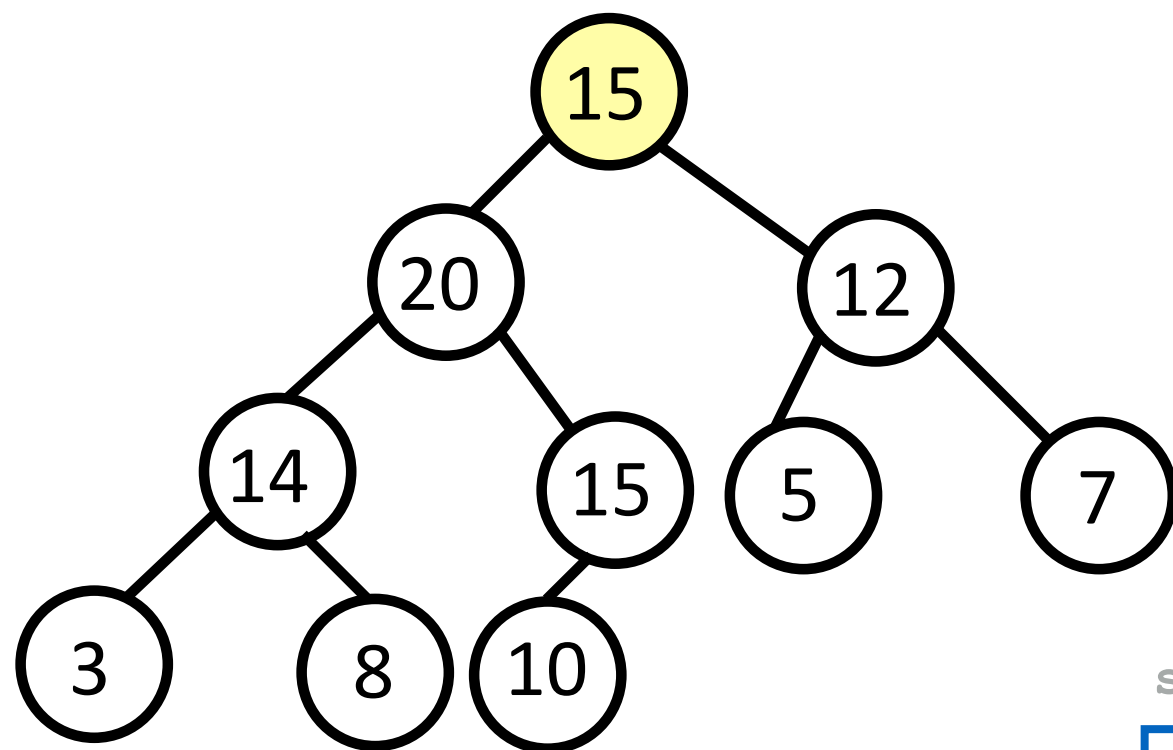
```
def largest_child(self, k):  
    # Check for only one child.  
    if 2*k == self.count or self.array[2*k] > self.array[2*k+1]:  
        return 2*k  
    else:  
        return 2*k+1
```

Make the item at index k sink to the correct position.

```
def sink(self, k):  
    while 2*k <= self.count:  
        child = self.largest_child(k)  
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        self.swap(child, k)  
        k = child
```

On subtle errors

- Errors like that are very easy to make, and hard to spot.
- Your armoury against them includes:
 - Thorough testing (recall us asking for test cases?)
 - Code review
 - Proofs of correctness



k=1

self.count = 10

self.array

	15	20	12	14	15	5	7	3	8	10
0	1	2	3	4	5	6	7	8	9	10

```

def largest_child(self, k):
    # Check for only one child.
    if 2*k == self.count or self.array[2*k] > self.array[2*k+1]:
        return 2*k
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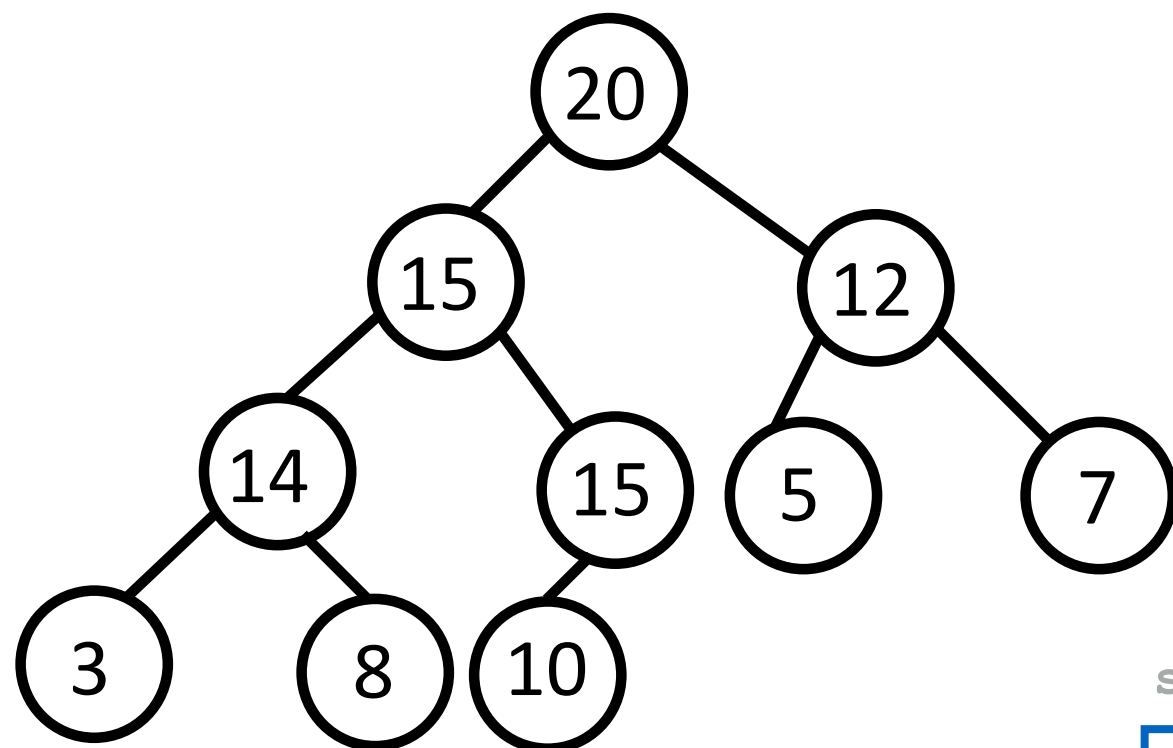
```

Make the item at index k sink to the correct position.

```

def sink(self, k):
    while 2*k <= self.count:
        child = self.largest_child(k)
        if self.array[k] >= self.array[child]:
            break
        self.swap(child, k)
        k = child

```

`child = 5`

`k = 2`

`self.count = 10`

`self.array`

	20	15	12	14	15	5	7	3	8	10
0	1	2	3	4	5	6	7	8	9	10

```

def largest_child(self, k):
    # Check for only one child.
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        k = child
  
```

best case: $O(1)$

worst case: $O(\log N)$

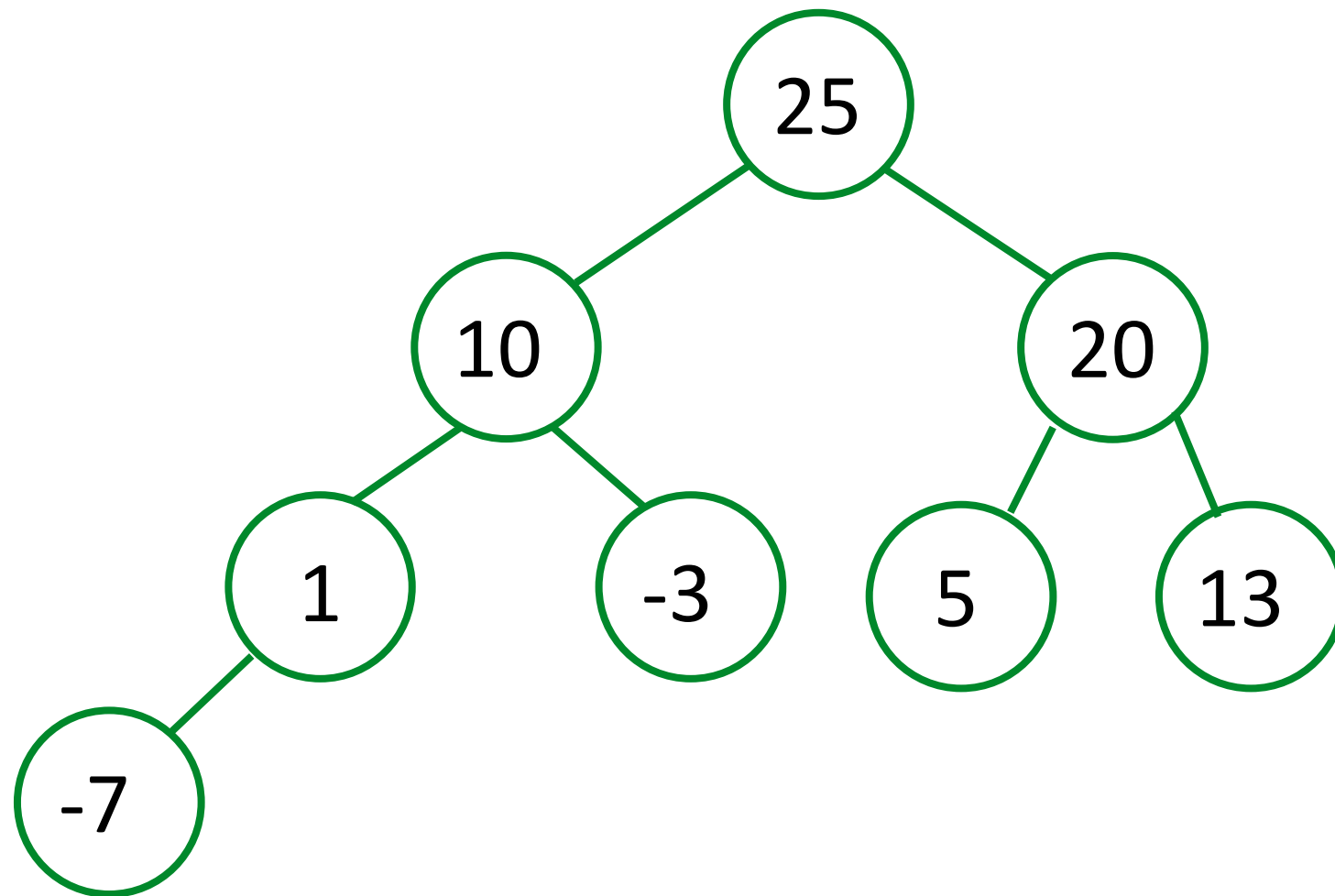
(may need to consider
comparison operations)

Complexity of get_max

- Loop in sink can iterate at most depth times $\approx \log(N)$
(after depth iterations, the new item is at the root)
- **Best case:** $O(1) * O_{\text{Compare}}$ when the item is larger or equal than largest children.
- **Worst case:** $O(\log N) * O_{\text{Compare}}$ when the item sinks all the way to the bottom.

Heap sort

[5, -7, 10, -3, 13, 20, 25, 1]



	25	10	20	1	-3	5	13	-7
0	1	2	3	4	5	6	7	8

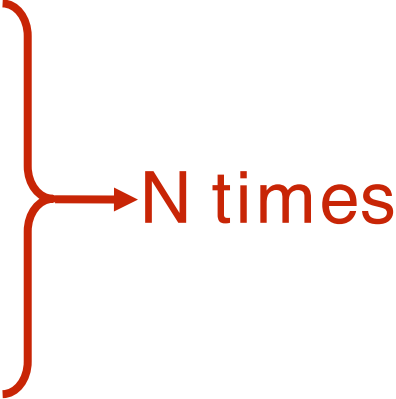
[5, -7, 10, -3, 13, 20, 25, 1]

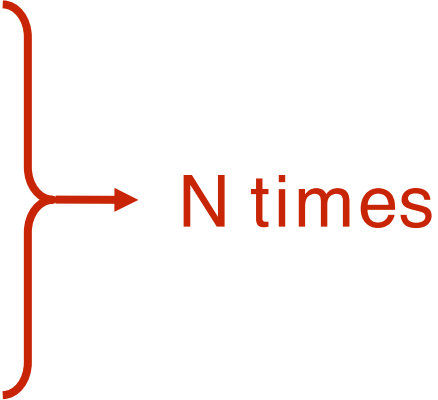
	-7	-3	1	5	10	13	20	25
0	1	2	3	4	5	6	7	8

These numbers are sorted!



Heap sort

- For each element in the array:
 - Add it to the Heap: $O(\log(N))$

N times
- While heap contains elements:
 - Get max item: $O(\log(N))$
 - Put it in the “hole” made by previous step: $O(1)$

N times

worst case: $O(N \log N)$

Summary

- A simple Heap implementation
 - rise
 - sink
 - largest_child
- Heap Sort