

# Lecture 6

# Memory in MIPS

FIT 1008  
Introduction to Computer Science



COMMONWEALTH OF AUSTRALIA

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# Objectives

- The need for **memory diagrams** and how to draw them
- How the system stack works and the role played by **\$sp** and **\$fp**
- How (and why) **local variables** are stored on the stack and how to access them
- How to use **addressing modes** to access variables

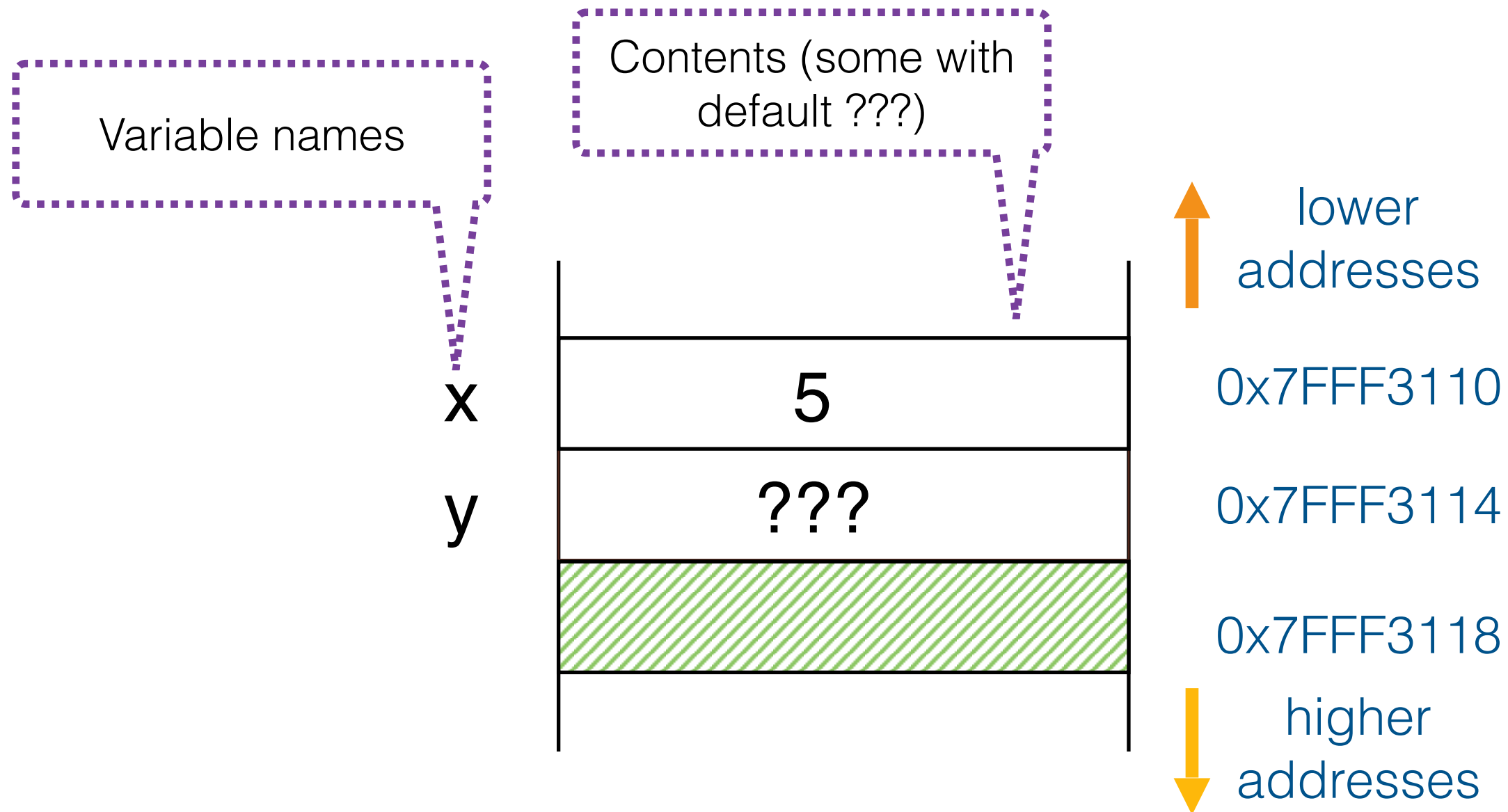
# What we have seen

- How to define and use global variables
- How to allocate memory on the Heap.
- How to use memory on the Heap.

# Memory diagrams

- Useful **for humans** to know how to access variables
- Show memory allocated to variables:
  - ➔ Addresses
  - ➔ Contents
  - ➔ Variable names

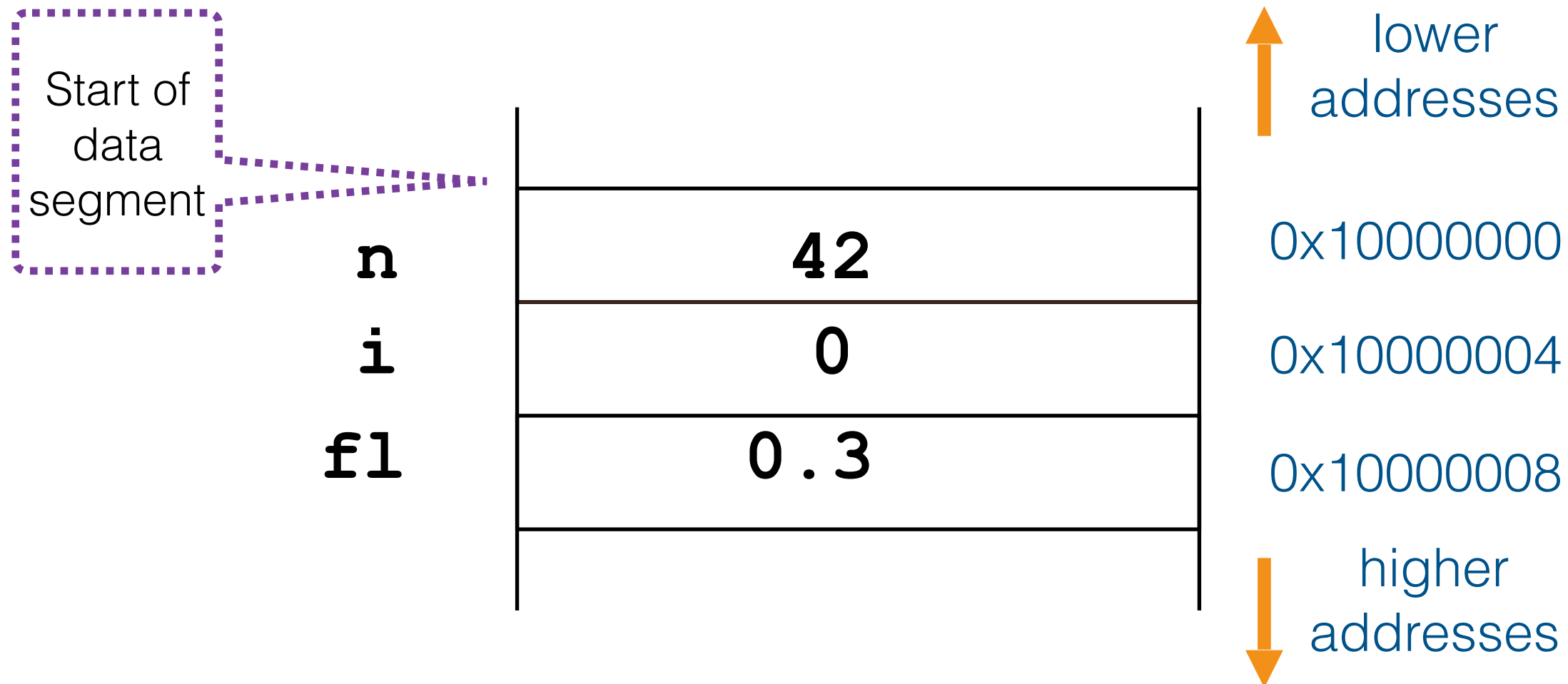
**Recall:** we assume numbers appear directly at the memory location (not true in Python, but true in C or Java) and occupy 4 bytes.



When variables contain addresses of other variables, helpful to draw arrow (pointer)

# Global variables

- Memory map **not** crucial for **global variables** (stored in **data segment**)
- Global variables: every variable has a **label** to identify it



# Local variables

- Why not store local variables in the data segment? Local variables do not have labels.
- **Properties of Data segment**
  - ➔ Accessible from all parts of the program
  - ➔ Labels must be **unique**
  - ➔ Each location can hold only one discrete value
- **Properties of Local Variables**
  - ➔ Accessible only within a function.
  - ➔ Several variables with same name (different scopes) within the same function
  - ➔ May have more than one version of the same function's variables existing (due to recursion)

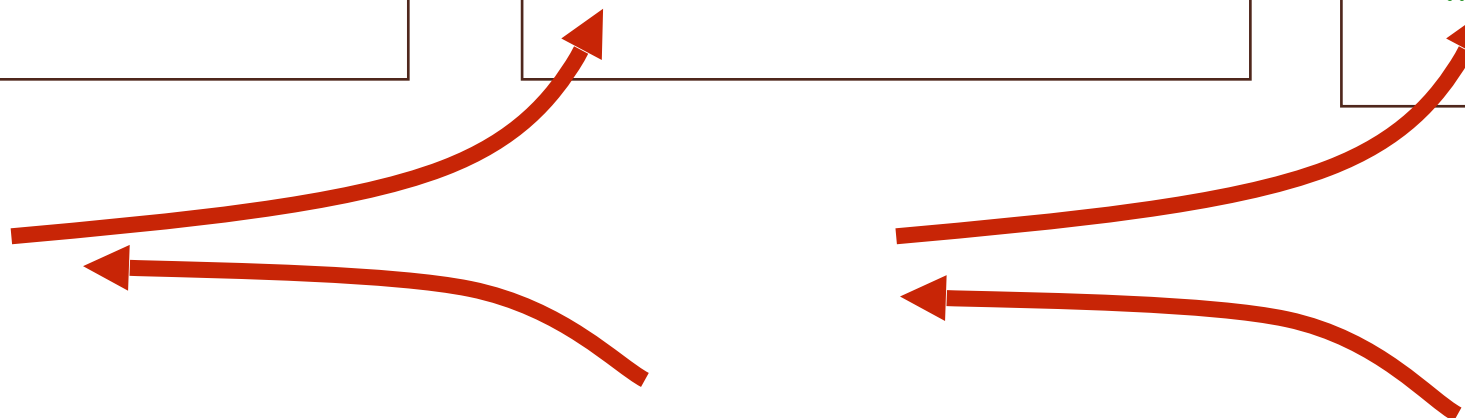
# Properties of local variables

- Must be created/**allocated** at function **entry**
- Must be destroyed/**deallocated** at function **exit**
- Other functions may be called in between, with the same rules

```
def a():  
    # create a_var  
    a_var = 0  
  
    b()  
    # delete a_var
```

```
def b():  
    # create b_var  
    b_var = 0  
  
    c()  
    # delete b_var
```

```
def c():  
    # create c_var  
    c_var = 0  
  
    ...  
    # delete c_var
```





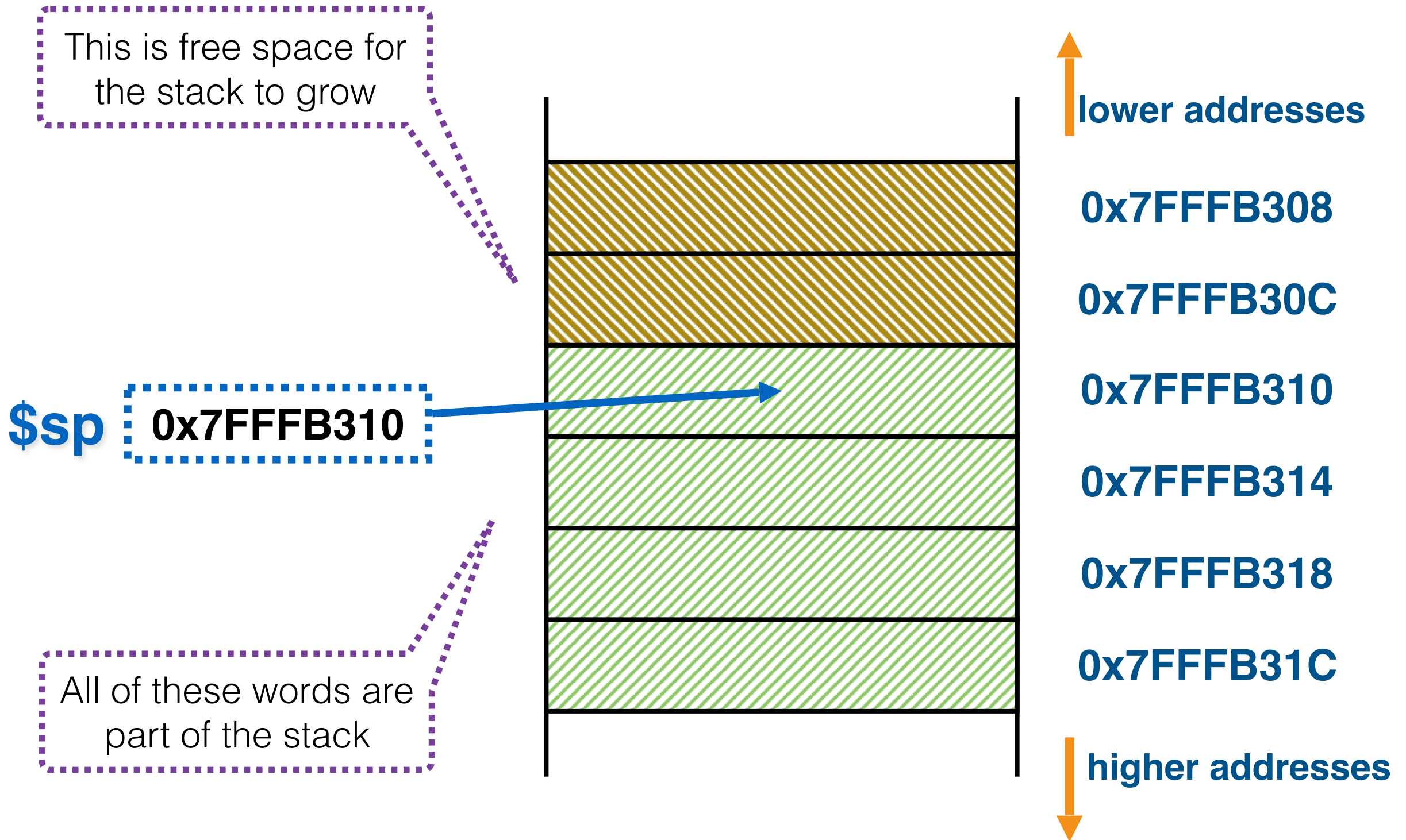
# Properties of local variables

- Allocation/deallocation is LIFO:  
**Last allocated, first deallocated**
- A stack data structure is ideal for storing local variables
  - **Allocate** = **push**
  - **Deallocate** = **pop**
- Most computers provide a memory stack for programs to use (initialised by OS): **system stack** or **runtime stack** or **process stack**
- The instruction set provides operations for pushing/popping off the system stack.

# System Segment

- Register **\$sp** (stack pointer) indicates the top of stack
  - ➔ Contains the address of the word of memory at the top of stack (i.e., with lowest address)
  - ➔ Its value changes during the execution of a function
- How do we **push** and **pop** variables?

# System stack



# System stack: **pushing**

This is the new word we want to push onto the stack

**\$sp** **0x7FFFB310**

**subtract 4 from  
stack pointer**



**lower addresses**

0x7FFFB308

0x7FFFB30C

0x7FFFB310

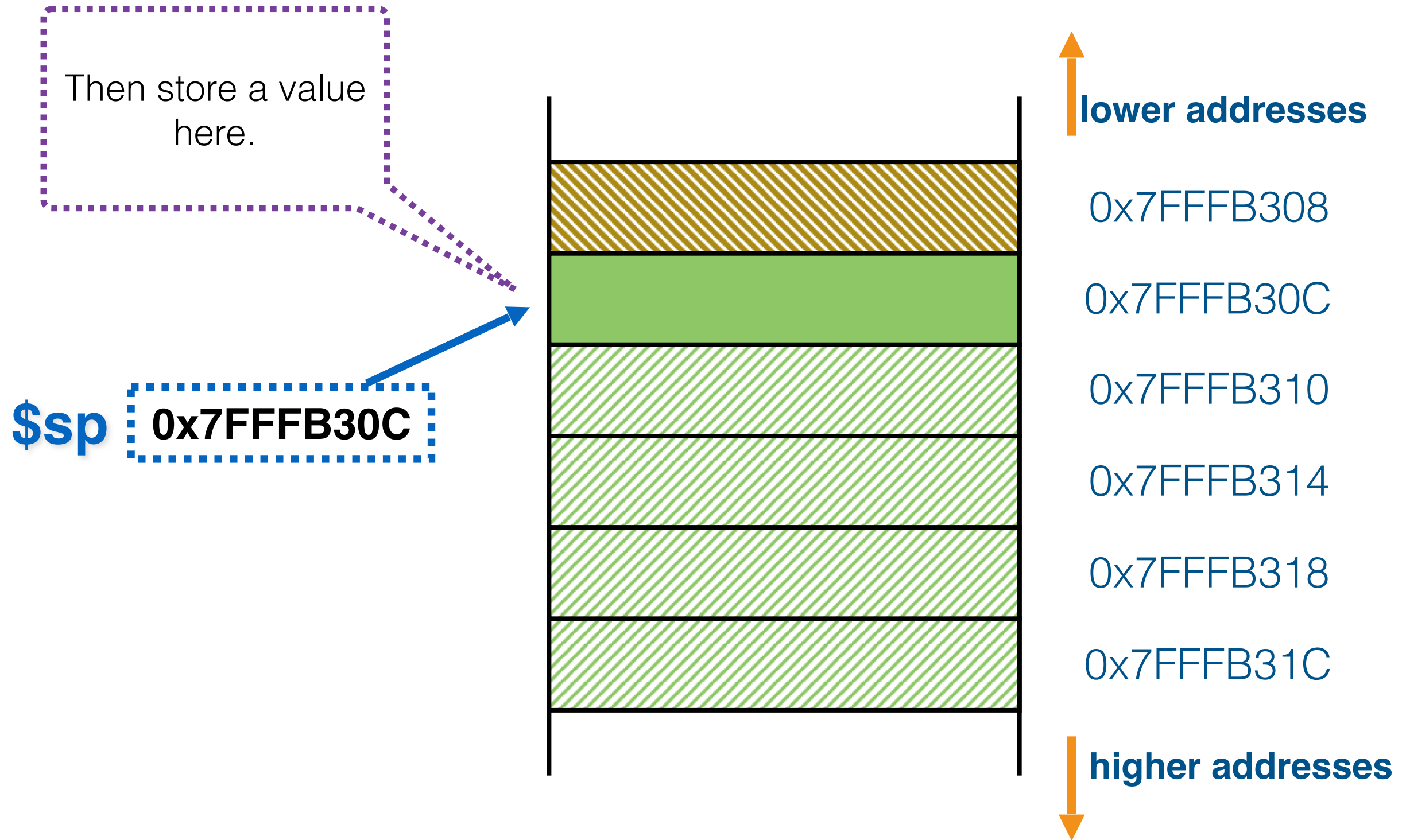
0x7FFFB314

0x7FFFB318

0x7FFFB31C

**higher addresses**

# System stack: **pushing**

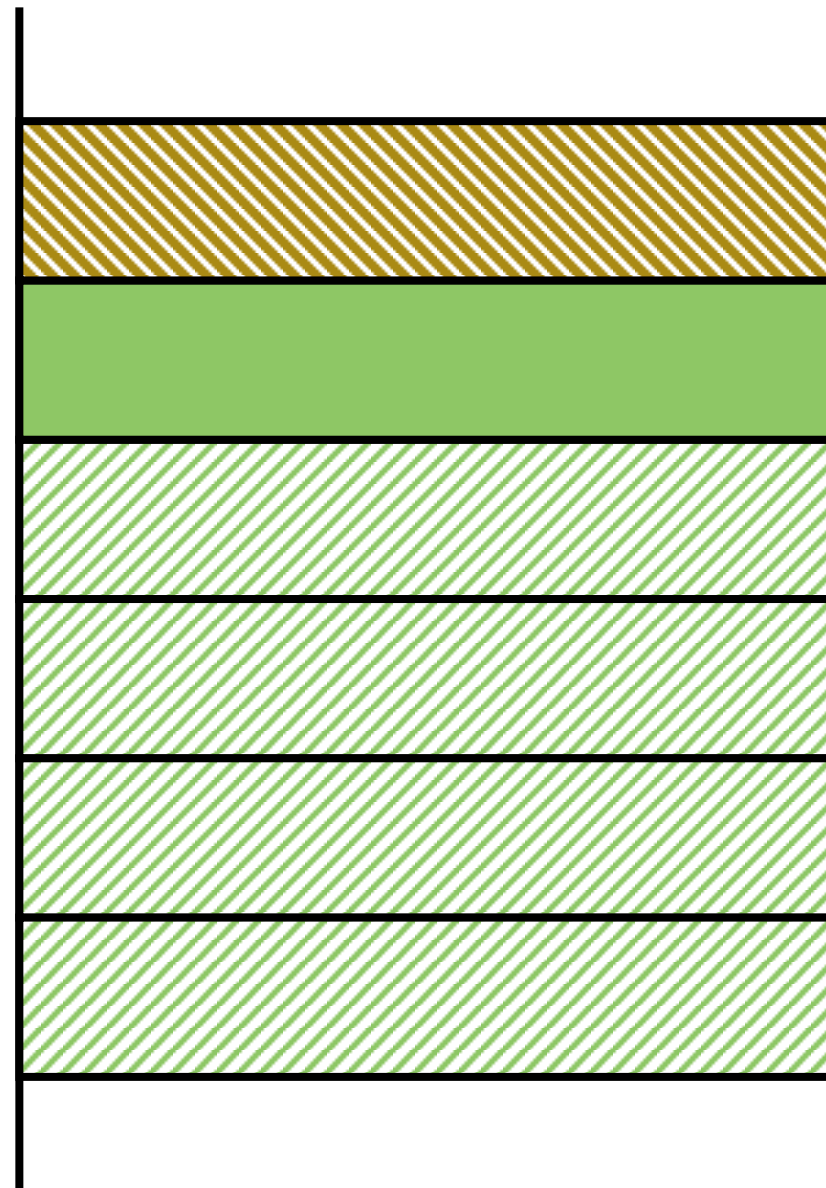


# System stack: **popping**

To **pop** this word:  
Fetch it  
into a register

**\$sp** 0x7FFFB30C

**then add 4 to the  
stack pointer**



↑ **lower addresses**

0x7FFFB308

0x7FFFB30C

0x7FFFB310

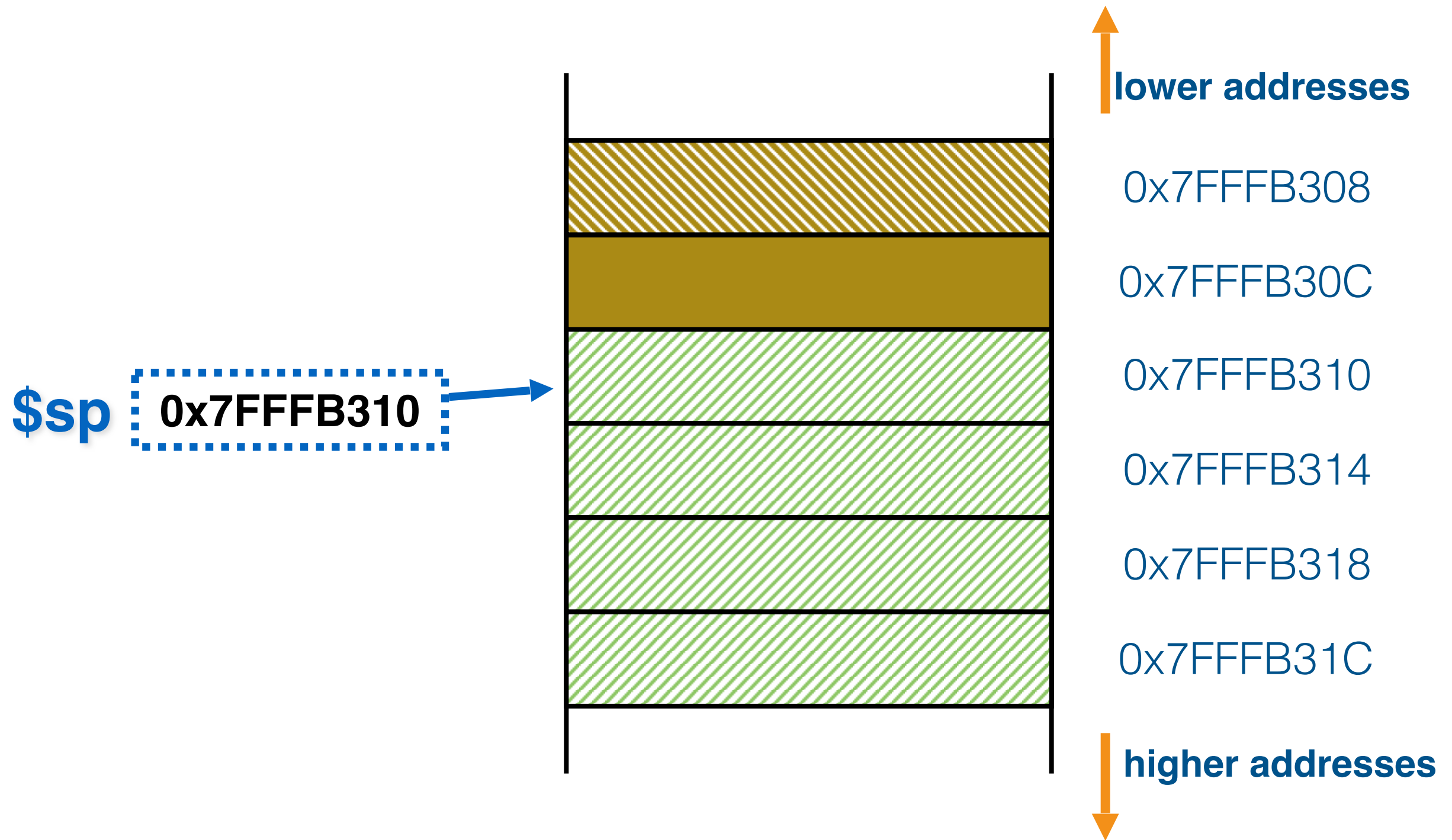
0x7FFFB314

0x7FFFB318

0x7FFFB31C

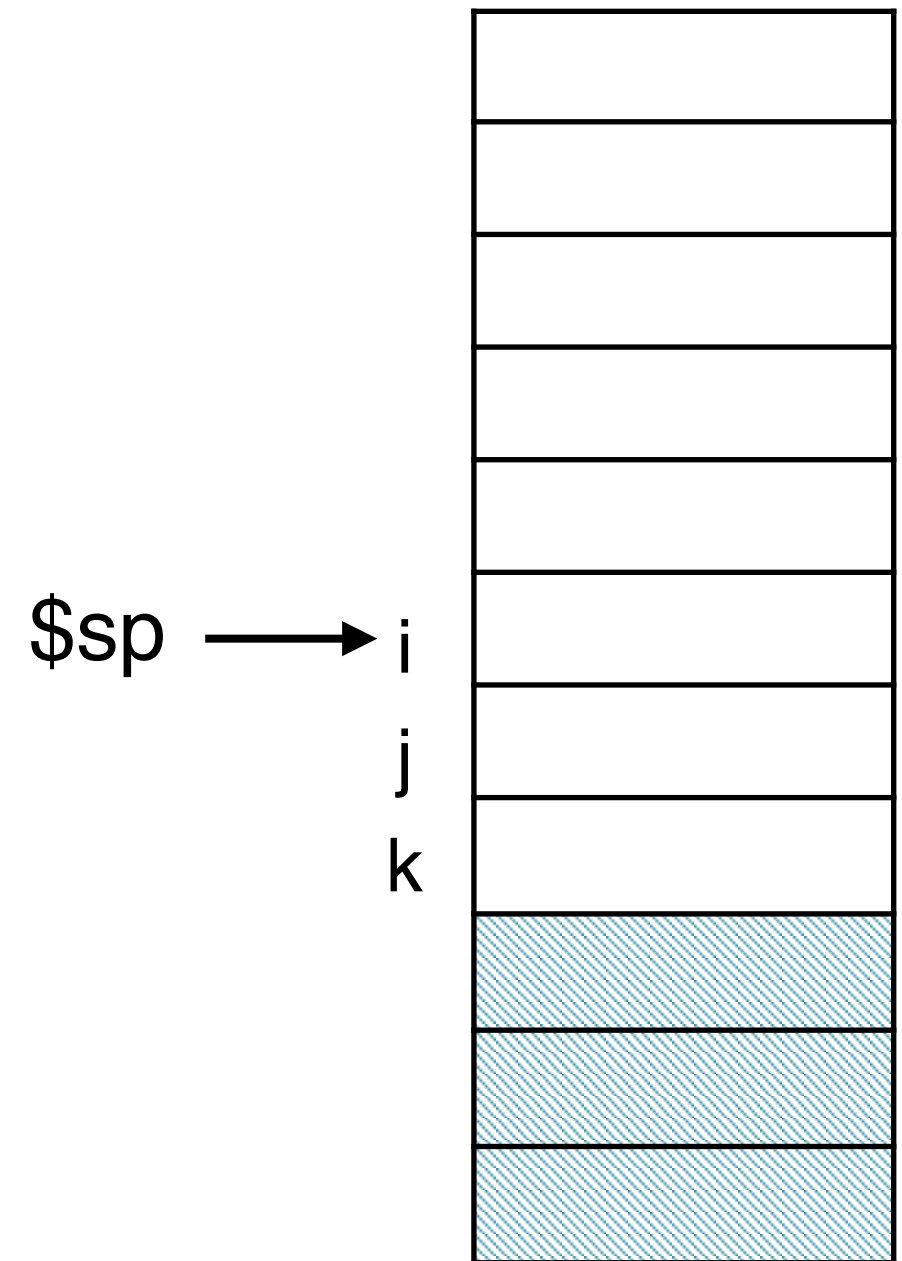
↓ **higher addresses**

# System stack: **popping**



# How does the system stack work?

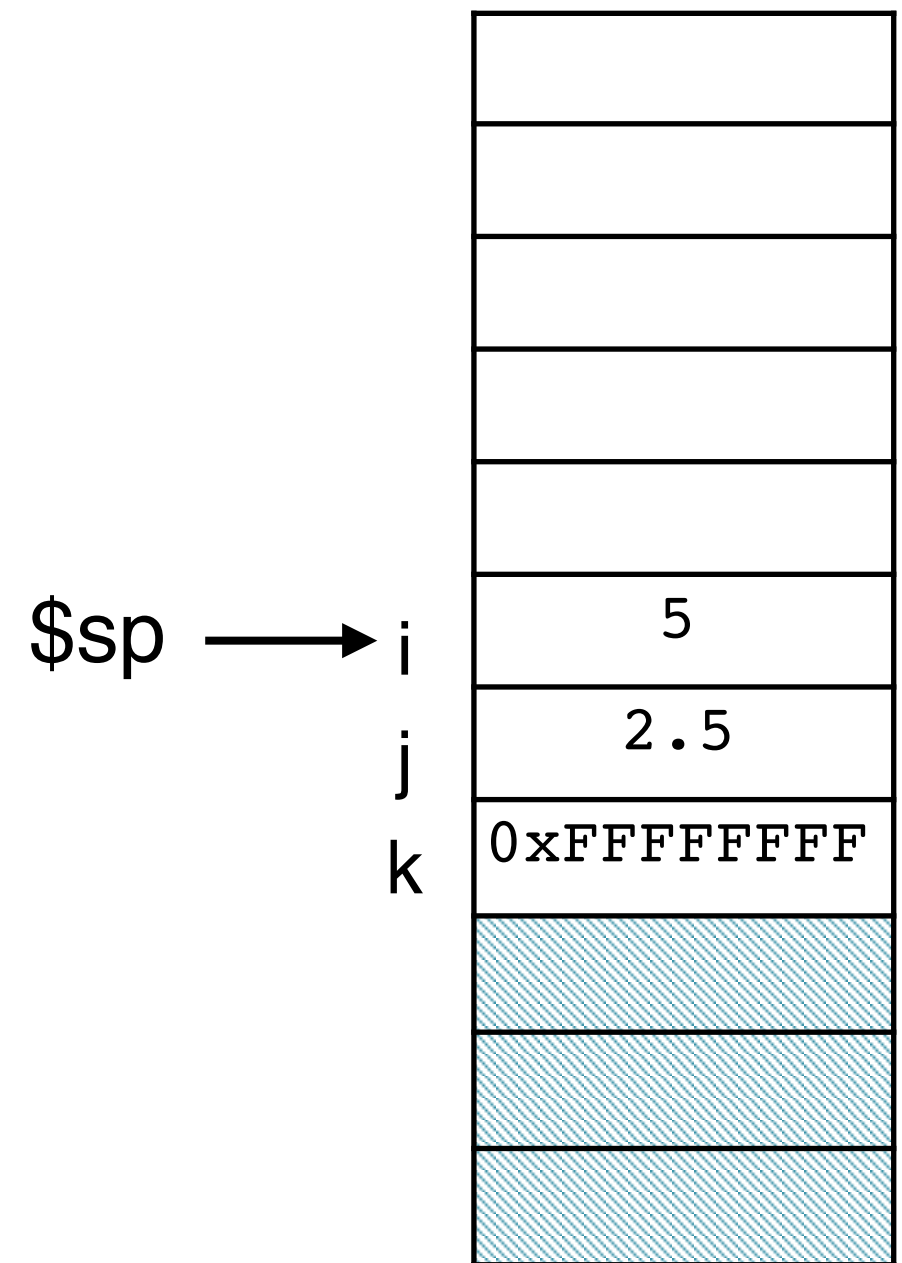
- At the beginning of a function
  - ➔ **Allocate** variables by **pushing** necessary space onto stack (subtract n bytes from **\$sp**)
  - ➔ Initialise space by storing values in newly allocated space





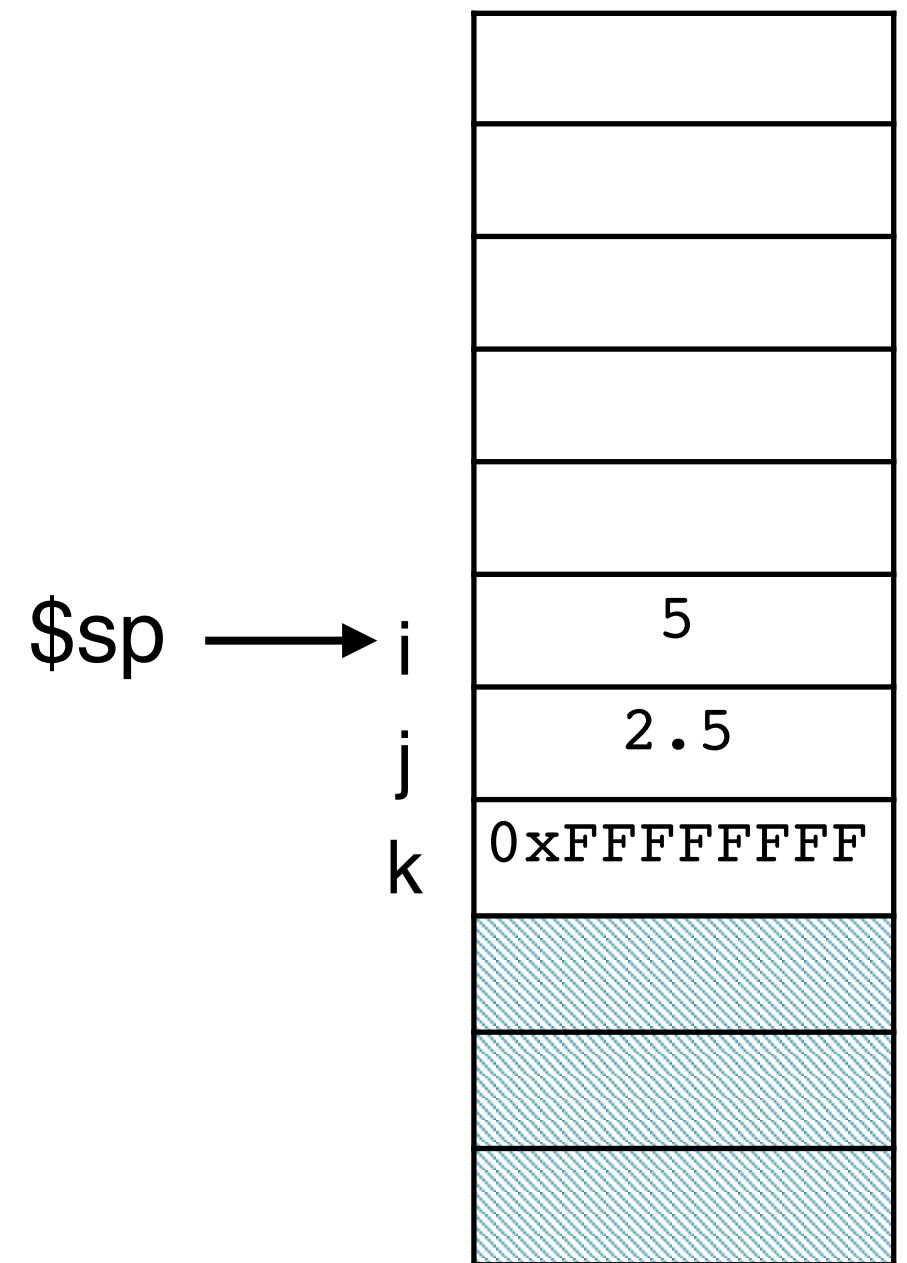
# How does the system stack work?

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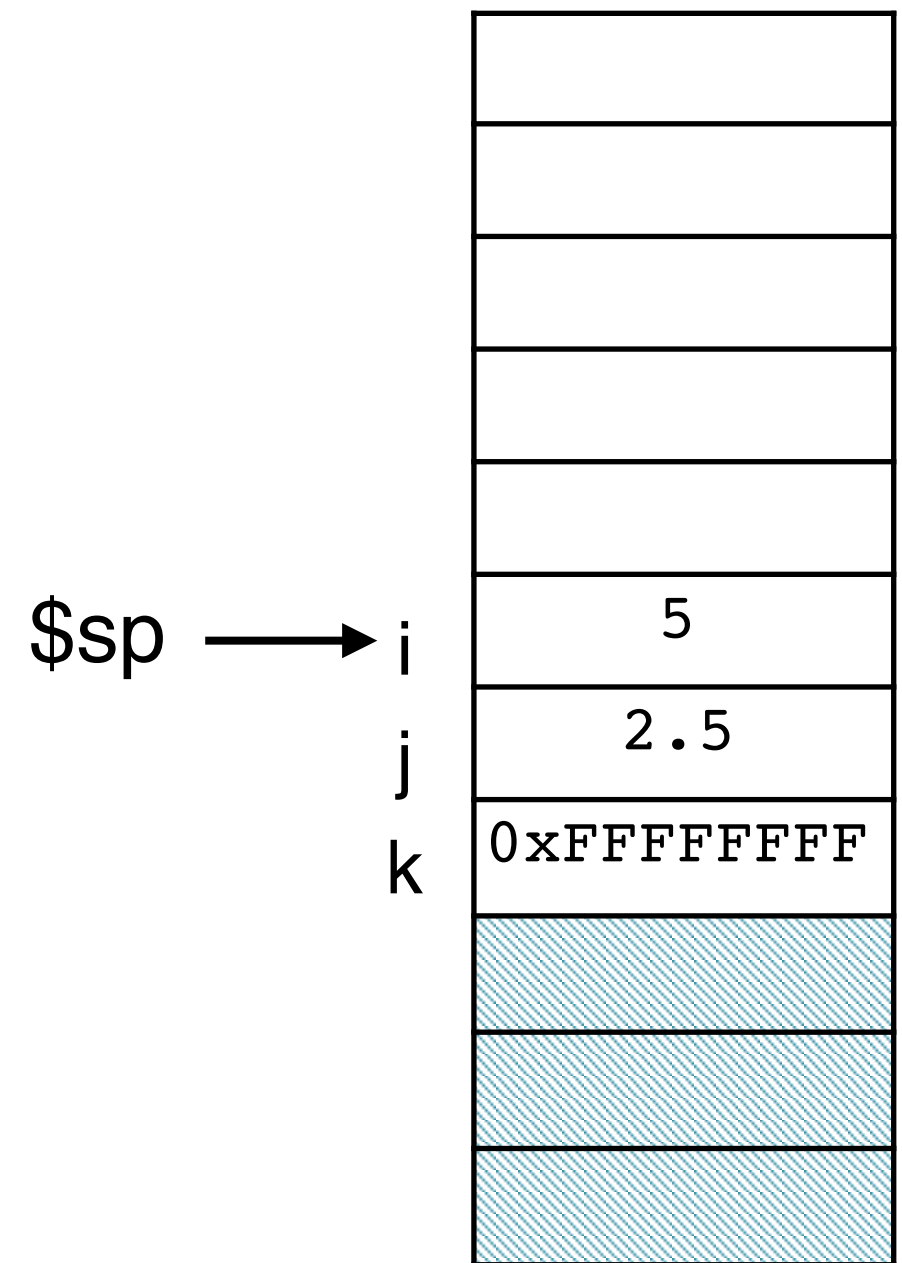
# How does the system stack work?

- At the beginning of a function
  - ➔ **Allocate** variables by **pushing** necessary space onto stack (subtract n bytes from **\$sp**)
  - ➔ Initialise space by storing values in newly allocated space
- During function:  
use variables using **lw/sw**



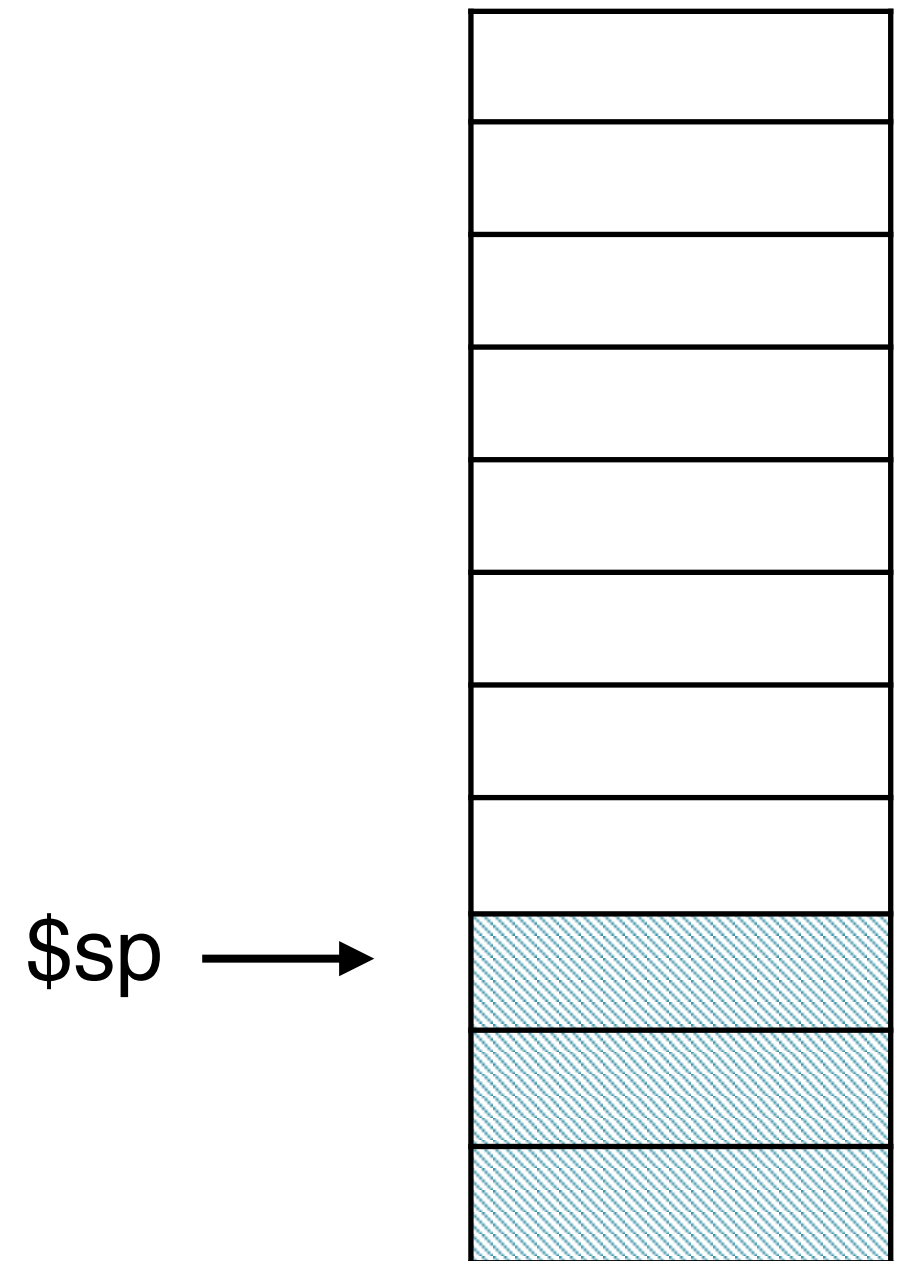
# How does the system stack work?

- At the beginning of a function
  - ➔ **Allocate** variables by **pushing** necessary space onto stack (subtract n bytes from **\$sp**)
  - ➔ Initialise space by storing values in newly allocated space
- During function:  
use variables using **lw/sw**
- At the end of the function:  
**Deallocate** variables by **popping** allocated space from stack (add n bytes to **\$sp**)



# How does the system stack work?

- At the beginning of a function
  - ➔ **Allocate** variables by **pushing** necessary space onto stack (subtract n bytes from **\$sp**)
  - ➔ Initialise space by storing values in newly allocated space
- During function:  
use variables using **lw/sw**
- At the end of the function:  
**Deallocate** variables by **popping** allocated space from stack (add n bytes to **\$sp**)



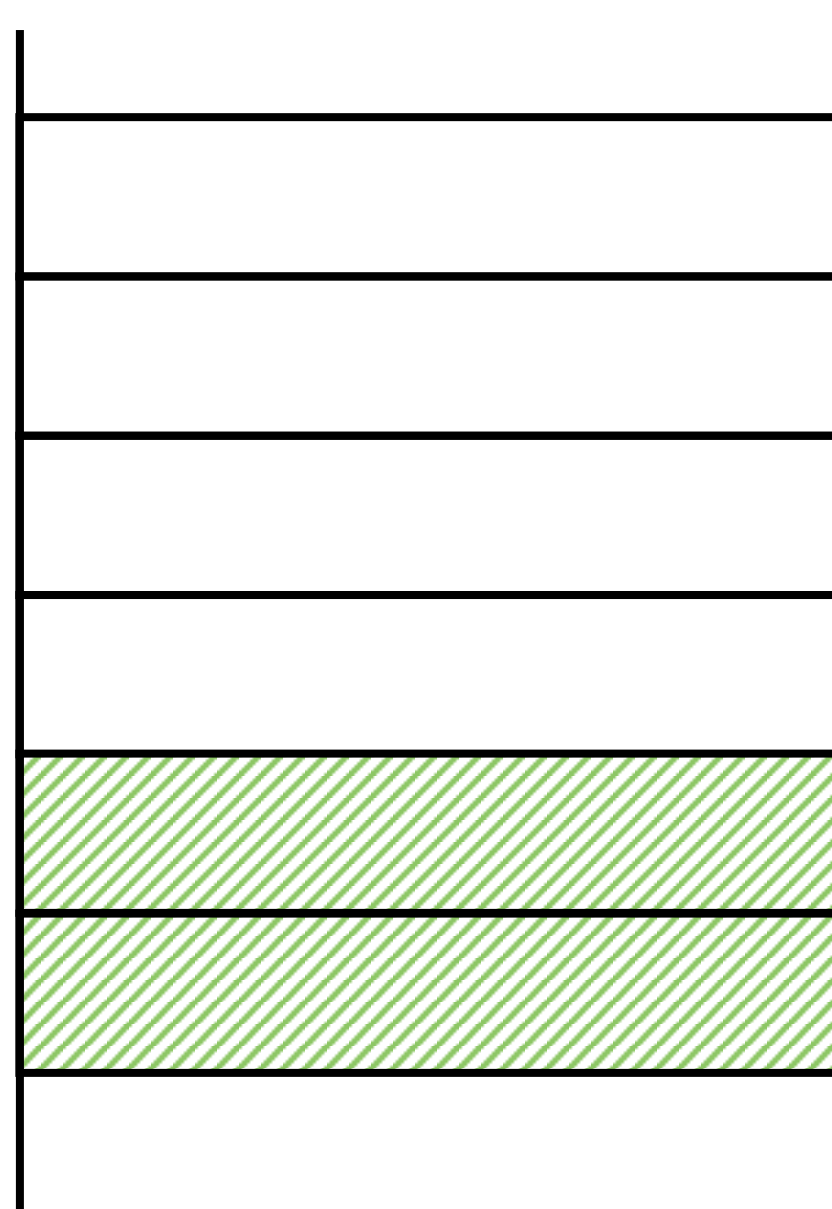
# Example

```
def a():  
    x = 5  
    y = 10  
    ...
```

**\$sp**

**0x7FFFB3118**

At the beginning of  
the function there  
may be data on the  
stack already



**lower addresses**

0x7FFF310C

0x7FFF3110

0x7FFF3114

0x7FFF3118

0x7FFF311C

**higher addresses**

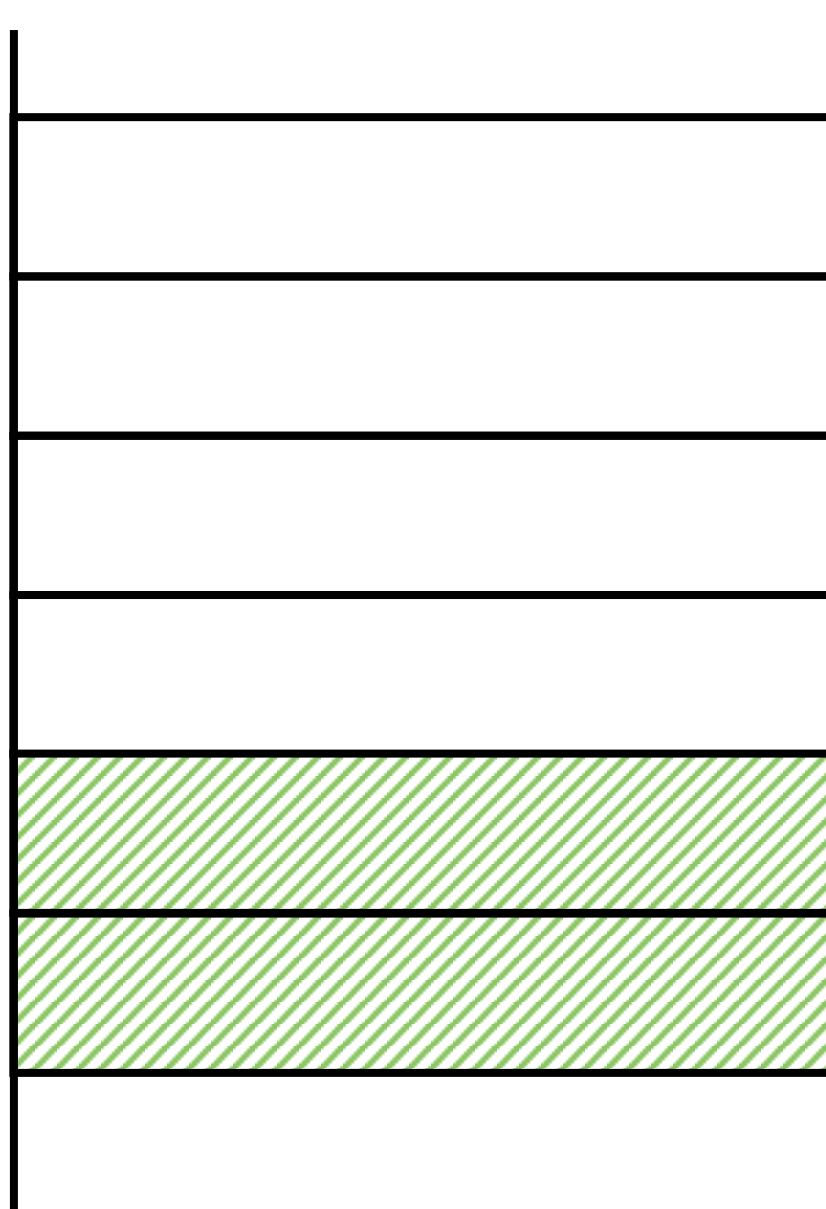
# Example

```
def a():  
    x = 5  
    y = 10  
    ...
```

**\$sp**

**0x7FFFB3118**

Allocate space  
4 bytes for x  
4 bytes for y  
\$sp = \$sp - 8



**lower addresses**

0x7FFF310C

0x7FFF3110

0x7FFF3114

0x7FFF3118

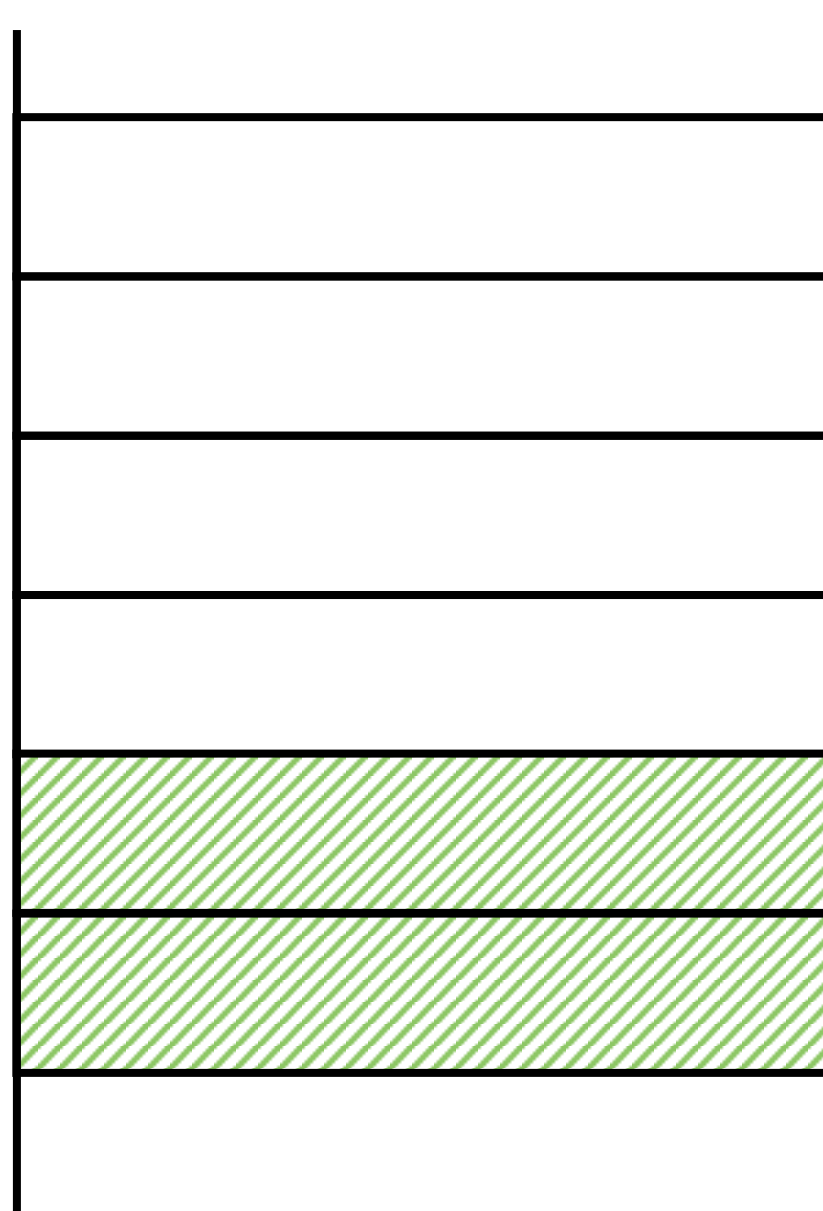
0x7FFF311C

**higher addresses**

# Example

```
def a():  
    x = 5  
    y = 10  
    ...
```

**\$sp** **0x7FFFB3110** →



↑  
**lower addresses**

0x7FFF310C

0x7FFF3110

0x7FFF3114

0x7FFF3118

0x7FFF311C

↓  
**higher addresses**

# Example

```
def a():  
    x = 5  
    y = 10  
    ...
```

**\$sp** **0x7FFFB3110** →

x

y

lower addresses

0x7FFF310C

0x7FFF3110

0x7FFF3114

0x7FFF3118

0x7FFF311C

higher addresses

Assign variables



# Example

```
def a():  
    x = 5  
    y = 10  
    ...
```

\$sp

0x7FFFB3110



x

5

y

10

Store initial values

lower addresses

0x7FFF310C

0x7FFF3110

0x7FFF3114

0x7FFF3118

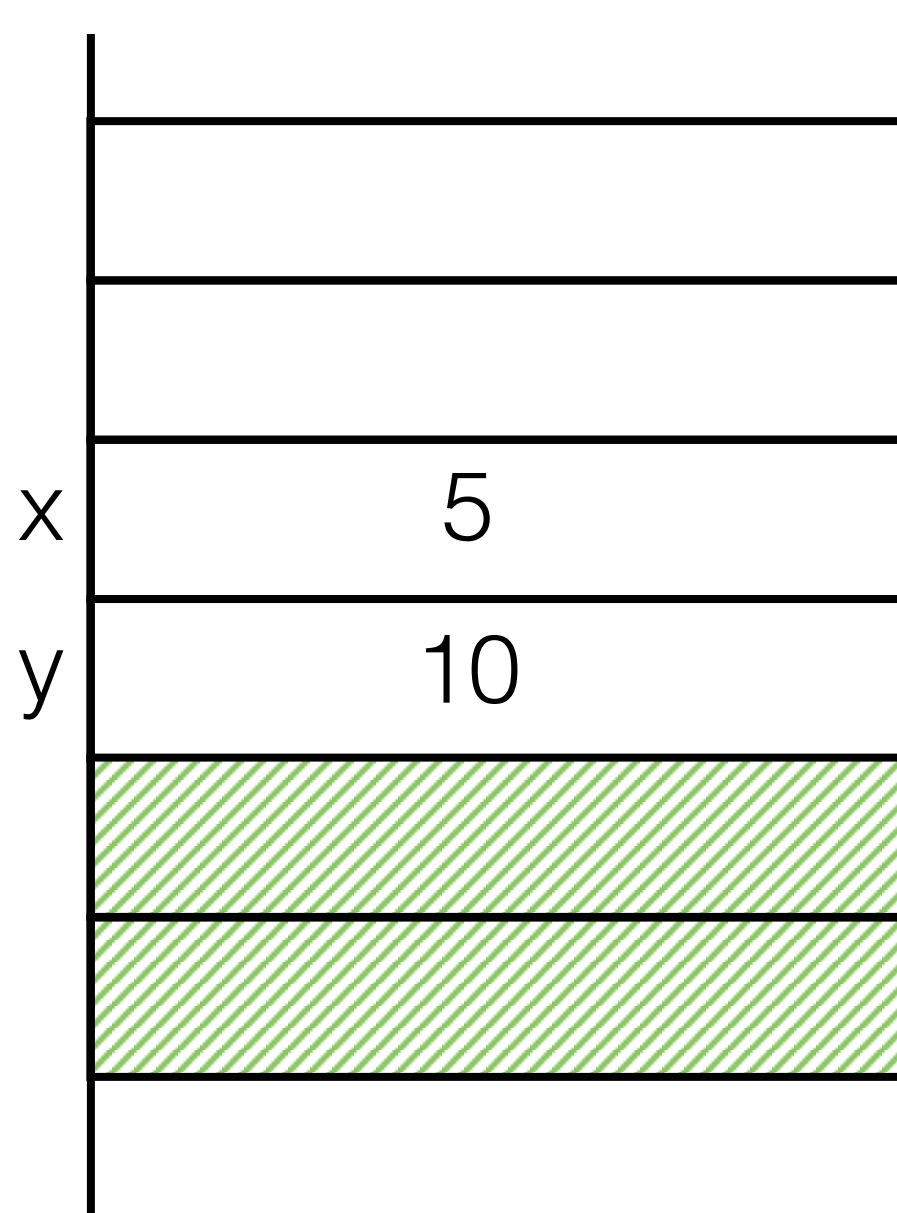
0x7FFF311C

higher addresses

# Example

```
def a():  
    x = 5  
    y = 10  
    ...
```

**\$sp** **0x7FFFB3110** →



lower addresses

0x7FFF310C

0x7FFF3110

0x7FFF3114

0x7FFF3118

0x7FFF311C

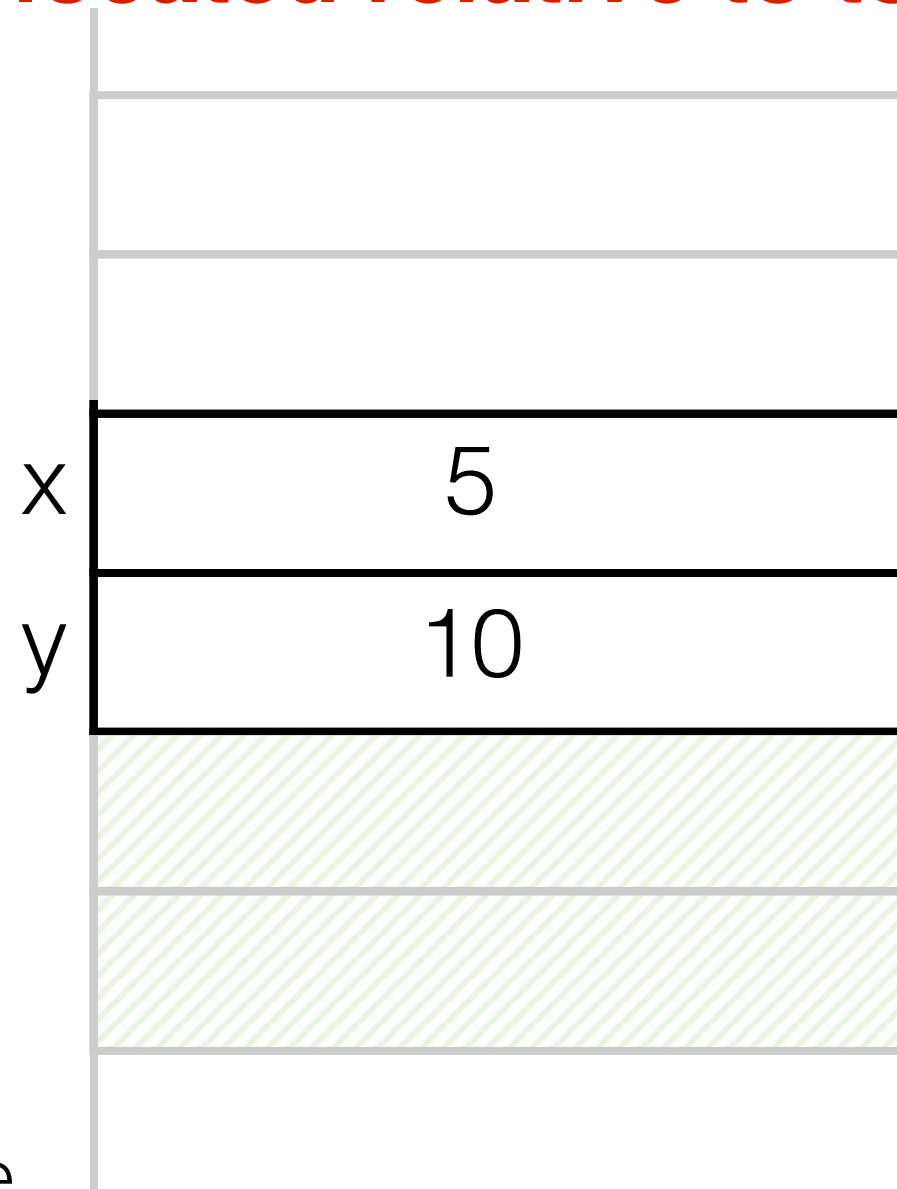
higher addresses

# How do we use these values or refer to them?

I can use `$sp` since variables are located relative to top of stack

```
def a():  
    x = 5  
    y = 10  
    ...
```

`$sp` `0x7FFFB3110` →



lower addresses

0x7FFF310C

0x7FFF3110

0x7FFF3114

0x7FFF3118

0x7FFF311C

higher addresses

**Labels?** Not possible.  
Labels are compile time,  
not run-time

**Addresses?** Not possible.  
stack may have a  
different depth every time

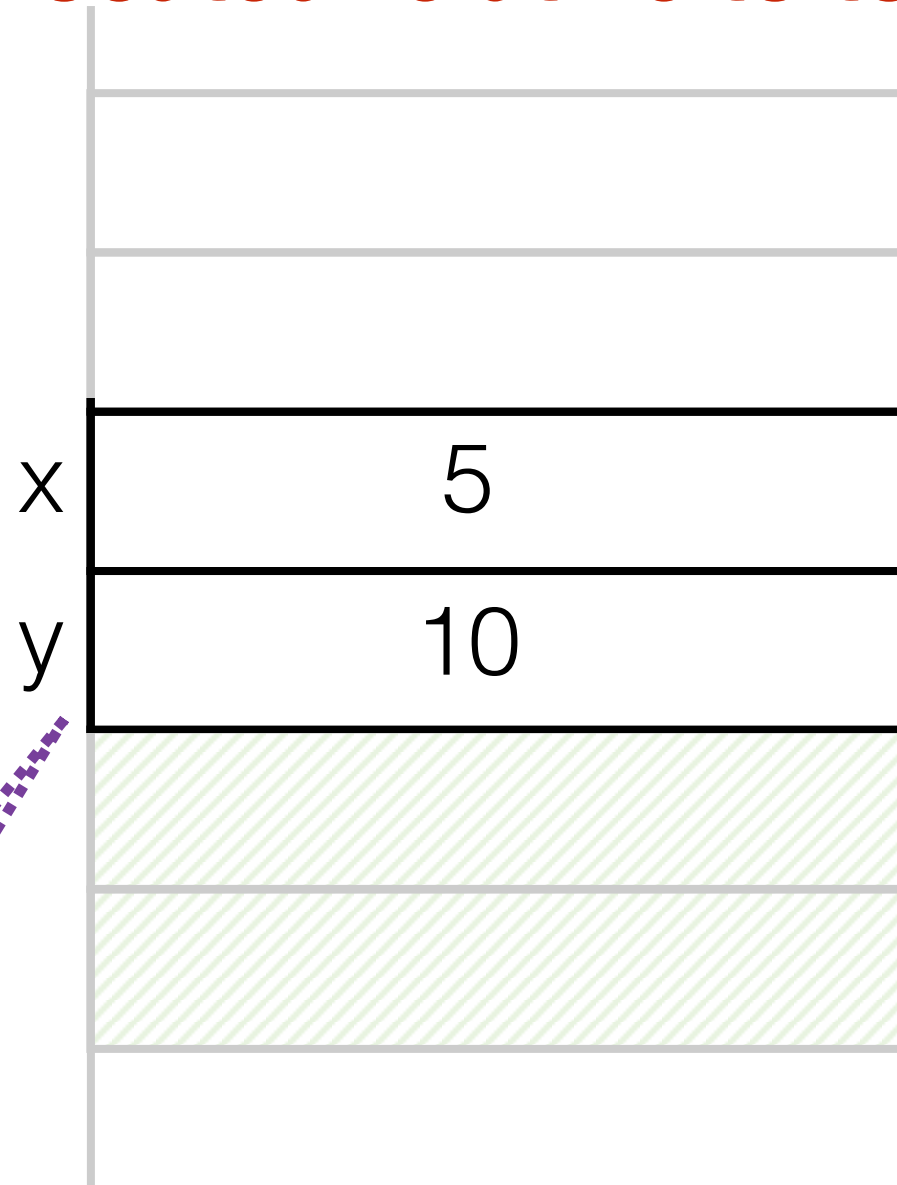
I can use `$sp` since variables are located relative to top of stack

```
def a():  
    x = 5  
    y = 10  
    ...
```

`$sp` **0x7FFFB3110** →

Store `x = 5` at  
address **`$sp + 0`**  
(0x7FFFB3110)

Store `y = 10` at  
address **`$sp + 4`**  
(0x7FFFB3114)



lower addresses

0x7FFF310C

0x7FFF3110

0x7FFF3114

0x7FFF3118

0x7FFF311C

higher addresses

# Reminder: addressing modes

**const** may be a label,  
signed number or  
expression at run time

**\$reg** is any GPR

**sw \$src, **const**(\$reg)**

**\$reg + const**  
Add const to the  
value of \$reg

# Examples of addressing modes

sw \$t0, 4(\$sp)

address is ( $\$sp + 4$ )

---

sw \$t0, -4(\$fp)

address is ( $\$fp - 4$ )

---

lw \$a0, 0(\$sp)

lw \$a0, (\$sp)



address is ( $\$sp + 0$ )

---

lw \$a0, var(\$zero)

lw \$a0, var



address is  
( $\$zero + \text{address of var}$ )

# Summary

- Memory diagrams.
- System stack:
  - Pushing and popping
  - **\$sp**
- Local variables:
  - Stored on stack
- Addressing: register + constant