Is It A Red-Black Tree

Zhang Youchao 3170100125

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Chapter 1: Introduction

There is a kind of balanced binary search tree named red-bla ck tree in the data structure. It has the following 5 properties:

- (1) Every node is either red or black.
- (2) The root is black.
- (3) Every leaf (NULL) is black.
- (4) If a node is red, then both its children are black.
- (5) For each node, all simple paths from the node to descend ant leaves contain the same number of black nodes.

For each given binary search tree, you are supposed to tell if it is a legal red-black tree. My algorithm is based on the property (2)(4)(5).

According to (3) The root should be black node.

According to (4): If a node is red, then both its children must be black.

According to (5): For each node, all simple paths from the node to descendant leaves contain the same number of black nodes. A red black tree can make the black node balanced, which means the left child a nd the right child must have the same depth(we just care about the black node)

Chapter 2: Algorithm Specification

```
1.the data structure
                                                          Element
/*define the node*/
                                                         Left
                                                                Right
typedef struct Node *tree;//define the tree
typedef tree ptr;
                             //define the pointer
struct Node
{
                                                                             Element
                                           Element
    int Element;
    ptr Left;
                                                                            Left
    ptr Right;
};
```

2. According to (3) The root should be black node.

```
Algorithm1:Check root
 Input: tree T
 Output
          0→wrong 1→right
     :
    1:
          function checkroot( tree T )
    2:
             if T==NULL return 1
    3:
             else if T→element>0 return 1
    4:
             else return 0
    5:
          end
/*According to 3 The root should be black node.*/
int checkroot(tree T)
{
    if (!T)return 1;
    else if (T->Element > 0) return 1;
    else return 0;
}
```

3.According to (4): If a node is red, then both its children must be black.

```
Algorithm2: Check children
Input: tree T
Output
           0 \rightarrow \text{wrong} \quad 1 \rightarrow \text{right}
    :
   1:
           function red_with_2black( tree T )
   2:
              if T!=NULL
   3:
                   if T→element<0
   4:
                        if T→left! =NULL and T→left→element<0
   5:
   6:
                        else if T→right! =NULL and T→right→element<0
   7:
                           return 0
                        end if
   8:
   9:
                        return red with 2black(T→Left)//recursion
   10:
                        return red with 2black(T→Right)
   11:
                else return 1
   12:
           end
```

```
/*According to 4 : If a node is red, then both its children must be black.*/
int red_with_2black(tree T)
{
    if (T)
    {
        if (T->Element < 0)
        {
        }
}</pre>
```

```
if (T->Left)
{
      if (T->Left->Element < 0)return 0;
}
if (T->Right)
{
      if (T->Right->Element < 0)return 0;
}
}
return red_with_2black(T->Left);//check leftchild
return red_with_2black(T->Right);//check rightchild
}
else
return 1;
}
```

4.According to (5): For each node, all simple paths from the node to descendant leaves contain the same number of black nodes. A red black tree can make the black node balanced, which means the left child and the right child must have the same depth(we just care about the black node)

```
Algorithm3: equal_black
            tree T
Input:
Output:
            0→wrong 1→right
    1:
            function equal_black(tree T)
    2:
                if (!T)
    3:
                   return 1
    4:
                int left_black = find_blackNode_depth(T->Left)
    5:
                int right_black = find_blackNode_depth(T->Right)
    6:
                if (left_black == right_black) return 1
                else return 0
    7:
    8:
                /*Recursion*/
    9:
                if (T->Left)
   10:
                   return equal_black(T->Left)
   11:
                if (T->Right)
   12:
                   return equal_black(T->Right)
   13:
            end
```

/*According to 5 :For each node, all simple paths from the node to descendant leaves contain the s ame number of black nodes. A red black tree can make the black node balanced, which means the leftchild and the rightchild must have the same depth(we just care about the black node)

```
int equal_black(tree T)
{
    if (!T)
        return 1;
  /*find the leftchild and rightchild black node depth*/
    int left_black = find_blackNode_depth(T->Left);
    int right_black = find_blackNode_depth(T->Right);
   /*the black node depth must equal to each other*/
    if (left black == right black)
        return 1;
    else
        return 0;
    /*Recursion*/
    if (T->Left)
        return equal_black(T->Left);
    if (T->Right)
        return equal_black(T->Right);
}
```

5. Find the depth of node T to the leaf. And the depth is calculate by the bl ack node. UsePostorder traversal, find the black node height of the curre nt tree.

```
Algorithm4: find_blackNode_depth
Input:
              tree T
Output:
              the number of black node on the deepth
     1:
              function find_blackNode_depth(tree T)
     2:
                if (!T)
     3:
                   return 0
     4:
                int left_black = find_blackNode_depth(T->Left)
     5:
                int right_black = find_blackNode_depth(T->Right)
     6:
                if (T->Element > 0) /*Black Node*/
     7:
                   if (left_black > right_black)
     8:
                     return left_black + 1
     9:
                   else
    10:
                     return right_black + 1
    11:
                else /*red node*/
    12:
                   if (left_black > right_black)
    13:
                     return left_black
    14:
                   else return right_black
    15:
              end
```

```
/*Find the deepth of node T to the leaf And the deepth is calculate by the black node. Post-order traversal, find the black node height of the current tree.*/
int find_blackNode_depth(tree T)
{
```

```
if (!T)
        return 0;
    int left_black = find_blackNode_depth(T->Left);
    int right black = find blackNode depth(T->Right);
    if (T->Element > 0) /*Black Node*/
        if (left_black > right_black) /*The bigger number is the depth,rather than t
he smaller, so we should return the bigger one*/
        {
             return left_black + 1;
        }
        else
             return right_black + 1;
        }
    }
    else /*red node*/
    {
        if (left_black > right_black)
             return left_black;
        }
        else
        {
             return right_black;
        }
    }
}
```

6.The main loop:

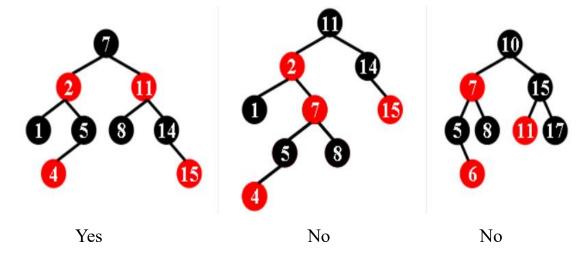
First read the number of binary trees, then read the number of nodes of the first binary tree, and then read the traversal sequence of the binary tree. Then read the number and sequence of nodes in turn and judge whether the binary tree is a legal red-black tree in real time. If it is, print "YES", if it is illegal, print "NO".

```
int main()
{
    int num = 0;//The number of testing tree
    scanf("%d", &num);
    for (int i = 0; i < num; i++)
    {
        tree T = NULL;
        int number_of_node = 0;
}</pre>
```

```
scanf("%d", &number_of_node);
for (int j = 0; j < number_of_node; j++)
{
    int data = 0;
    scanf("%d", &data);
    T = Insert(data, T);
}
if (checkroot(T) && red_with_2black(T) && equal_black(T))
//To judge the temp tree use the 3 algorithm.
    printf("YES\n");
else
    printf("NO\n");
    FreeTree(T); //To free the memory
}
system("pause");
return 0;
}</pre>
```

Chapter 3: Testing Results

Example:



Result table:

Test case	a brief description of the purpose	Expected res ult	Act ual beh avio r	Possible cause o f a bug	Current status
3	To test whether th	Yes	Yes	No	Pass
9	e program can det	No	No		
7 -2 1 5 -4 -11 8 14 -1	ect the input given				

		Γ	ı	I	
5	by the project	No	No		
9					
11 -2 1 -7 5 -4 8 14 -1					
5					
8					
10 -7 5 -6 8 15 -11 17					
1	To test whether th	No	No	No	Pass
1 _	e program can jud	The root m			
-1	ge by the root	ust be black			
	(property 2)				
30	To test whether th	No	No	No	Pass
30	e program will out				
1 2 3 430	of memory.	•••	•••		
		No	No		
30					
1 2 3 4 30					
1 2 3 4 30					
1	To test whether th	Yes	Yes	No	Pass
0	e program can jud	NULL is			1 433
	ge by the root	a kind of b			
NULL	(property 2)	lack node			
		lack hode			
1	To test whether th	No	No	No	Pass
3	e code can judge	12			
12	by the Property 5:				
-11	For each node, all	-11 13			
13	simple paths from	the path			
	the node to desce	from the			
	ndant leaves contai	node 12 to the leaf			
	n the same numbe	has			
12	r of black nodes.	different			
		black node			
11 13		12→-11→N			
		has 2 black			
		nodes			
		12→13→N			
		has 3 black			
		nodes			

1 4 11 -2 -14 -7 11 2 7	To test whether the code can judge by the Property 4: If a node is red, then both its children are black.	No In the tree, 11→2→7→ NULL both 2 and 7 are the red node, so according to property 4, the tree is not legal	No	NO	Pass
1		No	No	No	Pass
10 1 2 3 4 5 -6 -7 -8 -9 - 10					
2		No	No	No	Pass
5		No	No		
1 2 -3 -4 5		1.0			
7					
1 -4 -6 2 5 3 7					
2		No	No	No	Pass
5		No	No		
1 2 -3 -4 5					
7					
-1 -4 -6 2 5 3 7		No	No	NI-	Dana
1 8		NO	NO	No	Pass
-1 -4 -6 2 5 3 7 -9 -8					
1 7 0 2 3 3 7 -3 -0					
1		No	No	No	Pass
11					
1 2 3 4 5 6 7 8 9 10					
11					
1		No	No	No	Pass
11					
11 10 9 8 7 6 5 4 3 2					
1		27			
1		No	No	No	Pass
10					
-1 2 -3 4 -5 6 -7 8 -9 -					
10					

Chapter 4: Analysis and Comments

Analysis of the time and space complexities of the al gorithms. Comments on further possible improvements.

- 1. **Algorithm1** Check root, we can see that both the time complexity and space complexity is O(1), because we just need to judge the root of the tree.
- 2. **Algorithm2,** Check children, after analysis the code we can find that the function Check children is a kind of preorder traversal, we need to check every node, so t he time complexity is O(N). And we don't need extra space, so the space comple xity is O(1)
- 3. **Algorithm3** equal_black and **Algorithm4** find_blackNode_depth are used together in function equal_black. In the function we used find_blackNode_depth function for every node, and the time complexity of it is O(logN), so the total time complexity is **O(NlogN)**, And we don't need extra space, so the space complexity is O(1). When we build the tree the space complexity is O(N)
- 4. All in all the total time complexity is O(NlogN), the space complexity is O(N)

Further possible improvements

- 1. If the node→element is 0, we don't know whether it is red node or black node. So we should define it as black node.
- 2. The writing of function names and other code specifications can be improved, and when time permits, other methods can be devised to solve the practical problem.
- 3. Algorithm is implement by the recursion method, we can reimplement them by s ome basic data structure such as stack and queue. Because the efficiency of recurs ive programs is lower than that of loop programs

Appendix: Source Code (in C)

At least 30% of the lines must be commented. Otherwise the code will NOT be evaluated.

```
#include <stdio.h>
#include <stdlib.h>
/*define the node*/
typedef struct Node *tree;//define the tree
typedef tree ptr;  //define the pointer
struct Node
                          //define the node
    int Element;
    ptr Left;
    ptr Right;
};
/*calculate the absolute value */
int abs(int a)
{
    if (a < 0)
        a = -1 * a;
    return a;
}
/*to build the tree*/
tree Insert(int X, tree T)
{
    if (T == NULL)
    { /* Create and return a one-node tree */
        T = malloc(sizeof(struct Node));
        if (T == NULL)
            printf("Out of space!!!\n");
        else
            T \rightarrow Element = X;
           T->Left = T->Right = NULL;
        /* End creating a one-node tree */
    else /* If there is a tree */
        if (abs(X) < abs(T->Element))
        {
            T->Left = Insert(X, T->Left);
```

```
}
      else if (abs(X) > abs(T->Element))
          T->Right = Insert(X, T->Right);
      }
   }
   /* Else X is in the tree already; we'll do nothing */
   return T;
}
/*free tree by postorder*/
void FreeTree(tree T)
{
   if (T)
   {
      FreeTree(T->Left);
      FreeTree(T->Right);
      free(T);
      T = NULL;//let the pointer point to NULL
   }
}
****************
***************
/*According to 3 The root should be black node.*/
int checkroot(tree T)
{
   if (!T)//black node
      return 1;
   else if (T->Element > 0)//black node
      return 1;
   else//red node
      return 0;
}
/*
According to 4 : If a node is red, then both its children must be blac
k.
*/
int red_with_2black(tree T)
{
   if (T)
```

```
{
        if (T->Element < 0)//red node</pre>
            if (T->Left)
            {
                if (T->Left->Element < 0)//red node</pre>
                    return 0;
            }
            if (T->Right)
            {
                if (T->Right->Element < 0)//red node</pre>
                    return 0;
            }
        }
        return red_with_2black(T->Left);
        return red_with_2black(T->Right);
    }
    else//NULL black node
        return 1;
}
/*According to 5 :For each node, all simple paths from the node to desc
endant leaves contain the same number of black nodes.
So a red black tree can make the black node balanced, the leftchild an
d the rightchild must have the same depth(we just care about the blac
k node)*/
int equal_black(tree T)
{
    if (!T)
        return 1;
    /*find the leftchild and rightchild black node depth*/
    int left_black = find_blackNode_depth(T->Left);
    int right_black = find_blackNode_depth(T->Right);
    /*the black node depth must equal to each other*/
    if (left_black == right_black)
        return 1;
    else
        return 0;
    /*Recursion*/
    if (T->Left)
        return equal_black(T->Left);
    if (T->Right)
```

```
return equal_black(T->Right);
}
/*Find the deepth of node T to the leaf And the deepth is calculate b
y the black node.
Post-order traversal, find the black node height of the current tree.*/
int find_blackNode_depth(tree T)
   if (!T)
      return 0;
   int left_black = find_blackNode_depth(T->Left);
   int right_black = find_blackNode_depth(T->Right);
   if (T->Element > 0) /*Black Node*/
      if (left_black > right_black) /*The bigger number is the depth,
rather than the smaller, so we should return the bigger one*/
          return left_black + 1;
      }
      else
          return right_black + 1;
      }
   }
   else /*red node*/
   {
      if (left_black > right_black)
      {
          return left_black;
       }
      else
       {
          return right_black;
      }
   }
}
******************
***************
*/
/* to judge the tree */
int judge(tree T)
{
```

```
if (checkroot(T) && red_with_2black(T) && equal_black(T))
    {
        return 1;
    }
    else
    {
        return 0;
    }
}
int main()
{
    int num = 0;//The number of testing tree
    scanf("%d", &num);
    for (int i = 0; i < num; i++)</pre>
    {
        tree T = NULL;
        int number_of_node = 0;
        scanf("%d", &number_of_node);
        for (int j = 0; j < number_of_node; j++)</pre>
            int data = 0;
            scanf("%d", &data);
            T = Insert(data, T);
        }
        if (judge(T))//To judge the temp tree
            printf("YES\n");
        }
        else
        {
            printf("NO\n");
        FreeTree(T); //To free the memory
    }
    system("pause");
    return 0;
}
```

Declaration

I hereby declare that all the work done in this project titled "Is It A Red-Black Tree" is of my independent effort.