

Connected Graphs and Spanning Trees

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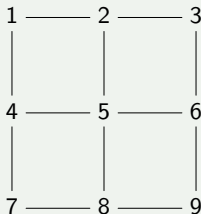
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Describing the problem I

$G = (V, E)$ is a graph, where

- ① V is a set of vertices
- ② E is a (multi)set of edges

Example 1



$$V = \{1, \dots, 9\}$$

$$E = \{[1, 2]; [1, 4]; [2, 3]; [2, 5]; [3, 6]; [4, 5]; [4, 7]; [5, 6]; [5, 8]; [6, 9]; [7, 8]; [8, 9]\}$$

Describing the problem II

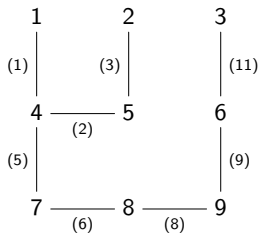
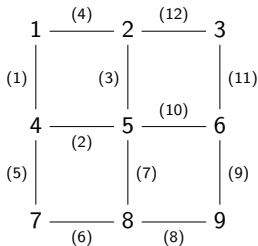
Let $G = (V, E)$ be a connected graph.

$T = (V, E')$ is spanning tree of G when

- ① T is tree, and
- ② $E' \subseteq E$.

Theorem 2

Every connected graph has a spanning tree.



Towards formalization

Specification

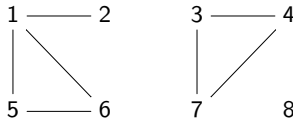
- ① Define data types to represent graphs
 - ▶ `connected`
- ② Define the following functions:
 - ▶ `nocycle`
 - ▶ `mktree`

Verification

- ① $\text{connected}(G) \Rightarrow \exists G' \subseteq G. \text{tree}(G')$
- ② $\text{connected}(G) \Rightarrow \text{connected}(\text{mktree}(G)) \wedge \text{nocycle}(\text{mktree}(G))$

Functions on graphs

$$G = (\{1, 2, 3, 4, 5, 6, 7, 8\}, [1, 2]; [1, 6]; [1, 5]; [3, 4]; [3, 7]; [4, 7]; [5, 6])$$



- $\text{mcc}(A, G) = \text{max. connected comp. of } A \text{ in } G.$
 $\text{mcc}(6, G) = \{1, 2, 5, 6\}, \text{mcc}(8, G) = \{8\}$
- $\#cc(G) = \text{no. of max. connected components}$
 $\#cc(G) = 3$
- $\text{nocycle}(G) = \text{false}$

Spanning forests

In the attempt of proving the desired properties we realized it is much easier to prove a more general result:

Every graph has a spanning forest!

Definition 3

A **spanning forest** of a graph is a subgraph that consists of a set of spanning trees, one for each maximal connected component of the initial graph.

We define the function `mktree` which returns the spanning forest of a graph.

$$G = (\{1, 2, 3, 4, 5, 6, 7, 8\}, [1, 2]; [1, 6]; [1, 5]; [3, 4]; [3, 7]; [4, 7]; [5, 6])$$



$$\text{mktree}(G) = (\{1, 2, 3, 4, 5, 6, 7, 8\}, [1, 2]; [1, 6]; [3, 4]; [3, 7])$$

Remark 4

The value is relative to the order chosen for the edges.

Properties to prove

- ① $\text{mcc}(A, G) = \text{mcc}(A, \text{mktree}(G))$
- ② $\#cc(G) = \#cc(\text{mktree}(G))$
- ③ $\text{nocycle}(\text{mktree}(G))$

Then we define $\text{connected}(G) := (\#cc(G) = 1)$ which implies

$$\text{connected}(G) \Rightarrow \text{connected}(\text{mktree}(G)) \wedge \text{nocycle}(\text{mktree}(G))$$