# **Connected Graphs and Spanning Trees**

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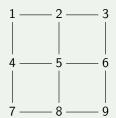
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## Describing the problem I

G = (V, E) is a graph, where

- V is a set of vertices
- E is a (multi)set of edges

### Example 1



- $V = \{1, \dots, 9\}$
- $E = \{[1,2], [1,4], [2,3], [2,5], [3,6], [4,5], [4,7], [5,6], [5,8], [6,9], [7,8], [8,9]\}$

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## Describing the problem II

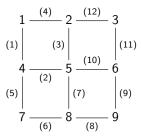
Let G = (V, E) be a connected graph.

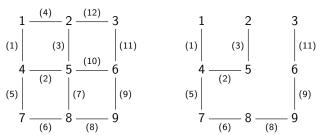
T = (V, E') is spanning tree of G when

- T is tree, and
- $\mathbf{Q}$   $E' \subseteq E$ .

#### Theorem 2

Every connected graph has a spanning tree.







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### **Towards formalization**

### Specification

- Define data types to represent graphs
  - connected
- ② Define the following functions:
- nocycle
- mktree

#### Verification

- $\bigcirc$  connected(G) $\Rightarrow \exists G' \subseteq G.tree(G')$
- $\bigcirc$  connected(G) $\Rightarrow$  connected(mktree(G)) $\land$  nocycle(mktree(G))

## **Functions on graphs**

- mcc(A,G) = max. connected comp. of A in G.
  mcc(6,G) = {1,2,5,6}, mcc(8,G) = {8}
- #cc(G) = no. of max. connected components#cc(G) = 3
- nocycle(G) = false



## **Spanning forests**

In the attempt of proving the desired properties we realized is much easier to prove a more general result:

Every graph has a spanning forest!

### **Definition 3**

A **spanning forest** of a graph is a subgraph that consists of a set of spanning trees, one for each maximal connected component of the initial graph.

We define the function mktree which returns the spanning forest of a graph.

#### Remark 4

The value is relative to the order chosen for the edges.

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### Properties to prove

- 1 mcc(A,G)=mcc(A,mktree(G))
- #cc(G)=#cc(mktree(G))
- 3 nocycle(mktree(G))

Then we define connected(G):= (#cc(G)=1) which implies

 $connected(G) \Rightarrow connected(mktree(G)) \land nocycle(mktree(G))$ 

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