

# KF04 project report

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# 1 Features

This chapter describes the features and functions, we have included in our Map of Denmark project, part 2, Spring 2011.

- Find shortest route
- Display path using blue coloured roads
- Multiple pins
- Total distance
- Estimated travel time
- Bike or car route
- Remove a single pin by clicking
- Remove all pins

## 1.1 Find shortest route

Our map is able to calculate the shortest route between two points. This is done by clicking the map. Clicking the map places a pin at the clicked location. If the user clicks on another part of the map, our map will show the shortest route between those two locations, if there is a route between them.

## 1.2 Display path using blue coloured roads

When the user has selected two or more locations, the map will show the shortest route highlighted in blue. The algorithm we use for finding this route (A\*) is quite fast, so a user will not feel annoyed by how long it takes to find a route.

## 1.3 Multiple pins

It is possible for a user to place multiple pins on the map. This is done similarly to finding the shortest route between two points. When any number of pins are placed, the user can add an additional one. This will calculate the route between the latest two pins placed. The route shown will be from pin  $1 \rightarrow 2 \rightarrow 3 \rightarrow 4..$

## **1.4 Total distance**

Our Map of Denmark program calculates the total distance of the route found, a functionality that assists the end-user in planning their trip. This was relatively simple to implement, but adds a big functionality.

## **1.5 Estimated travel time**

As with the total travel distance, our Map of Denmark calculates the estimated time it will take to travel the route found by our algorithm. This was relatively easy to calculate and is a huge benefit for the end user. We estimate the travel time of bikes to be 15 km/hour and with no breaks. This is optimistic, but we prefer optimistic to not having the feature at all.

## **1.6 Bike or car route**

It is possible for the user to choose between a car route and a bike route. The different options will make our algorithm consider what is possible to traverse for both cars and bikes. It will make sure not to show small paths for cars and not show highways for bikes.

## **1.7 Remove a single pin**

To make it easier to use our Map of Denmark, we have it possible to remove a pin by clicking on the location it was placed. This makes it possible for the end user to clear up a missclick without having to shut down the entire Map or clearing all the pins the user has placed.

## **1.8 Remove all pins**

In addition to enabling the user to remove a single pin by clicking on it, we have also made it possible to remove all the pins currently placed. This can be done by either pressing the “c” key on the keyboard or by clicking the button on the graphical interface. This means that the user will not have to click every single point in order to start a new route planning.

## 2 Implementation

This chapter describes our implementation of interesting methods and functions in our Map of Denmark project, part 2, Spring 2011.

### 2.1 Model

In order to implement the previously mentioned pathfinding the **Model** has been extended to include a list of **Edges**. These correspond to all the roads that are in the current route. Whenever it receives a new request to find a path from one **Node** to another it uses the static method in the **Dijkstra** class and adds the result to the list. When **Control** needs to draw the route it uses a method that returns the entire route as **Line** objects relative to the view. This approach is identical to what we use when the **Control** requests all roads that should be drawn as the map for a given view.

We have also included some statistics of the current route. These are calculated by the **Model** since it's the object that primarily handles the data. This information can be calculated in linear time relative to the length of the current route since the route is saved in memory when first found.

### 2.2 Dijkstra

Originally we intended to use Dijkstras algorithm to find the shortest path from one node in the graph to another. Dijkstras algorithm builds a **Shortest Path Tree** by always looking at the node closest to the source node. Because of this, whenever the destination node is reached, the shortest path to that node is found. When that happens we stop the algorithm, because we have found what we were looking for.

In order to always choose the node closest to the source, we use a **priority queue**, where the distance to the source indexes each node. Sometimes we want to decrease the index of a node, because we have found a route that is shorter than the previously indexed route. Because the **priority queue** in Java's library is slow when decreasing keys, we use an implementation from [1] called **IndexMinPQ**. Our Dijkstras algorithm is a version inspired from the implementation in [1].

We made some rather big changes to the Dijkstra algorithm to make it fit our requirements. First of all, we replaced the arrays that are used to hold *edgeTo* and *distTo* in the [1] with **Maps**, so we can use our objects instead of integers used in [1]. We have not changed anything in the **IndexMinPQ** implementation. It takes integers as values - this is not a problem, however,

because our `nodes` have unique index integers. Another major change we have made compared to the Dijkstra implementation from [1] is that ours is a static method instead of an object. This means we do not have to instantiate it before we use it.

To be able to compare the edges differently depending on the type of transportation the user has selected, we have created an `Evaluator`. An evaluator is given as argument to the pathfinding and is responsible for transforming the `KrakEdges` to a value that can be inserted into the priority queue.

We have made sure one-way roads are only traversable if you are coming from the right direction. This check is performed every time a road is going to be added to the priority queue. If the road cannot be traversed, it is skipped.

It is important to note that we have found an error in the data from `krak`. Some of the highways are oriented in the wrong direction, so the pathfinding cannot traverse a long stretch of highway. This will make it look like it wants the user to get on and off the highway in a very weird way.

We have included a “hack” to avoid this issue. This means that we can always travel on highways even if they are facing in the opposite direction. This works and gives accurate routes, because the entry and exit points for the highways are without error.

This gives accurate roads, but it is a bad fix. We have decided to use it, but it would be preferable if the data was correct. There are other errors in the data as well, where roads are not traversable at all, because they are one-way roads in both directions, which means they cannot be entered from either end.

### 2.2.1 A-star ( $A^*$ )

As mentioned in the start of the previous section, we originally intended to use Dijkstra. After we had implemented it, we learned of another, similar algorithm called A-star or  $A^*$ .

In order to make our path finding algorithm even faster, we used a variation called A-star. This implementation benefits from the fact that the roads, in contrast to a theoretical graph, are part of a geometric system. This means that we can calculate the fastest possible route from one point to another. This helps us determine which routes to look at first. The fastest possible route is always a straight line from the source point to the end point. The smaller the difference from the straight line is to the fastest possible route, the more we will benefit from the A-star technique. This difference will often

be bigger for the routes calculated for the car. This is because the fastest route would be a straight highway, which often is far from possible. We can therefore conclude that the A-star technique typically is more of a benefit when calculating routes for bicycling because they rely on the distance. In order to make the code work with our dynamic evaluations for the **Edges** we have created another method in the **Evaluator** that calculates the heuristic for at given **Node** relative to the target **Node**.

## 2.3 Control

A lot has changed in **Control** since our last version. We have changed a lot of the methods to work as static methods in separate **Tools** classes, so **Control** will feel less cluttered in this version. New logic has taken the place of the old logic however. Route calculation, pin placement, pin removal and route display all goes through **Control**.

When a user clicks on the map, **Control** uses our **Tools** classes to convert the screen coordinates to UTM coordinates. It then stores these in a **List**. If the list holds two or more elements, it asks the **Model** to calculate the shortest path between the first and second element, the second and third element and so on, if any additional elements exist beyond the first two. Every time **Control** calls it's internal repaint method, it clears all the pins currently in the **View**. It then recalculates the pin position(s) according to the new size and postion of the **View**.

When a user takes advantage of one of the options to clear all the pins currently on the Map of Denmark, **Control** empties the collection holding all the pin points, asks **Model** to clear the route and asks **View** to repaint.

If a user removes a single pin by clicking on the location of the pin, we run through our collection and remove one of the pins within a specific range of pixels. It then makes **Model** recalculate the route for all the pins. This is not an effective implementation - if there are a lot of pins, then removing one might make the recalculating take a while. We decided to use this implementation, for now at least, because we wanted a working functionality that we could later improve upon if nessecary.

## 2.4 Car/Bike routes

Another useful feature we have included is the ability to switch between multiple modes of pathfinding. We have named these two modes Car and Bike. In the Car mode the pathfinding will prioritize lower travel times, since we believe this to be most important when driving. Additionally the

Car mode will exclude small paths and pedestrian zones. The Bike mode prioritizes lower distances highest since intersections and speed-limits are of relatively low importance here. When using Bike mode the route will not include highways.

As described earlier these different ways of prioritizing the roads are achieved by using the **Evaluator** that is provided when asking the pathfinding method for a best path. We have written these two evaluators as static objects on the **Evaluator** class for easy use in the code, but the use of a separate object for the evaluations open up for later variation with little impact on the original code.

## 2.5 View

When looking at the changes in **View** relatively little has happened since most of the new code is located "behind the curtain". The **Canvas** has been extended to allow adding and drawing of the route. This is however, as mentioned earlier, done in exactly the same way as with the roads in the map itself. The **Canvas** also includes some points that is positions where the pins are placed. These are drawn on top of the map and (like the map) needs provided again if the map moves.

When looking at the menu to the left there are a few more options to the user. There is a new button to remove all pins from the map and two **RadioButtons** to toggle between Car mode and Bike mode in the pathfinding.

Beneath the tools for pathfinding there are displayed two pieces of information (and spare space for more). This information is changed by a simple public method in the **View** and the units will adjust will be adjusted accordingly.



## 3 Features that did not make it in

This chapter covers features we spend a lot of time on, but did not make it into the version that we hand in with this report.

### 3.1 Serialization

Everytime we start the program, we loop through the entire dataset given to us by Krak. This data is huge, and because of this it takes quite a lot of time to start the program. Because we need to load all these data, the user is presented with a blank screen for a long time, before all these data are loaded and the program starts.

We started looking for a way to speed up the loading process, so the user has a map in front of him or her quickly, when the user starts the program. What takes the most time is looping through the data and creating the needed datastructures (quadtrees and so on), so if we could skip these steps or speed them up, we could save a lot of time.

This is where serialization comes into play. By serializing an object, you transform your object into something that can be passed around, through streams and such. So by serializing objects, you can save them to files. If the object to be serialized contains references to other objects, these will also be serialized (if they are `Serializable` / implements `java.io.Serializable`). By doing this, we only need to build our datastructures the first time you start the program. After the objects have been created, they are been serialized and saved to files. The next time the user starts the program, we check whether the data has been changed. If it hasn't been changed, we load the objects that we saved, instead of making them all over.

Neither our serialization or our threading (described below) features made it into this version, because we had problems with using the serialized objects when making our Shortest Path Tree. It was always intended as a nice to have feature, and we hope to have it working in a later version of the project.

#### 3.1.1 Threading

By serializing our main objects, we cut several seconds of our load times. But we can do it even faster. We are using several objects, but only few of them are needed right from the start. So what we can do to speed it up even more, is loading the few necessary objects, and then load the rest in the background. We do this with threads. We load the few objects we need

from the start, then create a new thread to load the rest of the objects, and in the mean time, we create the window and draw the map.

The same goes for the first run. The user doesn't need to wait for the program to finish serializing and saving to files. By using threads, we can create the datastructures, and then immediately show the window to the user, while saving the objects to files in the background.

## References

- [1] Robert Sedgewick and Kevin Wayne. *Algorithms, Fourth Edition*. Preliminary Edition Fall 2010. Addison-Wesley 2010.