

Complexity

 Ω - omega - lower bound Θ - theta - constant O - big O - upper bound1) Compute $T(n)$ - number of steps

$\Omega(n)$

$\Theta(n^2)$

$O(n^2); O(n^3)$

$T(n) = 2n^2 + n$

1) $n^2 \in O(n^3)$ - true

6) $2^n \in O(n!)$ - true

2) $n^3 \in O(n^2)$ - false

7) $\lg n \in O(\log_2 n)$ - true

3) $2^{n+1} \in \Theta(2^n)$ - true

4) $2^{2n} \in \Theta(2^n)$ - false

5) $n^2 \in \Theta(n^3)$ - false

8) $O(n) + \Theta(n^2) = \Theta(n^2)$

9) $\Theta(n) + O(n^2) = O(n^2)$

10) $O(f) + O(g) = O(\max(f, g))$

11) $3^n \in O(2^n)$ - false

12) $\log_2 3^n \in O(\log_2 2^n)$ - true

13)

	Runtime				Space Complexity
	B.C.	W.C.	Avg	Total	
Linear Search	$\Theta(1)$	$\Theta(n)$	$\Theta(n)$	$O(n)$	$\Theta(1)$
Binary Search	$\Theta(1)$	$\Theta(\log n)$	$\Theta(\log n)$	$O(\log_2 n)$	$\Theta(1)$
Selection Sort	$\Theta(n^2)$	$\Theta(n^2)$	$\Theta(n^2)$	$\Theta(n^2)$	$\Theta(1)$ = in-place
Bubble Sort	$\Theta(n)$	$\Theta(n^2)$	$\Theta(n^2)$	$O(n^2)$	$\Theta(1)$
Quick Sort	$\Theta(n \log n)$	$\Theta(n^2)$	$\Theta(n \log n)$	$O(n^2)$	$\Theta(1)$
Merge Sort	$\Theta(n \log n)$	$\Theta(n \log n)$	$\Theta(n \log n)$	$\Theta(n \log n)$	$\Theta(n)$

14) Subalgorithm $S1(n)$ is:for $i \leftarrow 1, n$ execute:

1.1) subalgorithm is:

```

for i ← 1, n execute:
  j ← n
  while j ≠ 0 execute:
    j ← ⌊ $\frac{j}{2}$ ⌋
  end-while
end-for
end-subalgorithm

```

$\Theta(n \log n)$

$$T(n) = \Theta(n \log n)$$

Subalgorithm S1(n) is:

```

for i ← 1, n execute:
  j ← i
  while j ≠ 0 execute:
    j ← ⌊ $\frac{j}{2}$ ⌋
  end-while
end-for
end-subalgorithm

```

$$\sum_{i=1}^n \log_2 i = \log_2 1 + \log_2 2 + \dots + \log_2 n = \log_2 n!$$

$$\Theta(\log_2 n!)$$

$$T(n) = \Theta(\log(n!)) = \Theta(n \log n) \quad (\text{Sterling approximation})$$

Subalgorithm (x, n, a) is:

```

found ← false
for i ← 1, n execute:
  if  $x_i = a$  then:
    found ← true
  end-if
end-for
end-subalgorithm

```

$\Theta(n)$

subalgorithm (x, n, a) is:

```

found ← false
i ← 1
while i ≤ n execute:
    if xi = a then:
        found ← true
    end-if
end-for
end-subalgorithm

```

BC: $\Theta(1)$ Total: $\Theta(n)$

WC: $\Theta(n)$

AC: $\Theta(n)$

$$\begin{aligned}
 \text{AC: } T(n) &= C_1 P_{C_1} + C_2 P_{C_2} + \dots + C_n P_{C_n} = \\
 &= \frac{1}{n} (1 + 2 + \dots + n) = \frac{n(n+1)}{2n} = \frac{n+1}{2} \\
 &\Rightarrow \Theta(n)
 \end{aligned}$$