



Conversões entre gramáticas e autômatos finitos

- Se $G = (V, \Sigma, P, S)$ é uma gramática regular, então o NFA $N = \langle \Sigma, Q, q_0, \delta, F \rangle$, definido como segue, é tal que $\mathcal{L}(N) = \mathcal{L}(G)$:

$$Q = \begin{cases} V \cup \{Z\}, & \text{se } (A \rightarrow a) \in P, \text{ onde } Z \notin V; \\ V, & \text{caso contrário;} \end{cases}$$

$$q_0 = S;$$

$$\delta(A, a) = B, \text{ sempre que } A \rightarrow aB \in P;$$

$$\delta(A, a) = Z, \text{ sempre que } A \rightarrow a \in P;$$

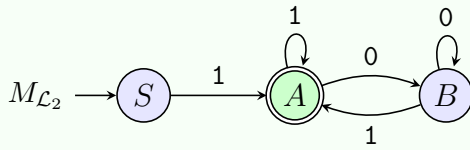
$$F = \begin{cases} \{A \mid A \rightarrow \varepsilon \in P\} \cup \{Z\}, & \text{se } Z \in Q; \\ \{A \mid A \rightarrow \varepsilon \in P\}, & \text{caso contrário.} \end{cases}$$

- Se $N = \langle \Sigma, Q, q_0, \delta, F \rangle$ é um NFA, então a gramática $G = (V, \Sigma, P, S) = (Q, \Sigma, \{q_i \rightarrow aq_j \mid \delta(q_i, a) = q_j\} \cup \{q_i \rightarrow \varepsilon \mid q_i \in F\}, q_0)$ é tal que $\mathcal{L}(G) = \mathcal{L}(N)$.

$$\mathcal{L}_1 = \{w \in \Sigma^* = \{0, 1\}^* \mid |w|_{01} > 0 \text{ ou } |w|_{10} > 0\}$$

$$\mathcal{L}_2 = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ representa um número binário ímpar (sem zeros à esquerda)}\}$$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_2

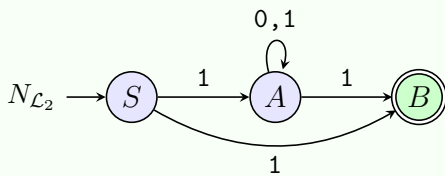


- Gramática G_1 que gera as cadeias de \mathcal{L}_2 :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 1A, \\ A \rightarrow 0B \mid 1A \mid \varepsilon, \\ B \rightarrow 0B \mid 1A \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_2 :



- Gramática G_2 que gera as cadeias de \mathcal{L}_2 :

$$G_2 = (V, \Sigma, P, S) = (\{A, B, S\}, \{0, 1\}, P, S), \text{ com:}$$

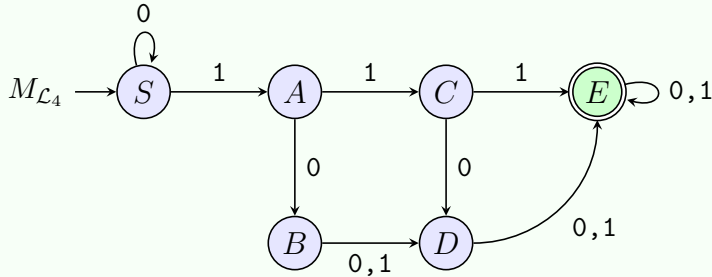
$$P = \left\{ \begin{array}{l} S \rightarrow 1A \mid 1B, \\ A \rightarrow 0A \mid 1A \mid 1B, \\ B \rightarrow \varepsilon \end{array} \right\}.$$



$$\mathcal{L}_3 = \{w \in \Sigma^* = \{0,1\}^* \mid w \text{ representa um número binário e } w \pmod{3} = 1\}$$

$$\mathcal{L}_4 = \{w \in \Sigma^* = \{0,1\}^* \mid w \text{ representa um número binário e } w \geq 7\}$$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_4

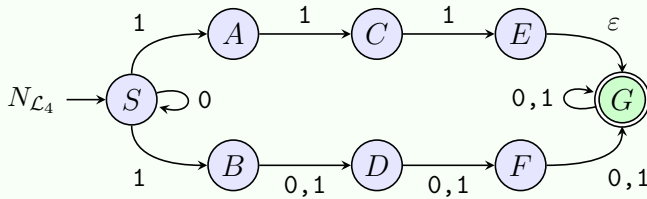


- Gramática G_1 que gera as cadeias de \mathcal{L}_4 :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, C, D, E, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0S \mid 1A, \\ A \rightarrow 0B \mid 1C, \\ B \rightarrow 0D \mid 1D, \end{array} \middle| \begin{array}{l} C \rightarrow 0D \mid 1E, \\ D \rightarrow 0E \mid 1E, \\ E \rightarrow 0E \mid 1E \mid \varepsilon \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_4 :



- Gramática G_2 que gera as cadeias de \mathcal{L}_4 :

$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, D, E, F, G, S\}, \{0, 1\}, P, S), \text{ com:}$$

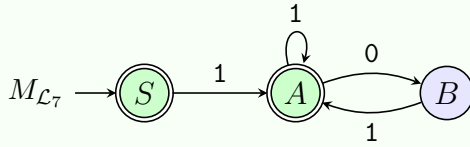
$$P = \left\{ \begin{array}{l} S \rightarrow 0S \mid 1A \mid 1B, \\ A \rightarrow 1C, \\ B \rightarrow 0D \mid 1D, \end{array} \middle| \begin{array}{l} C \rightarrow 1E, \\ D \rightarrow 0F \mid 1F, \\ E \rightarrow G, \end{array} \middle| \begin{array}{l} F \rightarrow 0G \mid 1G, \\ G \rightarrow 0G \mid 1G \mid \varepsilon \end{array} \right\}.$$

$$\mathcal{L}_5 = \{w \in \Sigma^* = \{0,1\}^* \mid w \text{ contém } 001 \text{ ou } 110\}$$

$$\mathcal{L}_6 = \{w \in \Sigma^* = \{0,1\}^* \mid w \text{ não contém } 001 \text{ ou não contém } 110\}$$

$$\mathcal{L}_7 = \{w \in \Sigma^* = \{0,1\}^* \mid \text{todo } 0 \text{ em } w \text{ é adjacente à esquerda e à direita a um } 1\}$$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_7

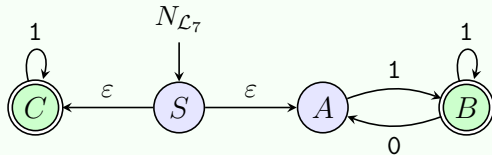


- Gramática G_1 que gera as cadeias de \mathcal{L}_7 :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 1A \mid \varepsilon, \\ A \rightarrow 0B \mid 1A \mid \varepsilon, \\ B \rightarrow 1A \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_7 :



- Gramática G_2 que gera as cadeias de \mathcal{L}_7 :

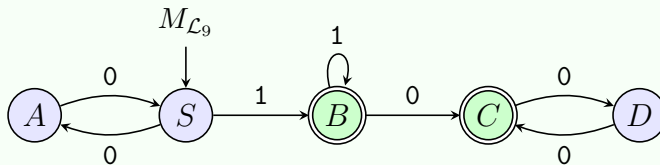
$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow A \mid C, \quad B \rightarrow 0A \mid 1B \mid \varepsilon \\ A \rightarrow 1B, \quad C \rightarrow 1C \mid \varepsilon \end{array} \right\}.$$

$$\mathcal{L}_8 = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ contém as subcadeias } 01 \text{ e } 10\}$$

$$\mathcal{L}_9 = \{w \in \Sigma^* = \{0, 1\}^* \mid w = xyz, \text{ com } x \in \{0\}^*, |x| = 2k, y \in \{1\}^+ \text{ e } z \in \{0\}^*, |z| = 0 \text{ ou } |z| = 2k' + 1; k, k' \in \mathbb{N}\}$$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_9

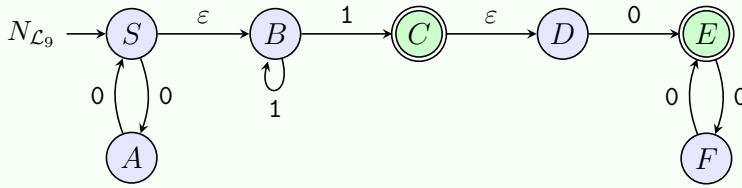


- Gramática G_1 que gera as cadeias de \mathcal{L}_9 :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, C, D, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1B, \quad C \rightarrow 0D \mid \varepsilon, \\ A \rightarrow 0S, \quad D \rightarrow 0C \\ B \rightarrow 0C \mid 1B \mid \varepsilon, \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_9 :



- Gramática G_2 que gera as cadeias de \mathcal{L}_9 :

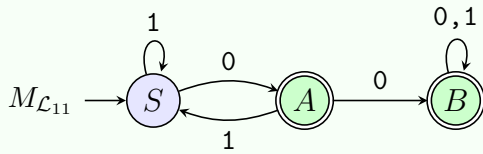
$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, D, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid B, \\ A \rightarrow 0S, \\ B \rightarrow 1B \mid 1C, \end{array} \left| \begin{array}{l} C \rightarrow D \mid \varepsilon, \\ D \rightarrow 0E, \\ E \rightarrow 0F \mid \varepsilon, \\ F \rightarrow 0E \end{array} \right. \right\}.$$

$$\mathcal{L}_{10} = \{w \in \Sigma^* = \{0, 1\}^* \mid w = x0y0z \text{ com } |y| = 2k \text{ ou } w = x1y1z \text{ com } |y| = 2k' + 1; x, y, z \in \Sigma^*; k, k' \in \mathbb{N}\}$$

$$\mathcal{L}_{11} = \{w \in \Sigma^* = \{0, 1\}^* \mid \text{pelo menos um } 0 \text{ em } w \text{ não é seguido de } 1\}$$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{11}

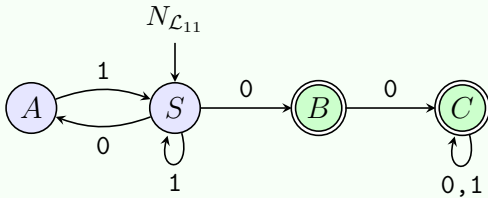


- Gramática G_1 que gera as cadeias de \mathcal{L}_{11} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1S, \\ A \rightarrow 0B \mid 1S \mid \varepsilon, \\ B \rightarrow 0B \mid 1B \mid \varepsilon \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{11} :



- Gramática G_2 que gera as cadeias de \mathcal{L}_{11} :

$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, S\}, \{0, 1\}, P, S), \text{ com:}$$

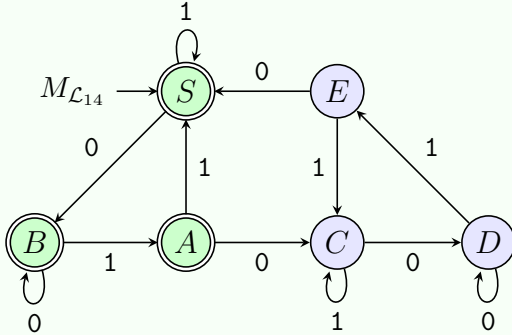
$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 0B \mid 1S, \\ A \rightarrow 1S, \end{array} \left| \begin{array}{l} B \rightarrow 0C \mid \varepsilon, \\ C \rightarrow 0C \mid 1C \mid \varepsilon \end{array} \right. \right\}.$$

$$\mathcal{L}_{12} = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ não contém } 101 \text{ e termina com } 1\}$$

$\mathcal{L}_{13} = \{w \in \Sigma^* = \{0,1\}^* \mid |w| \geq 3 \text{ e o terceiro e o penúltimo símbolos de } w \text{ não são } 1\}$

$\mathcal{L}_{14} = \{w \in \Sigma^* = \{0,1\}^* \mid w \text{ contém uma quantidade par da subcadeia } 010\}$

- $\mathcal{ER} = (1 \cup 0^+11)^*(0^+10(1 \cup 0^+11)^*0^+10(1 \cup 0^+11)^*)^*(0^*1^*)$.
- $\mathcal{ER} = (1 \cup 0^+11 \cup 0^+10(1 \cup 0^+11)^*0^+10)^*(\varepsilon \cup 0^+ \cup 0^+1)$.
- DFA mínimo que reconhece as cadeias de \mathcal{L}_{14}

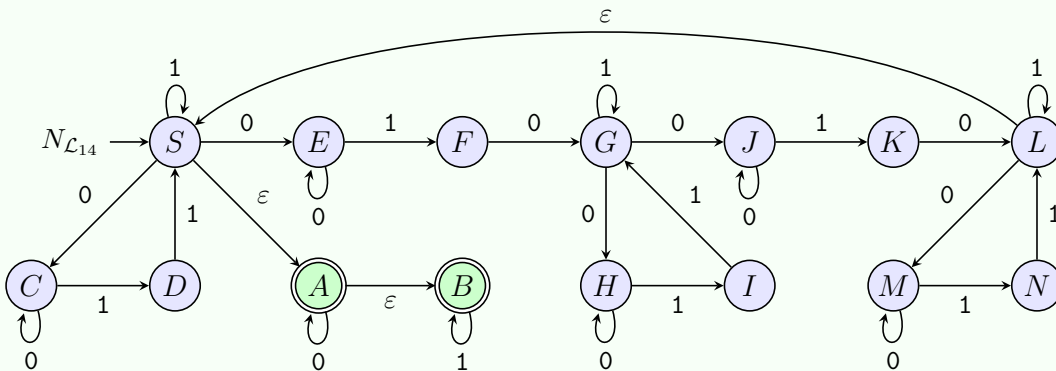


- Gramática G_1 que gera as cadeias de \mathcal{L}_{14} :

$G_1 = (V, \Sigma, P, S) = (\{A, B, C, D, E, S\}, \{0, 1\}, P, S)$, com:

$$P = \left\{ \begin{array}{ll} S \rightarrow 0B \mid 1S \mid \varepsilon, & C \rightarrow 0D \mid 1C, \\ A \rightarrow 0C \mid 1S \mid \varepsilon, & D \rightarrow 0D \mid 1E, \\ B \rightarrow 0B \mid 1A \mid \varepsilon, & E \rightarrow 0S \mid 1C \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{14} :



- Gramática G_2 que gera as cadeias de \mathcal{L}_{14} :

$G_2 = (V, \Sigma, P, S) = (\{A, B, C, D, E, F, G, H, I, J, K, L, M, N, S\}, \{0, 1\}, P, S)$, com:

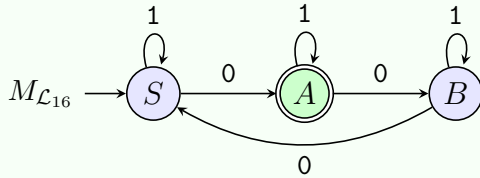
$$P = \left\{ \begin{array}{lll} S \rightarrow A \mid 0C \mid 0E \mid 1S, & E \rightarrow 0E \mid 1F, & J \rightarrow 0J \mid 1K, \\ A \rightarrow B \mid 0A \mid \varepsilon, & F \rightarrow 0G, & K \rightarrow 0L, \\ B \rightarrow 1B \mid \varepsilon, & G \rightarrow 0H \mid 0J \mid 1G, & L \rightarrow S \mid 0M \mid 1L, \\ C \rightarrow 0C \mid 1D, & H \rightarrow 0H \mid 1I, & M \rightarrow 0M \mid 1N, \\ D \rightarrow 1S, & I \rightarrow 1G, & N \rightarrow 1L \end{array} \right\}.$$



$\mathcal{L}_{15} = \{w \in \Sigma^* = \{0,1\}^* \mid w \text{ contém uma quantidade par da subcadeia } 000\}$

$\mathcal{L}_{16} = \{w \in \Sigma^* = \{0,1\}^* \mid |w|_0 \pmod 3 = 1\}$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{16}

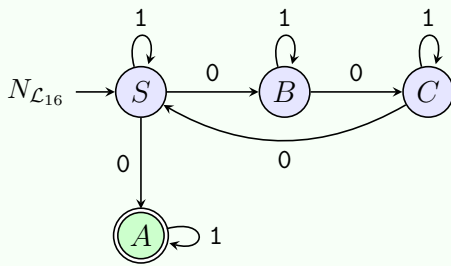


- Gramática G_1 que gera as cadeias de \mathcal{L}_{16} :

$G_1 = (V, \Sigma, P, S) = (\{A, B, S\}, \{0, 1\}, P, S)$, com:

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1S, \\ A \rightarrow 0B \mid 1A \mid \varepsilon, \\ B \rightarrow 0S \mid 1B \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{16} :



- Gramática G_2 que gera as cadeias de \mathcal{L}_{16} :

$G_2 = (V, \Sigma, P, S) = (\{A, B, C, S\}, \{0, 1\}, P, S)$, com:

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 0B \mid 1S, \\ A \rightarrow 1A \mid \varepsilon, \end{array} \middle| \begin{array}{l} B \rightarrow 0C \mid 1B \\ C \rightarrow 0S \mid 1C \end{array} \right\}.$$

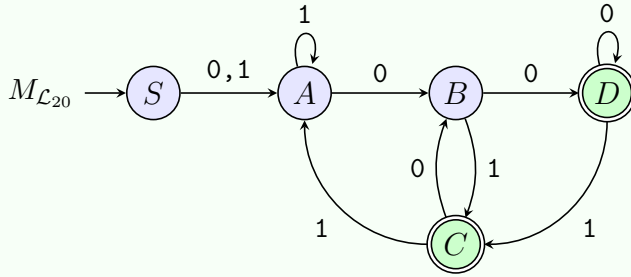
$\mathcal{L}_{17} = \{w \in \Sigma^* = \{0,1\}^* \mid |w|_0 \geq 3 \text{ e } |w|_1 \leq 2\}$

$\mathcal{L}_{18} = \{w \in \Sigma^* = \{0,1\}^* \mid |w|_0 \geq 3 \text{ ou } |w|_1 = 2, \text{ e } w \text{ não contém } 11\}$

$\mathcal{L}_{19} = \{w \in \Sigma^* = \{0,1\}^* \mid w \text{ contém exatamente uma ocorrência de } 00 \text{ ou de } 11\}$

$\mathcal{L}_{20} = \{w \in \Sigma^* = \{0,1\}^* \mid |w| \geq 3 \text{ e o penúltimo símbolo é } 0\}$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{20}

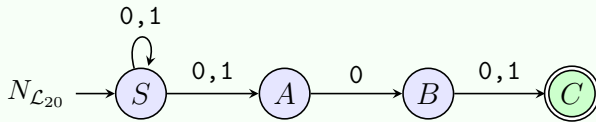


- Gramática G_1 que gera as cadeias de \mathcal{L}_{20} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, C, D, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1A, \\ A \rightarrow 0B \mid 1A, \\ B \rightarrow 0D \mid 1C, \end{array} \middle| \begin{array}{l} C \rightarrow 0B \mid 1A \mid \varepsilon, \\ D \rightarrow 0D \mid 1C \mid \varepsilon \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{20} :



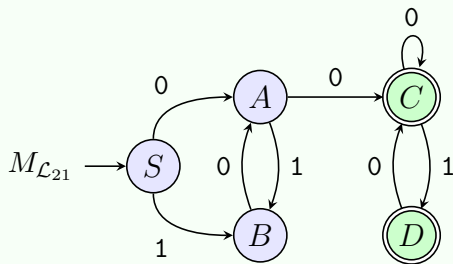
- Gramática G_2 que gera as cadeias de \mathcal{L}_{20} :

$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 0S \mid 1A \mid 1S, \\ A \rightarrow 0B, \end{array} \middle| \begin{array}{l} B \rightarrow 0C \mid 1C, \\ C \rightarrow \varepsilon \end{array} \right\}.$$

$$\mathcal{L}_{21} = \{w \in \Sigma^* = \{0, 1\}^* \mid |w|_{00} \geq 1 \text{ e } |w|_{11} = 0\}$$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{21}

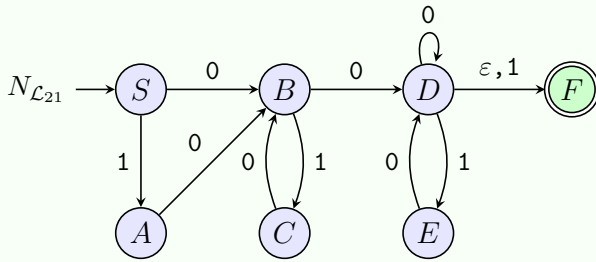


- Gramática G_1 que gera as cadeias de \mathcal{L}_{21} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, C, D, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1B, \\ A \rightarrow 0C \mid 1B, \\ B \rightarrow 0A, \end{array} \middle| \begin{array}{l} C \rightarrow 0C \mid 1D \mid \varepsilon, \\ D \rightarrow 0C \mid \varepsilon \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{21} :



- Gramática G_2 que gera as cadeias de \mathcal{L}_{21} :

$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, D, E, F, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l|l} S \rightarrow 0B \mid 1A, & D \rightarrow 0D \mid 1E \mid 1F \mid F, \\ A \rightarrow 0B, & E \rightarrow 0D, \\ B \rightarrow 0D \mid 1C, & F \rightarrow \varepsilon \\ C \rightarrow 0B, & \end{array} \right\}.$$

$\mathcal{L}_{22} = \{w \in \Sigma^* = \{0, 1\}^* \mid |w| \geq 2 \text{ e os dois primeiros símbolos de } w \text{ são iguais aos dois últimos}\}$

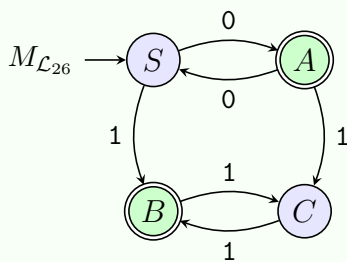
$\mathcal{L}_{23} = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ não começa com } 10, \text{ mas termina com } 10\}$

$\mathcal{L}_{24} = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ contém pelo menos um } 0 \text{ e pelo menos dois } 1\text{'s}\}$

$\mathcal{L}_{25} = \{w \in \Sigma^* = \{0, 1\}^* \mid w = 0u \text{ e } |w| \text{ é par ou } w = 1u' \text{ e } |u'| \text{ é par, com } u, u' \in \Sigma^*\}$

$\mathcal{L}_{26} = \{w \in \Sigma^* = \{0, 1\}^* \mid |w|_0 + |w|_1 = 2k + 1, k \in \mathbb{N} \text{ e } w \text{ não contém } 10\}$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{26}

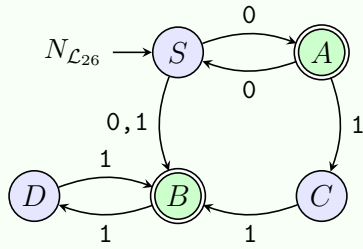


- Gramática G_1 que gera as cadeias de \mathcal{L}_{26} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, C, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l|l} S \rightarrow 0A \mid 1B, & B \rightarrow 1C \mid \varepsilon, \\ A \rightarrow 0S \mid 1C \mid \varepsilon, & C \rightarrow 1B \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{26} :



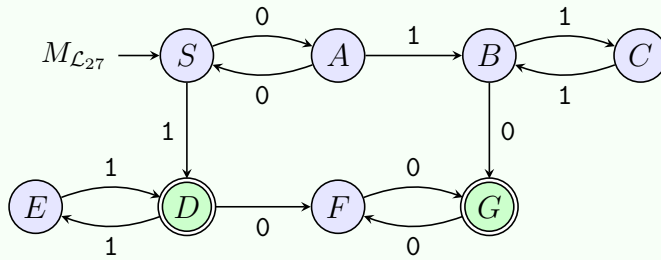
- Gramática G_2 que gera as cadeias de \mathcal{L}_{26} :

$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, D\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 0B \mid 1B, \\ A \rightarrow 0S \mid 1C \mid \varepsilon, \\ B \rightarrow 1D \mid \varepsilon, \end{array} \middle| \begin{array}{l} C \rightarrow 1B, \\ D \rightarrow 1B \end{array} \right\}.$$

$$\mathcal{L}_{27} = \{w \in \Sigma^* = \{0, 1\}^* \mid w = xyz, x, z \in \{0\}^*, y \in \{1\}^+; |x|_0 + |z|_0 = 2k, |y|_1 = 2k' + 1, k, k' \in \mathbb{N}\}$$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{27}

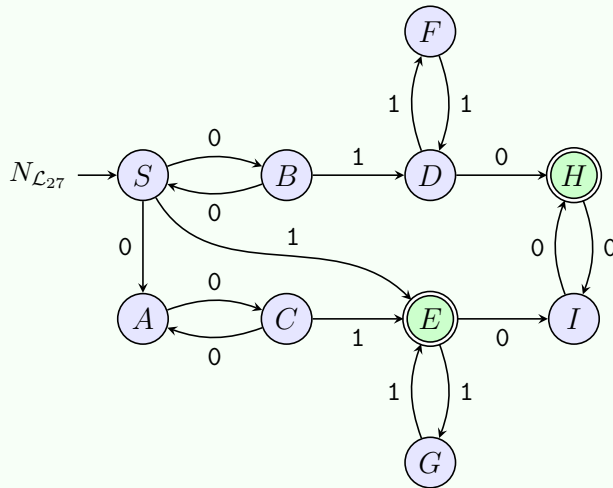


- Gramática G_1 que gera as cadeias de \mathcal{L}_{27} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, C, D, E, F, G, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1D, \\ A \rightarrow 0S \mid 1B, \\ B \rightarrow 0G \mid 1C, \\ C \rightarrow 1B, \end{array} \middle| \begin{array}{l} D \rightarrow 0F \mid 1E \mid \varepsilon, \\ E \rightarrow 1D, \\ F \rightarrow 0G, \\ G \rightarrow 0F \mid \varepsilon \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{27} :



- Gramática G_2 que gera as cadeias de \mathcal{L}_{27} :

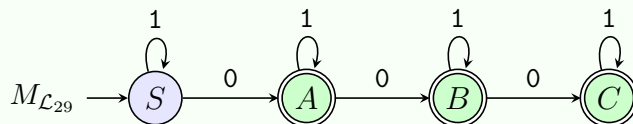
$G_1 = (V, \Sigma, P, S) = (\{A, B, C, D, E, F, G, H, I, S\}, \{0, 1\}, P, S)$, com:

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 0B \mid 1E, \\ A \rightarrow 0C, \\ B \rightarrow 0S \mid 1D, \\ C \rightarrow 0A \mid 1E, \end{array} \middle| \begin{array}{l} D \rightarrow 0H \mid 1F, \\ E \rightarrow 0I \mid 1G \mid \varepsilon, \\ F \rightarrow 1D, \\ G \rightarrow 1E, \end{array} \middle| \begin{array}{l} H \rightarrow 0I \mid \varepsilon \\ I \rightarrow 0H \end{array} \right\}.$$

$$\mathcal{L}_{28} = \{w \in \Sigma^* = \{0, 1\}^* \mid w = xycyz, c \in \Sigma, x, y, z \in \Sigma^*; |x| = 2k + 1, |z| = 2k', k, k' \in \mathbb{N}; |y| = 2\}$$

$$\mathcal{L}_{29} = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ contém uma, duas ou três ocorrências do símbolo } 0\}$$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{29}

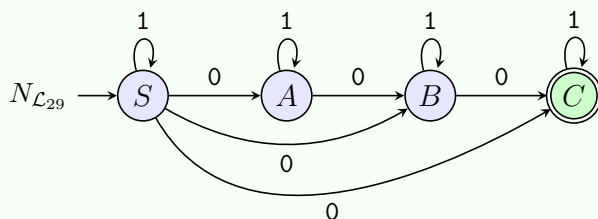


- Gramática G_1 que gera as cadeias de \mathcal{L}_{29} :

$G_1 = (V, \Sigma, P, S) = (\{A, B, C, S\}, \{0, 1\}, P, S)$, com:

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1S, \\ A \rightarrow 0B \mid 1A \mid \varepsilon, \end{array} \middle| \begin{array}{l} B \rightarrow 0C \mid 1B \mid \varepsilon, \\ C \rightarrow 1C \mid \varepsilon \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{29} :





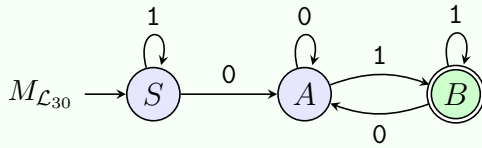
- Gramática G_2 que gera as cadeias de \mathcal{L}_{29} :

$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 0B \mid 0C \mid 1S, \\ A \rightarrow 0B \mid 1A, \end{array} \middle| \begin{array}{l} B \rightarrow 0C \mid 1B, \\ C \rightarrow 1C \mid \varepsilon \end{array} \right\}.$$

$$\mathcal{L}_{30} = \{w \in \Sigma^* = \{0, 1\}^* \mid w = u01^n, u \in \Sigma^*, n \in \mathbb{N}^+\}$$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{30}

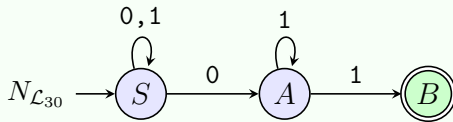


- Gramática G_1 que gera as cadeias de \mathcal{L}_{30} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1S, \\ A \rightarrow 0A \mid 1B, \\ B \rightarrow 0A \mid 1B \mid \varepsilon \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{30} :



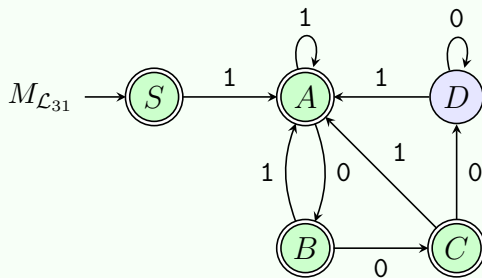
- Gramática G_2 que gera as cadeias de \mathcal{L}_{30} :

$$G_2 = (V, \Sigma, P, S) = (\{A, B, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 0S \mid 1S, \\ A \rightarrow 1A \mid 1B, \\ B \rightarrow \varepsilon \end{array} \right\}.$$

$$\mathcal{L}_{31} = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ não começa com } 0 \text{ e não termina com } 000\}$$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{31}



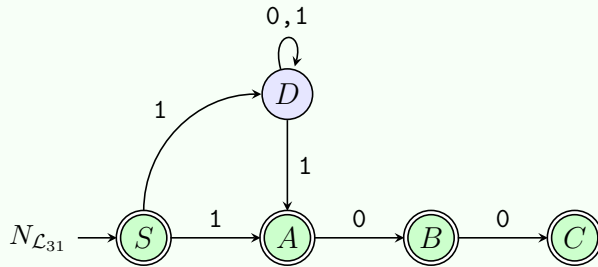


- Gramática G_1 que gera as cadeias de \mathcal{L}_{31} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, C, D, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 1A \mid \varepsilon, \\ A \rightarrow 0B \mid 1A \mid \varepsilon, \\ B \rightarrow 0C \mid 1A \mid \varepsilon, \end{array} \middle| \begin{array}{l} C \rightarrow 0D \mid 1A \mid \varepsilon, \\ D \rightarrow 0D \mid 1A, \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{31} :



- Gramática G_2 que gera as cadeias de \mathcal{L}_{31} :

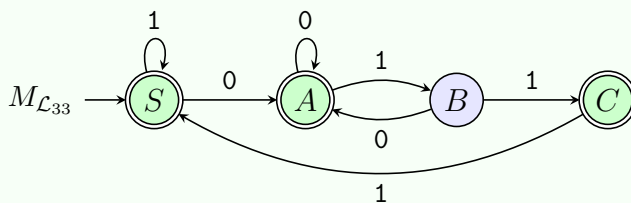
$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, D, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 1A \mid 1D \mid \varepsilon, \\ A \rightarrow 0B \mid \varepsilon, \\ B \rightarrow 0C \mid \varepsilon, \end{array} \middle| \begin{array}{l} C \rightarrow \varepsilon, \\ D \rightarrow 0D \mid 1A \mid 1D \end{array} \right\}.$$

$$\mathcal{L}_{32} = \{w \in \Sigma^* = \{0, 1\}^* \mid w = uc, u \in \Sigma^*, c \in \Sigma, |u|_c \leq 2\}$$

$$\mathcal{L}_{33} = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ não contém } 0110 \text{ e não termina com } 01\}$$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{33}

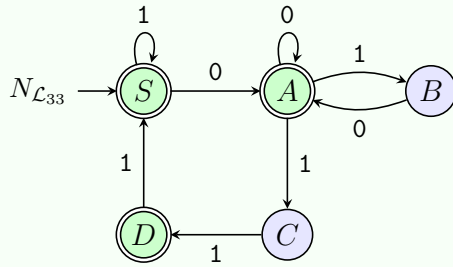


- Gramática G_1 que gera as cadeias de \mathcal{L}_{33} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, C, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1S \mid \varepsilon, \\ A \rightarrow 0A \mid 1B \mid \varepsilon, \end{array} \middle| \begin{array}{l} B \rightarrow 0A \mid 1C, \\ C \rightarrow 1S \mid \varepsilon \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{33} :



- Gramática G_2 que gera as cadeias de \mathcal{L}_{33} :

$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, D, S\}, \{0, 1\}, P, S), \text{ com:}$$

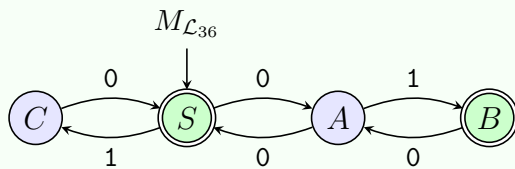
$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1S \mid \varepsilon, \\ A \rightarrow 0A \mid 1B \mid 1C \mid \varepsilon, \\ B \rightarrow 0A, \end{array} \left| \begin{array}{l} C \rightarrow 1D, \\ D \rightarrow 1S \mid \varepsilon, \end{array} \right. \right\}.$$

$\mathcal{L}_{34} = \{w \in \Sigma^* = \{0, 1\}^* \mid |w| \geq 4, \text{ começa com 0 e contém pelo menos um 1 do terceiro ao penúltimo símbolo}\}$

$\mathcal{L}_{35} = \{w \in \Sigma^* = \{0, 1\}^* \mid |w| = 2k + 1, k \in \mathbb{N}, w \text{ termina com 1 e contém pelo menos mais um 1}\}$

$\mathcal{L}_{36} = \{w \in \Sigma^* = \{0, 1\}^* \mid |w| = 2k, k \in \mathbb{N}, w \text{ não contém 11}\}$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{36}

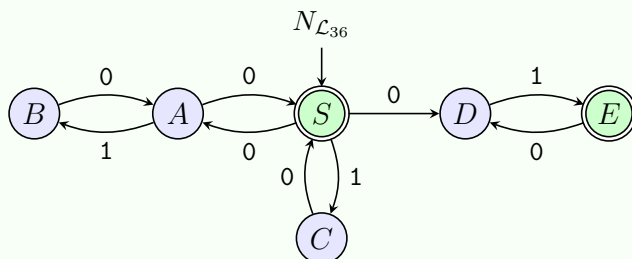


- Gramática G_1 que gera as cadeias de \mathcal{L}_{36} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1C \mid \varepsilon, \\ A \rightarrow 0S \mid 1B, \end{array} \left| \begin{array}{l} B \rightarrow 0A \mid \varepsilon \\ C \rightarrow 0S \end{array} \right. \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{36} :





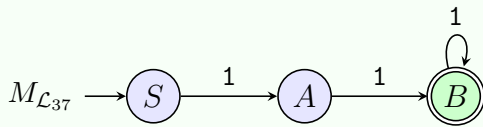
- Gramática G_2 que gera as cadeias de \mathcal{L}_{36} :

$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, D, E, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 0D \mid 1C \mid \varepsilon, \\ A \rightarrow 0S \mid 1B, \\ B \rightarrow 0A, \end{array} \left| \begin{array}{l} C \rightarrow 0S, \\ D \rightarrow 1E, \\ E \rightarrow 0D \mid \varepsilon \end{array} \right. \right\}.$$

$\mathcal{L}_{37} = \{w \in \Sigma^* = \{0, 1\}^* \mid w = u11, u \in \Sigma^* \text{ e todo } 0 \text{ em } u \text{ é seguido de um par de símbolos distintos}\}$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{37}

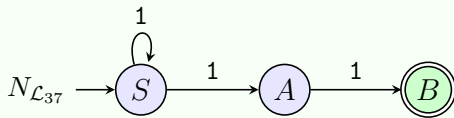


- Gramática G_1 que gera as cadeias de \mathcal{L}_{37} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 1A, \\ A \rightarrow 1B, \\ B \rightarrow 1B \mid \varepsilon \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{37} :



- Gramática G_2 que gera as cadeias de \mathcal{L}_{37} :

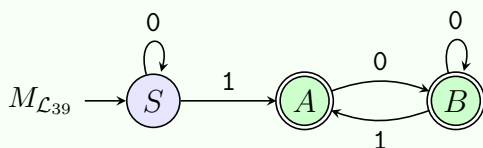
$$G_2 = (V, \Sigma, P, S) = (\{A, B, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 1A \mid 1S, \\ A \rightarrow 1B, \\ B \rightarrow \varepsilon \end{array} \right\}.$$

$\mathcal{L}_{38} = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ contém os símbolos } 0 \text{ e } 1, \text{ mas não contém } 00\}$

$\mathcal{L}_{39} = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ contém pelo menos um } 1, \text{ mas não contém } 11\}$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{39}



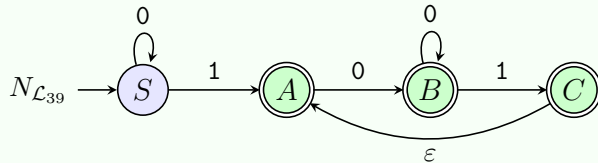


- Gramática G_1 que gera as cadeias de \mathcal{L}_{39} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0S \mid 1A, \\ A \rightarrow 0B \mid \varepsilon, \\ B \rightarrow 0B \mid 1A \mid \varepsilon \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{39} :



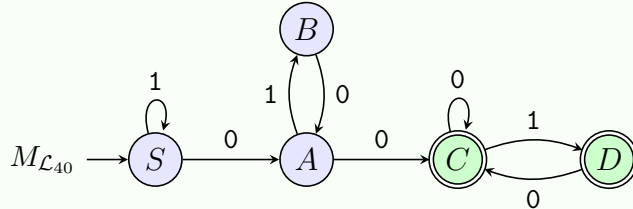
- Gramática G_2 que gera as cadeias de \mathcal{L}_{39} :

$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0S \mid 1A, \\ A \rightarrow 0B \mid \varepsilon, \\ B \rightarrow 0B \mid 1C \mid \varepsilon \\ C \rightarrow A \mid \varepsilon \end{array} \right\}.$$

$\mathcal{L}_{40} = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ contém } 00, \text{ mas não contém } 011\}$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{40}

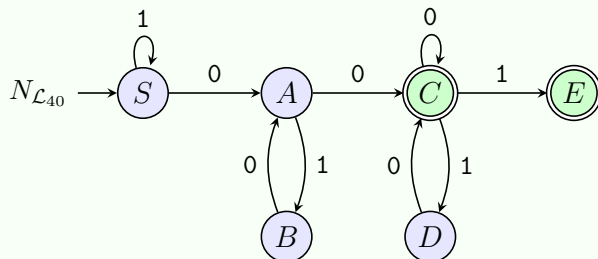


- Gramática G_1 que gera as cadeias de \mathcal{L}_{40} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, C, D, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1S, \\ A \rightarrow 0C \mid 1B, \\ B \rightarrow 0A, \\ C \rightarrow 0C \mid 1D \mid \varepsilon \\ D \rightarrow 0C \mid \varepsilon \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{40} :





- Gramática G_2 que gera as cadeias de \mathcal{L}_{40} :

$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, D, E, S\}, \{0, 1\}, P, S), \text{ com:}$$

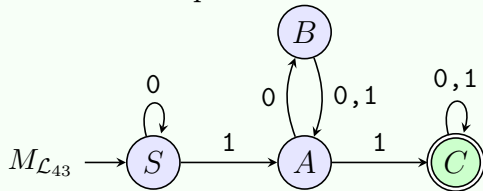
$$P = \left\{ \begin{array}{l|l} S \rightarrow 0A \mid 1S, & C \rightarrow 0C \mid 1D \mid 1E \mid \varepsilon \\ A \rightarrow 0C \mid 1B, & D \rightarrow 0C, \\ B \rightarrow 0A, & E \rightarrow \varepsilon \end{array} \right\}.$$

$$\mathcal{L}_{41} = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ contém pelo menos um } 00, \text{ mas não contém } 11\}$$

$$\mathcal{L}_{42} = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ começa com } 0 \text{ e contém } 010 \text{ ou } w \text{ começa com } 1 \text{ e contém } 101\}$$

$$\mathcal{L}_{43} = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ contém dois } 1\text{'s separados por uma quantidade par de símbolos}\}$$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{43}

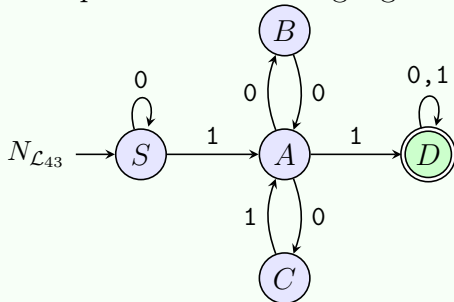


- Gramática G_1 que gera as cadeias de \mathcal{L}_{43} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, C, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l|l} S \rightarrow 0S \mid 1A, & B \rightarrow 0A \mid 1A, \\ A \rightarrow 0B \mid 1C, & C \rightarrow 0C \mid 1C \mid \varepsilon \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{43} :



- Gramática G_2 que gera as cadeias de \mathcal{L}_{43} :

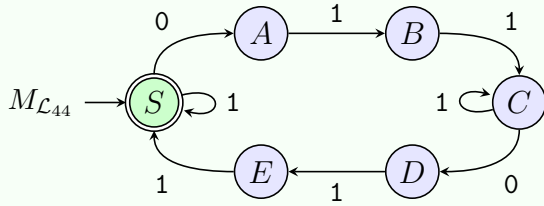
$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, D, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l|l} S \rightarrow 0S \mid 1A, & C \rightarrow 1A, \\ A \rightarrow 0B \mid 0C \mid 1D, & D \rightarrow 0D \mid 1D \mid \varepsilon \\ B \rightarrow 0A, & \end{array} \right\}.$$



$\mathcal{L}_{44} = \{w \in \Sigma^* = \{0,1\}^* \mid |w|_0 = 2k, k \in \mathbb{N}, \text{ e cada } 0 \text{ é seguido de pelo menos dois } 1\text{'s consecutivos}\}$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{44}

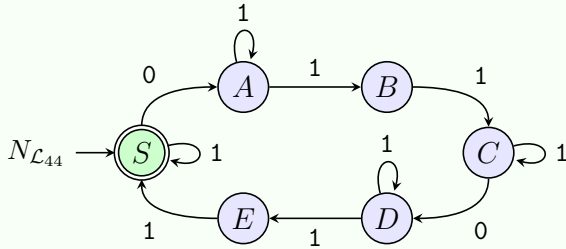


- Gramática G_1 que gera as cadeias de \mathcal{L}_{44} :

$G_1 = (V, \Sigma, P, S) = (\{A, B, C, D, E, S\}, \{0, 1\}, P, S)$, com:

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1S \mid \varepsilon, \\ A \rightarrow 1B, \\ B \rightarrow 1C, \end{array} \middle| \begin{array}{l} C \rightarrow 0D \mid 1C, \\ D \rightarrow 1E, \\ E \rightarrow 1S \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{44} :



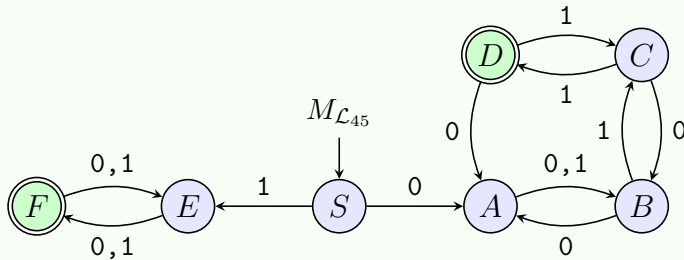
- Gramática G_2 que gera as cadeias de \mathcal{L}_{44} :

$G_2 = (V, \Sigma, P, S) = (\{A, B, C, D, E, S\}, \{0, 1\}, P, S)$, com:

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1S \mid \varepsilon, \\ A \rightarrow 1A \mid 1B, \\ B \rightarrow 1C, \end{array} \middle| \begin{array}{l} C \rightarrow 0D \mid 1C, \\ D \rightarrow 1D \mid 1E, \\ E \rightarrow 1S \end{array} \right\}.$$

$\mathcal{L}_{45} = \{w \in \Sigma^* = \{0,1\}^* \mid |w| = 2k, k \in \mathbb{N}, \text{ e } w \text{ começa com } 1 \text{ ou termina com } 11\}$

- DFA mínimo que reconhece as cadeias de \mathcal{L}_{45}



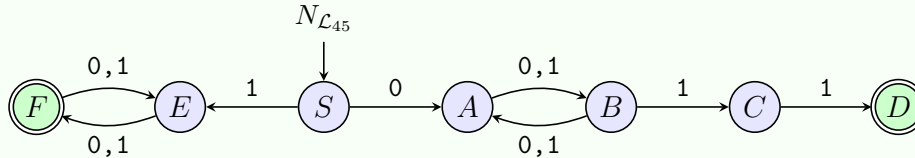


- Gramática G_1 que gera as cadeias de \mathcal{L}_{45} :

$$G_1 = (V, \Sigma, P, S) = (\{A, B, C, D, E, F, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1E, \\ A \rightarrow 0B \mid 1B, \\ B \rightarrow 0A \mid 1C, \end{array} \middle| \begin{array}{l} C \rightarrow 0B \mid 1D, \\ D \rightarrow 0A \mid 1C \mid \varepsilon, \\ \end{array} \middle| \begin{array}{l} E \rightarrow 0F \mid 1F \\ F \rightarrow 0E \mid 1E \mid \varepsilon \end{array} \right\}.$$

- NFA que reconhece a linguagem \mathcal{L}_{45} :



- Gramática G_2 que gera as cadeias de \mathcal{L}_{45} :

$$G_2 = (V, \Sigma, P, S) = (\{A, B, C, D, E, F, S\}, \{0, 1\}, P, S), \text{ com:}$$

$$P = \left\{ \begin{array}{l} S \rightarrow 0A \mid 1E, \\ A \rightarrow 0B \mid 1B, \\ B \rightarrow 0A \mid 1A \mid 1C, \end{array} \middle| \begin{array}{l} C \rightarrow 1D, \\ D \rightarrow \varepsilon, \end{array} \middle| \begin{array}{l} E \rightarrow 0F \mid 1F \\ F \rightarrow 0E \mid 1E \mid \varepsilon \end{array} \right\}.$$

$$\mathcal{L}_{46} = \{w \in \Sigma^* = \{0, 1\}^* \mid w \text{ é diferente de } 0, 00, 1, 11 \text{ e } 010\}$$

$$\mathcal{L}_{47} = \{w \in \Sigma^* = \{0, 1\}^* \mid |w|_0 = 2k \text{ e } |w|_1 = 3k', \ k, k' \in \mathbb{N}\}$$