

Lesson 1. Fundamentals of Graphs.

1. Definition and Graph Terminology.
2. Special Types of Graphs.
3. Vertex Degree.
4. Paths and Connectedness.
5. Matrix representation.

4. Paths and Connectivity

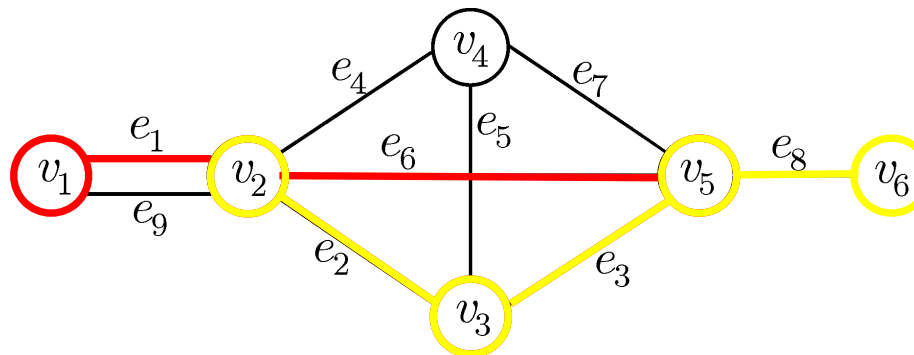
DEFINITIONS:

- Let x, y be (not necessarily distinct) vertices in an undirected graph $G = (V, E)$. An x - y **walk** in G is a (loop-free) finite alternating sequence

$$x = v_1, e_1, v_2, e_2, v_2, e_2, \dots, v_{n-1}, e_{n-1}, v_n = y$$

of vertices and edges from G , starting at vertex x and ending at vertex y and involving the $n-1$ edges $e_i = \{v_i, v_{i+1}\}$, where $1 \leq i \leq n-2$.

EXAMPLE:



$$C = \overline{v_1 e_1 v_2 e_2 v_3 e_3 v_5 e_6 v_2 e_2 v_3 e_3 v_5 e_8 v_6}$$

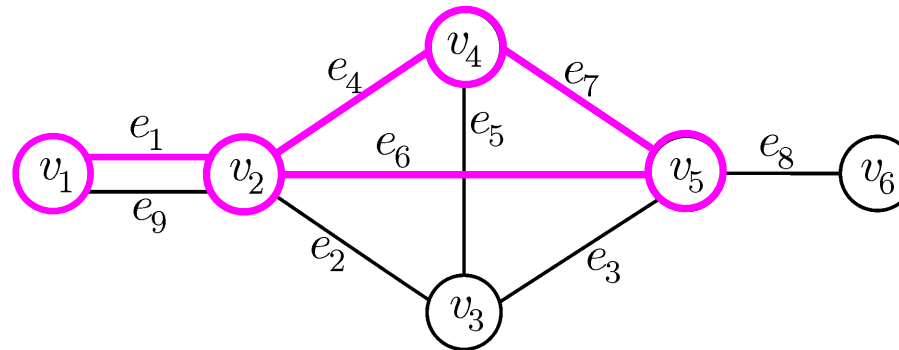
- The **length** of a walk is **the number of edges** in the walk. (When $n = 0$, there are no edges, $x = y$, and the walk is called trivial. These walks are not considered very much in our work.)

EXAMPLE: The length of the walk **C** is 7.

4. Paths and Connectivity

3. Any x - y walk where $x = y$ (and $n > 1$) is called a **closed walk**. Otherwise the walk is called open.

EXAMPLE:



The walk $C_4 = v_2 e_6 v_5 e_7 v_4 e_4 v_2$ is a closed walk.

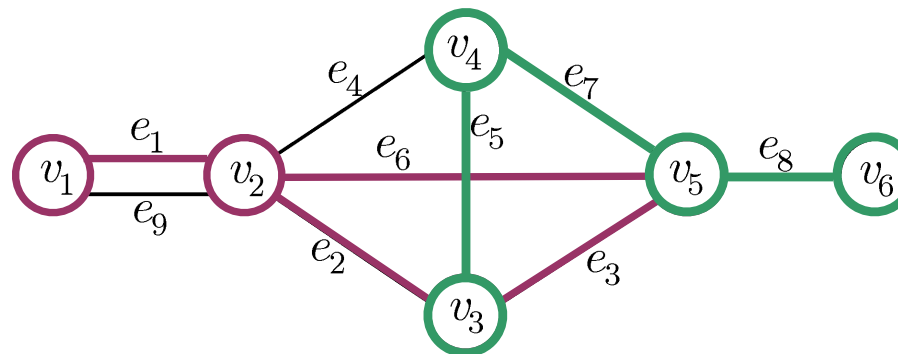
4. Paths and Connectivity

Note that a walk may repeat both vertices and edges.

Consider any x-y walk in an undirected graph $G = (V, E)$.

4. If **no edge** in the x-y walk is repeated, then the walk is called an x-y **trail**.
5. If **no vertex** of the x-y walk occurs more than once, then the walk is called an x-y **path**.
6. When $x = y$, the term **cycle** is used to describe such a closed path.

EXAMPLE:

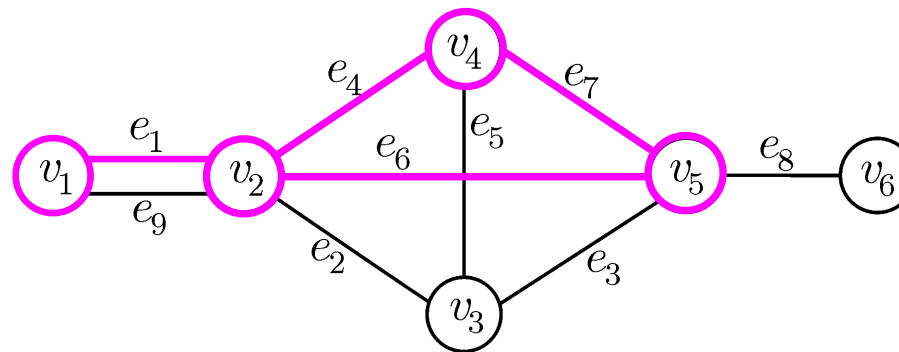


The walk $\mathbf{C}_2 = v_1 e_1 v_2 e_2 v_3 e_3 v_5 e_6 v_2$ is a trail of length 4.

The walk $\mathbf{C}_3 = v_6 e_8 v_5 e_7 v_4 e_5 v_3$ is a path.

4. Paths and Connectivity

EXAMPLE:



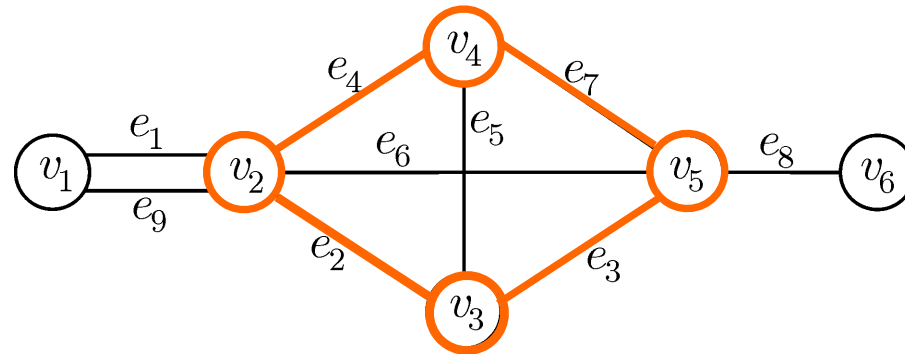
The walk $\mathbf{C}_4 = \mathbf{v}_2 \mathbf{e}_6 \mathbf{v}_5 \mathbf{e}_7 \mathbf{v}_4 \mathbf{e}_4 \mathbf{v}_2 \mathbf{e}_1 \mathbf{v}_1 \mathbf{e}_1 \mathbf{v}_2$ is a closed walk.

\mathbf{e}_1 is repeated, then \mathbf{C}_4 is not a trail.

\mathbf{v}_2 is repeated, then \mathbf{C}_4 is not a path.

4. Paths and Connectivity

EXAMPLE:



The walk $\mathbf{C} = v_2 e_4 v_4 e_7 v_5 e_3 v_3 e_2 v_2$ is a cycle.

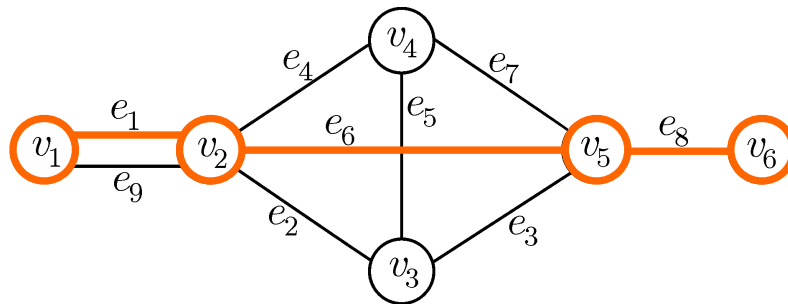
Remark.-

These concepts are defined similarly for directed graphs with the consideration that the **directions of the arcs must agree with the direction of the path or walk**. In some books directed cycles are called **circuits**.

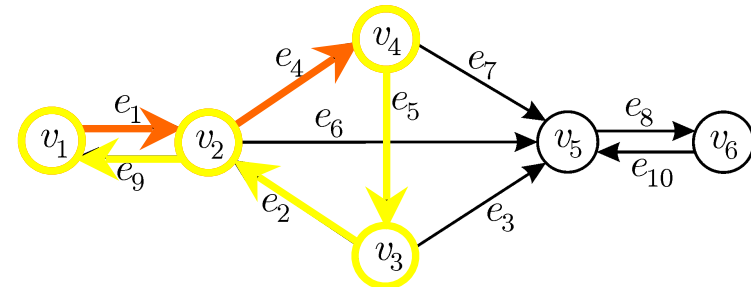
4. Paths and Connectivity

7. - Let $G=(V,A)$ be an undirected and let u, v two vertices of G .
 The **Connectedness** relation in V is defined in the following way:
 u and v are **connected** if there is a path from u to v .
- Let $G=(V,A)$ be an **directed** and let u, v two vertices of G .
 The **Connectedness** relation in V is defined in the following way:
 u and v are **connected** if there is: a path from u to v , and
 a path from v to u .

EXAMPLE:



Vertices v_1 and v_6 are connected through the path $C=v_1e_1v_2e_6v_5e_8v_6$.
 Every pair of distinct vertices are connected.



Vertices v_1 and v_4 are connected through the paths:

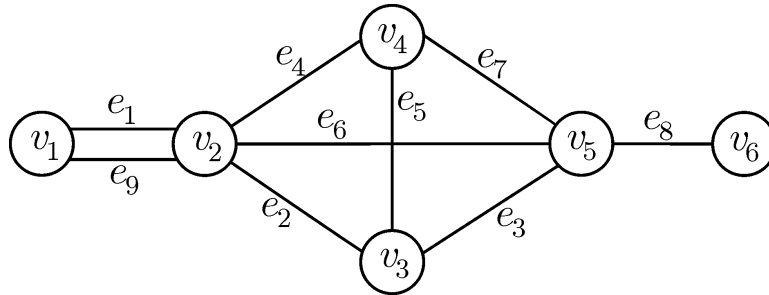
$$C_1=v_1e_1v_2e_4v_4$$

$$C_2=v_4e_5v_3e_2v_2e_9v_1$$

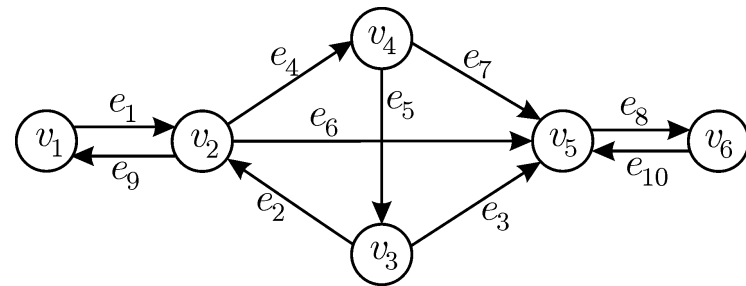
Vertices v_4 and v_5 are not connected.

4. Paths and Connectivity

8. A graph (undirected or directed) is called **connected** if every pair of distinct vertices of the graph are connected. A graph that is not connected is called **disconnected**.

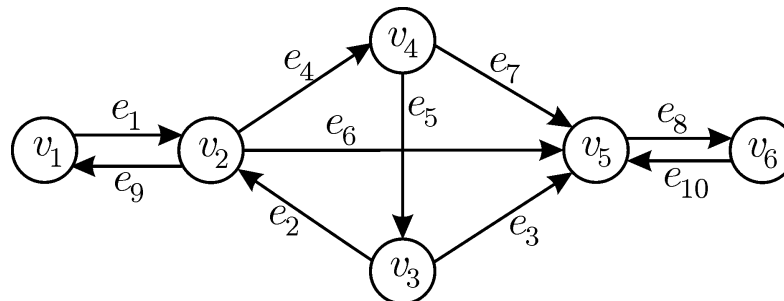


Connected graph



Disconneted graph

9. A directed graph is **weakly connected** if its associated undirected graph is connected.



4. Paths and Connectivity

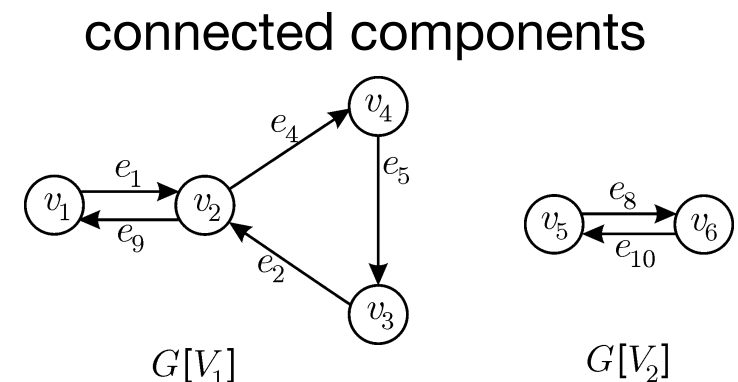
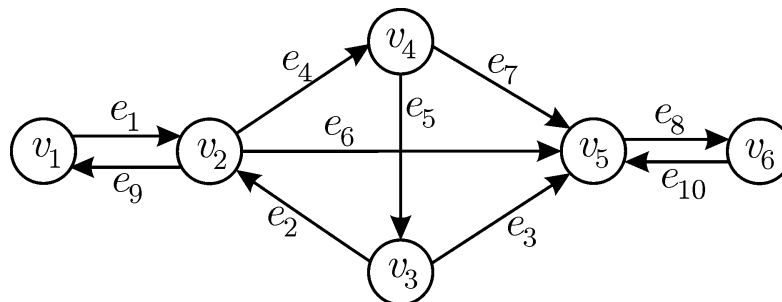
THEOREM

The connectedness relation is an equivalence relation and therefore determines a partition on the set of vertices. The elements of this partition are called **connected components** of the graph.

THEOREM

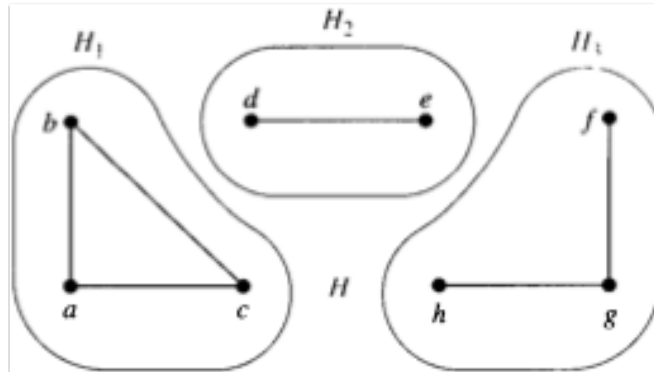
A graph is connected if and only if it has only one connected component.

EXAMPLE:



4. Paths and Connectivity

EXAMPLE:



Graph H and its Connected Components H_1 , H_2 , and H_3

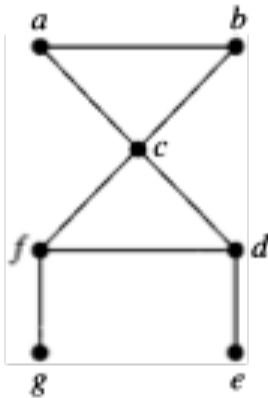
$$C_1 = \{a, b, c\}$$

$$C_2 = \{d, e\}$$

$$C_3 = \{f, g, h\}$$

4. Paths and Connectivity

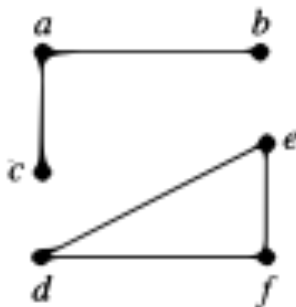
EXAMPLE:



G

- The graph G in Figure is connected, because for every pair of distinct vertices there is a path between them.
- There is only one connected component.

$$C=\{a,b,c,d,e,f,g\}$$

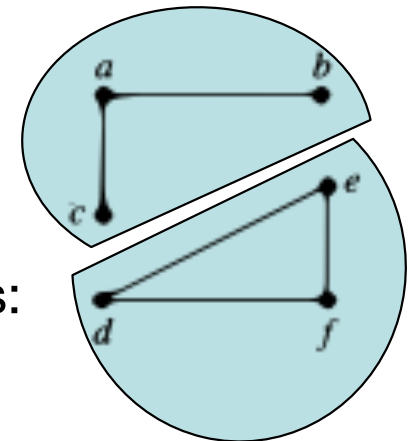


H

- The graph H is not connected. For instance, there is no path in H between vertices a and d.
- There are 2 connected components:

$$C_1=\{a,b,c\}$$

$$C_2=\{d,e,f\}$$



H

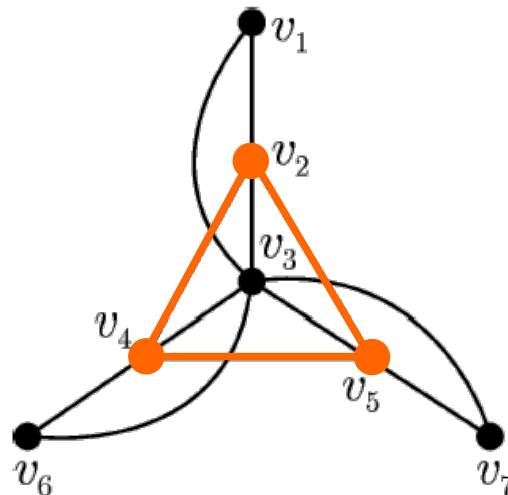
4. Paths and Connectivity

THEOREM (Undirected Graphs)

A graph is bipartite if and only if it does not contain a cycle of odd length.

EXAMPLE:

The following graph is NOT bipartite because it contains an odd cycle: $v_2v_4v_5$



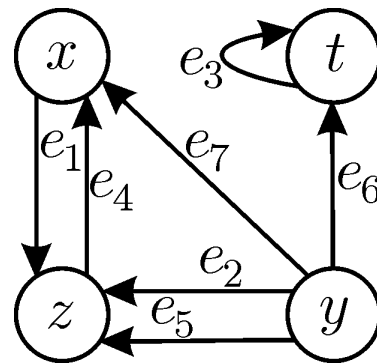
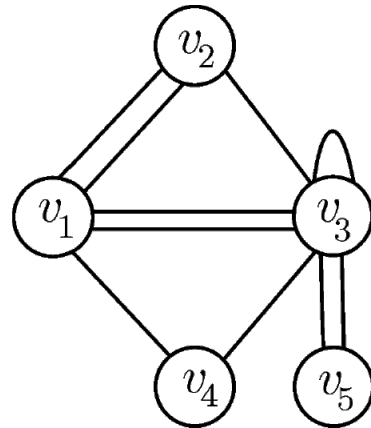
5. Matrix representation

DEFINITION:

Let G be a graph (directed or undirected) with n vertices $\{v_i\}_{i=1}^n$. The **adjacency matrix** A of G , with respect to this listing of the vertices, is the $n \times n$ matrix $A = [a_{ij}]$ where a_{ij} is the number of edges (arcs) from v_i to v_j . If G is undirected the loops are computed twice.

EXAMPLE:

$$A = \begin{bmatrix} 0 & 2 & 2 & 1 & 0 \\ 2 & 0 & 1 & 0 & 0 \\ 2 & 1 & 2 & 1 & 2 \\ 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 \end{bmatrix}$$



$$A = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 1 & 0 & 2 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \begin{matrix} x \\ y \\ z \\ t \end{matrix}$$

Remarks.

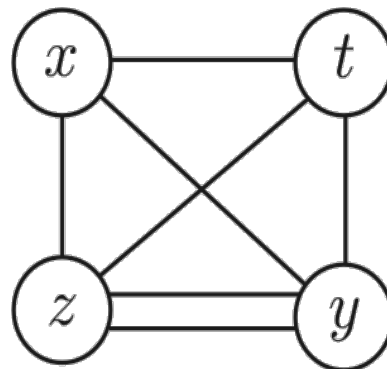
- If G is undirected then A is symmetric.
- If G is undirected, and the vertex v_i has a loop, then $a_{ii}=2$. If G is directed then $a_{ii}=1$.

5. Matrix representation

PROPERTIES OF THE ADJACENCY MATRIX:

1. Let G be an undirected graph with adjacency matrix A . Then, the sum of the elements of the row i (or column i) is equal to the degree of the vertex v_i .

EXAMPLE:



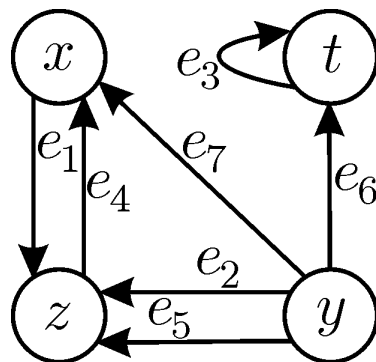
$$A = \begin{matrix} & \begin{matrix} x & y & z & t \end{matrix} \\ \begin{matrix} x \\ y \\ z \\ t \end{matrix} & \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 2 & 1 \\ 1 & 2 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix} \end{matrix} \begin{matrix} x \\ y \\ z \\ t \end{matrix} = 3$$

The degree of the vertex x is 3.

5. Matrix representation

2. Let G be a directed graph with adjacency matrix A . Then, the sum of the elements of the **row** i is equal to the **outdegree** of the vertex v_i , and the sum of the elements of the **column** j is equal to the **indegree** of the vertex v_j .

EXAMPLE:



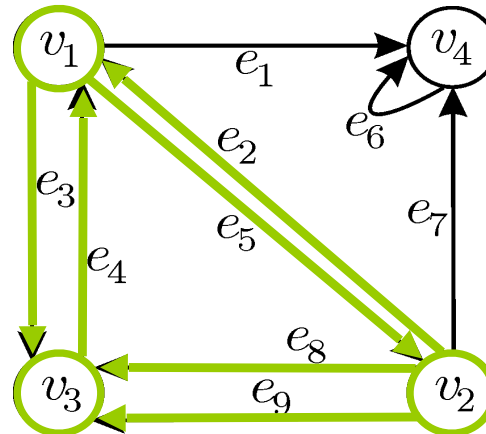
$$A = \begin{matrix} & \begin{matrix} x & y & z & t \end{matrix} \\ \begin{matrix} x \\ y \\ z \\ t \end{matrix} & \begin{bmatrix} 0 & 0 & 1 & 0 \\ 1 & 0 & 2 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \end{matrix} = 1 \quad d_{\text{out}}(x)=1$$

\downarrow
3
 $d_{\text{in}}(z)=3$

5. Matrix representation

3. Let G be a graph with adjacency matrix A . Then, the element (i,j) of the matrix A^r , $r \geq 1$, is equal to the number of walks from v_i to v_j of length r .

EXAMPLE:



$$A = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 2 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

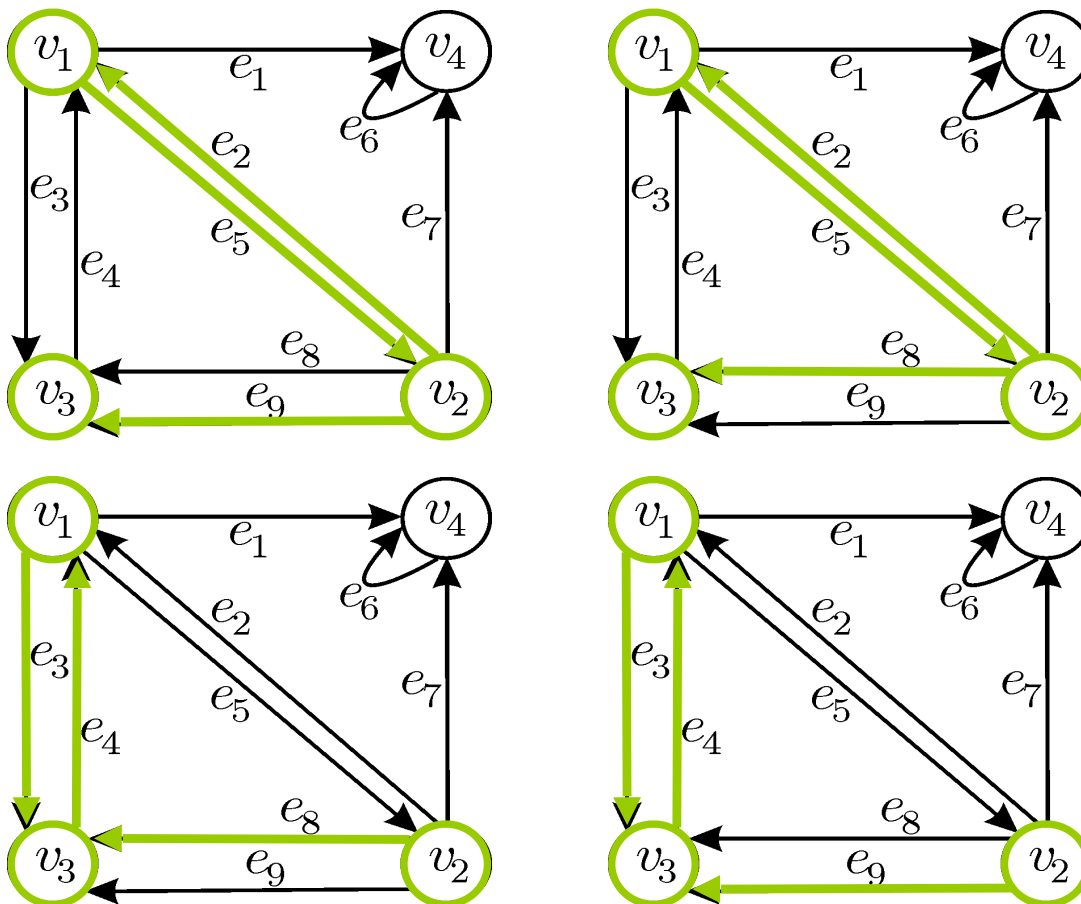
$$A^3 = \begin{bmatrix} 2 & 2 & 2 & 4 \\ 2 & 2 & 4 & 5 \\ 2 & 0 & 2 & 2 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

The number of walks of length 3 from v_2 to v_3 is 4.

5. Matrix representation

3. Let G be a graph with adjacency matrix A . Then, the element (i,j) of the matrix A^r , $r \geq 1$, is equal to the number of walks from v_i to v_j of length r .

EXAMPLE:



The number of walks of length 3 from v_2 to v_3 is 4.

$$A = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 2 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$$A^3 = \begin{bmatrix} 2 & 2 & 2 & 4 \\ 2 & 2 & 4 & 5 \\ 2 & 0 & 2 & 2 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

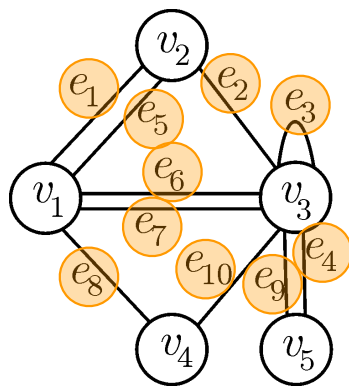
5. Matrix representation

DEFINITIONS:

- Let $G = (V, A)$ be an undirected graph with n vertices and m edges where $A = \{a_i\}_{i=1}^m$, and $V = \{v_i\}_{i=1}^n$. The **incidence matrix** of G is the $n \times m$ matrix

$$M = [m_{ij}] / m_{ij} = \begin{cases} 0 & \text{if } v_i \text{ is not incident with } a_j \\ 1 & \text{if } v_i \text{ is incident with } a_j \\ 2 & \text{if } a_j \text{ is a loop in } v_i \end{cases}$$

EXAMPLE:



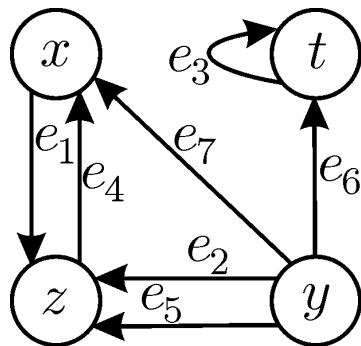
$$M = \begin{matrix} & e_1 & e_2 & e_3 & e_4 & e_5 & e_6 & e_7 & e_8 & e_9 & e_{10} \\ \begin{bmatrix} 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 2 & 1 & 0 & 1 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix} & \begin{matrix} v_1 \\ v_2 \\ v_3 \\ v_4 \\ v_5 \end{matrix} \end{matrix}$$

5. Matrix representation

2. Let $G = (V, A)$ be an directed graph with n vertices and m arcs where $V = \{v_i\}_{i=1}^n$, and $A = \{a_i\}_{i=1}^m$. The **incidence matrix** of G is the $n \times m$ matrix

$$M = [m_{ij}] / m_{ij} = \begin{cases} 0 & \text{if } v_i \text{ is not incident with } a_j \\ 1 & \text{if } v_i \text{ is the origin vertex of } a_j \\ -1 & \text{if } v_i \text{ is the terminating vertex of } a_j \\ 2 & \text{if } a_j \text{ is a loop in } v_i \end{cases}$$

EXAMPLE:



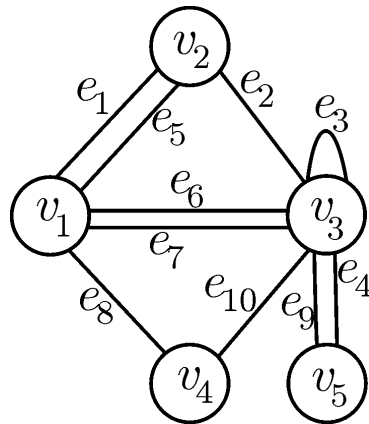
$$M = \begin{matrix} & \begin{matrix} e_1 & e_2 & e_3 & e_4 & e_5 & e_6 & e_7 \end{matrix} \\ \begin{bmatrix} 1 & 0 & 0 & -1 & 0 & 0 & -1 \\ 0 & 1 & 0 & 0 & 1 & 1 & 1 \\ -1 & -1 & 0 & 1 & -1 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 & -1 & 0 \end{bmatrix} & \begin{matrix} x \\ y \\ z \\ t \end{matrix} \end{matrix}$$

5. Matrix representation

PROPERTIES OF THE INCIDENCE MATRIX:

1. Let G be an undirected graph with incidence matrix M . Then, the sum of the elements of the row i of M is equal to the degree of the vertex v_i .

EXAMPLE:



$$M = \begin{bmatrix} 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 2 & 1 & 0 & 1 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix} \begin{matrix} = 5 \\ = 3 \\ = 8 \\ = 2 \\ = 2 \end{matrix}$$

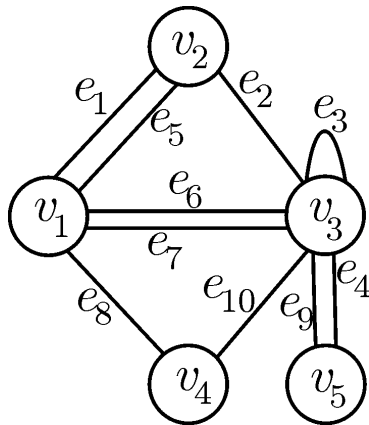
The degree of the vertices are:

$$d(v_1)=5, d(v_2)=3, d(v_3)=8, d(v_4)=2, d(v_5)=2$$

5. Matrix representation

2. Let G be an undirected graph with incidence matrix M . Then, the sum of the elements of each column of M is equal to 2.

EXAMPLE:



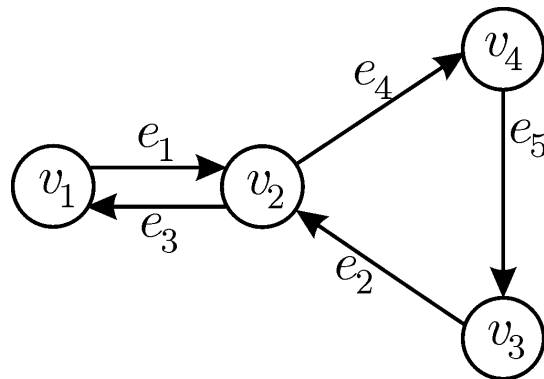
$$M = \begin{bmatrix} 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 2 & 1 & 0 & 1 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

2 2 2 2 2 2 2 2 2 2

5. Matrix representation

3. Let G be a directed graph without loops. Then, the sum of the elements of each column of the incidence matrix is equal to 0.

EXAMPLE:



$$M = \begin{bmatrix} 1 & 0 & -1 & 0 & 0 \\ -1 & -1 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 & -1 \\ 0 & 0 & 0 & -1 & 1 \end{bmatrix}$$

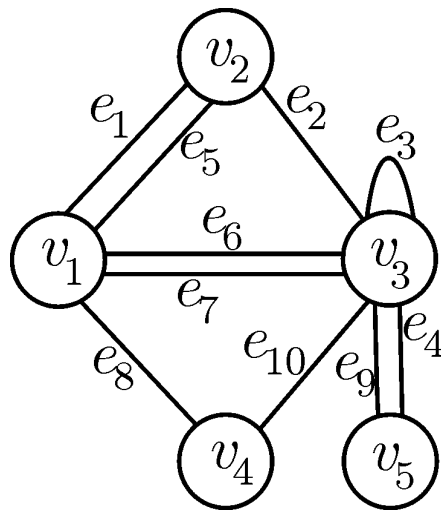
↓ ↓ ↓ ↓ ↓
0 0 0 0 0

5. Matrix representation

DEFINITIONS:

1. Let G be an undirected graph. Another way to represent a graph is to use **incidence lists**, which list the edges that are incident to each vertex of the graph.

EXAMPLE:



Incidence List

$v_1: e_1, e_5, e_6, e_7, e_8$

$v_2: e_1, e_2, e_5$

$v_3: e_2, e_3, e_4, e_6, e_7, e_9, e_{10}$

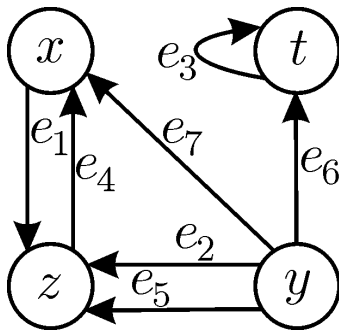
$v_4: e_7, e_8$

$v_5: e_3, e_4, e_9$

5. Matrix representation

2. Let G be a directed graph. The **List of Outward Arcs** of G is a table that lists, for each vertex v , the outward arcs from v . The **List of Inward Arcs** of G is the table that lists, for each vertex v , the inward arcs from v .

EXAMPLE:



Outward	Arcs	Inward	Arcs
x :	e_1	x :	e_4, e_7
y :	e_2, e_5, e_6, e_7	y :	
z :	e_4	z :	e_1, e_2, e_5
t :	e_3	t :	e_3, e_6