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Zero-sum stochastic games (1)

A zero-sum stochastic game is a 5-tuple (Ω, I, J, g, P) , where:

- Ω is a non-empty set of states;
- I is a non-empty set of actions of player 1;
- J is a non-empty set of actions of player 2;
- $g : I \times J \times \Omega \rightarrow \mathbb{R}$ is a payoff function of player 1;
- $P : I \times J \times \Omega \rightarrow \Delta(\Omega)$ is a transition probability function.

We assume that I, J, Ω are finite.

$\Delta(\Omega) :=$ the set of probability measures on Ω .

Zero-sum stochastic games (2)

A stochastic game (Ω, I, J, g, P) proceeds in stages as follows. At each stage n :

1. The players observe the current state ω_n ;
2. Players choose their mixed actions, $x_n \in \Delta(I)$ and $y_n \in \Delta(J)$;
3. Pure actions $i_n \in I$ and $j_n \in J$ are chosen according to $x_n \in \Delta(I)$ and $y_n \in \Delta(J)$;
4. Player 1 obtains a payoff $g_n = g(i_n, j_n, \omega_n)$, while player 2 obtains payoff $-g_n$;
5. The new state ω_{n+1} is chosen according to the probability law $P(i_n, j_n, \omega_n)$.

The above description of the game is known to the players.

Strategies

- Strategies σ, τ of players consist in choosing at each stage a mixed action;
- The players can take into account the previous actions of players, as well as the current and previous states.

λ -discounted game Γ^λ

- What is the goal of the players?
- Total payoff: $E_{\sigma,\tau}^\omega \left(\lambda \sum_{i=1}^{\infty} (1-\lambda)^{i-1} g_i \right)$;
- Depends on $\lambda \in (0, 1)$, initial state ω , and strategies of the players;
- Player 1 wants to maximize total payoff, while player 2 wants to minimize it;
- Value $v_\lambda : \Omega \rightarrow \mathbb{R}$:

$$\begin{aligned} v_\lambda(\omega) &= \sup_{\sigma} \inf_{\tau} E_{\sigma,\tau}^\omega \left(\lambda \sum_{i=1}^{\infty} (1-\lambda)^{i-1} g_i \right) = \\ &= \inf_{\tau} \sup_{\sigma} E_{\sigma,\tau}^\omega \left(\lambda \sum_{i=1}^{\infty} (1-\lambda)^{i-1} g_i \right). \end{aligned}$$

Limit of λ -discounted game Γ^λ

- $v_\lambda(\omega) = \sup_{\sigma} \inf_{\tau} E_{\sigma, \tau}^{\omega} \left(\lambda \sum_{i=1}^{\infty} (1 - \lambda)^{i-1} g_i \right);$
- One can ask: what happens if players become more and more patient? I.e., players are willing to wait a lot to obtain a big payoff;
- Mathematically, it means that $\lambda \rightarrow 0$;
- Thus, one is interested in the uniform (in ω) limit $\lim_{\lambda \rightarrow 0} v_\lambda(\omega)$.

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Kernel

- Kernel $Q : I \times J \times \Omega \rightarrow \Delta(\Omega)$

$$Q(i, j, \omega)(\omega') = \begin{cases} P(i, j, \omega)(\omega') & \text{if } \omega \neq \omega'; \\ P(i, j, \omega)(\omega') - 1 & \text{if } \omega = \omega'. \end{cases}$$

- Recall that $P(i, j, \omega)(\omega')$ is the probability that the next state is ω' , if the current state is ω and players' actions are (i, j) ;
- Hence the closer kernel Q is to 0, the more probable it is that the next state coincides with the current one.

Stochastic games with stage duration

- Consider a family of stochastic games G_h , parametrized by $h \in (0, 1]$;
- h represents stage duration;
- Players now play at times $h, 2h, 3h, \dots$, instead of playing at times $1, 2, 3, \dots$;
- State space Ω and action spaces I and J of player 1 and player 2 are independent of h ;
- Payoff function g_h of player 1 and kernel Q_h depend on h .

Stochastic games with stage duration

- Payoff $g_h = hg$;
- Kernel $Q_h = hQ$;
- $h = 1$: “Usual” stochastic game;
- When h small, g_h is close to zero (players receive almost nothing each turn), and Q_h is close to zero (the next state with a high probability will be the same).

Comparison (1)

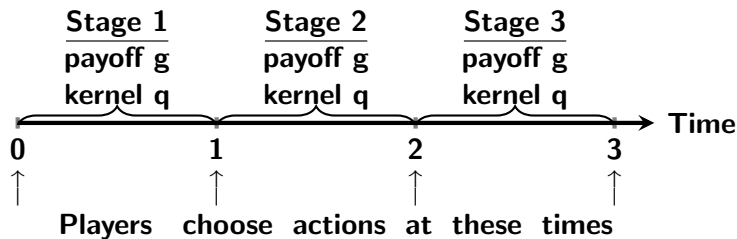


Figure: "Usual" stochastic game: duration of each stage is 1

Comparison (2)

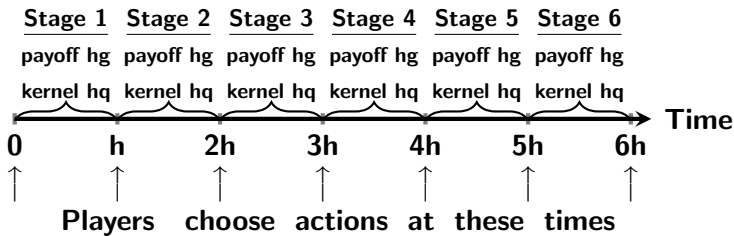


Figure: Stochastic game with stage duration h : stage payoff and kernel are proportional to h

Discounted games with stage duration

- For a game with stage duration h , the total payoff is (depending on the discount factor λ , initial state ω , and strategies σ, τ of players)

$$E_{\sigma, \tau}^{\omega} \left(\lambda \sum_{k=1}^{\infty} (1 - \lambda h)^{k-1} g_h^k \right);$$

- Why such a choice? Easy explanation:
- The total payoff is λ -discounted game with stage duration 1 is $E_{\sigma, \tau}^{\omega} \left(\lambda \sum_{k=1}^{\infty} (1 - \lambda)^{k-1} g^k \right)$. The total payoff of λ -discounted game with stage duration h is $E_{\sigma, \tau}^{\omega} \left(\sum_{k=1}^{\infty} \lambda h (1 - \lambda h)^{k-1} g^k \right)$;
- So, it may be seen as a game with discount factor λh . I.e., the discount factor is proportional to h , just as the payoff g and the kernel Q .

Real meaning behind the total payoff of the game with stage duration h (1)

- Total payoff: $E_{\sigma, \tau}^{\omega} \left(\lambda \sum_{k=1}^{\infty} (1 - \lambda h)^{k-1} g_h^k \right)$;
- When h is small, the total payoff of the λ -discounted stochastic game with stage duration h is close to the total payoff of the analogous λ -discounted continuous-time game;
- In a continuous-time game, players can choose actions at any time, and at each time t they receive instantaneous payoff g_t . The total payoff is (depending on the discount factor λ) $\int_0^{\infty} \lambda e^{-\lambda t} g_t dt$. The received payoff during a period of time $[(n-1)h, nh]$ is $\int_{(n-1)h}^{nh} \lambda e^{-\lambda t} g_t dt \underset{h \rightarrow 0}{\approx} \lambda h g_{(n-1)h} (1 - \lambda h)^{n-1}$.

Real meaning behind the total payoff of the game with stage duration h (2)

- Total payoff $E_{\sigma, \tau}^{\omega} \left(\lambda \sum_{k=1}^{\infty} (1 - \lambda h)^{k-1} g_h^k \right)$;
- Thus the total payoff of a continuous-time game is close to the total payoff of analogous discrete-time game;
- One can prove strictly that the values of λ -discounted games with stage duration h tend to the value of analogous continuous-time λ -discounted game when $h \rightarrow 0$.

Discounted games with stage duration (main properties)

- We denote by $v_{h,\lambda}$ the value of the game with total payoff $E_{\sigma,\tau}^{\omega} (\lambda \sum_{k=1}^{\infty} (1 - \lambda h)^{k-1} g_h^k)$;
- Main question: What happens with $v_{h,\lambda}$ when $h \rightarrow 0$?

Proposition (A. Neyman)

$\lim_{h \rightarrow 0} v_{h,\lambda}$ exists and is a unique solution of a functional equation.

Proposition (S. Sorin, G. Vigeral)

$\lim_{\lambda \rightarrow 0} \lim_{h \rightarrow 0} v_{h,\lambda}$ exists if and only if $\lim_{\lambda \rightarrow 0} v_{1,\lambda}$ exists.

- $\lim_{\lambda \rightarrow 0} v_{1,\lambda}$ should be considered as the limit value of the discrete-time stochastic game, whereas $\lim_{\lambda \rightarrow 0} \lim_{h \rightarrow 0} v_{h,\lambda}$ should be considered as the limit value of analogous continuous-time game.

Repeated Games with Stage Duration and Public Signals

Repeated Games with Public Signals (2)

A *zero-sum repeated game with symmetric information*, is a 7-tuple $(A, \Omega, f, I, J, g, P)$, where:

- A is a non-empty set of signals;
- Ω is a non-empty set of states;
- $f : \Omega \rightarrow A$ is a partition of Ω ;
- I is a non-empty set of actions of player 1;
- J is a non-empty set of actions of player 2;
- $g : I \times J \times \Omega \rightarrow \mathbb{R}$ is stage payoff function of player 1;
- $P : I \times J \times \Omega \rightarrow \Delta(\Omega)$ is the transition probability function.

We assume that I, J, Ω, A are finite.

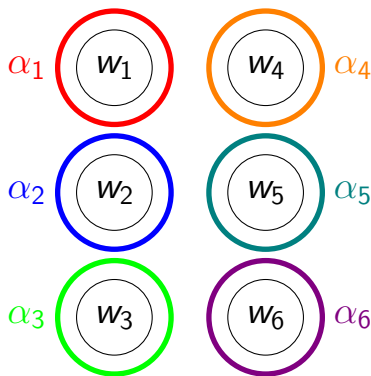
Repeated Games with Public Signals (3)

The game $(A, \Omega, f, I, J, g, P)$ proceeds in stages as follows. At each stage n :

1. The current state is ω_n . Players do not observe it, but they observe the signal $\alpha_n = f(\omega_n) \in A$ and the actions of each other at the previous stage;
2. Players choose their mixed actions, $x_n \in \Delta(I)$ and $y_n \in \Delta(J)$;
3. Pure actions $i_n \in I$ and $j_n \in J$ are chosen according to $x_n \in \Delta(I)$ and $y_n \in \Delta(J)$;
4. Player 1 obtains a payoff $g_n = g(i_n, j_n, \omega_n)$, while player 2 obtains payoff $-g_n$;
5. The new state ω_{n+1} is chosen according to the probability law $P(i_n, j_n, \omega_n)$. The new signal is $\alpha_{n+1} = f(\omega_{n+1})$.

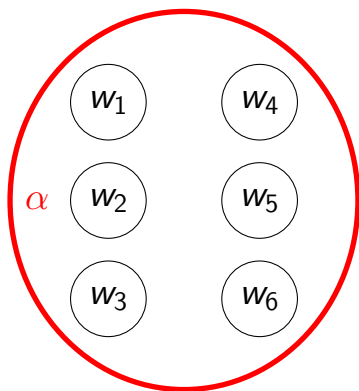
The above description of the game is known to the players.

Examples of the partition function f (1)



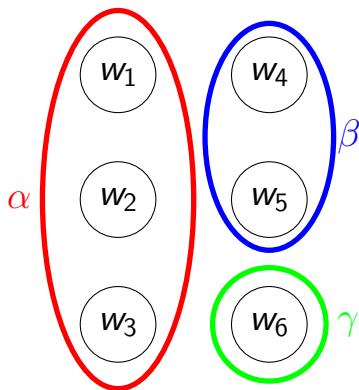
The perfect observation of the state, i.e. there are 6 signals $\alpha_1, \dots, \alpha_6$; and $f(w_i) := \alpha_i$.

Examples of the partition function f (2)



The state-blind case. There is only one signal α , and $f(w_i) := \alpha$

Examples of the partition function f (3)



Neither the perfect observation of the state nor the state-blind case. There are 3 signals, and $f(w_1) = f(w_2) = f(w_3) = \alpha$, $f(w_4) = f(w_5) = \beta$, $f(w_6) = \gamma$.

Stage duration

- We still can consider games with stage duration h in this new setting;
- Payoff $g_h = hg$;
- Kernel $Q_h = hQ$;
- State space Ω , signal set A , partition function f , and action spaces I and J of player 1 and player 2 are independent of h ;
- The total payoff is still $E_{\sigma,\tau}^{\omega} \left(\lambda \sum_{k=1}^{\infty} (1 - \lambda h)^{k-1} g_h^k \right)$;
- $v_{h,\lambda}$ is the value of the game with such a total payoff.

First result

Theorem

In the state-blind case, $\lim_{h \rightarrow 0} v_{h,\lambda}$ exists and is a unique viscosity solution of a partial differential equation.

- The main tool in the proof is the Shapley equation;
- Open question: can we say that $\lim_{h \rightarrow 0} v_{h,\lambda}$ in a general case of games with public signals; if yes, is it a unique viscosity solution of some PDE?

Second result

Theorem

There is a game in which uniform limit $\lim_{\lambda \rightarrow 0} \lim_{h \rightarrow 0} v_{h,\lambda}$ exists, but pointwise limit $\lim_{\lambda \rightarrow 0} v_{1,\lambda}$ does not exist.

Open question: can we say that

1. For any fixed $h \in (0, 1]$, the limit $\lim_{\lambda \rightarrow 0} v_{h,\lambda}$ does not exist?
2. We have $\left| \limsup_{\lambda \rightarrow 0} v_{h,\lambda}(p) - \liminf_{\lambda \rightarrow 0} v_{h,\lambda}(p) \right| \rightarrow 0$ as $h \rightarrow 0$, uniformly in p ?

Second result (proof)

Theorem

There is a game in which uniform limit $\lim_{\lambda \rightarrow 0} \lim_{h \rightarrow 0} v_{h,\lambda}$ exists, but pointwise limit $\lim_{\lambda \rightarrow 0} v_{1,\lambda}$ does not exist.

- $\lim_{\lambda \rightarrow 0} v_{1,\lambda}$ does not exist: the game in this case is equivalent to a game from “Hidden stochastic games and limit equilibrium payoffs” (Jérôme Renault and Bruno Ziliotto), which is known to not have the limit value;
- $\lim_{\lambda \rightarrow 0} \lim_{h \rightarrow 0} v_{h,\lambda}$ exists: use the first result.

Generalization: varying stage duration

- Now we allow different stage durations for different stages;
- There is a sequence $\{h_i\}_{i \in \mathbb{N}}$;
- Players act in times $h_1, h_1 + h_2, h_1 + h_2 + h_3, \dots$;
- i -th stage payoff is $h_i g$ and i -th stage kernel is $h_i Q$;
- The analogues of the above theorems hold in this more general model.

This is all.

Thank you!