

Regular Expressions for Languages over $\Sigma = \{0, 1\}$

(a) $\{w \in \Sigma^* \mid w \text{ begins with a 0 and ends with a 1}\}$

Regular Expression:

$0\Sigma^*1$

(b) $\{w \in \Sigma^* \mid w \text{ contains at least four 0s}\}$

Regular Expression:

$\Sigma^*0\Sigma^*0\Sigma^*0\Sigma^*0\Sigma^*$

(c) $\{w \in \Sigma^* \mid w \text{ contains the substring 1101}\}$

Regular Expression:

$\Sigma^*1101\Sigma^*$

(d) $\{w \in \Sigma^* \mid w \text{ has length at least 4 and its third symbol is a 0}\}$

Regular Expression: $\Sigma\Sigma0\Sigma^*$

(e)

$\{w \in \Sigma^* \mid w \text{ begins with a 0 and has even length, or begins with a 1 and has odd length}\}$

Regular expression:

$0(\Sigma\Sigma)^*1(\Sigma\Sigma)^*\Sigma$

(f) $\{\epsilon, 0\}$

Regular Expression:

$\epsilon|0$

(g) The empty set

Regular Expression: \emptyset

(h) All strings over Σ

Regular Expression: Σ^*

(i) All strings over Σ except the empty string

Regular Expression: Σ^+

2. Strings in Languages over $\Sigma = \{a, b\}$

For each of the following languages, give...

- two strings that **are** members and
- two strings that **are not** members
- **(a)** a^*b^*
 - members
 - ϵ
 - $bbaa$
 - non-members
 - ab
 - $baab$
- **(b)** $a(ab)^*b$
 - members
 - ab
 - abb
 - non-members
 - ba
 - aa
- **(c)** $a^+ \cup b^+$
 - members
 - a

- bb
- non-members
 - ϵ
 - ab
- **(d)** $(aba)^*$
 - members
 - aba
 - ϵ
 - non-members
 - $abaa$
 - ab
- **(e)** $\Sigma a \Sigma^* b \Sigma^* a \Sigma$
 - members
 - $aabaa$
 - $babab$
 - non-members
 - $bbbbaa$
 - $aaaaaa$
- **(f)** $aba \cup bab$
 - members
 - aba
 - bab
 - non-members
 - ab
 - ba
- **(g)** $b(\epsilon \cup a)b$
 - members
 - bb
 - bab
 - non-members
 - ba
 - ab

- (h) $\Sigma^*(a \cup ba \cup bb)\Sigma^*$
 - members
 - a
 - aaa
 - non-members
 - none; this language allows any string over Σ , as long as it contains at least one of a , ba , or bb somewhere in it.

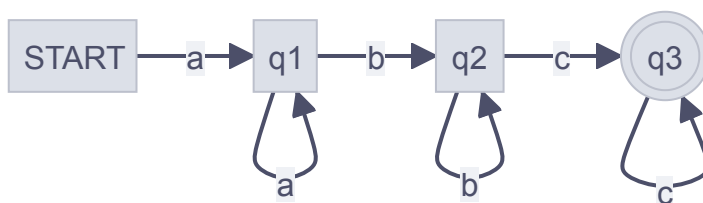
3. Regularity of Languages over $\Sigma = \{a, b, c\}$

(a) $\{a^{2i}b^j \mid i, j \geq 1\}$

- This language is **not** regular because it requires counting (to ensure the number of a 's are in multiples of 2)
- Counting indicates a non-regular language. Finite automata does not have memory for this.

(b) $\{a^n b^m c^k \mid n, m, k \geq 1\}$

- This language is regular because it doesn't require counting. (It only needs one of each member)



(c) $\{a^n b^{n+m} c^k \mid n, m, k \geq 1\}$

- This language is **not** regular. It requires counting the number of a 's and b 's to make sure that the number of b 's is at least as many as the number of a 's.
- Counting indicates a non-regular language. Finite automata does not have memory for this.