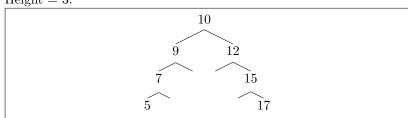
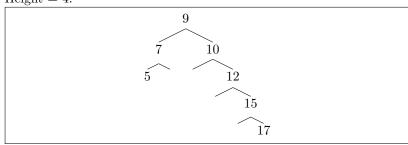
# 1. 6 points.

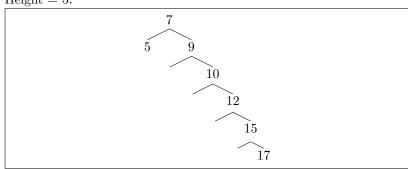
(a) Height = 3.



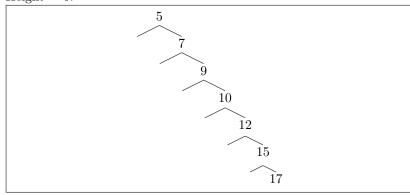
(b) Height = 4.



(c) Height = 5.



(d) Height = 6.



# 2. 24 points.

#### (a) 8 points.

No. This claim is not correct. Suppose A = all keys on the left of the search path, B = all keys on the search path, and C = all keys on the right of the search path.

In this example, we are searching for 6 and the sets  $A,\,B,\,$  and C are given by:

$$A = \{3\}, B = \{2, 8, 5, 6\}, C = \{9\}$$

$$(1)$$

As we can see,  $3 \in A$ ,  $2 \in B$  and  $3 \ge 2$ . Therefore, the property doesn't hold.

## (b) 8 points.

Suppose there is a node x that has two children and suppose its predecessor is p and its successor is s.

i. We will show by contradiction that the successor of x has no left child.

Suppose s has a left child. Then the key of s is > than left[s]. The key of s is also > than the key of x, and since s has a left child the key of left[s] is > than that of x.

Thus,  $key[s] \ge key[left[s]] \ge key[x]$ ,

which is a contradiction, since s is the successor of x. Hence the successor of x has no left child.

ii. We will show by contradiction that the predecessor of x has no right child.

Suppose p has a right child. Then the key of p is < than that of right[p]. The key of p is also < than the key of x, and since p has a right child the key of right[p] is < than x.

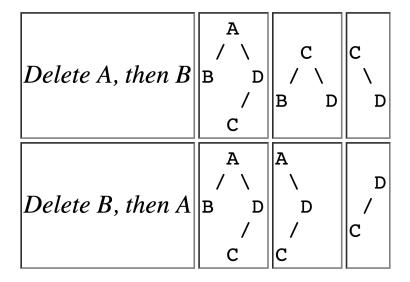
Thus  $key[p] \le key[right[p]] \le key[x]$ ,

which is a contradiction, since p is the predecessor of x. Hence the predecessor of x has no right child.

### (c) 8 points.

No, if we delete node x and then node y from a binary search tree,

we will **NOT** end up with the same resulting tree that we will get if we first delete node y and then node x. Deletion is not commutative. We will show this by an example:



- 3. 20 points.
  - (a) 5 points.

(b) 5 points.

(c) 5 points.

(d) 5 points.