

# Quantum interference channels

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# Overview

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- 3 Intro to Quantum Information Theory
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# Definitions

## Definition (Interference network)

A two party interference network

$(\mathcal{X}_1 \otimes \mathcal{X}_2, p(y_1, y_2 | x_1, x_2), \mathcal{Y}_1 \otimes \mathcal{Y}_2)$  is general model for communication networks with two inputs, two outputs and a probability transition matrix  $p(y_1, y_2 | x_1, x_2)$ .

## Definition (Interference channel)

A two party interference channel is a particular use of an interference network  $(\mathcal{X}_1 \otimes \mathcal{X}_2, p(y_1, y_2 | x_1, x_2), \mathcal{Y}_1 \otimes \mathcal{Y}_2)$  where messages  $M_1, M_2$  are independently encoded into a codewords  $X_1, X_2$  at rates  $(R_1, R_2)$  with  $Y_1, Y_2$  being the intended receiver.

# Best inner bound

## Han Kobayashi / Cheng Motani Garg inner bound

$$G_{HK} = G_{CMG} = \cup_{p(x_1|u_1)p(x_2|u_2)p(u_1)p(u_2)} \{ (R_1, R_2) | \text{Eqn. (1)} \}$$

$$R_1 \leq I(X_1; Y_1 | U_2 Q)$$

$$R_2 \leq I(X_2; Y_2 | U_1 Q)$$

$$R_1 + R_2 \leq I(X_1 U_2; Y_1 | Q) + I(X_2; Y_2 | U_1 U_2 Q)$$

$$R_1 + R_2 \leq I(X_1; Y_1 | U_1 U_2 Q) + I(X_2 U_1; Y_2 | Q) \quad (1)$$

$$R_1 + R_2 \leq I(X_1 U_2; Y_1 | U_1 Q) + I(X_2 U_1; Y_2 | U_2 Q)$$

$$2R_1 + R_2 \leq I(X_1 U_2; Y_1 | Q) + I(X_1; Y_1 | U_1 U_2 Q) + I(X_2 U_1; Y_2 | U_2 Q)$$

$$R_1 + 2R_2 \leq I(X_2; Y_2 | U_1 U_2 Q) + I(X_2 U_1; Y_2 | Q) + I(X_1 U_2; Y_1 | U_1 Q)$$

# First outer bound

$$G_{Naive} \triangleq \text{conv} \cup_{p(x_1)p(x_2)} \{ (R_1, R_2) \mid \text{Eqn. (3)} \} \quad (2)$$

$$\begin{aligned} R_1 &\leq I(X_1; Y_2 | X_2) \\ R_2 &\leq I(X_2; Y_2 | X_1) \end{aligned} \quad (3)$$

# Sato outer bound

$$G_{Sato} \triangleq \text{conv} \cup_{p(x_1)p(x_2)} \{ (R_1, R_2) \mid \text{Eqn. (5)} \} \quad (4)$$

$$\begin{aligned} R_1 &\leq I(X_1; Y_2 | X_2) \\ R_2 &\leq I(X_2; Y_2 | X_1) \\ R_1 + R_2 &\leq I(X_1 X_2; Y_1 Y_2) \end{aligned} \quad (5)$$

# Best outer bound

Consider the two random variables  $Z_1, Z_2$  such that

$Y_1$  is a degraded version of  $Z_1$

$Y_2$  is a degraded version of  $Z_2$

$Y_2$  is a degraded version of  $(X_1, Z_1)$

$Y_1$  is a degraded version of  $(X_2, Z_2)$

$$G_{Carl} \triangleq \text{conv} \cup_{p(x_1)p(x_2)} \{ (R_1, R_2) \mid \text{Eqn. (7)} \} \quad (6)$$

$$R_1 \leq I(X_1; Y_1 | X_2)$$

$$R_2 \leq I(X_2; Y_2 | X_1) \quad (7)$$

$$R_1 + R_2 \leq \min\{I(X_1 X_2; Z_1), I(X_1 X_2; Z_2)\}$$

# Quantum mechanics



In films and books by self-proclaimed life coaches:

- quantum mechanics is magic !
- stuff happens because you think about it ...

- quantum mechanics is linear algebra and probability

- quantum mechanics is linear algebra and probability
- replace bits<sup>01</sup> with  $\vec{q}$ ubits
- qubits live in a two dimensional complex vector space  $\mathcal{H}$ :

$$|x\rangle = \alpha|0\rangle + \beta|1\rangle \quad \alpha, \beta \in \mathbb{C} \quad (8)$$

- don't be fooled by the  $|\cdot\rangle$  notation...  
 $|\psi\rangle \equiv \vec{\psi}$  it is just a neat notational trick for dot products...
- Good for systems with finite degrees of freedom like electron spin, and light polarization
- Another *representation* exists for continuous systems via wave function

# Quantum source

- A quantum source mixture of qubits  $\{p(x), |\sigma_x\rangle\}$ , where state  $|\sigma_x\rangle$  is produced by the source with probability  $p(x)$ .
- We represent it as a density matrix

$$\rho = \sum_x p(x) |\sigma_x\rangle \langle \sigma_x| \quad (9)$$

- Models both classical uncertainty  $p(x)$  (which quantum state did the source produce?)
- and quantum uncertainty (given that the state produced was  $|x\rangle$  the probability of finding  $|0\rangle$  if measure is  $|\alpha|^2$ )

# Examples

- Pure quantum state  $|x\rangle = \alpha|0\rangle + \beta|1\rangle$  has a density matrix

$$\rho_x = |x\rangle\langle x| = \begin{bmatrix} |\alpha|^2 & \beta^* \alpha \\ \alpha^* \beta & |\beta|^2 \end{bmatrix}. \quad (10)$$

- Specific example  $|+\rangle = \frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$  becomes

$$\rho_+ = \begin{bmatrix} \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{2} \end{bmatrix}.$$

- A mixed state  $\rho = \{p(x), |\sigma_x\rangle\}$  with  $p(0) = \frac{1}{4}$ ,  $|\sigma_0\rangle = |0\rangle$  and  $p(1) = \frac{3}{4}$ ,  $|\sigma_1\rangle = |1\rangle$  gives a density matrix:

$$\rho = \frac{1}{4}|0\rangle\langle 0| + \frac{3}{4}|1\rangle\langle 1| = \frac{1}{4} \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} + \frac{3}{4} \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} \frac{1}{4} & 0 \\ 0 & \frac{3}{4} \end{bmatrix}.$$

- Density matrices are positive and have unit trace

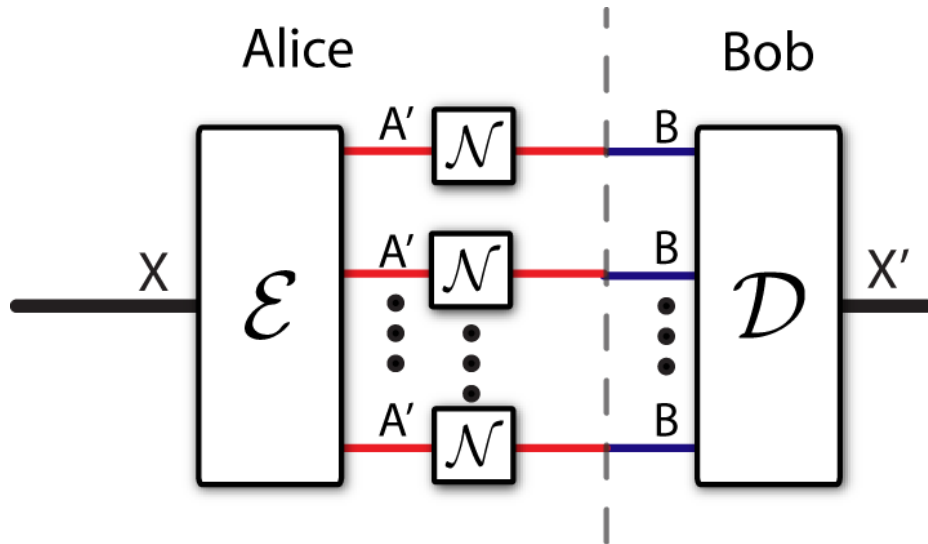
# Quantum channel

- A quantum channel  $\mathcal{N}: \mathcal{H}^A \rightarrow \mathcal{H}^B$  takes states to states

$$\mathcal{N}(\sigma^A) = \rho^B \quad (11)$$

- Completely positive, trace preserving (CPTP) map
- Analogue of cond. probability of a classical channel  $p(y|x)$

# Quantum channel



# Encoding

We can use a quantum channel to send either classical or quantum data.

For classical data, the encoding operation will be  $\mathcal{E}: \mathcal{X} \rightarrow \mathcal{H}^A$ ,

$$\mathcal{E}(x) = |\sigma_x\rangle^A \quad (12)$$

where  $|\sigma\rangle_x$  is some predetermined set of “signal states” that will be send over the channel.

To encode some quantum data  $\rho$  we can do any quantum operation  $\mathcal{E}: \mathcal{H}^S \rightarrow \mathcal{H}^A$ ,

$$\mathcal{E}(\rho^S) = \sigma^A. \quad (13)$$



# Decoding classical data

We want to extract the classical message from the received quantum state.

Measurements act on density matrices to produce a classical output as well as a possibly modified quantum state.

Measurements are sets of projection operators  $\{M_i\}$ .

The probability of outcome  $m$  occurring when the input system is  $\rho$  is

$$p(m) = \text{Tr}(M_m^\dagger \rho M_m) .$$

Our example from earlier

$$\Pr\{\text{measure } x = 0\} = \text{Tr} \left( \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} |\alpha|^2 & \beta^* \alpha \\ \alpha^* \beta & |\beta|^2 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix} \right) = |\alpha|^2$$

# Decoding quantum data

In the error criterion for classical information theory protocols is of the form  $P_e = \Pr\{M_1 \neq \hat{M}_1\}$ . We need an analogous criterion for quantum protocols.

The fidelity between two pure quantum states is the square of their inner product

$$F(|\varphi\rangle, |\psi\rangle) = |\langle\varphi|\psi\rangle|^2. \quad (14)$$

Two states that are very similar have fidelity close to 1. Success criterion is

$$F(|M\rangle, |\hat{M}\rangle) \geq 1 - \epsilon. \quad (15)$$

# Quantum information theory

# Quantum information theory

- Came from classical information theory
- Replace Shannon entropy with von Neumann entropy

## Definition (von Neumann Entropy)

A quantum system described by the density matrix  $\rho$  has von Neumann entropy  $S(\rho)$  given by the expression:

$$S(\rho) = -\text{Tr}(\rho \log \rho) \quad (16)$$

- from classical `import mutual_information,`  
`Fano_inequality,` `random_coding,`  
`channel_capacity,` `compression,` `etc..`

# Classical capacity of quantum channel

Holevo, Schumacher and Westmoreland found a multiletter formula for the classical capacity [H98,SW97].

## Theorem

*The classical capacity of a quantum channel  $\mathcal{N}$  is given by  $\mathcal{C}(\mathcal{N}) = \frac{1}{n} \bigcup_{n=1}^{\infty} C^{(1)}(\mathcal{N}^{\otimes n})$  where*

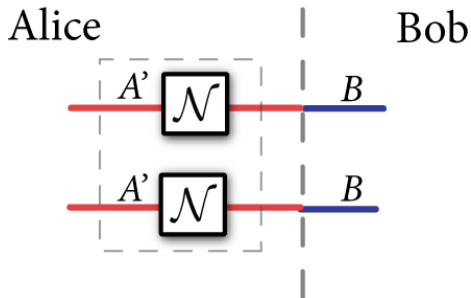
$$C^{(1)}(\mathcal{N}) = \max_{\rho} I(A; B) \quad (17)$$

*where  $\rho$  is the output state*

$$\rho = \sum_x p(x) |x\rangle \langle x|_A \otimes \mathcal{N}(\sigma_x)_B. \quad (18)$$

*and  $\{p(x), \sigma_x\}$  is the input distribution.*

# Limitations of these results



# Quantum multiple access channel

## Definition

The quantum multiple access channel (QMAC) is a map

$$\mathcal{M}^{A'B' \rightarrow C} \quad (19)$$

where systems  $A'$  and  $B'$  are the inputs and  $C$  is the output.

Three different capacities

- $\mathcal{C}$ : classical info
- $\mathcal{Q}$ : quantum info
- $\mathcal{C}_E$  or  $\mathcal{Q}_E$ : with additional resource *entanglement*

# Classical capacity of QMAC

## Theorem (Winter 01)

*The capacity of the quantum MAC  $\mathcal{M}$  to carry classical information is  $\mathcal{C}_{QMAC}(\mathcal{M}) = \frac{1}{n} \cup_{n=1}^{\infty} \mathcal{C}_{QMAC}^{(1)}(\mathcal{M}^{\otimes n})$  where*

$$\mathcal{C}_{QMAC}^{(1)} = \text{conv} \cup_{p(x_1)p(x_2)\sigma_x} \{ (R_1, R_2) \mid \text{Eqn. (21)} \} \quad (20)$$

$$\begin{aligned} R_1 &\leq I(A; C|B)_{\theta}, \\ R_2 &\leq I(B; C|A)_{\theta}, \\ R_1 + R_2 &\leq I(AB; C)_{\theta}, \end{aligned} \quad (21)$$

*where mutual informations are calculated with respect to*

$$\theta = \sum_x p(x_1)p(x_2) |x_1\rangle\langle x_1|_A \otimes |x_2\rangle\langle x_2|_B \otimes \mathcal{M}(\sigma_{x_1} \otimes \sigma_{x_2})_C. \quad (22)$$



# Quantum interference channel

## Definition

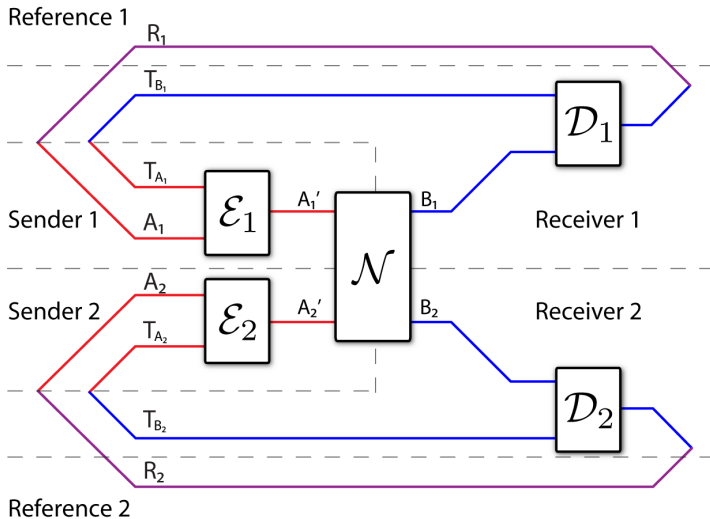
The quantum interference channel is a map

$$\mathcal{N}^{A'_1 A'_2 \rightarrow B_1 B_2} \quad (23)$$

Open problem.

can adapt some of the classical results...

# Entanglement assisted



thanks for your attention!

# Physicist's view

- Very similar to what we learn in signals and systems

# Physicist's view

- Very similar to what we learn in signals and systems
- A quantum state (signal) is represented by a *wavefunction*

$$\psi: \mathbb{R}^3 \rightarrow \mathbb{C} \quad (24)$$

- The wave function of Hydrogen's ground state ( $n=1$   $\ell=0$ )

$$\psi(r, \vartheta, \varphi) = \sqrt{\frac{1}{\pi a^3}} \exp\left(\frac{-r}{a}\right), \quad (25)$$

in the *position* basis. Apply Fourier transform to get the *momentum* basis

# Physicist's view (continued)

- The power of the signal is a probability density

$$\Pr\{\text{finding electron at } \vec{r}\} = |\psi(\vec{r})|^2 \quad (26)$$

- Verify that it is well normalized

$$\begin{aligned} P_{total} &= \iiint |\psi(\vec{r})|^2 d^3\vec{r} \\ &= \int_0^\infty \int_0^{2\pi} \int_0^\pi |\psi(r, \vartheta, \varphi)|^2 r^2 \sin \varphi d\varphi d\vartheta dr \\ &= \int_0^\infty \frac{4}{a^3} \exp\left(\frac{2r}{a}\right) r^2 dr = 1 \end{aligned} \quad (27)$$