

## Graded for Honest Effort

1.7.3 a) "someone did not get a large bonus" is a negation of "someone get a large bonus", and we can use the quantifier  $\exists$  to express someone, we can translate it into logic as

$$\exists x \neg B(x)$$

b) we use the quantifier  $\forall$  to express everyone. we can translate "Everyone got a large bonus" into logic as

$$\forall x B(x)$$

c) we use conjunction to express "even though" and "Sam did not get a large bonus" is a negation of "sam got a large bonus". we can translate it into logic as

$$\neg B(\text{sam}) \wedge T(x)$$

d) we use the quantifier  $\exists$  to express "someone" and "not on the executive team" is a negation of "on the executive team" and finally we use conjunction to express "not on the executive team and received large bonus"

$$\exists x (\neg T(x) \wedge B(x))$$

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- 1.7.3 e) We use the quantifier  $\forall$  to express "Everyone", and being a executive team member implies he/she got a large bonus

$$\forall x (T(x) \rightarrow B(x))$$

- 1.7.7 b) In every row, we can find at least  $D(x)$  or  $N(x)$  is true, therefore

True

Among all the group members, they either missed the deadline or they are a new employee.

- c) We can find that there exists a  $x$  that is both not  $D(x)$  and  $N(x)$ , for example, Melanie is not  $D(x)$  and  $N(x)$ , therefore

True

A group member did not miss the deadline and is a new employee.

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1.7.7 i) We find that Only AI is true on Both  $B(x)$  and  $N(x)$  which means

False

Every group member missed the deadline if and only if he is a new employee.

1.8.3 d) Translate to logic  
 $\exists x (P(x) \wedge C(x))$   
Apply negation  
 $\neg \exists x (P(x) \wedge C(x))$

$\equiv \forall x \neg (P(x) \wedge C(x))$  change quantifier to  $\forall$   
 $\equiv \forall x (\neg P(x) \vee \neg C(x))$  change conjunction to disjunction

Translate to English

Every student showed up either without a pencil or without a calculator

## Graded for Honest Effort

e) Translate to logic

$$\exists x (P(x) \vee C(x))$$

Apply negation

$$\neg \exists x (P(x) \vee C(x))$$

$$\equiv \forall x \neg (P(x) \vee C(x)) \quad \text{change the quantifier to } \forall$$

$$\equiv \forall x (\neg P(x) \wedge \neg C(x)) \quad \text{change disjunction to conjunction}$$

Translate to English

Every student showed up without a pencil and without a calculator.

f) Translate to logic

$$\forall x (P(x) \wedge C(x))$$

Apply negation

$$\neg \forall x (P(x) \wedge C(x))$$

$$\exists x \neg (P(x) \wedge C(x)) \quad \text{change quantifier to } \exists$$

$$\exists x (\neg P(x) \vee \neg C(x)) \quad \text{change conjunction to disjunction}$$

Translate to English

Some students showed up either without a pencil or without a calculator.

## Graded for Honest Effort

1.10.3 a)  $\forall x \exists y P(x, y)$  ,  $\exists x \forall y P(x, y)$

		y		
x	P	a	b	c
	a	F	F	F
	b	T	T	T
	c	F	F	T

$\forall x \exists y P(x, y)$  is false because  $P(1, y)$  are all false. It needs at least one true to make the expression true

$\exists x \forall y P(x, y)$  is true because  $P(2, y)$  are all true,

1.10.4 e) The reciprocal of every positive number less than one is greater than one

$R(x, y) \equiv x$  is the reciprocal of  $y$ .

$G(z) \equiv z$  is greater than one.

$$\forall x \exists y ((R(y, x) \wedge (x > 0) \wedge (x < 1)) \rightarrow G(y))$$

If every real number  $x$  that is greater than and less than 1, then there is a real number  $y$  that is greater than 1.

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1.10.4 f) There is no smallest number.

$S(x, y)$ :  $x$  is smaller than  $y$ .

$\exists x S(x)$ :  $x$  is the smallest number.

$$\forall y \exists x (S(x, y))$$

For every real number  $y$ , there is a real number  $x$

g) Every number other than 0 has a multiplicative inverse

$m(x)$ :  $x$  has a multiplicative inverse

$$\forall x (x \neq 0 \rightarrow m(x))$$

For every real number  $x$ , if  $x$  is not zero, it has a multiplicative inverse

## Graded for Honest Effort

- 1.11.3 b)  $p$  : he studied for the test  
 $q$  : he passed the test

$$\begin{array}{l} p \vee \neg q \\ \underline{q} \\ \therefore p \end{array}$$

$p$	$q$	$\neg q$	$p \vee \neg q$
T	T	F	T
T	F	T	T
F	T	F	F
F	F	T	T

The hypothesis are  $p \vee \neg q$  and  $q$ , so we only look at the row when both are true which is row 1. The conclusion which is  $p$  is true therefore it is **valid**.

- c)  $p$  :  $\sqrt{2}$  is an irrational number  
 $q$  :  $2\sqrt{2}$  is an irrational number

$$\begin{array}{l} p \rightarrow q \\ \underline{q} \\ \therefore p \end{array}$$

$p$	$q$	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

The hypothesis are  $p \rightarrow q$  and  $q$  so we only look at row 1 and 3, The conclusion at row 3 is false, therefore it is **invalid**.

Graded for Honest Effort and Feedback Given

1.8.4 b)  $\neg \forall x (\neg P(x) \rightarrow Q(x))$   
 $\equiv \exists x (\neg P(x) \wedge \neg Q(x))$

$\neg \forall x (\neg P(x) \rightarrow Q(x))$	start
$\exists x \neg (\neg P(x) \rightarrow Q(x))$	Demorgan's law
$\exists x \neg (\neg \neg P(x) \vee Q(x))$	conditional identity
$\exists x \neg (P(x) \vee Q(x))$	Double negation law
$\exists x (\neg P(x) \wedge \neg Q(x))$	Demorgan's law
$\therefore$ They are logically equivalent	

1.9.4 d)  $\exists x \forall y (P(x,y) \leftrightarrow P(y,x))$   
apply negation

$\neg \exists x \forall y (P(x,y) \leftrightarrow P(y,x))$	start
$\forall x \exists y \neg (P(x,y) \leftrightarrow P(y,x))$	Demorgan's law
$\forall x \exists y \neg ((P(x,y) \rightarrow P(y,x)) \wedge (P(y,x) \rightarrow P(x,y)))$	conditional identity
$\forall x \exists y \neg ((\neg P(x,y) \vee P(y,x)) \wedge (P(y,x) \rightarrow P(x,y)))$	conditional identity
$\forall x \exists y \neg ((\neg P(x,y) \vee P(y,x)) \wedge (\neg P(y,x) \vee P(x,y)))$	conditional identity
$\forall x \exists y (\neg(\neg P(x,y) \vee P(y,x)) \vee \neg(\neg P(y,x) \vee P(x,y)))$	Demorgan's law
$\forall x \exists y ((P(x,y) \wedge \neg P(y,x)) \vee \neg(\neg P(y,x) \vee P(x,y)))$	Demorgan's law
$\forall x \exists y ((P(x,y) \wedge \neg P(y,x)) \vee (P(y,x) \wedge \neg P(x,y)))$	Demorgan's law



## Graded for Honest Effort and Feedback Given

1.10.7

c) There is at least one new employee who missed the deadline.

We can translate it into logic as

$$\exists x (N(x) \wedge D(x))$$

We use quantifier  $\exists$  to express at least one, then we use conjunction to express a person is a new employee and missed the deadline.

d) Sam knows the phone number of everyone who missed the deadline.

We can translate it into logic as

$$\forall x (D(x) \rightarrow P(\text{Sam}, x))$$

We use quantifier  $\forall$  to express everyone, and we use conditional statement to express that if you missed the deadline then Sam knows your phone number.

## Graded for Honest Effort and Feedback Given

1.10.7 e) There is a new employee who knows everyone's phone number.

We can translate it into logic as

$$\exists x \forall y (N(x) \wedge P(x, y))$$

We use quantifier  $\exists$  to express a person, and use quantifier  $\forall$  to express everyone. Finally, we use conjunction to express there is a person who is a new employee and knows everyone's phone number.

f) Exactly one new employee missed the deadline. We can translate it into logic as

$$\exists x ((N(x) \wedge D(x)) \wedge \forall y ((x \neq y) \rightarrow \neg D(y)))$$

We use quantifier  $\exists$  to express there is one person, and this person is a new employee and missed the deadline. We use quantifier  $\forall$  to express everyone and if they are the same person  $x$  then they didn't miss the deadline which means exactly one person missed the deadline.

## Graded for Honest Effort and Feedback Given

1.10.10

$x$  = a set of student at a university.

$y$  = a set Math class offer at that university.

$T(x, y)$  = student  $x$  has taken class  $y$ .

c) Every student has taken at least one class other than Math 101.

$$\forall x \exists y ( (y \neq \text{Math 101}) \wedge T(x, y) )$$

We use quantifier  $\forall$  to express every student and  $\exists$  to express at least one class, then we say that this class is not Math 101 and every student has taken this Math class.

d) There is a student who has taken every math class other than Math 101.

$$\exists x \forall y ( (y \neq \text{Math 101}) \rightarrow T(x, y) )$$

We use quantifier  $\exists$  to express there is a student and  $\forall$  to express every math class. Then we say that if every math class that is not math 101 then there is a student has taken that class.

Graded for Honest Effort and Feedback Given

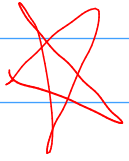
1.10.10 f) Sam has taken exactly two math classes.

$$\exists a \exists b \forall y (T(\text{Sam}, a) \wedge T(\text{Sam}, b) \wedge ((a \neq b) \wedge (a \neq y) \wedge (b \neq y)) \rightarrow \neg T(\text{Sam}, y))$$

We say that Sam has taken class  $a$  and  $b$ , and if class  $a$  and  $b$  are not the same class, then Sam didn't take any other classes.

## Graded for Honest Effort and Feedback Given

1.11.4



$p$ : 4 is a prime number (False)

$q$ : 5 is a prime number (True)

a)

$$\frac{q \quad p \vee q}{\therefore p}$$

$p$	$q$	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

The conclusion is a false statement and the first row in truth table shows the argument is valid.

b)

$$\frac{q \quad \neg p \wedge q}{\therefore p}$$

$p$	$q$	$\neg p$	$\neg p \wedge q$
T	T	F	F
T	F	F	F
F	T	T	T
F	F	T	F

The conclusion statement is true and the third row in truth table shows all the hypotheses are true but conclusion is false.