## Bow-Beck-Lurie thim

adjunction ~> dED + extra data Ex D ifaithfully flat descent) A (\*) B a may of commutative rings. T: BOAN -NOAB seitisfying "cocycle condition" If (\*) is faithfully flat => converse true 2 Morita theory R-module

generator =	finitely generated projective  Donverse true  R-Mod a EndelP-Mod
I. Monads =	monoids in the cat of endofunctors
l: cate	jory
Jef (Monad)	$T \in End(C)$ $\eta \in 1_C \rightarrow T$ $\mu \in T^2 \rightarrow T$ satisfying
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	$T^{2} \stackrel{HT}{/} T$ $T^{2} \stackrel{T}{/} T$ Burritos for the Hungry Mathematician
Ex D. bu	Ed Morehouse
· · · · · · · · · · · · · · · · · · ·	April 1, 2015
	Abstract  The advent of fast-casual Mexican-style dining establishments, such as
Shefting ()	The advent of fast-casual Mexican-style dining establishments, such as Chipotle and Qdoba, has greatly improved the productivity of research mathematicians and theoretical computer scientists in recent years. Still, many experience confusion upon encountering burritos for the first time.  Numerous burrito tutorials (of varying quality) are to be found on the Internet. Some describe a burrito as the image of a crêpe under the action of the new-world functor. But such characterizations merely serve to reindex the confusion contravariantly. Others insist that the only way to really understand burritos is to eat many different kinds of burrito,

to really understand burritos is to eat many different kinds of burrito, until the common underlying concept becomes apparent.

It has been recently remarked by Yorgey [9] that a burrito can be regarded as an instance of a universally-understood concept, namely, that of monad. It is this characterization that we intend to explicate here. To wit, a burrito is just a strong monad in the symmetric monoidal category of food, what's the problem?

O (free module). Tc is a module /m: Tc M Tc Ab R-mod  $T = U(-0R) = \emptyset R$  acts on Ab ID T-mod = R-mod In the adjunction setting HdED, aids is a Gof-module  $m: (aF) (ad) \xrightarrow{GE} a(d)$ Da factors Mody (l)  $e = \frac{1}{u} Mod_{T}(e) = \frac{1}{G} D$ a is monadic if a is eq

Remark: Everything has duch version.

III. Bourr-Beck th'm    Beck thesis Triples algebras and cohomologies   Machane & VI.7   Riehl & 5.6
is menadic if it satisfies  i) a is conservative > a(f) is an eq.
2) s (crude ver) Dadmits & a preserves reflexive coeq.  (precise ver) Dadimits & a preserves coeq of iff and in preserves coeq.
Reflexive coequalize:  X, \frac{1}{g7} \times 0 \frac{e}{s} \times 1  S common section
Split wegustize: $x_1 = \frac{1}{37} \times 0 = \frac{1}{2} \times 1$ pair $1_{x_0} = \frac{1}{5} = \frac{1}{2} \times 1$ $\frac{1}{5} = \frac{1}{5} = \frac{1}{5} \times 1$
G-split pair: a pair that is sent to a split pair by G

## Ex In Moniter theory

2 Morita theory

R: commutative ring

P: R-module S= EndR(P)

R-Mcd Homr(1,-) Ab

Q Homp (P,Q)

When P is finitely generated projective generator => converse true

R-Mod a Endell)-Mod

- D Homz(P,-) admits a left adjoint
- 2) P is generator => Home (P,-) is congenue
- 3) P is projective = D HomzlP,-) preserves reflexive veg

R-Mad ~ Mad\_(Ab) identified S-Mod

T(M) = HOMR (P, POM) T(Z) = Homp (P, P) = Endp(

IV higher categorical version DAG & 4.7
$ \begin{array}{cccc}                                  $
Mo mo e Ma mo
colort my F Cont my a
$\mathbb{Q}$ . Monad: $\mathbb{Z} : \mathbb{Z}^2 \to \mathbb{Z}$ , $\mathbb{Z} \to \mathbb{Z}$ + $\mathbb{Z}^2$ agramss.
com't encode infinite Wherene down
(F→FGF→F ~ id =D T²=T)
higher a monad is an algobj in End (2)
2) want a factors (Mod-(e)
ligher Ci E Fran (D, C)
Endic)
Got E (End(e) + an action on G) is
classifying obj
Da∈Mod_Fun(D,e) => G
In VAG Adj data trivial Fun (P,D)  ws End (P) > T
3 Thim conditions
(Crude version) reflexive (X, \(\sigma\) X_7

metric reclizations.  $(-\cdot, = X_2 \stackrel{?}{=} X_1 \stackrel{?}{=} X_0) \rightarrow X_{-1}$ cprecise version) split (X=X=)=X-) split simplicial diagrams  $\frac{1}{2} \left( \frac{1}{2} - \frac{1}{2} - \frac{1}{2} \right) = \frac{1}{2} \left( \frac{1}{2} - \frac{1}{2} - \frac{1}{2} \right) = \frac{1}{2} \left( \frac{1}{2} - \frac{1}{2} - \frac{1}{2} \right) = \frac{1}{2} \left( \frac{1}{2} - \frac{1}{2} - \frac{1}{2} - \frac{1}{2} \right) = \frac{1}{2} \left( \frac{1}{2} - \frac{1}{2} - \frac{1}{2} - \frac{1}{2} \right) = \frac{1}{2} \left( \frac{1}{2} - \frac{1$ 15-17-=3·3 [0] = 3 0 } [1] = {o < 1} [2] = 3001 (2] morph = [n] -> [m] = set of orderpresency maps  $[n] \ U + \infty \} \longrightarrow [m] \ U + -\infty \}$ sending  $\{-\infty\}$  to  $\{-\infty\}$ . A simplicial digram NOST) -> C split if it extends to NUDE). D-00 - d'aigrans give split creq in ordinary certegory. Thim (Bar-Beck-Lune)

adjunction

F: C = D:G

Then I monad T s.t G factors

C = ModT(C) = D

and G is an equivalence

if G is

1) conservative

2) (crude) D admits & G preserves
geometric realization

(pregise) = 1)

colim of G-split

simplicial diagrams.