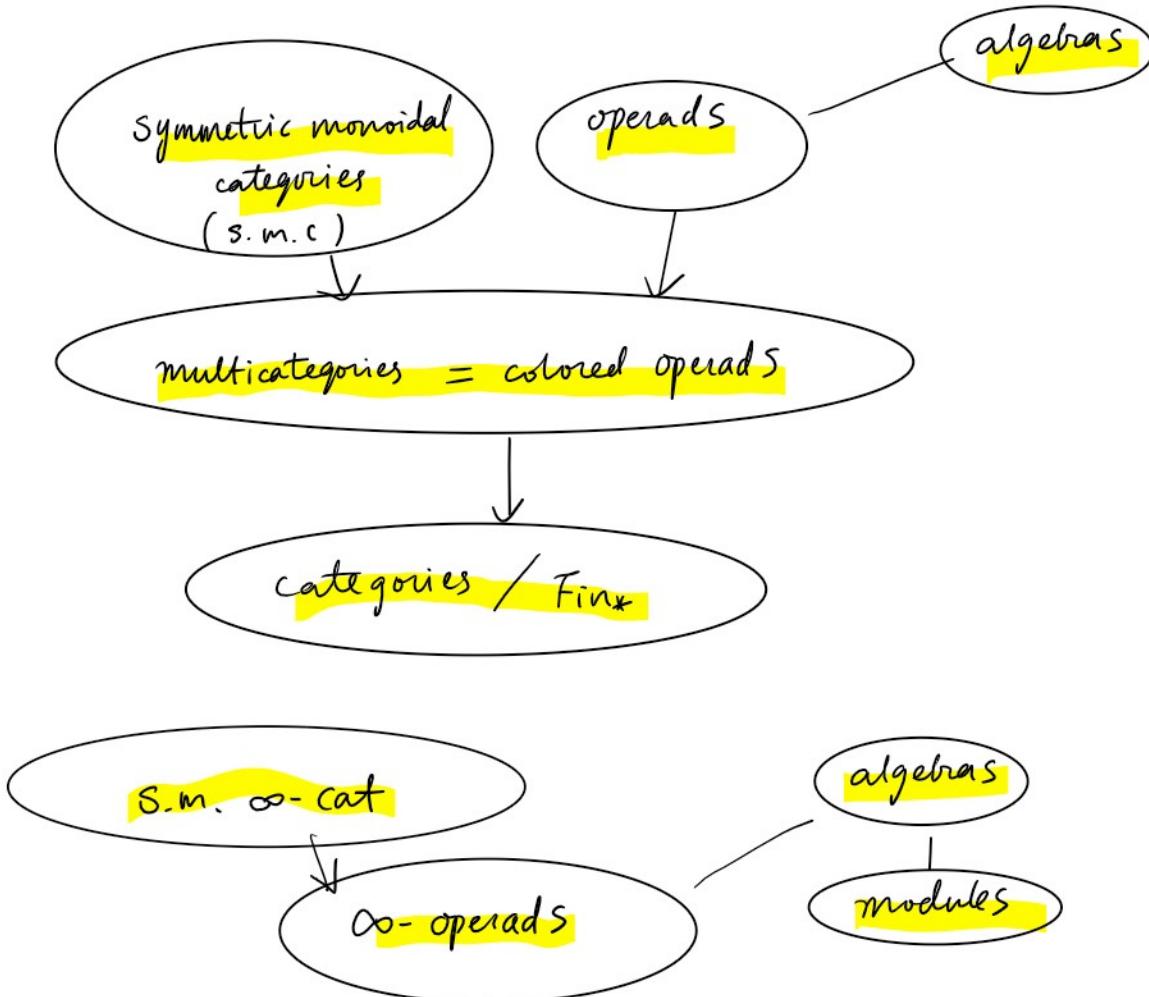
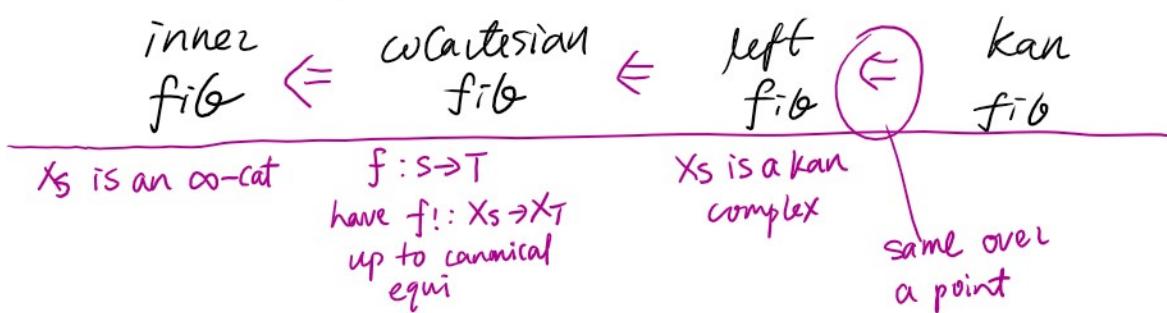


Monoidal structures, part 1



§ 0. recall $p : X \rightarrow S$



slogan.

- colax fib/S $\simeq \text{Fun}(S, \text{Cat}_{\infty})$
- lax fib/S $\simeq \text{Fun}(S^{\text{op}}, \text{Cat}_{\infty})$
- left fib/S $\simeq \text{Fun}(S, \text{Kan})$
- right fib/S $\simeq \text{Fun}(S^{\text{op}}, \text{Kan})$

§ 1. symmetric monoidal category $(\mathcal{C}, \otimes, I)$

§1. symmetric monoidal category $(\mathcal{C}, \otimes, I)$

🔗 Baez, some definitions everybody should know

<https://math.ucr.edu/home/baez/gg-fall2004/definitions.pdf>

Data: $\otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$

$I \in \text{ob}(\mathcal{C})$

$$\begin{cases} \alpha_x: I \otimes X \xrightarrow{\sim} X & (\text{left}) \text{ unit} \\ \beta_{X,Y}: X \otimes Y \xrightarrow{\sim} Y \otimes X & \text{commutator} \\ \gamma_{X,Y,Z}: (X \otimes Y) \otimes Z \xrightarrow{\sim} X \otimes (Y \otimes Z) & \text{associator} \end{cases}$$

natural
isomorphisms

Requirement: coherent diagrams,

including pentagon and hexagon diagrams

Examples:

- $(\text{Vect}_k, \otimes, k)$

- $(\text{Space}, \times, *)$

- $(\text{Sp}, \wedge, \$)$

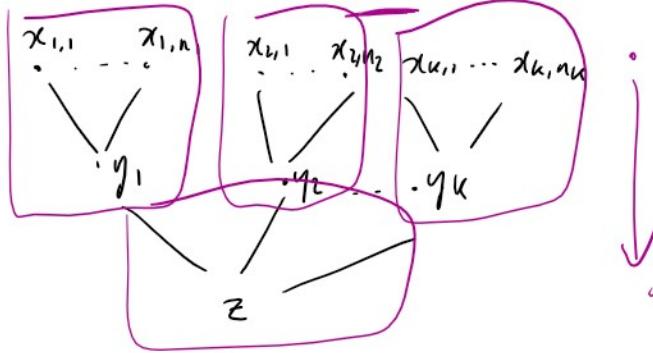
Multicategory \mathcal{M}

Objects: x, y

Multimorphisms: $\text{Mul}(x_1, \dots, x_n; y)$

$\sum n_i = n \geq 0$

compositions



Example: $\underline{\text{Mul}}(V_1, V_2; W) = \{ \text{bilinear maps } V_1 \times V_2 \rightarrow W \}$

example. $\underline{\text{Mul}(V_1, V_2; W)} = \left\{ \begin{array}{l} \text{bilinear maps } V_1 \times V_2 \rightarrow W \\ \text{linear maps } V_1 \otimes V_2 \rightarrow W \end{array} \right\}$

This is general:

$$\text{s.m.c} \longrightarrow \text{multi-cat}$$

- $(\mathcal{C}, \otimes, I)$ $\longrightarrow \mathcal{C}^\otimes$

What if we consider (x_1, \dots, x_n) as a single object?

$$\text{multi-cat} \longrightarrow \text{categories} / \text{Fin}_*$$

$$M \longrightarrow P: M^\otimes \rightarrow \text{Fin}_*$$

M^\otimes	obj (x_1, \dots, x_n) , $n \geq 0$, $x_i \in M$	mor $(x_1, \dots, x_n) \rightarrow (y_1, \dots, y_m)$
-------------	---	--

consists of

$$\left\{ \begin{array}{l} \alpha: \{0, 1, \dots, n\} \rightarrow \{0, 1, \dots, m\} \\ f_j \in \text{Mul}(x_{\alpha^{-1}(j)}; y_j) \\ (1 \leq j \leq m) \end{array} \right.$$

$$\text{Fin}_*: \text{obj} : \langle n \rangle = \{0, 1, \dots, n\}$$

mor: based maps

Why introduce the base point 0 ?

$\alpha: \langle n \rangle \rightarrow \langle m \rangle$
can "forget" elements
in $\langle n \rangle$.

• Important to forget points.

When and how to recover $(\mathcal{C}, \otimes, I)$ from $P: \mathcal{C}^\otimes \rightarrow \text{Fin}_*$?

Answer:



$\{\text{S.m.c up to symmetric monoidal equivalence}\}$



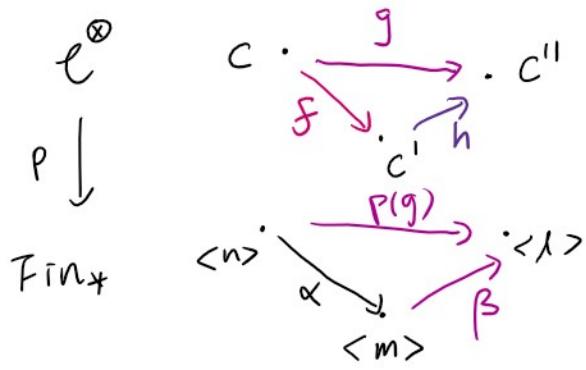
coCartesian condition

$\{P: \mathcal{D} \rightarrow \text{Fin}_\times \text{ satisfying } (M1) \text{ and } (M2)\}$

Segal condition

(M1) P is an op-fibration of categories

meaning that $\forall \alpha, \exists f$, such that $\forall c$, and c'' , such that $\forall c''$, and g, β with $P(g) = \beta \circ \alpha$



$\exists ! h$ with $P(h) = \beta$ and $h \circ f = g$.

f is "P-coCartesian"

(Content of M1) for $p: \mathcal{D} \rightarrow \text{Fin}_\times$

$\alpha: <n> \rightarrow <m>$ in Fin_\times

$\Rightarrow \alpha_!: \mathcal{D}_{<n>} \rightarrow \mathcal{D}_{<m>} \text{ well defined up to}$ Canonical isomorphism

Verification for $\mathcal{D} = \ell^\otimes$:

$$\ell_{<n>}^\otimes = \ell^n$$

$$\alpha: <n> \rightarrow <m>$$

projections to 1
injections

$\alpha_!$ induced by $\ell^k \rightarrow \ell$

$$(x_1, \dots, x_k) \mapsto x_1 \otimes x_2 \otimes \dots \otimes x_k$$

$$\text{and } * \rightarrow \ell$$

$$* \mapsto \underline{I}$$

(M2) Segal condition

- $\rho^i : \langle n \rangle \rightarrow \langle 1 \rangle$

$$\begin{array}{ccc} i & \mapsto & 1 \\ 0.w & \mapsto & 0 \end{array}$$

- $\mathcal{D}_{\langle n \rangle} \xrightarrow[\sim]{(p_1^1, \dots, p_1^n)} \mathcal{D}_{\langle 1 \rangle}^n$

$$e^n = e^n$$

Verification for $\mathcal{D} = e^\otimes$:

$$\begin{aligned} p_1^i : e_{\langle n \rangle}^\otimes = e^n &\longrightarrow e_{\langle 1 \rangle}^\otimes = e \\ (x_1, \dots, x_n) &\mapsto x_i \end{aligned}$$

Content of (M2) :

- $\mathcal{D}_{\langle n \rangle} \xrightarrow[\sim]{} \mathcal{D}_{\langle 1 \rangle}^n$

not equal, but equivalent by built in data

To recover from $p : \mathcal{D} \rightarrow \text{Fin}_X$ the s.m.c. e

$$(M2) \Rightarrow e = \mathcal{D}_{\langle 1 \rangle}$$

- $\mathcal{D}_{\langle 0 \rangle} \cong *$

$$\langle 0 \rangle \rightarrow \langle 1 \rangle \Rightarrow I \rightarrow e$$

- $\alpha : \langle 2 \rangle \rightarrow \langle 1 \rangle \quad 1, 2 \mapsto 1$

$$\Rightarrow \underline{\otimes} : \mathcal{D}_{\langle 1 \rangle}^2 \xleftarrow[\sim]{(M2)} \mathcal{D}_{\langle 2 \rangle} \xrightarrow{\alpha!} \mathcal{D}_{\langle 1 \rangle}$$

- $\beta_{x,y}$ is given by $\delta : \langle 2 \rangle \rightarrow \langle 2 \rangle$

... .

Summary

	classical	new
Data	<ul style="list-style-type: none"> ℓ \otimes, I α_x $\beta_{x,y}$ $\gamma_{x,y,z}$ 	$P: \mathcal{D} \rightarrow \text{Fin}_*$ (M1) op-fibration \Rightarrow (M2) segal
Requirement	<ul style="list-style-type: none"> diagrams commute (e.g. MacLane pentagons) 	remark active inert (Co)Cartesian fibration

Advantage of the new perspective:

- ① simpler data
- ② easier to generalize to ∞ -categories

generalization of (M1)

Def. A symmetric monoidal ∞ -category is a (Co)Cartesian fibration $P: \ell^{\otimes} \rightarrow N(\text{Fin}_*)$

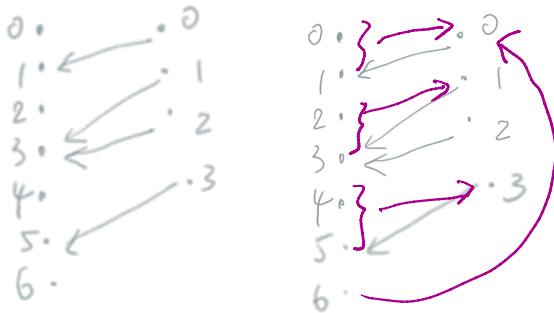
such that $(*) P_i^i$ induced by P^i determines an equivalence $\ell_{\langle n \rangle}^{\otimes} \simeq (\ell_{\langle 1 \rangle}^{\otimes})^n$.

Homework:

- ① Is $\text{id}: N(\text{Fin}_*) \rightarrow N(\text{Fin}_*)$ a symmetric monoidal ∞ -category? obsolete
- ② $\Delta^{\text{op}} \rightarrow \text{Fin}_*$ Segal's Γ -space in:
 $[m] \mapsto \langle m \rangle$ (categories and cohomology theories)

$[m] \mapsto \langle m \rangle$ (categories and cohomology theories)

$([m] \leftarrow [n]) \mapsto (\langle m \rangle \rightarrow \langle n \rangle)$ (In fact, this is s^1)



Is $N(\mathcal{O}^\otimes) \rightarrow N(Fin_*)$ a symmetric monoidal ∞ -category? An ∞ -operad?

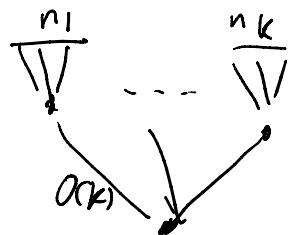
(No)

§2. Operads

data: sets $\mathcal{O}(n) \supseteq \sum_n$, $n \geq 0$

- $\eta: * \rightarrow \mathcal{O}(1)$ parameters $\text{id}: A \rightarrow A$
- $\circ: \mathcal{O}(k) \times \mathcal{O}(n_1) \times \dots \times \mathcal{O}(n_k) \rightarrow \mathcal{O}(n_1 + \dots + n_k)$

requirement: unital, equivariance, associativity diagrams.



Quick observation:

- an operad = a multi-category with one object.

$$\underline{\mathcal{O}(n)} = \text{Mul}(\underbrace{*; \dots; *}_n; *) \quad \text{including } n=0$$

\Rightarrow can form $p: \mathcal{O}^\otimes \rightarrow Fin_*$

$$\begin{cases} \text{obj: } \langle n \rangle = (*; \dots; *) & n \geq 0 \\ \text{mor: } \alpha: \langle m \rangle \rightarrow \langle n \rangle \text{ and } \left\{ x_j \in \mathcal{O}(n+1) \mid 1 \leq j \leq n \right\} \end{cases}$$

Algebra

Examples:

$$\text{Triv} \quad \text{Triv}(n) = \begin{cases} * & n=1 \\ \emptyset & n \neq 1 \end{cases}$$

everything

Triv

$\dots \circ \phi \circ \dots$

everything

Triv^{\otimes} obj : $\langle n \rangle$

mor : $\sigma \in \Sigma_n$

E_0

$E_0(n) = \begin{cases} * & n=0 \text{ or } 1 \\ \emptyset & \text{o.w.} \end{cases}$

based objects

E_0^{\otimes} obj : $\langle n \rangle$ (based)
mor : injections

Comm

Comm(n) = *

Comm^{\otimes} obj : $\langle n \rangle$
mor : based maps

Comm[⊗]
||
Fin*

comm monoids

Ass

Ass(n) = Σ_n

Ass[⊗] skipped

(associative) monoids

Algebra over Θ in $(\mathcal{C}, \otimes, I)$:

- $A \in \mathcal{C}$ with
- data : $\Theta(n) \times A^{\otimes n} \xrightarrow{\Sigma_n} A$
- requirement : unital and associativity diagrams

content : $\Theta(n)$ parametrizes n -ary operations on A .

If we allow $\Theta(n)$ to be spaces, we can parametrize
"unique n -ary operations up to contractible choices".

Remark : We have seen the need for coherence data already

in a symmetric monoidal category : there's

$\int I : * \rightarrow \mathcal{C}$ look like $e^0 \xrightarrow{\sim} e^1$

in a symmetric monoidal category

$$\left\{ \begin{array}{l} I : * \rightarrow \mathcal{C} \\ \otimes : \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C} \end{array} \right. \quad \text{look like} \quad \begin{array}{l} \mathcal{C}^0 \rightarrow \mathcal{C}' \\ \mathcal{C}^2 \rightarrow \mathcal{C}' \end{array}$$

$\Rightarrow \mathcal{C}$ is some sort of algebra in (Cat, \times , $*$)

If it's not a commutative algebra: we have

coherence data $\alpha_{x,y}, \beta_{x,y,z}, \gamma_{x,y,z}$.

In a symmetric monoidal ∞ -category, even more coherence data is introduced.

		cat	∞ -cat
		comm algebra	
-		Symmetric monoidal category	
-	+	Σ_∞ -algebra	Symmetric monoidal ∞ -category
			Comm algebra in S.m. ∞ -cat

Fact. $\underline{\mathcal{C}\text{Alg}(\mathcal{C}_\infty^\times)} \simeq \mathcal{C}_\infty^{\text{smon}}$

$$\mathcal{C}\text{Alg}(\mathcal{C}_\infty^\times) \simeq \mathcal{C}_\infty^{\text{smon}}$$

Homework.

③ \mathcal{C} : symmetric monoidal category. Find out the p-col cartesian edges in $p : \mathcal{C}^\otimes \rightarrow \text{Fin}_\infty$

④ Θ : operad. Do the same for $p : \Theta^\otimes \rightarrow \text{Fin}_\infty$.

§3. ∞ -operads.

Recall

Def. A symmetric monoidal ∞ -category is

a cocartesian fibration $p: \mathcal{C}^{\otimes} \rightarrow N(Fin_*)$

such that (*) P_i^i induced by p^i determines an equivalence $\mathcal{C}_{\langle n \rangle}^{\otimes} \simeq (\mathcal{C}_{\langle 1 \rangle}^{\otimes})^n$.

In Fin_*

Def. $\alpha: \langle m \rangle \rightarrow \langle n \rangle$ is inert if

$$|\alpha^{-1}\{i\}| = 1 \quad 1 \leq i \leq n$$



Example. $P^i: \langle n \rangle \rightarrow \langle 1 \rangle$ "coordinates"

Def. An ∞ -operad is $p: \mathcal{O}^{\otimes} \rightarrow N(Fin_*)$

between ∞ -categories such that

(1) A inert $\alpha: \langle m \rangle \rightarrow \langle n \rangle$, $C \in \mathcal{O}_{\langle m \rangle}^{\otimes}$,
exists p-cocartesian lift $\bar{\alpha}$.

(2) A $f: \langle m \rangle \rightarrow \langle n \rangle$, $C \in \mathcal{O}_{\langle m \rangle}^{\otimes}$, $C' \in \mathcal{O}_{\langle n \rangle}^{\otimes}$,
choose p-cocartesian lifts of $P^i: \langle n \rangle \rightarrow \langle 1 \rangle$

$$\bar{P}^i: C' \rightarrow C'_i \in \mathcal{O}_{\langle 1 \rangle}^{\otimes}$$

Then component over f comp_f component over $P^i \circ f$ $\text{comp}_{P^i \circ f}$

$$\text{comp}_f = \prod_{i=1}^n (\bar{P}^i \circ f)^{n_i} \quad \text{comp}_{P^i \circ f} = \prod_{i=1}^n P^i \circ f$$

$$\text{Map}_{\mathcal{O}^\otimes}^f(C, C') \xrightarrow{\sim} \prod_{i=1}^n \text{Map}_{\mathcal{O}^\otimes}^{p_i^* f}(C, C'_i)$$

is a homotopy equivalence.

(3) $\forall C_1, \dots, C_n \in \mathcal{O}_{\langle 1 \rangle}^\otimes$, $\exists C \in \mathcal{O}_{\langle n \rangle}$ and
 p -cocartesian morphisms $\bar{p}^i: C \rightarrow C_i$ covering $p^i: \langle n \rangle \rightarrow \langle 1 \rangle$.

Rem (3') p_i^* induces $\phi: \mathcal{O}_{\langle n \rangle}^\otimes \xrightarrow{\sim} (\mathcal{O}_{\langle 1 \rangle}^\otimes)^n \Rightarrow (3)$
(1) $\Rightarrow \phi$ well defined segal condition

(2) $\Rightarrow \phi$ is fully faithful

(3) $\Rightarrow \phi$ is essentially surjective

What's (2) ?

• \mathcal{O} : (ordinary) operad before, $p: \mathcal{O}^\otimes \rightarrow \text{Fin}_*$

$\rightsquigarrow p: N(\mathcal{O}^\otimes) \rightarrow N(\text{Fin}_*)$ ∞ -operad

Verify (2): $f: \langle 4 \rangle \rightarrow \langle 2 \rangle \xrightarrow{p^i} \langle 1 \rangle$

$0:$	$\{$	\rightarrow	0
$1:$	$\{$	\rightarrow	$1, 1$
$2:$	$\{$	\rightarrow	2
$3:$	$\{$	\rightarrow	2
$4:$	$\{$	\rightarrow	1

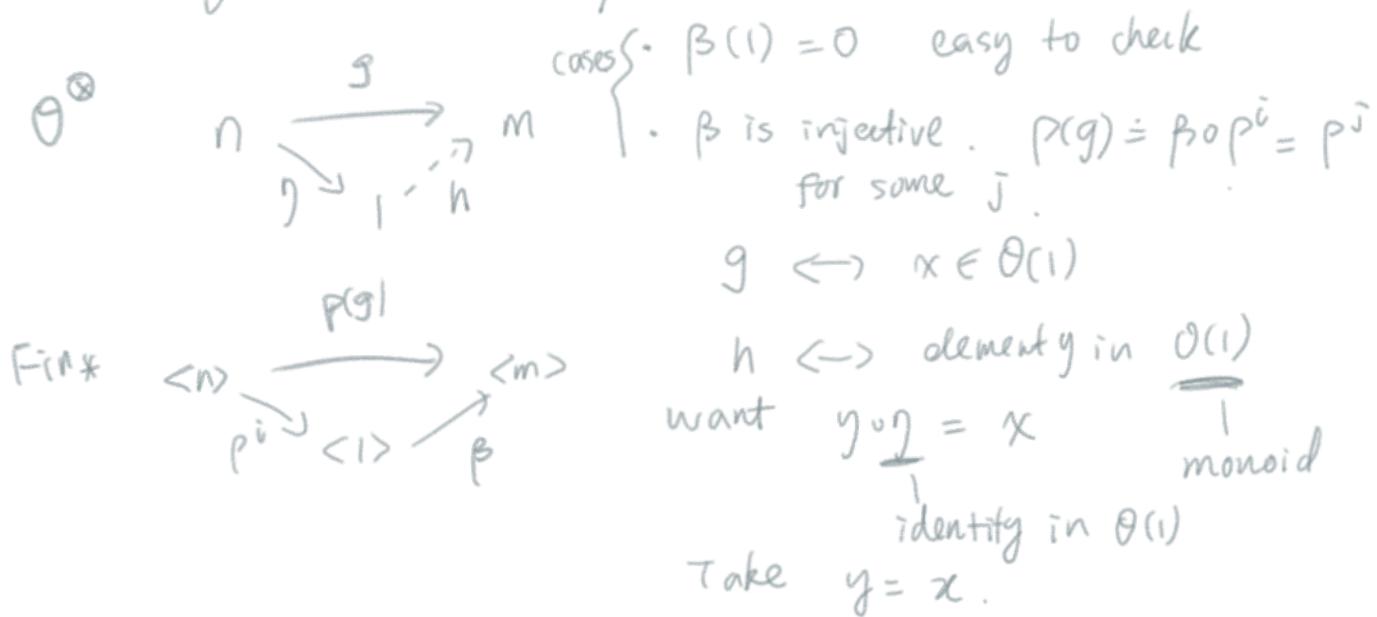
$\text{Map}_{\mathcal{O}^\otimes}^f(4, 2) = \mathcal{O}(1) \times \mathcal{O}(2)$
unique element in $\mathcal{O}_{\langle 4 \rangle}^\otimes$

$$\text{Map}_{\mathcal{O}^\otimes}^{p^i \circ f}(4, 1) = \begin{cases} \mathcal{O}(1) & i=1 \\ \mathcal{O}(2) & i=2 \end{cases}$$

Always pick $\bar{p}^i : n \rightarrow 1$ as $j \in \Theta(1)$.



* j is a cartesian edge



Then \bar{p}^{i_0} is projection onto the component.

(2) is saying that morphisms in Θ^\otimes all are made up by morphisms over $\langle n \rangle \rightarrow \langle 1 \rangle$.

Example. $\text{id}: N(\text{Fin}_\infty) \rightarrow N(\text{Fin}_\infty)$

$\begin{smallmatrix} \parallel \\ \text{Comm}^\otimes \end{smallmatrix}$

Example. $p: \mathcal{C}^\otimes \rightarrow N(\text{Fin}_\infty)$ a symmetric monoidal ∞ -cat.

$\Rightarrow p$ is also an ∞ -operad.

—End