

CS2100

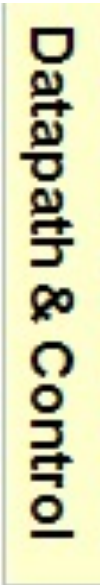
TUTORIAL #1

C AND NUMBER SYSTEMS

(PREPARED BY: AARON TAN)

You will
learn this
in week 5.

What is sign extension?



Q1. Sign extension – extending the sign bit to fill in the bit representation.

Example: from 4 bits to 8 bits.

$$(5)_{10} = (0101)_{2s} = (00000101)_{2s}$$

$$(-3)_{10} = (1101)_{2s} = (11111101)_{2s}$$

Note: Sign extension works for complement systems (1's complement and 2's complement.)

Does it work for sign-and-magnitude system?

Q1. Positive values – straightforward: $(5)_{10} = (0101)_{2s} = (00000101)_{2s}$

Negative values? Why does padding 1s in front work?

$$(-3)_{10} = (\mathbf{1}101)_{2s} = (\underbrace{\mathbf{11111}}_{\text{padding}}101)_{2s}$$

Value: $-2^3 = -8$

Value: $-2^7 + 2^6 + 2^5 + 2^4 + 2^3$
 $= -128 + 64 + 32 + 16 + 8 = -8$

$$-2^k$$

$$\begin{aligned} & -2^m + (2^{m-1} + \dots + 2^{k+1} + 2^k) \\ &= -2^m + 2^k(2^{m-k} - 1) \\ &= -2^m + 2^m - 2^k \\ &= -2^k \end{aligned}$$

Recall that sum of a GP = $\frac{a(r^n-1)}{r-1}$.

Here, $a = 2^k$, $n = m - 1 - k + 1 = m - k$, and $r = 2$.

Q2. Performing subtraction in 1's complement.

Strategy: Convert $A - B$ to $A + (-B)$.

Why not perform subtraction directly?

(a) $0101.11 - 010.0101$

Q2. Performing subtraction in 1's complement.

Strategy: Convert $A - B$ to $A + (-B)$.

(b) $010111.101 - 0111010.11$

Q3. Convert the following decimal numbers to fixed-point binary in 2's complement, with 4 bits for the integer portion and 3 bits for the fraction portion.

(a) 1.75

$$0.75 = 0.5 + 0.25 = \frac{1}{2} + \frac{1}{2^2} = (0.1)_2 + (0.01)_2 = (0.11)_2$$

$$\text{Therefore } 1.75 = (0001.110)_2 = \mathbf{(0001.110)_{2s}}$$

or

$0.75 \times 2 =$	1.5
$0.5 \times 2 =$	1.0
end	

Q3. Convert the following decimal numbers to fixed-point binary in 2's complement, with 4 bits for the integer portion and 3 bits for the fraction portion.

(b) -2.5

$$2.5 = (0010.100)_2$$

$$-2.5 = -(0010.100)_2$$

Invert bits in 0010.100 and add 0.001

$$= \mathbf{(1101.100)}_{2s}$$

$$\begin{array}{r} 1101.011 \\ + \quad \quad 1 \\ \hline 1101.100 \end{array}$$

Q3. Convert the following decimal numbers to fixed-point binary in 2's complement, with 4 bits for the integer portion and 3 bits for the fraction portion.

(c) 3.876

$$0.876 = (0.1110)_2 \approx (0.111)_2$$

$$\begin{aligned}\text{Therefore } 3.876 &= (0011.111)_2 \\ &= \mathbf{(0011.111)_{2s}}\end{aligned}$$

$$0.876 \times 2 = 1.752$$

$$0.752 \times 2 = 1.504$$

$$0.504 \times 2 = 1.008$$

$$0.008 \times 2 = 0.016$$

Q3. Convert the following decimal numbers to fixed-point binary in 2's complement, with 4 bits for the integer portion and 3 bits for the fraction portion.

(d) 2.1

$$0.1 = (0.0001)_2 \approx (0.001)_2$$

$$\begin{aligned}\text{Therefore } 2.1 &= (0010.001)_2 \\ &= \mathbf{(0010.001)}_{2s}\end{aligned}$$

$$0.1 \times 2 = 0.2$$

$$0.2 \times 2 = 0.4$$

$$0.4 \times 2 = 0.8$$

$$0.8 \times 2 = 1.6$$

Q3. Using the binary representations you have just derived, convert them back into decimal.

$$\begin{aligned} \text{(a) } 1.75 \\ &= (\mathbf{0001.110})_{2s} \\ &= (0001.110)_2 \end{aligned}$$

$$\begin{aligned} &(0001.110)_2 \\ &= 2^0 + 2^{-1} + 2^{-2} \\ &= 1.75 \end{aligned}$$

$$\begin{aligned} \text{(b) } -2.5 \\ &= (\mathbf{1101.100})_{2s} \\ &= -(0010.100)_2 \end{aligned}$$

$$\begin{aligned} &-(0010.100)_2 \\ &= -(2^1 + 2^{-1}) \\ &= -2.5 \end{aligned}$$

$$\begin{aligned} \text{(c) } 3.876 \\ &= (\mathbf{0011.111})_{2s} \\ &= (0011.111)_2 \end{aligned}$$

$$\begin{aligned} &(0011.111)_2 \\ &= 2^1 + 2^0 \\ &\quad + 2^{-1} + 2^{-2} + 2^{-3} \\ &= 3.875 \end{aligned}$$

$$\begin{aligned} \text{(d) } 2.1 \\ &= (\mathbf{0010.001})_{2s} \\ &= (0010.001)_2 \end{aligned}$$

$$\begin{aligned} &(0010.001)_2 \\ &= 2^1 + 2^{-3} \\ &= 2.125 \end{aligned}$$

Q4. How would you represent the decimal value **-0.078125** in the IEEE 754 single-precision representation? Express your answer in hexadecimal.

$$-0.078125 = -(0.000101)_2 = -(1.01)_2 \times 2^{-4}$$

$$\begin{aligned}\text{Exponent: } & -4 + 127 \\ & = 123\end{aligned}$$

$$= (01111011)_2$$



1 01111011 01000000000000000000000000000000

Below the bit string, curly braces group the bits into hexadecimal digits:

B D A 0 0 0 0 0

Q6. Trace the program manually and write out its output.

➔ `int a = 3, *b, c, *d, e, *f;`

➔ `b = &a;`

➔ `*b = 5;`

➔ `c = *b * 3;`

➔ `d = b;`

This does NOT make d point to b!
This copies the content of b into d.

➔ `e = *b + c;`

➔ `*d = c + e;`

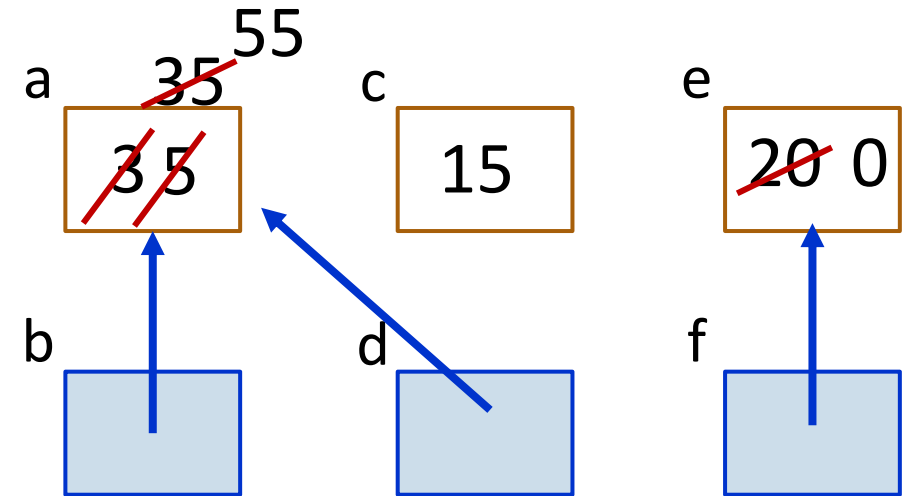
➔ `f = &e;`

➔ `a = *f + *b;`

➔ `*f = *d - *b;`

➔ `printf("a = %d, c = %d, e = %d\n", a, c, e);`

➔ `printf("*b = %d, *d = %d, *f = %d\n", *b, *d, *f);`



Output:

`a = 55, c = 15, e = 0`

`*b = 55, *d = 55, *f = 0`

END OF FILE

For reference.

1s complement

- Given a number x which can be expressed as an n -bit binary number, its negated value can be obtained in **2s-complement** representation using: $-x = 2^n - x - 1$
- Example: $x=12$; Assume 8 bits, $-x = 2^8 - 12 - 1 = 243 = (11110011)_{1s}$
- Technique: **invert the bits**. $12 = (00001100)_{1s} \rightarrow -12 = (11110011)_{1s}$

2s complement

- Given a number x which can be expressed as an n -bit binary number, its negated value can be obtained in **2s-complement** representation using: $-x = 2^n - x$
- Example: $x=12$; Assume 8 bits, $-x = 2^8 - 12 = 244 = (11110100)_{2s}$
- Technique: **invert the bits then plus 1**. $12 = (00001100)_{2s} \rightarrow -12 = (11110100)_{2s}$

For reference.

10.4 Comparisons

Important!

4-bit system

Positive values

Value	Sign-and-Magnitude	1s Comp.	2s Comp.
+7	0111	0111	0111
+6	0110	0110	0110
+5	0101	0101	0101
+4	0100	0100	0100
+3	0011	0011	0011
+2	0010	0010	0010
+1	0001	0001	0001
+0	0000	0000	0000

Negative values

Value	Sign-and-Magnitude	1s Comp.	2s Comp.
-0	1000	1111	-
-1	1001	1110	1111
-2	1010	1101	1110
-3	1011	1100	1101
-4	1100	1011	1100
-5	1101	1010	1011
-6	1110	1001	1010
-7	1111	1000	1001
-8	-	-	1000