
CS2040S Data Structures and Algorithms

Lecture Note #4

Sorting

Objectives

1

- To learn some classic sorting algorithms

2

- To analyse the running time of these algorithms

3

- To learn concepts such as in-place sorts and stable sorts

4

- Using Java methods to perform sorting

Programs used in this lecture

- SelectionSort.java
- BubbleSort.java, BubbleSortImproved.java
- InsertionSort.java
- MergeSort.java
- QuickSort.java
- Sort.java, Sort2.java
- Person.java, AgeComparator.java, NameComparator.java, TestComparator.java

Why Study Sorting?

- When an input is sorted by some **sort key**, many problems become easy (eg. searching, min, max, k^{th} smallest, etc.)

Q: What is a sort key? the field/attribute chosen to sort the items by

- Sorting has a variety of interesting algorithmic solutions, which embody many ideas:
 - **Internal** sort vs **external** sort
 - **Iterative** vs **recursive**
 - **Comparison** vs **non-comparison** based
 - **Divide-and-conquer**
 - **Best/worst/average** case time complexity bounds

Sorting applications

- Uniqueness testing
- Deleting duplicates
- Frequency counting
- Efficient searching
- Dictionary
- Telephone/street directory
- Index of book
- Author index of conference proceedings
- etc.

Outline

- *Comparison based and Iterative algorithms*
 1. Selection Sort
 2. Bubble Sort
 3. Insertion Sort
- *Comparison based and Recursive algorithms*
 4. Merge Sort
 5. Quick Sort
- *Non-comparison based*
 6. Radix Sort
- 7. Comparison of Sort Algorithms
 - In-place sort
 - Stable sort
- 8. Use of Java Sort Methods

Note: In the lecture, we consider only sorting in **ascending order** of data.

1 Selection Sort

1 Idea of Selection Sort

- Given an array of n items
 1. Find the **largest** item.
the largest item will be at the end of the array
 2. **Swap** it with the item at the **end** of the array.
 3. Go to step 1 by excluding the largest item from the array.

1 Selection Sort of 5 integers

29	10	14	37	13
----	----	----	----	----

37 is the largest, swap it with the last element, i.e. **13**.

Q: How to find the largest?

do a linear search, scan through to find the value and index of the largest

29	10	14	13	37
----	----	----	----	----

13	10	14	29	37
----	----	----	----	----

array of size n only need n-1 passes

13	10	14	29	37
----	----	----	----	----

10	13	14	29	37
----	----	----	----	----

Sorted!

1 Code of Selection Sort

```
public static void selectionSort(int[] a) {  
    for (int i = a.length-1; i >= 1; i--) {  
        int index = i; // i is the last item position and  
                       // index is the largest element position  
        // loop to get the largest element  
        for (int j = 0; j < i; j++) {  
            if (a[j] > a[index])  
                index = j; // j is the current largest item  
        }  
        // swap the largest item a[index] with the last item a[i]  
        int temp = a[index]; swap the last item with the end of the current array  
        a[index] = a[i];  
        a[i] = temp;  
    }  
}
```

SelectionSort.java

1 Analysis of Selection Sort

```
public static void selectionSort(int[] a)
{
    for (int i=a.length-1; i>=1; i--) {
        int index = i;
        for (int j=0; j<i; j++) {
            if (a[j] > a[index])
                index = j;
        }
        SWAP( ... )
    }
}
```

t_1 and t_2 = costs of statements in outer and inner blocks.

Number of times the statement is executed:

- $n-1$
- $n-1$
- $(n-1)+(n-2)+\dots+1$
 $= n \times (n-1)/2$ A.P

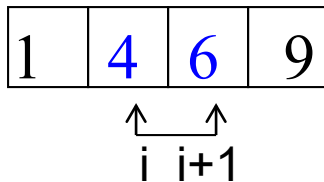
- $n-1$

$$\begin{aligned}\text{Total} &= t_1 \times (n-1) \\ &\quad + t_2 \times n \times (n-1)/2 \\ &= O(n^2)\end{aligned}$$

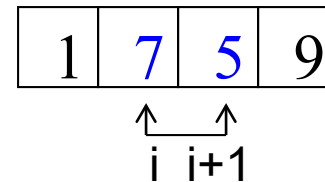
2 Bubble Sort

2 Idea of Bubble Sort

- “Bubble” down the largest item to the end of the array in each iteration by examining the **i-th** and **(i+1)-th** items
- If their values are not in the correct order, i.e. $a[i] > a[i+1]$, **swap** them.



// no need to swap



// not in order, need to swap

2 Example of Bubble Sort

- The first two passes of Bubble Sort for an array of 5 integers

(a) Pass 1

29	10	14	37	13
10	29	14	37	13
10	14	29	37	13
10	14	29	37	13
10	14	29	13	37

At the end of **pass 1**, the largest item **37** is at the last position.

(b) Pass 2

10	14	29	13	37
10	14	29	13	37
10	14	29	13	37
10	14	13	29	37

At the end of **pass 2**, the second largest item **29** is at the second last position.

No. of passes needed : $n - 1$

2 Code of Bubble Sort

```
public static void bubbleSort(int[] a) {  
    for (int i = 1; i < a.length; i++) {  
        for (int j = 0; j < a.length - i; j++) {  
            if (a[j] > a[j+1]) { // the larger item bubbles down (swap)  
                int temp = a[j];  
                a[j] = a[j+1];  
                a[j+1] = temp;  
            }  
        }  
    }  
}
```

BubbleSort.java

2 Analysis of Bubble Sort

- 1 iteration of the inner loop (test and swap) requires time bounded by a constant **c**
- Doubly nested loops:
 - **Outer loop:** exactly $n-1$ iterations
 - **Inner loop:**
 - When $i=1$, $(n-1)$ iterations
 - When $i=2$, $(n-2)$ iterations
 - ...
 - When $i=(n-1)$, 1 iteration
- Total number of iterations = $(n-1) + (n-2) + \dots + 1$
= $n \times (n-1) / 2$
- Total time = **c** $\times n \times (n-1) / 2 = O(n^2)$

```
public static void bubbleSort(int[] a) {  
    for (int i = 1; i < a.length; i++) {  
        for (int j = 0; j < a.length - i; j++) {  
            if (a[j] > a[j+1]) { // (swap)  
                int temp = a[j];  
                a[j] = a[j+1];  
                a[j+1] = temp;  
            }  
        }  
    }  
}
```


2 Bubble Sort can be improved

- Given a sorted input, Bubble Sort still requires $O(n^2)$ to sort.
- It does not make an effort to check whether the input has been sorted.
- Thus it can be improved by using a **flag**, **isSorted**, as follows (next slide):

2 Code of Bubble Sort (Improved version)

```
public static void bubbleSort2(int[] a) {  
    for (int i = 1; i < a.length; i++) {  
        → boolean isSorted = true; // isSorted = true if a[] is sorted  
        for (int j = 0; j < a.length-i; j++) {  
            if (a[j] > a[j+1]) { // the larger item bubbles up  
                int temp = a[j]; // and isSorted is set to false,  
                a[j] = a[j+1]; // i.e. the data was not sorted  
                a[j+1] = temp;  
                → isSorted = false; if there has been a swap between values,  
                                   means we cant determine if its sorted or  
                                   not yet so set isSorted to false  
            }  
        }  
        if there has been no swap between values, set isSorted to true  
        → if (isSorted) return; // why?  
    }  
}
```

BubbleSortImproved.java

2 Analysis of Bubble Sort (Improved version)

- **Worst case** 1. reverse sorted 2. smallest value at largest index and everything else sorted alr
 - ❑ Input in **descending order** (any other situation?)
 - ❑ How many iterations in the outer loop are needed?
Answer: **$n-1$** iterations
 - ❑ Running time remains the same: **$O(n^2)$**
- **Best case**
 - ❑ Input is already in **ascending order**
 - ❑ The algorithm returns after a **single iteration** in the outer loop. (Why?)
 - ❑ Running time: **$O(n)$** just do one pass through the array -> alr sorted

3 Insertion Sort

3 Idea of Insertion Sort

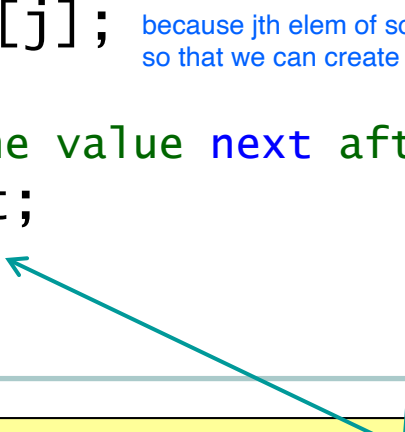
- Arranging a hand of poker cards
 - Start with one card in your hand
 - Pick the next card and **insert** it into its **proper sorted order**
 - Repeat previous step for all the rest of the cards

3 Example of Insertion Sort

- $n = 4$
 - Given a seq: 40 13 20 8
 - $i=1$ 13 40 20 8
 - $i=2$ 13 20 40 8
 - $i=3$ 8 13 20 40
- how to insert 20 in correct pos?
need to create a gap in between 13 and 40
how to create a gap?
take 40 and copy it to the right,
since $13 < 20$, copy 20 where 40 was previously
- n = no of items to be sorted
 - $S1$ = sub-array sorted so far
 - $S2$ = elements yet to be processed
 - In each iteration, how to insert the next element into $S1$ efficiently?

3 Code of Insertion Sort

```
public static void insertionSort(int[] a) {  
    for (int i=1;i<a.length;i++) { //Q: why i starts from 1?  
        // a[i] is the next data to insert index 0 is alr in S1  
        int next = a[i]; declare next cause a[i] needs to be stored as it changes with each iteration  
        // scan backwards to find a place. Q: why not scan forwards?  
        int j; // Q: why is j declared here?  
        // Q: what if a[j] <= next? dont do anything because it is > than largest element of sorted arr -> new largest  
        for (j=i-1; j>=0 && a[j]>next; j--)  
            a[j+1] = a[j]; because jth elem of sorted array is larger than next, we copy over to the right so that we can create a gap for next to be inserted  
        // Now insert the value next after index j at the end of loop  
        a[j+1] = next;  
    }  
}
```



InsertionSort.java

Q: Can we replace these two “next” with a[i]?

Ans: No! (Why?)

3 Analysis of Insertion Sort

- Outer loop executes exactly $n-1$ times
- Number of times inner loop executes depends on the inputs:
 - **Best case:** array already sorted, hence $(a[j] > \text{next})$ is always false
 - No shifting of data is necessary; Inner loop not executed at all.
 - **Worst case:** array reversely sorted, hence $(a[j] > \text{next})$ is always true
 - Need i shifts for $i = 1$ to $n-1$.
 - Insertion always occurs at the front.
- Therefore, the **best case** running time is $O(n)$. (Why?)
- The **worst case** running time is $O(n^2)$. (Why?)

```
... insertionSort(int[] a) {  
    for (int i=1; i<a.length; i++) {  
        int next = a[i];  
        int j;  
        for (j=i-1; j>=0 && a[j]>next; j--)  
            a[j+1] = a[j];  
  
        a[j+1] = next;  
    }  
}
```


4 Merge Sort

4 Idea of Merge Sort (1/3)

- Suppose we **only know how to merge** two sorted lists of elements into one combined list
- Given an unsorted list of n elements
- Since each element is a sorted list, we can repeatedly...
 - **Merge** each pair of lists, each list containing one element, into a sorted list of 2 elements.
 - **Merge** each pair of sorted lists of 2 elements into a sorted list of 4 elements.
 - ...
 - The final step **merges** 2 sorted lists of $n/2$ elements to obtain a sorted list of n elements.

4 Idea of Merge Sort (2/3)

- **Divide-and-conquer** method solves problem by three steps:
 - **Divide Step:** divide the larger problem into smaller problems.
 - **(Recursively)** solve the smaller problems.
 - **Conquer Step:** combine the results of the smaller problems to produce the result of the larger problem.

4 Idea of Merge Sort (3/3)

- **Merge Sort** is a divide-and-conquer sorting algorithm
 - **Divide Step:** Divide the array into two (equal) halves.
 - **(Recursively)** Merge sort the two halves.
 - **Conquer Step:** Merge the two sorted halves to form a sorted array.
- Q: What are the base cases?

4 Example of Merge Sort

7	2	6	3	8	4	5
---	---	---	---	---	---	---

Divide into
two halves

7	2	6	3
---	---	---	---

8	4	5
---	---	---

Recursively
sort the halves

2	3	6	7
---	---	---	---

4	5	8
---	---	---

Merge the halves

2	3	4	5	6	7	8
---	---	---	---	---	---	---

4 Code of Merge Sort

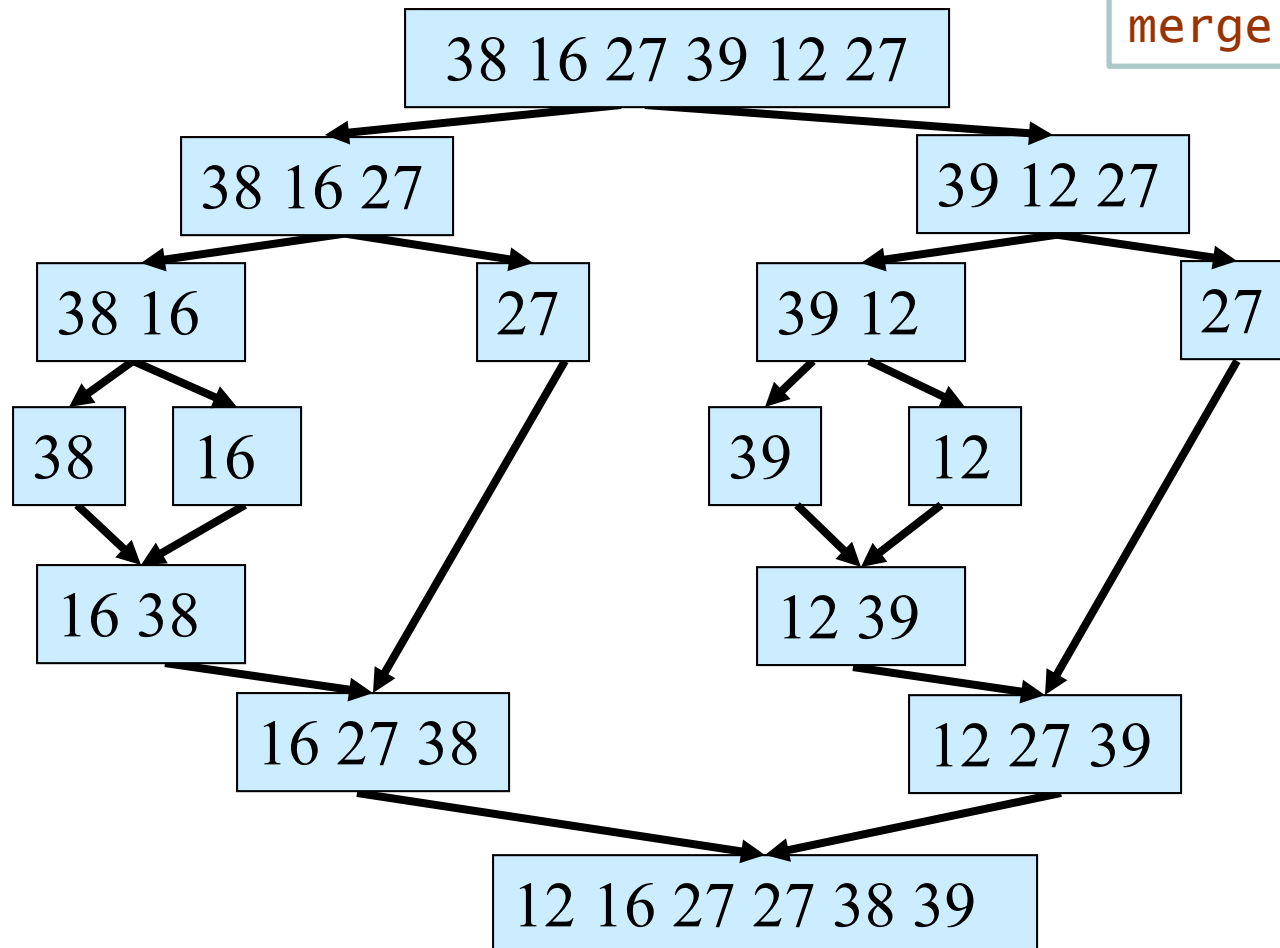
i = left most index, j = right most index

```
... mergeSort(int[] a, int i, int j) {  
    // to sort data from a[i] to a[j], where  $i < j$   
    if ( $i < j$ ) { // Q: what if  $i \geq j$ ? if left most = right most it means that array  
                                                only has one element -> nothing to sort  
        int mid = (i+j)/2; // divide  
        mergeSort(a, i, mid); // recursion  
        mergeSort(a, mid+1, j);  
        merge(a, i, mid, j); //conquer: merge  $a[i..mid]$  and  
                               // $a[mid+1..j]$  back into  $a[i..j]$   
    }  
}
```

MergeSort.java

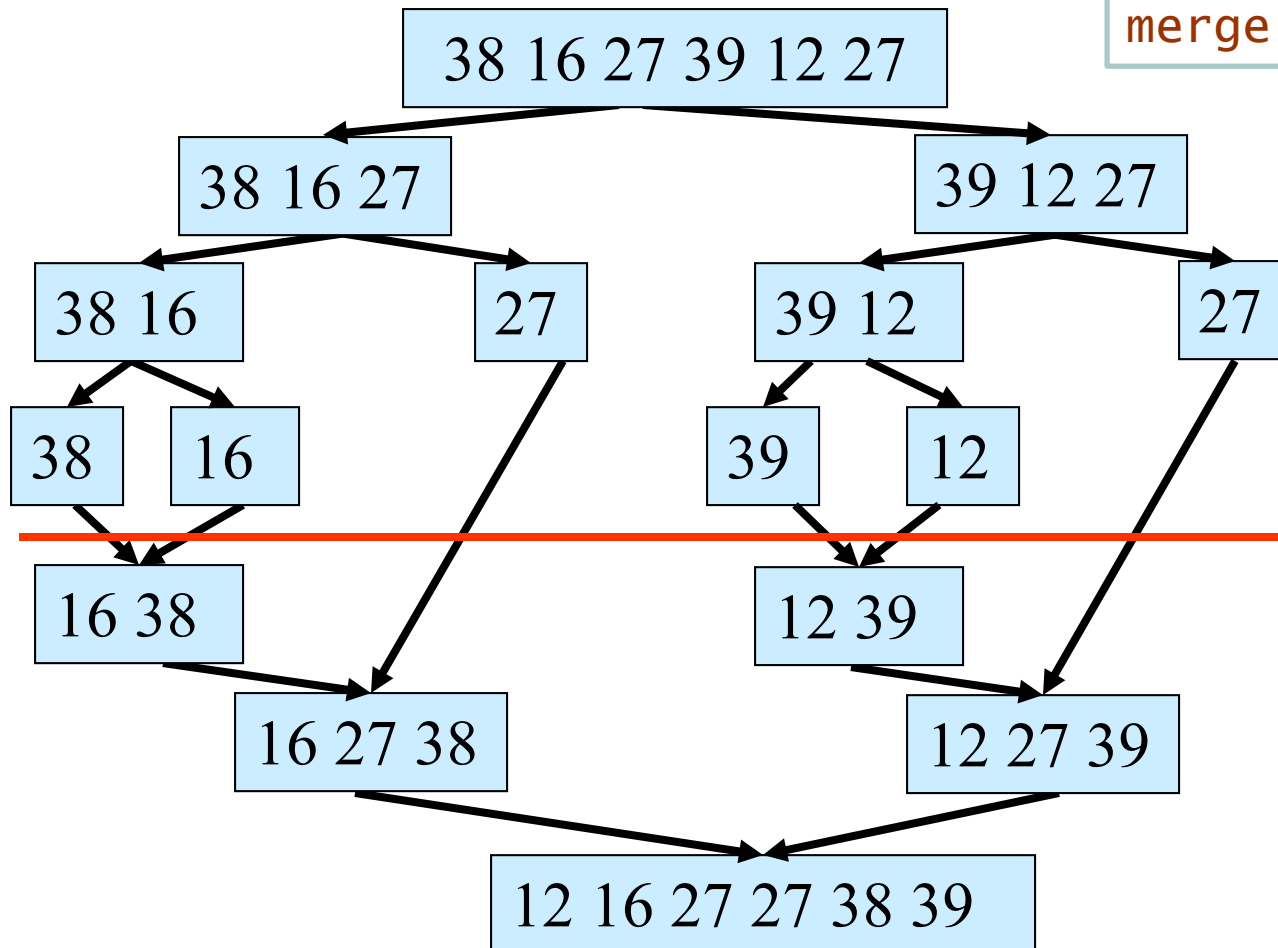
4 Merge Sort of a 6-element Array (1/2)

```
mergeSort(a,i,mid);  
mergeSort(a,mid+1,j);  
merge(a,i,mid,j);
```



4 Merge Sort of a 6-element Array (2/2)

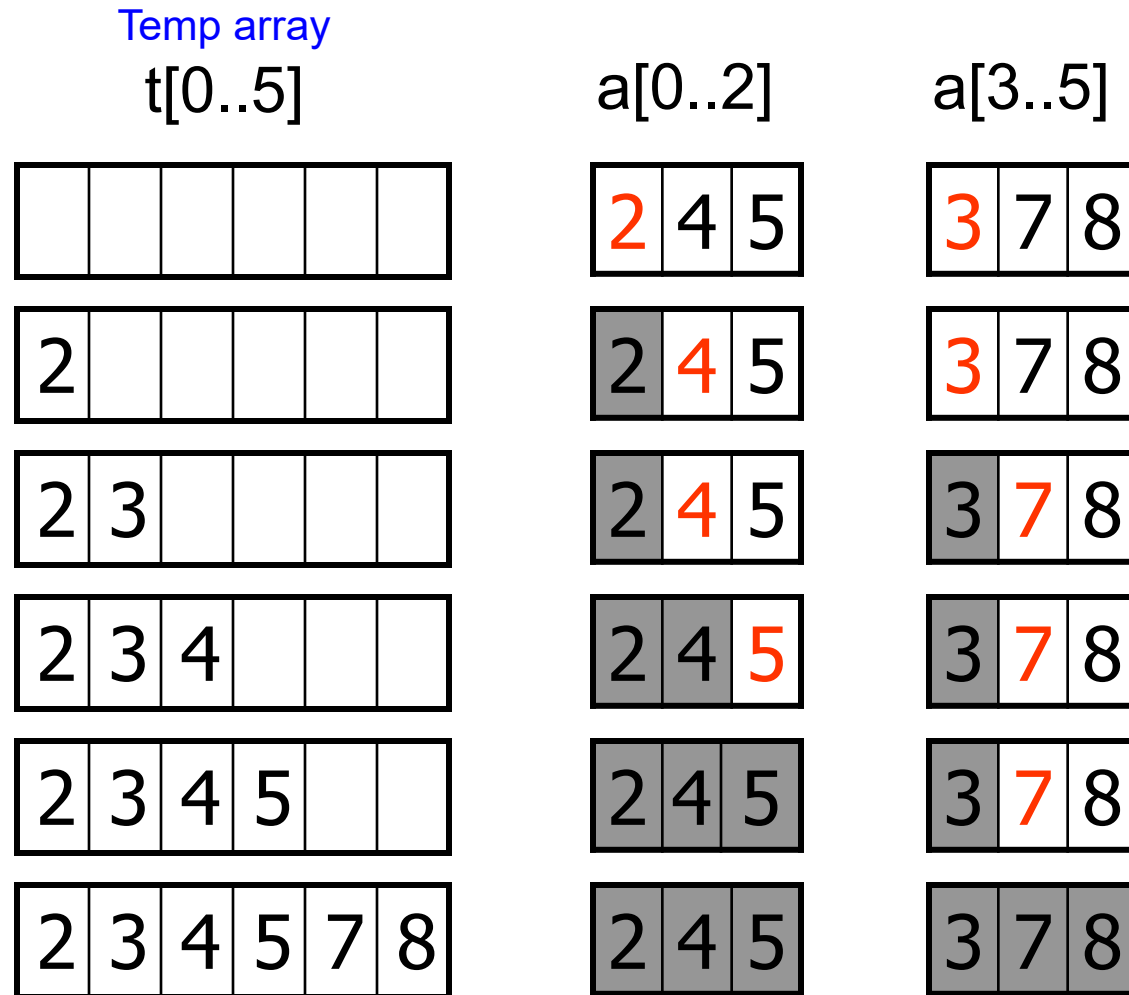
```
mergeSort(a,i,mid);  
mergeSort(a,mid+1,j);  
merge(a,i,mid,j);
```



Divide phase:
Recursive call to
mergeSort

Conquer phase:
Merge steps
The sorting is done
here

4 How to Merge 2 Sorted Subarrays?



4 Merge Algorithm (1/2)

```
... merge(int[] a, int i, int mid, int j) {  
    // Merges the 2 sorted sub-arrays a[i..mid] and  
    // a[mid+1..j] into one sorted sub-array a[i..j]  
  
    int[] temp = new int[j-i+1]; // temp storage  
    int left = i, right = mid+1, it = 0;  
    // it = next index to store merged item in temp[]  
    // Q: what are left and right?  
        left half: i — mid array, right half: mid — j array(index of first element of left and right halves)  
    while (left<=mid && right<=j) { // output the smaller  
        if (a[left] <= a[right]) // while we have not exhausted either one of the arrays  
            temp[it++] = a[left++];  
        else // adding the smallest element of the two into the temp array  
            // and incrementing the index to point to the next element  
            temp[it++] = a[right++];  
    }  
}
```

4 Merge Algorithm (2/2)

have not exhausted either left or right array

```
// Copy the remaining elements into temp. Q: why?  
while (left<=mid) temp[it++] = a[left++];  
while (right<=j)  temp[it++] = a[right++];  
// Q: will both the above while statements be executed?
```

no because the prev while loop only exhausts one of the sub arrays

```
// Copy the result in temp back into  
// the original array a  
for (int k = 0; k < temp.length; k++)  
    a[i+k] = temp[k];  
}
```

4 Analysis of Merge Sort (1/3)

- In Merge Sort, the bulk of work is done in the Merge step
`merge(a, i, mid, j)`
- Total number of items = $k = j - i + 1$
 - Number of comparisons $\leq k - 1$ (Q: Why not = $k - 1$?)
 - Number of moves from original array to temp array = k
 - Number of moves from temp array to original array = k
- In total, number of operations $\leq 3k - 1 = O(k)$
- How many times is `merge()` called?

```
... mergeSort(int[] a, int i, int j) {  
    if (i < j) {  
        int mid = (i+j)/2;  
        mergeSort(a, i, mid);  
        mergeSort(a, mid+1, j);  
        merge(a, i, mid, j);  
    }  
}
```

4 Analysis of Merge Sort (2/3)

Level 0:
Mergesort n items

Level 1:
2 calls to Mergesort $n/2$ items

Level 2:
4 calls to Mergesort $n/2^2$ items

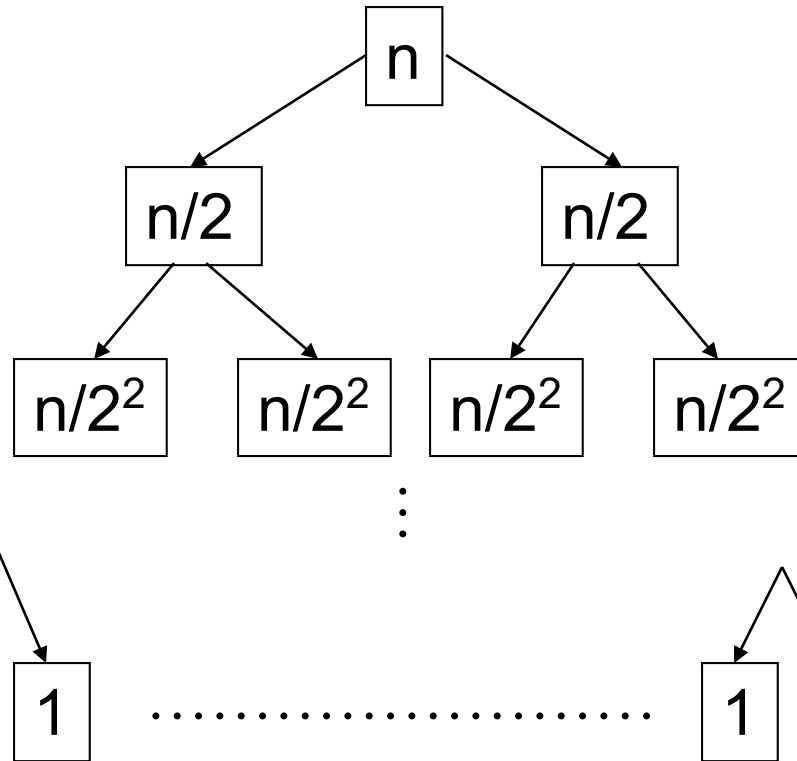
Level ($\log n$):
 n calls to Mergesort 1 item

Level 0:
0 call to Merge

Level 1:
1 calls to Merge

Level 2:
2 calls to Merge

Level ($\log n$):
 $2^{(\log n) - 1} (= n/2)$
calls to Merge



Let h be the maximum level, ie. Mergesort 1 item.
 $n/(2^h) = 1 \quad \rightarrow \quad n = 2^h \quad \rightarrow \quad h = \log n$

4 Analysis of Merge Sort (3/3)

- Level 0: 0 call to Merge
- Level 1: 1 call to Merge with $n/2$ items each,
 $O(1 \times 2 \times n/2) = O(n)$ time
- Level 2: 2 calls to Merge with $n/2^2$ items each,
 $O(2 \times 2 \times n/2^2) = O(n)$ time
- Level 3: 2^2 calls to Merge with $n/2^3$ items each,
 $O(2^2 \times 2 \times n/2^3) = O(n)$ time
- ...
- Level ($\log n$): $2^{(\log n)-1} (= n/2)$ calls to Merge with $n/2^{\log n}$
(= 1) item each,
 $O(n/2 \times 2 \times 1) = O(n)$ time
- In total, running time = $(\log n) \times O(n) = O(n \log n)$

4 Drawbacks of Merge Sort

- Implementation of merge() is not straightforward
- Requires **additional temporary arrays** and to copy the merged sets stored in the temporary arrays to the original array
- Hence, **additional** space complexity = $O(n)$

5 Quick Sort

5 Idea of Quick Sort

- Quick Sort is a **divide-and-conquer** algorithm
- **Divide Step:** Choose a **pivot** item **p** and partition the items of $a[i..j]$ into **2 parts** so that
 - Items in the first part are $< p$, and
 - Items in the second part are $\geq p$.
- **Recursively** sort the 2 parts
- **Conquer Step:** Do nothing! No merging is needed.
- What are the base cases? only one item left

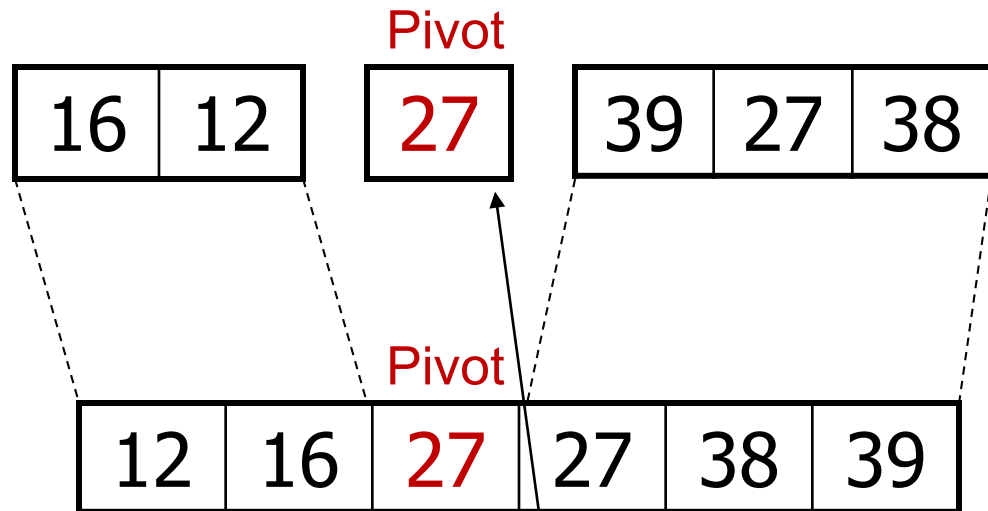
5 Example of Quick Sort

Choose the **1st** item as **pivot**

Pivot

27	38	12	39	27	16
-----------	----	----	----	----	----

Partition $a[]$ about
the pivot 27



Recursively sort
the two parts

Note that after the partition,
the pivot is moved to its **final position**!
No merge phase is needed.

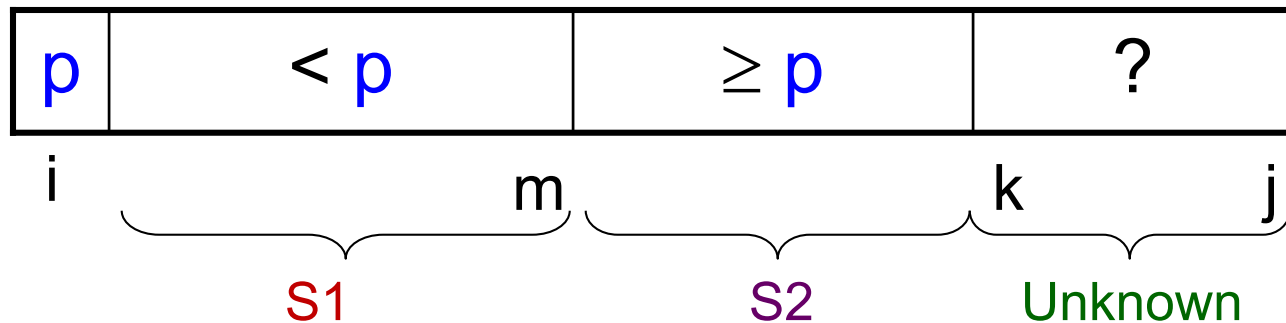
5 Code of Quick Sort

```
... quickSort(int[] a, int i, int j) {  
    if (i < j) { // Q: what if i >= j?  
        int pivotIdx = partition(a, i, j);  
        quickSort(a, i, pivotIdx-1);  
        quickSort(a, pivotIdx+1, j);  
        // No conquer part!  
    }  
}
```

QuickSort.java

5 Partition algorithm idea (1/4)

- To partition $a[i..j]$, we choose $a[i]$ as the **pivot** p .
 - Why choose $a[i]$? Are there other choices?
- The remaining items (i.e. $a[i+1..j]$) are divided into 3 regions:
 - **S1** = $a[i+1..m]$ where items $< p$
 - **S2** = $a[m+1..k-1]$ where item $\geq p$
 - **Unknown** (unprocessed) = $a[k..j]$, where items are yet to be assigned to S1 or S2.



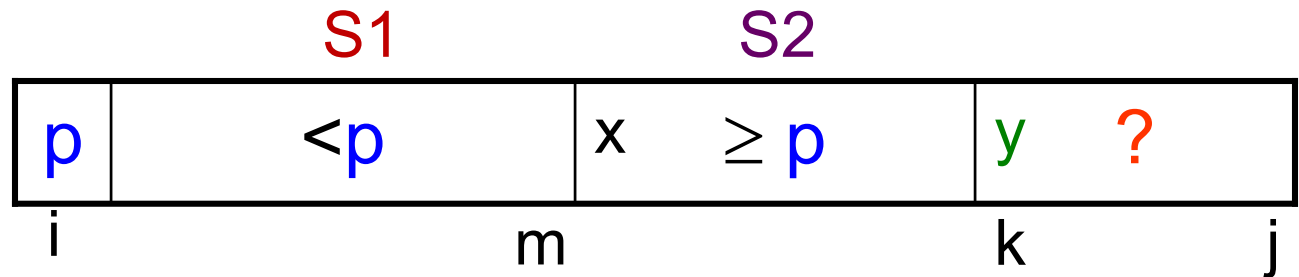
5 Partition algorithm idea (2/4)

- Initially, regions **S1** and **S2** are empty. All items excluding **p** are in the **unknown** region.
- Then, for each item $a[k]$ (for $k=i+1$ to j) in the **unknown** region, compare $a[k]$ with **p**:
 - If $a[k] \geq p$, put $a[k]$ into **S2**.
 - Otherwise, put $a[k]$ into **S1**.
- Q: How about if we change \geq to $>$ in the condition part?

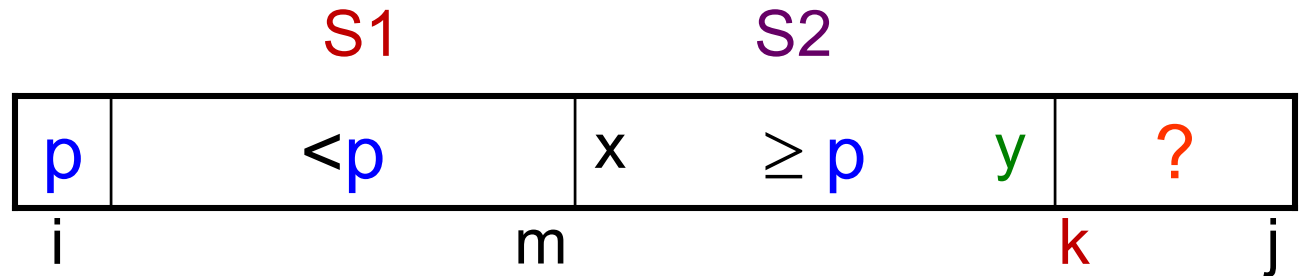
5 Partition algorithm idea (3/4)

■ Case 1:

If $a[k] = y \geq p$,



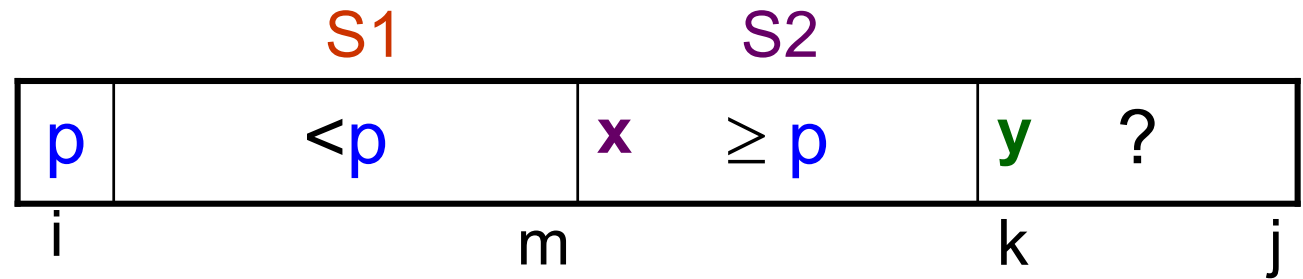
Increment k



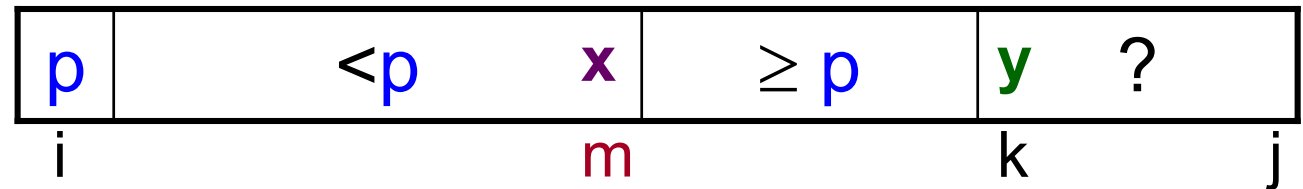
5 Partition algorithm idea (4/4)

■ Case 2:

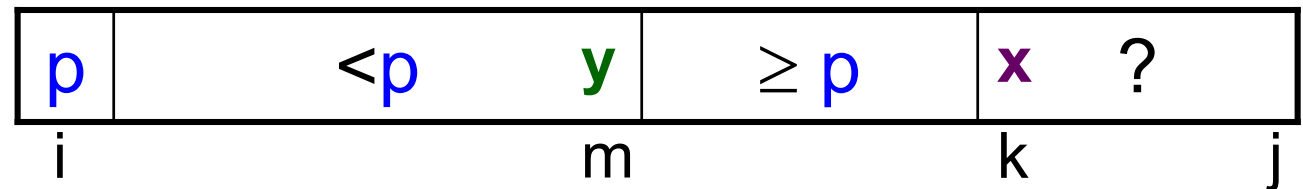
If $a[k]=y < p$



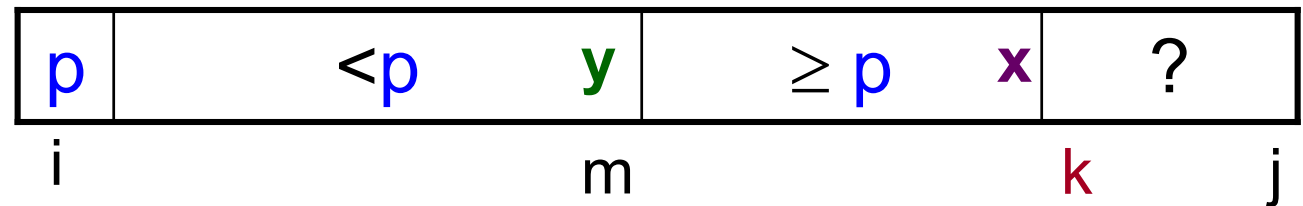
Increment *m*



Swap x and y



Increment *k*



5 Code of Partition Algorithm

```
... partition(int[] a, int i, int j) {  
    // partition data items in a[i..j]  
    int p = a[i]; // p is the pivot, the ith item  
    int m = i;    // Initially S1 and S2 are empty  
    for (int k=i+1; k<=j; k++) { //process unknown region  
        if (a[k] < p) { // case 2: put a[k] to S1  
            m++;  
            swap(a,k,m); // then we increment k but k is incremented in the for loop  
        } else { // case 1: put a[k] to S2. Do nothing!  
        } // else part should be removed since it is empty  
    }  
    swap(a,i,m); // put the pivot at the right place  
    return m;    // m is the pivot's final position  
}
```

- As there is only one 'for' loop and the size of the array is $n = j - i + 1$, so the complexity for partition() is $O(n)$

5 Partition Algorithm: Example

Pivot	Unknown				
27	38	12	39	27	16

Pivot	S_2	Unknown			
27	38	12	39	27	16



Pivot	S_1	S_2	Unknown		
27	12	38	39	27	16

Pivot	S_1	S_2	Unknown		
27	12	38	39	27	16

Pivot	S_1	S_2	Unknown		
27	12	38	39	27	16



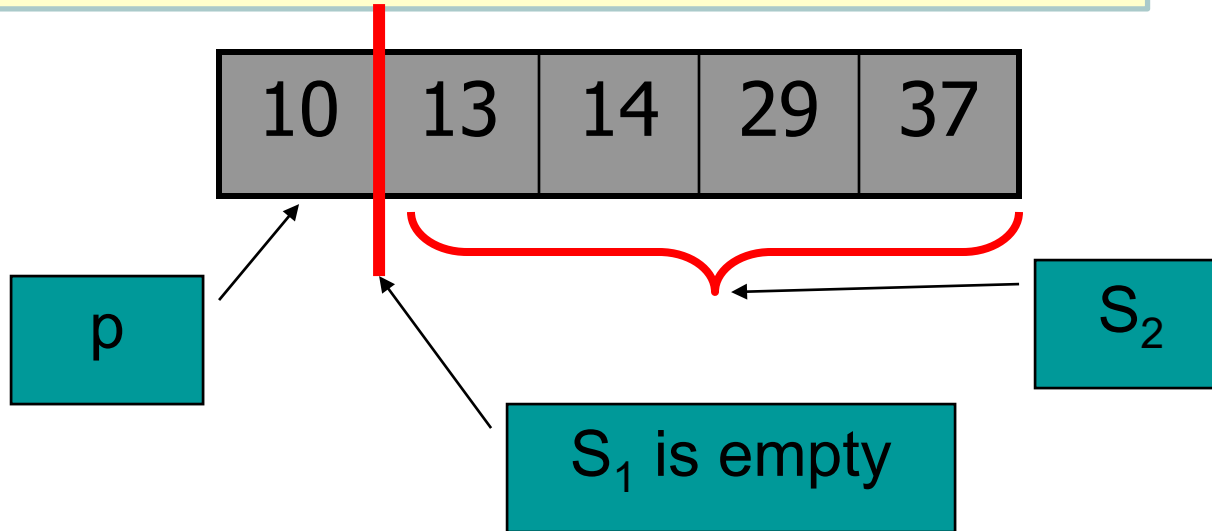
Pivot	S_1	S_2	Unknown		
27	12	16	39	27	38

S_1	Pivot	S_2	Unknown		
16	12	27	39	27	38

Same value, no need to swap them.

5 Analysis of Quick Sort: Worst Case (1/2)

When $a[0..n-1]$ is in increasing order:



What is the index returned by `partition()`?

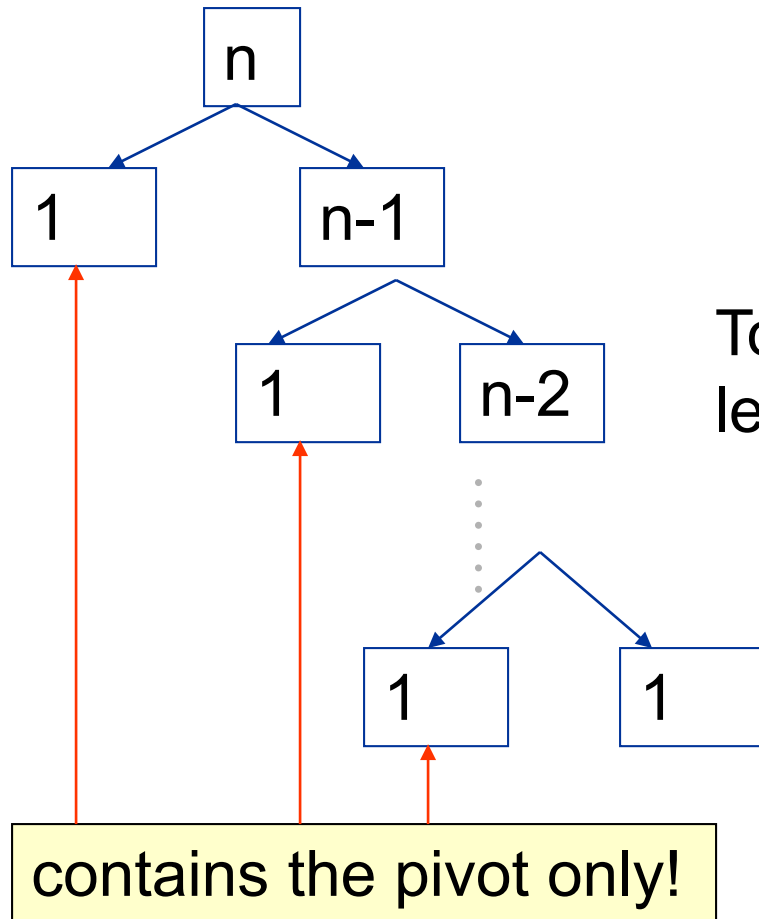
`swap(a,i,m)` will swap the pivot with itself!

The left partition (S_1) is **empty** and

The right partition (S_2) is the rest excluding the pivot.

What if the array is in decreasing order?

5 Analysis of Quick Sort: Worst Case (2/2)



Total no. of
levels = n

As each partition takes linear time, the algorithm in its worst case has n levels and hence it takes time $n + (n-1) + \dots + 1 = O(n^2)$

5 Analysis of Quick Sort: Best/Average case

- **Best case** occurs when partition always splits the array into **2 equal halves**
 - Depth of recursion is **$\log n$** .
 - Each level takes **n** or fewer comparisons, so the time complexity is **$O(n \log n)$**
- In practice, worst case is rare, and on the average, we get some good splits and some bad ones
- **Average time** is **$O(n \log n)$** → *especially true if using randomized pivoting (randomly choose the pivot rather than keep using the 1st element). Can get even better theoretical bounds if use median as pivot (outside scope of this course).

6 Radix Sort

6 Idea of Radix Sort

- Treats each data to be sorted as a **character string**.
- It is not using comparison, i.e., **no comparison** among the data is needed.
- Hence it is a **non-comparison based sort** (the preceding sorting algorithms are called comparison based sorts)
- In each iteration, organize the data into groups according to the **next** character in each data.

6 Radix Sort of Eight Integers

0123, 2154, 0222, 0004, 0283, 1560, 1061, 2150

Original integers

(156**0**, 215**0**) (106**1**) (022**2**) (012**3**, 028**3**) (215**4**, 000**4**)

Grouped by fourth digit

1560, 2150, 1061, 0222, 0123, 0283, 2154, 0004

Combined

(000**4**) (022**2**, 012**3**) (215**0**, 215**4**) (156**0**, 106**1**) (028**3**)

Grouped by third digit

0004, 0222, 0123, 2150, 2154, 1560, 1061, 0283

Combined

(000**4**, 106**1**) (012**3**, 215**0**, 215**4**) (022**2**, 028**3**) (156**0**)

Grouped by second digit

0004, 1061, 0123, 2150, 2154, 0222, 0283, 1560

Combined

(000**4**, 012**3**, 022**2**, 028**3**) (106**1**, 156**0**) (215**0**, 215**4**)

Grouped by first digit

0004, 0123, 0222, 0283, 1061, 1560, 2150, 2154

Combined (sorted)

6 Pseudocode and Analysis of Radix Sort

```
radixSort(int[] array, int n, int d) {  
    // Sorts n d-digit numeric strings in the array.  
    for (j = d down to 1) { // for digits in last position to 1st position  
        initialize 10 groups (queues) to empty // Q: why 10 groups?  
  
        for (i=0 through n-1) {  
            k = jth digit of array[i]  
            place array[i] at the end of group k  
        }  
        Replace array with all items in group 0, followed by all items  
        in group 1, and so on.  
    }  
}
```

Complexity is $O(d \times n)$ where d is the maximum number of digits of the n numeric strings in the array. Since d is fixed or bounded, so the complexity is $O(n)$.

7 Comparison of Sorting Algorithms

7 In-place Sort

- A sorting algorithm is said to be an **in-place** sort if it requires only a **constant amount, i.e. $O(1)$, of extra space** during the sorting process.
- Merge Sort is not in-place. (Why?) need an extra array
- How about Quick Sort? yes, no use of temp array

7 Stable Sort

- A sorting algorithm is **stable** if the **relative order of elements with the same key value is preserved** by the algorithm.
- Example 1 – An application of stable sort:
 - Assume that names have been sorted in alphabetical order.
 - Now, if this list is sorted again by tutorial group number, a **stable sort** algorithm would ensure that all students in the same tutorial groups still appear in alphabetical order of their names.
- Quick Sort and Selection Sort are not stable. (Why?)

7 Non-Stable Sort

- Example 2 – Quick Sort and Selection Sort are not stable:

Quick sort:

1285 150 5* 4746 602 5⁺ 8356 // pivot in bold

1285 (150 5* 602 5⁺) (4746 8356)

5⁺ 150 5* 602 **1285** 4746 8356 //pivot swapped with the last one in S1
// the **2 5's** are in different order of the initial list

Selection sort: select the largest element and swap with the last one

4746 1285 5* 602 **8356** 5⁺ 277

4746 1285 5* 602 277 5⁺ (8356)

5⁺ 1285 5* 602 277 (4746 8356)

// the **2 5's** are in different order of the initial list

7 Summary of Sorting Algorithms

	Worst Case	Best Case	In-place?	Stable?
Selection Sort	$O(n^2)$	$O(n^2)$	Yes	No
Insertion Sort	$O(n^2)$	$O(n)$	Yes	Yes
Bubble Sort	$O(n^2)$	$O(n^2)$	Yes	Yes
Bubble Sort 2 (improved with flag)	$O(n^2)$	$O(n)$	Yes	Yes
Merge Sort	$O(n \log n)$	$O(n \log n)$	No	Yes
Radix Sort (non-comparison based)	$O(n)$ (see notes 1)	$O(n)$	No	Yes
Quick Sort	$O(n^2)$	$O(n \log n)$	Yes	No

Notes: 1. $O(n)$ for Radix Sort is due to non-comparison based sorting.
2. $O(n \log n)$ is the best possible for comparison based sorting.

8 Use of Java Sort Methods

8 Java Sort Methods (in Arrays class)

```
static void sort(byte[] a)
static void sort(byte[] a, int fromIndex, int toIndex)
static void sort(char[] a)
static void sort(char[] a, int fromIndex, int toIndex)
static void sort(double[] a)
static void sort(double[] a, int fromIndex, int toIndex)
static void sort(float[] a)
static void sort(float[] a, int fromIndex, int toIndex)
static void sort(int[] a)
static void sort(int[] a, int fromIndex, int toIndex)
static void sort(long[] a)
static void sort(long[] a, int fromIndex, int toIndex)
static void sort(Object[] a)
static void sort(Object[] a, int fromIndex, int toIndex)
static void sort(short[] a)
static void sort(short[] a, int fromIndex, int toIndex)
static <T> void sort(T[] a, Comparator<? super T> c)
static <T> void sort(T[] a, int fromIndex, int toIndex,
                    Comparator<? super T> c)
```

8 To use `sort()` in Arrays

- The entities to be sorted must be stored in an **array** first.
- If they are stored in a **list**, then we have to use **`Collections.sort()`**
- If the data to be sorted are not primitive, then **`Comparator`** must be defined and used

Note: **`Collections`** is a Java public class and **`Comparator`** is a public interface. Comparators can be passed to a sort method (such as `Collections.sort()`) to allow precise control over the sort order.

8 Simple program using Collections.sort()

```
import java.util.*;
public class Sort {
    public static void main(String args[]) {
        List<String> list = Arrays.asList(args);
        Collections.sort(list);
        System.out.println(list);
    }
}
```

Sort.java

- Run the program:
`java Sort We walk the line`
- The following output is produced:
`[We, line, the, walk]`

Note: `Arrays` is a Java public class and `asList()` is a method of `Arrays` which returns a fixed-size list backed by the specified array.

8 Another solution using **Arrays.sort()**

```
import java.util.*;
public class Sort2 {
    public static void main(String args[]) {
        Arrays.sort(args);
        System.out.println(Arrays.toString(args));
    }
}
```

Sort2.java

- Run the program:
`java Sort2 We walk the line`
- The following output is produced:
`[We, line, the, walk]`

8 Example: class Person

```
class Person {  
    public String name;  
    public int age;  
  
    public Person(String name, int age) {  
        this.name = name;  
        this.age = age;  
    }  
    public String getName() { return name; }  
    public int getAge() { return age; }  
    public String toString() {  
        return name + " - " + age;  
    }  
}
```

Person.java

8 Comparator interface

java.util

Interface Comparator<T>

Type Parameters:

T - the type of objects that may be compared by this comparator

All Known Implementing Classes:

Collator, RuleBasedCollator

Functional Interface:

This is a functional interface and can therefore be used as the assignment target for a lambda expression or method reference.

Method Summary

All Methods	Static Methods	Instance Methods	Abstract Methods	Default Methods
Modifier and Type		Method and Description		
int		compare (T o1, T o2) Compares its two arguments for order.		
boolean		equals (Object obj) Indicates whether some other object is "equal to" this comparator.		

8 Comparator: AgeComparator

```
import java.util.Comparator;
class AgeComparator implements Comparator<Person> {
    public int compare(Person p1, Person p2) {
        // Returns the difference:
        // if positive, age of p1 is greater than p2
        // if zero, the ages are equal
        // if negative, age of p1 is less than p2
        return p1.getAge() - p2.getAge();
    }

    public boolean equals(Object obj) {
        // Simply checks to see if we have the same comparator object
        return this == obj;
    }
} // end AgeComparator
```

AgeComparator.java

Note: `compare()` and `equals()` are two methods of the interface `Comparator`. Need to implement them.

8 Comparator: NameComparator

```
import java.util.Comparator;
class NameComparator implements Comparator<Person> {

    public int compare(Person p1, Person p2) {
        // Compares its two arguments for order by name
        return p1.getName().compareTo(p2.getName());
    }

    public boolean equals(Object obj) {
        // simply checks to see if we have the same object
        return this == obj;
    }
} // end NameComparator
```

NameComparator.java

8 TestComparator (1/3)

```
import java.util.*;

public class TestComparator {

    public static void main(String args[]) {
        NameComparator nameComp = new NameComparator();
        AgeComparator ageComp = new AgeComparator();
        Person[] p = new Person[5];

        p[0] = new Person("Michael", 15);
        p[1] = new Person("Mimi", 9);
        p[2] = new Person("Sarah", 12);
        p[3] = new Person("Andrew", 15);
        p[4] = new Person("Mark", 12);
        List<Person> list = Arrays.asList(p);
    }
}
```

TestComparator.java

8 TestComparator (2/3)

```
System.out.println("Sorting by age:");
Collections.sort(list, ageComp);
System.out.println(list + "\n");

List<Person> list2 = Arrays.asList(p);
System.out.println("Sorting by name:");
Collections.sort(list2, nameComp);
System.out.println(list2 + "\n");

System.out.println("Now sort by name, then sort by age:");
Collections.sort(list2, ageComp); // list2 is already
                                   // sorted by name

System.out.println(list2);
} // end main

} // end TestComparator
```

TestComparator.java

8 TestComparator (3/3)

```
java TestComparator
```

Sorting by age:

[Mimi - 9, Sarah - 12, Mark - 12, Michael - 15, Andrew - 15]

Sorting by name:

[Andrew - 15, Mark - 12, Michael - 15, Mimi - 9, Sarah - 12]

Now sort by name, then sort by age:

[Mimi - 9, Mark - 12, Sarah - 12, Andrew - 15, Michael - 15]

8 Another solution using **Arrays.sort()**

We can replace the statements

```
List<Person> list = Arrays.asList(p);  
System.out.println("Sorting by age:");  
Collections.sort(list, ageComp);  
System.out.println(list + "\n");
```

with

```
System.out.println("Sorting by age using Arrays.sort():");  
Arrays.sort(p, ageComp);  
System.out.println(Arrays.toString(p) + "\n");
```

Summary

- We have introduced and analysed some classic sorting algorithms.
- Merge Sort and Quick Sort are in general faster than Selection Sort, Bubble Sort and Insertion Sort.
- The sorting algorithms discussed here are comparison based sorts, except for Radix Sort which is non-comparison based.
- $O(n \log n)$ is the best possible worst-case running time for comparison based sorting algorithms.
- There exist Java methods to perform sorting.

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