

CS2100

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COMPUTER ORGANISATION

Lecture #8

MIPS

Part II: More Instructions



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Lecture #8: MIPS Part 2: More Instructions

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1. Memory Organisation (General)

- The main memory can be viewed as a large, **single-dimension array** of memory locations.
- Each location of the memory has an **address**, which is an index into the array.
 - Given a k -bit address, the address space is of size 2^k .
- The memory map on the right contains one byte (8 bits) in every location/address.
 - This is called byte addressing

Address	Content
0	8 bits
1	8 bits
2	8 bits
3	8 bits
4	8 bits
5	8 bits
6	8 bits
7	8 bits
8	8 bits
9	8 bits
10	8 bits
11	8 bits
	:

1.1 Memory: Transfer Unit

- Using distinct memory address, we can access:
 - a single **byte** (**byte addressable**) or
 - a single **word** (**word addressable**)
- **Word** is:
 - Usually 2^n bytes
 - The common unit of transfer between processor and memory
 - Also commonly coincide with the register size, the integer size and instruction size in most architectures

1.2 Memory: Word Alignment

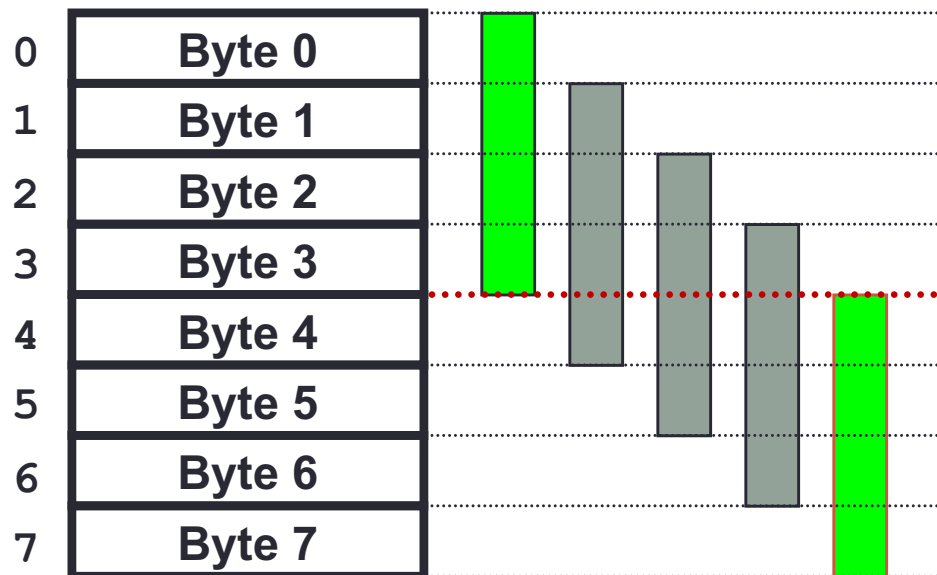
■ Word alignment:

- Words are aligned in memory if they begin at a byte address that is a multiple of the number of bytes in a word.

Example: If a word consists of 4 bytes, then:

Check address is
divisible by word
size

Address



Aligned word
Mis-aligned word

Question:


How do we quickly check whether a given memory address is word-aligned or not?

Processor can only access address 0 and 4 if it is a word addressable memory with 4 word size. This problem is called mis-aligned word.

2. MIPS Memory Instructions

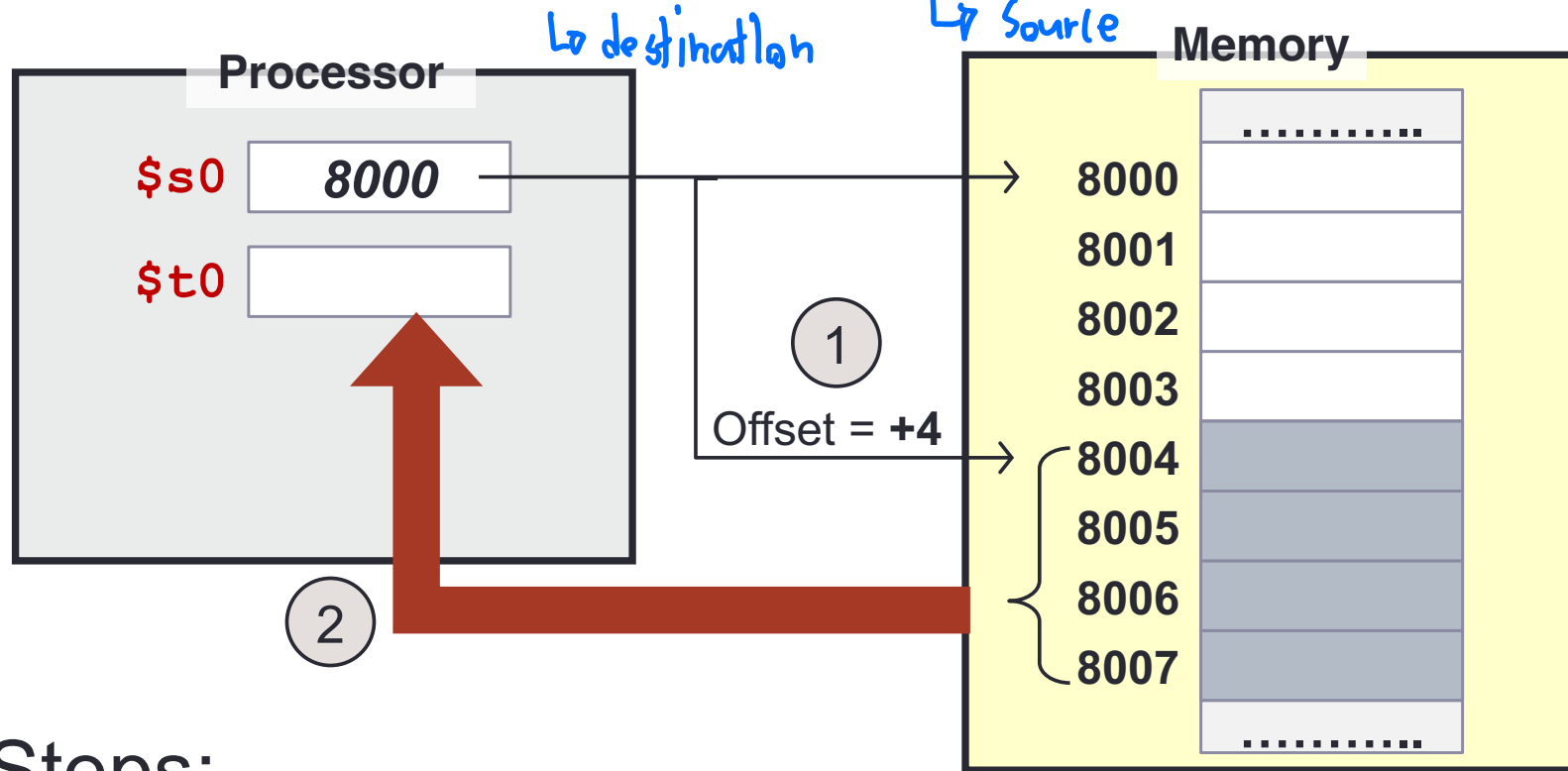
- MIPS is a load-store register architecture
 - 32 registers, each 32-bit (4-byte) long
 - Each word contains 32 bits (4 bytes)
 - Memory addresses are 32-bit long

NOTE:
Magic number
is 32

Name	Examples	Comments
32 registers	\$s0-\$s7, \$t0-\$t9, \$zero, \$a0-\$a3, \$v0-\$v1, \$gp, \$fp, \$sp, \$ra, \$at	Fast processor storage for data. In MIPS, data must be in registers to perform arithmetic.
2^{30} memory words 	Mem [0] , Mem [4] , ... , Mem [4294967292]	Accessed only by data transfer instructions. MIPS uses byte addresses, so consecutive words differ by 4. Memory holds data structures, such as arrays, and spilled registers, such as those saved on procedure calls.

2.1 Memory Instruction: Load Word

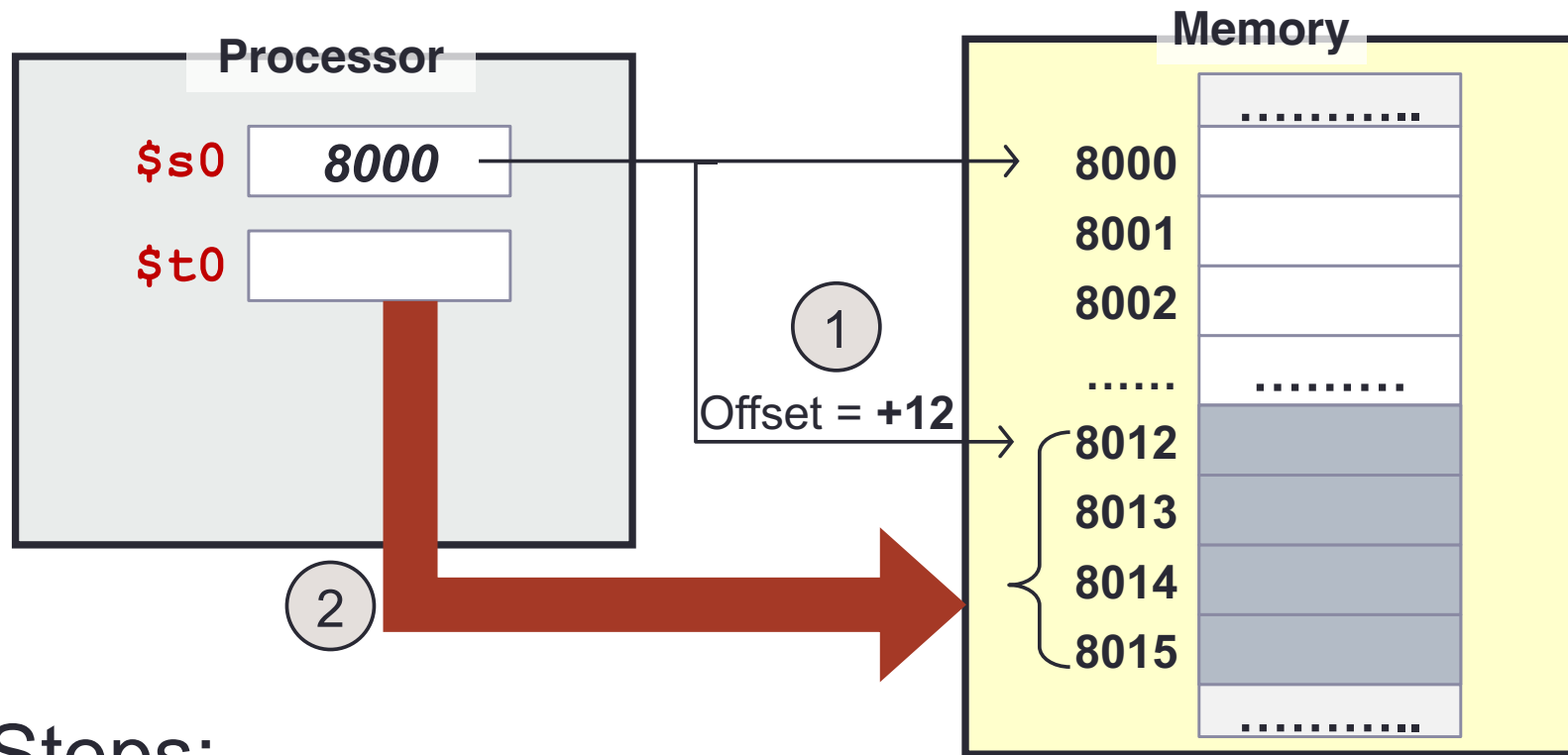
- Example: **lw** $\$t0$, 4 ($\$s0$)
↳ Load word *↳ displacement*
↳ destination *↳ source*



- Steps:
 1. Memory Address = $\$s0 + 4 = 8000 + 4 = 8004$
 2. Memory word at **Mem[8004]** is loaded into $\$t0$

2.2 Memory Instruction: Store Word

- Example: **sw** **\$t0**, 12 (**\$s0**)



- Steps:

1. Memory Address = **\$s0** + 12 = 8000 + 12 = **8012**
2. Content of **\$t0** is stored into word at **Mem[8012]**

2.3 Load and Store Instructions

- Only **load** and **store** instructions can access data in memory.
- Example: Each array element occupies a word.

C Code	MIPS Code
<pre>A[7] = h + A[10]; t0 = A[10]; t0 = h + t0; A[7] = t0;</pre>	<pre>lw \$t0, 40(\$s3) add \$t0, \$s2, \$t0 sw \$t0, 28(\$s3)</pre>

- Each array element occupies a word (4 bytes).
- \$s3** contains the **base address** (address of first element, A[0]) of array A. Variable **h** is mapped to **\$s2**.
- Remember arithmetic operands (for **add**) are registers, not memory!

2.4 Memory Instructions: Others (1/2)

- Other than load word (**lw**) and store word (**sw**), there are other variants, example:
 - load byte (**lb**)
 - store byte (**sb**)
- Similar in format:

lb **\$t1**, 12 (**\$s3**)
sb **\$t2**, 13 (**\$s3**)
- Similar in working except that one byte, instead of one word, is loaded or stored
 - Note that the offset no longer needs to be a multiple of 4

2.4 Memory Instructions: Others (2/2)

- MIPS disallows loading/storing unaligned word using **lw/sw**:
 - Pseudo-Instructions ***unaligned load word*** (**ulw**) and ***unaligned store word*** (**usw**) are provided for this purpose

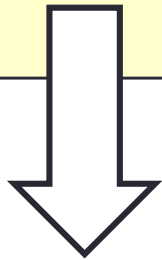
NOTE:

ulw/usw can be translated to sequence of lb/sb + other operations

- Other memory instructions:
 - **lh** and **sh**: load halfword and store halfword
 - **lwl**, **lwr**, **swl**, **swr**: load word left / right, store word left / right.
 - etc...

2.5 Example: Array (assume 4 bytes per element)

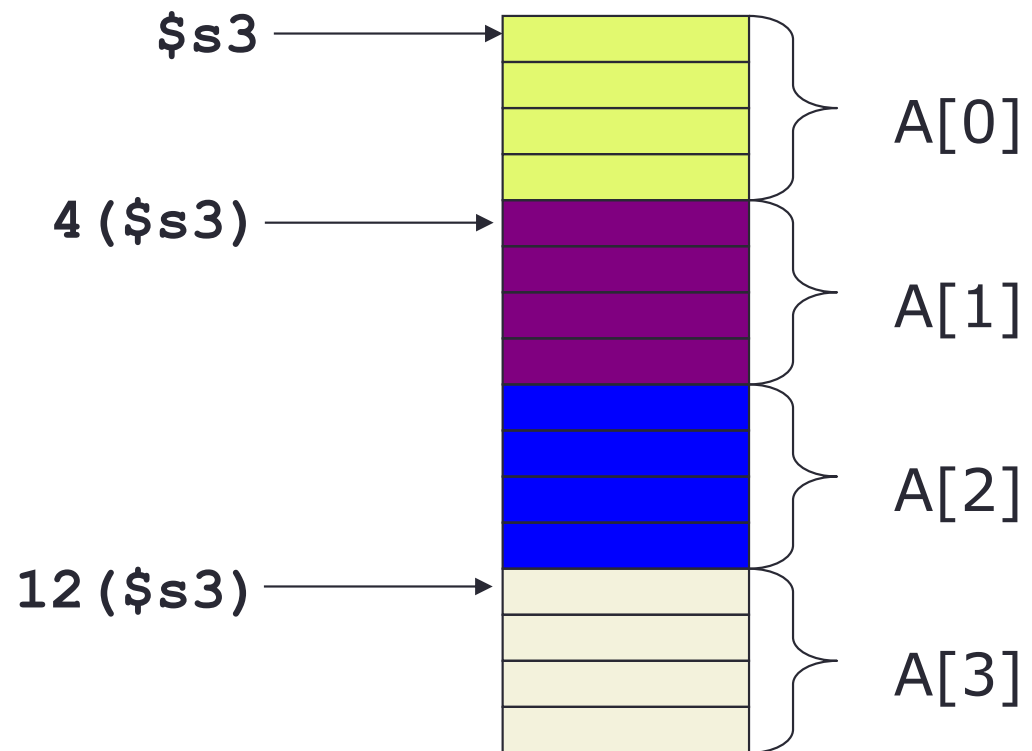
C Statement to translate	Variables Mapping
<code>A[3] = h + A[1];</code>	<code>h</code> \rightarrow <code>\$s2</code> base of <code>A[]</code> \rightarrow <code>\$s3</code>



```
lw  $t0, 4($s3)
```

```
add $t0, $s2, $t0
```

```
sw  $t0, 12($s3)
```



2.6 Common Questions: Address vs Value

Key concept:

Registers do NOT have types

- A register can hold any 32-bit number:
 - The number has no implicit data type and is interpreted according to the instruction that uses it
- Examples:
 - `add $t2, $t1, $t0`
 - `$t0` and `$t1` should contain data values
 - `lw $t2, 0($t0)`
 - `$t0` should contain a memory address

2.6 Common Questions: **Byte** vs **Word**

Important:

Consecutive **word addresses** in machines with **byte-addressing** do not differ by 1

- Common error:
 - Assume that the address of the next word can be found by incrementing the address in a register by 1 instead of by the word size in bytes
- For both **lw** and **sw**:
 - The sum of base address and offset must be a multiple of 4 (i.e. to adhere to word boundary)

2.7 Example: Swapping Elements

C Statement to translate	Variables Mapping
<pre> swap(int v[], int k) { int temp; temp = v[k] v[k] = v[k+1]; v[k+1] = temp; } </pre>	<p> $k \rightarrow \\$5$ Base address of $v[] \rightarrow \\$4$ $temp \rightarrow \\$15$ </p> <p> Example: $k = 3$; to swap $v[3]$ with $v[4]$. Assume base address of v is 2000. </p> <p> $\\$5 (k) \leftarrow 3$ $\\$4 (\text{base addr. of } v) \leftarrow 2000$ </p>
<pre> swap: sll \$2, \$5, 2 add \$2, \$4, \$2 lw \$15, 0(\$2) lw \$16, 4(\$2) sw \$16, 0(\$2) sw \$15, 4(\$2) </pre>	<p> $\\$2 \leftarrow 12$ $\\$2 \leftarrow 2012$ </p> <p> $\\$15 \leftarrow \text{content of mem. addr. 2012 (v[3])}$ $\\$16 \leftarrow \text{content of mem. addr. 2016 (v[4])}$ </p> <p> $\text{content of mem. addr. 2012 (v[3])} \leftarrow \\16 $\text{content of mem. addr. 2016 (v[4])} \leftarrow \\15 </p>

Note: This is simplified and may not be a direct translation of the C code.

Reading

- **Instructions: Language of the Computer**
 - Read up COD Chapter 2, pages 52-57. (3rd edition)
 - Read up COD Section 2.3 (4th edition)



3. Making Decisions (1/2)

- We cover only sequential execution so far:
 - Instruction is executed in program order
- To perform general computing tasks, we need to:
 - **Make decisions**
 - **Perform iterations** (in later section)
- Decisions making in high-level language:
 - **if** and **goto** statements
 - MIPS decision making instructions are similar to **if** statement with a **goto**
 - **goto** is discouraged in high-level languages but necessary in assembly 😊

3. Making Decisions (2/2)

- Decision-making instructions
 - Alter the control flow of the program
 - Change the next instruction to be executed
- Two types of decision-making statements in MIPS
 - **Conditional** (branch)
 - `bne $t0, $t1, label`
 - `beq $t0, $t1, label`
 - **Unconditional** (jump)
 - `j label`
- A label is an “anchor” in the assembly code to indicate point of interest, usually as branch target
 - Labels are NOT instructions!

NOTE:

Selection: branch/jump down

Repetition: branch/jump up

In both cases, it's a “goto” operation.

3.1 Conditional Branch: **beq** and **bne**

- Processor follows the branch only when the condition is satisfied (true)
- **beq** **\$r1**, **\$r2**, **L1**
 - Go to statement labeled **L1** if the value in register **\$r1** equals the value in register **\$r2**
 - **beq** is “bbranch if equal”
 - C code: **if** (**a** == **b**) **goto** **L1**
- **bne** **\$r1**, **\$r2**, **L1**
 - Go to statement labeled **L1** if the value in register **\$r1** does not equal the value in register **\$r2**
 - **bne** is “bbranch if not equal”
 - C code: **if** (**a** != **b**) **goto** **L1**

3.2 Unconditional Jump: j

- Processor **always** follows the branch
- **j** L1
 - Jump to label L1 unconditionally
 - C code: **goto** L1
- Technically equivalent to such statement

beq \$s0, \$s0, L1

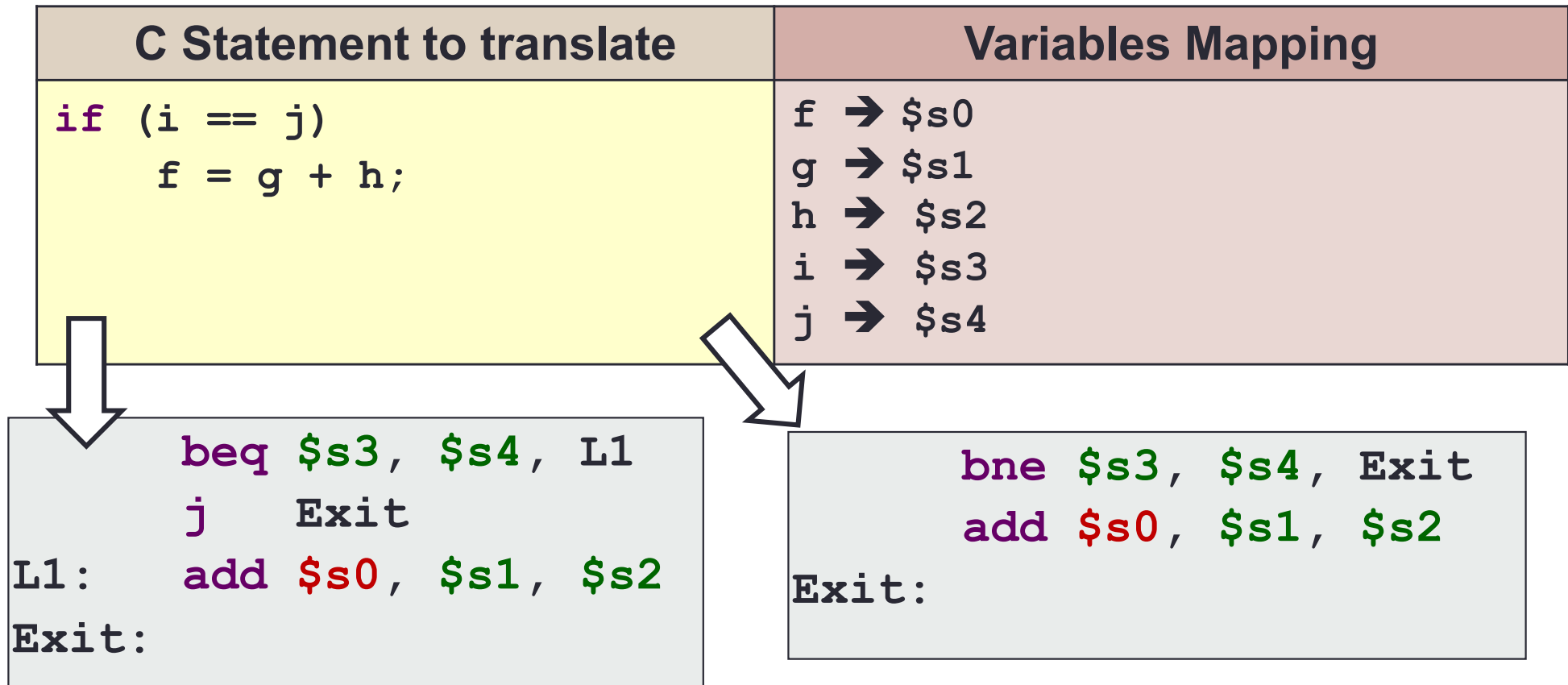
NOTE:

The general technique of compiling from C to MIPS with selection/repetition is to first remove the construct and replace it with either one of these:

```
if (x == y) goto L; // this is beq
if (x != y) goto L; // this is bne
```

Some tricks can be used (e.g., inversion) to generate fewer lines of codes.

3.3 IF statement (1/2)



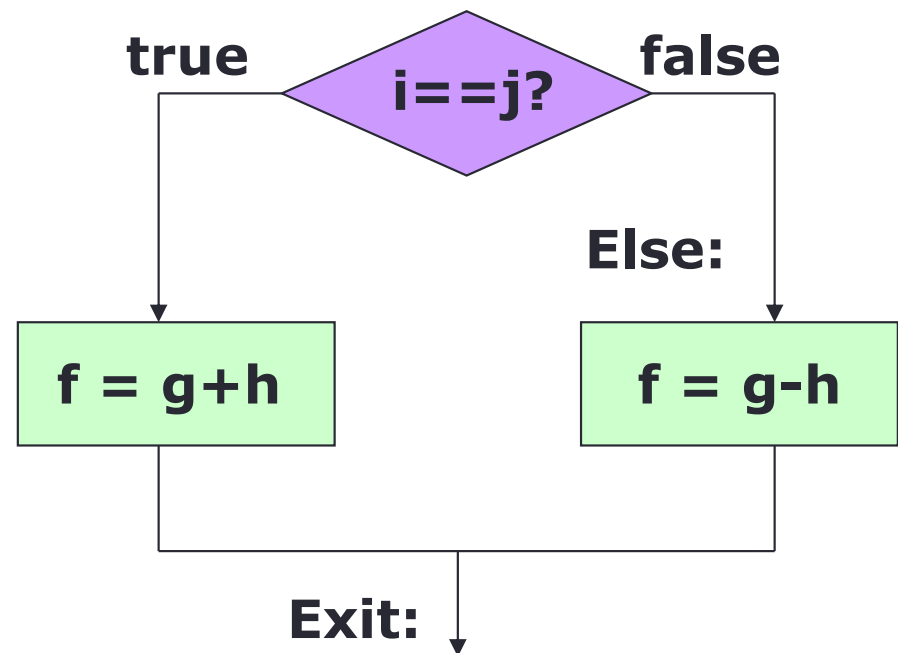
- Two equivalent translations:
 - The one on the right is more efficient
- Common technique: Invert the condition for shorter code

3.3 IF statement (2/2)

C Statement to translate	Variables Mapping
<pre>if (i == j) f = g + h; else f = g - h;</pre>	<pre>f → \$s0 g → \$s1 h → \$s2 i → \$s3 j → \$s4</pre>



```
bne $s3, $s4, Else
add $s0, $s1, $s2
j    Exit
Else: sub $s0, $s1, $s2
Exit:
```



- Question: Rewrite with **bneq**?

3.4 Exercise #1: IF statement

MIPS code to translate into C	Variables Mapping
<pre> beq \$s1, \$s2, Exit add \$s0, \$zero, \$zero Exit:</pre>	<pre>f → \$s0 i → \$s1 j → \$s2</pre>

- What is the corresponding high-level statement?

```
if (i != j) {
    f = 0;
}
```

NOTE:

Remember the inversion? Also works in “reverse” (or decompilation) process. beq becomes if (i != j).

4. Loops (1/2)

- C while-loop:

```
while (j == k)
    i = i + 1;
```



- Rewritten with goto

```
Loop:  if (j != k)
        goto Exit;
        i = i + 1;
        goto Loop;

Exit:
```

Key concept:

Any form of loop can be written in assembly with the help of conditional branches and jumps.

4. Loops (2/2)

C Statement to translate	Variables Mapping	
<pre>Loop: if (j != k) goto Exit; i = i+1; goto Loop; Exit:</pre>	<pre>i → \$s3 j → \$s4 k → \$s5</pre>	NOTE: This shows the process clearly: 1. Convert from while to if(...) goto 2. Convert from there to MIPS

- What is the corresponding MIPS code?

```
Loop: bne  $s4, $s5, Exit  # if (j!= k) Exit
      addi $s3, $s3, 1
      j    Loop           # repeat loop
Exit:
```

4.1 Exercise #2: FOR loop

- Write the following loop statement in MIPS

C Statement to translate	Variables Mapping
<pre>for (i=0; i<10; i++) a = a + 5;</pre>	<pre>i → \$s0 a → \$s2</pre>

```
        add    $s0, $zero, $zero  
        addi   $s1, $zero, 10  
Loop:   beq    $s0, $s1, Exit  
        addi   $s2, $s2, 5  
        addi   $s0, $s0, 1  
        j      Loop  
Exit:
```

NOTE:

Alternatively, if you know how to compile while-loop, then you can translate the for-loop into:

```
i = 0;  
while (i < 10) {  
    a = a + 5;  
    i++;  
}
```

4.2 Inequalities (1/2)

- We have **beq** and **bne**, what about branch-if-less-than?
 - There is no real **blt** instruction in MIPS
- Use **slt** (set on less than) or **slti**.

```
slt $t0, $s1, $s2
```

=

```
if ($s1 < $s2)  
    $t0 = 1;  
else  
    $t0 = 0;
```

NOTE:

If we use the “ternary operator” discussed in tutorial, then:

```
$t0 = ($s1 < $s2) ? 1 : 0;
```

4.2 Inequalities (1/2)

- To build a “**blt** \$s1, \$s2, L” instruction:

```
slt $t0, $s1, $s2  
bne $t0, $zero, L
```

==

```
if ($s1 < $s2)  
    goto L;
```

- This is another example of **pseudo-instruction**:
 - Assembler translates (**blt**) instruction in an assembly program into the equivalent MIPS (two) instructions

Reading

- **Instructions: Language of the Computer**
 - Section 2.6 Instructions for Making Decisions. (3rd edition)
 - Section 2.7 Instructions for Making Decisions. (4th edition)

NOTE:

If we know how to compile:

`if (a < b) ...`

Then we can compile:

`if (a < b) → if (b < a)`

`if (a >= b) → if (!(a < b))`

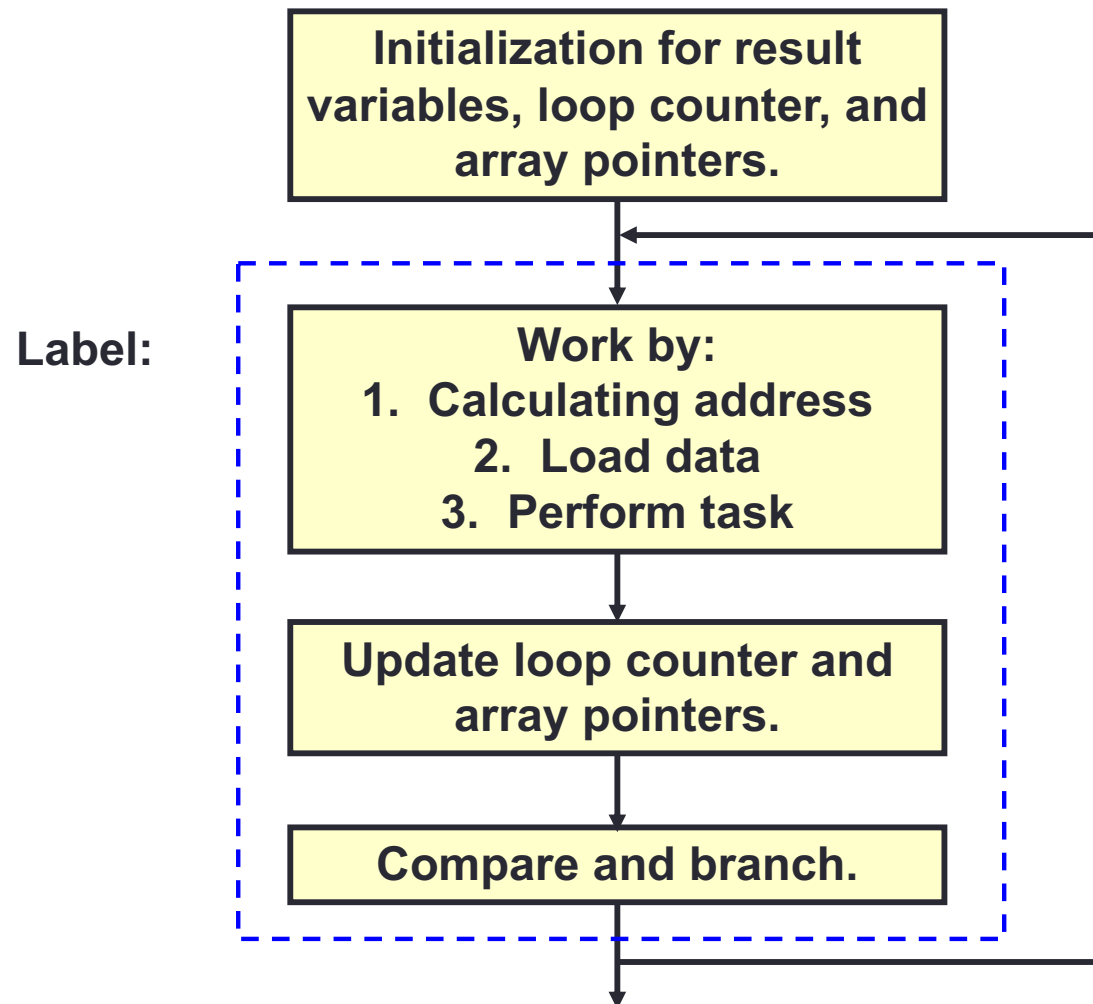
`if (a <= b) → if (b >= a) → if(!(b < a))`

So, with an addition of `slt`, we can do any other kind of inequalities.



5. Array and Loop

- Typical example of accessing array elements in a loop:



5. Array and Loop: Question

Count the number of zeros in an Array **A**

- **A** is word array with 40 elements
- Address of A[] → **\$t0**, Result → **\$t8**

Simple C Code

```
result = 0;
i = 0;
while ( i < 40 ) {
    if ( A[i] == 0 )
        result++;
    i++;
}
```

- Think about:
 - How to perform the right comparison
 - How to translate A[i] correctly

5. Array and Loop: Version 1.0

Address of A[] → \$t0 Result → \$t8 i → \$t1	Comments
<pre> addi \$t8, \$zero, 0 addi \$t1, \$zero, 0 addi \$t2, \$zero, 40 loop: bge \$t1, \$t2, end sll \$t3, \$t1, 2 add \$t4, \$t0, \$t3 lw \$t5, 0(\$t4) bne \$t5, \$zero, skip addi \$t8, \$t8, 1 skip: addi \$t1, \$t1, 1 j loop end: </pre>	<pre> # end point # i * 4 # \$t4 ← &A[i] # \$t5 ← A[i] # result++ # i++ </pre>

5. Array and Loop: Version 2.0

Address of A[] → \$t0 Result → \$t8 &A[i] → \$t1	Comments
<pre> addi \$t8, \$zero, 0 addi \$t1, \$t0, 0 addi \$t2, \$t0, 160 loop: bge \$t1, \$t2, end lw \$t3, 0(\$t1) bne \$t3, \$zero, skip addi \$t8, \$t8, 1 skip: addi \$t1, \$t1, 4 j loop end:</pre>	<p>NOTE: Consider the code using pointer arith:</p> <pre>result = 0; ptr = A; end = &A[40]; while (ptr < end) { if (*ptr == 0) result++; ptr++; }</pre> <p>The resulting MIPS looks like the code on the left.</p>

- Use of “pointers” can produce more efficient code!

6.1 Exercise #3: Simple Loop

- Given the following MIPS code:

```
        addi $t1, $zero, 10
        add  $t1, $t1, $t1
        addi $t2, $zero, 10
Loop:   addi $t2, $t2, 10
        addi $t1, $t1, -1
        beq  $t1, $zero, Loop
```

NOTE:

You may want to “decompile” into C so that you can reason in C rather than in MIPS.

Alternatively, you may just “trace” but it may take longer time.

So try to get familiar with a few “compilation” pattern for a few C construct (e.g., *while*, *do-while*, *for*, *if*, *if-else*, *switch-case*, etc).

- i. How many instructions are executed? **Answer: (a)**

(a) 6 (b) 30 (c) 33 (d) 36 (e) None of the above

- ii. What is the final value in **\$t2**? **Answer: (b)**

(a) 10 (b) 20 (c) 300 (d) 310 (e) None of the above

6.2 Exercise #4: Simple Loop II

- Given the following MIPS code:

```
    add    $t0, $zero, $zero
    add    $t1, $t0, $t0
    addi   $t2, $t1, 4
Again: add    $t1, $t1, $t0
        addi   $t0, $t0, 1
        bne    $t2, $t0, Again
```

- How many instructions are executed? **Answer: (c)**
(a) 6 (b) 12 (c) 15 (d) 18 (e) None of the above
- What is the final value in **\$t1**? **Answer: (c)**
(a) 0 (b) 4 (c) 6 (d) 10 (e) None of the above

6.3 Exercise #5: Simple Loop III (1/2)

- Given the following MIPS code accessing a word array of elements in memory with the starting address in `$t0`.

```
        addi $t1, $t0, 10
        add  $t2, $zero, $zero
Loop:   ulw  $t3, 0($t1) # ulw: unaligned lw
        add  $t2, $t2, $t3
        addi $t1, $t1, -1
        bne  $t1, $t0, Loop
```

- How many times is the **bne** instruction executed?
(a) 1 (b) 3 (c) 9 (d) 10 (e) 11 **Answer: (d)**
- How many times does the **bne** instruction actually branch to the label **Loop**?
(a) 1 (b) 8 (c) 9 (d) 10 (e) 11 **Answer: (c)**

6.3 Exercise #5: Simple Loop III (2/2)

- Given the following MIPS code accessing a word array of elements in memory with the starting address in `$t0`.

```
        addi $t1, $t0, 10
        add  $t2, $zero, $zero
Loop:   ulw  $t3, 0($t1) # ulw: unaligned lw
        add  $t2, $t2, $t3
        addi $t1, $t1, -1
        bne  $t1, $t0, Loop
```

- iii. How many instructions are executed?
(a) 6 (b) 12 (c) 41 (d) 42 (e) 46 **Answer: (d)**
- iv. How many **unique** bytes of data are read from the memory?
(a) 4 (b) 10 (c) 11 (d) 13 (e) 40 **Answer: (d)**

Summary: Focus of CS2100

- Basic MIPS programming
 - Arithmetic: among registers only
 - Handling of large constants
 - Memory accesses: load/store
 - Control flow: branch and jump
 - Accessing array elements
 - System calls (covered in labs)
- Things we are not going to cover
 - Support for procedures
 - Linkers, loaders, memory layout
 - Stacks, frames, recursion
 - Interrupts and exceptions

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