CS2100 Computer Organization

2022/23 Semester 2

Midterm Assessment

SOLUTIONS

1. The following C program fragment inverts the bits in a uint32_t integer. The uint32_t datatype is similar to the int32_t datatype in Assignment 1, except that it is 32-bits long and unsigned. The example below shows what this operation is supposed to do. For brevity we will use a 16-bit number instead of a 32-bit number:

Original: 0b1011 0010 0100 1100

Return from invert: 0b0100 1101 1011 0011

That is, all 0's are changed to 1's and all 1's are changed to 0's.

```
uint32_t invert(uint32_t x) {
    /* MISSING 1 */
}
```

Which ONE of the following C statements should be inserted into the space marked /* MISSING 1*/? (Note: You are not allowed to enter and run the code on your computer. Doing so constitutes cheating)

- a. return -x;
- b. return !x;
- c. return x^0b0;
- d. +return ~x;
- e. None of the above options are correct.

2. The following C program fragment reverses the bits in a uint32_t integer. The example below shows how this operation works. Again, for brevity we will show a 16-bit integer instead of 32-bits:

Original: 0b0010 1101 0110 1010 Return from reverse: 0b0101 0110 1011 0100

```
uint32_t reverse(uint32_t x) {
    uint32_t result = 0;

    for(int i=0; i<32; i++) {
        /* MISSING 2 */
    }

    return result;
}</pre>
```

Which ONE of the following C statements should be inserted into the space marked /* MISSING 2*/? (Note: You are not allowed to enter and run the code on your computer. Doing so constitutes cheating)

```
a. result |= (x & (0b1 << i)) << (31 - i);</li>
b. +result |= (x & (0b1 << i)) << (31 - 2*i);</li>
c. result |= (x & (0b1 >> i)) >> (31 - 2*i);
d. result &= (x & (0b1 << i)) << (31 - i);</li>
e. None of the above options a. to d. are correct.
```

3. We now call both the earlier invert and reverse functions in a function called "combine", to produce the behavior shown below (example this time is in 32-bits):

The combine function is shown here:

```
uint32_t combine(uint32_t x) {
    /* MISSING 3 */
}
```

- 4. Which ONE of the following C statements can be inserted into the space marked /* MISSING 3 */? (Note: You are not allowed to enter and run the code on your computer. Doing so constitutes cheating)
 - a. return !invert(reverse(x));
 - b. return !reverse(invert(x));
 - c. return ~invert(reverse(x));
 - d. return ~reverse(invert(x));
 - e. +None of the options a. to d. are correct.
- 5. Consider a 16-bit fixed-point number system, in sign and magnitude representation, with 8 bits in the integer portion (including sign bit) and 8 bits in the fraction portion. Fill in the following blanks:

Most positive number that can be represented: **127.99609**Most negative number that can be represented: **-127.99609**Smallest positive number that can be represented: **0.0039063**

Consider again the same 16-bit fixed point number system, in sign-and-magnitude representation with 8 bits in the integer portion (including sign bit) and 8 bits in the fraction portion. Find the ABSOLUTE ERROR for representing the numbers 3.625 and 9.151. When representing both numbers, consider possible rounding of the least significant bit.

```
3.625 = 00000011.10100000<sub>2</sub>

Absolute error = 0

9.151 = 00001001.001001101 = 00001001.00100111<sub>2</sub> = 9.1523438

Absolute error in representing 3.625: 0

Absolute error in representing 9.151: 0.001344
```

6. Given the same 16-bit fixed point number system in sign-and-magnitude representation with 8 bits for the integer portion (including sign bit) and 8 bits for the fraction portion, what is 0x95B0 in base-10?

```
1001 0101 1011 0000
```

- = 0x95B0
- = -21.6875
- 7. What is $1213_5 + 312_5$ in base 5?

```
1213_5 + 312_5 = 2030 in base 5.
```

8. What is $114_7 - 36_7$ in base-7? (Hint: Think about how you'd subtract two base-10 numbers, and apply the same idea here)

```
114_7 - 36_7 = 45 in base 7.
```

9. $312_x = 431_5$. What is the mystery base x?

x = 6

10. What is -19.1015625 in IEEE-754 format? Write your answer in hexadecimal.

```
19.1015625 = 10011.0001101 x 2°
```

1.00110001101 x 2⁴

Mantissa = 00110001101000000000000

Exponent = 4 + 127 = 131 = 10000011

Sign = 1

- = 0xC198D000

The MIPS code fragment below works on an integer array A whose base address is stored in \$s0 and size (number of elements) of the array in \$s1. The first instruction addi \$t1, \$s0, 0 is at address 0x00400034.

```
addi $t1, $s0, 0
                             # Inst1 @ 0x00400034
           $t2, 0($t1)
                             # Inst2
      lw
      addi $t0, $zero, 1
                            # Inst3
Loop: beq $t0, $s1, Exit
                            # Inst4
      addi $t1, $t1, 4
                             # Inst5
           $t3, 0($t1)
      lw
                             # Inst6
      add $t3, $t3, $t2
                             # Inst7
                            # Inst8
           $t3, 0($t1)
      sw
      addi $t2, $t3, 0
                             # Inst9
      addi $t0, $t0, 1
                             # Inst10
      Ė
           Loop
                             # Inst11
Exit:
```

11. Below is an equivalent code in C. Let the array name be A and the size (number of elements) of A be size. int i;

```
for (i = 1; i < size; i++) {
          A[i] = A[i-1] + A[i];
}</pre>
```

12.

a. Assuming that *size* is 8 and array A contains the following values: 3, 5, -2, 4, 0, 7, -1, 6 (A[0]=3 and A[7]=6). What is the final value of **\$t3**?

```
Answer: 22 (Array A after execution: 3, 8, 6, 10, 10, 17, 16, 22.)
```

b. For the instruction **beq \$t0, \$s1, Exit**, write out its hexadecimal representation.

```
Answer: 1 1 1 1 0 0 0 7
opcode = 4; $t0 = $8; $s1 = 17; Exit = 7; beq $t0, $s1, 7 \rightarrow 4, 8, 17, 7
4, 8, 17, 7 \rightarrow 0001/00, 01/000, 1/0001/, 0000/0, 000/00, 00/0111
```

 \rightarrow 11110007

c. For the instruction **sw \$t3, 0(\$t1)**, write out its hexadecimal representation.

Answer: A D 2 B 0 0 0 0

opcode = 2b; \$t3 = \$11; \$t1 = 9; sw \$t3, $0($t1) \rightarrow 2b$, 9, 11, 0

2b, 9, 11, 0 \rightarrow 1010/11, 01/001, 0/1011/, 0000/0, 000/00, 00/0000

\rightarrow AD2B0000

d. For the instruction **j Loop**, write out its hexadecimal representation.

Answer: 08100010

opcode = 2; address of Inst4: 0x00400040

 $0000/10 \ 0000 \ 00/00 \ 01/00 \ 00/00 \ 00/00 \ 01/00 \ 0000$

 \rightarrow 08100010

Expected wrong answer: address of Inst4: 0x00400037

0000/10 0000 00/00 01/00 00/00 00/00 00/11 0111 \rightarrow 0 8 1 0 0 0 0 D

Expected wrong answer:

 $0000/10\ 00/00\ 00/00\ 01/00\ 00/00\ 00/00\ 01/00\ 0000$ \rightarrow 0 8 0 1 0 0 0 1

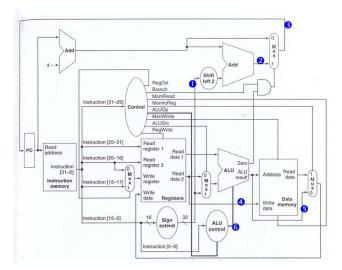
Expected wrong answer: use relative displace (as in branch) -8

-8 = 11 1111 1111 1111 1111 1111 1000 (-8 in 2's complement)

000010 11 1111 1111 1111 1111 1111 1000 → 0 B F F F F F 8

Consider the MIPS datapath with the values of the registers shown in the table below, where all values are in hexadecimal. Suppose the instruction **sw \$a3, -20(\$s2)** at address 100_{10} is currently being executed.

```
R0
    (r0) = 0x00000000 \mid R1
                              (at) = 0x00002000
R2
    (v0) = 0x00000321 \mid R3
                              (v1) = 0x0000000A
                              (a1) = 0x7FFFF000
    (a0) = 0x00000005 \mid R5
R4
R6
    (a2) = 0x7FFFF004 \mid R7
                              (a3) = 0x700003AC
R8
    (t0) = 0x0000001 \mid R9
                              (t1) = 0x00000c00
R10 (t2) = 0x000000000 | R11 (t3) = 0xfffffff0
R12 (t4) = 0xF0000000 | R13 (t5) = 0x000000fff
R14 (t6) = 0x00006200 \mid R15 (t7) = 0x00000e00
R16 (s0) = 0x00300000 | R17 (s1) = 0x000000000
R18 (s2) = 0xCC040200 | R19 (s3) = 0x000E5003
R20 (s4) = 0xCC030200 | R21 (s5) = 0x10000000
R22 (s6) = 0x00055000 | R23 (s7) = 0xF0000000
R24 (t8) = 0x00000005 | R25 (t9) = 0x0000D624
R26 (k0) = 0x00000000 | R27 (k1) = 0x00000000
R28 \text{ (gp)} = 0x10008000 \mid R29 \text{ (sp)} = 0x7FFFEff4
R30 (s8) = 0x1000000F \mid R31 (ra) = 0x00400018
```



13. What is the value at **①** (output from sign-extend)? Your answer must be in hexadecimal.

Answer: **0xFFFFFEC**

 $20_{10} = 0000\ 0000\ 0000\ 0000\ 0000\ 0001\ 0100_{2s}$

-20₁₀ = 1111 1111 1111 1111 1111 1111 1110 1100_{2s}

- A. 0x0000014
- B. 0x0000FFEC
- C. 0xFFFFFEB
- D. +0xFFFFFEC
- E. None of the above.
- 14. What is the value at **2**? Your answer must be in hexadecimal. (Taker

Answer: 0x0000018

Note: Answer without leading zeroes is acceptable.

Immediate value = -20_{10} = 1111 1111 1111 1111 1111 1110 1100_{2s} (from part a)

Immed << 2 = 1111 1111 1111 1111 1111 1011 0000

 $PC+4 = 100_{10} + 4 = 68_{16} = 0x0000 0000 0000 0000 0000 0110 1000$

Adding, we have 0x0000 0000 0000 0000 0000 0001 1000

- A. +0x0000018
- B. 0x0000050
- C. 0x00000054
- D. 0x80000054
- E. None of the above.

15. What is the value at **3**? Your answer must be in hexadecimal. (Taken)

Answer: 0x00000068

Note: Answer without leading zeroes is acceptable.

 $PC+4 = 100_{10} + 4 = 68_{16} = 0x0000 0000 0000 0000 0000 0110 1000$

- A. 0x00000064
- B. **+0x00000068**
- C. 0x00000100
- D. 0x00000104
- E. None of the above.
- 16. What is the value at **6** (ALU control output)? Your answer must be in binary. (Taken)

Answer: 0b0010

sw requires ALU to perform an add operation. The 4-bit ALUcontrol for add is 0010.

- A. 0b0000
- B. 0b0001
- C. **+0b0010**
- D. 0b0110
- E. None of the above.