

# DIGITAL IMAGE PROCESSING

## IMAGE COMPRESSION

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## IMAGE COMPRESSION

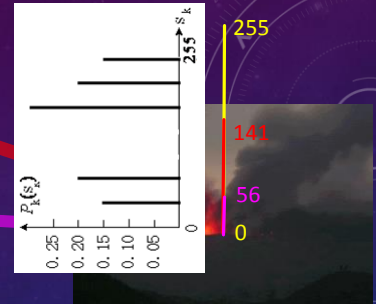
Calculate the amount of data required to represent a two-hour high definition(HD) television movie using  $1920 \times 1080 \times 3 \times 8$  bit pixel arrays, at rates of 30 fps(frames per second)

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$$1920 \times 1080 \times 8 \times 3 \times 30 = 1.493 \text{ Gbit/s} = 186.624 \text{ MB/s}$$

$$186.624 \text{ MB/s} \times 3600 \text{ s} / h \times 2 \text{ hrs} = 1343.7 \text{ GB}$$

IMAGE



- fundamentals

- Entropy of a single event

$$I(E) = \log \frac{1}{P(E)} = -\log P(E)$$

- Entropy of a zero-memory source

$$H = -\sum_i P(a_i) \log P(a_i)$$

- Entropy of a zero-memory intensity source(image)

- Shannon's first theorem

- Entropy of Markov(finite memory) source

## IMAGE COMPRESSION

$$-\log_2 P(E) \sim (\text{bit})$$

- fundamentals

- Entropy of a single event

$$I(E) = \log \frac{1}{P(E)} = -\log P(E)$$

source symbol

- Entropy of a zero-memory source

$$H = -\sum_i P(a_i) \log P(a_i)$$

- Entropy of a zero-memory intensity source(image)

$$\tilde{H} = -\sum_{k=0}^{L-1} P_r(r_k) \log_2 P_r(r_k)$$

- Shannon's first theorem

$$H = \lim_{n \rightarrow \infty} \frac{L_{\text{avg},n}}{n}$$

- Entropy of Markov(finite memory) source

$$H_N = -\sum_{a_0} \sum_{a_1} \cdots \sum_{a_N} P(a_0, a_1, \dots, a_N) \log_2 P(a_0 | a_1, \dots, a_N)$$

$$H_0 \geq H_1 \geq \cdots \geq H_{N-1} \geq H_N = H_{N+1} = H_{N+2} = \cdots = H_\infty$$

# IMAGE COMPRESSION

- fundamentals

- Entropy
- average code length
- Coding efficiency
- Compression ratio
- SNR

$$\frac{H(x)}{R(x)}$$

$$\tilde{H} = - \sum_{k=0}^{L-1} P_r(r_k) \log_2 P_r(r_k)$$

$$R(x) = \sum_i P(x_i) L_i$$

$$c = \frac{b_1}{b_2} = \frac{R_{before}(x)}{R_{after}(x)} = \frac{ceil(\log_2(L))}{R_{after}(x)}$$

$$10 \log \frac{\sigma_x^2}{\sigma_e^2}$$

# IMAGE COMPRESSION

- Image data redundancy

- Coding redundancy
- Spatial and temporal redundancy
- Irrelevant information

$$R(x) = \sum_i P(x_i) L_i$$

IMAGE COMPRESSION

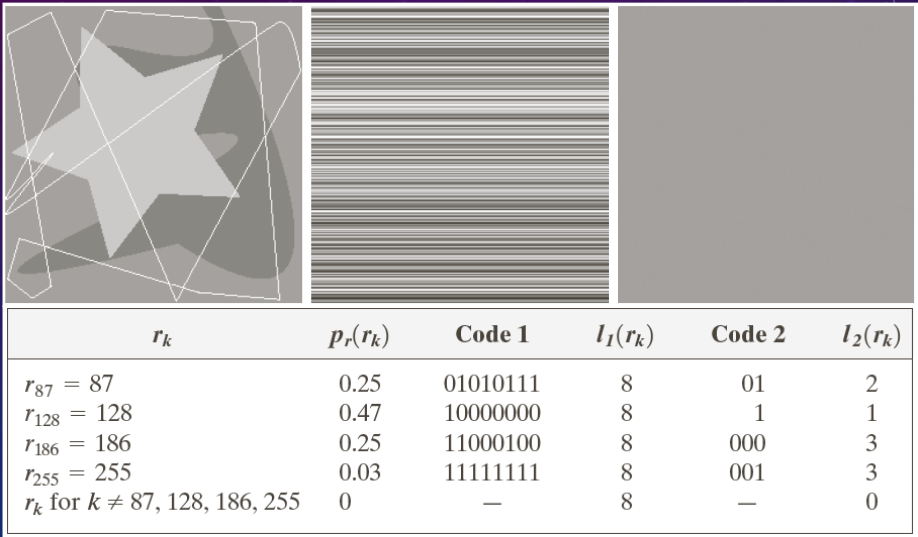
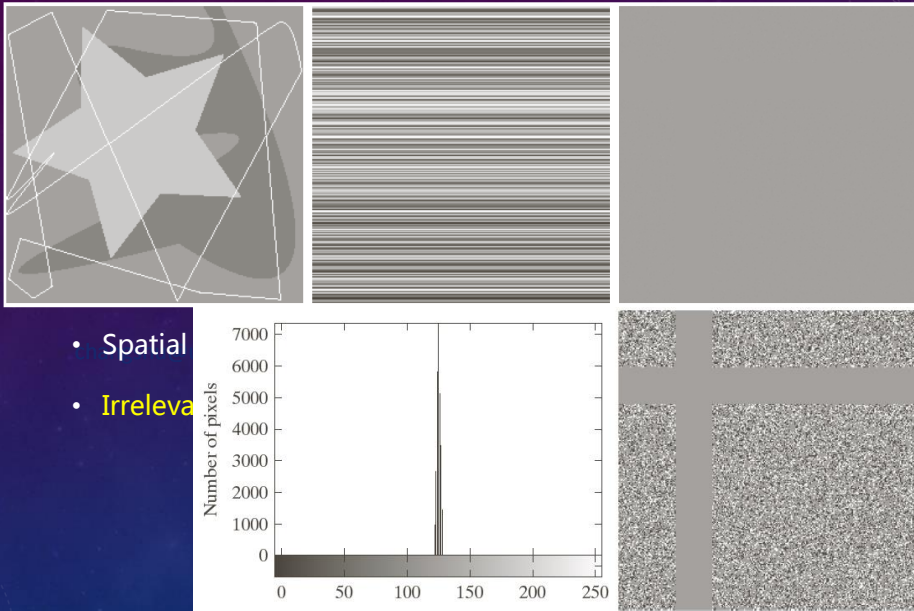


IMAGE COMPRESSION

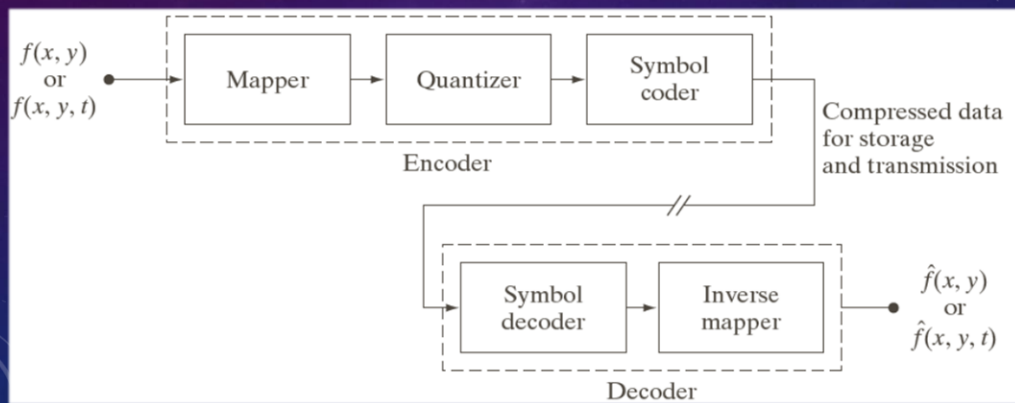
0.631	0.642	0.641	0.628	0.288	0.325	0.375	0.419
0.700	0.716	0.721	0.710	0.435	0.493	0.568	0.630
0.711	0.811	0.842	0.845	0.558	0.635	0.742	0.835





## IMAGE COMPRESSION

- Image compression models

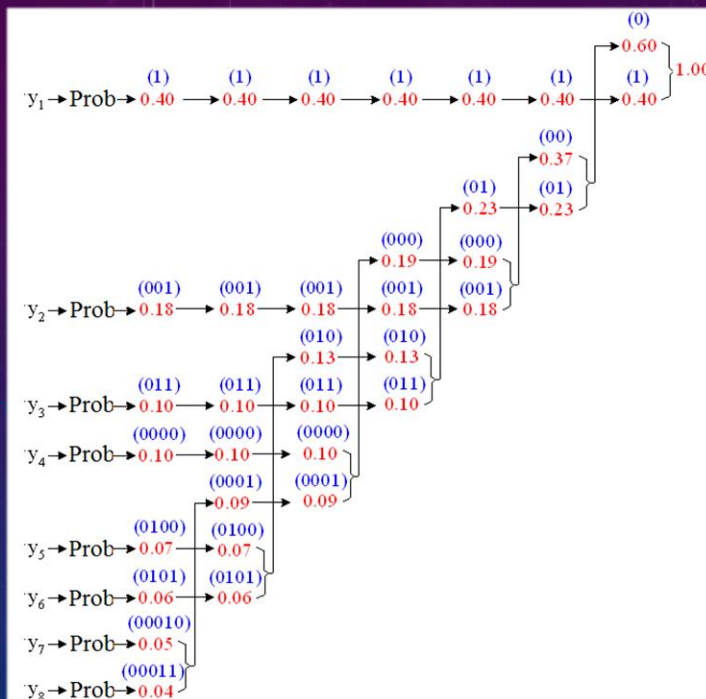




# IMAGE COMPRESSION

Huffman's procedure:

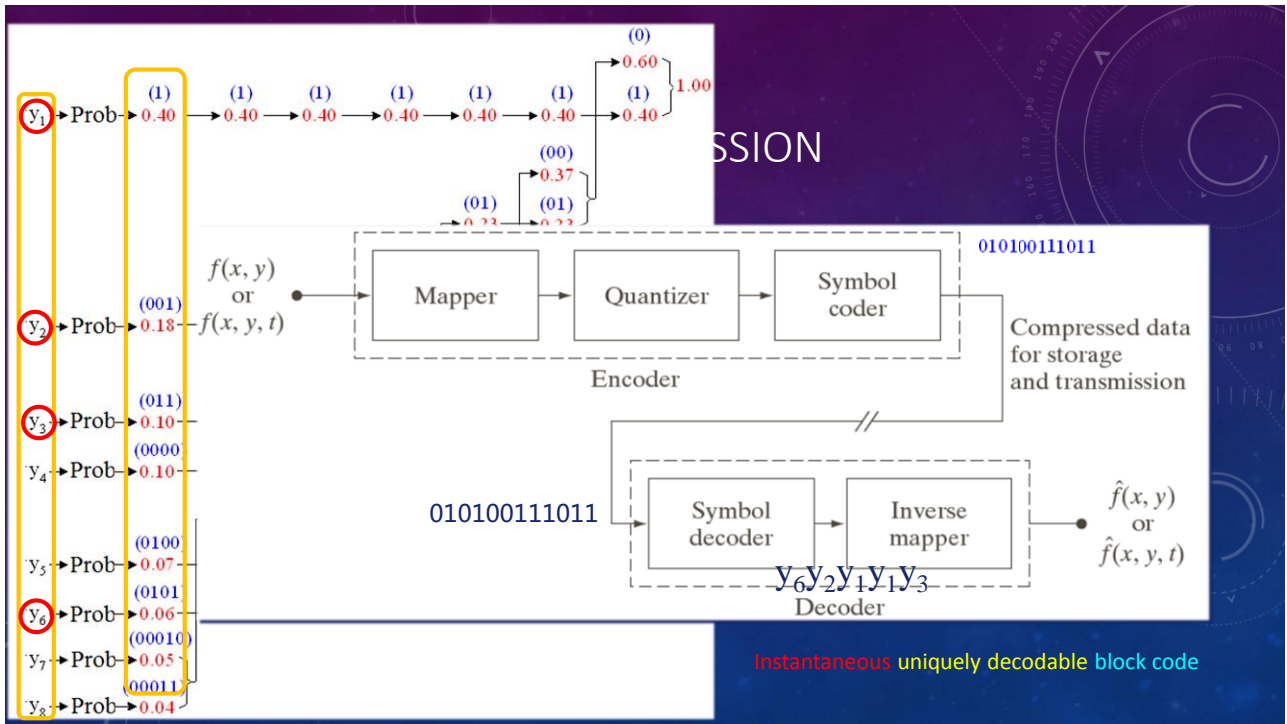
- (1) Create a series of source reductions by **ordering the probabilities of the symbols** under consideration and **combining the lowest probability symbols into a single symbol** that replaces them in the next source reduction
- (2) Code each reduced source, starting with the smallest source and working back to the original source.





What is the code symbols for the encoded string 010100111011 :

7



## IMAGE COMPRESSION

- Huffman code is an instantaneous uniquely decodable block code
  - block code --- each source symbol is mapped into a fixed sequence of code symbols
  - instantaneous --- each code word in a string of code symbols can be decoded without referencing succeeding symbols
  - uniquely decodable --- any string of code symbols can be decoded in only one way

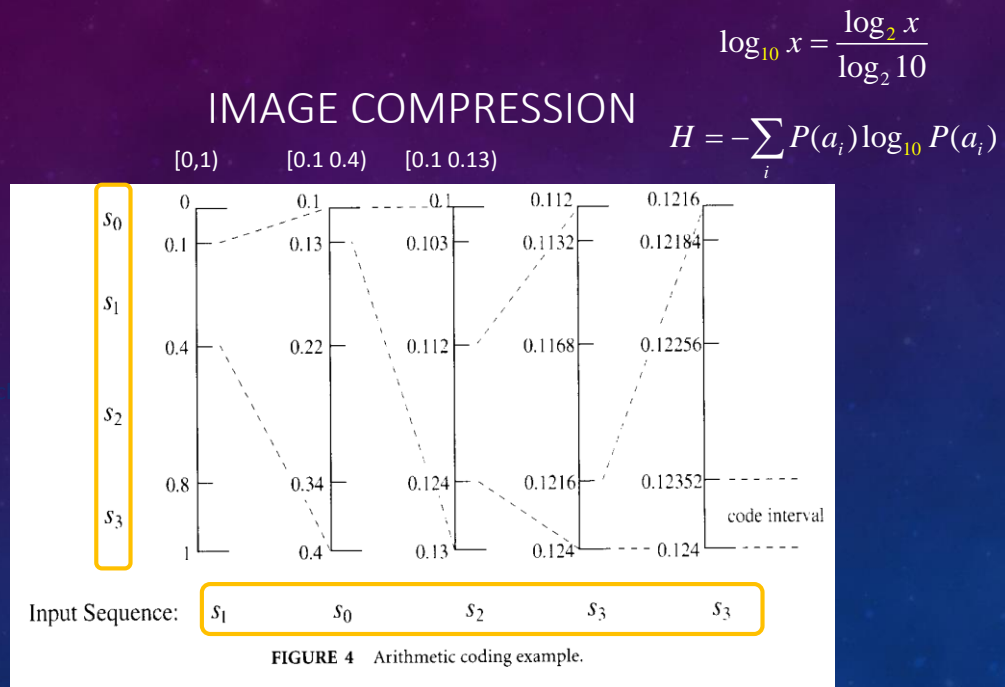


## IMAGE COMPRESSION

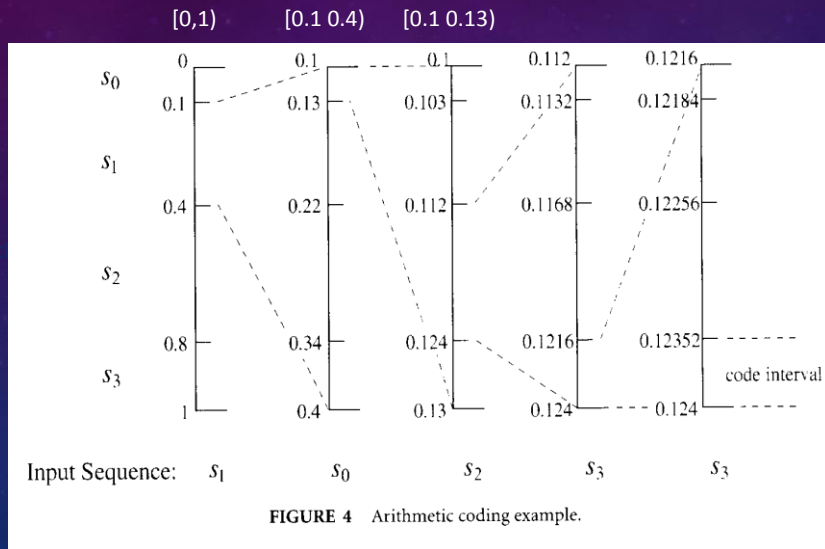
### ➤ Arithmetic coding :

An **entire sequence** of source symbols is assigned a **single** arithmetic code word. The code word itself defines an interval of real numbers between 0 and 1.

As the number of symbols in the message increases, the interval used to represent it becomes smaller and the number of information units required to represent the interval becomes larger.



# IMAGE COMPRESSION



Question:  
What's the  
input  
sequence for a  
given  
arithmetically-  
coded  
message 0.618  
if S3 is the  
end-of-  
message  
indicator?

$S_2S_2S_1S_3$