# M482 Homework 2

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# 1 Practice with definitions.

 $(G_1)$  [image here]

Formally,  $G_1$  is the pairing of the set of states  $S = \{a, b, c, d\}$  with the relation  $R = \{(a, a), (b, b), (c, c), (d, d), (b, a), (b, c), (c, d)\}.$ 

reflexive True, all nodes a, b, c, d bear R to themselves.

symmetric False, counter-example: bRc but cRb does not exist.

anti-symmetric Yes, since there is no symmetry, the only time aRb and bRa holds is when a=b. We know the graph is reflexive, so a=b holds for all nodes.

transitive False, counter-example: bRc and cRd but bRd does not exist.

euclidean False, in the case of b when bRa and bRb, aRb does not exist.

equivalence relation False, only reflexive, i.e. not symmetric and not transitive.

partial function False, counter-example: bRa and bRc but  $a \neq c$ .

 $(G_2)$  [image here]

Formally,  $G_2$  is the pairing of the set of states  $S = \{a, b, c\}$  with the relation  $R = \{(a, b), (a, c), (b, c)\}.$ 

reflexive False, there is no node  $s \in S$  such that sRs.

symmetric False, counter-example: aRb but bRa does not exist.

anti-symmetric True, vacuously because  $xRy \wedge yRx$  is false since no such nodes x,y exist.

transitive True, aRb and bRc and aRc, and  $bRc \wedge cRb$  is false so vacuously true, and similar case for c.

euclidean  $\forall a, b, c. \ aRb \land aRc \Rightarrow bRc.$ 

x	$\mid y \mid$	z	xRy	xRz	$xRy \wedge xRz : (P)$	yRz:(Q)	$P \Rightarrow Q$
a	b	c	T	T	T	T	T
a	c	b	T	T	T	T	T
b	a	c	F	T	F	T	T
b	c	a	T	F	F	T	T
c	a	b	F	F	F	T	T
c	b	a	F	F	F	F	T

#### equivalence relation

# partial function

 $(G_3)$  [image here]

Formally,  $G_3$  is the pairing of the set of states  $S = \{a, b\}$  with the relation  $R = \{(a, a), (a, b), (b, a)\}.$ 

reflexive

symmetric

anti-symmetric

transitive

euclidean

euclidean  $\forall a, b, c. \ aRb \land aRc \Rightarrow bRc.$ 

x	y	z	xRy	xRz	$xRy \wedge xRz : (P)$	yRz:(Q)	$P \Rightarrow Q$
a	b	c	T	T	T	T	T
$\mid a \mid$	c	b	T	T	T	T	T
b	a	c	F	T	F	T	T
b	c	a	T	F	F	T	T
c	a	b	F	F	F	T	T
c	b	a	F	F	F	F	T

# equivalence relation

partial function

 $(G_4)$  [image here]

Formally,  $G_4$  is the pairing of the set of states  $S = \{a\}$  with the relation  $R = \{(a, a)\}.$ 

reflexive

symmetric

anti-symmetric

transitive

euclidean

euclidean  $\forall a, b, c. \ aRb \land aRc \Rightarrow bRc.$ 

$\boldsymbol{x}$	y	z	xRy	xRz	$xRy \wedge xRz : (P)$	yRz:(Q)	$P \Rightarrow Q$
a	b	c	T	T	T	T	T
a	c	b	T	T	T	T	T
b	a	c	F	T	F	T	T
b	c	a	T	F	F	T	T
c	a	b	F	F	F	T	T
c	b	a	F	F	F	F	T

# equivalence relation

#### partial function

 $(G_5)$  [image here] Formally,  $G_5$  is the pairing of the set of states  $S = \{a, b, c\}$  with the relation  $R = \{(a, b), (b, a), (b, c), (c, b)\}.$ 

reflexive

 $\operatorname{symmetric}$ 

anti-symmetric

transitive

euclidean

euclidean  $\forall a, b, c. \ aRb \land aRc \Rightarrow bRc.$ 

x	$\mid y \mid$	z	xRy	xRz	$xRy \wedge xRz : (P)$	yRz:(Q)	$P \Rightarrow Q$
a	b	c	T	T	T	T	T
$\mid a \mid$	c	b	T	T	T	T	T
b	a	c	F	T	F	T	T
b	c	a	T	F	F	T	T
c	a	b	F	F	F	T	T
c	b	a	F	F	F	F	T

#### equivalence relation

# partial function

 $(G_6)$  [image here]

Formally,  $G_6$  is the pairing of the set of states  $S = \{a, b, c\}$  with the relation  $R = \{(a, a), (b, b), (c, c), (a, b), (b, a), (b, c), (c, b)\}.$ 

reflexive

symmetric

anti-symmetric

transitive

euclidean  $\forall a, b, c. \ aRb \land aRc \Rightarrow bRc.$ 

$\boldsymbol{x}$	y	z	xRy	xRz	$xRy \wedge xRz : (P)$	yRz:(Q)	$P \Rightarrow Q$
a	b	c	T	T	T	T	T
a	c	b	T	T	T	T	T
b	a	c	F	T	F	T	T
b	c	a	T	F	F	T	T
c	a	b	F	F	F	T	T
c	b	a	F	F	F	F	T

#### equivalence relation

# partial function

# $(G_7)$ [image here]

Formally,  $G_7$  is the pairing of the set of states  $S = \{a, b, c\}$  with the relation  $R = \{(a, a), (b, b), (c, c), (b, a), (b, c)\}.$ 

reflexive

symmetric

anti-symmetric

transitive

euclidean  $\forall a, b, c. \ aRb \land aRc \Rightarrow bRc.$ 

	$\boldsymbol{x}$	y	z	xRy	xRz	$xRy \wedge xRz : (P)$	yRz:(Q)	$P \Rightarrow Q$
ſ	a	b	c	T	T	T	T	T
	a	c	b	T	T	T	T	T
	b	a	c	F	T	F	T	T
	b	c	a	T	F	F	T	T
	c	a	b	F	F	F	T	T
	c	b	a	F	F	F	F	T

#### equivalence relation

# partial function

# $(G_8)$ [image here]

Formally,  $G_8$  is the pairing of the set of states  $S = \{a, b, c\}$  with the relation  $R = \{(a, b), (a, c), (b, b), (b, c), (c, c), (c, b)\}.$ 

reflexive

symmetric

anti-symmetric

transitive

euclidean  $\forall a, b, c. \ aRb \land aRc \Rightarrow bRc.$ 

$\boldsymbol{x}$	y	z	xRy	xRz	$xRy \wedge xRz : (P)$	yRz:(Q)	$P \Rightarrow Q$
a	b	c	T	T	T	T	T
a	c	b	T	T	T	T	T
b	a	c	F	T	F	T	T
b	c	a	T	F	F	T	T
c	a	b	F	F	F	T	T
c	b	a	F	F	F	F	T

equivalence relation

partial function

# 2 True or false, and why?

The statements below are about relations on a fixed set X. For each statement below, say whether it is true or false. If it is true, give a short proof. If it is false, give a counter-example.

- 1. Every reflexive relation on X is serial.
- 2. Every reflexive relation on X which is reflexive and transitive is also symmetric.
- 3. Every reflexive relation on X which is Euclidean, symmetric and transitive is reflexive.
- 4. Every reflexive relation on X which is Euclidean and reflexive is symmetric.

# 3 Multiply two graphs

# 4 True or false, and why?

For each of the following assertions, tell whether it is true or false. (As always in a problem like this, "true" means *true for all graphs*. So "false" means *false for some graphs*.) Give a short proof of the true ones, and a counter-example for the false ones.

- 1. The product of two reflexive graphs is reflexive.
- 2. The product of two euclidean graphs is reflexive.
- 3. If the product of two graphs is reflexive, then both of the given graphs are reflexive.

- $4.\$  If the product of two graphs is transitive, then both of the given graphs are transitive.
- 5. Let G be a one-point graph with no arrows. Then for all  $H,G\times H$  looks like H.

# 5 Modal semantics with boxes and diamonds rather than K-operators

Attached.