

Modal Logic, Winter 2019

Homework 8

due Tuesday, March 19

1 Sentences from the past

We shall use the following notation at several points later on.

$$\begin{array}{ll}
 \alpha_1 = \Box F \wedge p & \alpha_5 = \Box \neg p \wedge \neg \Diamond p \wedge p \\
 \alpha_2 = \Box F \wedge \neg p & \alpha_6 = \Box \neg p \wedge \neg \Diamond p \wedge \neg p \\
 \alpha_3 = \Box p \wedge \Diamond p \wedge p & \alpha_7 = \Diamond p \wedge \Diamond \neg p \wedge p \\
 \alpha_4 = \Box p \wedge \Diamond p \wedge \neg p & \alpha_8 = \Diamond p \wedge \Diamond \neg p \wedge \neg p
 \end{array} \tag{1}$$

Check that $\alpha_1, \dots, \alpha_8$ are satisfiable and hence consistent in K .

[Hint: it's not hard to do this directly by drawing pictures. But if you wish, you might look back at Homework 2, problem 5.]

Recall the α_i sentences from Exercise 1 above. Fill in the following chart, putting a \checkmark in the box if the given sentence is consistent in the given logic, and an \times if it is not. A few of the entries are given for you.

	α_1	α_2	α_3	α_4	α_5	α_6	α_7	α_8
K	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark
KT	\times	\times	\checkmark	\times	\times	\checkmark	\checkmark	\checkmark
KD	\times	\times	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark
KB	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark
$K4$	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark
$S4$	\times	\times	\checkmark	\times	\times	\checkmark	\checkmark	\checkmark
$S5$	\times	\times	\checkmark	\times	\times	\checkmark	\checkmark	\checkmark

Reasons for above are attached on scratch paper.

2 An unprovable sentence

Show that the B axiom $p \rightarrow \Box \Diamond p$ is not provable in $KD45$. [The way to start is to find a graph which is serial, transitive, and euclidean, and at the same time is not symmetric. Then one needs to find a point and put a valuation on the graph so that the point does not satisfy the given B -axiom.]

Consider the following model $M = (\{x, y\}, \{(x, y), (y, y)\}, Val)$ where $Val(x) = \{p\}$ and $Val(y) = \emptyset$. This model is serial, transitive, and euclidean while also not symmetric. At point x , p is true but $\Box \Diamond p$ is not. By this counter-example, we know the B axiom is not provable in $KD45$.

3 The logics KB and KB'

Recall that the B axioms are the sentences of the form $\varphi \rightarrow \Box\Diamond\varphi$. Let's call the B' axioms the sentences of the form $(\varphi \wedge \Diamond\psi) \rightarrow \Diamond(\psi \wedge \Diamond\varphi)$. So we get a logical system KB' by adding the B' axioms to K .

Prove that $KB' \leq KB$.

1.	$\varphi \wedge \Diamond\psi$	Assume
2.	$\Diamond\psi$	$\wedge_e, 1$
3.	φ	$\wedge_e, 1$
4.	$\Box\Diamond\varphi$	B axiom, 3
5.	$\Diamond\psi \wedge \Box\Diamond\varphi$	$\wedge_i, 2, 4$
6.	$\Diamond\psi \wedge \Box\Diamond\varphi \rightarrow \Diamond(\psi \wedge \Diamond\varphi)$	Powerpoint
7.	$\Diamond(\psi \wedge \Diamond\varphi)$	$\rightarrow_e, 5, 6$
8.	$(\varphi \wedge \Diamond\psi) \rightarrow \Diamond(\psi \wedge \Diamond\varphi)$	$\rightarrow_i, 1 - 9$

4 The Converse

Continuing with the last problems, show that $KB \leq KB'$.

1.	φ	Assume
2.	$\neg \Box \Diamond \varphi$	Assume
3.	$\Diamond \neg \Diamond \varphi$	$\neg \Box \equiv \Diamond \neg$
4.	$\Diamond \Box \neg \varphi$	$\neg \Diamond \equiv \Box \neg$
5.	$\varphi \wedge \Diamond \Box \neg \varphi$	$\wedge_i, 1, 4$
6.	$(\varphi \wedge \Diamond \Box \neg \varphi) \rightarrow \Diamond(\Box \neg \varphi \wedge \Diamond \varphi)$	B' axiom, 5
7.	$\Diamond(\Box \neg \varphi \wedge \Diamond \varphi)$	$\rightarrow_e, 5, 6$
8.	$\Box \neg \varphi \wedge \Diamond \varphi$	Assume
9.	$\Box \neg \varphi$	$\wedge_e, 8$
10.	$\Diamond \varphi$	$\wedge_e, 8$
11.	F	$F_i, 9, 10$
12.	$\neg(\Box \neg \varphi \wedge \Diamond \varphi)$	$\neg_i, 8 - 11$
13.	$\Box \neg(\Box \neg \varphi \wedge \Diamond \varphi)$	$*_i, 8 - 12$
14.	F	$F_i, 7, 13$
15.	$\neg \neg \Box \Diamond \varphi$	$\neg_i, 2 - 14$
16.	$\Box \Diamond \varphi$	PR
17.	$\varphi \rightarrow \Box \Diamond \varphi$	$\rightarrow_i, 1 - 16$

It follows that the two logical systems KB and KB' can prove each other's axioms, and hence prove all the same sentences. That is, if $\vdash \varphi$ in KB , then $\vdash \varphi$ in KB' ; and vice-versa.

5 A result on KB

1. Assume that $\psi \wedge \Diamond\varphi$ is inconsistent in KB' . (This means that in KB' , $\vdash \neg(\psi \wedge \Diamond\varphi)$.) Show that in KB' , $\varphi \wedge \Diamond\psi$ is also inconsistent. [Hint: you will need to use the following axiom of KB' :

$$(\varphi \wedge \Diamond\psi) \rightarrow \neg\Box\neg(\psi \wedge \Diamond\varphi).$$

(I have “unabbreviated” the \Diamond on the right.)]

1.	$\varphi \wedge \Diamond\psi$	Assume
2.	$\neg(\psi \wedge \Diamond\varphi)$	Given as assumption
3.	$\Box\neg(\psi \wedge \Diamond\varphi)$	$*_i, 2$
4.	$(\varphi \wedge \Diamond\psi) \rightarrow \neg\Box\neg(\psi \wedge \Diamond\varphi)$	B' axiom
5.	$\neg\Box\neg(\psi \wedge \Diamond\varphi)$	$\rightarrow_e, 4, 1$
6.	F	$F_i, 3, 5$
7.	$\neg(\varphi \wedge \Diamond\psi)$	$\neg_i, 1 - 6$

2. Using the last part, prove the following: if $\varphi \wedge \Diamond\psi$ is consistent in KB' , then $\psi \wedge \Diamond\varphi$ is also consistent in KB' .

We prove the contrapositive. We assume $\psi \wedge \Diamond\varphi$ is inconsistent in KB' and prove $\varphi \wedge \Diamond\psi$ is inconsistent in KB' . For this, we have just constructed a proof above of $\vdash \neg(\varphi \wedge \Diamond\psi)$ which is the definition of inconsistent. Hence we're done.

3. Show that KB and KB' have the same consistent sentences.

Fix a sentence φ . Assume about φ that φ is consistent in KB . We prove φ is consistent in KB' . So in KB , $\not\vdash \neg\varphi$. Suppose, towards a contradiction, that φ is inconsistent in KB' ; this means in KB' , $\vdash \neg\varphi$. By the previous result (5.3), $\vdash \neg\varphi$ in KB , a contradiction! Hence φ is consistent in KB' .

4. Putting things together, show that if $\varphi \wedge \Diamond\psi$ is consistent in KB , then $\psi \wedge \Diamond\varphi$ is also consistent in KB .

Assume $\varphi \wedge \Diamond\psi$ is consistent in KB . We prove $\psi \wedge \Diamond\varphi$ is consistent in KB . By part 3, we know that $\varphi \wedge \Diamond\psi$ is also consistent in KB' . By part 2, we know that if $\varphi \wedge \Diamond\psi$ is consistent in KB' then $\psi \wedge \Diamond\varphi$ is consistent in KB' . By part 3 again, we know that $\psi \wedge \Diamond\varphi$ is consistent in KB . Done.