

Homework 2

Divide-and-Conquer Algorithms, Sorting Algorithms, Greedy Algorithms

Deadline: March 14, 11:59pm.

Available points: 110. Perfect score: 100.

Homework Instructions:

Teams: Homeworks should be completed by teams of students - three at most. No additional credit will be given for students that complete a homework individually. Please inform Athanasios Krontiris **if your team has changed from the previous assignment.** (email: ak979 AT cs.rutgers.edu).

Submission Rules: Submit your solutions electronically as a PDF document through sakai.rutgers.edu. Do not submit Word documents, raw text, or hardcopies etc. Make sure to generate and submit a PDF instead. Each team of students should submit only a single copy of their solutions and indicate all team members on their submission. Failure to follow these rules will result in lower grade in the assignment.

Late Submissions: No late submission is allowed. If you don't submit a homework on time, you get 0 points for that homework.

Extra Credit for L^AT_EX: You will receive 10% extra credit points if you submit your answers as a typeset PDF (using L^AT_EX, in which case you should also submit electronically your source code). Resources on how to use L^AT_EX are available on the course's website. There will be a 5% bonus for electronically prepared answers (e.g., on MS Word, etc.) that are not typeset.

25% Rule: For any homework problem (same will hold for exam questions), you can either attempt to answer the question, in which case you will receive between 0 and 100% credit for that question (i.e., you can get partial credit), or you can write "I don't know", in which case you receive 25% credit for that question. Leaving the question blank is the same as writing "I don't know." You can and will get less than 25% credit for a question that you answer erroneously.

Handwritten Reports: If you want to submit a handwritten report, scan it and submit a PDF via Sakai. We will not accept hardcopies. If you choose to submit handwritten answers and we are not able to read them, you will not be awarded any points for the part of the solution that is unreadable.

Precision: Try to be precise. Have in mind that you are trying to convince a very skeptical reader (and computer scientists are the worst kind...) that your answers are correct.

Collusion, Plagiarism, etc.: Each team of students must prepare its solutions independently from other teams, i.e., without using common notes or worksheets with other students or trying to solve problems in collaboration with other teams. You must indicate any external sources you have used in the preparation of your solution. Do not plagiarize online sources and in general make sure you do not violate any of the academic standards of the course, the department or of the university (the standards are available through the course's website). Failure to follow these rules may result in failure in the course.

Part A (20 points)

Problem 1: The more general version of the Master Theorem is the following. Given a recurrence of the form:

$$T(n) = aT\left(\frac{n}{b}\right) + f(n)$$

where $a \geq 1$ and $b > 1$ are constants and $f(n)$ is an asymptotically positive function, there are 3 cases:

1. If $f(n) = O(n^{\log_b a - \epsilon})$ for some constant $\epsilon > 0$, then $T(n) = \Theta(n^{\log_b a})$.
2. If $f(n) = \Theta(n^{\log_b a} \log^k n)$ with $k \geq 0$, then $T(n) = \Theta(n^{\log_b a} \log^{k+1} n)$. In most cases, $k = 0$.
3. If $f(n) = \Omega(n^{\log_b a + \epsilon})$ with $\epsilon > 0$, and $f(n)$ satisfies the regularity condition, then $T(n) = \Theta(f(n))$. The regularity condition specifies that $af(\frac{n}{b}) \leq cf(n)$ for some constant $c < 1$ and all sufficiently large n .

Give asymptotic bounds for the following recurrences. Assume $T(n)$ is constant for $n = 1$. Make your bounds as tight as possible, and justify your answers.

A. $T(n) = 2T(\frac{n}{4}) + n^{0.51}$

From $T(n) = 2T(\frac{n}{4}) + n^{0.51}$, we have $a = 2$, $b = 4$, and $f(n) = n^{0.51}$. Therefore, $n^{\log_b a} = n^{\log_4 2} = n^{0.5}$. We compare $n^{0.51}$ to $n^{0.5}$, and we know that $n^{0.51} > n^{0.5}$, so it satisfies case 3 because $f(n) = n^{0.51}$, which is bounded below by $n^{0.5}$. Now we have to check the regularity condition, so $2f(\frac{n}{4}) \leq cf(n)$

B. $T(n) = 16T(\frac{n}{4}) + n!$

C. $T(n) = \sqrt{2}T(\frac{n}{2}) + \lg n$

D. $T(n) = T(n-1) + \lg n$

E. $T(n) = 5T(n/5) + \frac{n}{\lg n}$

Part B (25 points)

Problem 2: You are in the HR department of a technology firm, and here is a job for you. There are n different projects, and n different programmers.

Every project has its unique payoff when completed and level of difficulty (which are uniform, regardless which programmer will work on the project). Every programmer has a unique skill set as well as expectations for compensation (which are uniform, regardless the project the programmer will work on). You cannot directly collect information that allows you to compare the payoff or difficulty level of two projects, or the capability or expectations for compensation of two programmers.

Instead, you can arrange a meeting between each project manager and programmer. In each meeting, the project manager will give the programmer an interview to see whether the programmer can do the project; the programmer can ask the project manager about the compensation to see whether it meets her expectations. After the meeting, you can get a result based on the feedback of the project manager and the programmer. The result can be:

1. The programmer can't do the project.
2. The programmer can do the project, but the compensation of the project doesn't meet her expectation.
3. The programmer can do the project, and the compensation for the project matches her expectations. At this time, we say the project and the programmer *match* with each other.

Assume that the projects and programmers match one to one. Your goal is to match each programmer to a project.

- A. Show that any algorithm for this problem must need $\Omega(n \log n)$ meetings in the worst case.
- B. Design a randomized algorithm for this problem that runs in expected time $O(n \log n)$.

Part C (40 points)

Problem 3: A nation-wide programming contest is held at k universities in North America. The i^{th} university has m_i participants. The total number of participants is n , i.e., $n = \sum_{i=1}^k m_i$. In the contest, participants have to write programs to solve 6 problems. Each problem contains 10 test cases, each test case is worth 10 points. Participants aim to maximize their collected points.

After the contest, each university sorts the scores of participants belonging to it and submits the grades to the organizer. Then the organizer has to collect the sorted scores of participants and provide a final sorted list for all participants.

1. For each university, how do they sort the scores of participants belonging to it? Please briefly describe a comparative sorting algorithm that is appropriate for this purpose and a non-comparative sorting algorithm that works in this setup.
2. How does the organizer sort the scores of participants given k files, where each file includes the sorted scores of participants from a specific university? Please describe an algorithm with a $O(n \log k)$ running time and justify its time complexity.
3. Suppose the organizer want to figure out the participants of ranking r . Given k sorted files, how does the organizer find the r th largest scores without sorting the scores of participants? Please describe an algorithm in $O(k(\sum_{i=1}^k \log m_i))$ time and justify its time complexity.

Could you do this in $O(k + \log k \sum_{i=1}^n \log m_i)$ time?

[Hint: If r is larger than $\frac{n}{2}$, the elements that have at least $\frac{n}{2}$ elements larger than them should not be possible answer. The problem is how to identify these elements.]

Problem 4: You have a collection of n New York Times crossword puzzles from 01/01/1943 until 12/31/2012 stored in a database. The only operations that you can perform to the database are the following:

- crossword_puzzle $x \leftarrow \text{getPuzzle}(\text{int index})$; where the index is between 1 and n ; the puzzles are *not* sorted in the database in terms of the date they appeared.
- $\text{getDay}(\text{crossword_puzzle } x)$; which returns a number between 1 to 31.
- $\text{getMonth}(\text{crossword_puzzle } x)$; which returns a number between 1 to 12.
- $\text{getYear}(\text{crossword_puzzle } x)$; which returns a number between 1943 to 2012.

All of the above queries can be performed in constant time. You have found out that the number of puzzles is less than the number of days in the above period (from 01/01/1943 until 12/31/2012) by one, i.e., one crossword puzzle was not included in the database. We need to identify the date of the missing crossword puzzle.

Design a linear-time algorithm that minimizes the amount of space that it is using to find the missing date. Ignore the effect of leap years.

1. Start by summing up all the years from 1943 to 2012, and save that value into a variable, totalYearSum.
2. Go through the entries in the database using $\text{getPuzzle}()$ starting at zero, and perform a $\text{getYear}()$ on all of them. For every puzzle, take totalYearSum and subtract from it the value you get from $\text{getYear}()$, and save the new value into totalYearSum. By the time all the entries are iterated through and the subtraction is done, you'll end up with totalYearSum equalling the year that you're missing.

3. Sum up the months from 1 (January) to 12 (December), and save that value into a variable, totalMonthSum.
4. Go through the entries in the database again and do a getYear() on them. If the year matches the year of the missing puzzle, do a getMonth() on it. For every puzzle that you do a getMonth() on, take totalMonthSum and subtract from it the value you get from getMonth(), and save the new value into totalMonthSum. By the time all the entries are iterated through and the subtraction is done, you'll end up with totalMonthSum equalling the month that you're missing.
5. Since you know what month you're missing, you know how many days are in that month (ignoring the fact that it could be a leap year). Sum up from 1 to the number of days there are in the month, and save that value into a variable, totalDaySum.
6. Go through the entries in the database again and do a getYear() on them. If the year matches the year of the missing puzzle, do a getMonth() on it. If the month matches the month of the missing puzzle, do a getDay() on it. For every puzzle that you do a getDay() on, take totalDaySum and subtract from it the value you get from getDay(), and save the new value into totalDaySum. By the time all the entries are iterated through and the subtraction is done, you'll end up with the totalDaySum equalling the day that you're missing.
7. Now you have the date of the missing puzzle, and this was done with constant space.

Part D (25 points)

Problem 5: You are running a promotional event for a company during which the plan is to distribute n gifts to the participants. Consider that each gift i is worth an integer number of dollars a_i . There are m people participating in the event, where $m < n$. The j -th person is satisfied if he receives gifts that are worth at least s_j dollars each. The task is to satisfy as many people as possible given that you have a knowledge of the gift amounts a_i and the satisfaction requirements of each person s_j . Give an approximation algorithm for assigning rewards to people with a running time of $O(m \log m + n)$. What is the approximation ratio of your algorithm and why?

Let n be the number of gifts and m be the number of people participating in the event, where $m < n$. To satisfy as many people as possible given the gift amounts a_i and the satisfaction requirements of each person s_j , we will sort gifts in terms of their prices from lowest price to highest price, using radix sort. Because there are n gifts, and radix sort is linear, we have n as our running time now. To continue, we will build a min heap with the people's satisfaction requirements, and building a min heap is linear time, so we have a running time of m in addition to the n from sorting the gifts. To proceed, we will take the minimum satisfaction requirement from the heap, which requires a worst case running time of $\log(m)$ to traverse the heap. We will take the minimum satisfaction requirement and compare it to first value – smallest value – in sorted gifts. If that gift value is greater than or equal to that person's satisfaction requirement, then we remove that person from the heap, and we remove the gift from the list. If that gift value is not greater than or equal to that person's satisfaction requirement, then we continue traversing the sorted gifts, and whenever we hit a gift that satisfies the person, then we remove the person from the heap and the gift from the list. Continue removing the minimum value from the heap until each person gets a gift or until no gift satisfies that person. Therefore, this algorithm will satisfy as many people as possible with the running time of $O(m \log m + n)$.