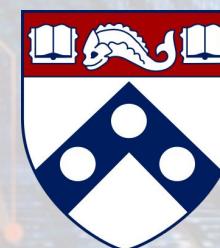


Learning to Reason about Programs

Mayur Naik

Associate Professor
Computer & Information Science

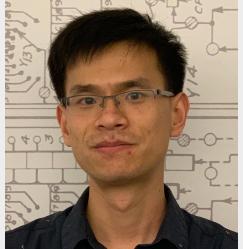


Penn
UNIVERSITY OF PENNSYLVANIA

Acknowledgments



Halley Young



Xujie Si



Sulekha Kulkarni



Mukund
Raghethaman

Postdocs



Kihong
Heo



Pardis
Pashakhanloo



Jiani Huang



Elizabeth
Dinella



Yuan Yang
(PhD Student)



Hanjun Dai
(PhD Student)

Collaborators



Osbert Bastani
(Faculty)



Le Song
(Faculty)

PhD Students

Where is ML for Programming?

DeepMind's AlphaZero AI is the new champion in chess, shogi, and Go

Dec 2018



Forget chess and Go - artificial intelligence can now beat humans in StarCraft II



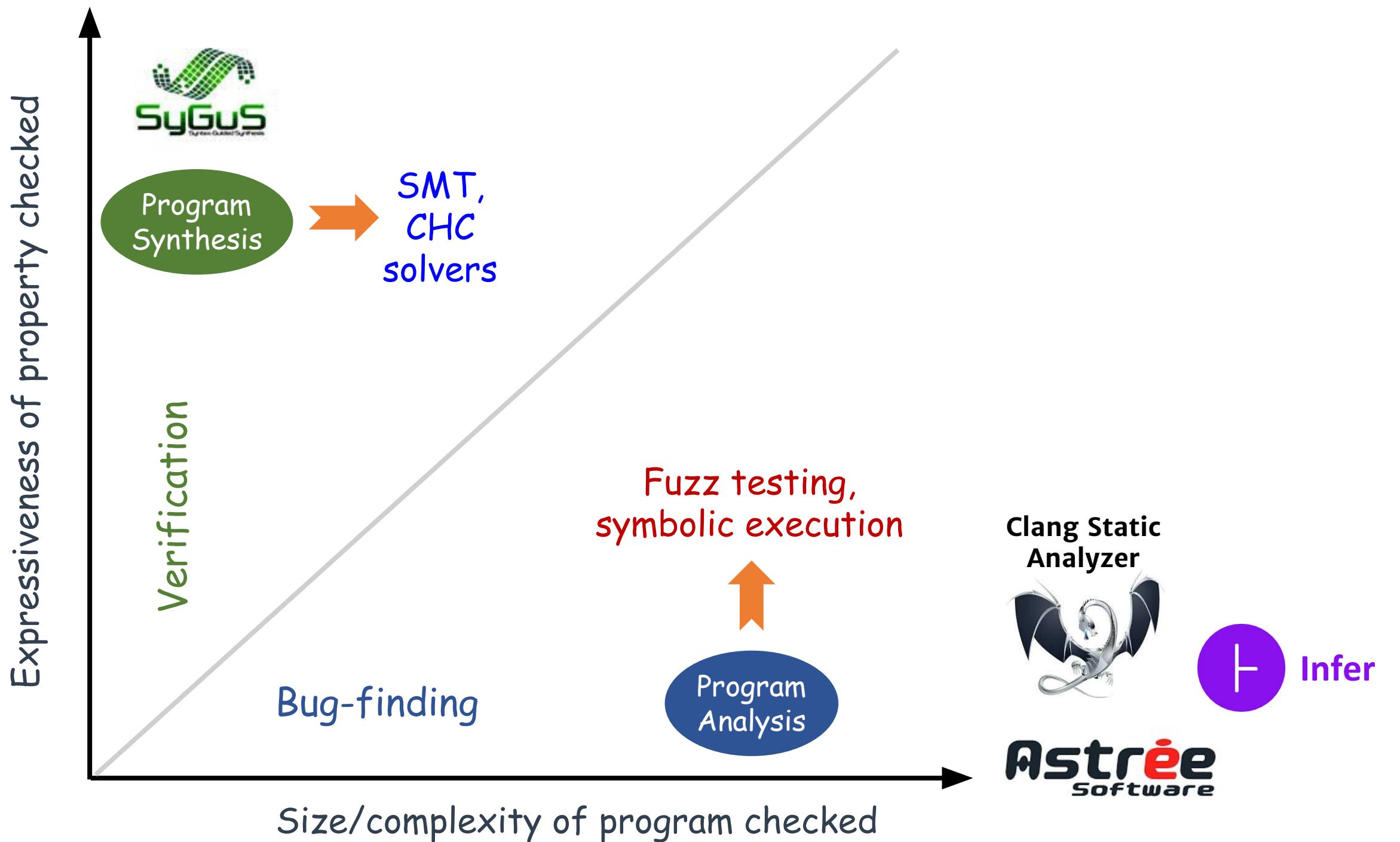
Jan 2019

A.I. Is Flying Drones
(Very, Very Slowly)

Artificial intelligence has bested top players in chess, Go and even StarCraft. But can it fly a drone faster than a pro racer? More than \$1 million is on the line to find out.



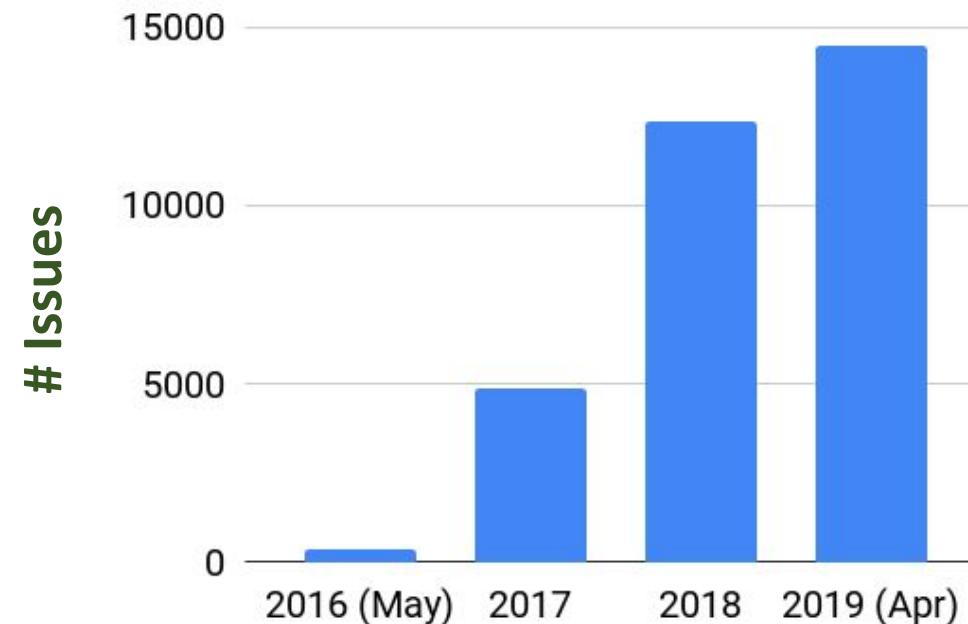
March 2019



State of the Practice: Bug-Finding

OSS-Fuzz: Continuous Fuzzing of Open-Source Software

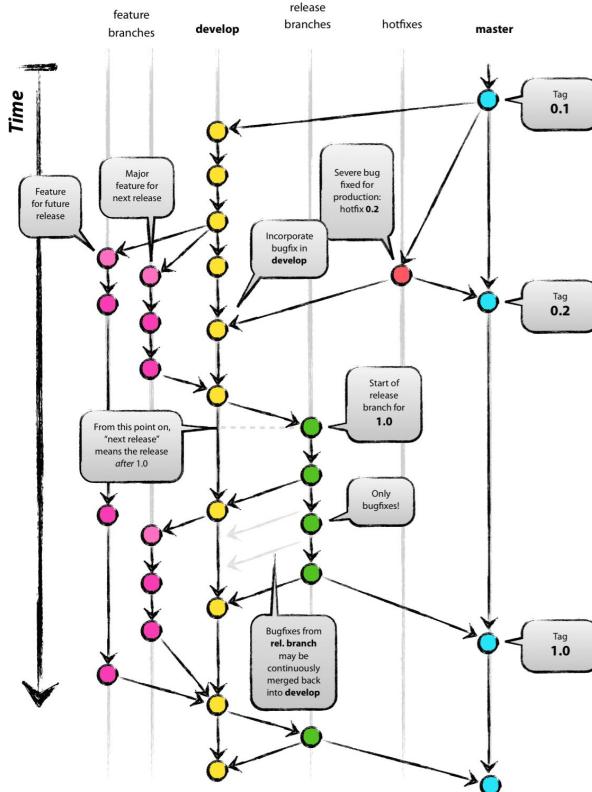
- heap buffer overflows
- global buffer overflows
- stack buffer overflows
- use after frees
- uninitialized memory
- stack overflows
- timeouts
- ooms
- leaks
- ubsan
- unknown crashes
- other (e.g. assertions)



Project	# MLOC
jsc	5.05
gnutls	2.31
llvm	2.18
solidity	2.10
grpc	1.82

(top 5 projects by size)

A Challenge in Bug-Finding



- **Heartbleed**
- **Buffer over-read** bug in OpenSSL's Heartbeat implementation
- Introduced in 2011, discovered in 2014
- Estimated to affect nearly 66% of all web servers

The coding mistake that caused Heartbleed can be [traced to a single line of code](#):

```
memcpy(bp, pl, payload);
```

`memcpy()` is the command that copies data. `bp` is the place it's copying it to, `pl` is where it's being copied from, and `payload` is the length of the data being copied. The problem is that there's never any attempt to check if the amount of data in `pl` is equal to the value given of `payload`.

The most ironic thing here is that OpenSSL is open source software. Anyone could look at the code, and presumably hundreds did, but nobody noticed the fairly elementary coding error.

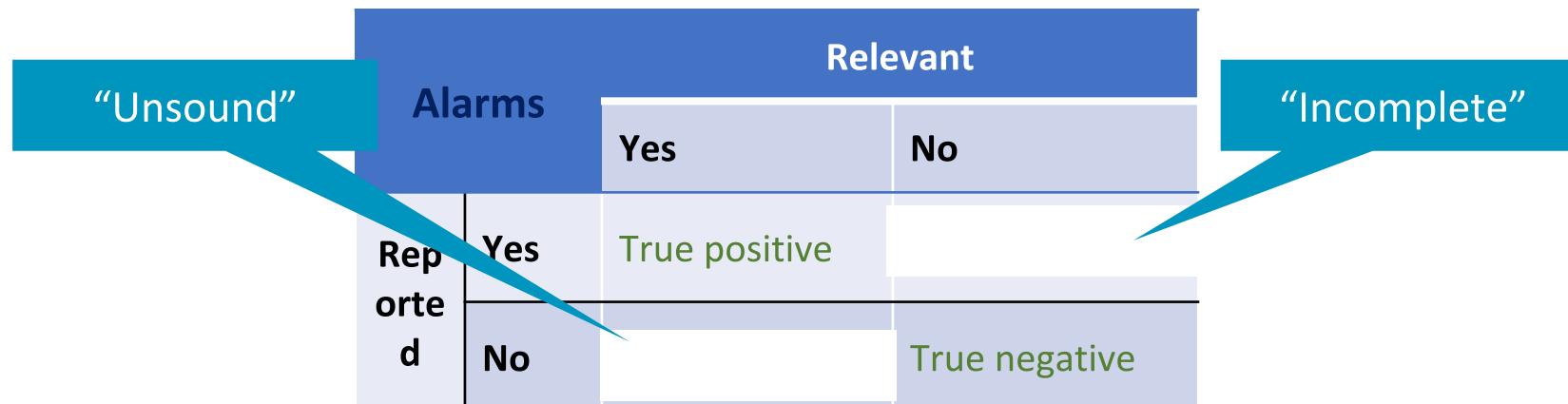
<https://gizmodo.com/how-heartbleed-works-the-code-behind-the-internets-se-1561341209>

Why did program analysis tools not discover the Heartbleed bug?

Approximations in Program Analysis

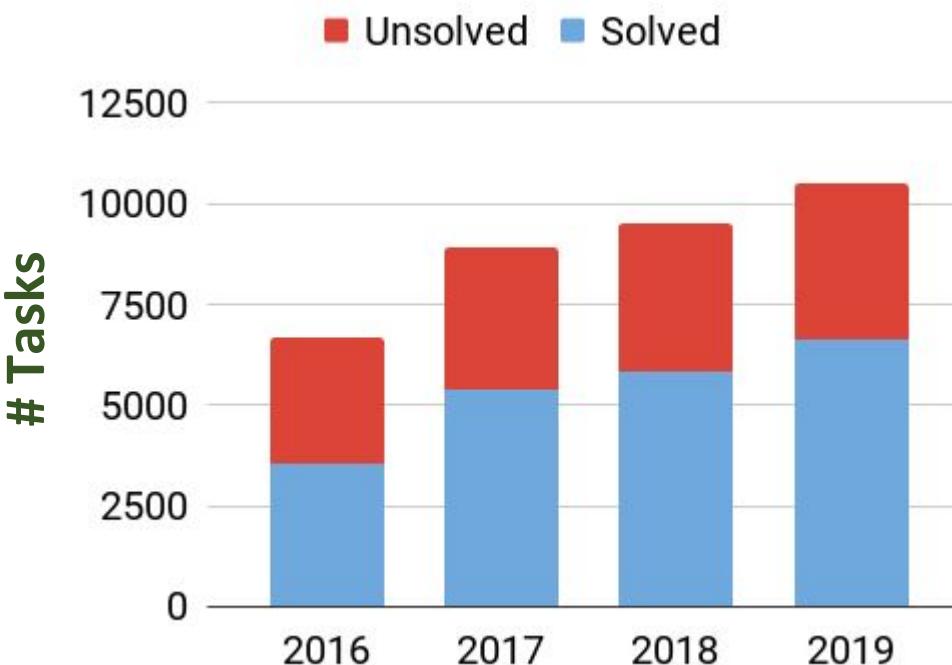
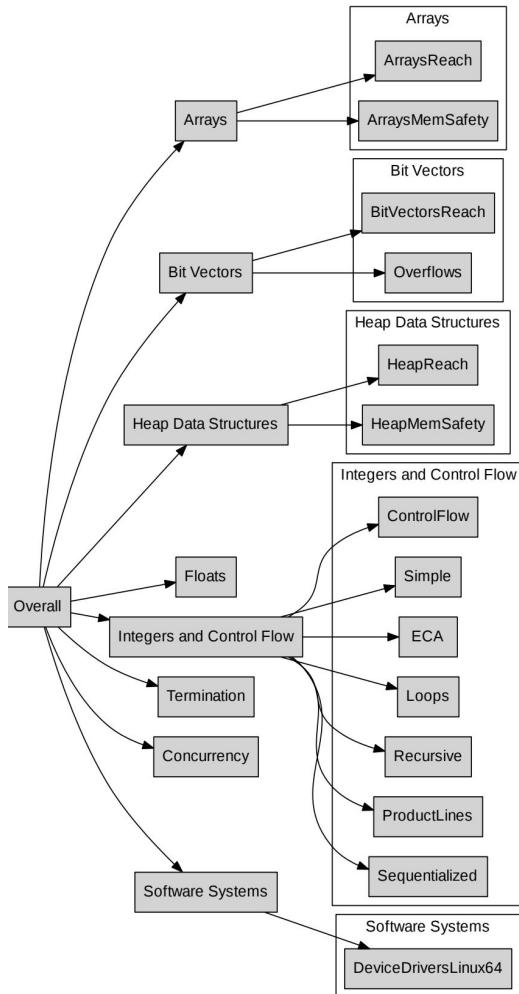
“... can be difficult to do without introducing large numbers of **false positives**, or scaling **performance** exponentially poorly. In this case, **balancing these** and other factors in the analysis design caused us to miss the defect.”

— Coverity, On Detecting Heartbleed with Static Analysis, 2014



State of the Practice: Verification

SV-COMP: International Software Verification Competition



Task	# KLOC
linux-4.2-rc1.tar.xz-32_7a-drivers	228
linux-4.2-rc1.tar.xz-08_1a-drivers	205
eca-programs/Problem103_label59	185
eca-programs/Problem103_label58	185
eca-programs/Problem103_label57	185

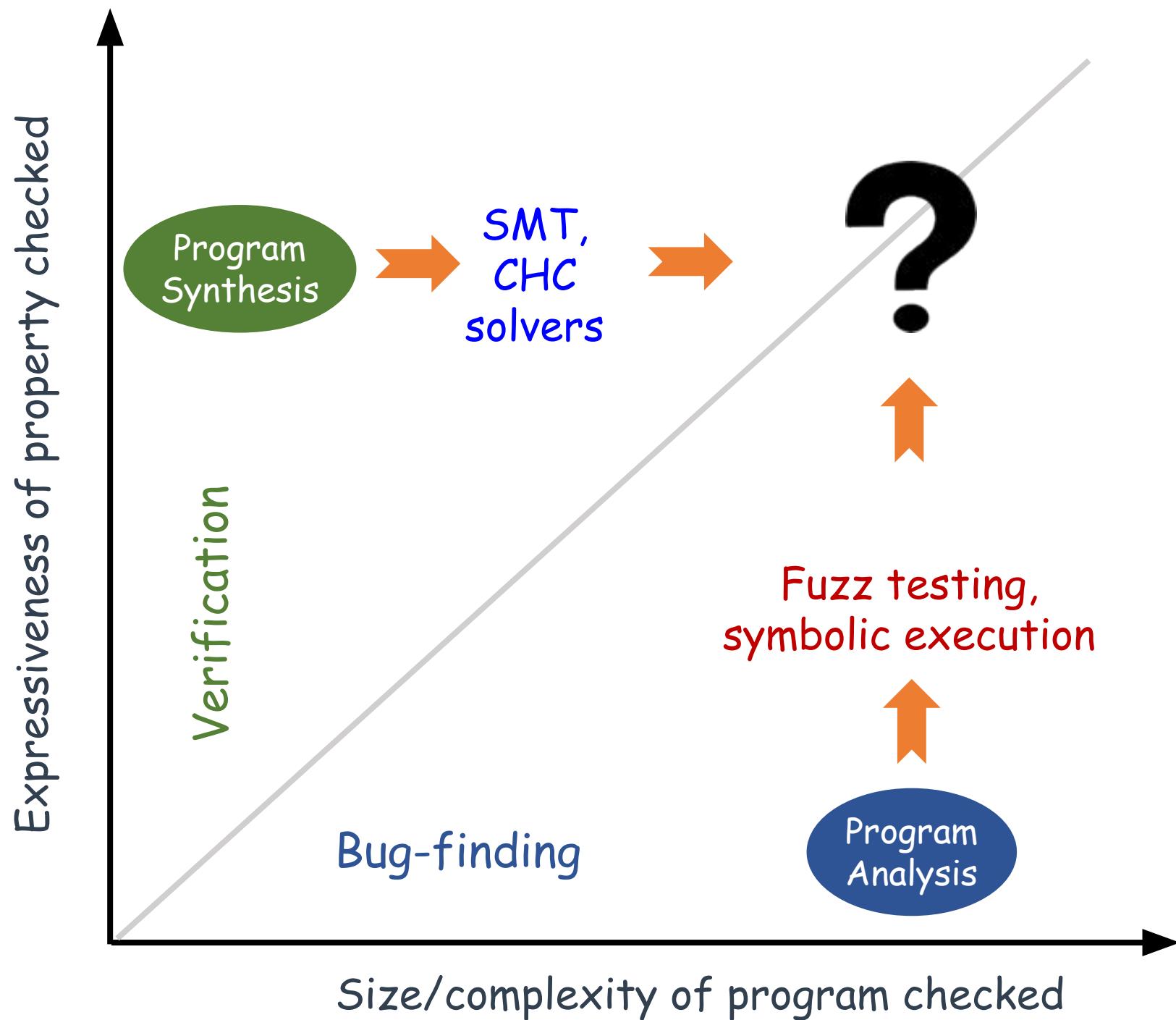
(top 5 solved tasks by size)

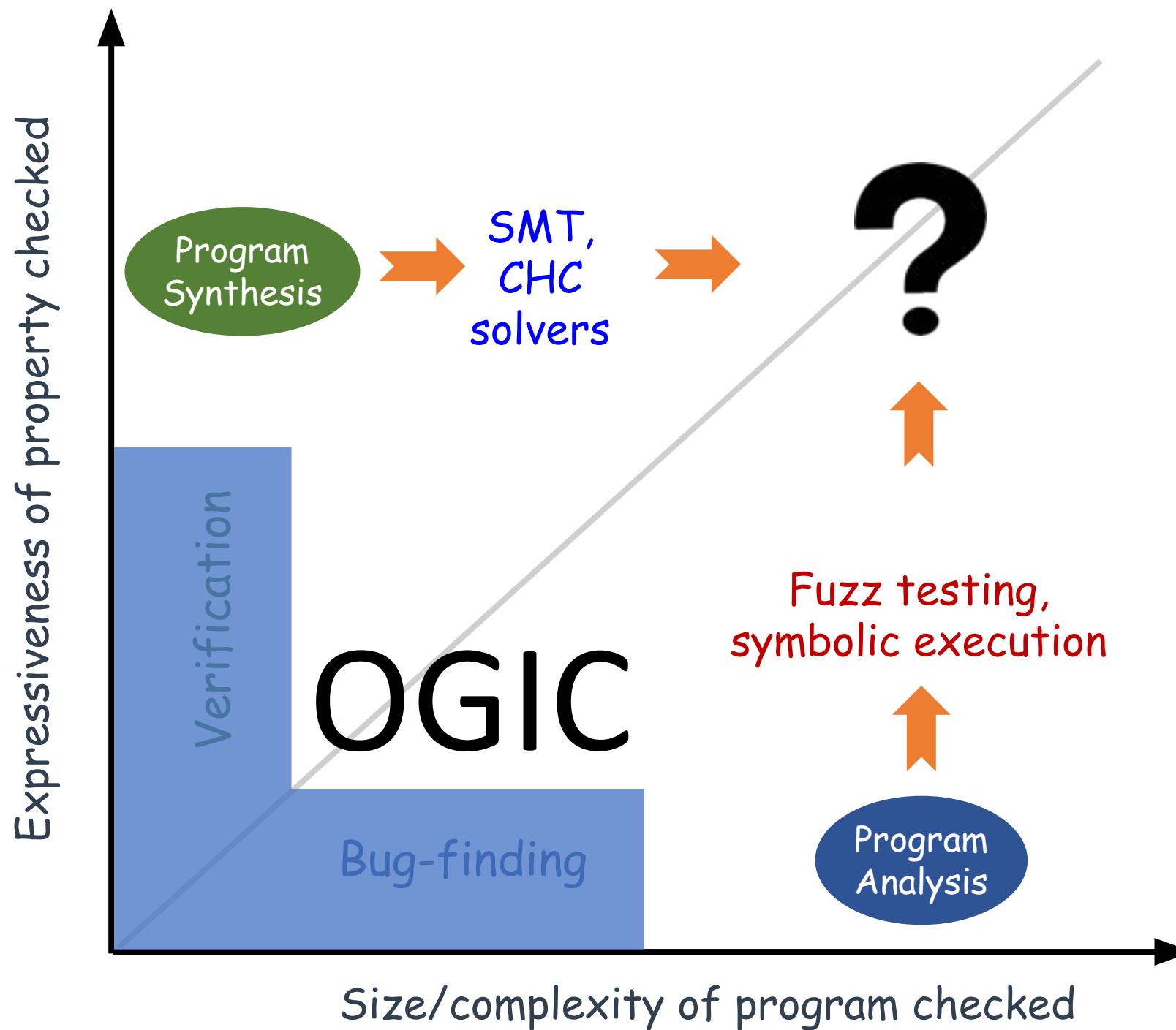
A Challenge in Verification

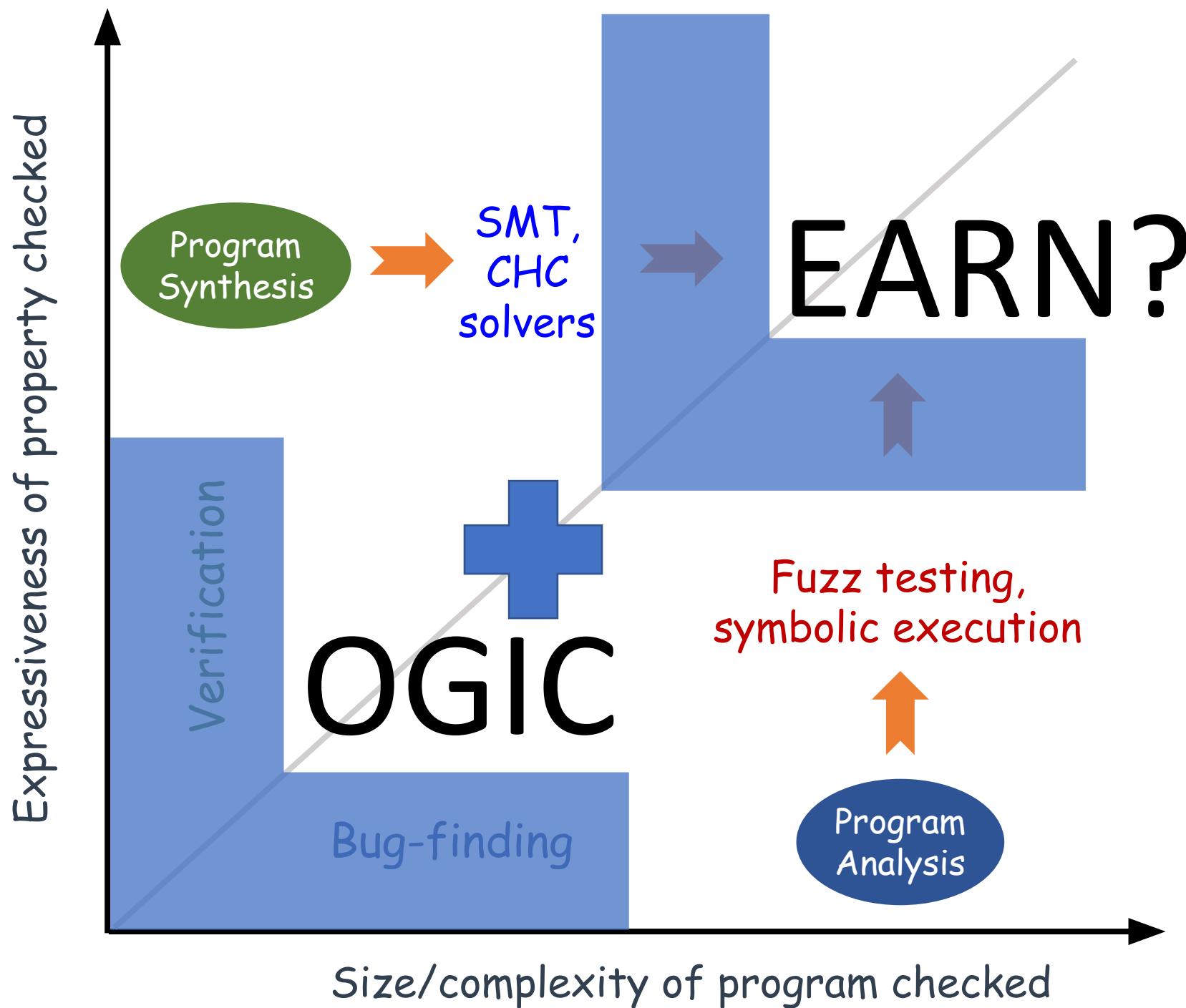
```
1 procedure main(n : int)
2 {
3     var x : int, m : int;
4
5     x := 0;
6     m := 0;
7
8     while (x < n)
9         invariant m <= 0 || (x > 0 && n > m);
10    {
11        if (*) {
12            m := x;
13        }
14        x := x + 1;
15    }
16    if (n > 0) {
17        assert(m < n);
18    }
19 }
```

Task	# LOC
sorting_bubblesort_false-unreach-call_ground	31
array_mul_init_true-unreach-call_true-termination	33
array_shadowinit_true-unreach-call_true-termination	33
nr5_true-unreach-call	44
revcpyswp2_true-unreach-call	52

(smallest 5 unsolved tasks)







Talk Outline

- Motivation

- User-Guided Program Reasoning Using Bayesian Inference. **PLDI 2018**.
- Continuously Reasoning about Programs using Differential Bayesian Inference. **PLDI 2019**.

- Learning for Bug-Finding

- Learning for Verification

- Learning Loop Invariants for Program Verification. **NeurIPS 2018**.
- Learning a Meta-Solver for Syntax-Guided Program Synthesis. **ICLR 2019**.

Talk Outline

- Motivation
- Learning for Bug-Finding
- Learning for Verification

A Static Analysis in Datalog

Analysis inputs:

next(p_1, p_2), mayAlias(p_1, p_2), guarded(p_1, p_2)

Analysis

parallel

program
immedi

Analysis

parallel

p_1 :
happ

If p_1 & p_2 may happen in parallel,
and they may access the same memory location,
and they are not guarded by the same lock,
then p_1 & p_2 may have a datarace.

parallel(p_1, p_2) :- $P_1 \leftarrow \dots, P_2 \leftarrow \dots, P_1 \wedge P_2$.

race(p_1, p_2) :- parallel(p_1, p_2), mayAlias(p_1, p_2), \neg guarded(p_1, p_2).

...

Why Datalog?

... 50 pages!

```
if (CONTEXT_SENSITIVE || THREAD_SENSITIVE || OBJECT_SENSITIVE) {
    if (HC_BITS > 0) {
        varorder = "N_F_Z_I_I2_M2_M_T1_V2xV1_V2xv1c_H2xH2c_T2_H1c";
    } else {
        //varorder = "N_F_Z_I_I2_M2_M_T1_V2xV1_V2cxV1c_B2_T2_H1";
        varorder = "N_F_I_I2_M2_M_Z_V2xV1_V2cxV1c_T1_H2_T2_H1";
    }
} else if (CARTESIAN_PRODUCT && false) {
    varorder = "N_F_Z_I_I2_M2_M_T1_V2xV1_H2_T2_H1";
    for (int i = 0; i < V1c.length; ++i) {
        varorder += "x1c"+i+"xV1c"+i;
    }
} else {
    //varorder = "N_F_Z_I_I2_M2_M_T1_V2xV1_H2_T2_H1";
    varorder = "N_F_I_I2_M2_M_Z_V2xV1_T1_H2_T2_H1";
}

System.out.println("Using variable ordering "+varorder);
int[] ordering = bdd.makeVarOrdering(reverseLocal, varorder);
bdd.setVarOrder(ordering);
```

VS.

```
### Context-sensitive inclusion-based pointer analysis using cloning
#
# Calculates the numbering based on the call graph relation.
#
# Author: John Whaley
#
#include "fielddomains.pa"
.bddnodes 10000000
.bddcache 1000000
.bddvarorder N0_F0_I0_M1_M0_V1_V0_VC1xVC0_T0_Z0_T1_H0_H1

mI0(m,i) :- mI(m,i,_).
IEnum(i,m,vc2,vc1) :- roots(m), mI0(m,i), IE0(i,m). number

V1toV2 = bdd.makePair();
V1toV2.set(V1c, V2c);
V1toV2.set(V1c, V2c);
V1c2tov2c1c.set(V1c, V2c);
V1c2tov2c1c.set(V2c, V1c);
if (OBJECT_SENSITIVE) {
    V1tob1c2tov2c1c = bdd.makePair();
    V1tob1c2tov2c1c.set(V1c, V2c);
    V1tob1c2tov2c1c.set(H1c, V2c);
}
V1tob2c2tov2c1c = bdd.makePair();
V1tob2c2tov2c1c.set(V1c, V2c);
V1tob2c2tov2c1c.set(H1c, V2c);

T2toV1 = bdd.makePair();
T2toV1.set(T1, V1);
T2toV1.set(T2, V1);
V1toV2.set(V1, V2);
V1toV2.set(V1c, V2c);
V2toV1 = bdd.makePair();
V2toV1.set(V2, V1);
V2toV1.set(V2c, V1c);
V2tob1c2tov2c1c = bdd.makePair();
V2tob1c2tov2c1c.set(V2, V1);
V2tob1c2tov2c1c.set(H2c, V1);
V1tob1c2tov2c1c = bdd.makePair();
V1tob1c2tov2c1c.set(V1, V2);
V1tob1c2tov2c1c.set(H1, H2);
V1tob1c2tov2c1c.set(H1c, H2c);
V1tob2c2tov2c1c = bdd.makePair();
V1tob2c2tov2c1c.set(V1, V2);
V1tob2c2tov2c1c.set(H1, H2);
V1tob2c2tov2c1c.set(V1c, V2c);
V1tob2c2tov2c1c.set(H1c, H2c);
V2tob2c2tov2c1c = bdd.makePair();
V2tob2c2tov2c1c.set(V2, V1);
V2tob2c2tov2c1c.set(H2, H1);
V2tob2c2tov2c1c.set(V2c, V1c);
V2tob2c2tov2c1c.set(H2c, H1c);

IE0(m,i) :- mI(m,i,_).
IE0(i,m,vc2,vc1) :- roots(m), mI0(m,i), IE0(i,m). number

cvP(_,v,h) :- vP0(v,h).
ca(_,v1,_) :- A(v1,v2).
IE0(vc2,i,vc1,m) :- IE0(i,m), IEnum(i,m,vc2,vc1).
vPfilter(v,h) :- vT(v, tv), aT(tv, th), hT(h, th).
ca(vc1,v1,vc2,v2) :- formal(m,z,v1), IEcs(vc2,i,vc1,m), actual(i,z,v2).
ca(vc2,v2,vc1,v1) :- Mret(m,v1), IEcs(vc2,i,vc1,m), Iret(i,v2).
ca(vc2,v2,vc1,v1) :- Mthr(m,v1), IEcs(vc2,i,vc1,m), Ithr(i,v2).

cvP(vc1,v1,h) :- cA(vc1,v1,vc2,v2), cvP(vc2,v2,h), vPfilter(v1,h).
hP(h1,f,h2) :- S(v1,f,v2), cvP(vc1,v1,h1), cvP(vc1,v2,h2).
cvP(vc1,v2,h2) :- L(v1,f,v2), cvP(vc1,v1,h1), hP(h1,f,h2), vPfilter(v2,h2). split
IE(i,m) :- IEcs(_,i,_,m).
```

- Fewer bugs
- Extensible
- Runs faster

Applying the Analysis to a Program

```
1 public class RequestHandler {  
2     Request request;  
3     FtpWriter writer;  
4     BufferedReader reader;  
5     Socket controlSocket;  
6     boolean isConnectionClosed;  
7     ...  
8     public Request getRequest() {  
9         return request;  
10    }  
11    public void close() {  
12        synchronized (this) {  
13            if (isClosed)  
14                return;  
15            isClosed = true;  
16        }  
17        request.clear();  
18        request = null;  
19        writer.close();  
20        writer = null;  
21        reader.close();  
22        reader = null;  
23        controlSocket.close();  
24        controlSocket = null;  
25    }  
}
```

R1

Code snippet of concurrent program **Apache FTP Server**

Applying the Analysis to a Program

```
1 public class RequestHandler {  
2     Request request;  
3     FtpWriter writer;  
4     BufferedReader reader;  
5     Socket controlSocket;  
6     boolean isConnectionClosed;  
7     ...  
  
8     public Request getRequest() {  
9         return request;  
10    }  
  
11    public void close() {  
12        synchronized (this) {  
13            if (isClosed)  
14                return;  
15            isClosed = true;  
16        }  
17        request.clear();  
18        request = null;  
19        writer.close();  
20        writer = null;  
21        reader.close();  
22        reader = null;  
23        controlSocket.close();  
24        controlSocket = null;  
25    }  
}
```

Code snippet of concurrent program **Apache FTP Server**

Applying the Analysis to a Program

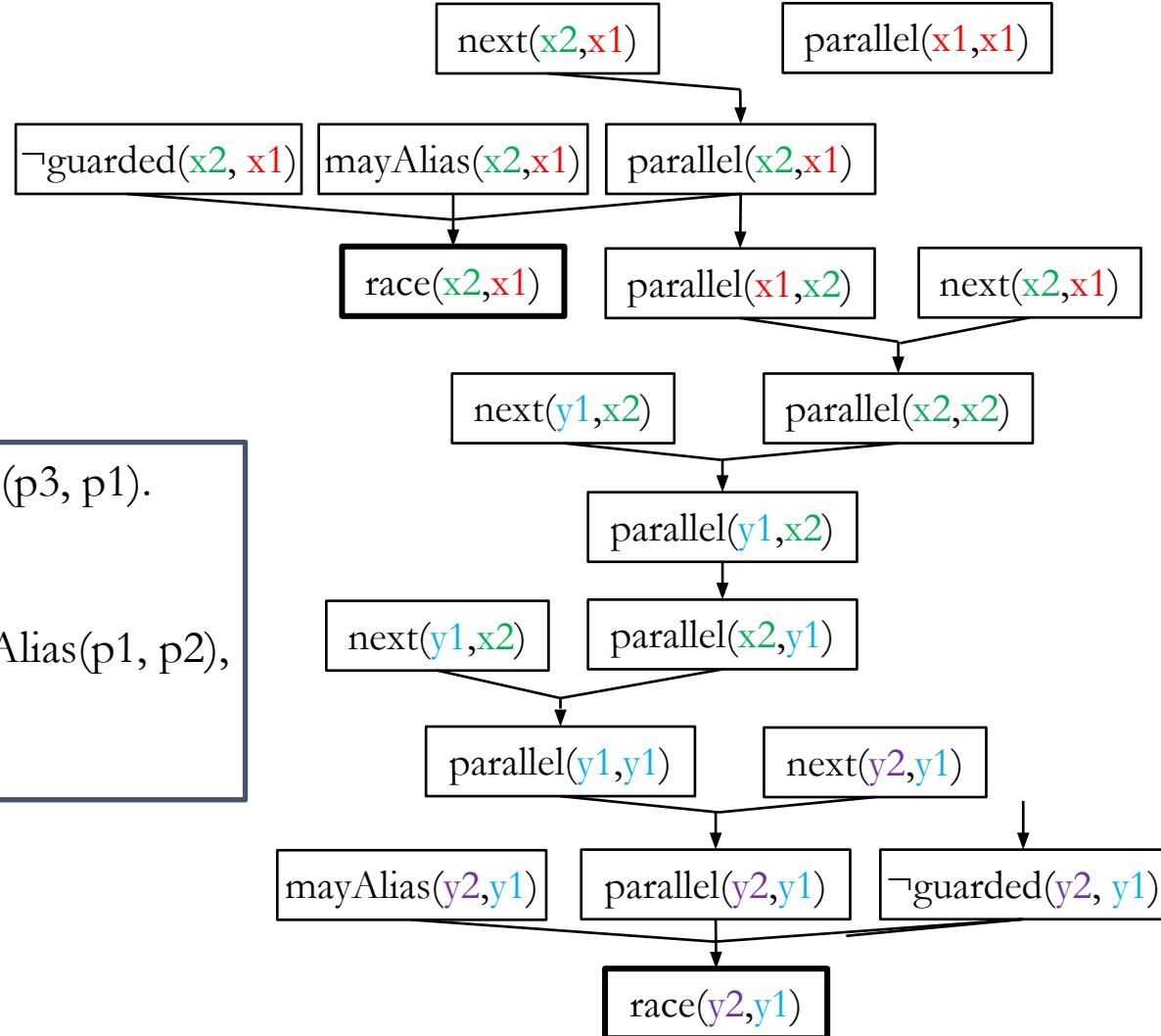
```
1 public class RequestHandler {  
2     Request request;  
3     FtpWriter writer;  
4     BufferedReader reader;  
5     Socket controlSocket;  
6     boolean isConnectionClosed;  
7     ...  
  
8     public Request getRequest() {  
9         return request;  
10    }  
  
11    public void close() {  
12        synchronized (this) {  
13            if (isClosed)  
14                return;  
15            isClosed = true;  
16        }  
17        request.clear();      // x1  
18        request = null;      // x2  
19        writer.close();       // y1  
20        writer = null;       // y2  
21        reader.close();  
22        reader = null;  
23        controlSocket.close();  
24        controlSocket = null;  
25    }
```

Code snippet of concurrent program **Apache FTP Server**

How Does Datalog Work?

```
...  
request.clear(); // x1  
request = null; // x2  
writer.close(); // y1  
writer = null; // y2  
...
```

```
parallel(p3, p2) :- parallel(p1, p2), next (p3, p1).  
parallel(p1, p2) :- parallel(p2, p1).  
race(p1, p2) :- parallel(p1, p2), mayAlias(p1, p2),  
     $\neg$ guarded(p1, p2).  
...
```



How To Go From This ...

Detected Races

R1: Race on field `org.apache.ftpserver.RequestHandler.request`

<code>org.apache.ftpserver.RequestHandler: 9</code>	<code>org.apache.ftpserver.RequestHandler: 18</code>
---	--

R2: Race on field `org.apache.ftpserver.RequestHandler.request`

<code>org.apache.ftpserver.RequestHandler: 17</code>	<code>org.apache.ftpserver.RequestHandler: 18</code>
--	--

R3: Race on field `org.apache.ftpserver.RequestHandler.writer`

<code>org.apache.ftpserver.RequestHandler: 19</code>	<code>org.apache.ftpserver.RequestHandler: 20</code>
--	--

R4: Race on field `org.apache.ftpserver.RequestHandler.reader`

<code>org.apache.ftpserver.RequestHandler: 21</code>	<code>org.apache.ftpserver.RequestHandler: 22</code>
--	--

R5: Race on field `org.apache.ftpserver.RequestHandler.controlSocket`

<code>org.apache.ftpserver.RequestHandler: 23</code>	<code>org.apache.ftpserver.RequestHandler: 24</code>
--	--

Eliminated Races

E1: Race on field `org.apache.ftpserver.RequestHandler. isClosed`

<code>org.apache.ftpserver.RequestHandler: 13</code>	<code>org.apache.ftpserver.RequestHandler: 15</code>
--	--

... To This?

Detected Races

R1: Race on field `org.apache.ftpserver.RequestHandler.request`

`org.apache.ftpserver.RequestHandler: 9`

`org.apache.ftpserver.RequestHandler: 18`

Eliminated Races

E1: Race on field `org.apache.ftpserver.RequestHandler. isClosed`

`org.apache.ftpserver.RequestHandler: 13`

`org.apache.ftpserver.RequestHandler: 15`

E2: Race on field `org.apache.ftpserver.RequestHandler.request`

`org.apache.ftpserver.RequestHandler: 17`

`org.apache.ftpserver.RequestHandler: 18`

E3: Race on field `org.apache.ftpserver.RequestHandler.writer`

`org.apache.ftpserver.RequestHandler: 19`

`org.apache.ftpserver.RequestHandler: 20`

E4: Race on field `org.apache.ftpserver.RequestHandler.reader`

`org.apache.ftpserver.RequestHandler: 21`

`org.apache.ftpserver.RequestHandler: 22`

E5: Race on field `org.apache.ftpserver.RequestHandler.controlSocket`

`org.apache.ftpserver.RequestHandler: 23`

`org.apache.ftpserver.RequestHandler: 24`

An Idea: Mixed Hard and Soft Rules

Analysis inputs:

`next(p1, p2)`, `mayAlias(p1, p2)`, `guarded(p1, p2)`

Analysis outputs:

`parallel(p1, p2)`, `race(p1, p2)`

```
if (num_threads == 1) { // p1
    x := x + 1           // p3
}
```

```
x := x + 1 // p2
```

Analysis rules:

`parallel(p3, p2) :- parallel(p1, p2), next (p3, p1).` **prob. 0.9**

`parallel(p1, p2) :- parallel(p2, p1).`

`race(p1, p2) :- parallel(p1, p2), mayAlias(p1, p2), \neg guarded(p1, p2).`

`...`

"Soft" Rule

"Hard" Rule

A Long History

```
parallel(p3, p2) :- parallel(p1, p2), next (p3, p1).
```

prob. 0.9

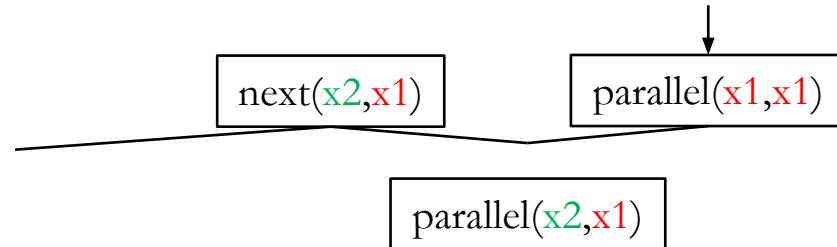
Logic

+

Probability

- 1988: Bayesian Networks [Pearl]
- 1996: Stochastic Logic Programs (SLP) [Muggleton]
- 1999: Probabilistic Relational Models (PRM) [Koller]
- 2005: Bayesian Logic (BLOG) [Milch et al.]
- 2006: Markov Logic Network (MLN) [Richardson & Domingos]
- 2007: Probabilistic Prolog (ProbLog) [De Raedt et al.]
- ...

From Derivation Trees to Bayesian Networks



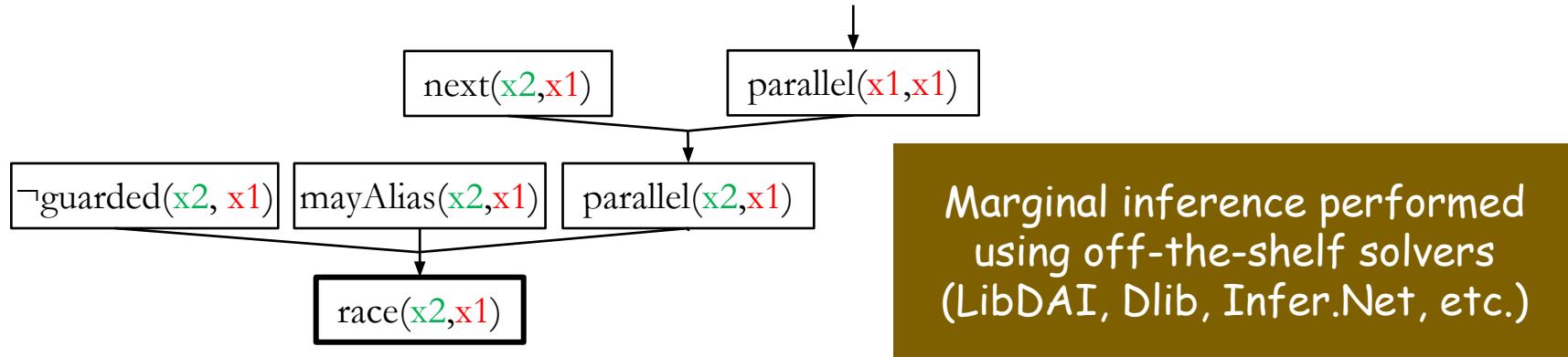
parallel(p3,p2) :- parallel(p1,p2), next(p3,p1). **prob. 0.9**

parallel(x1,x1)	next(x2,x1)	P(parallel(x2,x1) parallel(x1,x1), next(x2,x1))
-----------------	-------------	---

True	True	0.9
True	False	0
False	True	0
False	False	0

parallel(x2,x1) may only hold if
parallel(x1,x1) and next(x2,x1)
are true.

Marginal Inference in Bayesian Networks

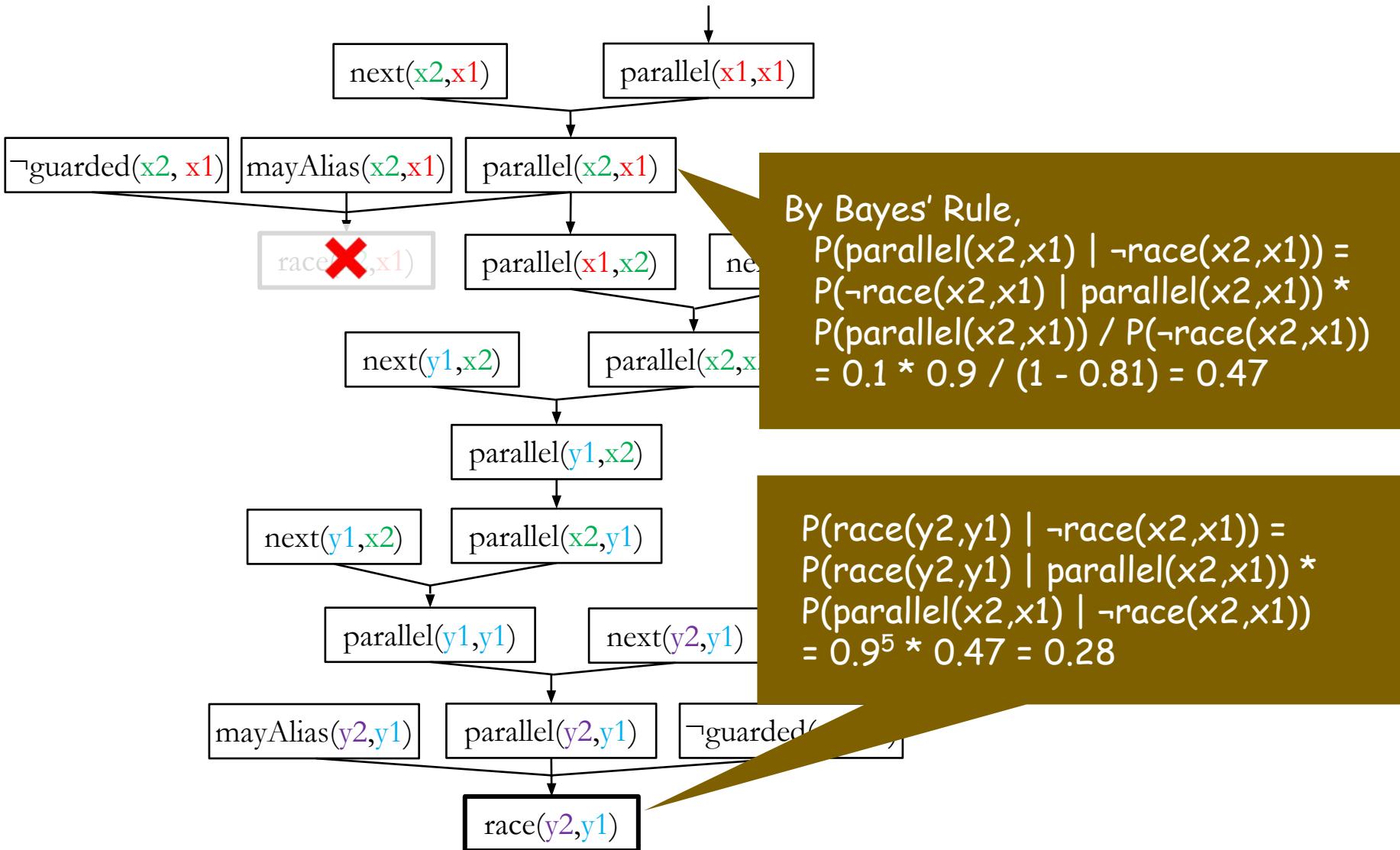


$$\begin{aligned} P(\text{race}(x_2, x_1)) &= P(\text{race}(x_2, x_1), \neg\text{guarded}(x_2, x_1), \text{mayAlias}(x_2, x_1), \\ &\quad \text{parallel}(x_2, x_1)) \\ &= P(\text{race}(x_2, x_1)) \neg\text{guarded}(x_2, x_1), \text{mayAlias}(x_2, x_1), \text{parallel}(x_2, x_1)) * \\ &+ P(\text{race}(x_2, x_1), \neg\text{guarded}(x_2, x_1), \text{mayAlias}(x_2, x_1), \neg\text{parallel}(x_2, x_1)) \\ &+ P(\text{race}(x_2, x_1), \neg\text{guarded}(x_2, x_1), \neg\text{mayAlias}(x_2, x_1), \text{parallel}(x_2, x_1)) * \\ &+ \dots \\ &+ P(\text{parallel}(x_2, x_1) | \text{next}(x_2, x_1), \text{parallel}(x_1, x_1)) * \\ &+ P(\text{parallel}(x_2, x_1), \text{guarded}(x_2, x_1), \neg\text{mayAlias}(x_2, x_1), \\ &\quad \neg\text{parallel}(x_2, x_1)) \end{aligned}$$

0. If any of the antecedents fail, then the race cannot happen.

Confidence	Detected Races	
0.81	R2: Race on field org.apache.ftpserver.RequestHandler.request org.apache.ftpserver.RequestHandler:17	org.apache.ftpserver.RequestHandler:18
0.53	R3: Race on field org.apache.ftpserver.RequestHandler.writer org.apache.ftpserver.RequestHandler:19	org.apache.ftpserver.RequestHandler:20
0.35	R4: Race on field org.apache.ftpserver.RequestHandler.reader org.apache.ftpserver.RequestHandler:21	org.apache.ftpserver.RequestHandler:22
0.30	R1: Race on field org.apache.ftpserver.RequestHandler.request org.apache.ftpserver.RequestHandler:9	org.apache.ftpserver.RequestHandler:18
0.23	R5: Race on field org.apache.ftpserver.RequestHandler.controlSocket org.apache.ftpserver.RequestHandler:23	org.apache.ftpserver.RequestHandler:24





Confidence	Detected Races	
0.81 0	R2: Race on field org.apache.ftpserver.RequestHandler.request org.apache.ftpserver.RequestHandler:17	org.apache.ftpserver.RequestHandler:18
0.53 0.28	R3: Race on field org.apache.ftpserver.RequestHandler.writer org.apache.ftpserver.RequestHandler:19	org.apache.ftpserver.RequestHandler:20
0.35 0.18	R4: Race on field org.apache.ftpserver.RequestHandler.reader org.apache.ftpserver.RequestHandler:21	org.apache.ftpserver.RequestHandler:22
0.30 0.30	R1: Race on field org.apache.ftpserver.RequestHandler.request org.apache.ftpserver.RequestHandler:9	org.apache.ftpserver.RequestHandler:18
0.23 0.12	R5: Race on field org.apache.ftpserver.RequestHandler.controlSocket org.apache.ftpserver.RequestHandler:23	org.apache.ftpserver.RequestHandler:24

$P(R_i \mid \neg R_2)$

Confidence	Detected Races	
0	R2: Race on field <code>org.apache.ftpserver.RequestHandler.request</code>	
	<code>org.apache.ftpserver.RequestHandler:17</code>	<code>org.apache.ftpserver.RequestHandler:18</code>
0.28	R3: Race on field <code>org.apache.ftpserver.RequestHandler.writer</code>	
	<code>org.apache.ftpserver.RequestHandler:19</code>	<code>org.apache.ftpserver.RequestHandler:20</code>
0.18	R4: Race on field <code>org.apache.ftpserver.RequestHandler.reader</code>	
	<code>org.apache.ftpserver.RequestHandler:21</code>	<code>org.apache.ftpserver.RequestHandler:22</code>
0.30	R1: Race on field <code>org.apache.ftpserver.RequestHandler.request</code>	
	<code>org.apache.ftpserver.RequestHandler:9</code>	<code>org.apache.ftpserver.RequestHandler:18</code>
0.12	R5: Race on field <code>org.apache.ftpserver.RequestHandler.controlSocket</code>	
	<code>org.apache.ftpserver.RequestHandler:23</code>	<code>org.apache.ftpserver.RequestHandler:24</code>

$P(R_i \mid \neg R_2)$

Experimental Setup

- **Analyses:**

58 input relations
44 output relations
102 rules

Race conditions checker

52 input relations
25 output relations
62 rules

Information flow checker

- **Programs:**

Concurrent Java programs

(~ 50-550 KB in size)

Symantec Android apps

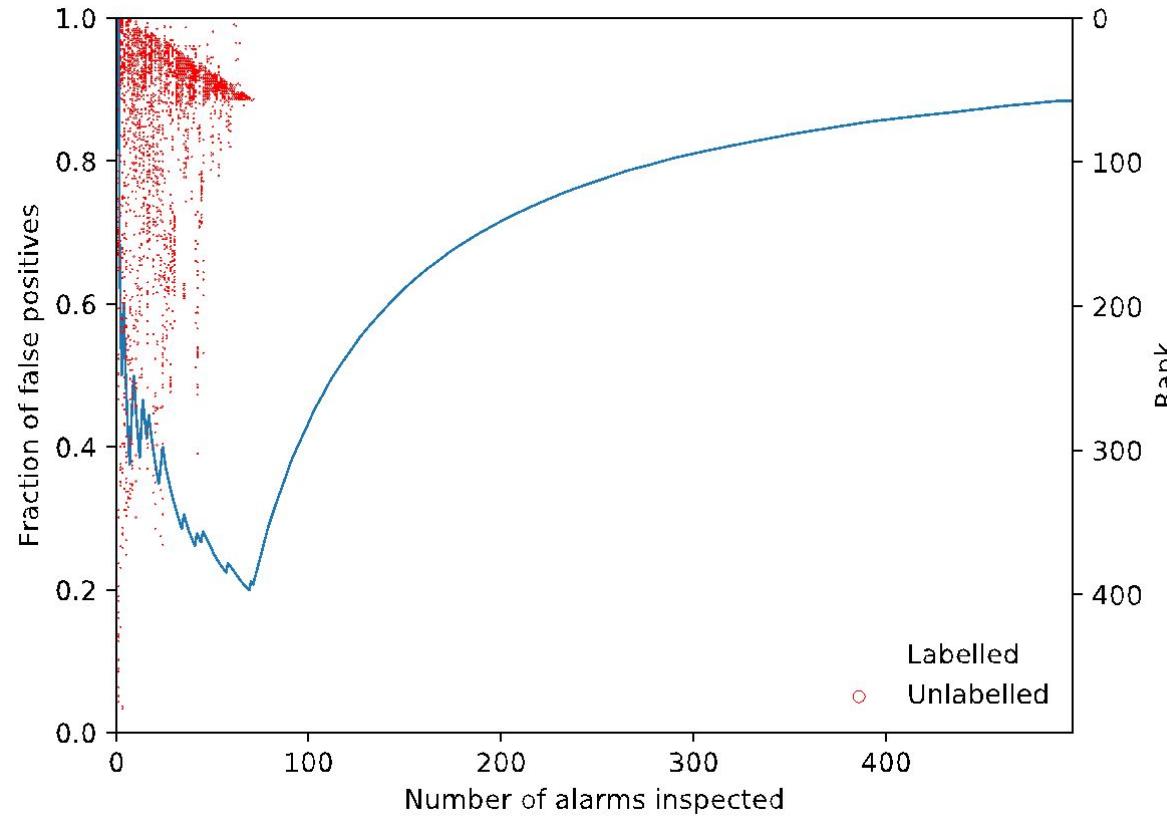
(~ 68-81 KB in size)

Empirical Results

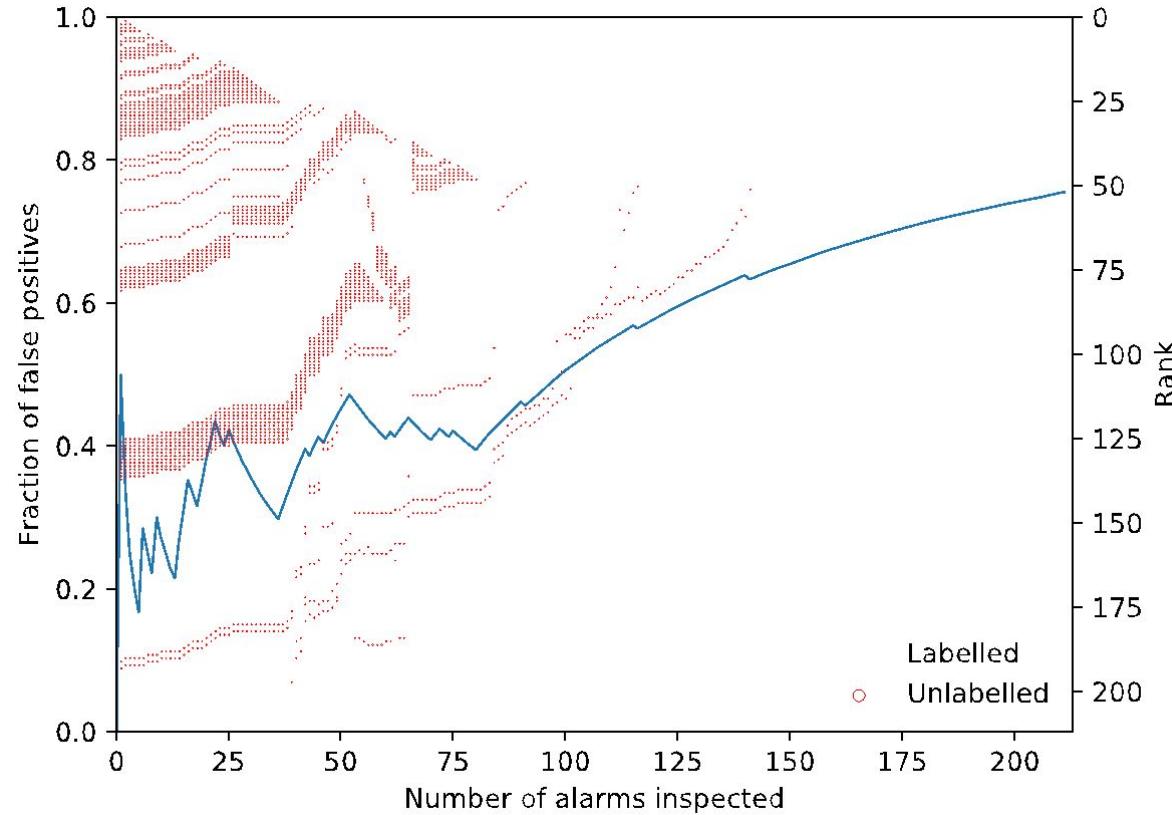
.90 - 1	=	excellent	(A)
.80 - .90	=	good	(B)
.70 - .80	=	fair	(C)
.60 - .70	=	poor	(D)
.50 - .60	=	fail	(F)

	Graph size		Alarms		FP rate	AUC
	Tuples	Clauses	Total	Bugs		
Race conditions checker						
weblech	2.5K	1.5K	188	55	71%	0.88
hedc	12K	10K	152	9	94%	0.71
jspider	45K	45K	257	7	97%	0.87
ftpserver	110K	112K	522	75	86%	0.97
Information flow checker						
AndorsTrail	2.7K	3.2K	156	7	96%	0.99
kQm-LO	12K	18K	817	160	81%	0.94
gingermaster	15K	20K	437	87	80%	0.88
iNJ-Cw	17K	24K	1,012	248	76%	0.91

Ranking Quality: Race Conditions Checker



Ranking Quality: Information Flow Checker



Taxonomy of Research Directions

Balancing Analysis Tradeoffs

- Analysis Accuracy vs. Soundness
- Analysis Accuracy vs. Cost

Tailoring Analysis Results

- Unguided vs. Interactive
- Batch vs. Continuous Reasoning
- Alarm Clustering vs. Ranking

Analysis Specification and Implementation

- Synthesizing Analyses from Data
- Expressiveness of Analysis Language
- Capabilities of Analysis Solvers

More details in SAS 2019 paper “Rethinking Static Analysis by Combining Discrete and Continuous Reasoning”

Talk Outline

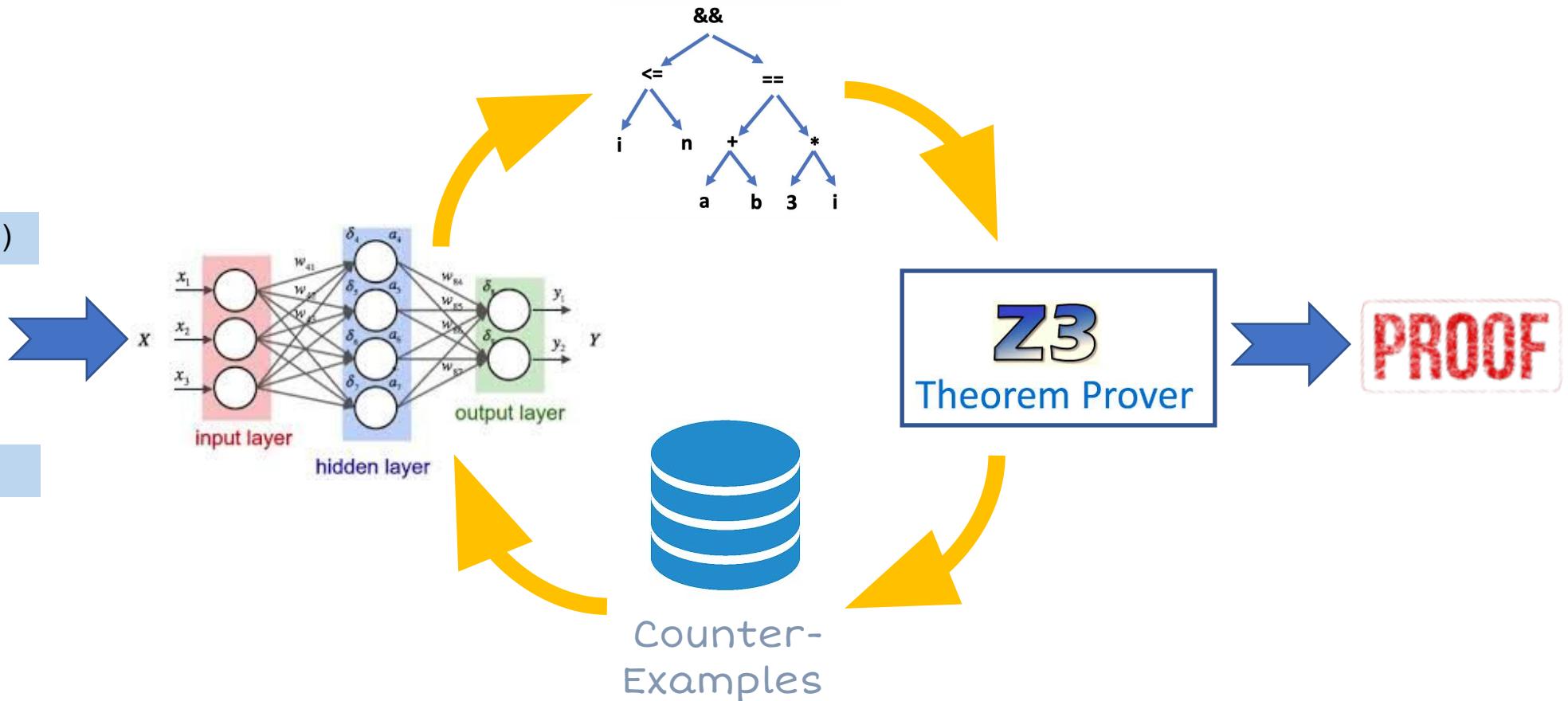
- Motivation
- Learning for Bug-Finding
- Learning for Verification

Example: Loop Invariants

```
void main(int n) {  
    int x = 0; int m = 0;  
  
    human expert:  
    (m == 0 || m < x) && (n <= 0 || x <= n)  
  
generated:  
    (m <= 0 || x > 0) && (m <= 0 || n > m)  
    while (x < n) {  
        if (*) { m = x; }  
        x = x + 1;  
    }  
    if (n > 0) assert(m < n);  
}
```

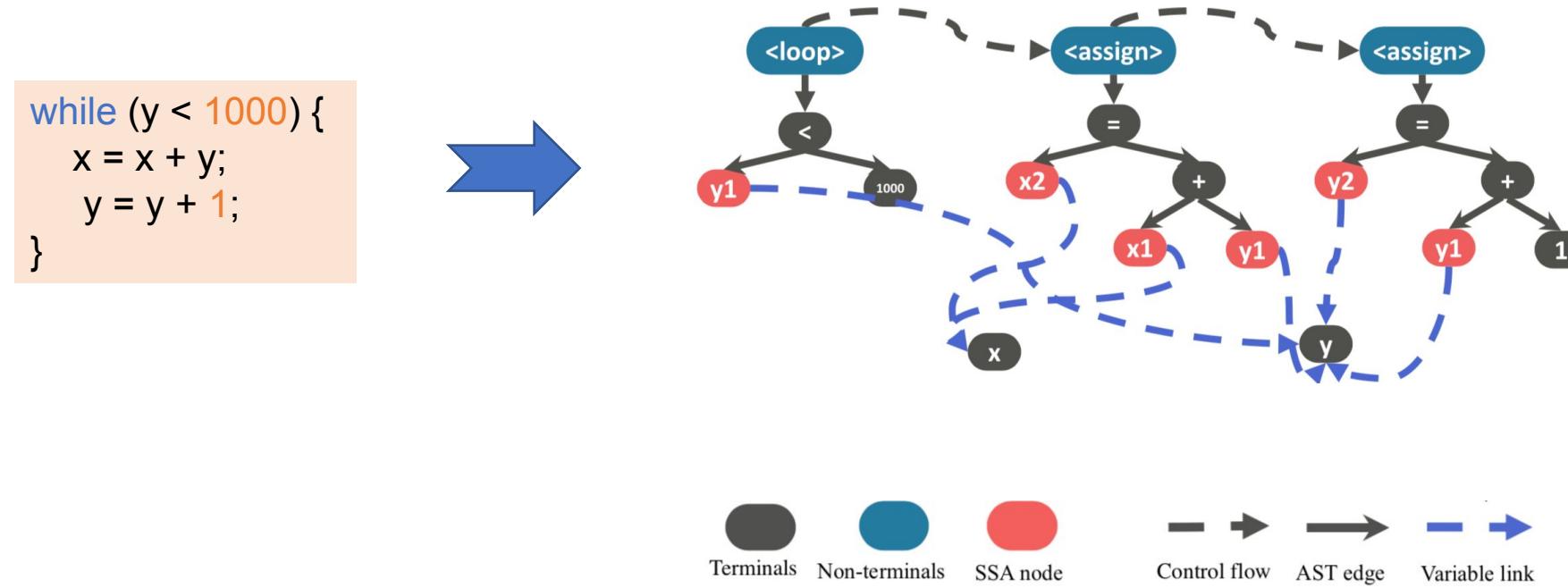
Architecture of code2inv

```
assume (a == 1 ∧ b == 1)
while (b < 1000) {
    a = a + b;
    b = b + 1;
}
assert (a >= 1600)
```



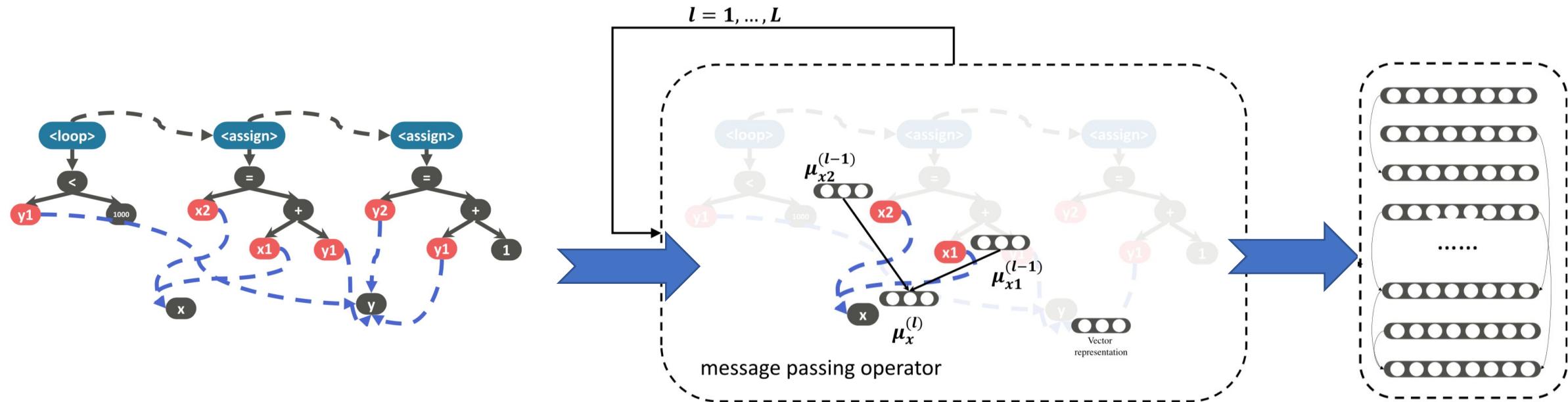
Step 1: Representing Program as Graph

- Encode the program as a graph that captures its rich structure



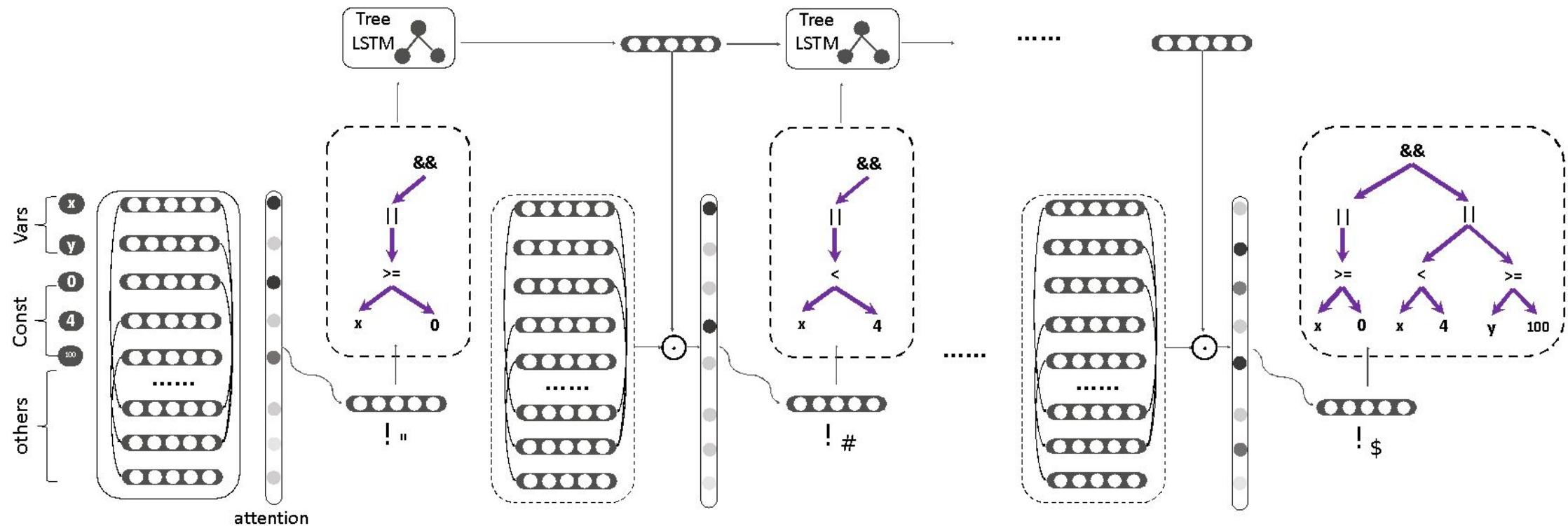
Step 2: Converting Graph to Vector

- Convert the graph to a vector representation using a [graph neural network](#)



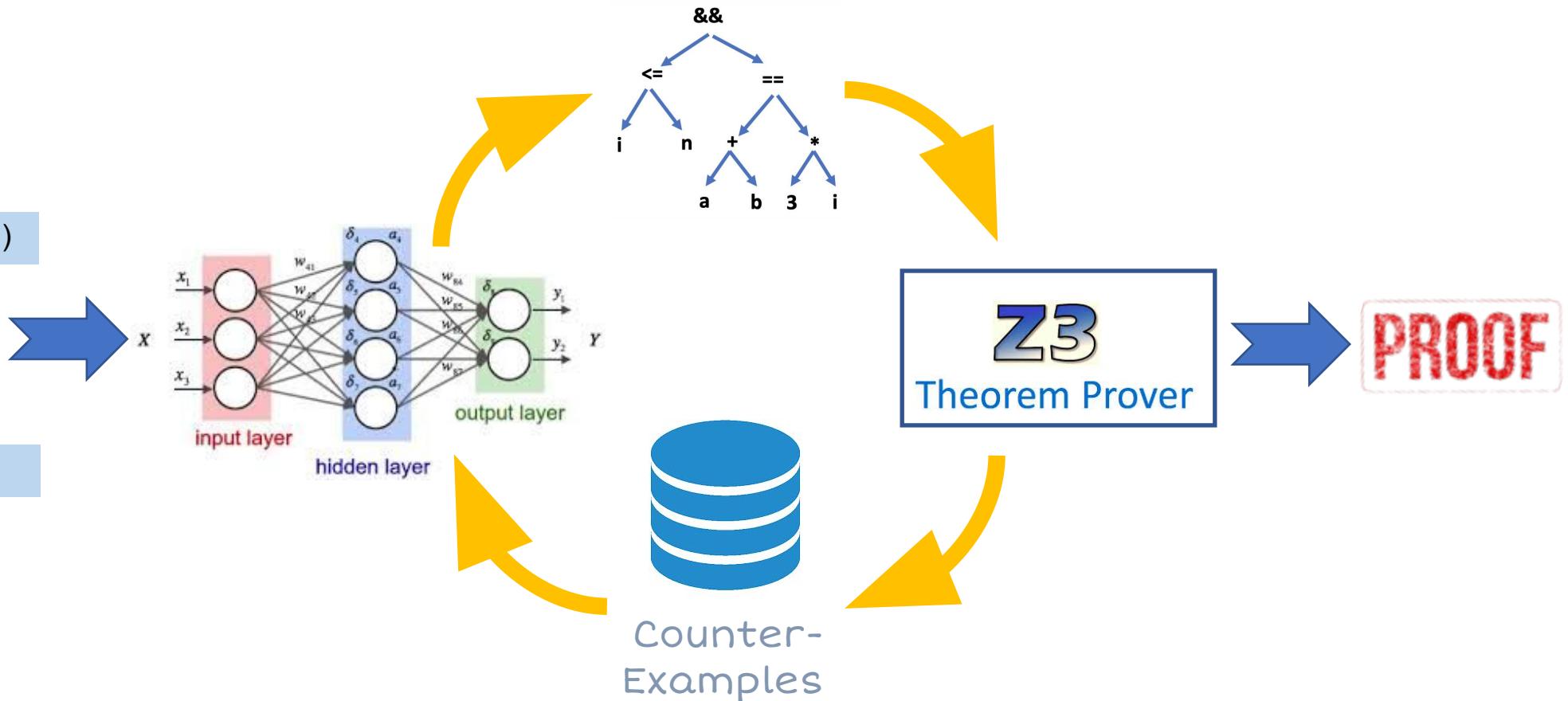
Step 3: Predicting Loop Invariant

- Model loop invariant generation as a **multi-step decision making** process

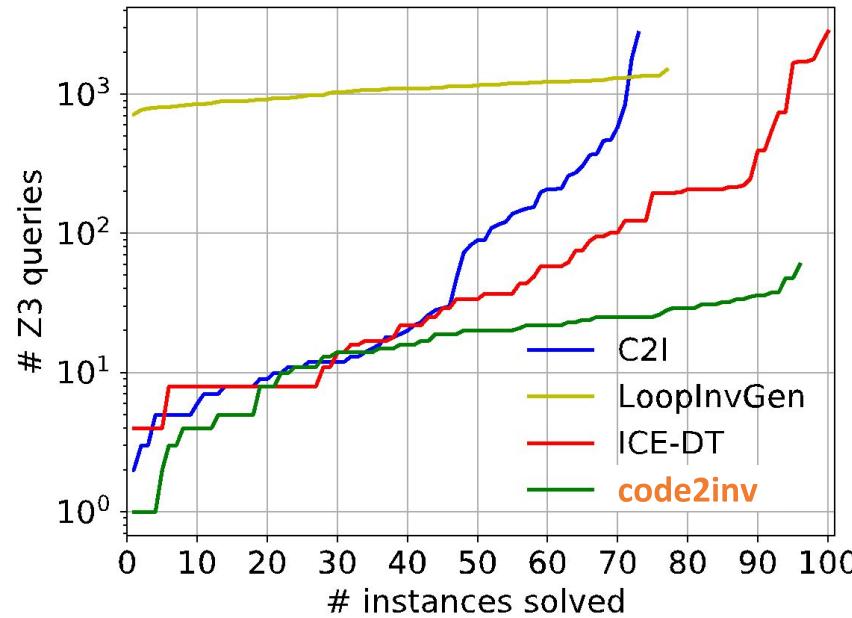


Architecture of code2inv

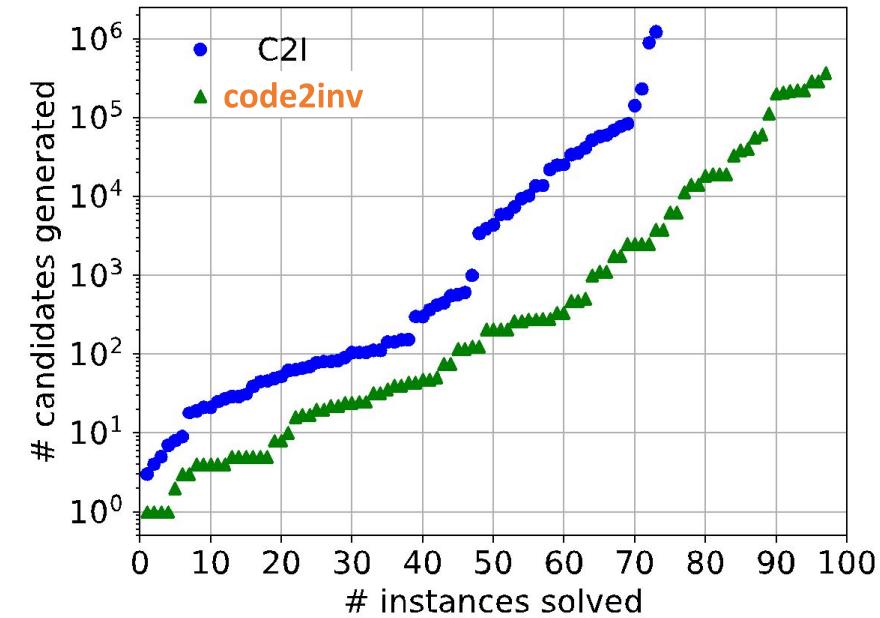
```
assume (a == 1 ∧ b == 1)
while (b < 1000) {
    a = a + b;
    b = b + 1;
}
assert (a >= 1600)
```



Comparison to State-of-the-Art



Verification Cost



Sample Complexity

Other Applications

- Verification: Program \rightarrow Invariant
- Bug-Finding: Program \rightarrow Counterexample
- Repair: Program \rightarrow Edit Sequence

Example: Bug Detection and Repair for JS

```
2 helpers.js
@@ -5,7 +5,7 @@ exports.parseTitle = title => {
 5     if (matches) {
 6         return {
 7             episode: Number(matches.groups.episode),
 8 -             hosts: matches.groups.hosts.split(/[,&]+|\sand\s/).map(el => S(el).trim().s)
 9         };
10     }
11     return false;
12
13     if (matches) {
14         return {
15             episode: Number(matches.groups.episode),
16             hosts: matches.groups.hosts.split(/[,&]+|\sand\s/).map(el => S(el).trim().s)
17         };
18     }
19     return false;
20
21     const S = require ('string');
22
23     exports.parseTitle = title => {
24         const matches = title.match (/ Joe Rogan Experience # (? <episode> \ d *) (\ \ s? [-] {0,} \ s?) (? <hosts>. *)? (? <part ?> \ s \ (share \ s \ d {1} \)) $ / i);
25         if (matches) {
26             return {
27                 episode: Number (matches.groups.episode),
28                 hosts: matches.groups.hosts.split (/ [, &] + | \ sand \ s /). map (el => S (el) .trim () .s)
29             };
30         }
31         return false;
32     };
33 }
```

Intended Goal:

Split a string based on delimiter that matches regex: [, &] + | \sand\s

Input: " and "

Output of buggy code:

[' ', ' and ', ' ']

Output of fixed code:

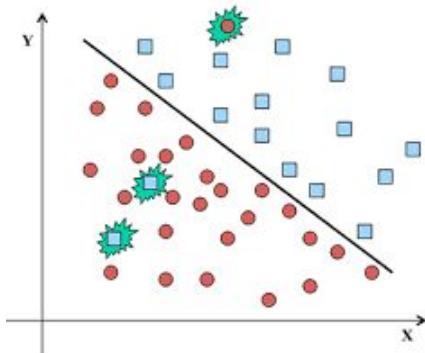
[' ', ' ']

Limits of Training Data



Stats of Data Crawled from Github per Week:

Downloaded JS files:	9,425,472
Valid AST files and diffs:	4,712,736
ASTs with a single diff:	585,984
Valid data points:	47,040



Sampling 50 data points in test set reveals
21 real bugs and 29 non-bugs

Conclusions

- Logical Reasoning: Discrete → Continuous
- Which machine learning models worked?
 - Bug-finding: Bayesian networks, MLNs, SLPs, ...
 - Relies on good human-engineered features
 - Verification: graph neural network
 - Suitable program representation is critical
- Challenge: How to obtain training data?
 - Bug-finding: supervised learning
 - Leverage continuously growing open-source datasets: OSS-Fuzz, GHArchive, ...
 - Verification: reinforcement learning
 - Leverage formal methods tools evolved over decades: SMT/CHC solvers, Coq, ...