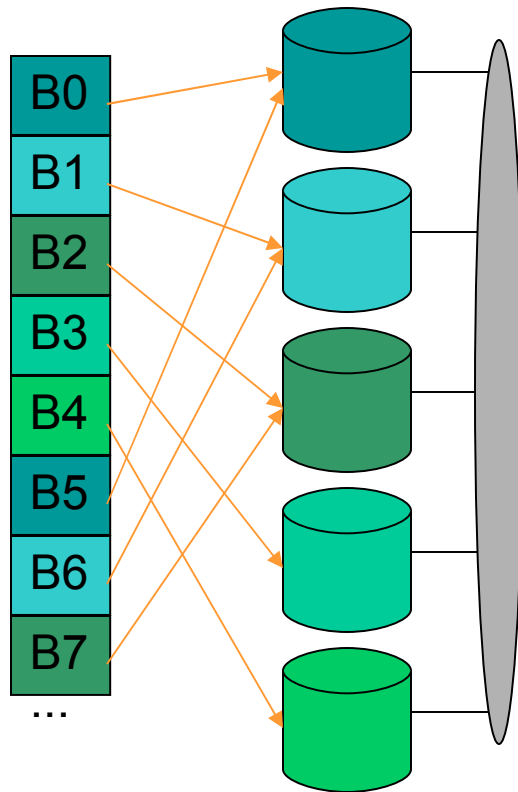


A Case for Heterogeneous Disk Arrays

Toni Cortes and Jesús Labarta
Departament d'Arquitectura de Computadors
Univeritat Politècnica de Catalunya - Barcelona

Disk Arrays (RAIDs)



n Group several disks

- | Single address space
- | High capacity
- | Improved performance
- | *Low cost*

n Heterogeneous RAID

- | Not all disks are equal

n Heterogeneous disk arrays are becoming a common configuration

- | Replacing a new disk
- | Adding new disks

n Current solution

- | All disks are treated as equal
 - ü No performance gain is obtained
 - ü No capacity gain is obtained

n **AdaptRaid0**

- | Block-distribution policy
- | Take advantage of the *goodies* of each disk

n **Target Environment**

- | Scientific and general purpose
- | Not multimedia
 - ü Solutions have already been presented
 - ü Very dependent on some characteristics
- | Disk arrays level 0 (*RAID0*)
 - ü Level 5 is under development

n Multimedia Systems

- | Random distribution with replication (*Santos98*)
- | Policy based on logical disks (*Zimmerman98*)
- | Use fast disk for hot data (*Dan95*)
- | **Differences:**
 - ü Large blocks, only reads, and sustained bandwidth

n General purpose

- | HP AutoRaid (*Wilkes95*)
- | Disc-Cache Disk (*Hu98*)
- | **Differences:**
 - ü Do not adapt to the existent hardware

Disk Arrays and Parallelism

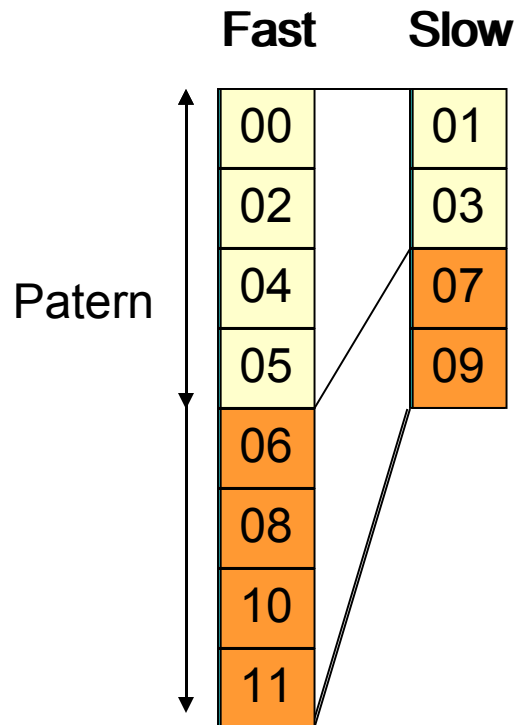
n Parallelism within a request

- | Requests have to be large
 - ü The sub-request of each disk has to be large
 - ü Seek + search + transfer in all disks

n Parallelism between requests

- | The number of disks has to be large
 - ü Compared to the average number of disks used in a request

AdaptRaid0: An Example



n Basic idea

- | Load each disk depending on its characteristics

n Example

- | 1 fast disk
 - ü Size = S
 - ü Performance = P
- | 1 slow disk
 - ü Size = S/2
 - ü Performance = P/2

AdaptRaid0: The Parameters

n Utilization factor (UF)

- | One factor per disk
 - ü Larger disks have more blocks?
 - ü Faster disks have more blocks?

n Lines in pattern (LIP)

- | We define a pattern using the UF
 - ü Large patterns allow more requests with good disks
 - ü Small patterns allow a better distribution

AdaptRaid0: The Algorithm

n Algorithm

- | Decide LIP and Ufd
- | Compute number of blocks per disk in the pattern
 - ü $\text{Blocksd} = \text{int}(\text{UFd} * \text{LIP})$
- | Distribute blocks in a round-robin way
 - ü Use the available disks
 - ü A disk becomes unavailable when Blocksd have already been placed in it
- | Repeat step 3 until one disk becomes full

n Parameters

- | UF based on the size of the disk
- | Lines in pattern
 - ü 100 lines for 8-disk arrays
 - ü 10 lines for 32-disk arrays

n Simulation

- | Simulator: HRaid (*Cortes99*)
- | Workload from HP labs (*1999*)

n Reference systems

- | Raid0 and OnlyFast

n Disks

- | Fast disk
 - ü Seagate Barracuda 4LP (4.339 Gbytes)
- | Slow disk
 - ü Seagate Cheetah 4LP (2.061 Gbytes)

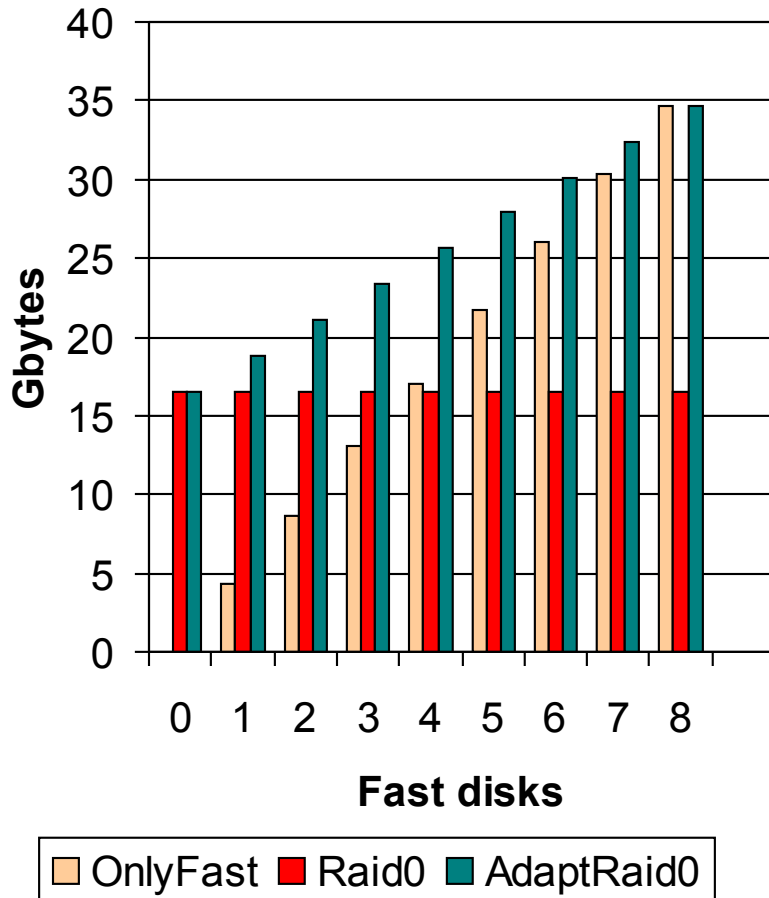
n Bus

- ü 10us latency
- ü 100Mbit/s bandwidth

n File system

- ü 10 requests in parallel

Capacity Evaluation



n **Raid0**

| Constant capacity

ü Small

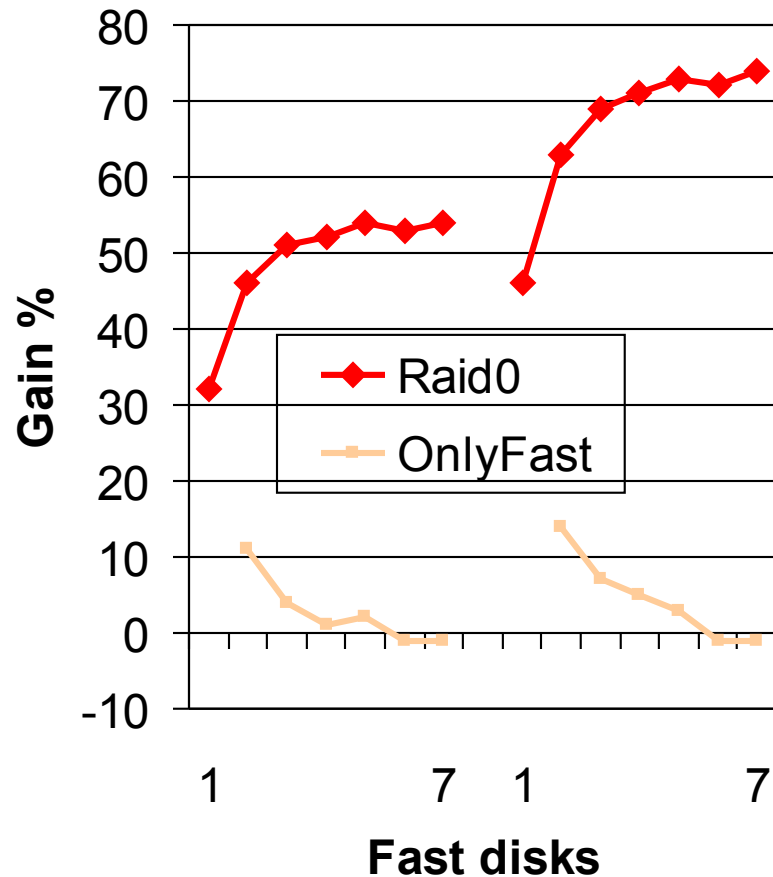
n **OnlyFast**

| Small capacity with few disks

n **AdaptRaid0**

| Offers the best size

Performance Evaluation (8 disks)



n Raid0

| Does not use

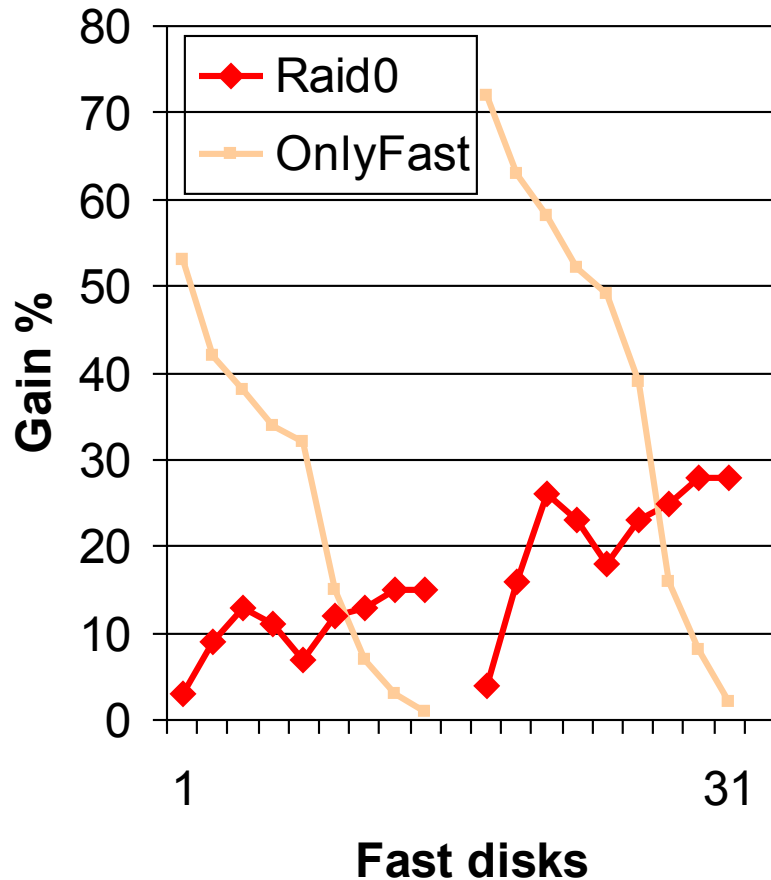
ü Characteristics of good disks

n OnlyFast

| Does not use

ü Parallelism between requests

Performance Evaluation (32 disks)



n Raid0

- | Does not use
 - ü Characteristics of good disks
- | It uses
 - ü Parallelism between requests

n OnlyFast

- | Does not use
 - ü Parallelism between requests

n AdaptRaid0

- | Performance**

- ü It knows how to use the disks

- ü Allows parallelism

- | Size**

- ü It uses all the available capacity

n Solve the same problem for Raid5

- | Problem of parity blocks
- | Less scalable
 - ü No parallelism among requests