

1. Error Analysis

Definition:

let x is a value, \tilde{x} is a estimated value

(1) absolute error, $E_a = |x - \tilde{x}|$

(2) relation error, $E_r = \left| \frac{x - \tilde{x}}{x} \right|$

(3) percentage error, $E_p = 100 \times \left| \frac{x - \tilde{x}}{x} \right|$

$\exists \epsilon > 0, |x - \tilde{x}| < \epsilon$, Then ϵ is upper limit of the absolute error measures the absolute accuracy.

1.1. Error in Implementation of Numerical Methods.

- (1) Round-off Error
- (2) Overflow & Underflow
- (3) Floating Point Arithmetic and Error Propagation
- (4) Truncation Error
- (5) Machine eps (Epsilon)

(3) Floating Point Arithmetic and Error Propagation.

Let x_1, x_2 are values, E_1, E_2 are error of x_1, x_2 , We want to check the change of error in
" + ", " - ", " * ", " / "
" + "

Let $x = x_1 + x_2$, error of x is E

Then $x + E = x_1 + x_2 + E_1 + E_2 \implies E = E_1 + E_2$

by triangle inequality

Absolute Error = $|E| \leq |E_1| + |E_2|$

Relative Error = $\frac{|E|}{|x|} \leq \frac{|E_1|}{|x|} + \frac{|E_2|}{|x|}$

" - " (Similar " + ")

”*”

Let $x = x_1 * x_2$

Then $x + E = (x_1 + E_1)(x_2 + E_2) = x_1x_2 + E_2x_1 + E_1x_2 + E_1E_2$

Absolute Error = $|E| \leq |x_2E_1| + |x_1E_2|$

Relative Error = $\frac{|E_1|}{|x|} \leq \frac{|E_1|}{|x_1|} + \frac{|E_2|}{|x_2|}$

”/”

Let $x = x_1/x_2$

$x + E_x = \frac{x_1 + E_1}{x_2 + E_2} \left(\frac{x_2 - E_2}{x_2 - E_2} \right) = \frac{x_1x_2 + E_1x_2 - x_1E_2}{x_2^2 - E_2^2} + E_1E_2$

Absolute Error = $|E_x| = \left| \frac{E_1x_2 - x_1E_2}{x_2^2} \right| \leq \frac{|E_1|}{|x_2|} + \frac{|x_1E_2|}{x_2^2}$

Relative Error = $\frac{|E_x|}{|x|} \leq \frac{|E_1|}{|x_1|} + \frac{|E_2|}{|x_2|}$

(4) Truncation Error. Cause by approximation infinite with its finite terms.

Use Taylor series ($f(x) \in P(C)$) as example

Let $x = a, f(x) = f(a) + f'(a)(x - a) + f''(a)\frac{(x - a)^2}{2!} + \dots + \frac{(x - a)^n}{n!}f^n(a) + \dots + R_n$

$R_n = \int_a^x \frac{(x - t)^n}{n!} f^{(n+1)}(t) dt$

Thm 1(First Mean Value Theorem)

If g is continuous on $[a, x]$, then $\exists \xi$ between a and x s.t.

$$\int_a^x g(t) dt = g(\xi)(x - a)$$

Thm 2(Second Mean Value Theorem)

If g, h is differentiable and integrable on $[a, x]$, h does not change sign on $[a, x]$ then $\exists \xi$ that $a \leq \xi \leq x$ s.t.

$$\int_a^x g(t)h(t) dt = g(\xi) \int_a^x h(t) dt$$

since $t \in [a, x], h(t) = (x - t)^n \frac{1}{n!}, f^{(n+1)}(t)$ is continuous

$\exists \xi \in [a, x], R_n = \frac{f^{(n+1)}(\xi)}{(n+1)!} f^{(x+1)}(\xi), \xi \in [a, a + h]$

(Ref. Violin page:799)

since power series convergent, $R_n(x) \rightarrow 0, as_n \rightarrow \infty$

Definition

Given $\{a_n\} \{b_n\}$, $b_n \geq 0$, $\forall n \geq 1$
 $a_n = O(b_n)$ if $\exists M > 0 \rightarrow |a_n| \leq Mb_n \forall n \geq 1$
 $R_n(x) = O(h^{n+1})$

1.2. Condition & Stability.

Condition number is sensitivity of the function

Stability is used to describe the sensitivity of the process

Condition number of the $f(x)$

$$\text{CN} = \frac{\left| \frac{f(x) - f(\tilde{x})}{x - \tilde{x}} \right|}{\left| \frac{x - \tilde{x}}{x} \right|} = \left| \frac{f(x) - f(\tilde{x})}{x - \tilde{x}} \right| \cdot \left| \frac{x}{f(x)} \right| = \left| \frac{x}{f(x)} \cdot f'(x) \right|$$

by Mean Value Theorem,

$$\frac{f(x) - f(\tilde{x})}{x - \tilde{x}} \approx f'(x)$$

when $\text{CN} \leq 1$ is **well condition**, other is **ill condition**

when the function is more sensitive to change, the condition number will be more big.

2. Methods for $f(x) = 0$

we have four way to deal this problem

- (1) Direct analytical Method
- (2) Graphical
- (3) Trial and Error Method
- (4) Iterative Method

Thm. 3(Mean Value Theorem)

Let f be a continuous function on $[a, b] = I$ (connected),
if $f(a) \leq c \leq f(b)$ that $\exists \xi \in [a, b] \rightarrow f(\xi) = c$

Corollary

Let f be a continuous function on $[a, b] = I$ (connected)
i.e. $f(a) \cdot f(b) < 0 \Rightarrow \exists c \in (a, b) \Rightarrow f(c) = 0$
 c is a root of $f(t)$

Iterative Method

2.1. Bisection Method.

Let a, b be fixed satisfying Thm.3

$\therefore f(a) \cdot f(b) < 0, f$ is continuous on $[a, b]$. The first approximation is $x_0 = \frac{a+b}{2}$
if $f(a) \cdot f(x_0) \leq 0$, then By Thm. 3 the root will lie on (a, x_0) and $x_1 = \frac{a+x_0}{2}$
continue the process, let $x_{n-3}, x_{n-2}, x_{n-1}$ be same step, then nth approximation
if $f(x_{n-1}) \cdot f(x_{n-3}) \leq 0$, then $x_n = \frac{x_{n-1} + x_{n-3}}{2}$
else $f(x_{n-1}) \cdot f(x_{n-3}) \geq 0$, then $x_n = \frac{x_{n-1} + x_{n-2}}{2}$
we shall label the interval by algorithm

$$[a, b] = [a_0, b_0], [a_1, b_1], [a_2, b_2], \dots$$

by construction $b_n a_n = \frac{1}{2}(b_{n-1} - a_{n-1})$, Hence $b_n - a_n = \frac{1}{2^n}[b_0 - a_0], \forall n \geq 1$

Clearly $a_0 \leq a_1 \leq \dots \leq b, b_0 \geq b_1 \geq \dots \geq a, \{a_n\}, \{b_n\}$ is bdd and monotonic

$$\lim_{n \rightarrow \infty} a_n = \lim_{n \rightarrow \infty} b_n = f(r)$$

by assumption $f(a_n)f(b_n) < 0, \lim_{n \rightarrow \infty} f(a_n) = f(\lim_{n \rightarrow \infty} a_n) = f(r)$

$\therefore f(b_n) = f(r), 0 \leq [f(r)]^2 \leq 0 \implies f(r) = 0$

The process is called **nested internal property**

Let $\{C_k\}_{k=1}^\infty$ is a \downarrow sequence of nonempty closed compact subset of X , then $\cap_k C_k \neq \emptyset$ if $C_k \rightarrow 0$, then $\cap_k C_k = \{r\}$

Let ξ be the solution $f(x) = 0$, then $\{x_0 - \xi\} \leq \frac{b-a}{2}, \dots, \{x_n - \xi\} \leq \frac{b-a}{2^{n+1}}$

Definition(p-order-convergence)

$\{x_n\} : \text{seq}, x_n \rightarrow z, s_n \rightarrow \infty$, define $\epsilon_n = z - x_n$, if $\exists c > 0, p \geq 1$

$$\lim_{n \rightarrow \infty} \frac{|\epsilon_{n+1}|}{|\epsilon_n|^p} = c$$

we call $\{x_n\}$ is p order convergence

if $c \leq 1$, then it's good(only check this when it's a first order convergence)

Let ϵ_n be the error i.e. $\epsilon_n = |x_n - \xi|$, $\epsilon_n \leq \frac{b-a}{2^{n+1}} \leq \epsilon$, i.e. $h \geq \frac{\ln(b-a) - \ln \epsilon}{\ln 2} - 1$

$$\epsilon_n = |x_n - \xi| \leq \frac{1}{2} \left(\frac{b-a}{2^n} \right) \approx \frac{1}{2} \epsilon_n - 1 \implies \lim_{n \rightarrow \infty} \left| \frac{\epsilon_n}{\epsilon_n - 1} \right| = \frac{1}{2}$$

Then Bisection Method is first order convergence

2.2. Newton-Taphson Method.

observation:

Let x_0 be an initial approximate to the root of $f(x) = 0$, then $x_0 + h$ is the exact root of $f(x) = 0$, i.e. $f(x_0 + h) = 0$, from Taylor series, $f(x_0 + h) = f(x_0) + h \cdot f'(x_0) + \dots$
i.e. $x_0 \approx x_0 + h$

the first order approximation, $f(x_0 + h) = f(x_0) + h \cdot f'(x_0) = 0 \implies h = \frac{-f(x_0)}{f'(x_0)}$

Let $x_1 = x_0 + h$ be the next approximation to the root, $x_1 = x_0 - \frac{f(x_0)}{f'(x_0)}$

In general $x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)} \forall n \geq 1$

Example

Consider the $f(x) = x^2 - M = 0 (M > 0)$

$$x_{n+1} = x_n - \frac{x_n^2 - M}{2x_n} = \frac{1}{2} \left(x_n + \frac{M}{x_n} \right) (\star)$$

In general, also can obtain for the k th root of M , i.e. $\sqrt[k]{M}$ with $f(x) = x^k - M = 0$ if $x_1 > \sqrt[k]{M}$, and define x_2, \dots by the interaction formula (\star) , then

(1) $\{x_n\}$ is \downarrow (trivial) (2) $\{x_n\}$ is bounded above ($x_{n+1} = \frac{1}{2} \left(x_n + \frac{M}{x_n} \right) \geq \sqrt{x_n \left(\frac{M}{x_n} \right)} = \sqrt{M}$)

By (1)(2), $\lim_{n \rightarrow \infty} x_n = \sqrt{M}$ exists.

observation

let $(x_0, f(x_0))$ be any point on the curve

$y = f(x)$, then $y - f(x_0) = f'(x_0)(x - x_0)$

Thm. 4(The NR method is 2 order convergence)

Let x denote the exact value of the root of $f(x) = 0$

x_n, x_{n+1} be two approximation S to the exact root $a, (f(a) = 0)$

if $\epsilon_n, \epsilon_{n+1}$ corresponding error S , then $x_n = a + \epsilon_n, x_{n+1} = a + \epsilon_{n+1}$

by(NR)

$$\begin{aligned}
 a + \epsilon_{n+1} &= a + \epsilon_n - \frac{f(a + \epsilon_n)}{f'(a + \epsilon_n)} \\
 \epsilon_{n+1} &= S_n - \frac{f(a) + \epsilon_n f'(a) + \frac{\epsilon_n^2}{2!} f''(a) + \dots}{f'(a) + \epsilon_n f''(a) + \frac{\epsilon_n^2}{2!} f'''(a) + \dots} \\
 &= \epsilon_n - \frac{\epsilon_n \left(f'(a) + \epsilon_n f''(a) + \frac{\epsilon_n^2}{2!} f'''(a) + \dots \right)}{f'(a) + \epsilon_n f''(a) + \frac{\epsilon_n^2}{2!} f'''(a) + \dots} \\
 &= \frac{\epsilon_n [f'(a) + \epsilon_n f''(a) + \frac{\epsilon_n^2}{2!} f'''(a) + \dots - [f'(a) + \frac{\epsilon_n}{2!} f''(a) + \dots]]}{f'(a) + \epsilon_n f''(a) + \frac{\epsilon_n^2}{2!} f'''(a) + \dots} \\
 &= \frac{\epsilon_n [\frac{\epsilon_n}{2} f''(a) + \frac{\epsilon_n^2}{3} f'''(a) + \dots]}{f'(a) + \epsilon_n f''(a) + \frac{\epsilon_n^2}{2!} f'''(a) + \dots} \\
 &= \frac{\epsilon_n^2 [\frac{1}{2} f''(a) + \frac{\epsilon_n}{3} f'''(a) + \dots]}{f'(a) [1 + \epsilon_n \frac{f''(a)}{f'(a)} + \frac{\epsilon_n^2}{2!} \frac{f'''(a)}{f''(a)} + \dots]} \\
 \Rightarrow \frac{\epsilon_{n+1}}{\epsilon_n^2} &= \frac{\frac{1}{2} f''(a) + \frac{\epsilon_n}{3} f'''(a) + \dots}{f'(a) (1 + \epsilon_n \frac{f''(a)}{f'(a)} + \dots)}
 \end{aligned}$$

$$\lim_{n \rightarrow \infty} \left| \frac{\epsilon_{n+1}}{\epsilon_n^2} \right| > \frac{1}{2} \left| \frac{f''(a)}{f'(a)} \right| < +\infty$$

Remark: if $f(x)$ has double root S

3. Eigen Problem

3.1. Review eigenvalue & eigenvector.

$$A \in M_{n \times n}(\mathbb{R}/\mathbb{C}), AX = \lambda X = \lambda(IX) = \lambda IX \implies (A - \lambda I)X = 0$$

it's a homogeneous system of n linear equation, it determinate is 0

$$p(\lambda) = \det(A - \lambda I) = 0, \deg(p(\lambda)) = n$$

Define $\lambda = \begin{pmatrix} \lambda_1 \\ \vdots \\ \lambda_n \end{pmatrix}$, $X = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix}$, X is a eigen vector of A , λ is a eigenvalue of A

the normalized eigenvector $\hat{X} = \frac{1}{\|X\|} \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$ where $\|X\| = (X^T X)^{\frac{1}{2}} = (x_1^2 + \dots + x_n^2)^{\frac{1}{2}}$

if T is diagonalizable, then \exists order basis β , $\beta \ni [T]_{\beta} = D$, which is a diagonal matrix
similarly A is diagonalizable if L_A is diagonalizable

diagonalizable

$$\left\{ \begin{array}{l} \text{the c.p split} \left\{ \begin{array}{l} n \text{ distinct eigenvalue} \\ \text{other} \left\{ \begin{array}{l} \text{algebraic multiplicity} = \text{geometric multiplicity} \\ \text{algebraic multiplicity} \neq \text{geometric multiplicity} (\text{not diagonalizable}) \end{array} \right. \end{array} \right. \\ \text{the cp does not split (not diagonalizable)} \end{array} \right.$$
 (c.p. is charateristic polynomial)

E_{λ} is subspace, $E_{\lambda} = N(T - \lambda I)$, E_{λ} is T -invariant, i.e. $T(E_{\lambda}) \subseteq E_{\lambda}$, $1 \leq \dim(E_{\lambda}) \leq m$
if T is diagonalizable, then

$$V = E_{\lambda_1} \oplus E_{\lambda_2} + \dots + E_{\lambda_n} \Leftrightarrow V = k\lambda_1 \oplus \dots \oplus k\lambda_n$$

Let any eigenvalue λ be repeated r times with k linearly independent eigenvector
 r is algebraic multiplicity, k is geometric multiplicity

3.2. some introduction.

we will learn ODE and PDE next time

$$\frac{dX}{dt} = AX, X = \begin{pmatrix} x_1 \\ x_2 \end{pmatrix}, A = \begin{pmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{pmatrix}, \frac{dx_1}{dt} = a_{11}x_1 + a_{12}x_2, \frac{dx_2}{dt} = a_{21}x_1 + a_{22}x_2$$

$X = \chi e^{\lambda t}$ is the solution of system, χ is column vector, λ is parameter to be determined

$$\frac{d\chi e^{\lambda t}}{dt} = \lambda \chi e^{\lambda t} \implies \lambda \chi e^{\lambda t} = A \chi e^{\lambda t} \implies \lambda \chi = A \chi$$

Definition

The spectrum of A , radius p of the smallest circle with center at the origin and contains all the spectral radius

3.3. Power Method.**Definition**

Let $A \in M_{n \times m}(\mathbb{C})$, for $1 \leq i, j \leq n$

define $p_i(A)$ to be the sum of the abs-values of the entries of row i of A and $r_i(A)$ to be the sum of the abs-values of the entries of column j of A

$$p_i(A) = \sum_j ||A_{ij}||, \quad r_j(A) = \sum_i ||A_{ij}||$$

$$e(A) = \max(p_i(A)), \quad r(A) = \max(r_j(A)), \quad 1 \leq i, j \leq n$$

Definition

an $n \times n$ matrix A , we define the i th Geisg disk c_i to be the disk in the complex plain with center A_{ii} an radius $r_i = p_i(A) - |A_{ii}|$, $c_i = \{ z \in \mathbb{C} \mid |z - A_{ii}| < r_i \}$

Theorem(Geisg Disk Theorem 1)

Let $A \in M_{n \times n}(\mathbb{C})$, then every eigenvalue of A is contained in a Geisg Disk

pf: Let λ be eigenvalue of A r.t. eigenvector $v = \begin{pmatrix} v_1 \\ \vdots \\ v_n \end{pmatrix}$, clearly $Av = \lambda v$

Then $I_j^n = A_{ij}v_j = \lambda_{ri}$, $1 \leq i \leq n(\star)$

suppose v_k is the coordinate of V having the largest absolute, ($v_k \neq 0$)

claim $\lambda \in C_k$, i.e. $|\lambda - A_{kk}| \leq r_k$ For $i = k$, by (\star)

$$\begin{aligned} |\lambda v_k - A_{kk}v_k| &= \left| \sum_{j=1}^n A_{kj}v_j - A_{kk}v_k \right| \\ &= \left| \sum_{j \neq k} A_{kj}v_j \right| \\ &\leq \sum_{j \neq k} |A_{kj}| |v_j| \\ &\leq \sum_{j \neq k} |A_{kj}| |v_k| = r_k |v_k| \end{aligned}$$

Corollary 1

Let λ be any eigenvalue of $A \in M_{n \times n}(\mathbb{C})$, then $|\lambda| \leq p(A) = \max(p_i(A))$

pf: by Thm. $|\lambda - A_{kk}| \leq r_k$ for some k , $1 \leq k \leq n$

$|\lambda| = |\lambda - A_{kk}| + |A_{kk}| \leq r_k + |A_{kk}| = p_k(A) \leq p(A)$

Corollary 2

$A^T \in M_{n \times n}(\mathbb{C})$, $|\lambda| \leq r(A) = \max(r_j(A))$

Corollary 3

Let λ be eigenvalue of $A \in M_{n \times n}(\mathbb{C})$, $|\lambda| \leq \min \{ p(A), r(A) \}$

by corollary 1 & 2, we are done.

Theorem(Geisg Disk Theorem 2)

Let $A \in M_{n \times n}(\mathbb{C})$, k of the disks are disjoint from the others, then exactly k eigenvalue are contained in the union of these disks.

pf: the gumltprinciple

Ref: Matrix Analysis 2/e (Horn/Johnson) P.388,389

Rayleigh Power Method

Let $\lambda_1, \dots, \lambda_n$ be the eigenvalue of matrix, $|\lambda_1| > |\lambda_2| > \dots > |\lambda_n|$
 our goal is to find $|\lambda_1|$

Let x_1, \dots, x_n be eigenvectors, r.t. $\lambda_1, \dots, \lambda_n, \implies Ax_i = \lambda_i x_i, \forall 1 \leq i \leq n$
 if the matrix A (which is diagonalizable) has n linearly independent eigenvectors
 then $x = c_1 x_1 + c_2 x_2 + \dots + c_n x_n$ for some $c_i \in \mathbb{C}$

$$\begin{aligned} Ax &= A(c_1 x_1 + \dots + c_n x_n) \\ &= c_1 A x_1 + \dots + c_n A x_n \\ &= c_1 \lambda_1 x + \dots + c_n \lambda_n x \\ &= \lambda_1 \left(c_1 x + c_2 \left(\frac{\lambda_2}{\lambda_1} \right) x + \dots + c_n \left(\frac{\lambda_n}{\lambda_1} \right) x \right) \end{aligned}$$

$$\begin{aligned} A^2 x &= A \left(\lambda_1 \left(c_1 x + c_2 \left(\frac{\lambda_2}{\lambda_1} \right) x + \dots + c_n \left(\frac{\lambda_n}{\lambda_1} \right) x \right) \right) \\ &= \lambda_1 \left(c_1 A x + c_2 \left(\frac{\lambda_2}{\lambda_1} \right) A x + \dots + c_n \left(\frac{\lambda_n}{\lambda_1} \right) A x \right) \\ &= \lambda_1^2 \left(c_1 x + c_2 \left(\frac{\lambda_2}{\lambda_1} \right)^2 x + \dots + c_n \left(\frac{\lambda_n}{\lambda_1} \right)^2 x \right) \end{aligned}$$

Continue process

$$\begin{aligned} A^k x &= \lambda_1^k \left(c_1 x + c_2 \left(\frac{\lambda_2}{\lambda_1} \right)^k x + \dots + c_n \left(\frac{\lambda_n}{\lambda_1} \right)^k x \right) \\ A^{k+1} x &= \lambda_1^{k+1} \left(c_1 x + \dots + c_n \left(\frac{\lambda_n}{\lambda_1} \right)^{k+1} x \right) \\ \lim_{k \rightarrow \infty} \frac{A^{k+1} x}{A^k x} &= \lambda_1 \end{aligned}$$

A stepwise procedure

- (i) $X^{(0)}$ is initial vector
- (ii) $Y^{(0)} = AX^{(0)}$
- (iii) $\lambda^{(1)}$ is the absolutely largest element, common from the vector $Y^{(0)}$
Let the remainly vector be $X^{(1)}$, $Y^{(0)} = \lambda^{(1)}X^{(1)}$
- (iv) repeating (ii) and (iii), $Y^{(k)} = \lambda^{(k+1)}X^{(k+1)}$
- (v) $|\lambda^{(k+1)}|, x^{(k+1)}$ is goal

Example

$$A = \begin{pmatrix} 0 & 2 & 4 \\ 1 & 1 & -2 \\ -2 & 0 & 5 \end{pmatrix}, \quad X^{(0)} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}, \quad Y^0 = AX^0 = \begin{pmatrix} 6 \\ 0 \\ 3 \end{pmatrix},$$

$$\lambda^{(1)} = 6, \quad Y^{(0)} = 6 \begin{pmatrix} 1 \\ 0 \\ \frac{1}{2} \end{pmatrix} = \lambda^{(1)}X^{(1)} \implies Y^{(1)} = AX^{(1)} = A^{(1)} = \begin{pmatrix} 2 \\ 0 \\ 0.5 \end{pmatrix}, \quad \lambda^{(2)} = 2$$

Inverse Power Method

Let λ_i be an eigenvalue of matrix A , then $\frac{i}{\lambda_i}$ is eigenvalue of the matrix A^{-1} , The eigenvector of A^{-1} is X_i
 pf: $Ax_i = \lambda_i x_i \implies \frac{1}{\lambda_i} (Ax_i) = x_i \implies \frac{1}{\lambda_i} x_i = A^{-1}x_i$

Shifted Power Method

Let λ_i be an eigenvalue of matrix A , then $(\lambda_i - k)$ is an eigenvalue of the matrix $A - kI$ with the same eigenvector as that matrix A
 pf: $Ax_i = \lambda_i x_i \implies (A - kI)x_i = AX_i - kX_i = \lambda_i x_i - kx_i = (\lambda_i - k)x_i$

4. Review Linear Algebra

4.1. Lagrange polynomials. Let $T : P_n(F) \rightarrow F^{n+1}$ be linear transform defined by $T(f) = (f(c_0), \dots, f(c_n))$, which c_0, c_1, \dots, c_n are distinct scalars in an infinite field F , β be the stander order basis for $P_n(F)$, γ be the stander order basis for F^{n+1}

Claim 1:

$$[T]_{\beta}^{\gamma} = M = \begin{bmatrix} 1 & c_0 & c_0^2 & \cdots & c_0^n \\ 1 & c_1 & c_1^2 & \cdots & c_1^n \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & c_n & c_n^2 & \cdots & c_n^n \end{bmatrix}, \beta = \{1, x, \dots, x^n\}, \gamma = \{(1, \dots, 0), \dots, (0, \dots, 1)\}$$

$$T(1) = (1, \dots, 1), T(x) = (c_0, \dots, c_n), \dots, T(x^n) = (c_0^n, \dots, c_n^n)$$

M is called a Vandemonde Matrix

Claim2: $\det(M) \neq 0$

$\because \dim(P_n(F)) = \dim(F^{n+1}) = n+1$, T is linear **check**, T is one-to-one **check**,

$\therefore T$ is invertible, $\therefore [T]_{\beta}^{\gamma}$ is invertible $\Leftrightarrow \det([T]_{\beta}^{\gamma}) \neq 0 \implies \det(M) \neq 0$

Claim3: $\det(M) = \prod_{0 \leq i < j \leq n} (c_j - c_i)$

Proof. we use the induction on $n = \deg(P_n(F))$

$$n = 1, \det \begin{bmatrix} 1 & c_0 \\ 1 & c_1 \end{bmatrix} = c_1 - c_0$$

Suppose the statement holds for n

$$\begin{aligned} \det \begin{pmatrix} 1 & c_0 & c_0^2 & \cdots & c_0^n \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & c_n & c_n^2 & \cdots & c_n^n \end{pmatrix} &= \det \begin{pmatrix} 1 & c_0 & c_0^2 & \cdots & c_0^n \\ 0 & c_1 - c_0 & c_1^2 - c_0^2 & \cdots & c_1^n - c_0^n \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & c_n - c_0 & c_n^2 - c_0^2 & \cdots & c_n^n - c_0^n \end{pmatrix} \\ &= \det \begin{pmatrix} 1 & 0 & 0 & \cdots & 0 \\ 0 & c_1 - c_0 & c_1^2 - c_1 c_0 & \cdots & c_1^n - c_0 c_1^{n-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & c_n - c_0 & c_n^2 - c_0 c_n & \cdots & c_n^n - c_0 c_n^{n-1} \end{pmatrix} = \det \begin{pmatrix} c_1 - c_0 & c_1(c_1 - c_0) & \cdots & c_1^{n-1}(c_1 - c_0) \\ \vdots & \vdots & \ddots & \vdots \\ c_n - c_0 & c_n(c_n - c_0) & \cdots & c_n^{n-1}(c_n - c_0) \end{pmatrix} \\ &= (c_1 - c_0) \cdots (c_n - c_0) \cdot \det \begin{pmatrix} 1 & c_1 & \cdots & c_1^{n-1} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & c_n & \cdots & c_n^{n-1} \end{pmatrix} = (c_1 - c_0) \cdots (c_n - c_0) \prod_{1 \leq i < j \leq n} (c_j - c_i) \\ &= \prod_{0 \leq i < j \leq n} (c_j - c_i) \quad \blacksquare \end{aligned}$$

Let $P_n(X) = a_0 + a_1x + \cdots + a_nx^n$, where $a_0, \dots, a_n \in F$

$P_n(X)$ is a polynomial s.t. it interpolated the $n+1$ points

$$P_n(x_0) = a_0 + a_1x_0 + a_2x_0^2 + \cdots + a_nx_0^n = y_0$$

\vdots

$$P_n(x_n) = a_0 + a_1x_n + a_2x_n^2 + \cdots + a_nx_n^n = y_n$$

In matrix form

$$\begin{bmatrix} 1 & x_0 & x_0^2 & \cdots & x_0^n \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 1 & x_n & x_n^2 & \cdots & x_n^n \end{bmatrix} \begin{bmatrix} a_0 \\ \vdots \\ a_n \end{bmatrix} = \begin{bmatrix} y_0 \\ \vdots \\ y_n \end{bmatrix}$$

Now we define the Lagrange polynomials of degree $l_0(x), \dots, l_n(x)$ as

$$l_i(x_j) = \begin{cases} 1 & \text{if } i = j \\ 0 & \text{if } i \neq j \end{cases}$$

The $P_n(x) = y_0 l_0(x) + \cdots + y_n l_n(x) = \sum_{i=0}^n y_i l_i(x)$

$l_i(x)$ is an n -degree polynomial with roots says

$$l_i(x) = c_i(x - x_0) \cdots (x - x_n) = c_i \prod_{j \neq i} (x - x_j)$$

$$l_i(x_i) = 1 = 1c_i \prod_{j \neq i} (x_i - x_j), \quad c_i = \frac{1}{\prod_{j \neq i} (x_i - x_j)}, \quad l_i(x) = \frac{\prod_{j \neq i} (x - x_j)}{\prod_{j \neq i} (x_i - x_j)}$$

4.2. special matrix.

Theorem (Shor's Lemma). *Let T is a linear operator on V which is a finite dimension inner product space, Suppose the characteristic polynomial splits, Then \exists order normal basis $\beta \implies [T]_\beta$ is uppertriangle.*

Note

normal : $AA^* = A^*A$ ($TT^* = T^*T$)
 self-adjoint : $A^* = A$ ($T^* = T$)

Theorem (Spectral Theorem). *Let T be a linear operator on V which is a finite dimensional inner product space*

$\mathbb{C} : T$ is normal $\Leftrightarrow \exists$ order normal basis β containing eigenvectors $\Leftrightarrow T$ is diagonal over \mathbb{C}

$\mathbb{R} : T$ is self-adjoint $\Leftrightarrow \exists$ order normal basis β containing eigenvector $\Leftrightarrow T$ is diagonal over \mathbb{R}

T is diagonal $\implies \exists$ order normal basis $\implies [T]_\beta$ is diagonal, T is normal (over \mathbb{C})
 $\implies [T]_\beta^* [T]_\beta = [T]_\beta [T]_\beta^* \implies [T^*]_\beta [T]_\beta = [T]_\beta [T^*]_\beta \implies [T^* T]_\beta = [T T^*]_\beta$

T is diagonalizable $\begin{cases} (C) \Leftrightarrow T \text{ is normal (unitary equivalent) (by Shur's Lemma)} \\ (R) \Leftrightarrow T \text{ is self-adjoint (orthogonal equivalent) (eigenvalue is real + Shur's)} \end{cases}$

Property

T is unitary \Leftrightarrow every row(column) vectors is orthonormal basis

unitary equivalent $A \sim B \Leftrightarrow \exists$ unitary matrix $Q \implies A = Q^* B Q$

orthogonally equivalent $A \sim B \Leftrightarrow \exists$ orthogonal matrix $P \implies A = P^* B P$

Define: Let V be a vector space, $W_1, W_2 \leq V \implies V = W_1 \oplus W_2$

A function $T : V \rightarrow V$ is called projection on W_1 along W_2 if for $x = x_1 + x_2$,
 $x_1 \in W_1, x_2 \in W_2, T(x) = x_1$

Property

$$R(T) = W_1, N(T) = W_2, V = R(T) \oplus N(T)$$

Proof. Claim: $R(T) = W_1$

$$(\supseteq) x \in W_1 \implies T(x) = x \in R(T)$$

$$(\subseteq) x \in R(T) \implies \exists y \in V \implies T(y) = x$$

$$\because V = W_1 \oplus W_2 \therefore y = x_1 + x_2 \text{ for some } x_1 \in W_1, x_2 \in W_2 \implies T(y) = x_1 = x \in W_1$$

Claim: $N(T) = W_2$ **exercise**

Claim: $V = R(T) \oplus N(T), \because (1), (2)$, it's trivial ■

Property

$$T \text{ is projection} \Leftrightarrow T = T^2$$

Proof.

$$(\implies) T \text{ is projection} (V = W_1 \oplus W_2)$$

$$\text{given } y \in V, \because V = W_1 \oplus W_2, \therefore \exists x_1 \in W_1, x_2 \in W_2 \ni y = x_1 + x_2$$

$$\therefore T(y) = T(x_1 + x_2) = x_1 = T(x_1) = T(T(y))$$

$$(\Leftarrow) T = T^2 (\text{use the previous proposition to build } R(T), N(T))$$

$$\implies V = R(T) \oplus N(T) \implies T \text{ is projection}$$

Given $x \in V$

$$(1) x = T(x) + [x - T(x)]$$

$$(2) T^2(x) = T(x) (\text{assumption})$$

$$\begin{cases} (i) T(x) \in R(T) \\ (ii) x - T(x) \in N(T) \\ (iii) R(T) \cap N(T) = \{0\} \\ (iv) V = R(T) + N(T) \end{cases}$$

$$(i) T(T(x)) = T(x) \in R(T), (ii) T(x - T(x)) = T(x) - T(T(x)) = T(x) - T(x) = 0,$$

$$(iv) \text{trivial} (\because (1))$$

$$(iii) \text{ Suppose } N(T) \cap R(T) = \{v\}, v \neq 0, v \in N(T) \implies T(v) = 0$$

$$v \in R(T) \implies \exists y \in V \implies T(y) = v \implies T(T(y)) = T(v) = 0, y \in N(T), v = 0 \rightarrow \leftarrow$$
 ■

Property: every projections is uniquely determined by the range & kernal

$$\text{Let } T, U : V \rightarrow V, R(T) = R(U) = W_1, N(T) = N(U) = W_2$$

$$\forall x \in V, \text{ let } y' = T(x), y = U(x) \in W_1$$

$$T(x - y') = T(x) - T(y') = y' - y' = 0, x - y' \in W_2 \implies x \in y' + W_2 (\text{coset})$$

$$\implies \exists z' \in W_2 \ni x = y' + z' \implies y = U(x) = U(y' + z') = U(y') + U(z') = y' + 0 =$$

$$y' = T(x)$$

Theorem. orthogonal projection T , $V = R(T) \oplus N(T)$, $R(T)^\perp = N(T)$, $N(T)^\perp = R(T)$
 T is orthogonal projection $\Leftrightarrow T = T^2 = T^*$

Theorem. Let matrix A normal(\mathbb{C}), self-adjoint(\mathbb{R})

A is unitary equivalent to a diagonal matrix

u_1, \dots, u_n :eigenvectors(orthonormal), $\lambda_1, \dots, \lambda_n$:eigenvalues

$$\begin{bmatrix} u_1 & u_2 & \cdots & u_n \end{bmatrix} \begin{bmatrix} \lambda_1 & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & \lambda_n \end{bmatrix} \begin{bmatrix} u_1^T \\ \vdots \\ u_n^T \end{bmatrix} = A = \sum_{i=1}^n \lambda_i u_i u_i^T \text{ (spectral decomposition)}$$

Check A is normal

$$A = u_i u_i^*, L_A = u_i u_i^*, L_A L_A = L_A^2 = (u_i u_i^*)(u_i u_i^*) = u_i u_i^* = L_A$$

$$L_A^* = (L_A)^* = (u_i u_i^*)^* = u_i^* u_i = u_i u_i^* = L_A$$

example

$$A = \begin{bmatrix} 1 & 2 \\ 2 & -2 \end{bmatrix}, \text{c.p. of } A = (1-t)(-2-t) - 4 = t^2 + t - 6 = (t+3)(t-2)$$

$$\text{the eigenvector are } \begin{pmatrix} -1 \\ 2 \end{pmatrix}, \begin{pmatrix} 2 \\ 1 \end{pmatrix}$$

\therefore they are distinct eigenvalue \Rightarrow orthogonal

$$r_1 = \frac{1}{\sqrt{5}} \cdot \begin{pmatrix} -1 \\ 2 \end{pmatrix}, r_2 = \frac{1}{\sqrt{5}} \cdot \begin{pmatrix} 2 \\ 1 \end{pmatrix}$$

$$\begin{bmatrix} r_1 & r_2 \end{bmatrix} \begin{bmatrix} -3 & - \\ 0 & 2 \end{bmatrix} \begin{bmatrix} r_1^T \\ r_2^T \end{bmatrix} = \begin{bmatrix} 1 & 2 \\ 2 & -2 \end{bmatrix}$$

$$A = -3r_1 r_1^T + 2r_2 r_2^T$$