# 1. Error Analysis

### **Definition:**

let x is a value,  $\tilde{x}$  is a estimated value

(1) absolute error,  $E_a = |x - \tilde{x}|$ 

(2) relation error,  $E_r = \left| \frac{x - \tilde{x}}{x} \right|$ 

(3) percentage error,  $E_p = 100 \times \left| \frac{x - \tilde{x}}{x} \right|$ 

 $\exists \epsilon > 0, |x - \tilde{x}| < \epsilon$ , Then  $\epsilon$  is upper limit of the absolute error measures the absolute accuracy.

# 1.1. Error in Implementation of Numerical Methods.

- (1) Round-off Error
- (2) Overflow & Underflow
- (3) Floating Point Arithmetic and Error Propagation
- (4) Truncation Error
- (5) Machine eps (Epsilon)

# (3) Floating Point Arithmetic and Error Propagation.

Let  $x_1, x_2$  are values,  $E_1, E_2$  are error of  $x_1, x_2$ , We want to check the change of error in " + ", " - ", " \* ", "/"

Let 
$$x = x_1 + x_2$$
, error of  $x$  is  $E$ 

Then 
$$x + E = x_1 + x_2 + E_1 + E_2 \implies E = E_1 + E_2$$

by triangle inequality

Absolute Error = 
$$|E| \le |E_1| + |E_2|$$
  
Relative Error =  $\frac{|E|}{|x|} \le \frac{|E_1|}{|x|} + \frac{|E_2|}{|x|}$ 

" \*"

Let 
$$x = x_1 * x_2$$
  
Then  $x + E = (x_1 + E_1)(x_2 + E_2) = x_1x_2 + E_2x_1 + E_1x_2 + E_1E_2$   
Absolute Error  $= |E| \le |x_2E_1| + |x_1E_2|$   
Relative Error  $= \frac{|E_1|}{|x|} \le \frac{|E_1|}{|x_1|} + \frac{|E_2|}{|x_2|}$ 

"/"

Let 
$$x = x_1/x_2$$

$$x + E_x = \frac{x_1 + E_1}{x_2 + E_2} \left( \frac{x_2 - E_2}{x_2 - E_2} \right) = \frac{x_1 x_2 + E_1 x_2 - x_1 E_2}{x_2^2 - E_2^2} + E_1 E_2$$
Absolute Error =  $|E_x| = \left| \frac{E_1 x_2 - x_1 E_2}{x_2^2} \right| \le \frac{|E_1|}{|x_2|} + \frac{|x_1 E_2|}{x_2^2}$ 
Relative Error =  $\frac{|E_x|}{|x|} \le \frac{|E_1|}{|x_1|} + \frac{|E_2|}{|x_2|}$ 

(4) Truncation Error. Cause by approximation infinite with its finite terms.

Use Taylor series  $(f(x) \in P(C))$  as example

Let 
$$x = a$$
,  $f(x) = f(a) + f'(a)(x - a) + f''(a)\frac{(x - a)^2}{2!} + \dots + \frac{(x - a)^n}{n!}f^n(a) + \dots + Rn$ 

$$Rn = \int_a^x \frac{(x - t)^n}{n!}f^{(n+1)}(t)dt$$

Thm 1(First Mean Value Theorem)

If g is continuous on [a, x], then  $\exists \xi$  between a and x s.t.

$$\int_{a}^{x} g(t) dt = g(\xi)(x - a)$$

# Thm 2(Second Mean Value Theorem)

If g, h is differentiable and integrable on [a, x], h does not change sign on [a, x] then  $\exists \xi$  that  $a \leq \xi \leq x$  s.t.

$$\int_{a}^{x} g(t)h(t) dt = g(\xi) \int_{a}^{x} h(t) dt$$

since 
$$t \in [a, x], h(t) = (x - t)^n \frac{1}{n!}, f^{(n+1)}(t)$$
 is continuous  $\exists \xi \in [a, x], R_n = \frac{f^{(n+1)}(\xi)}{(n+1)!} f^{(x+1)}(\xi), \xi \in [a, a+h]$  (Ref. Violin page:799)

since power series convergent,  $R_n(x) \to 0$ ,  $as_n \to \infty$ 

Definition

Given 
$$\{a_n\}\{b_n\}, b_n \ge 0, \forall n \ge 1$$
  
 $a_n = O(b_n) \text{ if } \exists M > 0 \to |a_w| \le Mb_n \forall n \ge 1$   
 $R_n(x) = O(h^{n+1})$ 

## 1.2. Condition & Stability.

Condition number is sensitivit of the function Stability is used to describle the sensitity of the process

Condition number of the f(n)

$$CN = \frac{\left|\frac{f(x) - f(\tilde{x})}{x - \tilde{x}}\right|}{\left|\frac{x - \tilde{x}}{x}\right|} = \left|\frac{f(x) - f(\tilde{x})}{x - \tilde{x}}\right| \cdot \left|\frac{x}{f(x)}\right| = \left|\frac{x}{f(x)} \cdot f'(x)\right|$$

by Mean Value Theorem,

$$\frac{f(x) - f(\tilde{x})}{x - \tilde{x}} \approx f'(x)$$

when  $CN \leq 1$  is well condition, other is ill condition

when the function is more sensitive to change, the condition number will be more big.

# **2.** Methods for f(x) = 0

we have four way to deal this problem

- (1) Direct analytical Method
- (2) Graphical
- (3) Trial and Error Method
- (4) Iterative Method

#### Thm. 3(Mean Value Theorem)

Let f be a continuous function on [a, b] = I (connected), if  $f(a) \le c \le f(b)$  that  $\exists \xi \in [a, b] \to f(\xi) = c$ 

#### Corollary

Let f be a continuous function on [a,b]=I (connected) i.e.  $f(a)\cdot f(b)<0\ \ni\ \exists c\in(a,b)\ \ni\ f(c)=0$  c is a root of f(t)

#### Iterative Method

#### 2.1. Bisection Method.

Let a, b be fixed satisfying Thm.3

 $f(a) \cdot f(b) < 0$ , f is continuous on [a,b]. The first approximation is  $x_0 = \frac{a+b}{2}$  if  $f(a) \cdot f(x_0) \le 0$ , then By Thm. 3 the root will lie on  $(a,x_0)$  and  $x_1 = \frac{a+x_0}{2}$  continue the process, let  $x_{n-3}, x_{n-2}, x_{n-1}$  be same step, then nth approximation if  $f(x_n-1) \cdot f(x_{n-3}) \le 0$ , then  $x_n = \frac{x_{n-1}+x_{n-2}}{2}$  else  $f(x_n-1) \cdot f(x_{n-3}) \ge 0$ , then  $x_n = \frac{x_{n-1}+x_{n-3}}{2}$  we shall label the interval by algorithm

$$[a,b] = [a_0,b_0], [a_1,b_1][a_2,b_2], \cdots$$

by construction  $b_n a_n = \frac{1}{2}(b_{n-1} - a_{n-1})$ , Hence  $b_n - a_n = \frac{1}{2^n}[b_0 - a_0]$ ,  $\forall n \ge 1$ Clearly  $a_0 \le a_1 \le \cdots \le b$ ,  $b_0 \ge b_1 \ge \cdots \ge a$ ,  $\{a_n\}$ ,  $\{b_n\}$  is bdd and monotonic

$$\lim_{n \to \infty} a_n = \lim_{n \to \infty} b_n = f(r)$$

by assumption 
$$f(a_n)f(b_n) < 0$$
,  $\lim_{n \to \infty} f(a_n) = f(\lim_{n \to \infty} a_n) = f(r)$   
 $\therefore f(b_n) = f(r), 0 \le [f(r)]^2 \le 0 \implies f(r) = 0$ 

The process is called **nested internal property** 

Let  $\{C_k\}_{k=1}^{\infty}$  is a  $\downarrow$  sequence of nonempty closed compact subset of X, then  $\cap k \subset k \neq \emptyset$  if  $c_k \to 0$ , then  $\cap_k c_k = \{r\}$ 

Let  $\xi$  be the solution f(x) = 0, then  $\{x_0 - \xi\} \le \frac{b - a}{2}, \dots, \{x_n - \xi\} \le \frac{b - a}{2^{n+1}}$ 

# Definition(p-order-convergence)

$$\{x_n\}$$
: seq,  $x_n \to z$ ,  $s_n \to \infty$ , define  $\epsilon_n = z - x_n$ , if  $\exists c > 0, p \ge 1$ 

$$\lim_{n \to \infty} \frac{|\epsilon_{n+1}|}{|\epsilon_n|^p} = c$$

we call  $\{x_n\}$  is p order convergence

if  $c \leq 1$ , then it's good(only check this when it's a first order convergence)

Let 
$$\epsilon_n$$
 be the error i.e.  $\epsilon_n = |x_n - \xi|$ ,  $\epsilon_n \le \frac{b-a}{2^{n+1}} \le \epsilon$ , i.e.  $h \ge \frac{\ln(b-a) - \ln\epsilon}{\ln 2} - 1$   
 $\epsilon_n = |x_n - \xi| \le \frac{1}{2} (\frac{b-a}{2^n}) \approx \frac{1}{2} \epsilon_n - 1 \implies \lim_{n \to \infty} \left| \frac{\epsilon_n}{\epsilon_n - 1} \right| = \frac{1}{2}$   
Then Bisection Method is first order convergence

# 2.2. Newton-Taphson Method.

observation:

Let  $x_0$  be an initial approximate to the root of f(x) = 0, then  $x_0 + h$  is the exact root of f(x) = 0, i.e.  $f(x_0 + h) = 0$ , from Taylor series,  $f(x_0 + h) = f(x_0) + h \cdot f(x_0) + \cdots$  i.e.  $x_0 \approx x_0 + h$ 

the first order approximation,  $f(x_0 + h) = f(x_0) + h \cdot f'(x_0) = 0 \implies h = \frac{-f(x_0)}{f'(x_0)}$ 

Let  $x_1 = x_0 + h$  be the next approximation to the root,  $x_1 = x_0 - \frac{f(x_0)}{f'(x_0)}$ 

Ingeneral  $x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)} \ \forall n \ge 1$ 

# Example

Consider the  $f(x) = x^2 - M = 0(M > 0)$ 

$$x_{n+1} = x_n - \frac{x_n^2 - M}{2x_n} = \frac{1}{2}(x_n + \frac{M}{x_n})(\star)$$

In general, also can obtain for the kth root of M, i.e.  $\sqrt[k]{M}$  with  $f(x) = x^k - M = 0$  if  $x_1 > \sqrt{M}$ , and define  $x_2, \cdots$  by the interaction formula  $(\star)$ , then

$$(1)\{x_n\}$$
 is  $\downarrow$  (trivial)  $(2)\{x_n\}$  is bounded above  $(x_{n+1} = \frac{1}{2}(x_n + \frac{M}{x_n}) \ge \sqrt{x_n(\frac{M}{x_n})} = \sqrt{M})$ 

By 
$$(1)(2)$$
,  $\lim_{n\to\infty} x_n = \sqrt{M}$  exists.

observation

let  $(x_0, f(x_0))$  be any point on the curve

y = f(x), then  $y - f(x_0) = f'(x_0)(x - x_0)$ 

Thm. 4(The NR method is 2 order convergence)

Let x denote the exact value of the root of f(x) = 0 $x_n, x_{n+1}$  be two approximation S to the exact root a, (f(a) = 0)if  $\epsilon_n, \epsilon_{n+1}$  corresponding error S, then  $x_n = a + \epsilon_n, x_{n+1} = a + \epsilon_{n+1}$ by (NR)

$$a + \epsilon_{n+1} = a + \epsilon_n - \frac{f(a - \epsilon)}{f'(a + \epsilon_n)}$$

$$\epsilon_{n+1} = S_n - \frac{f(a) + \epsilon_n f'(a) + \frac{\epsilon_n^2}{2!} f''(a) + \cdots}{f'(a) + \epsilon_n f''(a) + \frac{\epsilon_n^2}{2!} f'''(a) + \cdots}$$

$$= \epsilon_n - \frac{\epsilon_n \left( f'(a) + \epsilon_n f''(a) + \frac{\epsilon_n^2}{2!} f'''(a) + \cdots \right)}{f'(a) + \epsilon_n f''(a) + \frac{\epsilon_n^2}{2!} f''(a) + \cdots}$$

$$= \frac{\epsilon[f'(a) + \epsilon_n f''(a) + \frac{\epsilon_n^2}{2!} f(a) + \cdots - [f'(a) + \frac{\epsilon_n}{2!} f'''(a) + \cdots]]}{f'(a) + \epsilon_n f''(a) + \frac{\epsilon_n^2}{3} f''(a) + \cdots}$$

$$= \frac{\epsilon_n \left[ \frac{\epsilon_n}{2} f'(a) + \frac{\epsilon_n^2}{3} f''(a) + \cdots \right]}{f'(a) + \epsilon_n f''(a) + \frac{\epsilon_n^2}{3} f'''(a) + \cdots}$$

$$= \frac{\epsilon_{n+1}^2}{\epsilon_n^2} = \frac{\frac{\epsilon_n f''(a) + \frac{\epsilon_n}{3} f'''(a) + \cdots}{f'(a)(1 + \epsilon_n \frac{f''(a)}{f'(a)} + \cdots)}$$

$$\Rightarrow \frac{\epsilon_{n+1}}{\epsilon_n^2} = \frac{\frac{1}{2} f''(a) + \frac{\epsilon_n}{3} f'''(a) + \cdots}{f'(a)(1 + \epsilon_n \frac{f''(a)}{f'(a)} + \cdots)}$$

$$\lim_{n \to \infty} \left| \frac{\epsilon_{n+1}}{\epsilon^n} \right| > \frac{1}{2} \left| \frac{f''(a)}{f'(a)} \right| < +\infty$$

Remark: if f(x) has double root S

# 3. Eigen Problem

#### 3.1. Review eigenvalue & eigenvector.

$$A \in M_{n \times n}(R/C), \ AX = \lambda X = \lambda (IX) = \lambda IX \implies (A - \lambda I)X = 0$$

it's a homogeneous system of n linear equation, it determinate is 0

$$p(\lambda) = det(A - \lambda I) = 0, \ deg(p(\lambda)) = n$$

Define 
$$\lambda = \begin{pmatrix} \lambda_1 \\ \vdots \\ \lambda_n \end{pmatrix}$$
,  $X = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix}$ ,  $X$  is a eigen vector of A,  $\lambda$  is a eigenvalue of A

Define 
$$\lambda = \begin{pmatrix} \lambda_1 \\ \vdots \\ \lambda_n \end{pmatrix}$$
,  $X = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix}$ ,  $X$  is a eigen vector of A,  $\lambda$  is a eigenvalue of A the normalized eigenvector  $\hat{X} = \frac{1}{||X||} \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$  where  $||X|| = (X^T X)^{\frac{1}{2}} = (x_1^2 + \dots + x_n^2)^{\frac{1}{2}}$ 

if T is diagonalizable, then  $\exists$  order basis  $\beta$ ,  $\beta \ni [T]_{\beta} = D$ , which is a diagonal matrix similarly A is diagonalizable if  $L_A$  is diagonalizable

the c.p split  $\begin{cases} n \text{ distinct eigenvalue} \\ \text{other} \end{cases}$  algebraic multiplicity = geometric multiplicity  $\text{algebraic multiplicity} \neq \text{geometric multiplicity}(\mathbf{not \ diagonalizable})$  the cp does not split (not diagonalizable) (c.p. is charateristic polynomial)

 $E_{\lambda}$  is subspace  $E_{\lambda} = N(T - \lambda I)$ ,  $E_{\lambda}$  is T-invariant, i.e.  $T(E_{\lambda}) \subseteq E_{\lambda}, 1 \leq \dim(E_{\lambda}) \leq m$ if T is diagonalizable, then

$$V = E_{\lambda_1} \oplus E_{\lambda_2} + \dots + E_{\lambda_n} \Leftrightarrow V = k\lambda_1 \oplus \dots \oplus k\lambda_n$$

Let any eigenvalue  $\lambda$  be repeated r times with k linearly independent eigenvector r is algebraic multiplicity, k is geometric multiplicity

#### 3.2. some introduction.

we will learn ODE and PDE next time  $\frac{dX}{dt} = AX$ ,  $X = \begin{pmatrix} x_1 \\ x_2 \end{pmatrix}$ ,  $A = \begin{pmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{pmatrix}$ ,  $\frac{dx_1}{dt} = a_{11}x_1 + a_{12}$ ,  $\frac{dx_2}{dt} = a_{21}x_1 + a_{22}x_2$ 

 $X = \chi e^{\lambda t}$  is the solution of system,  $\chi$  is column vector,  $\lambda$  is parameter to be determind  $\frac{d\chi e^{\lambda t}}{dt} = \lambda \chi e^{\lambda t} \implies \lambda \chi e^{\lambda t} = A\chi e^{\lambda t} \implies \lambda \chi = A\chi$ 

#### **Definition**

The spectrum of A, radius p of the smallest circle with center at the origin and contains all the spactual radius

#### 3.3. Power Method.

#### **Definition**

Let  $A \in M_{n \times m}(\mathbb{C})$ , for  $1 \le i, j \le n$ 

define  $p_i(A)$  to be the sum of the abs-values of the entries of row i of A and  $r_i(A)$  to be the sum of the abs-values of the entries of column j of A

$$p_i(A) = \sum_{j=1}^{n} ||A_i j||, \quad r_j(A) = \sum_{i=1}^{n} ||A_i j||$$
$$e(A) = \max(p_i(A)), \quad r(A) = \max(r_j(A)), \quad 1 \le i, j \le n$$

#### **Definition**

an  $n \times n$  matrix A, we define the *i*th Geisg disk  $c_i$  to be the disk in the complex plain with center  $A_{ii}$  an radius  $r_i = p_i(A) - |A_{ii}|$ ,  $c_i = \{ z \in \mathbb{C} \mid |z - A_{ii}| < r_i \}$ 

# Theorem (Geisg Disk Theorem 1)

Let  $A \in M_{n \times n}(\mathbb{C})$ , then every eigenvalue of A is contained in a Geisg Disk

pf: Let 
$$\lambda$$
 be eigenvalue of  $A$  r.t. eigenvector  $v = \begin{pmatrix} v_1 \\ \vdots \\ v_n \end{pmatrix}$ , clearly  $Av = \lambda v$ 

Then 
$$I_j^n = A_{ij}v_j = \lambda_{ri}$$
,  $1 \le i \le n(\star)$ 

suppose  $v_k$  is the coordinate of V having the largest abs-solute,  $(v_k \neq 0)$  claim  $\lambda \in C_k$ , i.e.  $|\lambda - A_{kk}| \leq r_k$  For i = k, by  $(\star)$ 

$$|\lambda v_k - A_{kk}v_k| = |\sum_{j=1}^n A_{kj}v_j - A_{kk}v_k|$$

$$= |\sum_{j\neq k} A_{kj}v_j|$$

$$\leq \sum_{j\neq k} |A_{kj}||v_j|$$

$$\leq \sum_{j\neq k} |A_kj||v_k| = r_k|v_k|$$

## Corollary 1

Let 
$$\lambda$$
 be any eigenvalue of  $A \in M_{n \times n}(C)$ , then  $|\lambda| \leq p(A) = \max(p_i(A))$  pf: by Thm.  $|\lambda - A_{kk}| \leq r_k$  for some  $k, 1 \leq k \leq n$   $|\lambda| = |\lambda - A_{kk}| + |A_{kk}| \leq r_k + |A_{kk}| = p_k(A) \leq p(A)$ 

## Corollary 2

$$A^T \in M_{n \times n}(C), |\lambda| \le r(A) = \max(r_j(A))$$

#### Corollary 3

Let  $\lambda$  be eigenvalue of  $A \in M_{n \times n}(\mathbb{C})$ ,  $|\lambda| \leq \min \{ p(A), r(A) \}$  by corollary 1 & 2, we are done.

## Theorem (Geisg Disk Theorem 2)

Let  $A \in M_{n \times n}(C)$ , k of the disks are disjoint from the others, then exactly k eigenvalue are contained in the union of these disks.

pf: the gumltprinciple

Ref:Matrix Analysis 2/e (Horn/Johnson) P.388,389

#### Rayleign Power Method

Let  $\lambda_1, \dots, \lambda_n$  be the eigenvalue of matrix,  $|\lambda_1| > |\lambda_2| > \dots > |\lambda_n|$  our goal is to find  $|\lambda_1|$ 

Let  $x_1, \dots, x_n$  be eigenvectors, r.t.  $\lambda_1, \dots, \lambda_n$ ,  $\Longrightarrow Ax_i = \lambda_i x_i$ ,  $\forall 1 \leq i \leq n$  if the matrix A(which is diagonalizable) has n linearly independent eigenvectors then  $x = c_1 x_1 + c_2 x_2 + \dots + c_n x_n$  for some  $c_i \in \mathbb{C}$ 

$$Ax = A(c_1x_1 + \dots + c_nx_n)$$

$$= c_1Ax_1 + \dots + c_nAx_n$$

$$= c_1\lambda_1x + \dots + c_n\lambda_nx$$

$$= \lambda_1\left(c_1x + c_2\left(\frac{\lambda_2}{\lambda_1}\right)x + \dots + c_n\left(\frac{\lambda_n}{\lambda_1}\right)\right)$$

$$A^{2}x = A\left(\lambda_{1}\left(c_{1}x + c_{2}\left(\frac{\lambda_{2}}{\lambda_{1}}\right)x + \dots + c_{n}\left(\frac{\lambda_{n}}{\lambda_{1}}\right)\right)\right)$$

$$= \lambda\left(c_{1}Ax + c_{2}\left(\frac{\lambda_{2}}{\lambda_{1}}\right)Ax + \dots + c_{n}\left(\frac{\lambda_{n}}{\lambda_{1}}\right)Ax\right)$$

$$= \lambda_{1}^{2}\left(ax + c_{2}\left(\frac{\lambda_{2}}{\lambda_{1}}\right)^{2}x + \dots + c_{n}\left(\frac{\lambda_{n}}{\lambda_{1}}\right)^{2}x\right)$$

Continue process

$$A^{k}x = \lambda_{1}^{k} \left( c_{1}x_{1} + c_{2} \left( \frac{\lambda_{2}}{\lambda_{1}} \right)^{k} x_{2} + \dots + c_{n} \left( \frac{\lambda_{n}}{\lambda_{1}} \right)^{k} x_{n} \right)$$

$$A^{k+1}x = \lambda_{1}^{k+1} \left( c_{1}x_{1} + \dots + c_{n} \left( \frac{\lambda_{n}}{\lambda_{1}} \right)^{k+1} x_{n} \right)$$

$$\lim_{k \to \infty} \frac{A^{k+1}x}{A^{k}x} = \lambda_{1}$$

# A stepwise procedure

- (i)  $X^{(0)}$  is initial vector
- (ii)  $Y^{(0)} = AX^0$
- (iii)  $\lambda^{(1)}$  is the absolutely largest element, common from the vector  $Y^{(0)}$ Let the remainly vector be  $X^1, Y^{(0)} = \lambda^{(1)}X^{(1)}$
- (iv) reapeating (ii) and (iii),  $Y^{(k)} = \lambda^{(k+1)} X^{k+1}$
- (v)  $|\lambda^{(k+1)}|, x^{(k+1)}$  is goal

#### Example

$$A = \begin{pmatrix} 0 & 2 & 4 \\ 1 & 1 & -2 \\ -2 & 0 & 5 \end{pmatrix}, \quad X^{(0)} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}, \quad Y^0 = AX^0 = \begin{pmatrix} 6 \\ 0 \\ 3 \end{pmatrix},$$
$$\lambda^{(1)} = 6, \quad Y^{(0)} = 6 \begin{pmatrix} 1 \\ 0 \\ \frac{1}{2} \end{pmatrix} = \lambda^{(1)}X^{(1)} \implies Y^{(1)} = AX^{(1)} = A^{(1)} = \begin{pmatrix} 2 \\ 0 \\ 0.5 \end{pmatrix}, \lambda^{(2)} = 2$$

#### **Inverse Power Method**

Let  $\lambda_i$  be an eigenvalue of matrix A, then  $\frac{i}{\lambda_i}$  is eigenvalue of the matrix  $A^{-1}$ , The eigenvector of  $A^{-1}$  is  $X_i$ 

pf: 
$$Ax_i = \lambda_i x_i \implies \frac{1}{\lambda_i} (Ax_i) = x_i \implies \frac{1}{\lambda_i} x_i = A^{-1} x_i$$

#### Shifted Power Method

Let  $\lambda_i$  be an eigenvalue of matrix A, then  $(\lambda_i - k)$  is an eigenvalue of the matrix A - kI with the same eigenvector as that matrix A

pf: 
$$Ax_i = \lambda_i x_i \implies (A - kI)x_i = AX_i - kX_i = \lambda_i x_i - kx_i = (\lambda_i - k)x_i$$

# 4. Review Linear Algebra

**4.1. Lagrange polynomials.** Let  $T: P_n(\mathcal{F} \to \mathcal{F}^{n_1})$  be linear transform defined by  $T(f) = (f(c_0), \dots, f(c_n))$ , which  $c_0, c_1, \dots, c_n$  are distinct scalars in an infinite field  $\mathcal{F}, \beta$  be the stander order basis for  $P_n(\mathcal{F}), \gamma$  be the stander order basis for  $\mathcal{F}^{n+1}$ 

Claim 1:

$$[T]_{\beta}^{\gamma} = M = \begin{bmatrix} 1 & c_0 & c_0^2 & \cdots & c_0^n \\ 1 & c_0 & c_1^2 & \cdots & c_1^n \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 1 & c_n & c_n^2 & \cdots & c_n^n \end{bmatrix}, \ \beta = \{1, x, \cdots, x^n\}, \ \gamma = \{(1, \cdots, 0), \cdots, (0, \cdots, 1)\}$$

 $I(1) = (1, \dots, 1), I(x) = (c_0, \dots, c_n), \dots, I(x^n) = (c_0^n, \dots, c_n)$ 

M is called a Vandemonde Matrix

Claim2:  $det(M) \neq 0$ 

 $:: \dim(P_n(F)) = \dim(F^{n+1}) = n+1, T \text{ is linear } \mathbf{check}, T \text{ is one-to-one } \mathbf{check},$ 

T is invertible, T is invertible  $\Leftrightarrow \det([T]^{\gamma}_{\beta} \neq 0 \implies \det(M) \neq 0$ 

Claim3: 
$$det(M) = \prod_{0 \le i < j \le n} (c_j - c_i)$$

*Proof.* we use the induction on  $n = \deg(P_n(F))$ 

$$n = 1, \det \begin{bmatrix} 1 & c_0 \\ 1 & c_1 \end{bmatrix} = c_1 - c_0$$

Suppose the statement holds for n

$$\det\begin{pmatrix} 1 & c_{0} & c_{0}^{2} & \cdots & c_{0}^{n} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 1 & c_{n} & c_{n}^{2} & \cdots & c_{n}^{n} \end{pmatrix} = \det\begin{pmatrix} 1 & c_{0} & c_{0}^{2} & \cdots & c_{0}^{n} \\ 0 & c_{1} - c_{0} & c_{1}^{2} - c_{0}^{2} & \cdots & c_{1}^{n} - c_{0}^{n} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & c_{n} - c_{0} & c_{n}^{2} - c_{0}^{2} & \cdots & c_{n}^{n} - c_{0}^{n} \end{pmatrix}$$

$$= \det\begin{pmatrix} 1 & 0 & 0 & \cdots & 0 \\ 0 & c_{1} - c_{0} & c_{1}^{2} - c_{1}c_{0} & \cdots & c_{1}^{n} - c_{0}c_{1}^{n-1} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & c_{n} - c_{0} & c_{n}^{2} - c_{0}c_{n} & \cdots & c_{n}^{n} - c_{0}c_{n}^{n-1} \end{pmatrix} = \det\begin{pmatrix} c_{1} - c_{0} & c_{1}(c_{1} - c_{0}) & \cdots & c_{1}^{n-1}(c_{1} - c_{0}) \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ c_{n} - c_{0} & c_{n}(c_{n} - c_{0}) & \cdots & c_{n}^{n-1}(c_{n} - c_{0}) \end{pmatrix}$$

$$= (c_{1} - c_{0}) \cdots (c_{n} - c_{0}) \cdot \det\begin{pmatrix} 1 & c_{1} & \cdots & c_{n}^{n-1} \\ \vdots & \vdots & \vdots & \vdots \\ 1 & c_{n} & \cdots & c_{n}^{n-1} \end{pmatrix} = (c_{1} - c_{0}) \cdots (c_{n} - c_{0}) \prod_{1 \leq i < j \leq n} (c_{j} - c_{i})$$

$$= \prod_{0 \leq i \leq j \leq n} (c_{j} - c_{i})$$

Let  $P_n(X) = a_0 + a_1x + \cdots + a_nx^n$ , where  $a_0, \dots, a_n \in F$ 

 $P_n(X)$  is a polynomial s.t. it interpolated the n+1 points

$$P_n(x_0) = a_0 + a_1x_0 + a_2x_0^2 + \dots + a_nx_0^n = y_0$$

$$P_n(x_n) = a_0 + a_1 x_n + a_2 x_n^2 + \dots + a_n x_n^n = y_n$$

In matrix form

$$\begin{bmatrix} 1 & x_0 & x_0^2 & \cdots & x_0^n \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 1 & x_n & x_n^2 \cdots & x_n^n \end{bmatrix} \begin{bmatrix} a_0 \\ \vdots \\ a_n \end{bmatrix} = \begin{bmatrix} y_0 \\ \vdots \\ y_n \end{bmatrix}$$

Now we define the Lagrange polynomials of degree  $l_0(x), \dots, l_n(x)$  as

$$l_i(x_j) = \begin{cases} 1 \text{ if } i = j \\ 0 \text{ if } i \neq j \end{cases}$$

The 
$$P_n(x) = y_0 l_0(x) + \dots + y_n l_n(x) = \sum_{i=0}^n y_i l_i(x)$$

 $l_i(x)$  is an n-degree polynomial with roots says

$$l_i(x) = c_i(x - x_0) \cdots (x - x_n) = c_i \prod_{i \in I} (x - x_j)$$

$$l_i(x_i) = 1 = 1c_i \prod_{j \neq i} (x_i - x_j), \ c_i = \frac{1}{\prod_{j \neq i} (x_i - x_j)}, \ l_i(x) = \frac{\prod_{j \neq i} (x - x_j)}{\prod_{j \neq i} (x_i - x_j)}$$

#### 4.2. special matrix.

**Theorem** (Shor's Lemma). Let T is a linear operator on V which is a finite dimension inner product space, Suppose the characteristic polynomial splits, Then  $\exists$  order normal basis  $\beta \implies [T]_{\beta} \text{ is uppertriangle.}$ 

#### Note

normal :  $AA^* = A^*A(TT^* = T * T)$ self-adjoint :  $A* = A(T^* = T)$ 

**Theorem** (Spectral Theorem). Let T be a linear operator on V which is a finite dimensional inner product space

 $\mathbb{C}: T \text{ is normal} \Leftrightarrow \exists \text{ order normal basis } \beta \text{ containing eigenvectors} \Leftrightarrow T \text{ is diagonal over } \mathbb{C}$  $\mathbb{R}: T \text{ is self-adjoint} \Leftrightarrow \exists \text{ order normal basis } \beta \text{ containing eigenvector} \Leftrightarrow T \text{ is diagonal over}$  $\mathbb{R}$ 

T is diagonal 
$$\Longrightarrow \exists$$
 order normal basis  $\Longrightarrow [T]\beta$  is diagonal,  $T$  is normal(over C)  $\Longrightarrow [T]_{\beta}^{*}[T]_{\beta} = [T]_{\beta}[T]^{*} \Longrightarrow [T^{*}]_{\beta} = [T]_{\beta}[T^{*}]_{\beta} \Longrightarrow [T^{*}T]_{\beta} = [TT^{*}]_{\beta}$ 
T is diagonalizable 
$$\begin{cases} (C) \Leftrightarrow T \text{ is normal(unitary equivalent)(by Shur's Lemma)} \\ (R) \Leftrightarrow T \text{ is self-adjoint(orthogonal equivalent)(eigenvalue is real + Shur's)} \end{cases}$$

#### Property

T is unitory  $\Leftrightarrow$  every row(column) vectors is orthonormal basis

unitary equivalent  $A \sim B \Leftrightarrow \exists$  unitary matrix  $Q \implies A = Q^*BQ$  orthogonally equivalent  $A \sim B \Leftrightarrow \exists$  orthogonal matrix  $P \implies A = P^*BP$ 

Define: Let V be a vector space,  $W_1, W_2 \leq V \implies V = W_1 \oplus W_2$ A function  $T: V \to V$  is called projection on  $W_1$  along  $W_2$  if for  $x = x_1 + x_2$ ,  $x_1 \in W_1, x_2 \in W_2, T(x) = x_1$ 

#### **Property**

$$R(T) = W_1, \ N(T) = W_2, \ V = R(T) \oplus N(T)$$

*Proof.* Claim:  $R(T) = W_1$ 

$$(\supset)x \in W_1 \implies T(x) = x \in R(T)$$

$$(\subseteq)x \in R(T) \implies \exists y \in V \implies T(y) = x$$

$$V = W_1 \oplus W_2$$
  $y = x_1 + x_2$  for some  $x_1 \in W_1$ ,  $x_2 \in W_2 \implies T(y) = x_1 = x \in W_1$ 

Claim:  $N(T) = W_2$  exercise

Claim: 
$$V = R(T) \oplus N(T)$$
, ::(1),(2), it's trivial

#### **Property**

T is projection  $\Leftrightarrow T = T^2$ 

Proof.

$$(\Rightarrow) \ T \text{ is projection}(V = W_1 \oplus W_2)$$
 given  $y \in V$ ,  $\because V = W_1 \oplus W_2$ ,  $\therefore \exists x_1 \in W_1, x_2 \in W_2 \ni y = x_1 + x_2$   $\therefore T(y) = T(x_1 + x_2) = x_1 = T(x_1) = T(T(y))$   $(\Leftarrow) \ T = T^2 \text{ (use the previous proposition to build } R(T), N(T)$   $\Rightarrow V = R(T) \oplus N(T) \implies T \text{ is projection})$  Given  $x \in V$   $(1) \ x = T(x) + [x - T(x)]$   $(2)T^2(x) = T(x) \text{ (assumption)}$  
$$\begin{cases} (i)T(x) \in R(T) \\ (ii)x - T(x) \in N(T) \\ (iii)R(T) \cap N(T) = \{0\} \\ (iv)V = R(T) + N(R) \\ (i)T(T(x)) = T(x) \in R(T), \ (ii)T(x - T(x)) = T(x) - T(x) = 0, \\ (iv)\text{trivial}(\because 1)) \\ (iii) \ \text{Suppose } N(T) \cap R(T) = \{v\}, v \neq 0, v \in N(T) \implies T(v) = 0 \ \rightarrow \leftarrow V \in R(T) \implies \exists y \in V \implies T(y) = v \implies T(T(y)) = T(v) = 0, y \in N(T), v = 0 \rightarrow \leftarrow V \end{cases}$$

**Property**: every projections is uniquely determined by the range & kernal Let  $T, U : V \to V, R(T) = R(U) = W_1, \ N(T) = N(U) = W_2$  $\forall x \in V, \text{ let } y' = T(x), y = U(x) \in W_1$ 

$$T(x - y') = T(x) - T(y') = y' - y' = 0, \ x - y' \in W_2 \implies x \in y' + W_2(\text{coset})$$
  
 $\implies \exists z' \in W_2 \ni x = y' + z' \implies y = U(x) = U(y' + z') = U(y') + U(z') = y' + 0 = y' = T(x)$ 

**Theorem.** orthogonal projection T,  $V = R(T) \oplus N(T)$ ,  $R(T)^{\perp} = N(T)$ ,  $N(T)^{\perp} = R(T)$  T is orthogonal projection  $\Leftrightarrow T = T^2 = T^*$ 

**Theorem.** Let matrix A normal( $\mathbb{C}$ ), self-adjoint( $\mathbb{R}$ )

A is unitary equivalent to a diagonal matrix

 $u_1, \dots, u_n$ : eigenvectors(orthonormal),  $\lambda_1, \dots, \lambda_n$ : eigenvalues

$$\begin{bmatrix} u_1 & u_2 & \cdots & u_n \end{bmatrix} \begin{bmatrix} \lambda_1 & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & \lambda_n \end{bmatrix} \begin{bmatrix} u_1^T \\ \vdots \\ u_n^T \end{bmatrix} = A = \sum_{i=1}^n \lambda_i u_i u_i^T (spectral \ decomposition)$$

Check A is normal

$$A = u_i u_i^*, L_A = u_i u_i^*, L_A L_A = L_A^2 = (u_i u_i^*)(u_i u_i^*) = u_i u_i^* = L_A$$
  
$$L_A^* = (L_A)^* = (u_i u_i^*)^* = u_I^{**} u_i^* = u_i u_i^* = L_A$$

#### example

$$A = \begin{bmatrix} 1 & 2 \\ 2 & -2 \end{bmatrix}, \text{c.p. of } A = (1-t)(-2-t) - 4 = t^2 + t - 6 = (t+3)(t-2)$$
the eigenvector are  $\begin{pmatrix} -1 \\ 2 \end{pmatrix}, \begin{pmatrix} 2 \\ 1 \end{pmatrix}$ 

$$\therefore \text{ they are distinct eigenvalue} \implies \text{ orthogonal}$$

$$r_1 = \frac{1}{\sqrt{5}} \cdot \begin{pmatrix} -1 \\ 2 \end{pmatrix}, r_2 = \frac{1}{\sqrt{5}} \cdot \begin{pmatrix} 2 \\ 1 \end{pmatrix}$$

$$\begin{bmatrix} r_1 & r_2 \end{bmatrix} \begin{bmatrix} -3 & - \\ 0 & 2 \end{bmatrix} \begin{bmatrix} r_1^T \\ r_2^T \end{bmatrix} = \begin{bmatrix} 1 & 2 \\ 2 & -2 \end{bmatrix}$$

$$A = -3r_1r_1^T + 2r_2r_2^T$$

# 5. Integral & Differential

we have two usually problem in math

ODE (ordinary differential equation)

PDE (partial value problem)

Example 
$$\begin{cases} y' = f(x, y) \\ f(x_0) = y_0 \end{cases}$$

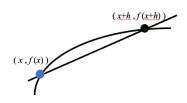
Example 
$$\left\{ \frac{\partial^2 f}{\partial x^2} + \frac{\partial f}{\partial y^2} = 0 \right\}$$
  
where  $\Delta = \frac{\partial^2}{\partial x^2} + \frac{\partial^2}{\partial y^2}$ 

## 5.1. Differentials.

 $f'(x) = \lim_{h \to 0} \frac{f(x+h) - f(x)}{h}, \text{ correction: } h > 0 \text{(time step)}$  The symbol example:  $D_{a\ b}^{\ c}$  where a is converge nation, b is method, c is first order

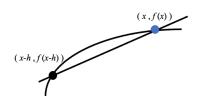
#### **Forward**

$$D_1^{\ 1}_f(f,x,h) = \frac{f(x+h) - f(x)}{h}$$



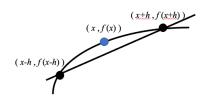
#### **Backward**

$$D_1^{\ 1}{}_b(f, x, h) = \frac{f(x) - f(x - h)}{h}$$



#### Centr

$$D_{2c}^{1}(f,x,h) = \frac{f(x+h) - f(x-h)}{2h}$$



#### Consider convergence from Taylor formula(expansion)

$$f(x+h) = f(x) + \frac{f'(x)}{1!} \cdot h + \frac{f''(x)}{2} \cdot h^2 + \cdots$$

$$= f(x) + \frac{f'(x)}{1!} \cdot h + \frac{f''(\xi)}{2!} \cdot h^2, \text{ for some } \xi \in (x, x+h), h \text{ small}$$

$$\implies f(x+h) - f(x) + \frac{f''(\xi)}{2!} \cdot h^2 = h \cdot f'(x)$$

$$\implies \frac{f(x+h) - f(x)}{h} + \frac{f''(\xi)}{2!} = f'(x)$$

which  $\frac{f(x+h)-f(x)}{h}$  is approximate value,  $\frac{f''(\xi)}{2!}$  is truncation error, f'(x) is true error

$$D_1^{\ 1}_f(f,x,h) - f'(x) = \frac{f''(\xi)}{2!} \cdot h = O(h)$$

# Convergence of Central Method

(1) 
$$f(x+h) = f(x) + f'(x) + \frac{1}{2}f''(x)h^2 + \frac{1}{6}f^{(3)}(\xi_1)h^3$$

(2) 
$$f(x-h) = f(x) - f'(x) \cdot h + \frac{1}{2}f''(x) \cdot h^2 - \frac{1}{6}f^{(3)}(\xi_2)h^3$$

(1) - (2) 
$$\implies f(x+h) - f(x-h) = 2f'(x) \cdot h + \frac{1}{6} [f^{(3)}(\xi_1) + f^{(3)}(\xi_2)]h^3$$

$$\implies 2f'(x) \cdot h + \frac{1}{6} \cdot 2f^{(3)}(\xi) \cdot h^3$$
, where  $\xi$  is closed to  $\xi_1, \xi_2$ 

$$\frac{f(x+h) - f(x-h)}{2h} - f'(x) = \frac{1}{3}f^3(\xi) \cdot h^2 = O(h^2)$$

#### Total error of Differentials

total error = rounding + truncation

$$f'(x) \approx \frac{f(x+h) - f(x)}{h} - \frac{f''(\xi)}{2!} \cdot h$$

we can change f(x+h) to  $y(x+h)+e_1$ , f(x) to  $y(x)+e_2$ , where  $e_1,e_2$  are error

$$f'(x) \approx \frac{y(x+h) + e_1 - [y(x) + e_2]}{h} - \frac{f''(\xi)}{2!} \cdot h$$
$$= \frac{[y(x+h) - y(x)]}{h} + \left(\frac{e_1 - e_2}{h} + \frac{f''(\xi)}{2!}\right)$$

and the error is a polynomial  $\rightarrow$  differ to find minimum

$$|E(h)| = \left| \frac{e_1 - e_2}{h} - \frac{f''(\xi)}{2!} \cdot h \right| \le \left| \frac{e_1 - e_2}{h} \right| + \left| \frac{f''(\xi)}{2!} h \right|$$

$$\le \frac{2e}{h} + \left| \frac{f''(\xi)}{2!} \cdot h \right| \equiv T(h) \text{ where } e = \max\{ |e_1|, |e_2| \}$$

$$T'(h) = \frac{-2e}{h^2} + \frac{f''(\xi)}{2!} = 0$$

$$h^2 = \frac{4e}{|f''(\xi)|}, h = 2\sqrt{\frac{e}{|f''(\xi)|}} \approx 10^{-8}$$

But need to find  $\xi$ 

## 5.2. Integral.

# Middle point

Let 
$$x_0 = a$$
,  $x_n = 0$ ,  $x_i = x_0 + ih$ ,  $h = \frac{b-a}{h}$ ,  $x_i^* = \frac{x_{i-1} + x_i}{2}$ 
$$I_m(f: a = b: h) = \sum_{i=1}^n h \cdot f(x_i^*)$$

## Trapezoid

Let 
$$x_0 = a$$
,  $x_n = 0$ ,  $x_i = x_0 + ih$ ,  $h = \frac{b - a}{h}$ ,  $x_i^* = \frac{x_{i-1} + x_i}{2}$ 

$$I_i(f, a, b, h) = \frac{1}{2}(f(x_0) + f(x_1)) \cdot h + \dots + \frac{1}{2}[f(x_{n-1}) + f(x_n)] \cdot h$$

$$= \frac{1}{2}f(x_0) + f(x_1) \cdot h + \dots + f(x_n - 1)h + \frac{1}{2}f(x_n)h$$

$$= \sum_{i=1}^{n} w_i f(x_i), \text{ where } w_1 = \begin{cases} \frac{1}{2}, & i = 0, n \\ 1, & i \neq 0, n \end{cases}$$

 $\textbf{Simpson Rule}(\sim O(h^2))$ 

Simpson Rule(
$$\sim O(h^2)$$
)
Let  $P_2(x) = ax^2 + bx + c$ , 
$$\begin{cases} f_0 = c \\ f_1 = ah^2 + bh + f_0 \\ f_{-1} = ah^2 - bh + f_0 \end{cases}$$
,  $b = \frac{f_1 - f_2}{h}$ ,  $a = \frac{f_1 + f_{-1} - 2f_0}{2h^2}$ 
Consider the integration, we get that

$$\int_{-h}^{h} ax^{2} + bx + c = \int_{-h}^{h} ax^{2} dx + \int_{-h}^{h} bx dx + \int_{-h}^{h} c dx$$
$$= \frac{1}{3} 2ah^{3} + 2ch$$
$$= \frac{n}{3} (f_{-1} + 4f_{0} + f_{1})$$

#### Newton Cotes Quadrature Formula

Let  $x_0 = a$ ,  $x_n = b$ ,  $x_i = x_0 + ih$ ,  $\forall 1 \le i \le n$ , then

$$I = \int_{x_0}^{x_n} f(x)dx = \int_{x_0}^{x_n} (P_n(x) + \epsilon_n(x))dx, \text{ where } \epsilon_n = \frac{f^{(n+1)}(\xi)}{(n+1)!}$$

$$\prod_{n+1} (x) = (x - x_0) \cdots (x - x_n) 
= \int_{x_0}^{x_n} f(x_0) L_0(x) + f(x_1) L_1(x) + \cdots + f(x_n) L_n(x) dx + \int_{x_0}^{x_n} \epsilon_n(x) dx 
= \int_{x_0}^{x_n} \int_{i=1}^n f(x_i) L_i(x) dx + \int_{x_0}^{x_n} \epsilon_n(x) dx 
= \sum_{i=0}^n f(x_i) \int_{x_0}^{x_n} L_i(x) dx + \int_{x_0}^{x_n} \epsilon_n(x) dx$$

 $P_n(x)$  is the Lagrange polynomial (interpolating polynomial)

 $\lambda_i$  are the constants to be determined

Let  $x = x_0 + sh$ ,  $x_i = x_0 + ih$ ,  $\forall \ 1 \le i \le n$ , then

$$L_{i}(x) = \frac{(x - x_{0})(x - x_{1}) \cdots (x - x_{i}) \cdots (x - x_{n})}{(x_{i} - x_{0}) \cdots x_{i} - x_{i} \cdots (x_{i} - x_{n})}$$

$$= \frac{(sh)((s - 1) \cdot h) \cdots ((s - n) \cdot h)}{(ih)((i - 1)h) \cdots ((i - n) \cdot h)}$$

$$= \frac{s(s - 1) \cdots (s - i + 1)(s - i - 1) \cdots (s - n)}{i(i - 1) \cdots (i - n)}$$

$$= \frac{s(s - 1) \cdots (s - i + 1)(s - i - 1) \cdots (s - n)}{i!(-1)^{n-1}(n - i)!}$$

$$\implies \lambda_{i} = \int_{x_{0}}^{x_{n}} \frac{s(s - 1) \cdots (s - i + 1)(s - i - 1) \cdots (s - n)}{i!(-1)^{n-1}(x - i)!} \cdot hds$$

we use this formula, we have

1. Simpson's 
$$\frac{1}{3}$$
 (Rule  $n=2$ )

2. Simpson's 
$$\frac{3}{8}$$
 (Rule  $n = 3$ )

3. boole rule 
$$(n=4)$$

4. Weddle rule 
$$(n = 6)$$