

*Instituto Superior Técnico, 1st December 2014*

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# Securing Client Side Web Applications

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Team INDES

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# Problem

- ❖ Client Side Web Applications:
  - ❖ Combine data / code from different origins to create a new services



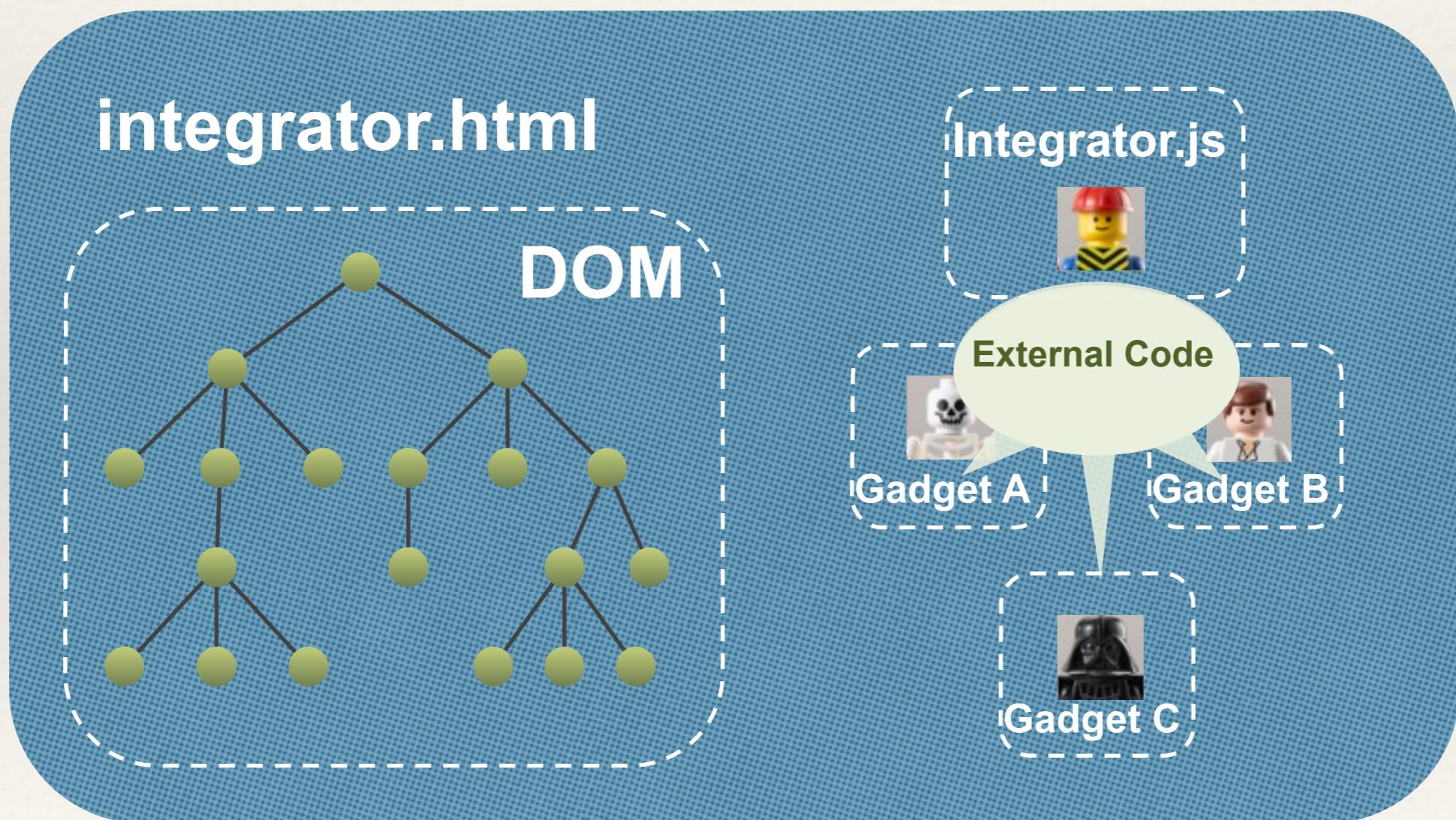
# Problem

- ❖ Client Side Web Applications:
  - ❖ Combine data / code from different origins to create a new services
  - ❖ Are implemented in JavaScript



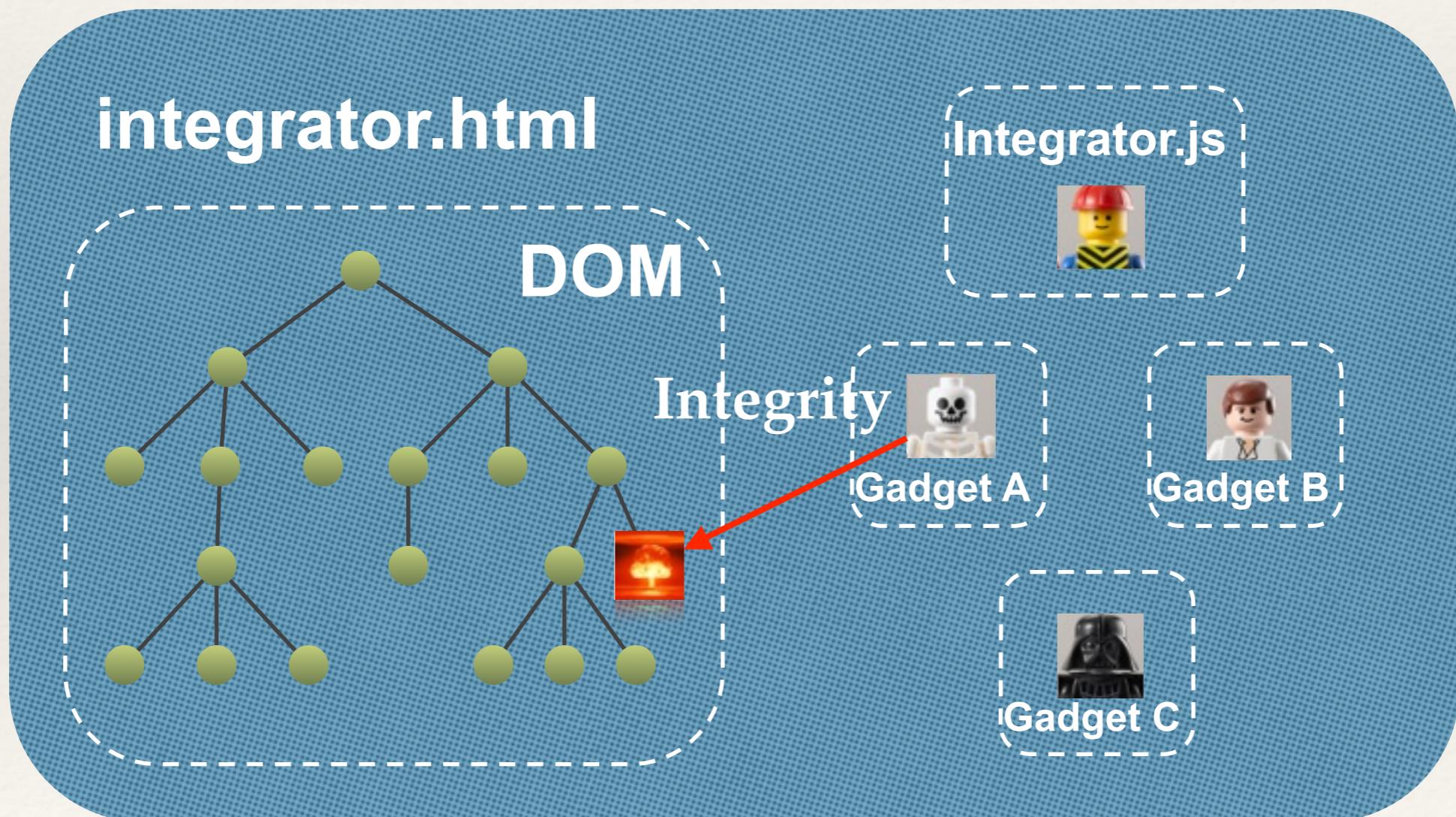
# Problem

Data and code from different origins



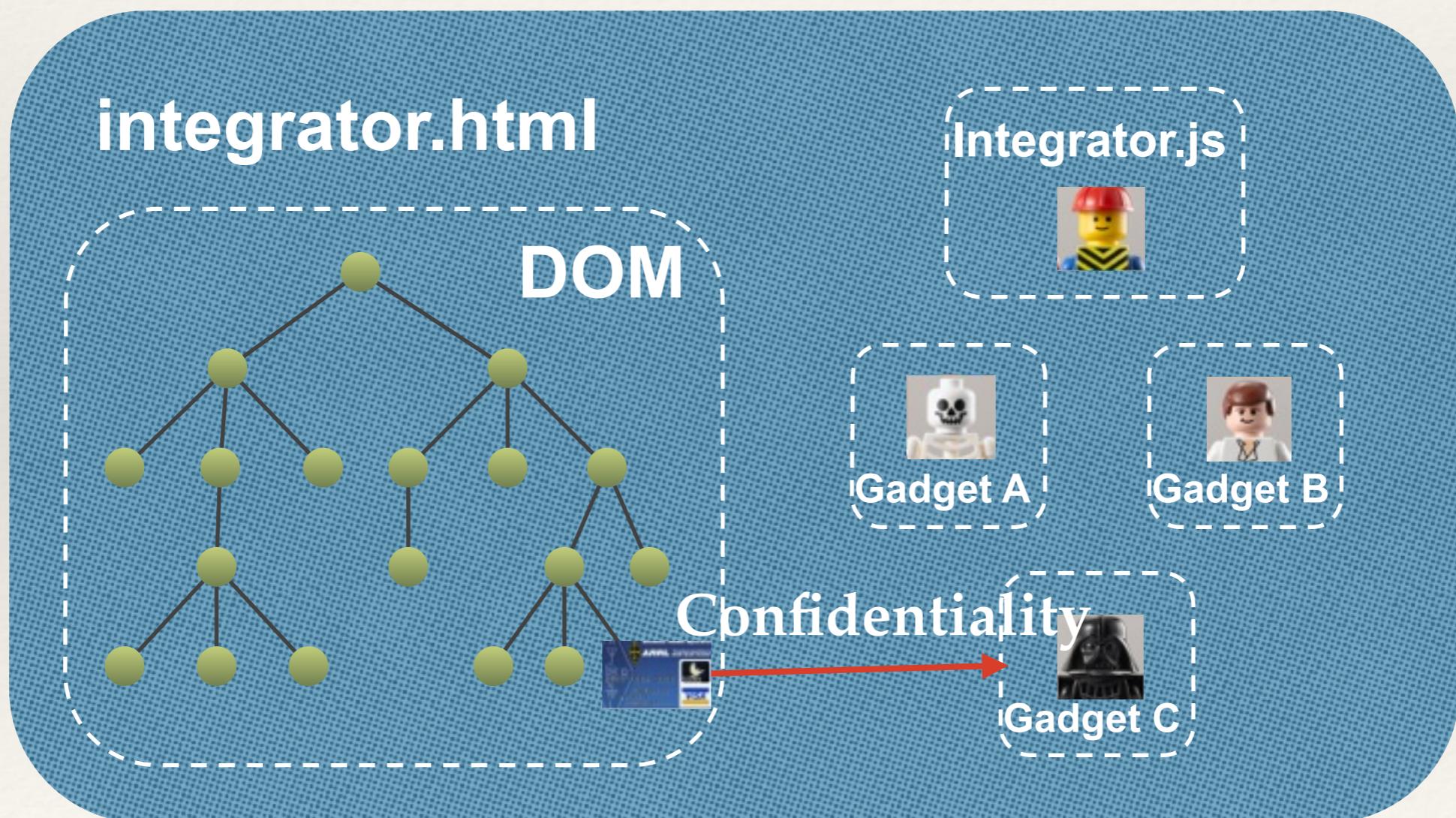
# Problem

Untrusted scripts can modify **high integrity** content



# Problem

Untrusted scripts can read high confidentiality content



# Concrete Threats in the Browser

# Common Threats in the Browser

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- **Cookie Stealing:** External scripts have unlimited access to the `cookie` property of the `window` object
- **Location Hijacking:** External scripts can either influence the document's location directly or the string variables that are forming a URL
- **History Sniffing:** Malicious scripts can check whether the user has ever visited a specific URL
- **Behaviour Tracking:** Malicious scripts can gather precise information about the user's mouse clicks and movements, scrolling behaviour

# Cookie Stealing

- External scripts have unlimited access to the `cookie` property of the `window` object
- External scripts can subsequently **transmit** the value of the cookies to an arbitrary **remote website**

```
document.cookie = "cookieval=" +  
    encodeURIComponent(val) + ";" ;
```

# Location Hijacking

- External scripts can either influence the document's location directly or the string variables that are forming a URL
- External scripts can navigate the page to a malicious website without the knowledge of the user.

```
document.location // document.URL  
window.history
```

```
history.back()  
history.forward()  
history.go()
```

# History Sniffing

- Malicious scripts can check whether the user has ever visited a specific URL
- JavaScript code can create an invisible link to the target URL and then use the browser's interface to check how this link is displayed. TRICK: the browser displays visited and unvisited links in different colours

```
$(“#link1”).css(“color”);
```

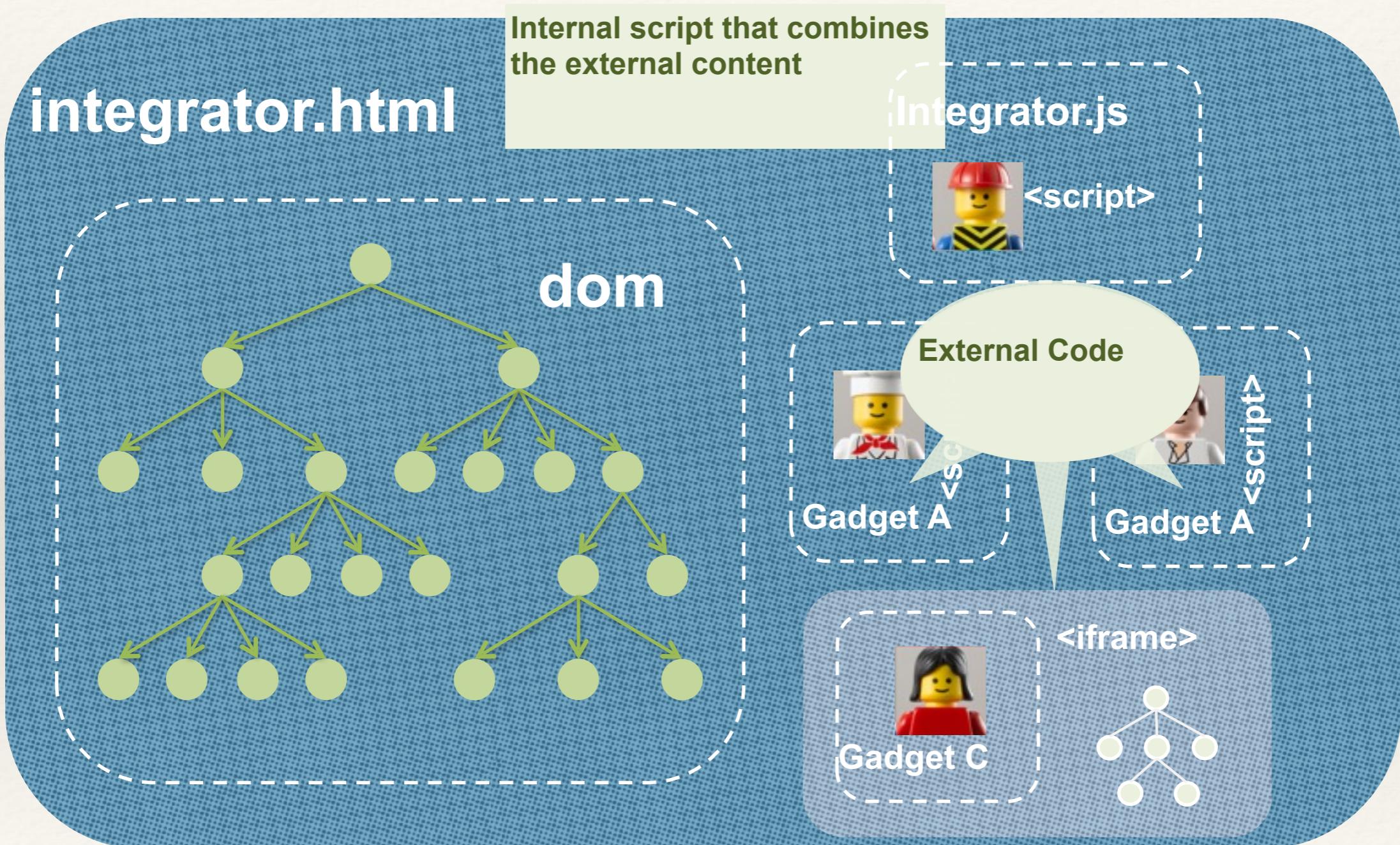
# Behaviour Tracking

- Malicious scripts can gather precise information about the user's mouse clicks and movements, scrolling behaviour

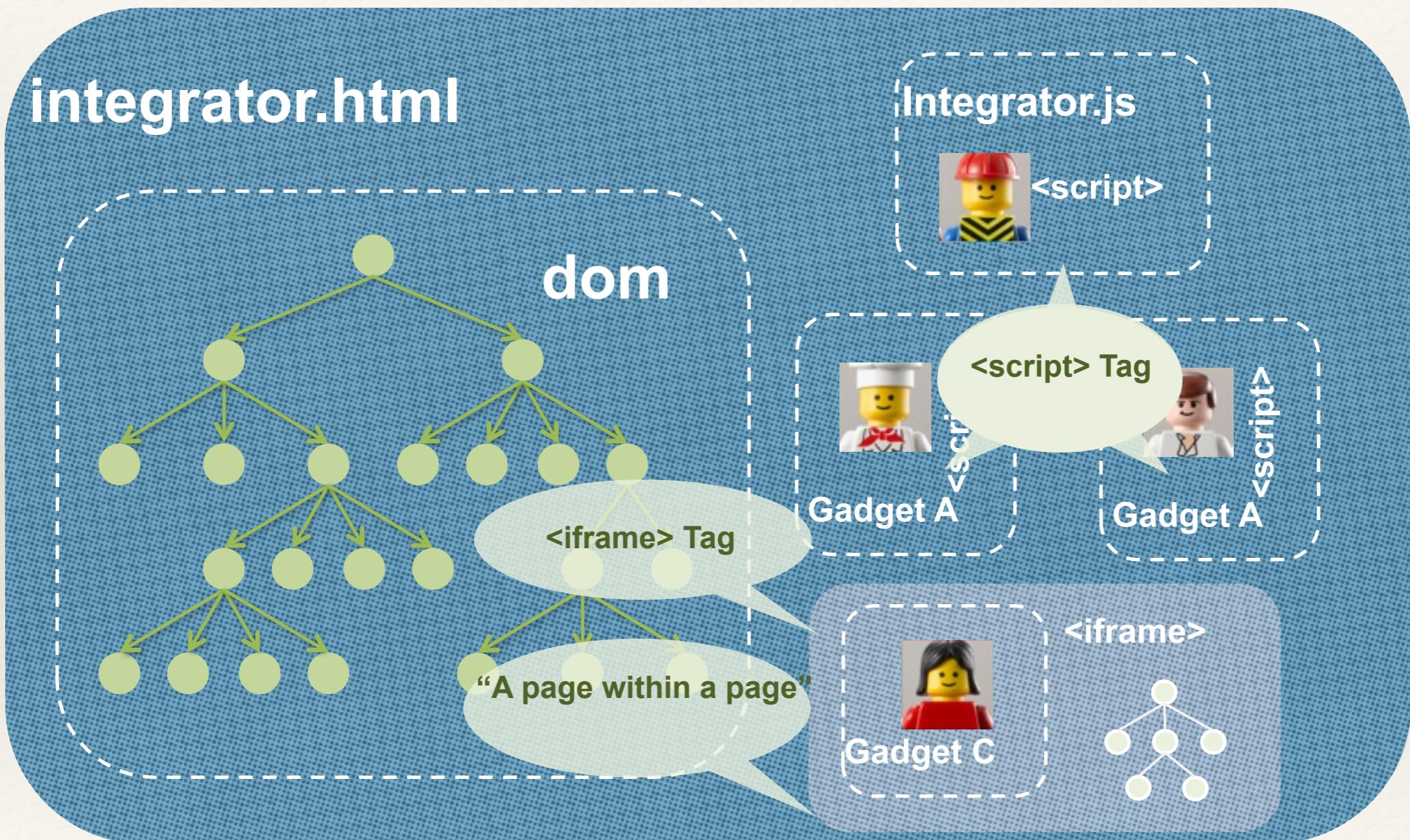
```
mydiv.addEventListener("mousedown",  
    handleMouseDown, true);
```

# The Browser Security Model

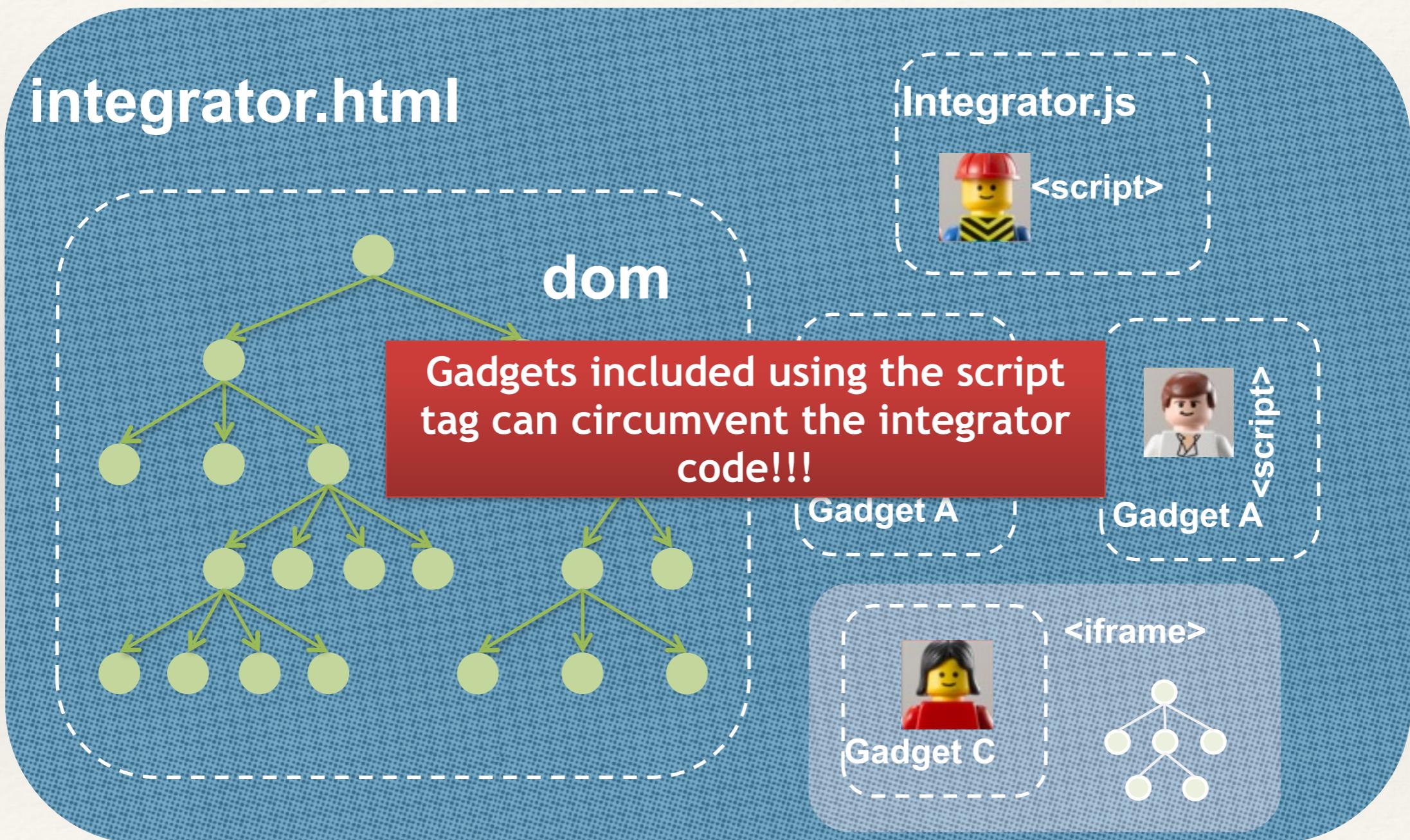
# Including External Gadgets



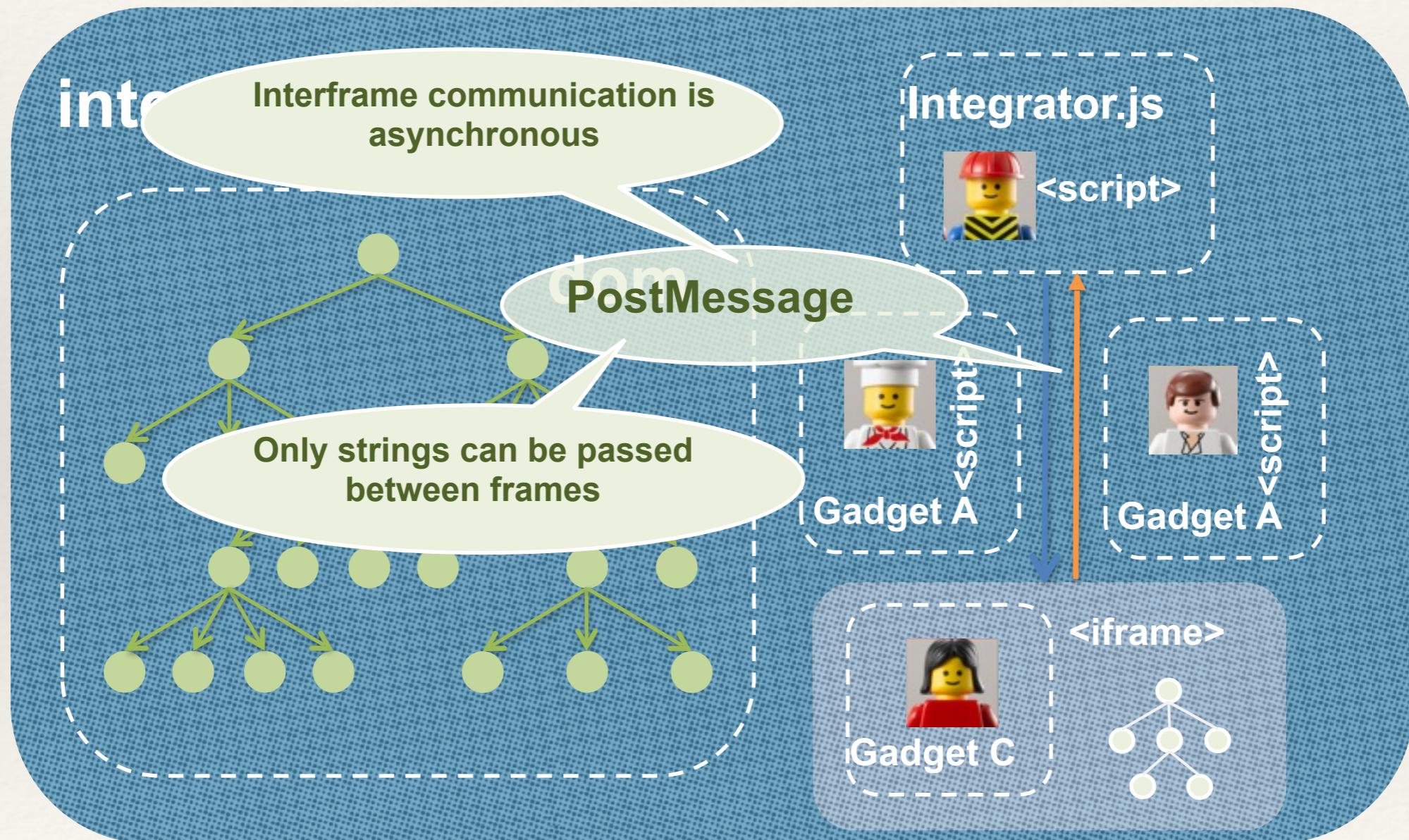
# Including External Gadgets



# Including External Gadgets



# <iframe> and the PostMessage API



# <script> versus <iframe>

## Gadgets with the script tag



**Security Issues**



**Communication**

## Gadgets with the iframe tag

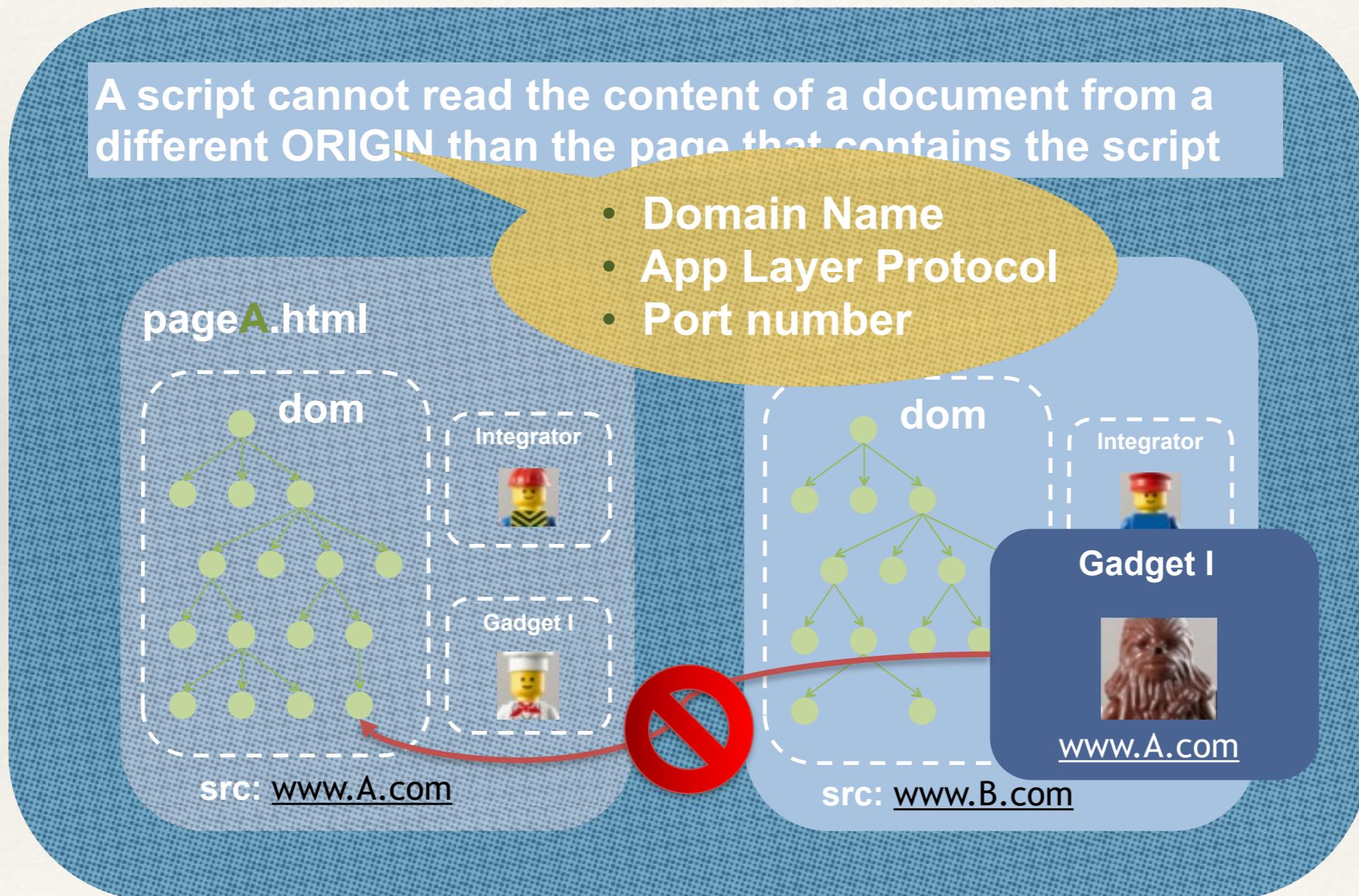


**Communication**



**Security**

# The Same Origin Policy (SOP)



# Simulating Private Fields

# Simulating Private Fields

- JavaScript does not allow the specification of accessors for object properties

```
Person = function (name, id) {  
    this.name = name;  
    this.id = id;  
}
```

```
p = new Person("John Doe", 1);  
p.id
```

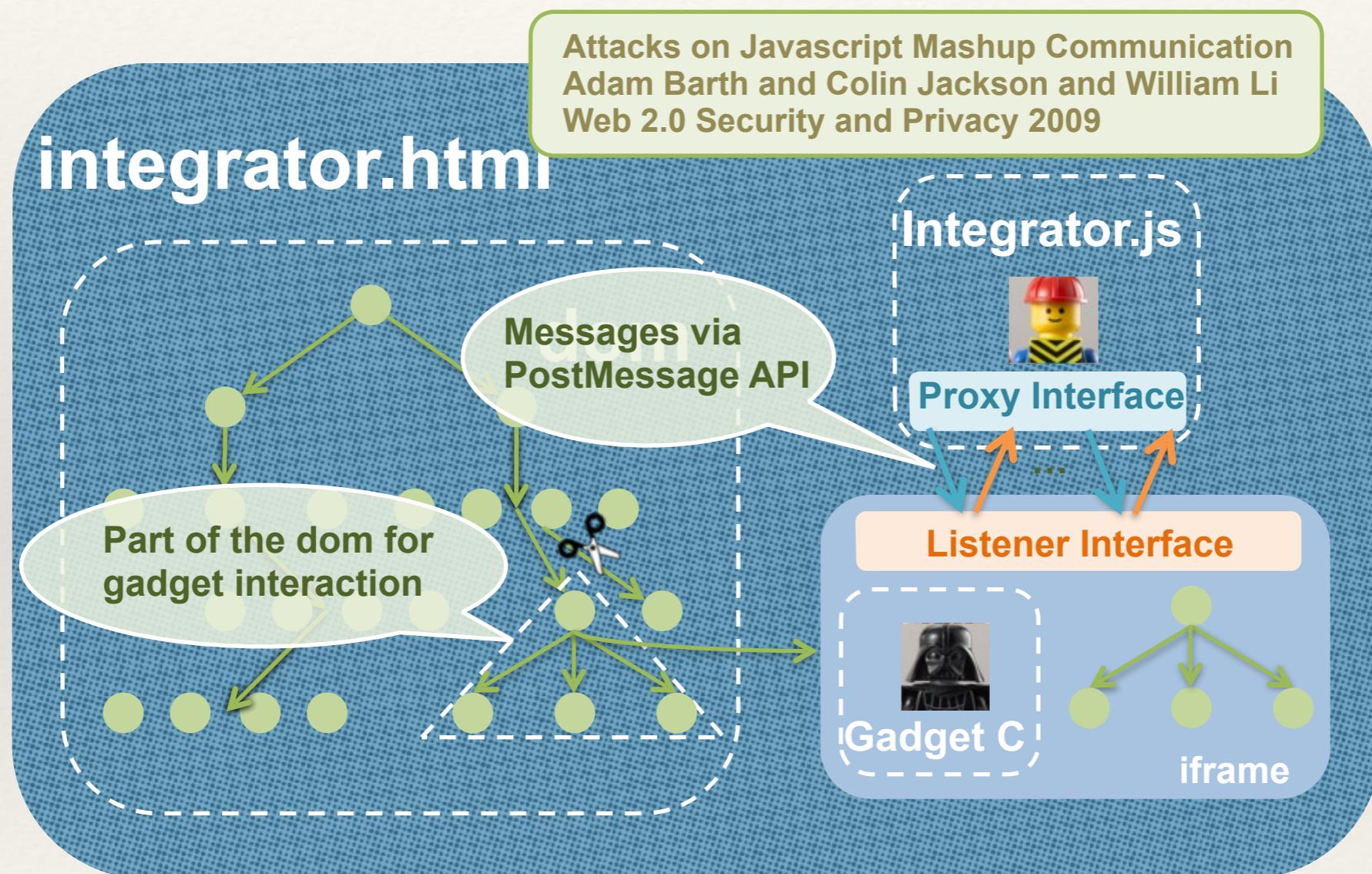
# Simulating Private Fields

- JavaScript does not allow the specification of accessors for object properties

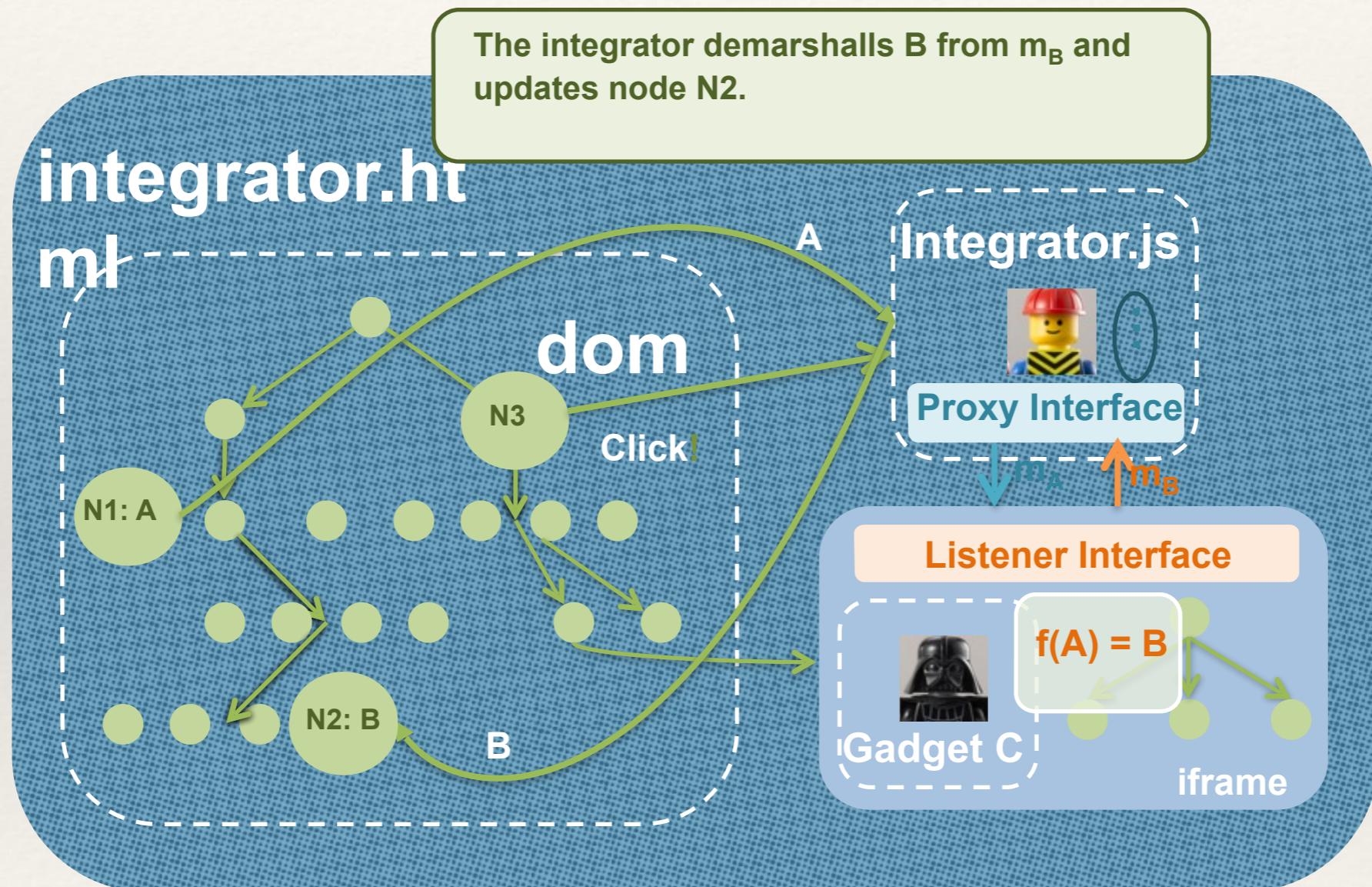
```
Person = function (name, id) {  
    var internal_id = id;  
    this.name = name;  
    this.getId = function (key) {  
        // verify the key  
        return internal_id;  
    }  
}
```

# Automatic Sandboxing of External Code

# Sandboxing External Scripts: A recipe



# Sandboxing External Scripts: A recipe



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# Security By Compilation

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Rewrite a program  $M$  as  $M'$  so that:

- ❖ The external scripts in  $M'$  are sandboxed in frames
- ❖ The semantics of  $M'$  is contained in the semantics of  $M$

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# Security By Compilation

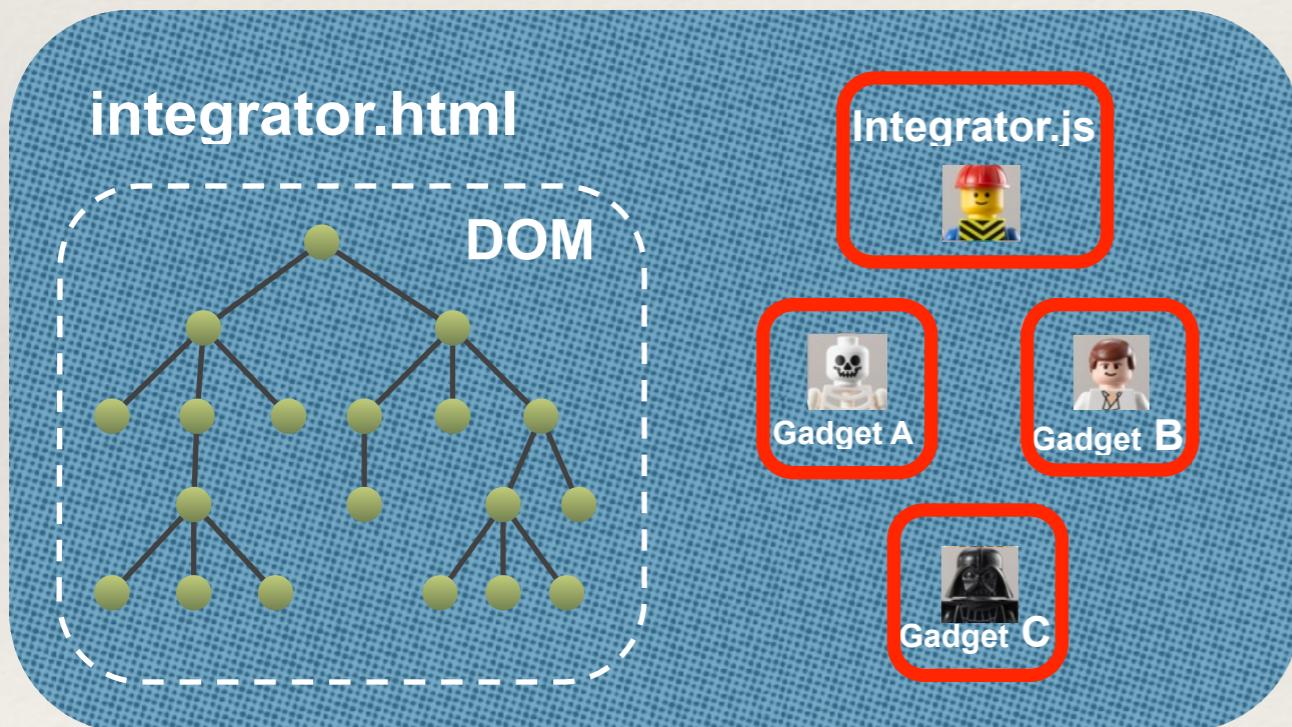
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Rewrite a program  $M$  as  $M'$  so that:

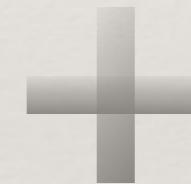
- ❖ Use a CPS transformation
- ❖ Use Opaque Object Handlers

# Monitoring Secure Information Flow in Core JavaScript

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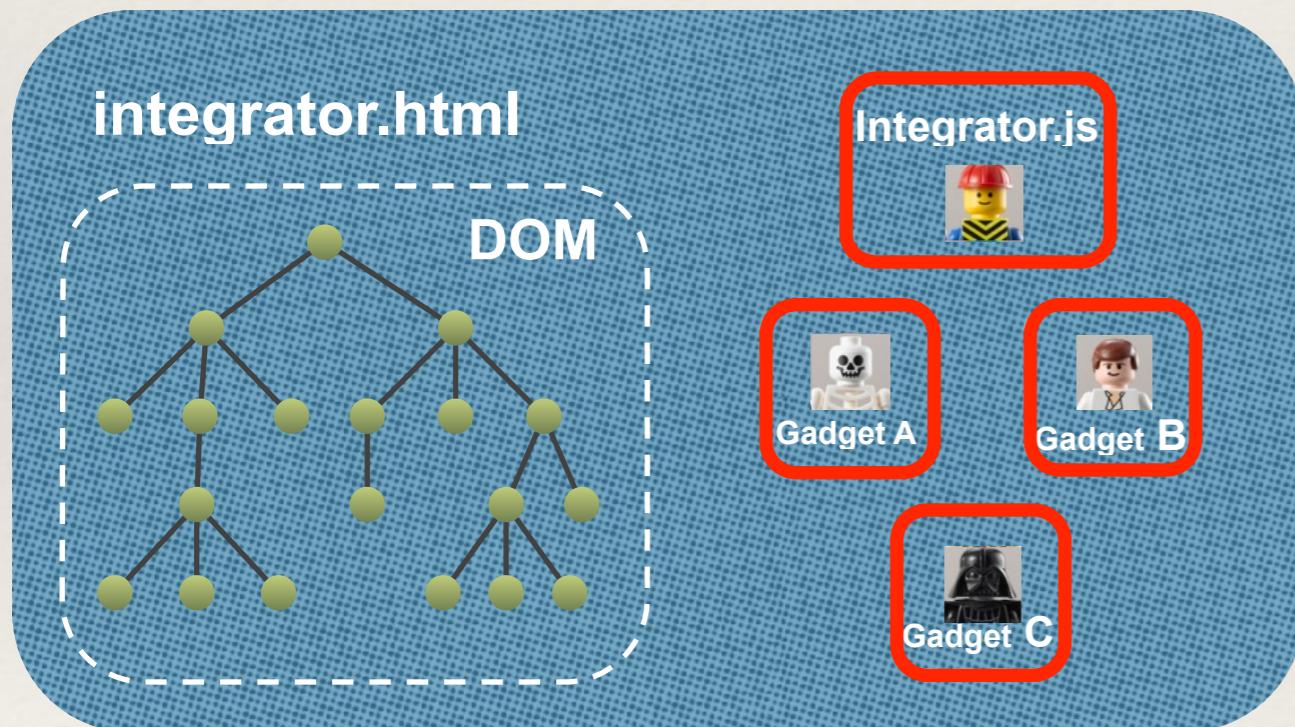


Information flow Monitor



Inlining Transformation

# Monitoring Secure Information Flow in Core JavaScript



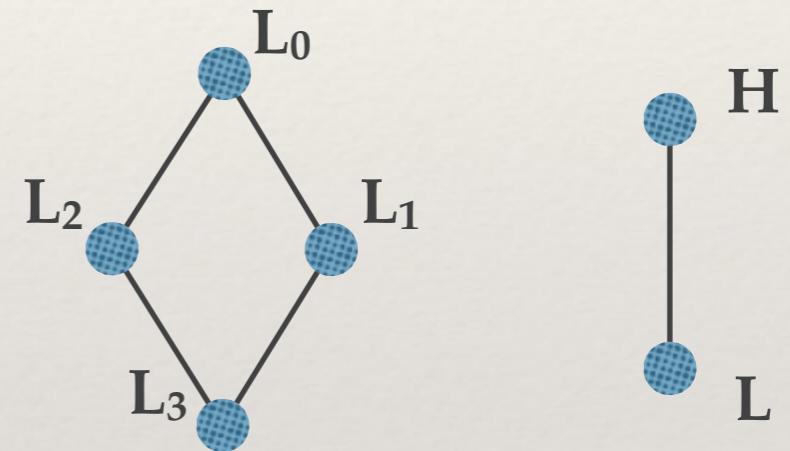
## Information flow Monitor and Inlining Transformation

- ❖ Purely Dynamic
- ❖ Flow Sensitive
- ❖ No-Sensitive Upgrade

# Labeling Program Resources

Public Outputs (Low) may NOT depend on Private Inputs (High)

- ❖ Establish a **lattice of security levels**
- ❖ Label **resources** with security levels



$\Sigma : \text{Resources} \rightarrow \text{Levels}$

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# Labeling JavaScript Objects

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What do we mean by RESOURCES?

How should we label these RESOURCES?

- ❖ What can we **know** about a JavaScript memory?
- ❖ How can we **use** the language to learn it?

Resources: Object Properties and Variables

# Labeling JavaScript Objects

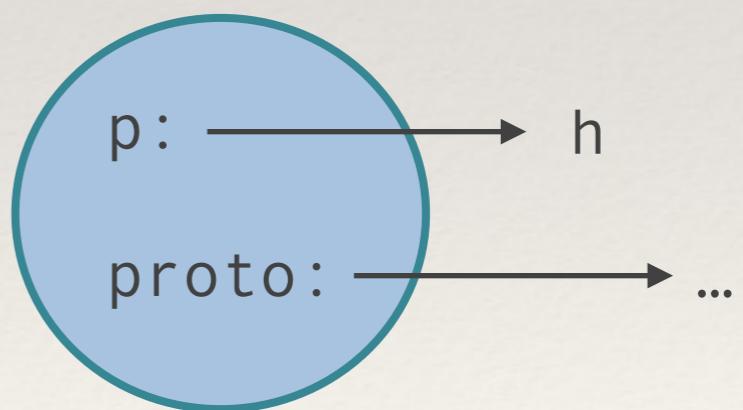
Check the **value** of a property:

```
o = {};  
o.p = h;  
l = o.p
```

Check whether a given  
property **exists**:

```
o = {};  
if (h) {  
    o.p = h  
}  
l = "p" in o
```

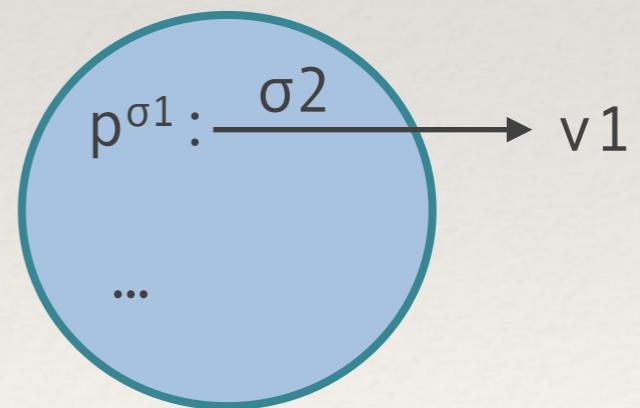
- ❖ What can we **know** about a JavaScript object?
- ❖ How can we **use the language** to learn it?



# Labeling JavaScript Objects

Associate each **property** with:

- ❖ A **value level** corresponding to its value
- ❖ An **existence level** corresponding to its existence
- ❖ What can we **know** about a JavaScript object?
- ❖ How can we **use the language** to learn it?



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# Indistinguishable JavaScript Memories

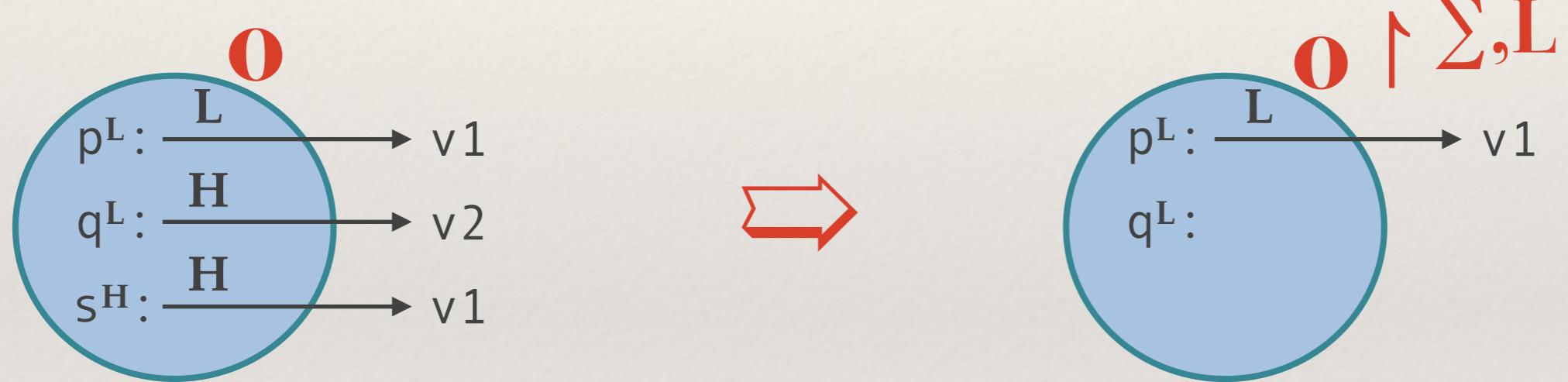
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What can an attacker see at level  $\sigma$ ?

- The existence of properties whose existence levels are  $\leq \sigma$
- The values associated with properties whose position levels are  $\leq \sigma$
- The values associated with variables whose security levels are  $\leq \sigma$

# Indistinguishable JavaScript Memories

What can an attacker see at level L?



# Indistinguishable JavaScript Memories

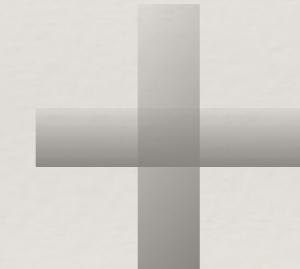
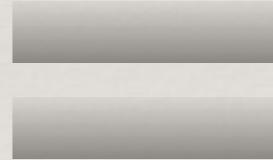
What can an **attacker observe** at a given level  $\sigma$ ?

Observable Part

Low Projection



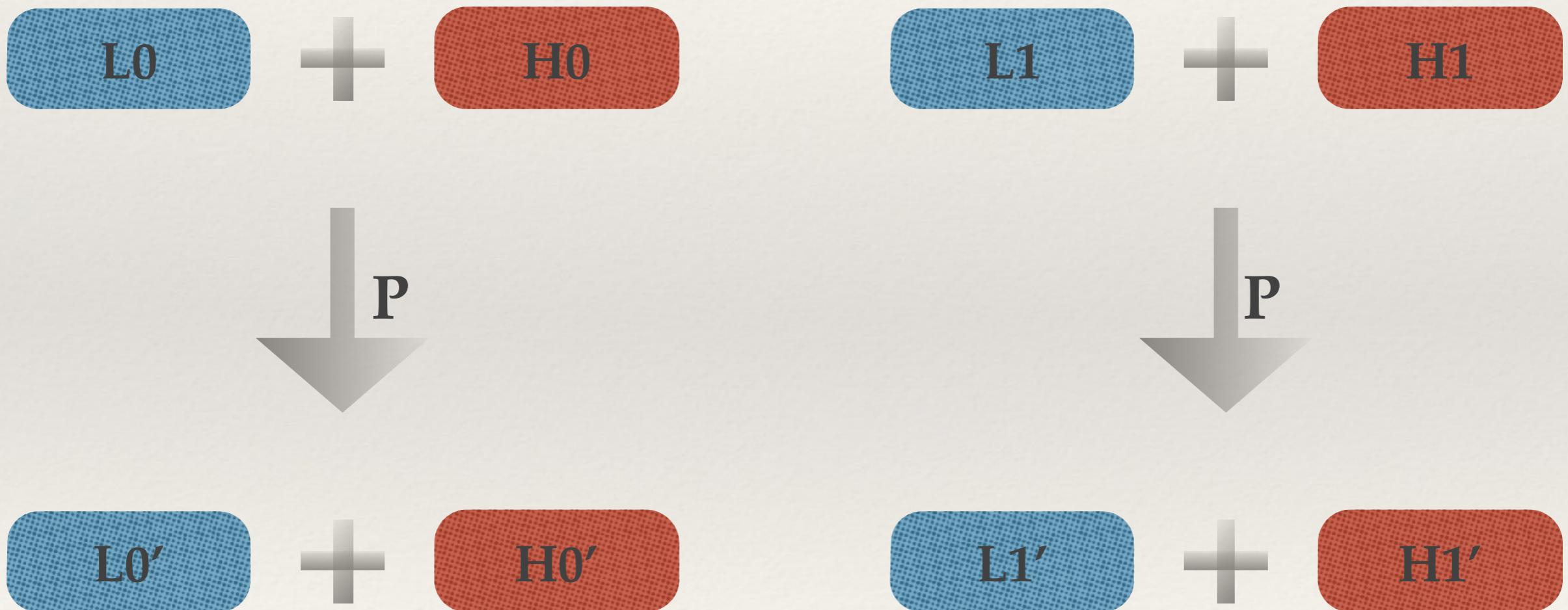
Memory



High Projection

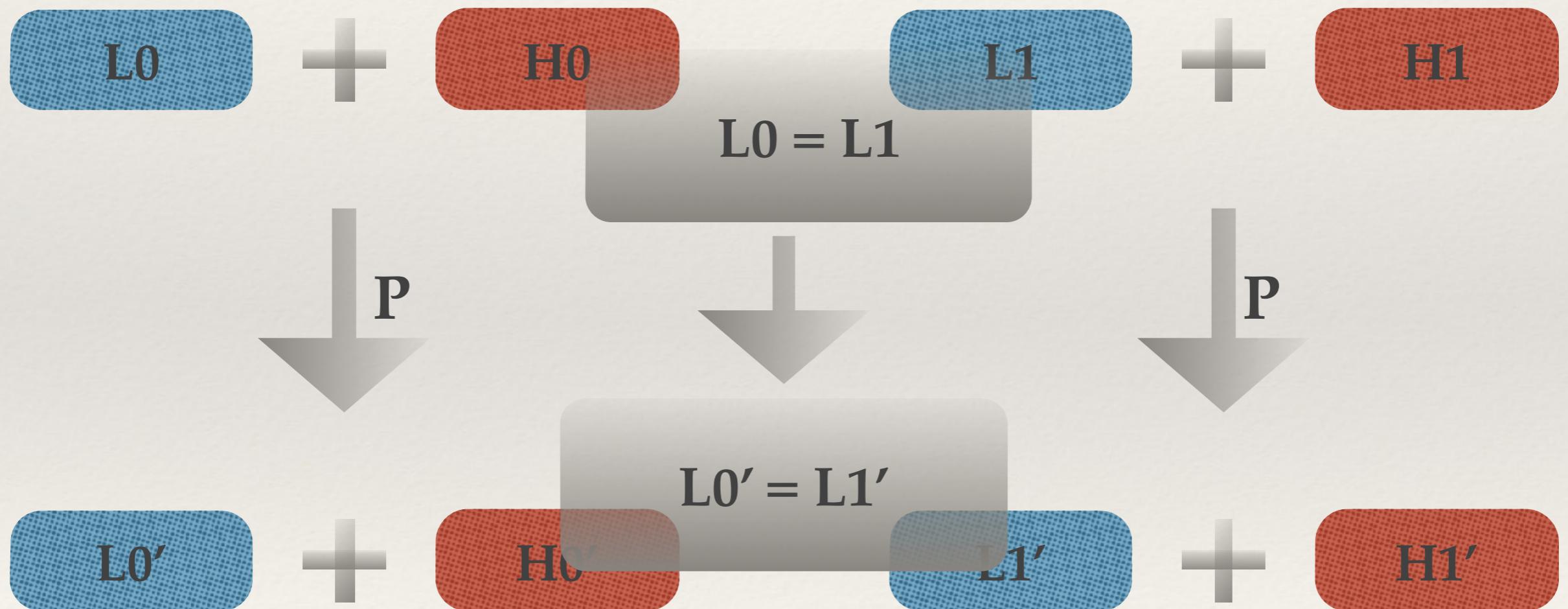
# Noninterference

When should a problem P allowed to execute?



# Noninterference

When should a problem P allowed to execute?

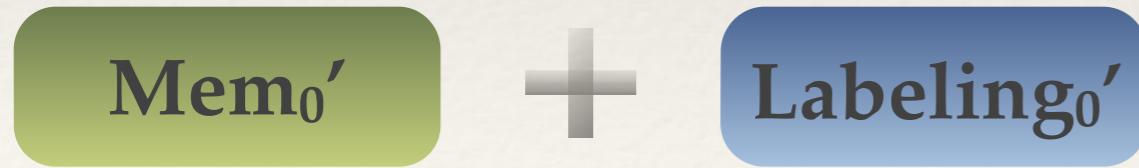


# Information Flow Monitoring

Standard Execution + Abstract Execution on the  
Labeling + Constraints

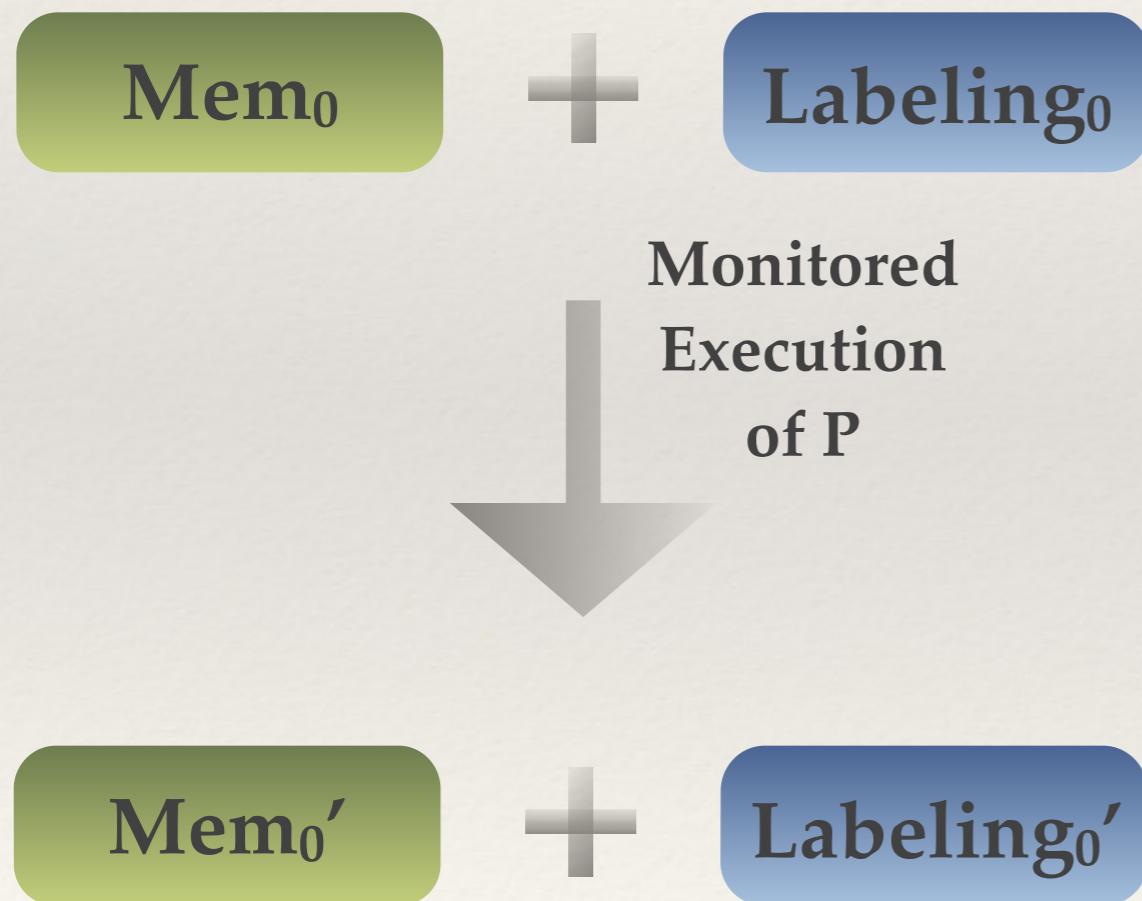


Monitored  
Execution  
of P



# Information Flow Monitoring

Standard Execution + Abstract Execution on the  
Labeling + Constraints



## Constraints:

- ❖ **No-sensitive upgrades**
- ❖ Visible resources cannot be upgraded in invisible contexts
- ❖ Example: One cannot update a property with a visible value level in an invisible context

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# Security By Compilation

---

Rewrite a program  $P$  as  $P'$  so that:

- ❖  $P'$  only executes if the execution of  $P$  is secure
- ❖ The semantics of  $P'$  is contained in the semantics of  $P$

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# Security By Compilation

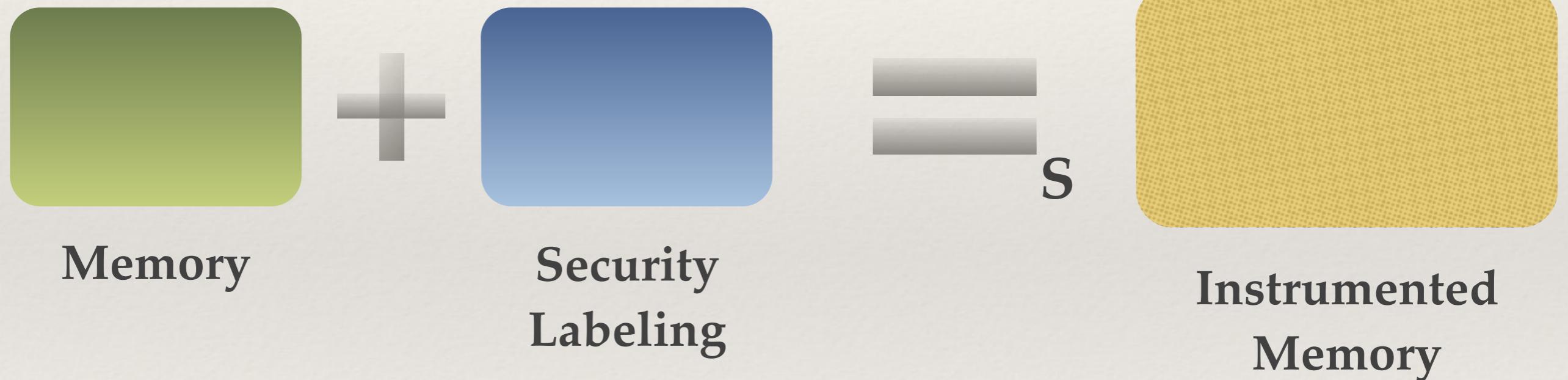
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A **recipe** for designing inlining transformations:

- ❖ Specify a monitored semantics
- ❖ Prove the monitored semantics secure
- ❖ Specify an program transformation that inlines the monitor
- ❖ Prove that the inlining compiler is correct

# Security By Compilation

Encode the Security Labeling in Memory



# Security By Compilation

$$\text{Comp}\langle P \rangle = P'$$



# PhD Grants at INRIA Sophia Antipolis

starting date between March and October 2015 (3 years)

- ❖ Formal Analysis and Implementations for JavaScript Security and Privacy:
  - ❖ **Tamara Rezk:** [tamara.rezk@inria.fr](mailto:tamara.rezk@inria.fr)
  - ❖ **Nataliia Bielova:** [nataliia.bielova@inria.fr](mailto:nataliia.bielova@inria.fr)
- ❖ Reactive Web Programming:
  - ❖ **Manuel Serrano:** [manuel.serrano@inria.fr](mailto:manuel.serrano@inria.fr)

Thank You