

Algorithm Project On Patience Sort

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History of Patience Sort

Patience Sort is a sorting Algorithm which rather than being well known for sorting is well known for finding the Longest Increasing Subsequence of a given set of data say in form of sequence of numbers in an array or just a given sequence of numbers. The inspiration behind this algorithm was from the card game **patience** where the rules were to keep making piles of cards which will remain in a sorted order in the pile.

Patience sort has many and rather unclear traces of its origins as original publications were hard to find regarding the very first implementation of the the Algorithm. However The name Patience sort was given by C.L Mallows and he credited its invention to A.S.C Ross in the 1960s. So we conclude it was invented in the 1960s.

Aldous and Diacons mentioned that it was used as finding the longest increasing subsequence by Hammersley.

Much Later Robert W.Floyd recognized this as a sorting algorithm which was analyzed by Mallows. The floyds game was developed by Floyd and Donald Knuth according to available sources based on the same principle of patience sort.

A small Overview overall and a bit about how it works

As mentioned in the history section Patience sort was inspired by the card game patience. Now to provide a short summary the way it works as a sorting algorithm is say we have a few numbers like 7 2 4 9 5 3 6 12 11 8

We start making piles and the rules of making piles is such that we can make a pile using the first number we see in this case 7 and on top of it we can keep all numbers less than 7, if it is not less than 7 we make a new pile with that greater number and keep following this process and keep stacking up numbers less then the first one in that pile like so

```
    3   8
2   5  11
7  4  9  6 12
```

Now as we see the piles here on top we use either priority queue after forming the piles or any other method to find the smallest numbers in piles, removing it from the pile and storing it in a list or an array of sorts, and doing so we will get a sorted sequence of numbers from the piles for in this case which is 2 3 4 5 6 7 8 9 11 12 . Note that a special rule we will always try to fill the leftmost pile first which implies it uses the greedy strategy.

```

IMPORT bisect and heapq
DEFINE FUNCTION sort(inputarray):
    SET piles TO [ ]
    FOR x IN inputarray:
        SET new_pile TO [x]
        SET i TO bisect.bisect_left(piles, new_pile)
        IF i != len(piles):
            piles[i].insert(0, x)
        ELSE:
            piles.append(new_pile)

    FOR i IN range(len(inputarray)):
        SET small_pile TO piles[0]
        SET inputarray[i] TO small_pile.pop(0)
        IF small_pile:
            heapq.heappreplace(piles, small_pile)
        ELSE:
            heapq.heappop(piles)
    CONTINUE till not piles

```

Pseudocode

Firstly here we will be using two python library modules named `heapq` and `bisect`. Where `heapq` will make an implementation of the priority queue algorithm and the `bisect` module will help in locating an insertion point in which if we pass a data the list will be sorted.

The `bisect` module will help us in creating the piles and `heapq` will help in retrieving the smallest number from the pile and then we will put them in the list which will finally make a sorted list after placing all of them.

Implementation based on psuedocode

Based On pseudocode the code will be divided and explained into two parts. One where I create the piles and next where i find the minimum numbers from the pile and sort them in a list. The code is given below : -

We take demo input `input = [5,12.5,-16.3,5,-4,6]`

Pile Creating Part

```
import bisect
```

```
import heapq
```

```
def patiencesort(inparr):
```

```
    piles = [ ]
```

```
    for p in inparr:
```

```
        new_pile = [p]
```

```
        q = bisect.bisect_left(piles, new_pile)
```

```
        if q!=len(piles):
```

```
            piles[q].insert(0, p)
```

```
        else:
```

```
            piles.append(new_pile)
```

Sorting part

```
    for q in range(len(inparr)):
```

```
        small_pile = piles[0]
```

```
        inparr[q] = small_pile.pop(0)
```

```
    if small_pile:
```

```
        heapq.heapreplace(piles, small_pile)
```

```
    else:
```

```
        heapq.heappop(piles)
```

```
    assert not piles
```

Passing data

```
    input = [5,12.5,-16.3,5,-4,6]
```

```
    patiencesort(input)
```

```
    print(input)
```

Explaining the pile creation part

```
1 import bisect
2 import heapq
3 def patiencesort(inparr):
4     piles = []
5     for p in inparr:
6         new_pile = [p]
7         q = bisect.bisect_left(piles, new_pile)
8         if q != len(piles):
9             piles[q].insert(0, p)
10        else:
11            piles.append(new_pile)
12
```

```
23 input = [5,12.5,-16.3,5,-4,6]
24 patiencesort(input)
```

Here we take the python's list named `input` which is going to be our set of numbers which we will sort

We pass the input to our function `patiencesort` where it is passed to the variable named `inparr` which now has the input list

list-inparr

0	1	2	3	4	5
5	12.5	-16.3	5	-4	6

```

1 import bisect
2 import heapq
3 def patiencesort(inparr):
4     piles = []
5     for p in inparr:
6         new_pile = [p]
7         q = bisect.bisect_left(piles, new_pile)
8         if q != len(piles):
9             piles[q].insert(0, p)
10        else:
11            piles.append(new_pile)
12

```

In line 4 we create a new list named piles which will be an empty list for now

We run a for loop p upto the last place of inparr which will serially iterate and store the value of inparr list starting from index 0 to 5 to the new list called new_pile. In the first run new_pile stores 5 as it is the data of the first index in inparr

list-inparr					
0	1	2	3	4	5
5	12.5	-16.3	5	-4	6

Note that python list is very beneficial as it not only can store single values but can also store another list in its list whose concept we are using to make the piles. Note that python list is also dynamic so we can add remove delete at will.

Now remember from last part we stored the value of p in a new list named new pile like so

```
6 new_pile = [p]
```

We have now the new list new_pile

new_pile - 0

5

Now I move On to line 7 where we will be using module bisect and use bisect.bisect_left to find and return the position of the list at which if the element is inserted a sorted order will be maintained .

More explanation next page-

```
1 import bisect
2 import heapq
3 def patiencesort(inparr):
4     piles = []
5     for p in inparr:
6         new_pile = [p]
7         q = bisect.bisect_left(piles, new_pile)
8         if q!=len(piles):
9             piles[q].insert(0, p)
10        else:
11            piles.append(new_pile)
12
```

To understand how bisect works take an example code -

```
main.py  [ ] [ ] Run Shell
1 import bisect
2 list = [1,2,4,6]
3 print(bisect.bisect_left(list,2))
```

```
1
>
```

Here say we have a list 1 2 4 6 having index position of 0 1 2 3 respectively

Say we want to insert 2 now where in the left most region possible would we put 2 in the list so that the list remains sorted

The answer is at position 1 as it will give 1 2 2 4 6 having index 0 1 2 3 4

This concept was based on binary search

Back To code

Now back in line 7 we have piles and new_pile in bisect.bisect_left with piles being where we want to insert and new_pile being the element we want to insert.

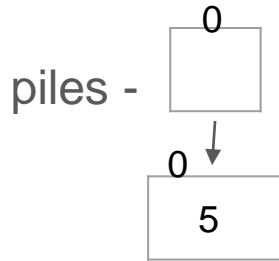
If we remember, piles was initially created as empty list so bisect.left will find that if new_pile is placed in 0th index the pile will be sorted

So $q = 0$

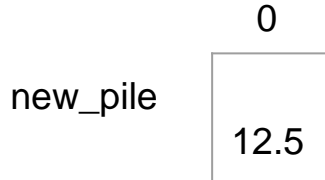
Now since piles is still an empty list so its size is 0 so in line 8 the else loop will initiate And by piles.append(new_piles)
We will add new_pile which had the value 5 into the list piles

Now piles has one data which is 5

```
1  import bisect
2  import heapq
3  def patiencesort(inparr):
4      piles = []
5      for p in inparr:
6          new_pile = [p]
7          q = bisect.bisect_left(piles, new_pile)
8          if q!=len(piles):
9              piles[q].insert(0, p)
10         else:
11             piles.append(new_pile)
12
23  input = [5,12.5,-16.3,5,-4,6]
24  patiencesort(input)
```



Now we continue the loop now `p` has value 12.5 from `inparr`
We now have `new_pile=12.5`

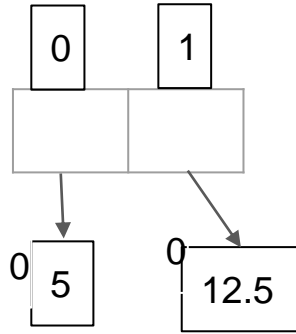


We now execute the line 7 `bisect.bisect left` and check **where would we place 12.5 so that the list named piles remains sorted ?**

Since 12.5 is greater than 5 it will be place in position / index one(1) so `q = 1`

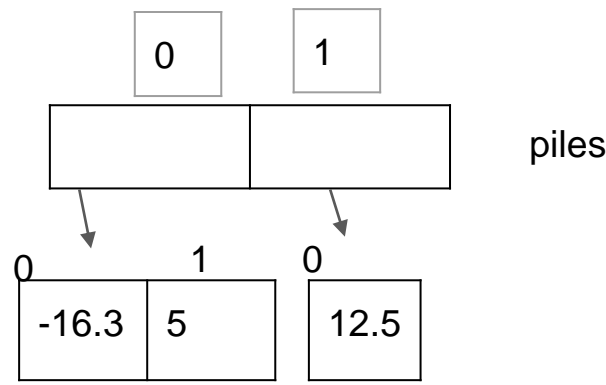
Now since $q = 1$ and length of piles is also one the if condition `if q!=len(piles):` will not execute, else condition will proceed and add 12.5 at the end of the list named piles using append

piles -



Now we continue the loop again now p denotes to -16.3

Store p in new_pile
Check where -16.3 can be place so that
piles will be sorted
Since $-16.3 < 5$ So at index 0
So $q = 0$ and $\text{len}(\text{piles}) = 2$
If loop will execute insert $p = -16.3$ at 0 th
index of $\text{piles}[q]$ i.e $\text{piles}[0].\text{insert}(0, -16.3)$



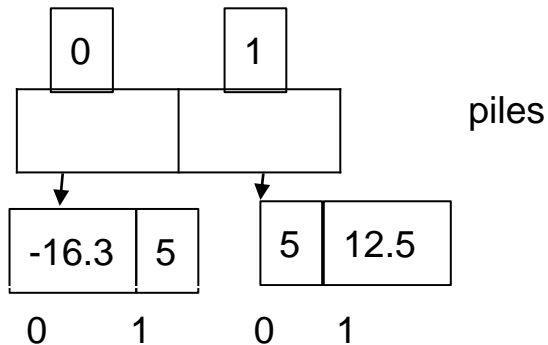
```
1 import bisect
2 import heapq
3 def patienceSort(inparr):
4     piles = []
5     for p in inparr:
6         new_pile = [p]
7         q = bisect.bisect_left(piles, new_pile)
8         if q != len(piles):
9             piles[q].insert(0, p)
10        else:
11            piles.append(new_pile)
12
23 input = [5,12.5,-16.3,5,-4,6]
24 patienceSort(input)
```

Now moving on p denotes to 5.

We store in new_pile.

Check using bisect.left and find we can place it on index 1 so q =1 and length of piles len(piles) = 2

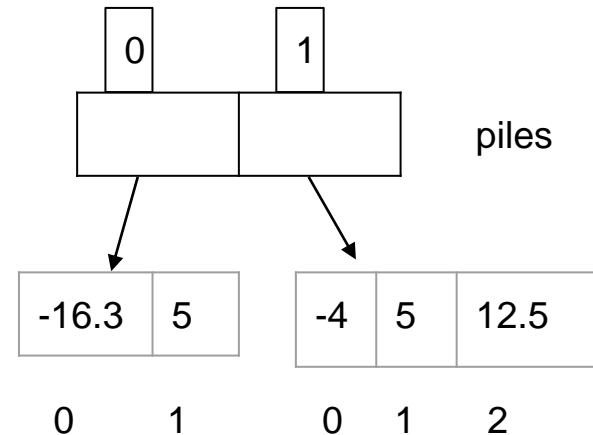
If loop will execute and I insert data at piles[1].insert(0,5)



```
1 import bisect
2 import heapq
3 def patienceSort(inparr):
4     piles = []
5     for p in inparr:
6         new_pile = [p]
7         q = bisect.bisect_left(piles, new_pile)
8         if q != len(piles):
9             piles[q].insert(0, p)
10        else:
11            piles.append(new_pile)
12
23 input = [5,12.5,-16.3,5,-4,6]
24 patienceSort(input)
```

Move on to p denotes to -4 of inparr
We store -4 in new_pile

Check in bisect we get $q = 1$
 $\text{len}(\text{piles}) = 2$ so they are not equal
If condition will execute and i will insert -4 into $\text{piles}[1].\text{insert}(0, -4)$ at 0th index



```
1 import bisect
2 import heapq
3 def patienceSort(inparr):
4     piles = []
5     for p in inparr:
6         new_pile = [p]
7         q = bisect.bisect_left(piles, new_pile)
8         if q != len(piles):
9             piles[q].insert(0, p)
10        else:
11            piles.append(new_pile)
12
23 input = [5, 12.5, -16.3, 5, -4, 6]
24 patienceSort(input)
```


Move on to the last number 6

`new_pile = 6`

Check `bisect_left`, we get `q = 2`

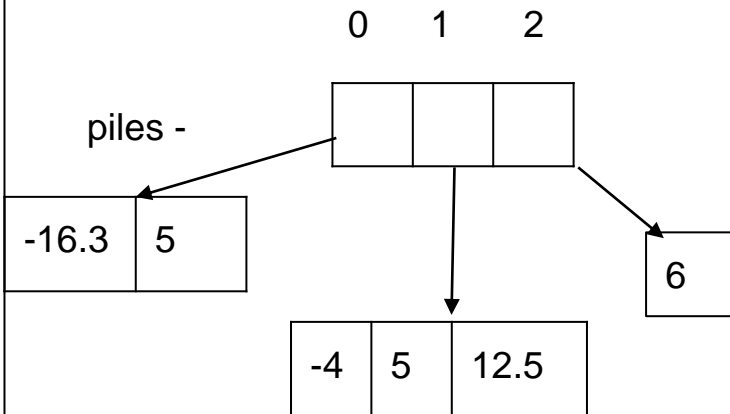
And `len(piles) = 2`

So if condition wont check out

We go to else loop and append the

`new_pile = 6` at the end of piles

So we get



```
1 import bisect
2 import heapq
3 def patienceSort(inparr):
4     piles = []
5     for p in inparr:
6         new_pile = [p]
7         q = bisect.bisect_left(piles, new_pile)
8         if q != len(piles):
9             piles[q].insert(0, p)
10        else:
11            piles.append(new_pile)
12
23 input = [5,12.5,-16.3,5,-4,6]
24 patienceSort(input)
```

There are no more elements in the list the first loop terminates and the pile creation is done.

Now for the sorting part.....

```
12
13     for q in range(len(inparr)):
14         small_pile = piles[0]
15         inparr[q] = small_pile.pop(0)
16         if small_pile:
17             heapq.heapreplace(piles, small_pile)
18         else:
19             heapq.heappop(piles)
20     assert not piles
21
```

Here we are using `heapq.heapreplace` and `heapq.heappop()`

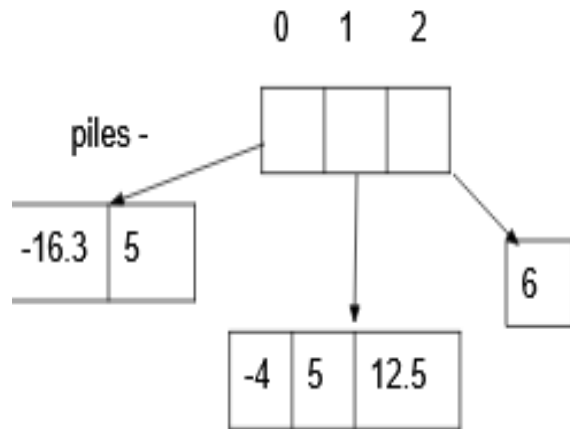
The `heapq` module which applies the priority queue algorithm where and follows minheap property i.e smallest element is at the root.

`Heapq.heapreplace` is used here to return the smallest item form heap, here heap is `piles` and to push the item on heap for in this case is `small_pile` back on the heap

`Heap.pop` pushes the item to the heap

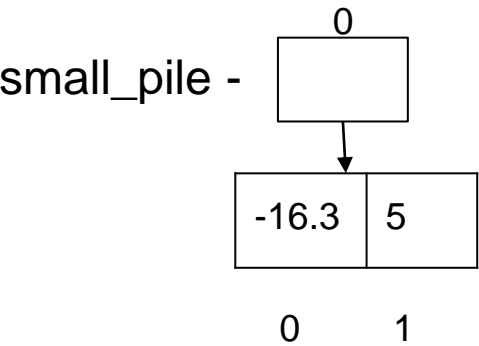
```
12
13 for q in range(len(inparr)):
14     small_pile = piles[0]
15     inparr[q] = small_pile.pop(0)
16     if small_pile:
17         heapq.heapreplace(piles, small_pile)
18     else:
19         heapq.heappop(piles)
20 assert not piles
21
```

We run a loop q equal to the length of our `inparr` list
We take a variable `small_pile` and store the data of the 0th index of `piles` i.e `piles[0]`

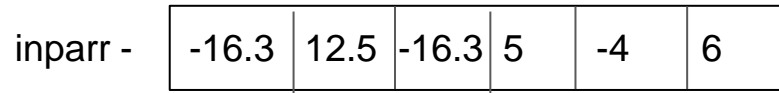
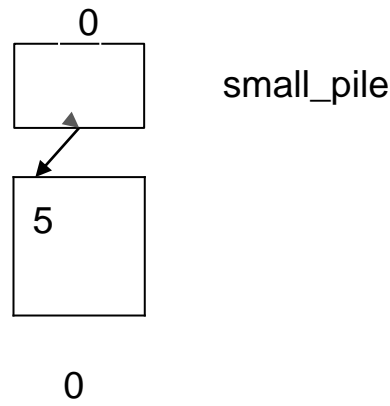


```
12
13 for q in range(len(inparr)):
14     small_pile = piles[0]
15     inparr[q] = small_pile.pop(0)
16     if small_pile:
17         heapq.heapreplace(piles, small_pile)
18     else:
19         heapq.heappop(piles)
20 assert not piles
21
```

So now small_pile has



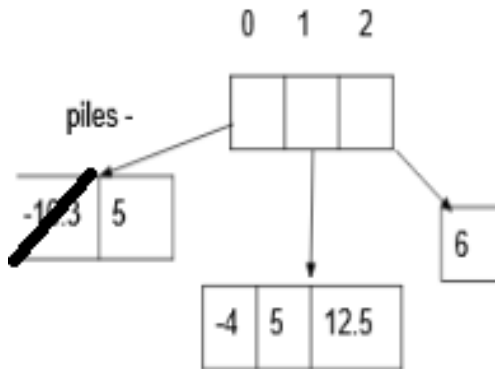
Now at `inparr[q]` where initially `q = 0` we store `small_pile.pop(0)` we remove the 0th index store it in the 0th index of `inparr`



```
12
13 for q in range(len(inparr)):
14     small_pile = piles[0]
15     inparr[q] = small_pile.pop(0)
16     if small_pile:
17         heapq.heapreplace(piles, small_pile)
18     else:
19         heapq.heappop(piles)
20 assert not piles
21
```

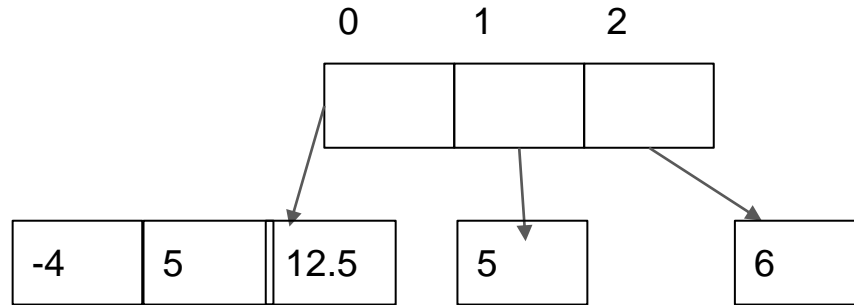
Now heapreplace finds the smallest item and locates in which pile it is

The next smallest item is in 1th index of piles



```
12
13+   for q in range(len(inparr)):
14       small_pile = piles[0]
15       inparr[q] = small_pile.pop(0)
16+       if small_pile:
17           heapq.heapreplace(piles, small_pile)
18+       else:
19           heapq.heappop(piles)
20   assert not piles
21
```

Now since we removed the first smallest digit the heap has changed and the new pile in index 0 will be

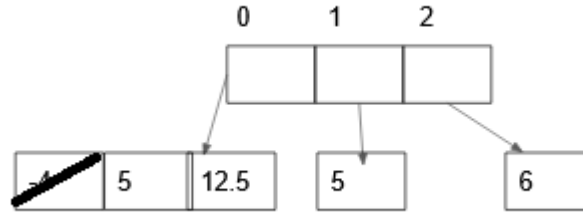


Now the smallest value is again in the root

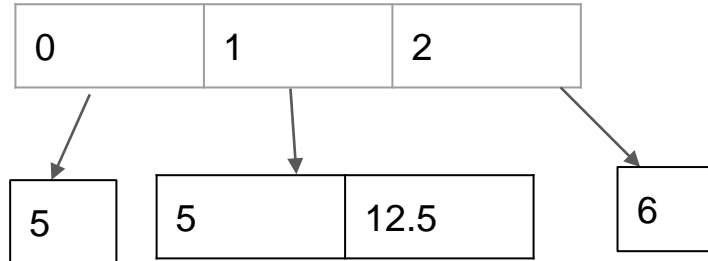
We then run the loop with remaining length and get

-16.3	-4	-16.3	5	-4	6
-------	----	-------	---	----	---

Now we again find the next smallest number in pile using `heapq.heapreplace` and find the least number to be in



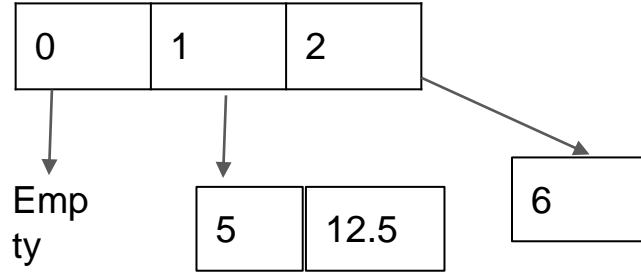
The 1th index of pile
We then bring it to root



We then again store the `pile[0]` in `small_pile` , pop `small_pile` and store in `inparr[q]`
Which gives us

-16.3	-4	5	5	-4	6
-------	----	---	---	----	---

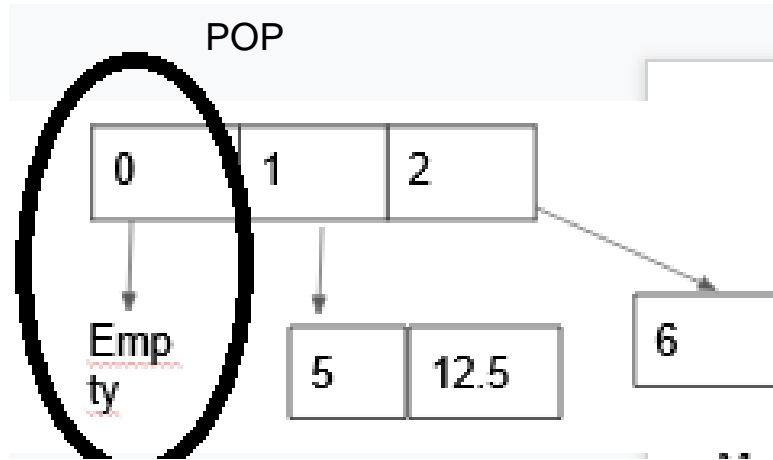
Moving on we get here an empty list now



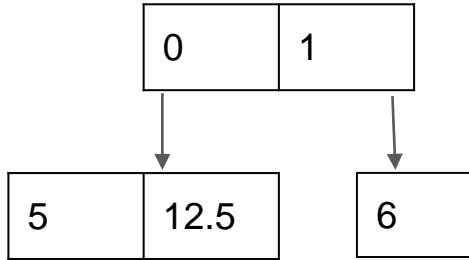
Since empty the else loop will initiate

```
if small_pile:
    heapq.heapreplace(piles, small_pile)
else:
    heapq.heappop(piles)
```

What this else loop will do is pop the empty list with 0th index



Then we will have



We thus continue this process till there are no more piles left and while we keep storing pop of small_pile in inp array we will get a sorted list

-16.3	-4	5	5	6	12.5
-------	----	---	---	---	------

```
IMPORT bisect and heapq
```

```
DEFINE FUNCTION sort(inputarray):
```

```
  SET piles TO [ ]
```

```
  FOR x IN inputarray:
```

```
    SET new_pile TO [x]
```

```
    SET i TO bisect.bisect_left(piles, new_pile)
```

```
    IF i != len(piles):
```

```
      piles[i].insert(0, x)
```

```
    ELSE:
```

```
      piles.append(new_pile)
```

```
  FOR i IN range(len(inputarray)):
```

```
    SET small_pile TO piles[0]
```

```
    SET inputarray[i] TO small_pile.pop(0)
```

```
    IF small_pile:
```

```
      heapq.heappreplace(piles, small_pile)
```

```
    ELSE:
```

```
      heapq.heappop(piles)
```

```
  CONTINUE till not piles
```

n times

logn times as it uses
binary search

nlogn

n times

O(1)

logn for priority queue

nlong

So the time complexity stands as $n\log n + n\log n$
 $= 2n\log n$
 $= O(n\log n)$

Points To note

- Not stable , does not maintain the order the input is given in it generally depends whether that number entered into the first / leftmost piles first
- It is inplace as at most the piles creation , the running in sequence and storing takes the same space as sample input
- It is online as it does not know what type of data it will deal with before executing
- Yes it is adaptive , depending on the sortedness or size of input ,the size or shape of piles and runtime may change
- Can work with positive , negative ,floating(neg and pos) , integer numbers (neg and pos)
- Follows greedy paradigm

Pros and Cons

Pros

- has running time of $n \log n$ and sometimes n when input is sorted
- used in finding longest increasing subsequence
- is used in playing the card game patience
- can sort data from a piles of data

Cons

- hard to implement
- too many variations and not worth the hassle
- better and easier codes available
- not that much used

Practical Uses

- This algorithm was used as process control according to wiki and further research led to this algorithm being used in a site called Bazaar as version control. Besides this according to a paper in microsoft a modified version of the patience sort is being used in data processing systems and applications - reference <https://www.microsoft.com/en-us/research/uploads/prod/2018/04/impatience-icde18.pdf>


```
1 import bisect
2 import heapq
3
4
5 def patience_sort(inparr):
6     piles = []
7     for p in inparr:
8         new_pile = [p]
9         q = bisect.bisect_left(piles, new_pile)
10        if q != len(piles):
11            piles[q].insert(0, p)
12        else:
13            piles.append(new_pile)
14
15    for q in range(len(inparr)):
16        small_pile = piles[0]
17        inparr[q] = small_pile.pop(0)
18        if small_pile:
19            heapq.heapreplace(piles, small_pile)
20        else:
21            heapq.heappop(piles)
22    assert not piles
23
24
25 input = [5, 12.5, -16.3, 5, -4, 6]
26 patience_sort(input)
27 print(input)
```

```
[-16.3, -4, 5, 5, 6, 12.5]
```

```
> |
```