

# Find Firsts

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- **Introduction**

In syntax analysis, the concept of "First Set" or "Firsts" is crucial in determining the first terminal symbol(s) that can be derived from a given non-terminal symbol or sequence of symbols in a grammar. The First Set plays a fundamental role in parsing and building the parser tables in *compiler design*.

- **Purpose**

The purpose of the "Find Firsts" algorithm is to compute the First Set for each non-terminal symbol in the grammar. By calculating the First Sets, the parser can determine the viable options for expanding non-terminal symbols during parsing.

## \* Algorithm

1. Initialize the First Set for all terminal symbols as themselves.
2. For each non-terminal symbol A:
  - If  $A \rightarrow \epsilon$  is a production rule, add  $\epsilon$  to  $\text{First}(A)$ .
  - For each production rule  $A \rightarrow X_1X_2...X_k$ :
    - If  $\epsilon$  belongs to  $\text{First}(X_1), \text{First}(X_2), \dots, \text{First}(X_i)$  for  $1 \leq i < k$ , add  $\text{First}(X_i)$  to  $\text{First}(A)$ .
    - If  $\epsilon$  does not belong to  $\text{First}(X_j)$  for any  $1 \leq j < k$ , add  $\text{First}(X_k)$  to  $\text{First}(A)$ , and stop further expanding  $\text{First}(A)$ .
3. Repeat step 2 until no more symbols can be added to any First Set.

## • Example

Consider the following grammar:

...

$S \rightarrow AB$

$A \rightarrow a \mid \epsilon$

$B \rightarrow b \mid \epsilon$

...

**The First Set calculations:**

**- First(S) = { a }**

**- First(A) = { a,  $\epsilon$  }**

**- First(B) = { b,  $\epsilon$  }**

- **Conclusion**

In conclusion, the **"Find Firsts"** algorithm is a key component in syntax analysis for compilers. By accurately computing the First Sets for **non-terminal symbols**, the parser can efficiently determine the next steps in parsing input code. Understanding and implementing this algorithm is essential for building a robust and efficient parser in **compiler design**.

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