



Header

Location: *fsptr

Header contains needed information about the filesystem and root directory attributes.

Will implement an integrity check to see if the filesystem has been initialized.

```
struct _fsHeader{
    size_t fsSize;
    size_t blockSize;
    int totalBlocks;
    int freeBlocks;
    fileAttr rtAttr;
    block_t firstFree;
    block_t lastFree;
}
```

```
struct _fileAttr{
    offset_t name;
    block_t iNode;
    off_t size;
    nlink_t nlink;
    time_t atim;
    time_t mtim;
}
```

Root Directory

Location: *fsptr + blockSize

Type: _dirBlock

Directory Block

A directory block holds file attributes for files and directories within that directory. If a directory contains more than n files, a new block is allocated of struct _dirOverflowBlock, which looks just like the _dirBlock with only the files array and overflow pointer; creating a linked list of directory blocks.



```
struct _dirBlock{
    struct _fileAttr files[(blockSize-(2*sizeof(block_t))-
        sizeof(int))/sizeof(_fileAttr)];
    int fileCount;
    block_t overflow = 0;
    block_t filenames = 0;
}

struct _dirOverflowBlock{
    struct _fileAttr files[(blockSize-sizeof(block_t))/sizeof(_fileAttr)];
    block_t overflow = 0;
}
```

Free Space

Free space will be managed using a linked list located on the free space itself. Each free block contains a struct _freeBlock, which has the number of the next free block.

Upon initialization, each block will simply point to its adjacent block.

Free space is allocated to files by block.

When a block is freed, a _freeBlock struct is created on that block and the linked list is updated accordingly.

Free blocks are zeroed out upon allocation.

```
struct _freeBlock{
    block_t next;
}
```

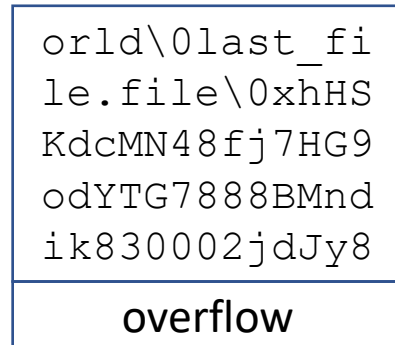
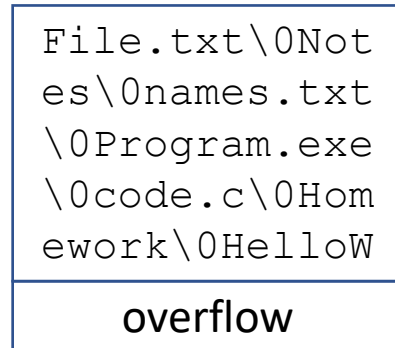
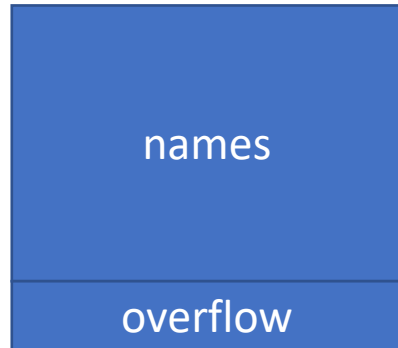
Filename

Each Directory Block contains an offset to a block containing a char array of names separated by \0 characters. The filesystem uses this block in conjunction with the offset in a file's attributes to find a file's name.

Some housekeeping must be done when files are renamed or removed.

The last bytes of the block gives an offset to a new block where the string is continued; creating a linked list of name blocks.

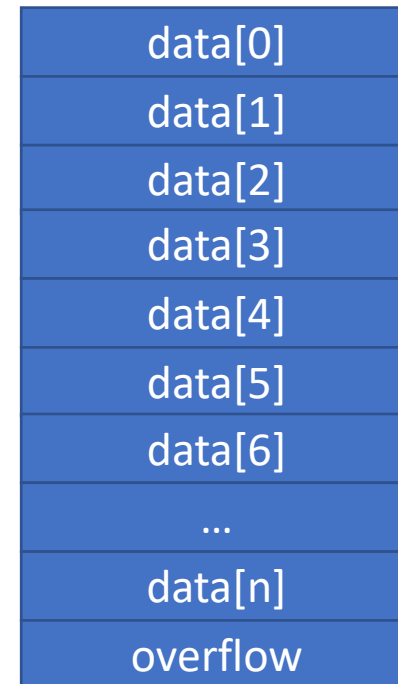
```
struct _filenameBlock{
    char names[blocksize-sizeof(block_t)];
    block_t overflow = 0;
```



Files

Directory blocks contain file attributes which contain the starting I-Node Block for each file. The file's I-Node Block has a data array containing offsets to all blocks the file uses (in sequential order). If the data array is full, the overflow offset designates a block where the array is continued; creating a linked list of i-nodes.

```
struct _fileBlock{
    offset_t data[(blkSize-sizeof(_fileBlock*))];
    block_t overflow = 0;
```



File System Attributes

Block Size = 4096 bytes = 4KB

Max filename length = N/A

Allocating free block = O(1)

Freeing used block = O(n)

Finding read/write offsets = O(n)

Locating filenames = O(n)