

Problem Solving By Computer - Project 1

1 Cubic Taxicab Number

Definition 1.1. A **cubic taxicab number** is a positive integer that can be expressed as the sum of two positive cubic numbers in at least two distinct ways.

Example 1.2. The first cubic taxicab number is 1729, since $1729 = 1^3 + 12^3 = 9^3 + 10^3$.

Problem 1.3. Given a positive integer N , how can we determine the smallest cubic taxicab number greater than or equal to N ? In particular, how can we write a function `CubicTaxicabNum(N)` such that `ctn = CubicTaxicabNum(N)` returns the smallest cubic taxicab number greater or equal to N ?

1.1 Approach

Suppose that we have a function `IsCubicTaxicabNum(N)` that determines whether a positive integer N is a cubic taxicab number, then we are able to deduce the smallest cubic taxicab number greater or equal to N by checking the integers greater than or equal to N in turn, until we find a cubic taxicab number.

```
function ctn = CubicTaxicabNum(N)
% CUBICTAXICABNUM    Returns the smallest cubic taxicab number greater
%                    than or equal to N

ctn = N;
while (~IsCubicTaxicabNum(ctn))
    ctn = ctn + 1;
end
end
```

All that remains to do is to implement the function `IsCubicTaxicabNum(N)`, such that the function returns 1 if N is a taxicab number and 0 otherwise.

1.2 Is it a Cubic Taxicab Number?

Before looking at the implementation of the function `IsCubicTaxicabNum(N)`, we will first make some observations that will help us derive an implementation for `IsCubicTaxicabNum(N)`.

Claim 1.4. Suppose that t is a positive integer such that $t = x^3 + y^3$, where x, y are positive integers, then $x, y \leq \lfloor \sqrt[3]{t} \rfloor$.

Proof. Let x, y, t be positive integers such that $t = x^3 + y^3$. Suppose by contradiction, we have $x > \lfloor \sqrt[3]{t} \rfloor$. Note that $\lceil \sqrt[3]{t} \rceil \geq \sqrt[3]{t}$. As x is an integer, we must have $x \geq \lceil \sqrt[3]{t} \rceil$. Together, we have $x \geq \lceil \sqrt[3]{t} \rceil \geq \sqrt[3]{t}$ and so $x^3 \geq \lceil \sqrt[3]{t} \rceil^3 \geq t$. In particular, for $t = x^3 + y^3$ to hold, we must have $t = x^3$, but this means that we must have $y = 0$, contradicting the choice of y being a positive integer. \square

Observation 1.5. Suppose that t is a positive integer such that $t = x^3 + y^3$, where x, y are positive integers, then there is no positive integer $a \neq x$, such that $t = a^3 + y^3$. Similarly, there is no positive integer $b \neq y$, such that $t = x^3 + b^3$.

Proof. Let a, x, y, t be positive integers such that $x^3 + y^3 = t$ and $a^3 + y^3 = t$. We have $x^3 + y^3 = t = a^3 + y^3$, and so $x^3 = a^3$. As both x and a are positive integers, we must have $x = a$. By symmetry, the latter statement also holds. \square

Observation 1.6. Suppose that t is a positive integer such that $x^3 + y^3 < t$, where x, y are positive integers, with $x \leq y$. Then, there is no positive integer $z < y$, with $x^3 + z^3 = t$.

Proof. Let x, y, t be positive integers with $x \leq y$ and $x^3 + y^3 < t$. Let z be any positive integer with $z < y$, then we must have $z^3 < y^3$. It follows that $x^3 + z^3 < x^3 + y^3 < t$ and so the result holds. \square

Observation 1.7. Suppose that t is a positive integer such that $x^3 + y^3 > t$, where x, y are positive integers, with $x \leq y$. Then, there is no integer $z > x$, with $z^3 + y^3 = t$.

Proof. Let x, y, t be positive integers with $x \leq y$ and $x^3 + y^3 > t$. Let z be any positive integer with $z > x$, then we must have $z^3 > x^3$. It follows that $z^3 + y^3 > x^3 + y^3 > t$ and so the result holds. \square

1.2.1 Implementation

The implementation of IsCubicTaxicabNum(N) in MATLAB is as follows.

```

function isTaxicab = IsCubicTaxicabNum(N)
% ISCUBICTAXICABNUM Returns isTaxicab = 1, if N is a taxicab number
% Returns isTaxicab = 0, otherwise

nPairs = 0;
left = 1;
right = floor(nthroot(N, 3)); % Claim 1.4
while (left ≤ right)
    total = left^3 + right^3;
    if (total == N)
        nPairs = nPairs + 1;
        if (nPairs == 2) % If we have found the
            isTaxicab = true; % second pair, we are done
            return;
        end
        left = left + 1; % Observation 1.5
        right = right - 1; % Observation 1.5
    elseif (total < N)
        left = left + 1; % Observation 1.6
    else
        right = right - 1; % Observation 1.7
    end
end
isTaxicab = false;
end

```

1.3 Analysis

Example 1.8. We can verify that the first cubic taxicab number is indeed 1729 by calling CubicTaxicabNum(1).

Example 1.9. The smallest cubic taxicab number greater or equal to 36032 is 39312. Indeed, 39312 is a cubic taxicab number since $39312 = 2^3 + 34^3 = 15^3 + 33^3$.

1.3.1 Complexity

The implementation of IsCubicTaxicabNum(N) makes use of the *two-pointer technique*, in which we have two pointers, `left` and `right`, and we insist that `left ≤ right`. We compute the sum of the cubes, $S = \text{left}^3 + \text{right}^3$ and if $S < N$, then we increment `left` by 1, and if $S > N$, then we decrement `right` by 1. If $S = N$, then we increment `left` by 1 and decrement `right` by 1. It follows that we have at most $\sqrt[3]{N}$ iterations of the **while** loop, and the operations inside the **while** loop run in constant time, so the time complexity of IsCubicTaxicabNum(N) is $O(\sqrt[3]{N})$.

2 Catalan's Constant

Definition 2.1. The **Catalan's constant** is a mathematical constant named after Eugène Charles Catalan, and is defined as

$$G = \sum_{k=0}^{\infty} \frac{(-1)^k}{(2k+1)^2} = \frac{1}{1^2} - \frac{1}{3^2} + \frac{1}{5^2} + \cdots \approx 0.915965594177219$$

Although the Catalan's constant G can be expressed in terms of the above sum of series, it is not known whether G is irrational or not.

Problem 2.2. In many situations, it would be more convenient to approximate the Catalan's constant as a ratio of two positive integers. Given a positive integer N , what is the best rational approximation p/q of the Catalan's constant, subject to the constraint $p + q \leq N$. In particular, how can we write a function `RatAppCat(N)` such that `[p, q] = RatAppCat(N)` returns the pair of integers p and q for the best rational approximation p/q of G , such that $p + q \leq N$?

2.1 Approximating Catalan's Constant

Observation 2.3. Given $G = 0.915965594177219$ and an integer q , we can find p such that $p/q = G$, by computing $p \approx Gq$. Of course, p needs not to be an integer; however, we can find the integer p_0 such that p_0/q is closest to G among all integers by $p_0 = \text{round}(p)$.

Claim 2.4. Suppose we have integers p_0, p_1, q_0, q_1 such that $p_0/q_0 = G$ and $p_1/q_1 = G$. If we have $q_1 > q_0$, then it must be the case $p_1 > p_0$.

Proof. Let p_0, p_1, q_0, q_1 be integers such that $p_0/q_0 = G$, $p_1/q_1 = G$ and $q_0 < q_1$. We have $p_0 = Gq_0$ and $p_1 = Gq_1$, but as $q_0 < q_1$, we have $Gq_0 < Gq_1$. It follows that $p_0 < p_1$. \square

2.1.1 Approach

As a result of Observation 2.3, we can enumerate over values of q in ascending order, computing the *perfect* integer p for each denominator q , and checking whether p/q is a better approximation than any previous pairs of (p, q) considered. Further, when we encounter a pair (p, q) such that $p + q > N$, then we are done. Indeed, as we are enumerating over values of q in ascending order, any subsequent pairs we consider, say (p_0, q_0) , we will have $q_0 > q_1$, and so $p_0 > p_1$ by Claim 2.4. In particular, $p_0 + q_0 > p + q > N$.

2.2 Implementation

The implementation of `RatAppCat(N)` in MATLAB is as follows.

```
function [p, q] = RatAppCat(N)
% RATAPPCAT    Returns the best approximation p / q of the Catalan's
%              constant, among all pairs of (p, q) such that p + q ≤ N.

G = 0.915965594177219;
p = 0; q = 1;                                % The best (p, q) pair so far
minDelta = abs(G - p / q);                    % The difference between p / q and G

for q0 = 1 : N
    p0 = round(G * q0);                       % Observation 2.3
    if (p0 + q0 > N)                           % Claim 2.4
        return
    end
    delta = abs(G - p0 / q0);
    if (delta < minDelta)                       % Update if current pair is better
        minDelta = delta;                     % than the best pair we have seen
        p = p0; q = q0;
    end
end
end
```

2.3 Analysis

Example 2.5. For $N = 2018$, the pair of integers 109 and 119 provides the best rational approximation of the Catalan's constant, with $109/119 = 0.915966386554622$.

2.3.1 Complexity

The implementation of `RatAppCat(N)` loops over possible denominators from 1 to N , and for each possible denominator, we compute the perfect numerator, and check whether this pair gives a better approximation than any previously encountered pairs. The operations in each iteration are run in constant time, and so the time complexity of `RatAppCat(N)` is $O(N)$, as the loop is run at most N times. The space complexity of `RatAppCat(N)` is $O(1)$, as we only store a constant number of variables.

3 Sum of Reciprocal Squares with Prime Factors

Definition 3.1. The **sum of reciprocal squares**, also known as the Basel problem, asks for the precise summation of the reciprocal squares of the natural numbers. That is,

$$\sum_{n=1}^{\infty} \frac{1}{n^2} = \frac{1}{1^2} + \frac{1}{2^2} + \frac{1}{3^2} + \frac{1}{4^2} + \dots$$

which was shown to be exactly $\pi^2/6$ by Leonhard Euler.

Definition 3.2. The **prime omega function (with multiplicity)** $\Omega(n)$ counts the number of prime factors of a natural number n , counting multiplicity. By convention, we have $\Omega(1) = 0$.

Example 3.3. $\Omega(12) = 3$ as $12 = 2^2 \cdot 3^1$, $\Omega(25) = 2$ as $25 = 5^2$ and $\Omega(47) = 1$ as 47 is a prime number. Moreover, $\Omega(p) = 1$ for any prime number p , as prime numbers only have one prime factor, by definition.

Definition 3.4. The **sum of reciprocal squares with prime factors** asks for the precise summation of the following sum.

$$\sum_{n=1}^{\infty} \frac{(-1)^{\Omega(n)}}{n^2} = \frac{1}{1^2} - \frac{1}{2^2} - \frac{1}{3^2} + \frac{1}{4^2} + \dots$$

Problem 3.5. How can we derive a reasonable approximation of the sum of the reciprocal squares with prime factors? In particular, can we make any observations about the answer by truncating a finite number of terms?

3.1 Approximating the Sum of Reciprocal Squares with Prime Factors

Observation 3.6. For any $a, b \in \mathbb{N}$ with $a < b$, we have

$$-\sum_{n=a}^b \frac{1}{n^2} \leq \sum_{n=a}^b \frac{(-1)^{\Omega(n)}}{n^2} \leq \sum_{n=a}^b \frac{1}{n^2} \quad \text{and} \quad -\sum_{n=a}^{\infty} \frac{1}{n^2} \leq \sum_{n=a}^{\infty} \frac{(-1)^{\Omega(n)}}{n^2} \leq \sum_{n=a}^{\infty} \frac{1}{n^2}$$

Observation 3.7. Let $x, y, z \in \mathbb{R}$ with $x \leq y \leq z$. Let $n \in \mathbb{N} \cup \{0\}$. Let **round**(a , n) be a function that rounds a to n decimal places. If **round**(x , n) = **round**(z , n) = c , then **round**(y , n) = c . That is, if x and z round to the same value for n decimal places, say c , then y rounded to n decimal places is also c .

Notation 3.8.

$$S(a, b) = \sum_{n=a}^b \frac{(-1)^{\Omega(n)}}{n^2} \quad \text{and} \quad T(a, b) = \sum_{n=a}^b \frac{1}{n^2}$$

We know that $T(1, \infty) = \pi^2/6$ and our goal is to approximate $S(1, \infty)$.

Observation 3.9. We can compute $T(k+1, \infty)$ by $T(k+1, \infty) = T(1, \infty) - T(1, k) = \pi^2/6 - T(1, k)$. That is, we can compute $T(k+1, \infty)$ by subtracting the first k terms from $\pi^2/6$.

3.1.1 Approach

Let $k \in \mathbb{N}$. It follows from Observation 3.6 that we have

$$S(1, k) - T(k+1, \infty) \leq \underbrace{S(1, k) + S(k+1, \infty)}_{S(1, \infty)} \leq S(1, k) + T(k+1, \infty)$$

In particular, for some $d \in \mathbb{N}$, if we are able to find k such that the above inequality holds, with the lower and upper bound rounding to the same value (up to d decimal places), then by Observation 3.7, this value is an accurate approximation for $S(1, \infty)$ up to d decimal places.

3.2 Implementation

3.2.1 Prime Omega Function

First, we need to implement the function `PrimeOmega(N)` such that `n = PrimeOmega(N)` returns the number of prime factors of N (counting multiplicity). In fact, we can make use of the `factor(N)` function provided by MATLAB, which returns a vector containing the prime factors of N . It follows that `PrimeOmega(N)` can be implemented by counting the length of the vector returned by `factor(N)` and can be implemented as an anonymous function, as follows.

```
PrimeOmega = @(x) (x ~= 1) * length(factor(x));
```

As `factor(1) = [1]`, we use `(x ~= 1)` to handle the case $\Omega(1) = 0$. That is, if $x \neq 1$, then `PrimeOmega(x)` returns `length(factor(x))`, otherwise it returns `0`.

3.2.2 Finding the value of k

Let d be the number of decimal places which we wish to get an approximation for $S(1, \infty)$. Let $k = 0$, $S = 0$ and $T = \pi^2/6$. We can enumerate over values of k in turn, adding $(-1)^{\Omega(k)}/k^2$ to S , and subtracting $1/k^2$ from T . If $S - T$ and $S + T$ round to the same value up to d decimal places, then we have an approximation for $S(1, \infty)$ up to d decimal places, by considering the first k terms.

3.2.3 Sum of Reciprocal Squares with Prime Factors

The implementation of SumPF in MATLAB is as follows.

```
function SumPF
% SUMPF    Finds an approximation of the sum of reciprocal squares
%          with prime factors

S = 0;
T = pi^2 / 6;
k = 0; d = 1;
PrimeOmega = @ (n) (n ~= 1) * length(factor(n));

fprintf('%-8s %-8s %8s\n', 'Terms', 'Value', 'Accuracy');
while (k < 1000000)
    if (round(S - T, d) == round(S + T, d))
        fprintf('%-8d %8f %-8d\n', k, round(S, d), d);
        d = d + 1;
    end
    k = k + 1;
    S = S + (-1)^PrimeOmega(k) / k^2;
    T = T - 1 / k^2;
end
end
```

3.3 Analysis

```
>> SumPF
Terms    Value    Accuracy
123      0.700000    1
334      0.660000    2
2102     0.658000    3
42353    0.658000    4
728565   0.657970    5
```

Example 3.10. By considering the first 2102 terms of the sum of reciprocal squares with prime factors, we get that $S = 0.6580$ to 4 decimal places. Similarly, we get that $S = 0.65797$ to 5 decimal places, by considering 728565 terms.

3.3.1 Alternate Approach

Our implementation of SumPF makes use of the sum of reciprocal squares, which was shown to be $\pi^2/6$. As a result, our implementation would not be suitable to compute an approximation for an arbitrary sum. Alternatively, we could derive an approach that deduces such bounds without knowledge of another summation.