



Network Layer

Part 8

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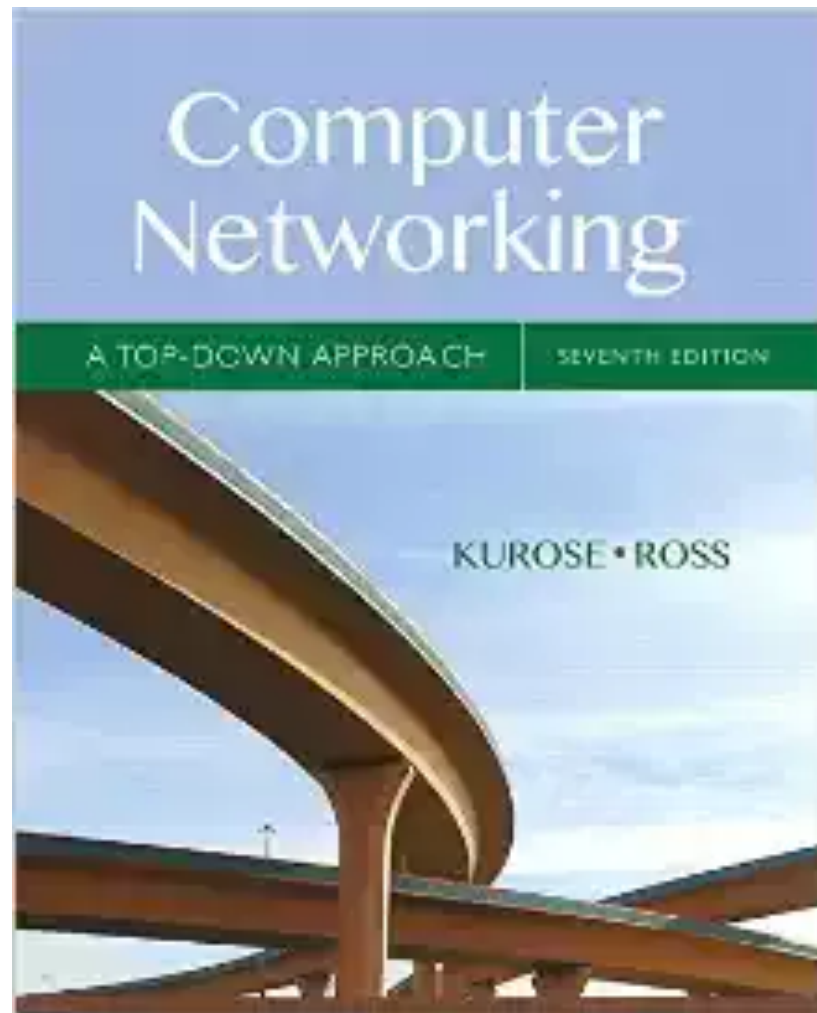
“my, my, my, ... these are lawless times ...”

These slides are more-or-less directly from the slide set developed by Jim Kurose and Keith Ross for their book “Computer Networking: A Top Down Approach, 5th edition”.

The slides have been lightly adapted for Mark Allman’s EECS 325/425 Computer Networks class at Case Western Reserve University.

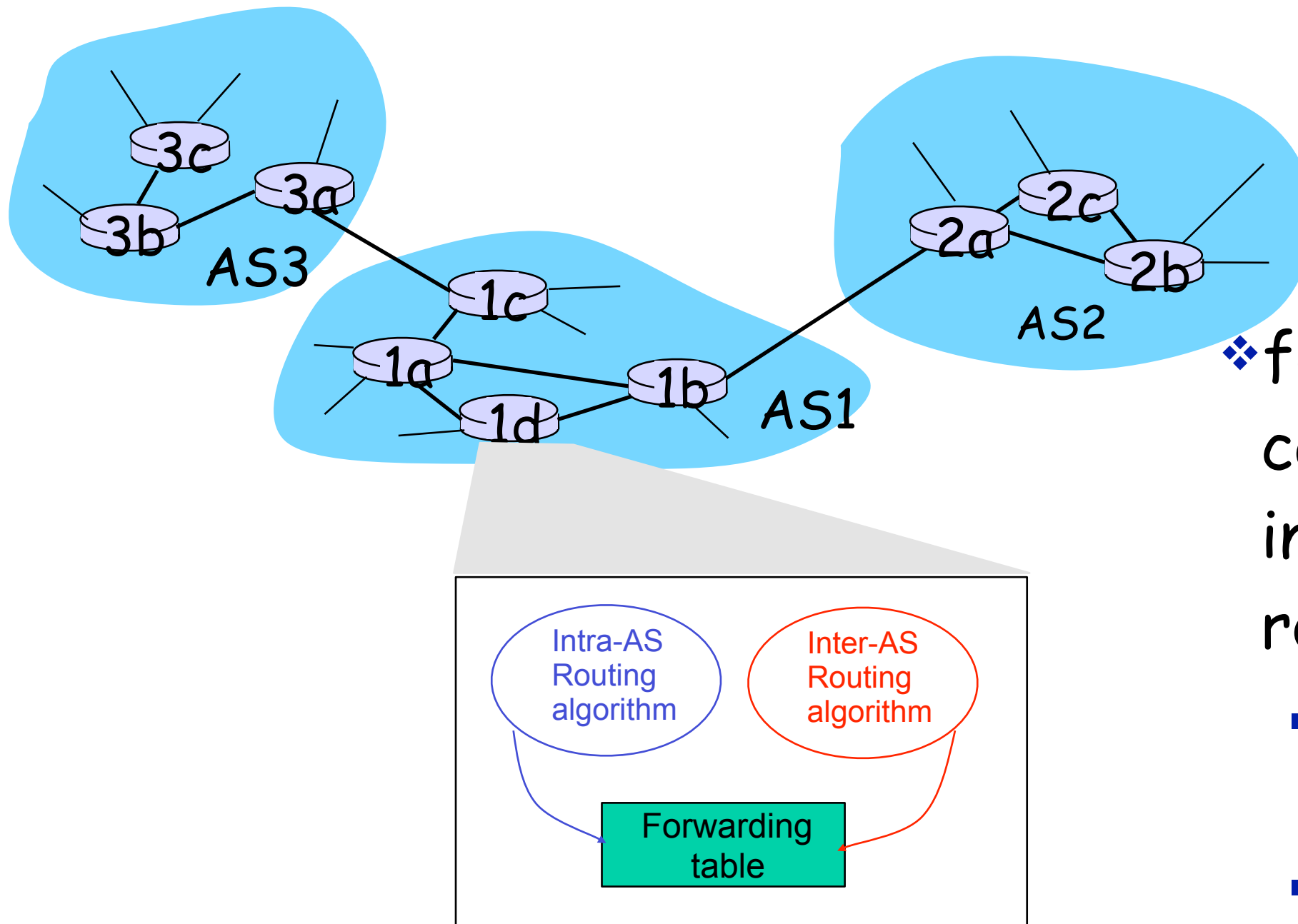
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Reading Along ...



- Network layer is chapters 4 & 5
- Hierarchical Routing

Interconnected ASes



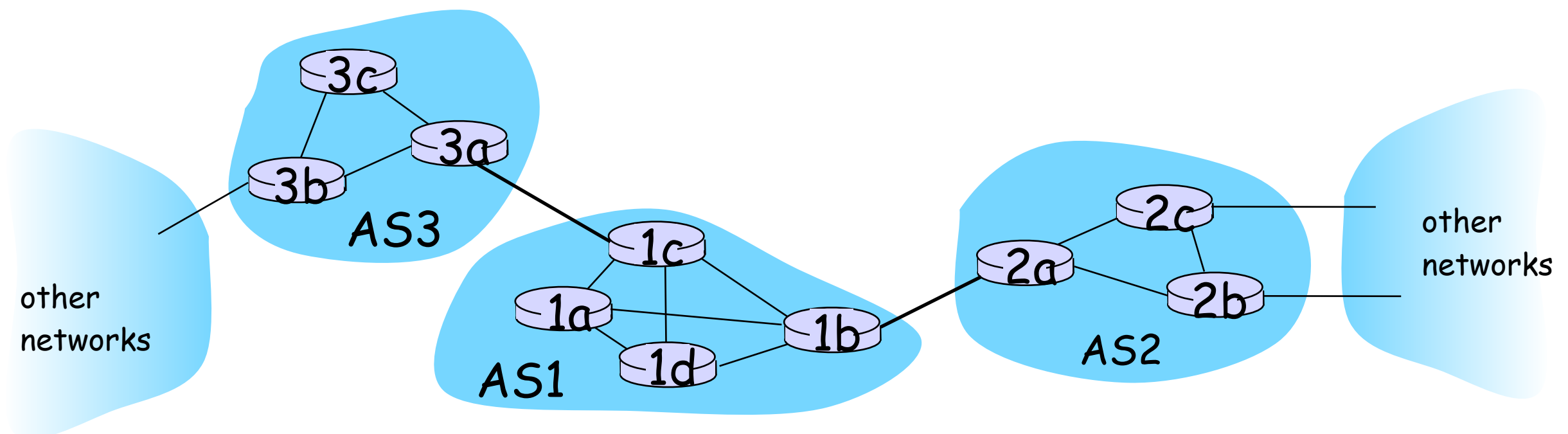
❖ forwarding table configured by both intra- and inter-AS routing algorithm

- intra-AS sets entries for internal dests
- inter-AS & intra-As sets entries for external dests

Inter-AS tasks

❖ suppose router in AS1 receives datagram destined outside of AS1:

- router should forward packet to gateway router, but which one?



Inter-AS tasks

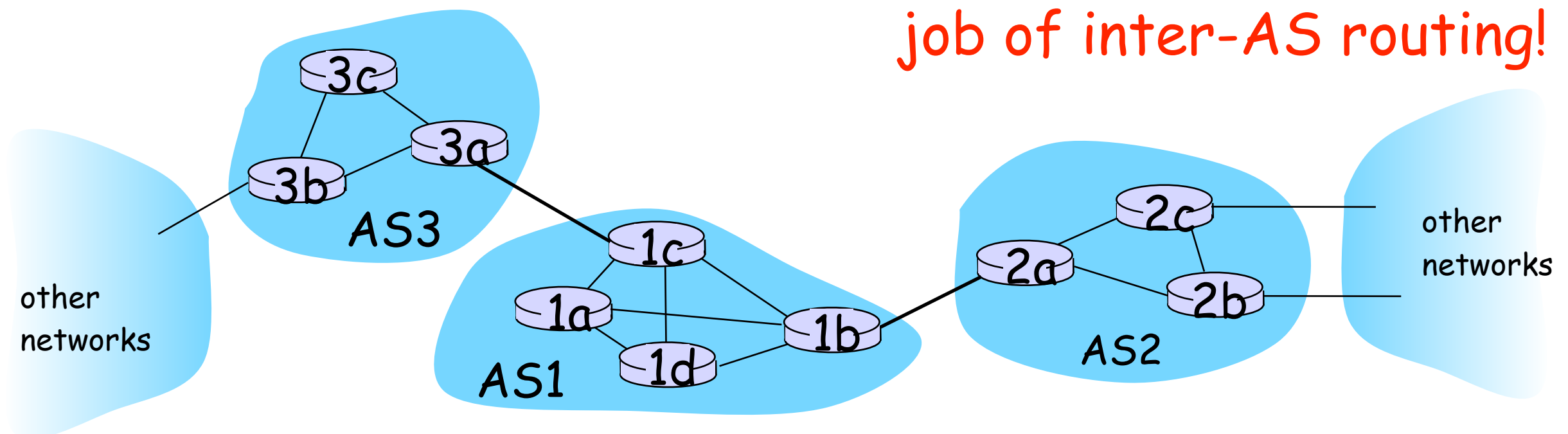
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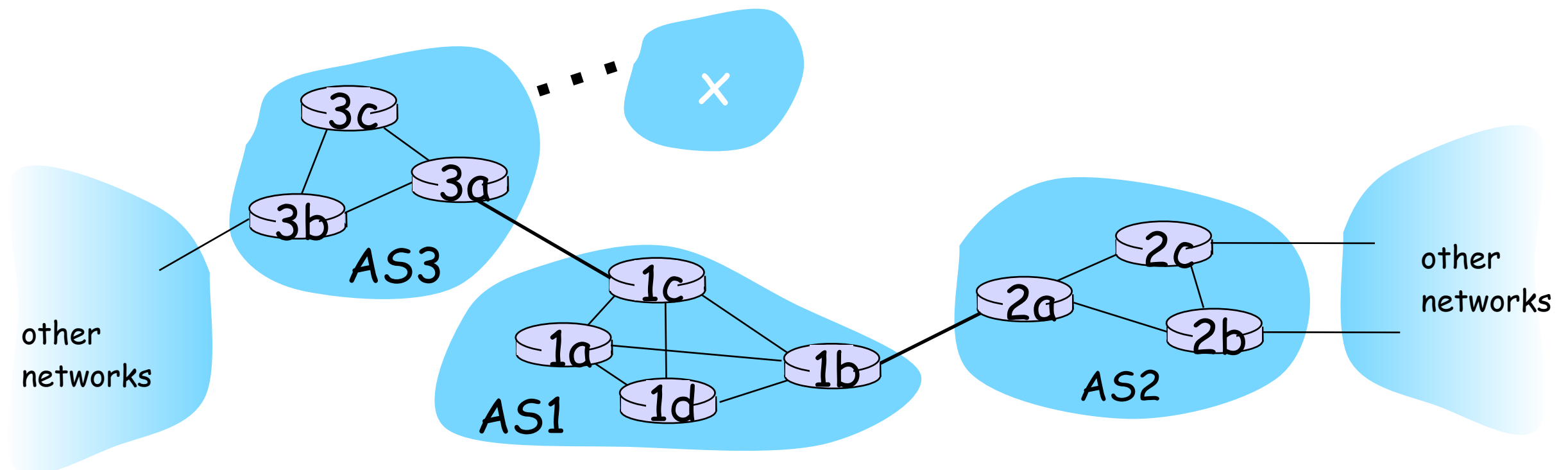
AS1 must:

1. learn which destds are reachable through AS2, which through AS3
2. propagate this reachability info to all routers in AS1

job of inter-AS routing!

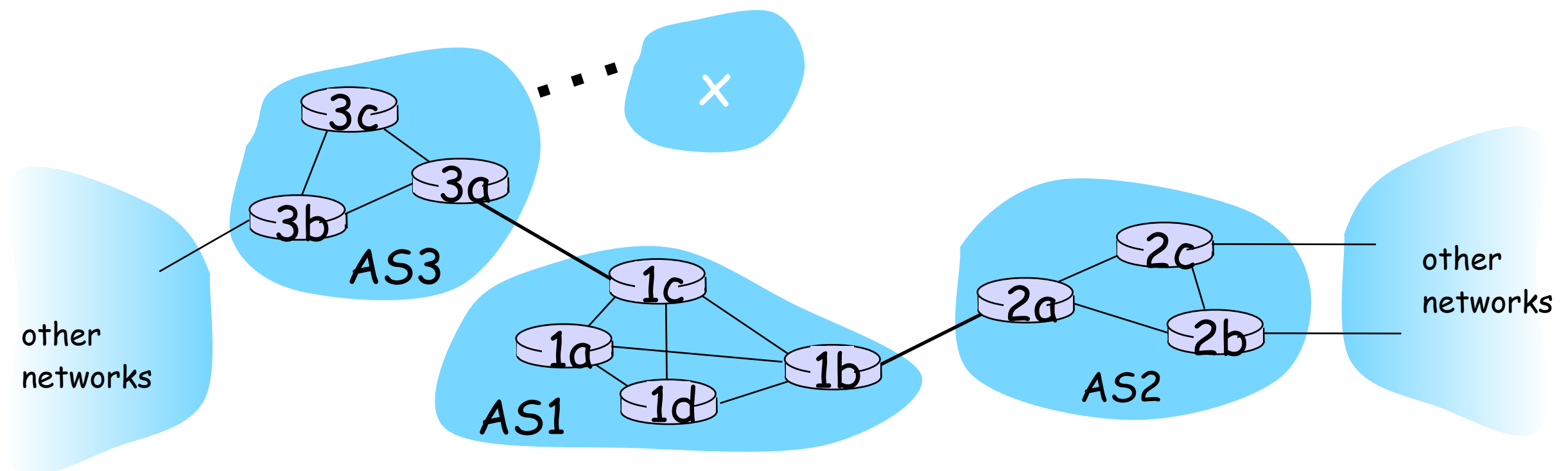


Example: Setting forwarding table in router 1d



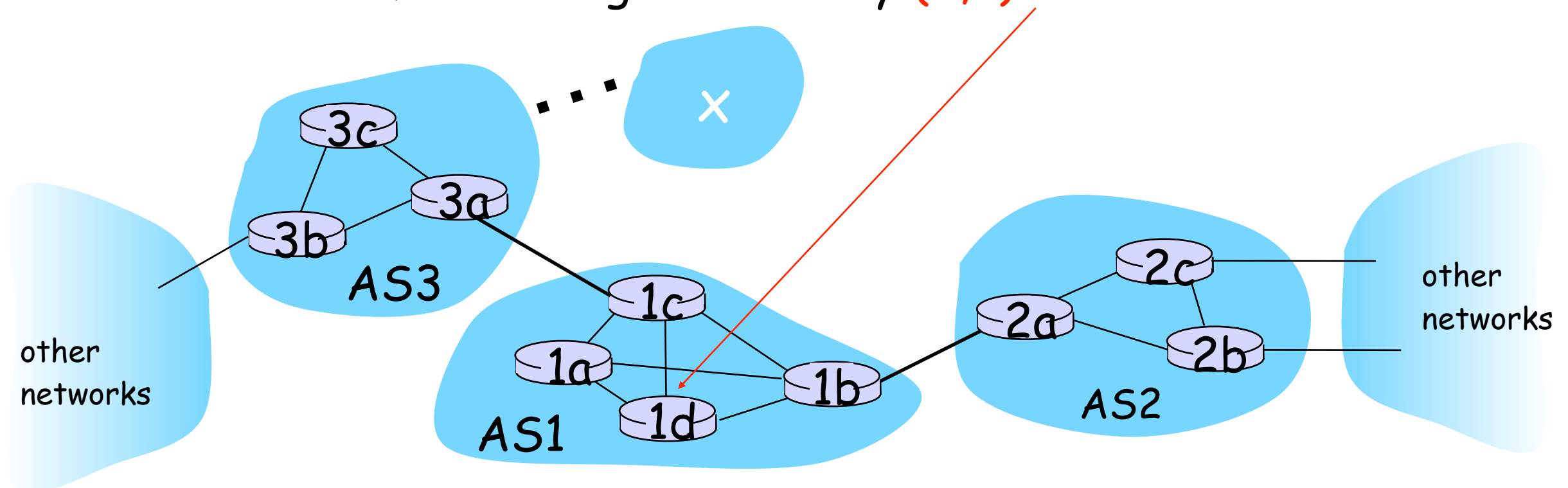
Example: Setting forwarding table in router 1d

- ❖ suppose AS1 learns (via inter-AS protocol) that subnet **x** reachable via AS3 (gateway 1c) but not via AS2.
 - inter-AS protocol propagates reachability info to all internal routers



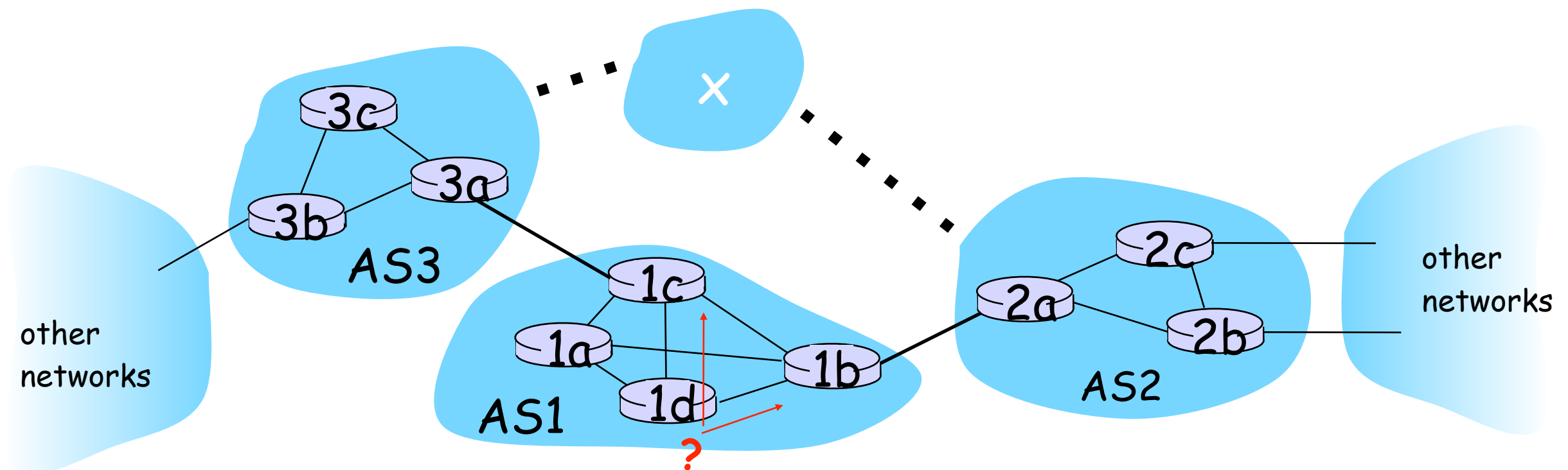
Example: Setting forwarding table in router 1d

- ❖ suppose AS1 learns (via inter-AS protocol) that subnet **x** reachable via AS3 (gateway 1c) but not via AS2.
 - inter-AS protocol propagates reachability info to all internal routers
- ❖ router 1d determines from intra-AS routing info that its interface **I** is on the least cost path to 1c.
 - installs forwarding table entry **(x,I)**



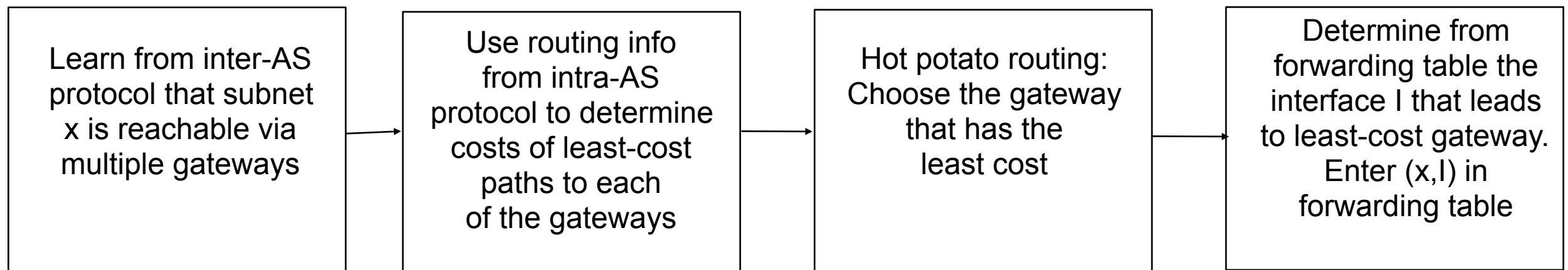
Example: Choosing among multiple ASes

- ❖ now suppose AS1 learns from inter-AS protocol that subnet **x** is reachable from AS3 and from AS2.
- ❖ to configure forwarding table, router 1d must determine which gateway it should forward packets towards for dest **x**
 - this is also job of inter-AS routing protocol!



Example: Choosing among multiple ASes

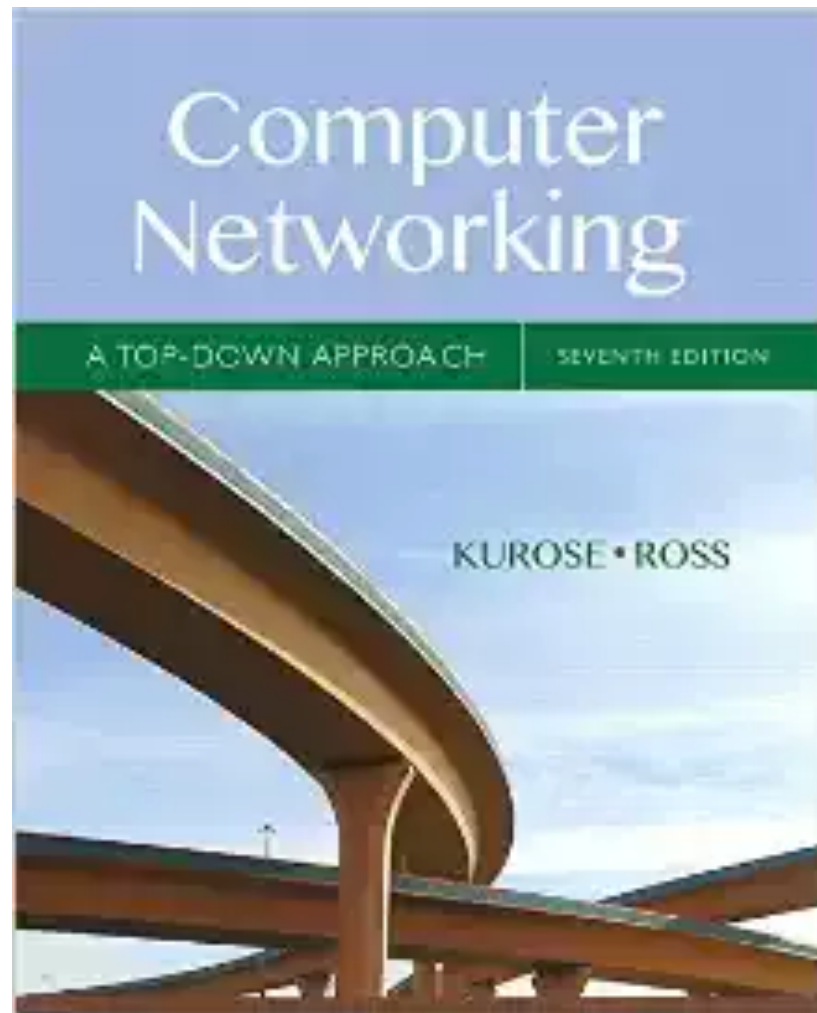
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- ❖ to configure forwarding table, router 1d must determine which gateway it should forward packets towards for dest **x**
 - this is also job of inter-AS routing protocol!
- ❖ **hot potato routing**: send packet towards closest of two routers.



Intra-AS Routing

- ❖ also known as **Interior Gateway Protocols (IGP)**
- ❖ most common Intra-AS routing protocols:
 - RIP: Routing Information Protocol
 - OSPF: Open Shortest Path First
 - IGRP: Interior Gateway Routing Protocol (Cisco proprietary)

Reading Along ...



- Network layer is chapters 4 & 5
- Routing Among ISPs

Internet inter-AS routing: BGP

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- ❖ **BGP (Border Gateway Protocol):** the de facto inter-domain routing protocol
 - “glue that holds the Internet together”

Internet inter-AS routing: BGP

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 - "glue that holds the Internet together"
- ❖ allows subnet to advertise its existence to rest of Internet: "I am here"

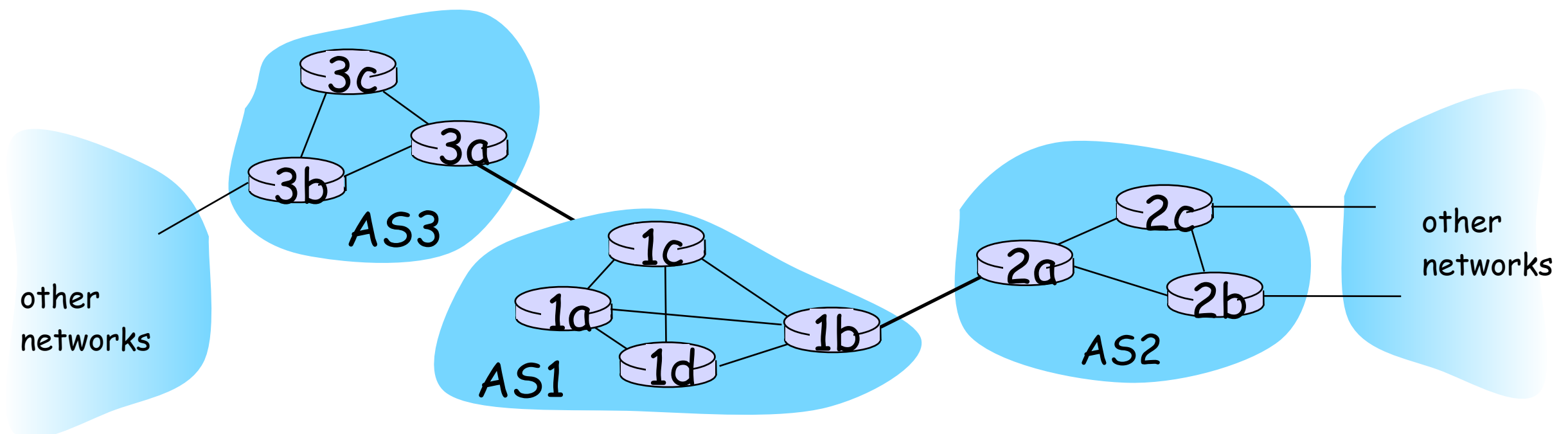
Internet inter-AS routing: BGP

- ❖ **BGP (Border Gateway Protocol):** the de facto inter-domain routing protocol
 - “glue that holds the Internet together”
- ❖ allows subnet to advertise its existence to rest of Internet: “I am here”
- ❖ BGP provides each AS a means to:
 - **eBGP:** obtain subnet reachability information from neighboring ASs.
 - **iBGP:** propagate reachability information to all AS-internal routers.
 - determine “good” routes to other networks based on reachability information and policy.

BGP basics

❖ **BGP session:** two BGP routers ("peers") exchange BGP messages:

- advertising **paths** to different destination network prefixes ("path vector" protocol)
- exchanged over semi-permanent TCP connections



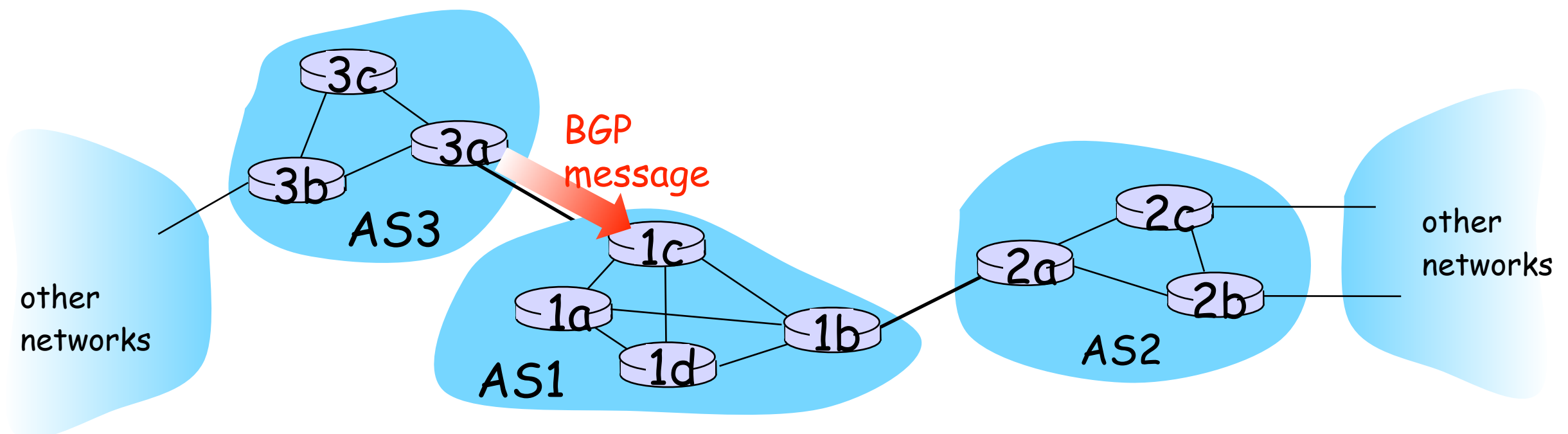
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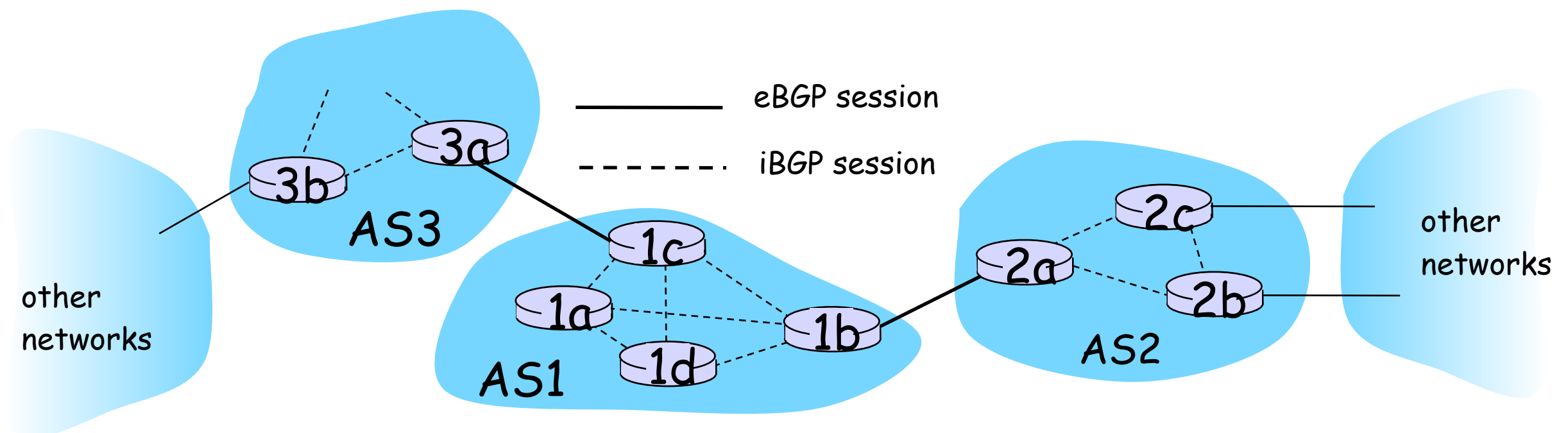
- advertising **paths** to different destination network prefixes ("path vector" protocol)
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❖ when AS3 advertises a prefix to AS1:

- AS3 **promises** it will forward datagrams towards that prefix
- AS3 can aggregate prefixes in its advertisement

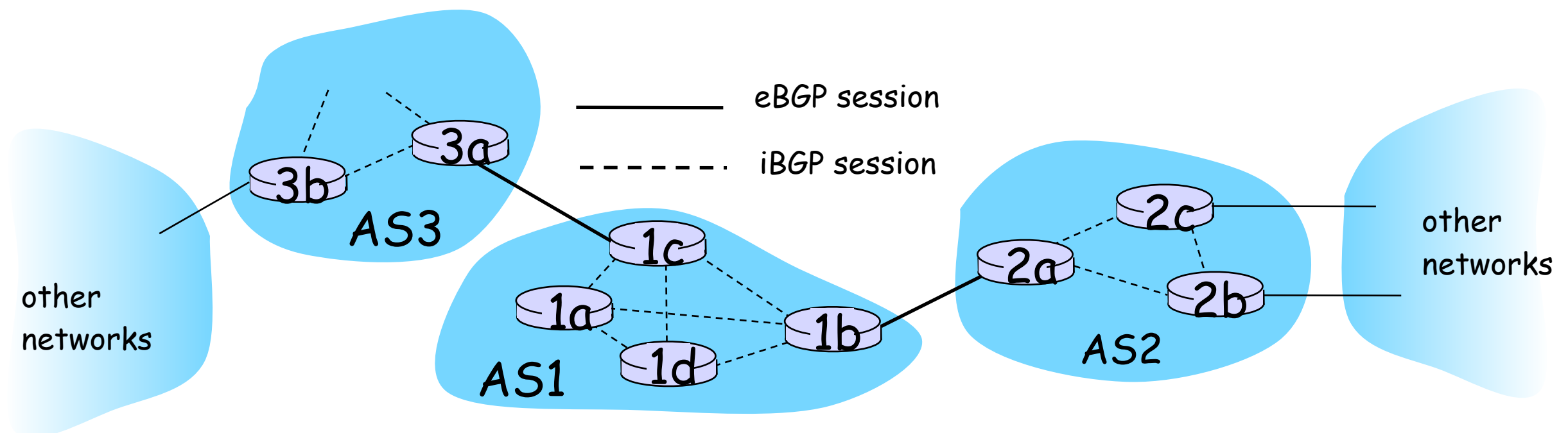


BGP basics: distributing path information



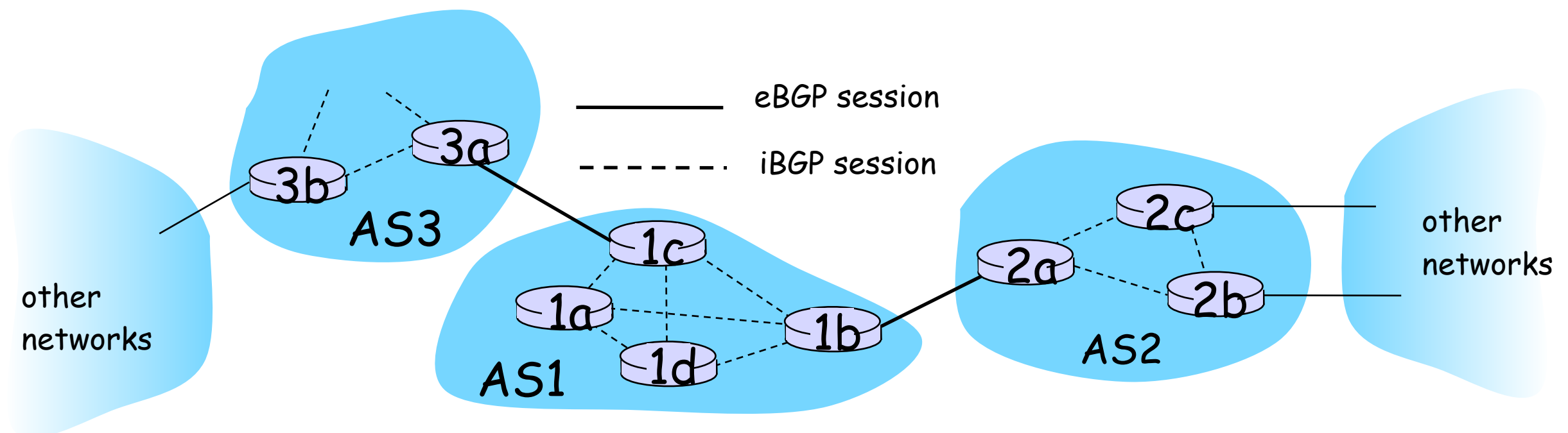
BGP basics: distributing path information

- ❖ using eBGP session between 3a and 1c, AS3 sends prefix reachability info to AS1.



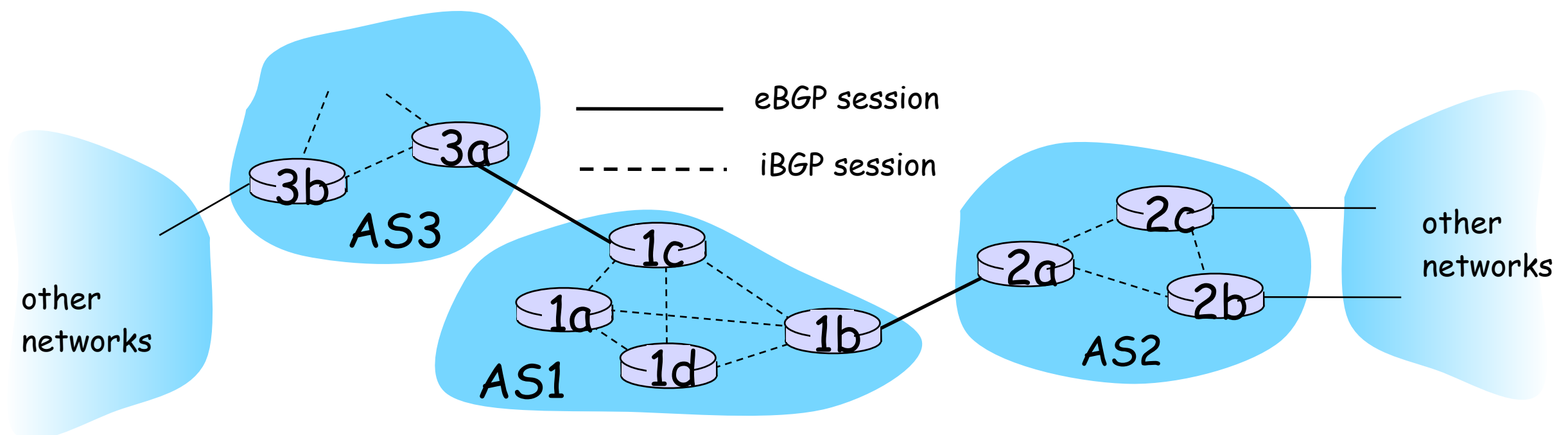
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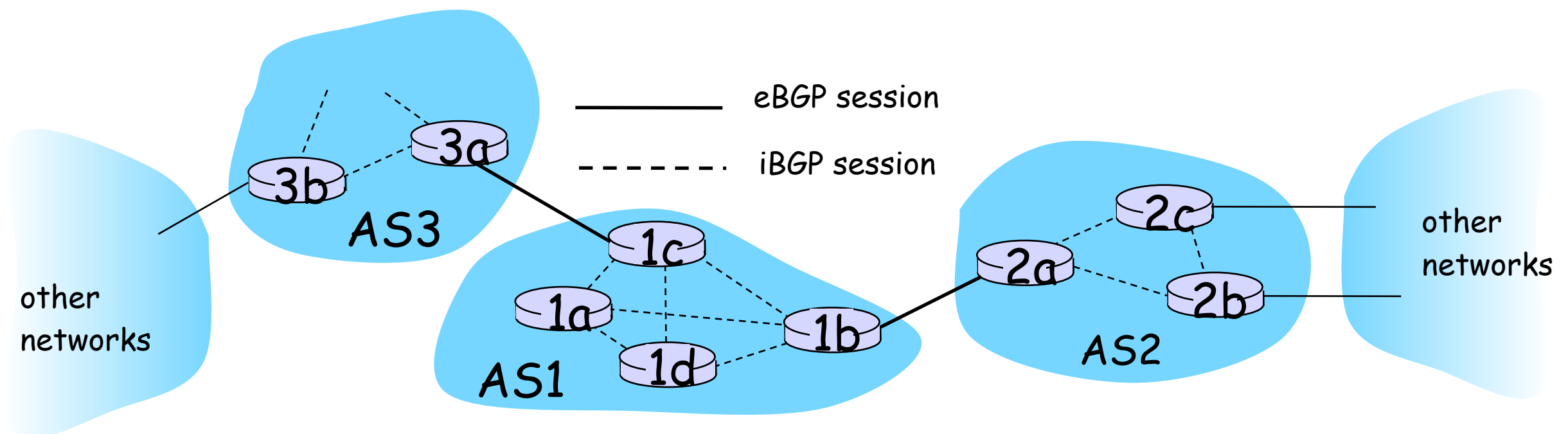
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BGP basics: distributing path information

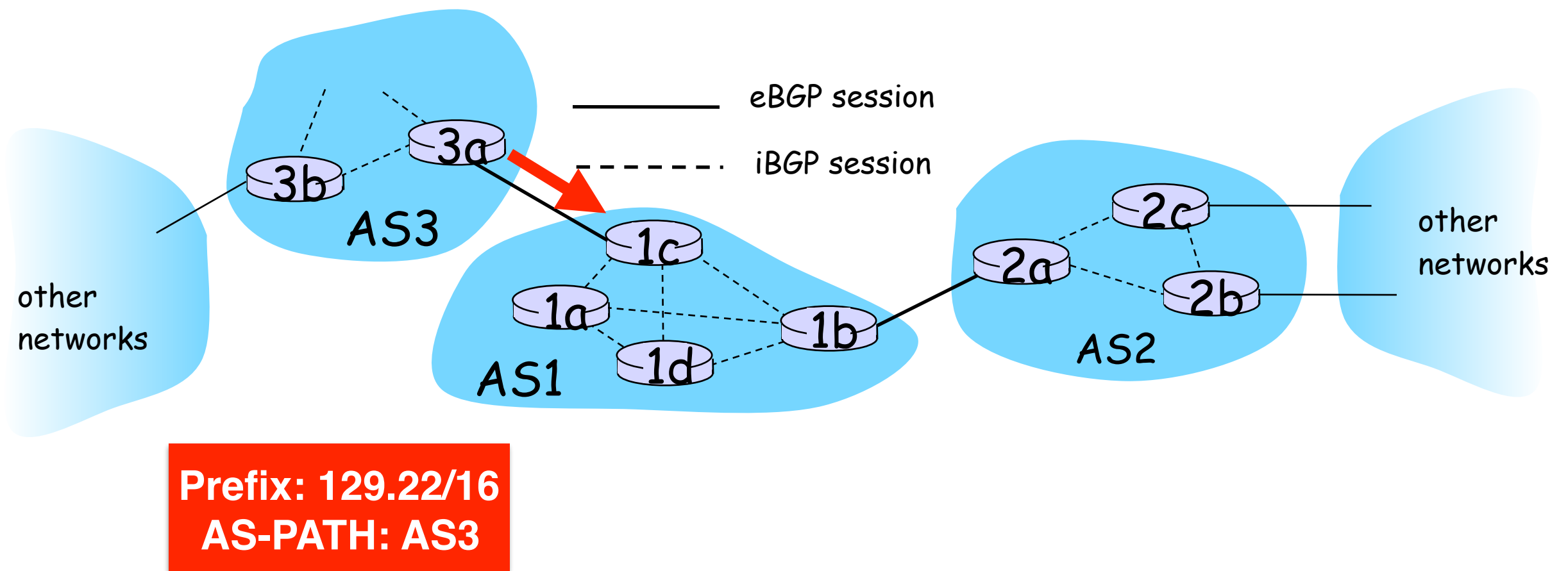
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 - 1c can then use iBGP to distribute new prefix info to all routers in AS1
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- ❖ when router learns of new prefix, it creates entry for prefix in its forwarding table.



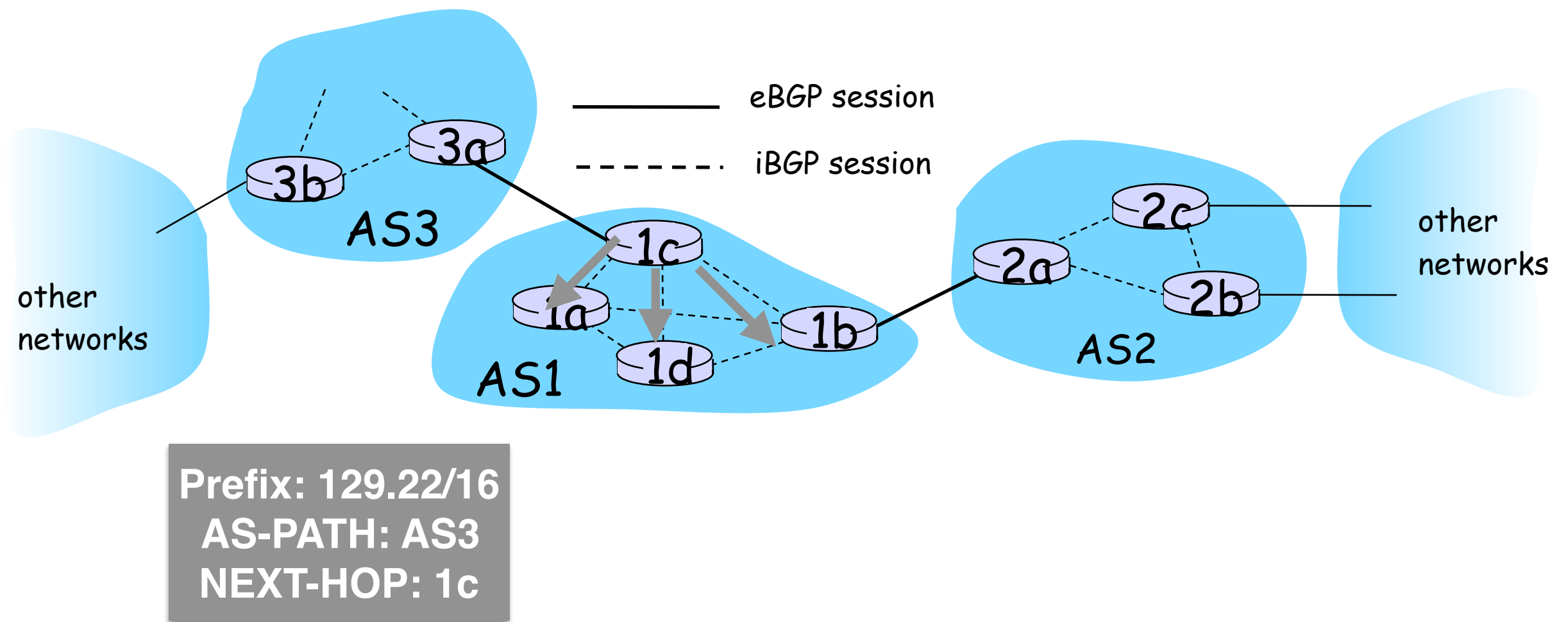
Path attributes & BGP routes

- ❖ advertised prefix includes BGP attributes
 - prefix + attributes = "route"
- ❖ two important attributes:
 - **AS-PATH**: contains ASs through which prefix advertisement has passed: e.g., AS 67, AS 17
 - **NEXT-HOP**: indicates specific internal-AS router to next-hop AS. (may be multiple links from current AS to next-hop-AS)

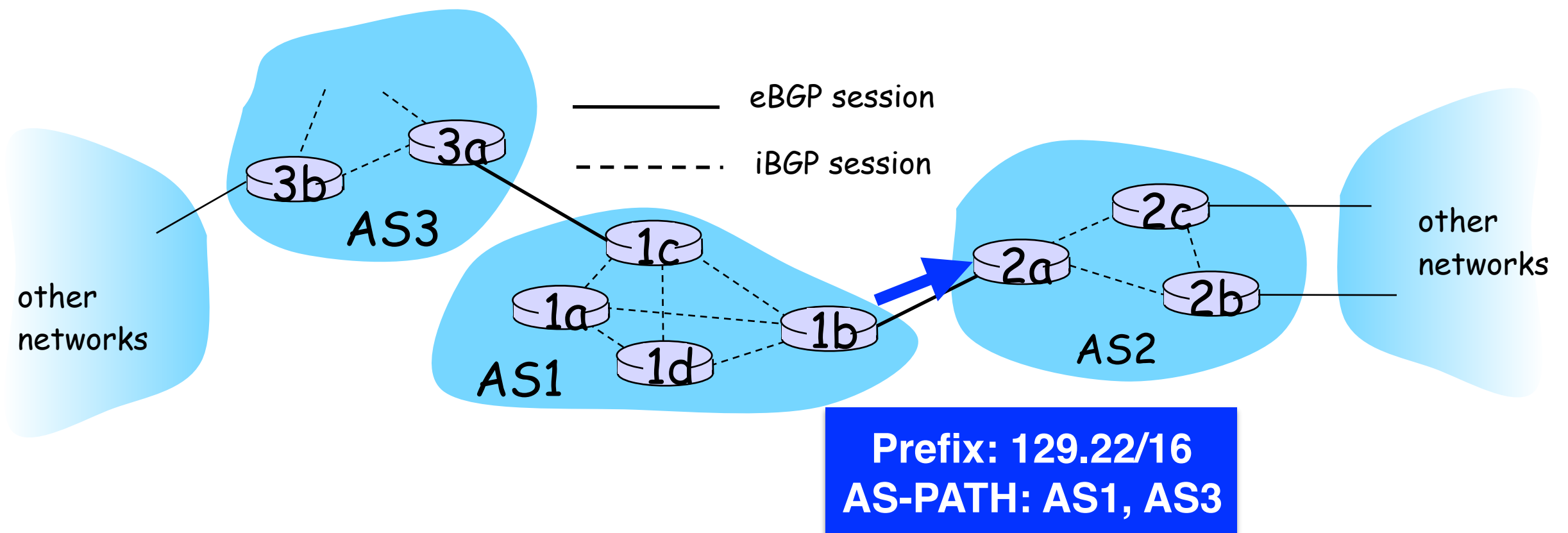
Path attributes & BGP routes



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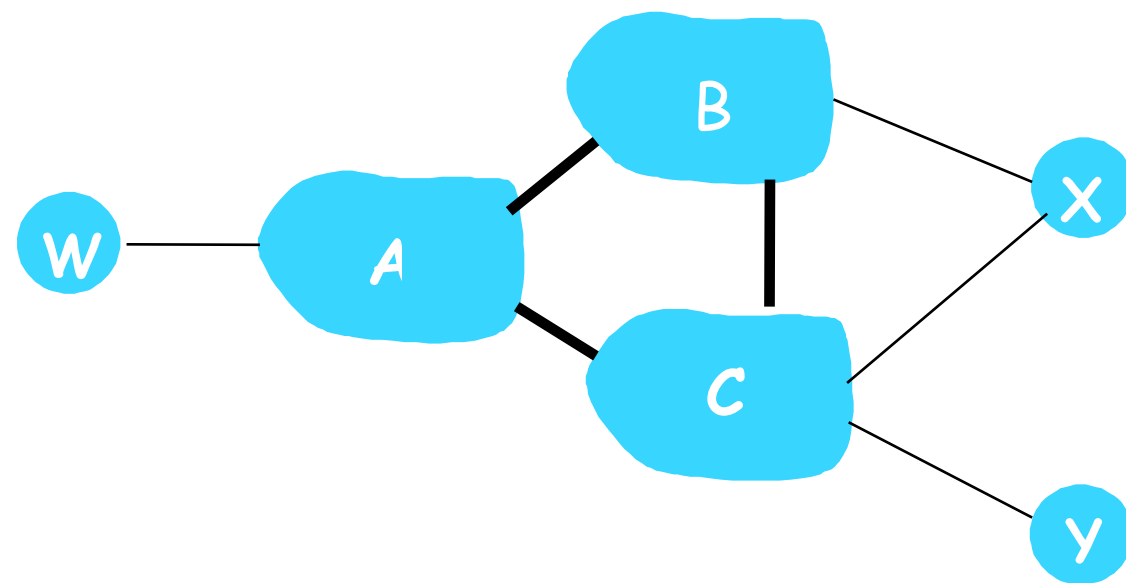
Route Import



- ❖ gateway router receiving route advertisement uses **import policy** to accept/decline
 - e.g., never route through AS x
 - **policy-based** routing

BGP route selection

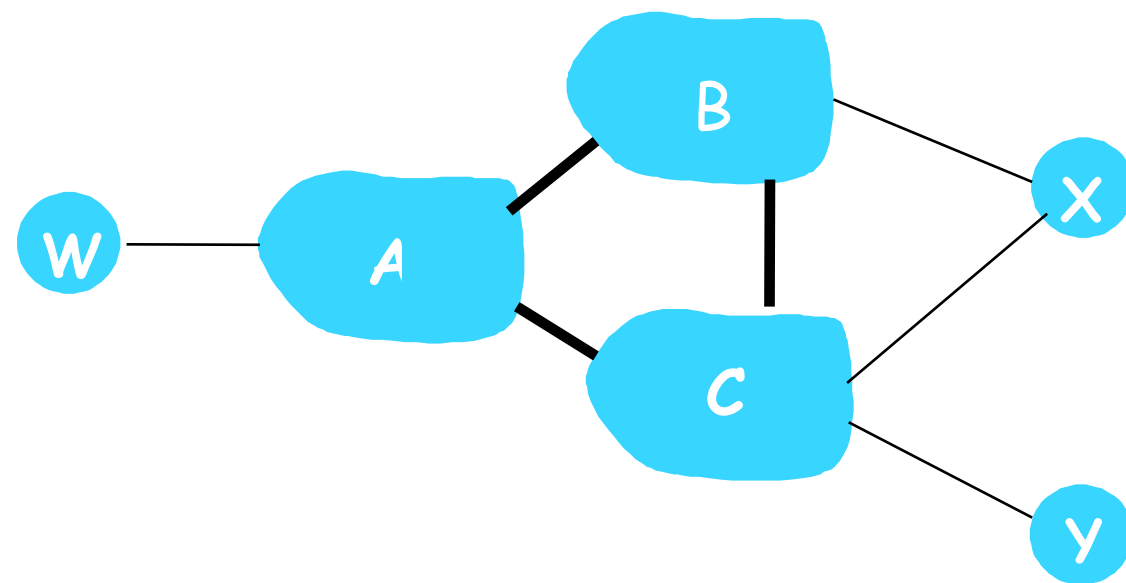
- ❖ router may learn about more than 1 route to destination AS, selects route based on:
 1. local preference value attribute: policy decision
 2. shortest AS-PATH
 3. closest NEXT-HOP router: hot potato routing
 4. additional criteria

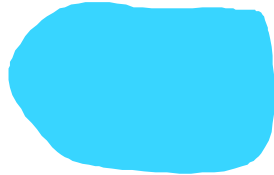

BGP routing policy



legend:  provider network
 customer network:

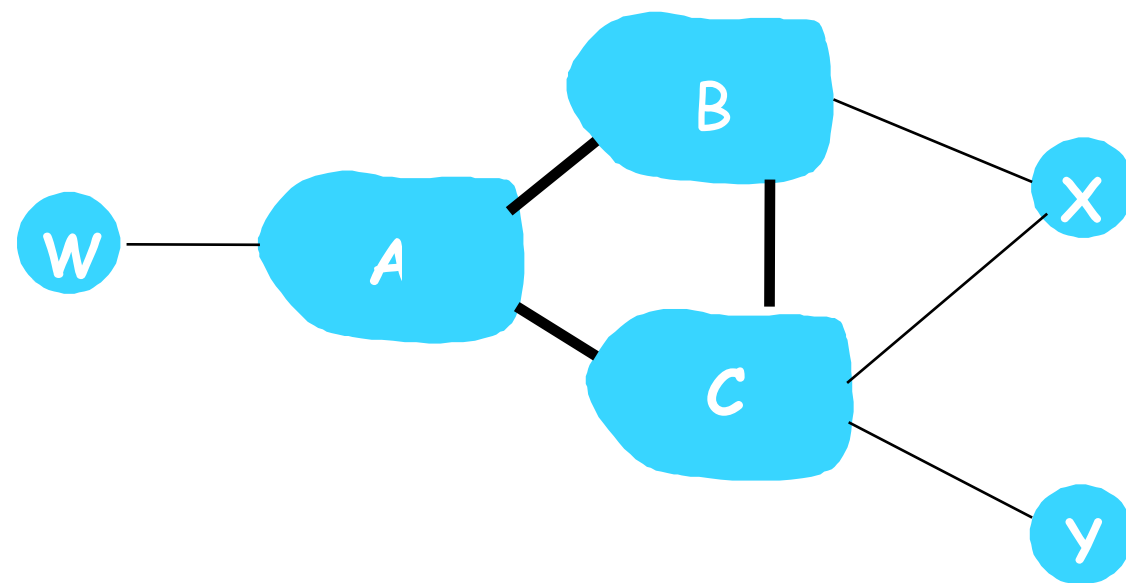
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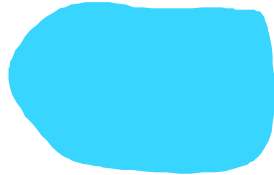



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- ❖ A,B,C are **provider networks**
- ❖ X,W,Y are customers (of provider networks)

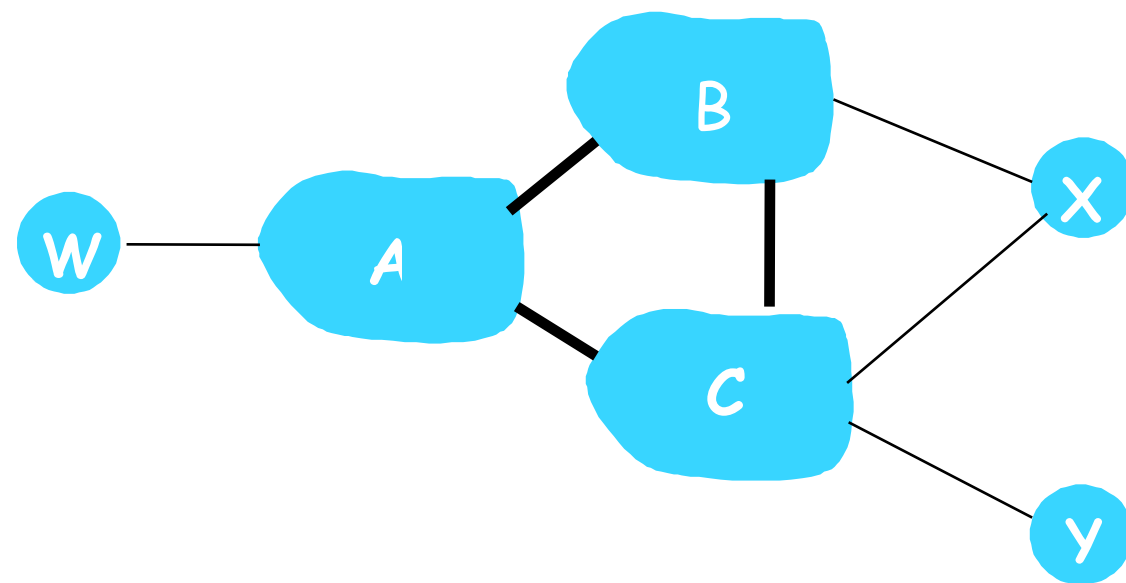
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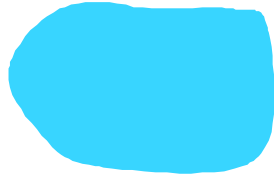



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- ❖ X is **dual-homed**: attached to two networks

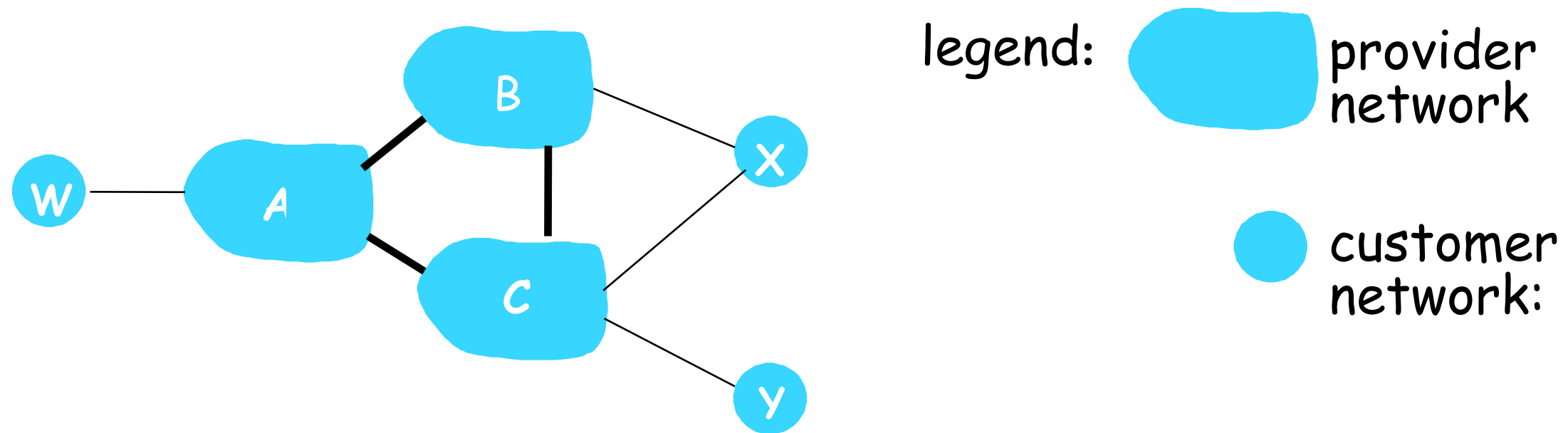
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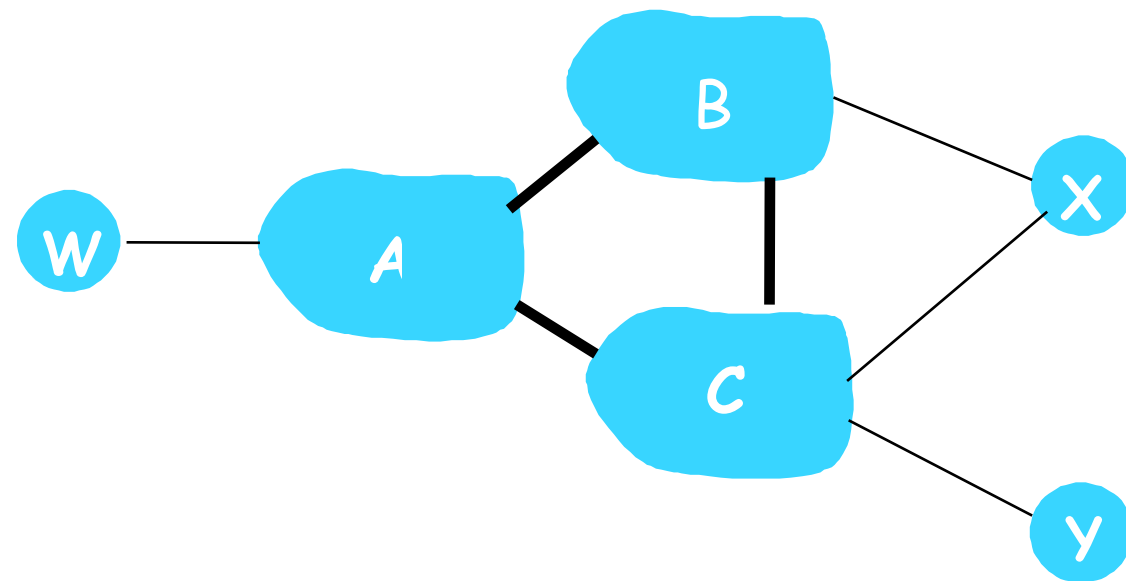
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 - X does not want to route traffic from B to C

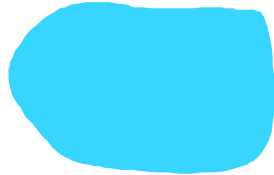

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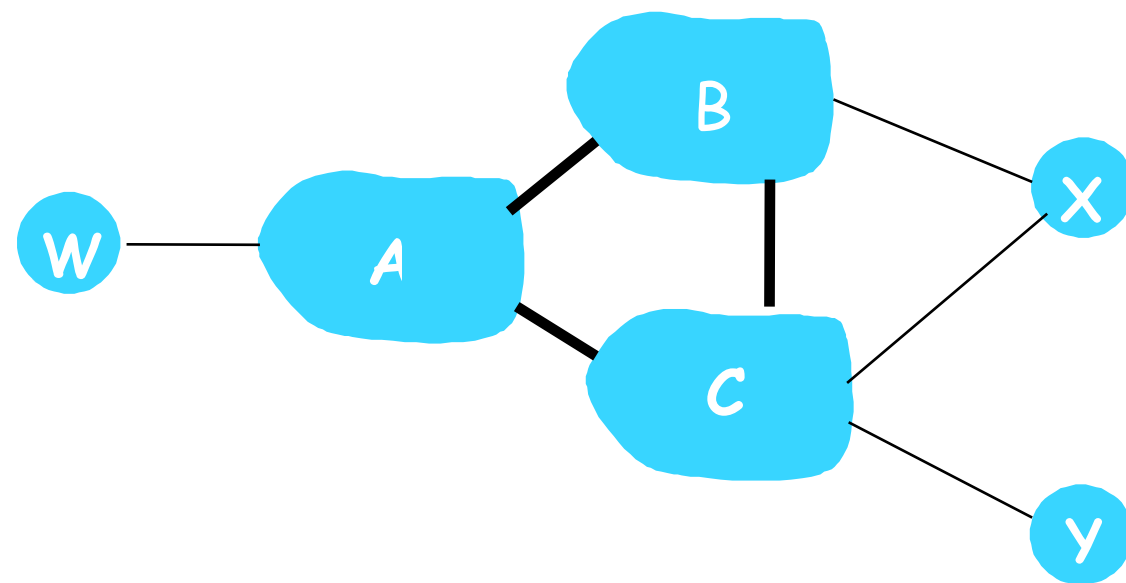
- ❖ A,B,C are **provider networks**
- ❖ X,W,Y are customers (of provider networks)
- ❖ X is **dual-homed**: attached to two networks
 - X does not want to route traffic from B to C
 - .. so X will not advertise a route to C in messages to B

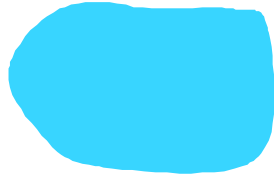

BGP routing policy (2)



legend:  provider network
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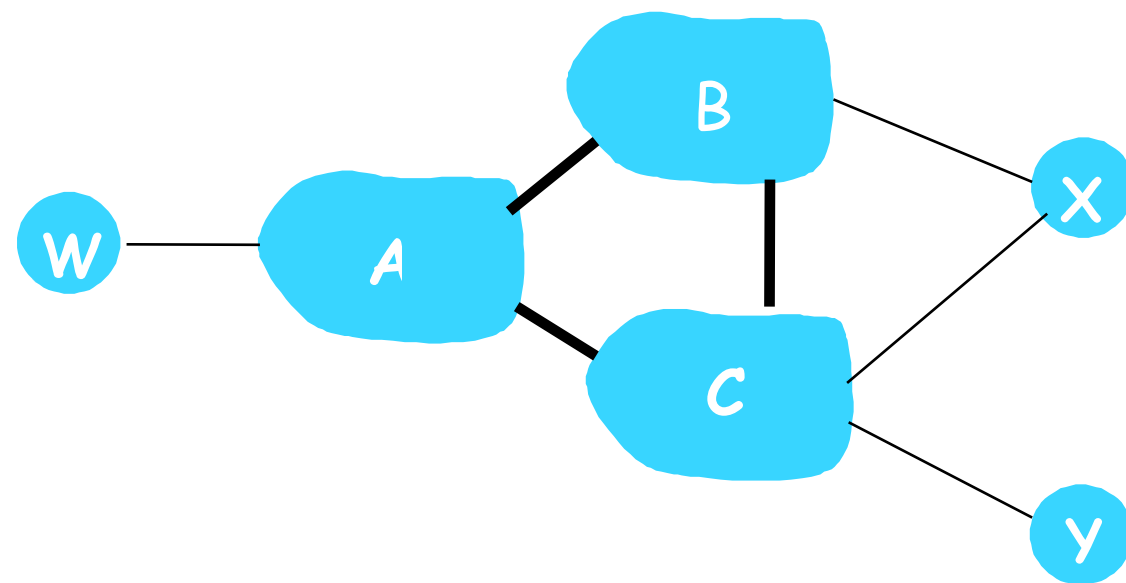
BGP routing policy (2)



legend:  provider network
 customer network:

- ❖ A advertises path *AW* to B
- ❖ B advertises path *BAW* to X

BGP routing policy (2)



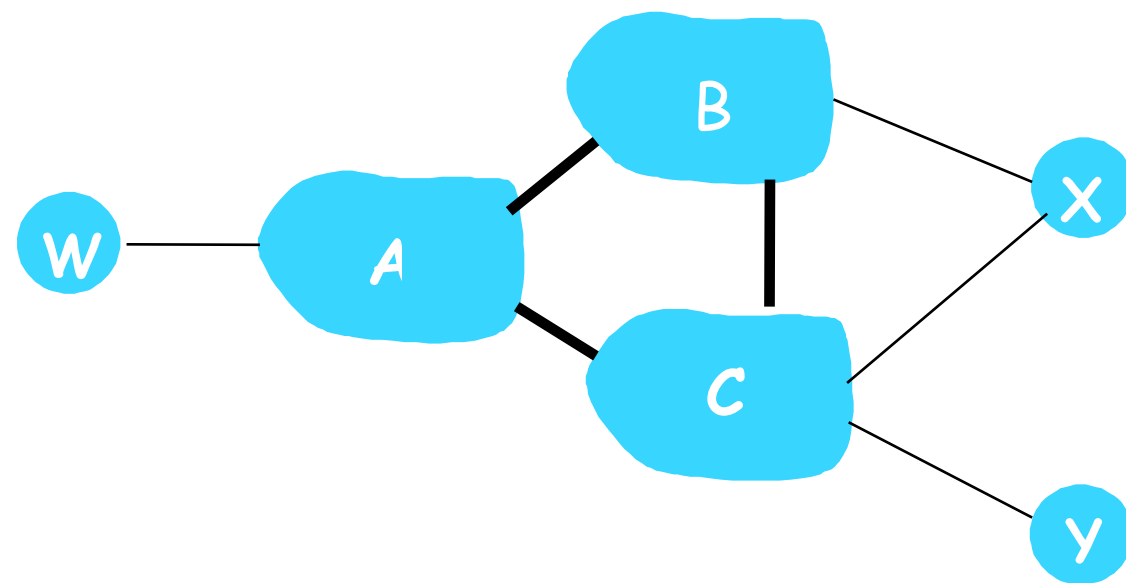
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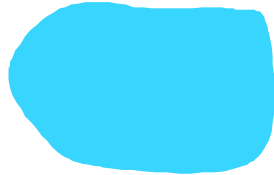

 provider network

 customer network:

- ❖ A advertises path AW to B
- ❖ B advertises path BAW to X
- ❖ Should B advertise path BAW to C?

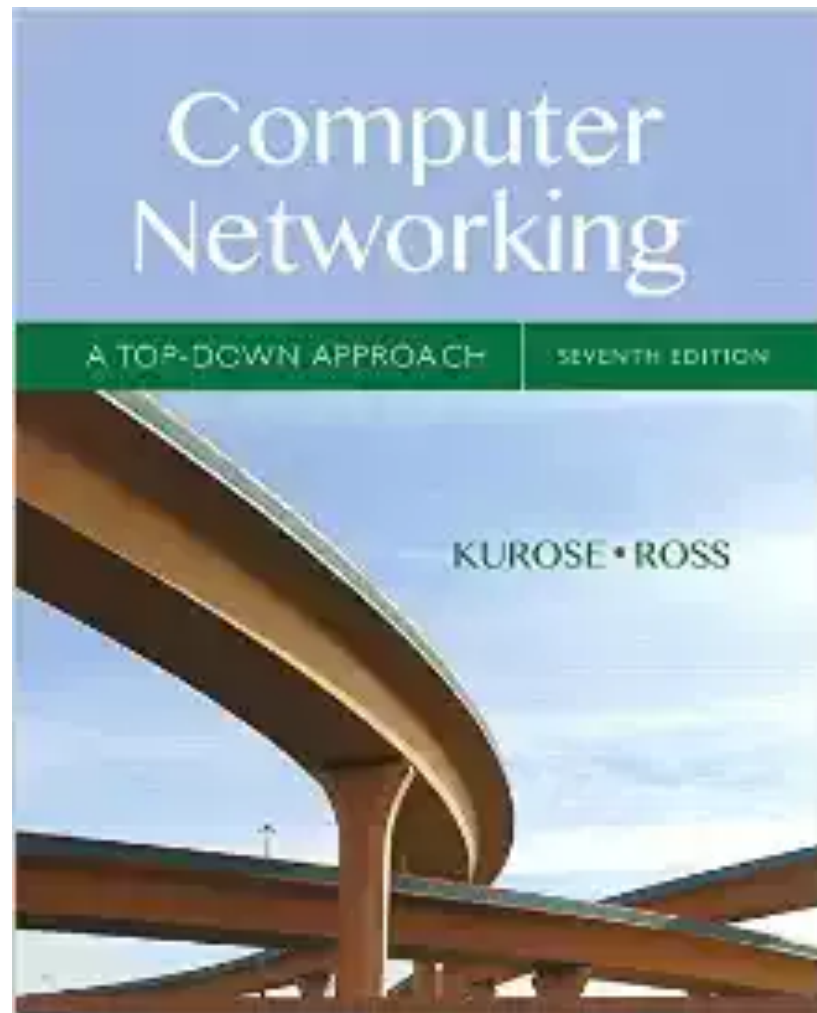
BGP routing policy (2)



legend:  provider network
 customer network:

- ❖ A advertises path AW to B
- ❖ B advertises path BAW to X
- ❖ Should B advertise path BAW to C?
 - No way! B gets no "revenue" for routing CBAW since neither W nor C are B's customers
 - B wants to force C to route to w via A
 - B wants to route **only** to/from its customers!

Reading Along ...



- Network layer is chapters 4 & 5
- Transmission Classes

Transmission Classes

Transmission Classes

❖ Unicast

- send message to a single recipient

Transmission Classes

❖ Unicast

- send message to a single recipient

❖ Broadcast

- send same message to everyone

Transmission Classes

Transmission Classes

❖ Multicast

Transmission Classes

❖ Multicast

- when sending the same content to multiple destinations
- ... but not everyone!
- e.g., radio station broadcast

Transmission Classes

Transmission Classes

❖ Anycast

Transmission Classes

❖ Anycast

- for replication of services, any one of which will work
- e.g., root DNS servers