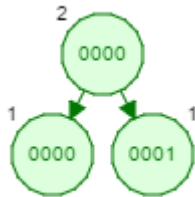


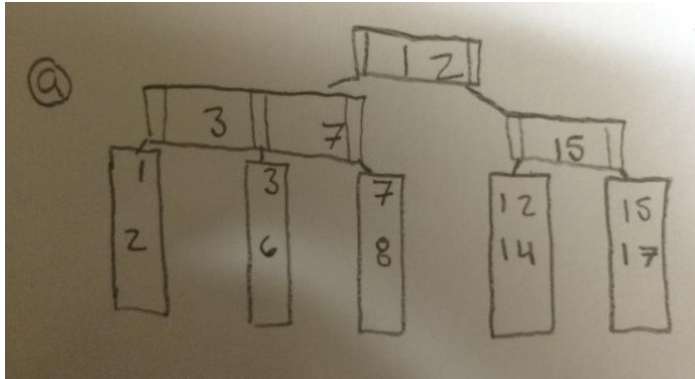
- 1) Our counter Example will include arbitrary functions $f(n) = \sin(n)$ and $g(n) = \cos(n)$. These two functions were chosen because they are not monotonic. These two functions have periods of increasing and decreasing. Therefore, at many points the two functions will switch between which is larger. However, by adding a constant, one of these functions can dominate over the other and the theory of Big – O is held.

If we drop the constant, like the procedure we discussed in class, one of the two functions above will no longer be dominating throughout the entire domain. Big – O is no longer preserved. The constant in such non-monotonic functions is important and cannot be dropped.

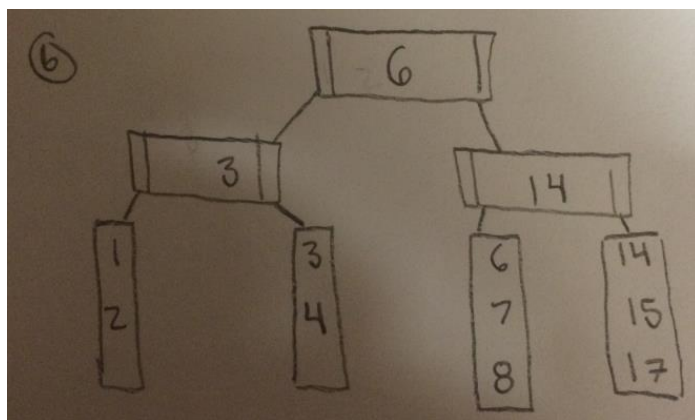
- 2) If we add value 0 into our tree, then a number greater than 0, and then again zero, then we will obtain the desired AVL tree. The right child of the first node will be the number greater than zero. The left node to this child will be zero. The AVL tree will need to rearrange to balance out and will produce the tree below.



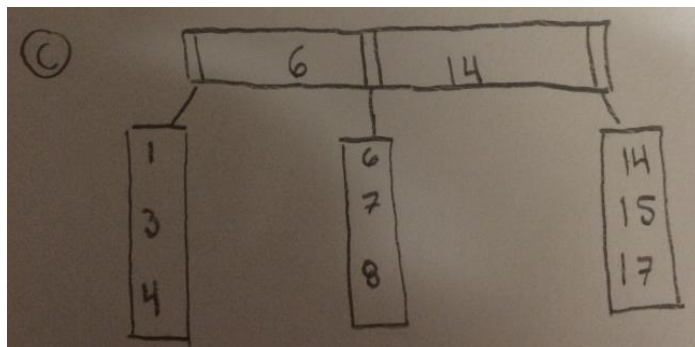
3)



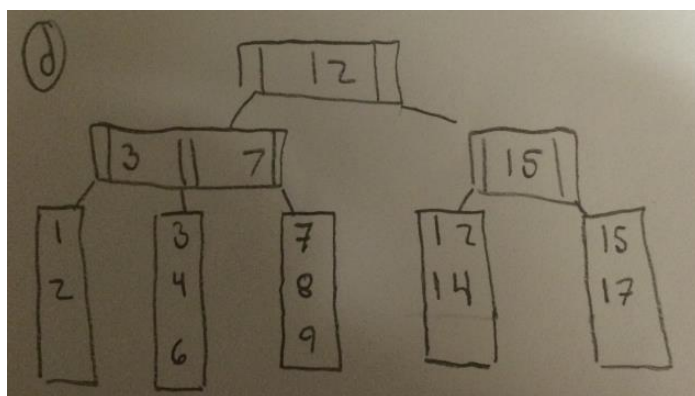
This is a simple problem of removing the number four and pulling from the next consecutive leaf so the under loaded leaf is satisfied. Re arrange the pointers.



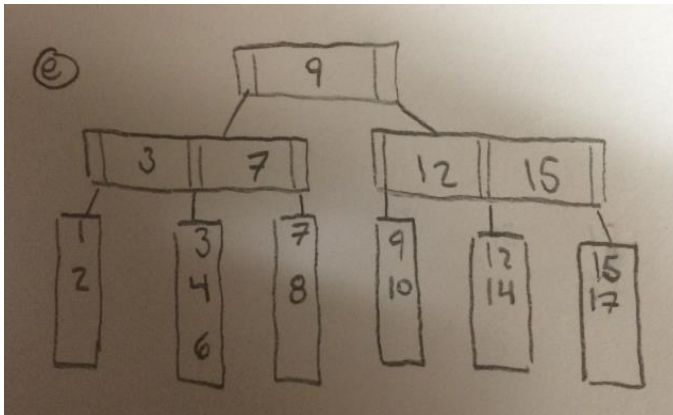
You must first remove the value 12 from the leaf. That leaf becomes under loaded, so you merge 14 into the next leaf. Now the tree is unbalanced so you must pull a leaf over and rearrange the pointers.



By removing 2, the leaf becomes under loaded. The only way to satisfy this is to merge with the next leaf and remove our top pointer which means we must merge our remaining pointers.

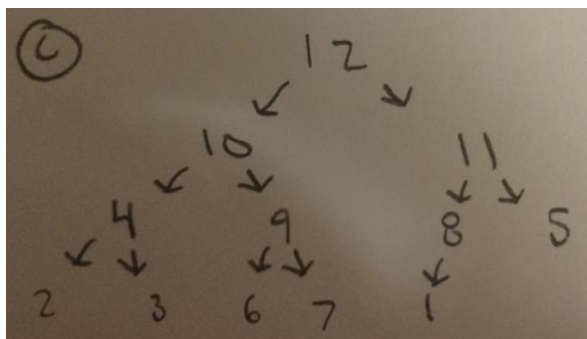
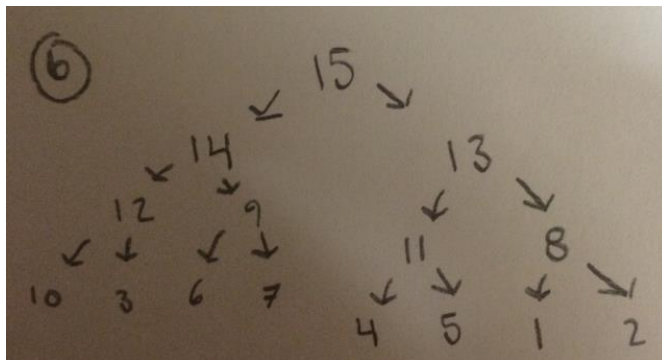
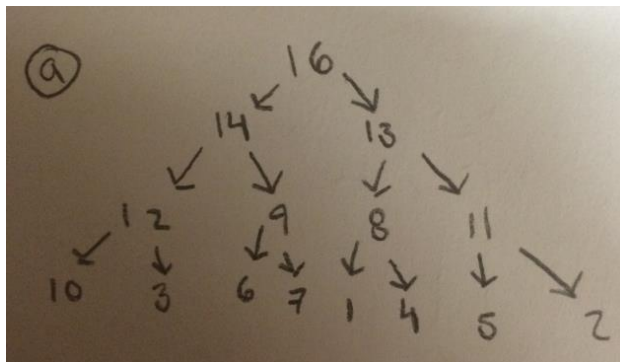


We must first add 9, but the leaf will then be over loaded. So we move the least value to the previous leaf. And then change our pointers.



When adding ten, the leaf becomes over loaded, so we must move values. However its adjacent sibling is full. So then it decides to split the overflowing node to the other, not full branch. The pointers must be changed accordingly.

4)



```

1  public class HW {
2
3      //array to store heap values
4      T[] array = new T[];
5
6      public void insert(T item) {
7          //adds item at the end of the array
8          array[size] = item;
9          //increment the size so sift up includes the new value
10         size++;
11         //sifts the array and sorts in order of
12         array.siftUp(size);
13     }
14
15     public T removeMax() {
16         //variable to store max value
17         T variable = array[0];
18         //setting the first heap value
19         array[0] = array[size - 1];
20         //change size for exclusion the last heap value
21         size--;
22         //sifting to reheapify
23         array.siftDown(size);
24         //max value
25         return variable;
26     }
27
28     public void siftUp(int i){
29         //if we arent trying to sift up the first node
30         if(i != 0) {
31             //if the parent is greater than or equal to its child
32             if(array[(i - 1)/2].compareTo(array[i]) == 1 || array[(i - 1)/2].compareTo(array[i]) == 0) {
33                 //do nothing
34             }
35             //if parent is less than child
36             else {
37                 //keep a variable to store value of parent
38                 T variable = array[(i - 1)/2];
39                 //parent equals child
40                 array[(i - 1)/2] = array[i];
41                 //child equals parents value
42                 array[i] = variable;
43                 //recursively call sift up
44                 array.siftUp((i-1)/2);
45             }
46         }

```

```

47     }
48
49     public void siftDown(int i) {
50         //if parent has two children
51         if(2i + 1 < size && 2i + 2 < size) {
52             //if left child is greater than right
53             if(array[2i + 1].compareTo(array[2i + 2]) == 1) {
54                 //keep a variable to store value of left child
55                 T variable = array[2i + 1];
56                 //left child equals parent
57                 array[2i + 1] = array[i];
58                 //parent equals left child value
59                 array[i] = variable;
60                 //recursively call sift down on left child
61                 array.siftDown(2i + 1);
62             }
63             //if right child is greater than or equal to left
64             else {
65                 //keep a variable to store value of right child
66                 T variable = array[2i + 2];
67                 //right child equals parent
68                 array[2i + 2] = array[i];
69                 //parent equals right child value
70                 array[i] = variable;
71                 //recursively call sift down on right child
72                 array.siftDown(2i + 2);
73             }
74         }
75         //if parent has only one child (it must be a left child)
76         else if(2i + 1 < size) {
77             //keep a variable to store value of left child
78             T variable = array[2i + 1];
79             //left child equals parent
80             array[2i + 1] = array[i];
81             //parent equals left child value
82             array[i] = variable;
83             //recursively call sift down on left child
84             array.siftDown(2i + 1);
85         }
86     }
87 }

```