

The Root of the Matter: A Discussion of DNS Security

Mark Allman
International Computer Science Institute

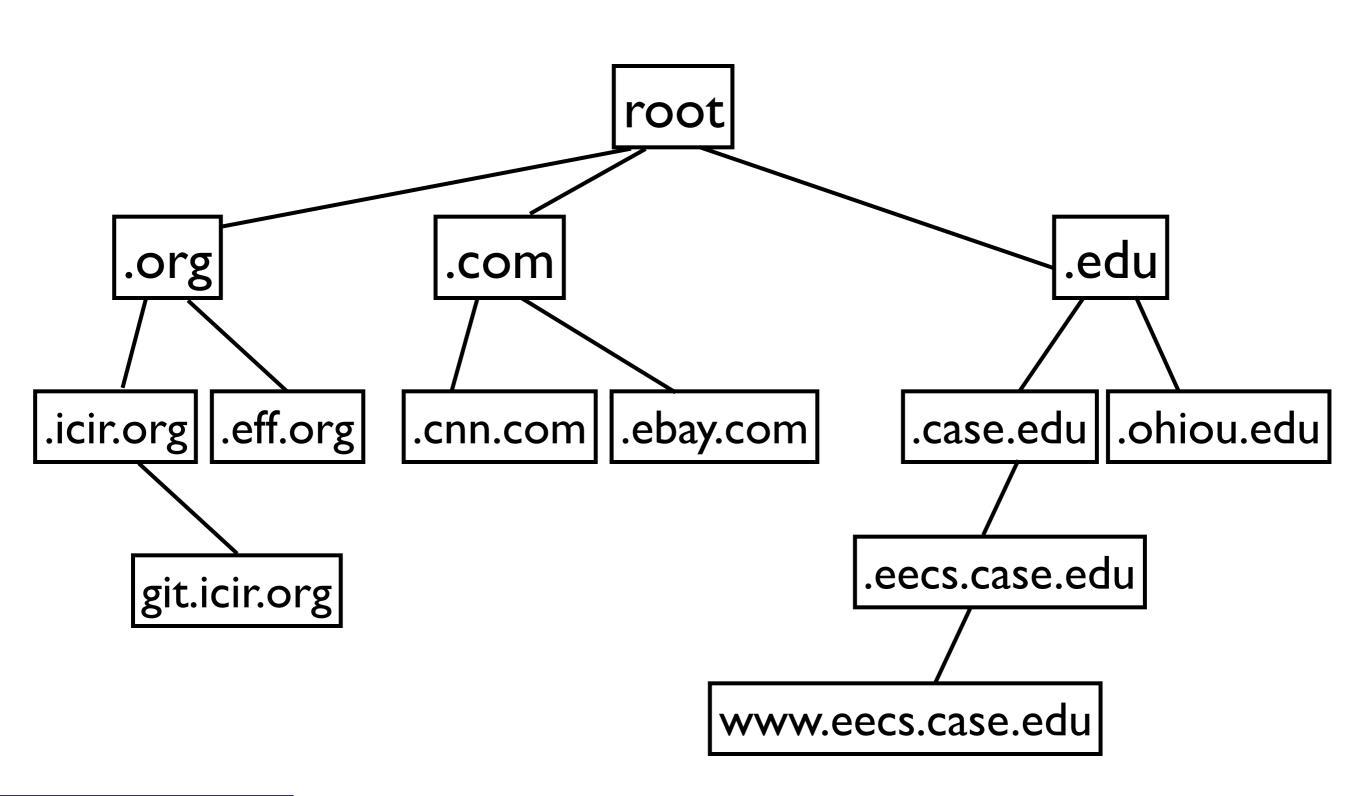
EECS 325 / 425 November 2018

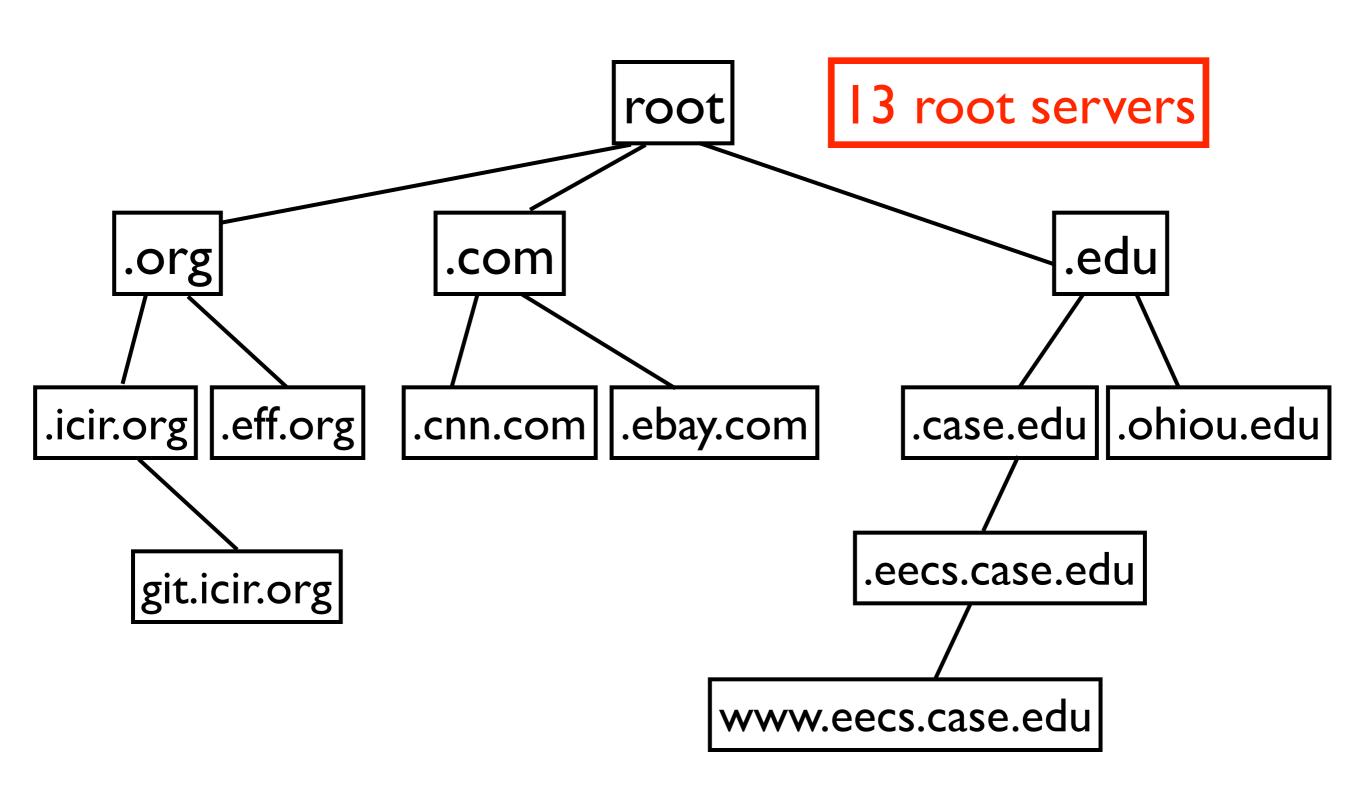
"Like a preacher stealin' hearts in a travelin' show ..."

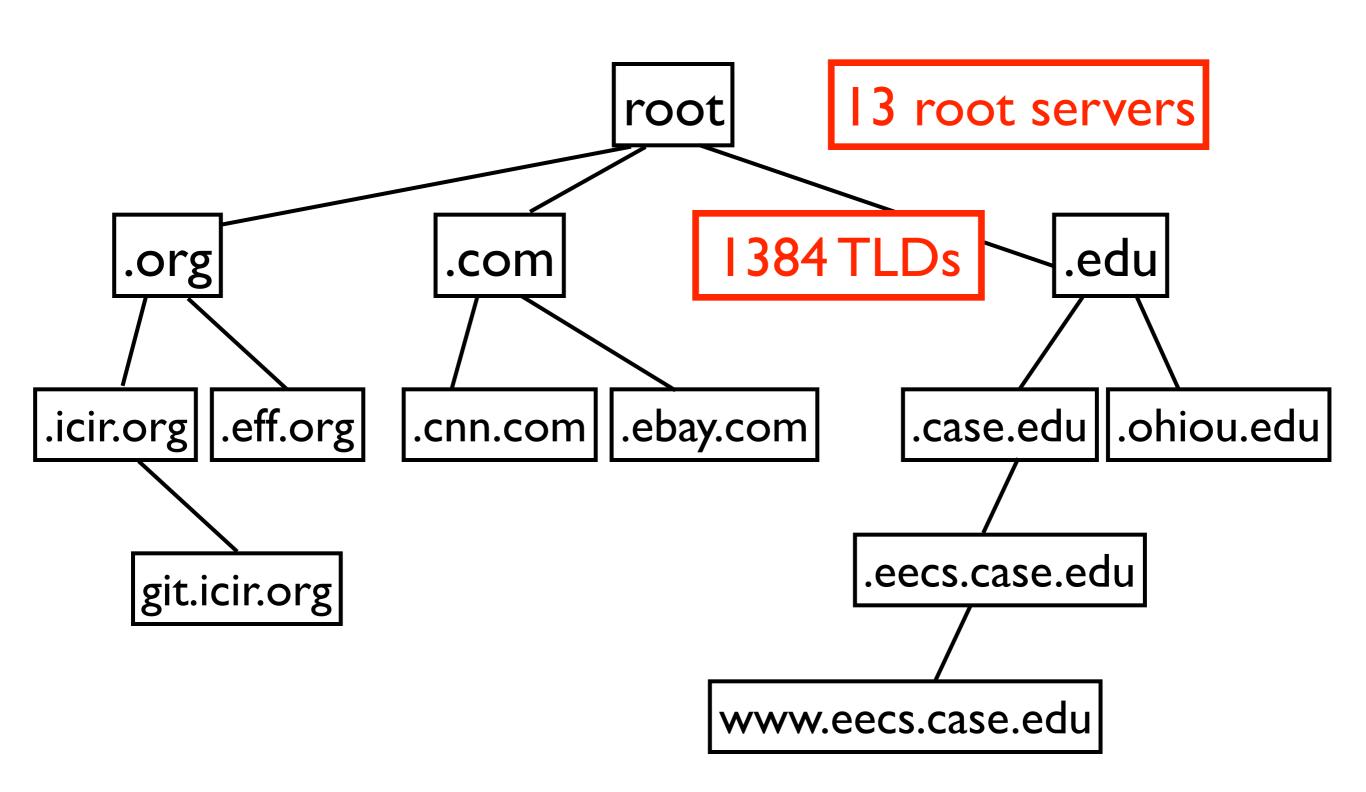
Collaborators

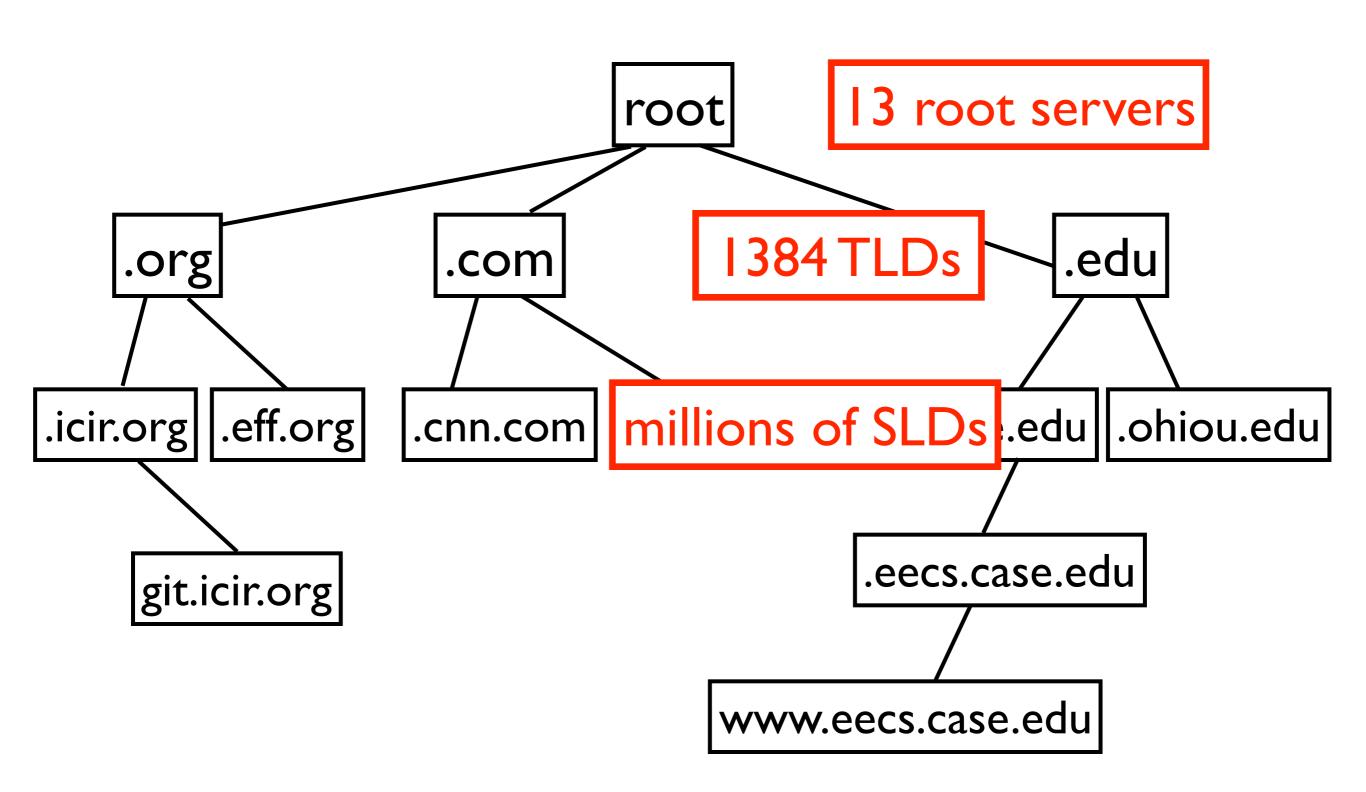
- Michael Bailey, Illinois/UMich
- Owen Bell, Case
- Tom Callahan, Case
- Jake Czyz, UMich
- Nick Feamster, Princeton
- Scott lekel-Johnson, Arbor
- Ben Jones, Princeton
- Andrew Kalafut, Grand Valley

- Eric Osterweil, Verisign
- Vern Paxson, ICSI & UCB
- Michael Rabinovich, Case
- Kyle Schomp, Case
- Craig Shue, WPI
- Nicholas Weaver, ICSI
- Jing Zhang, UMich







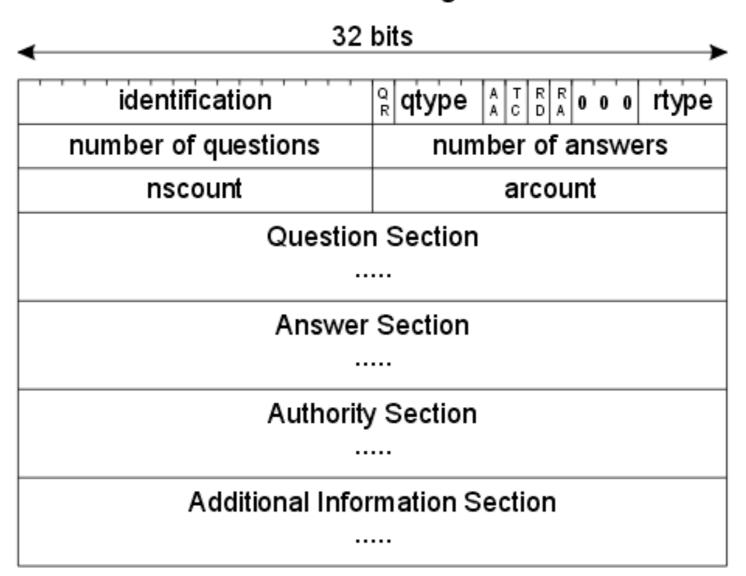


- DNS, as we use it:
 - connectionless UDP transport
 - single packet request
 - single packet response

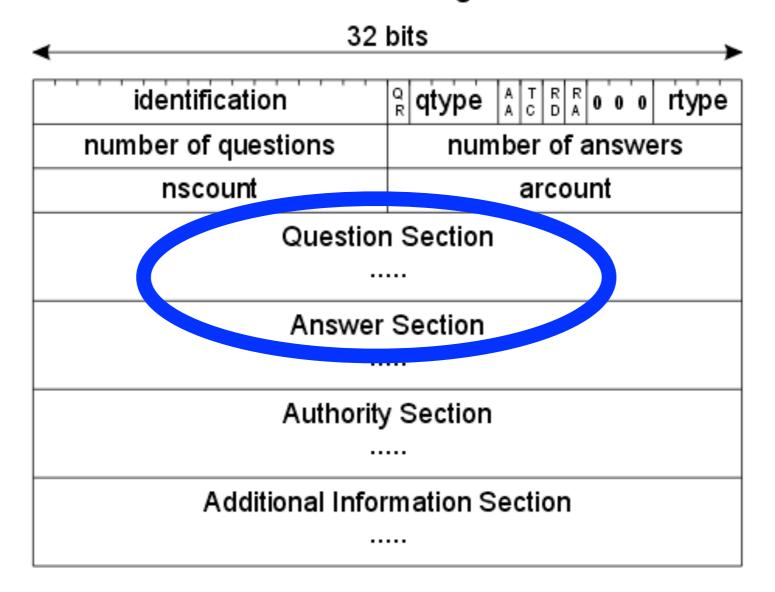
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 However, name resolution is structurally complex and mysterious with many hidden components

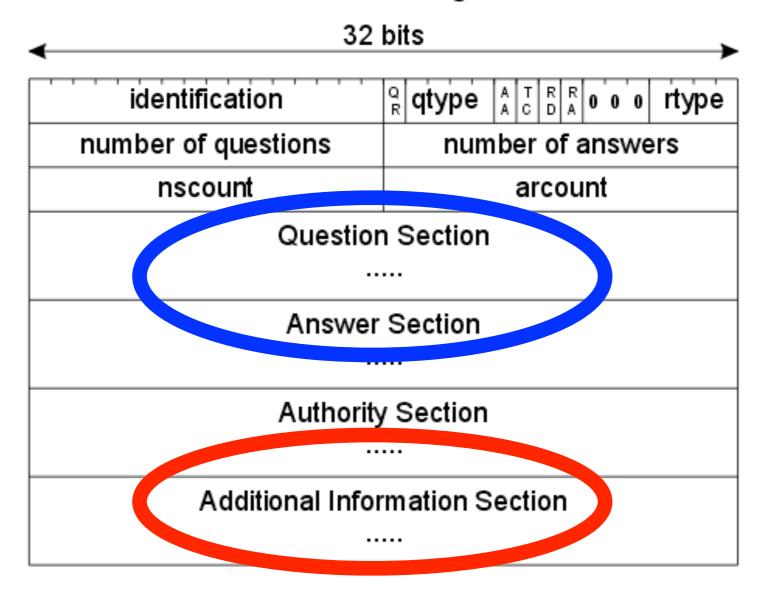
DNS server message format



DNS server message format



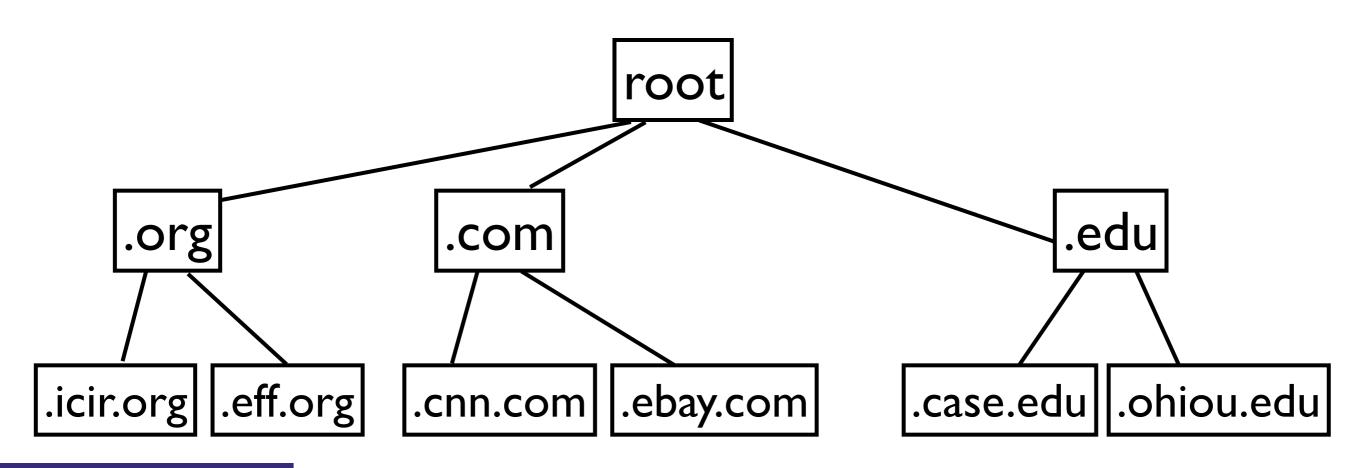
DNS server message format

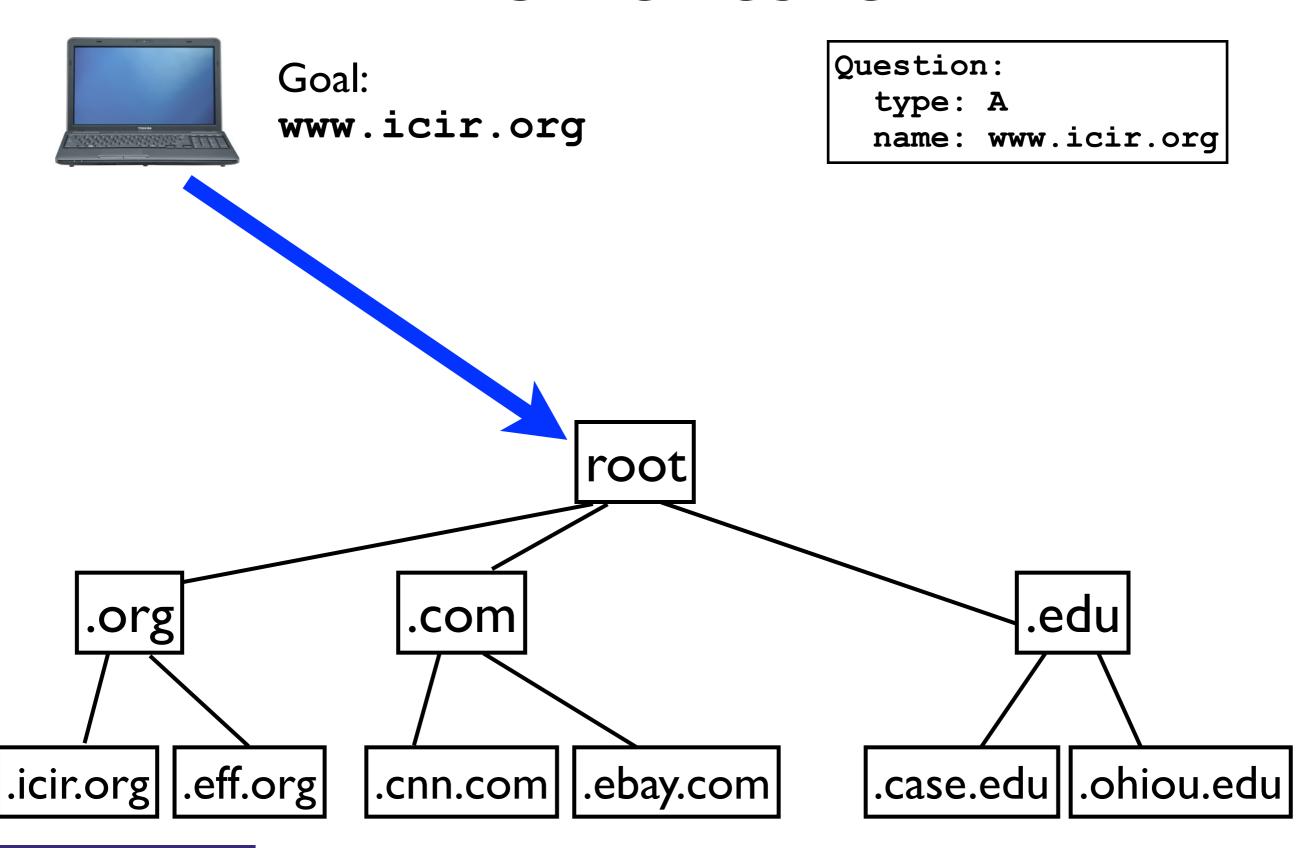


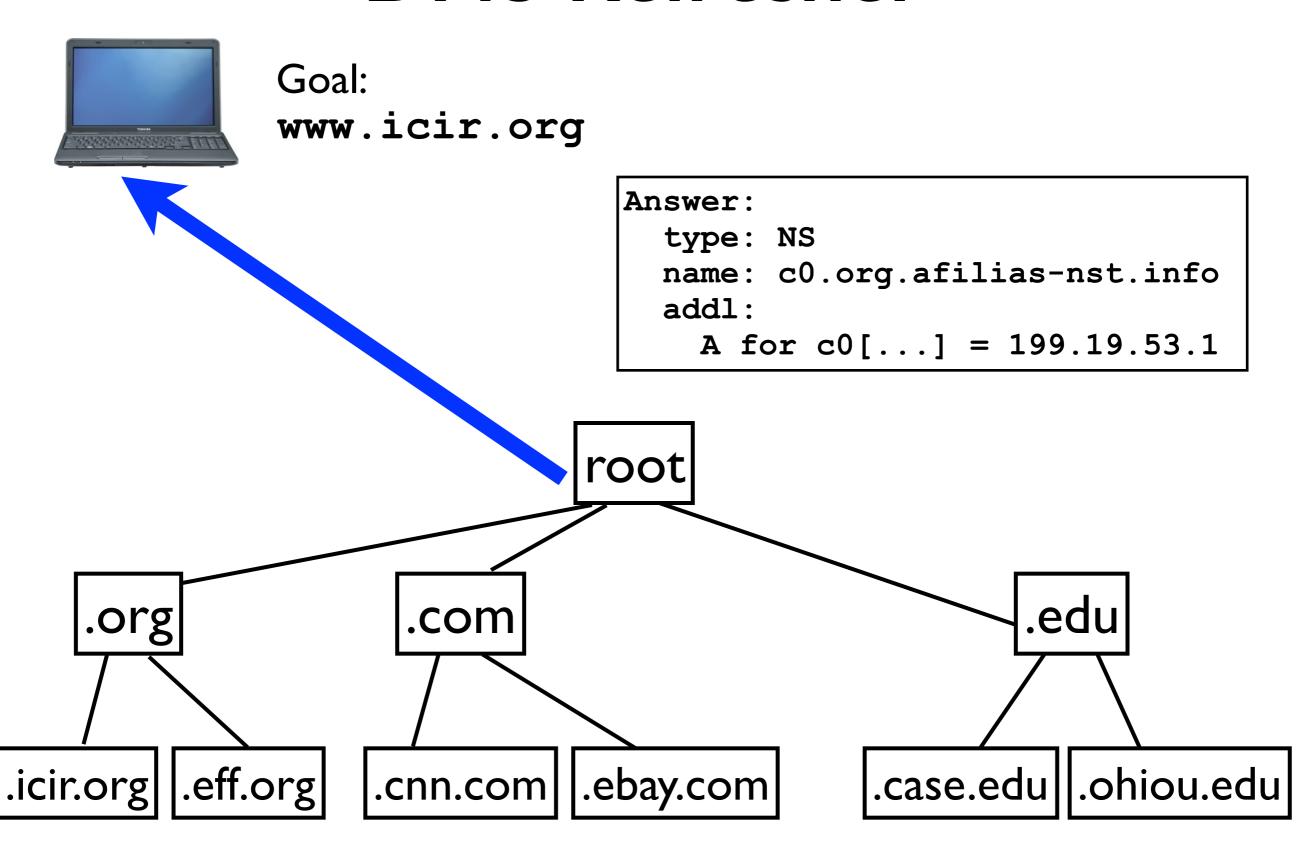


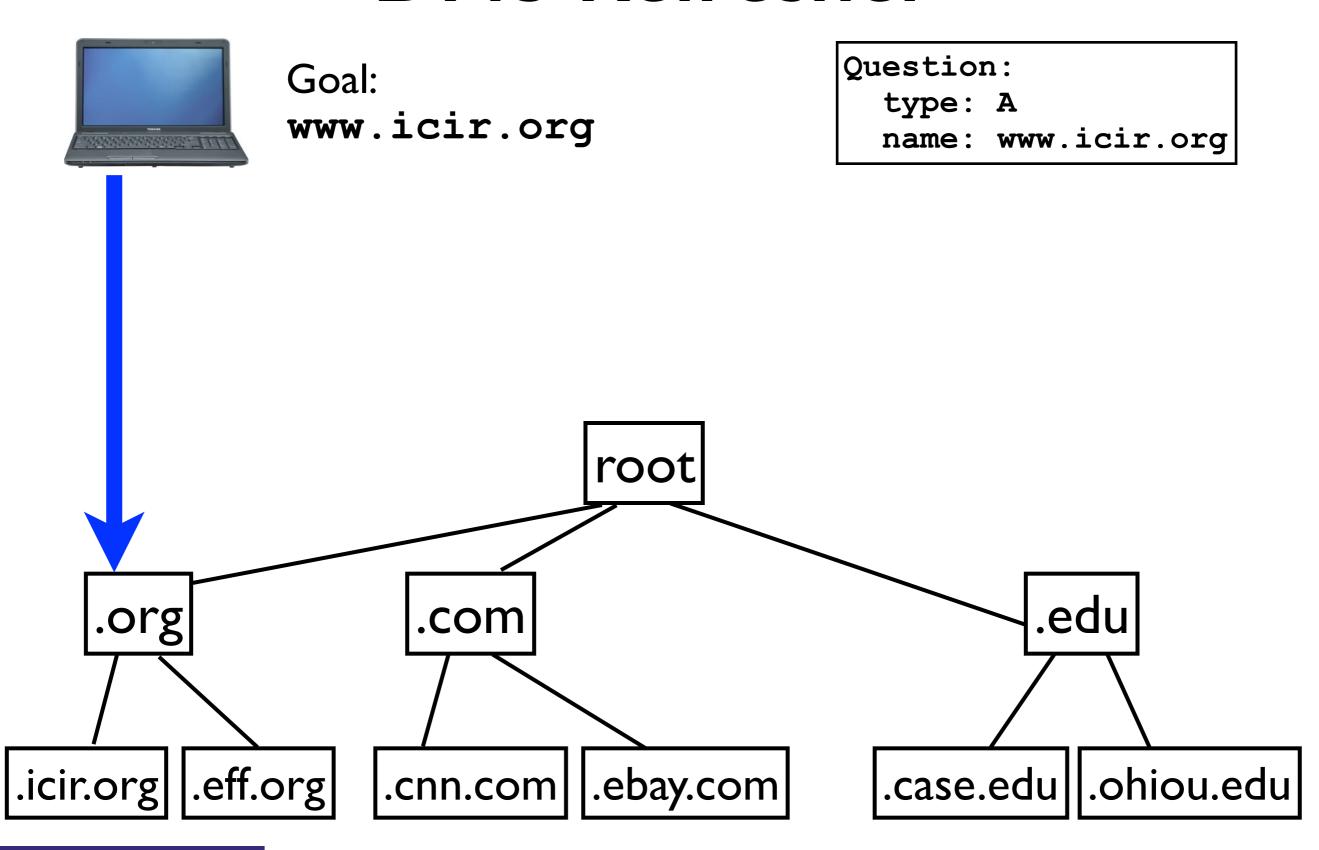
Goal:

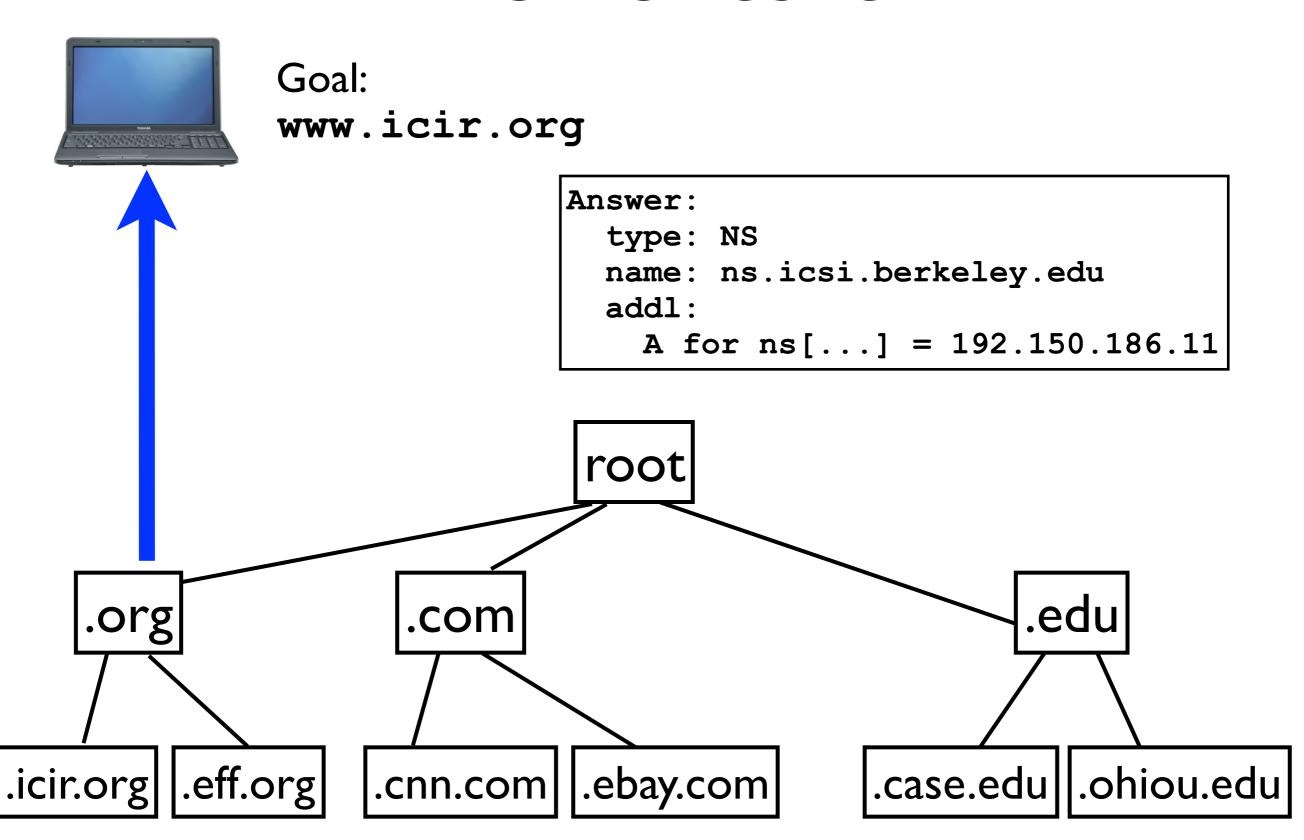
www.icir.org

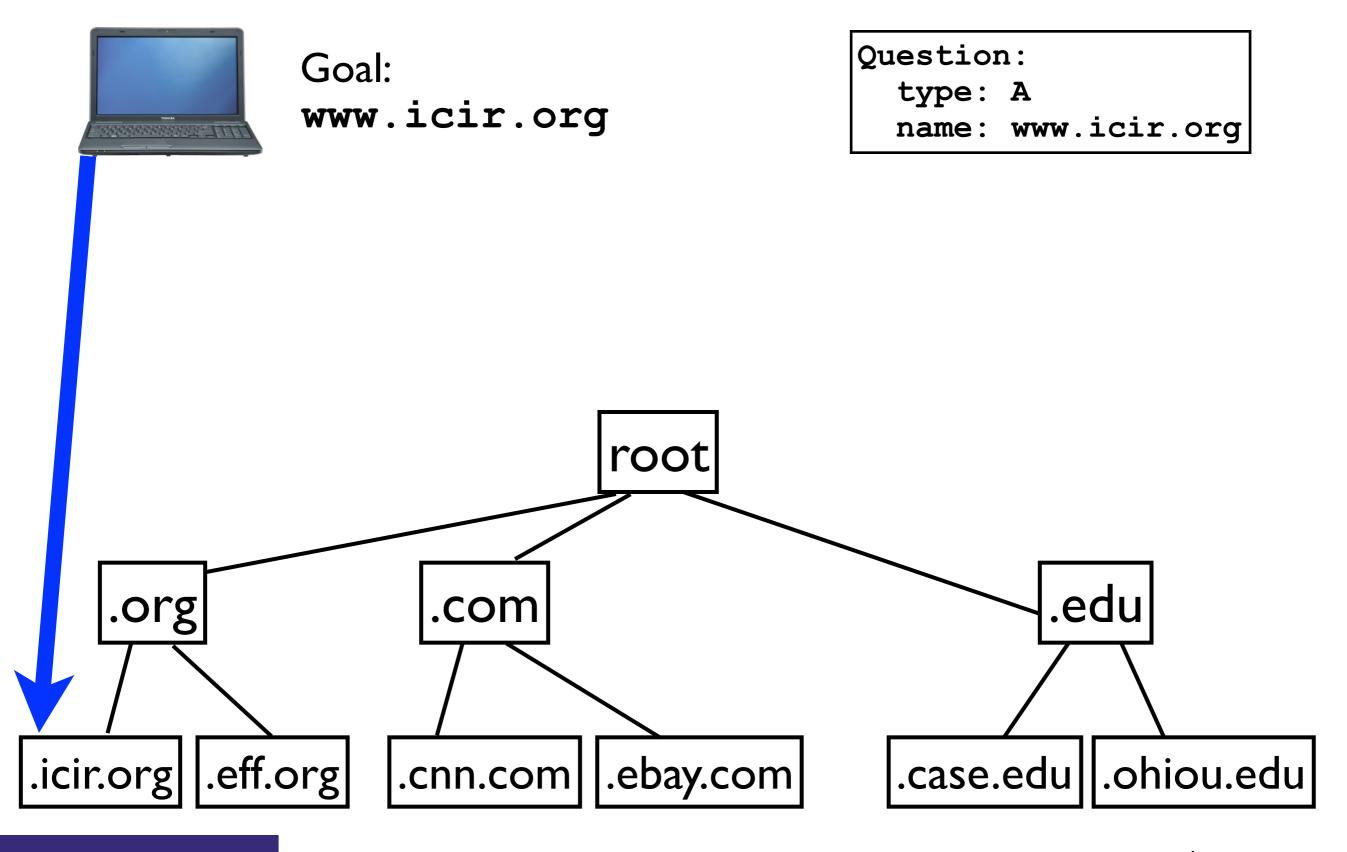


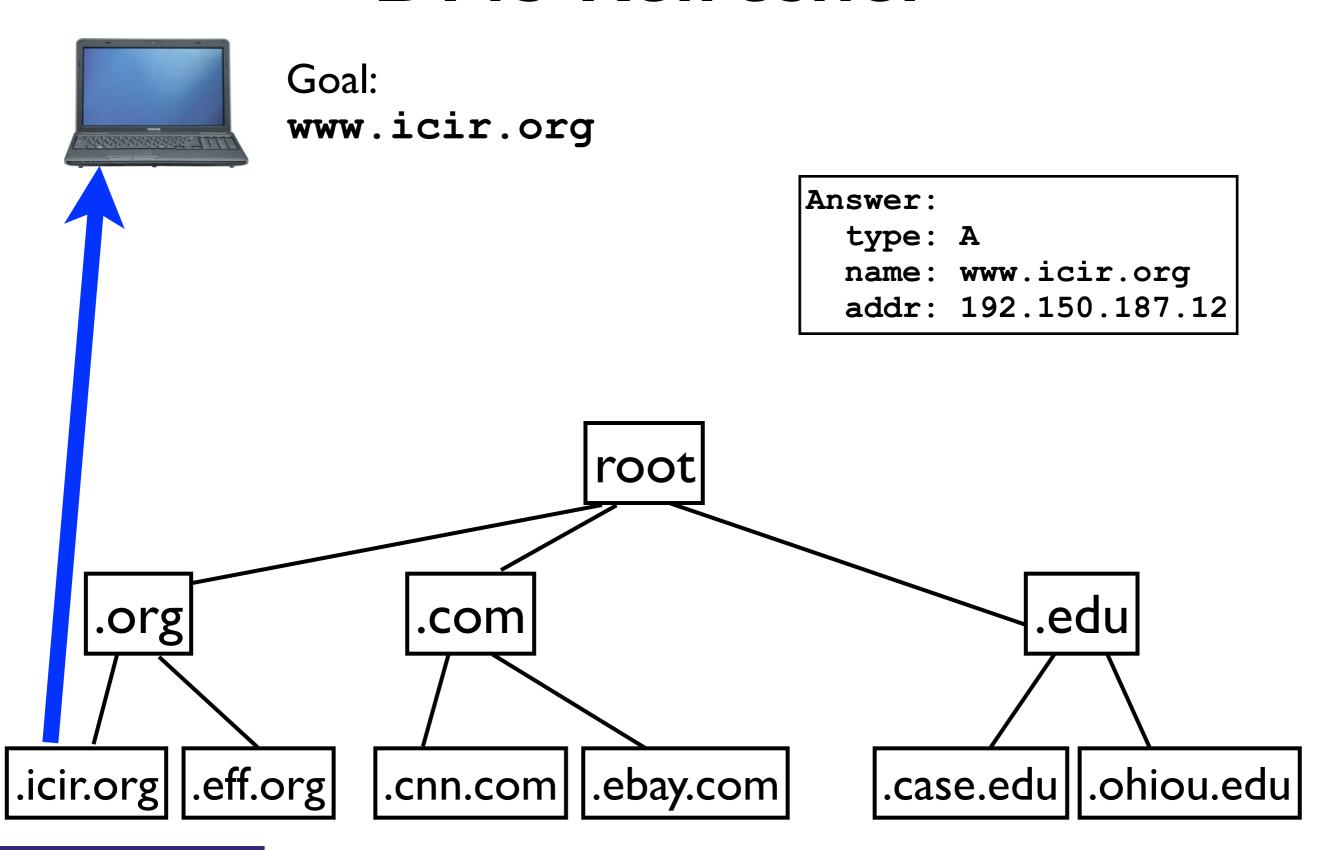




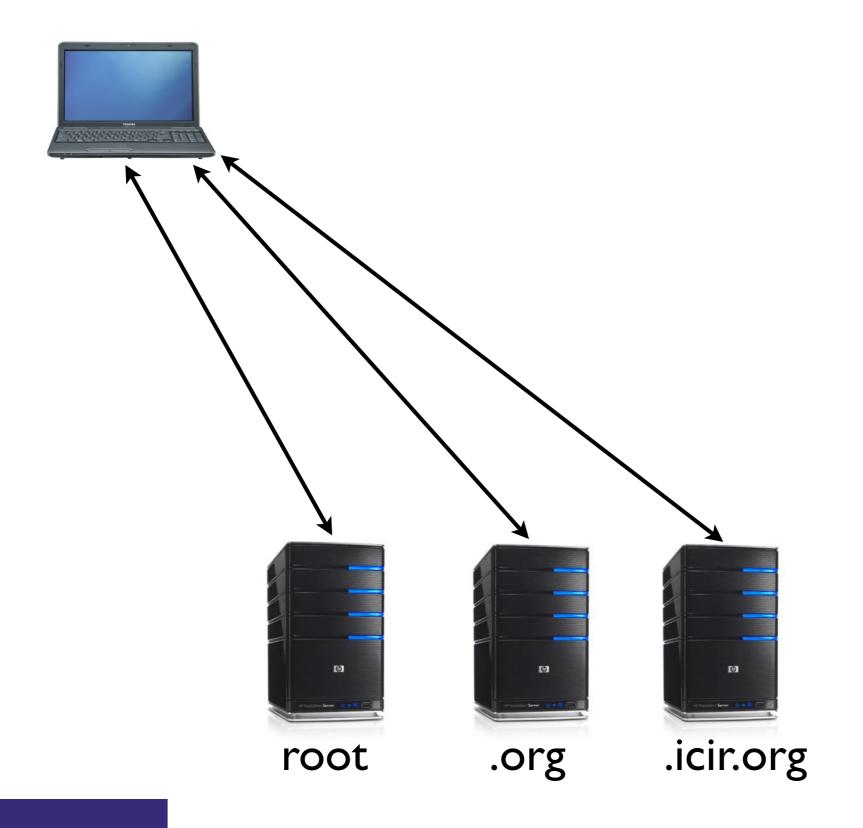




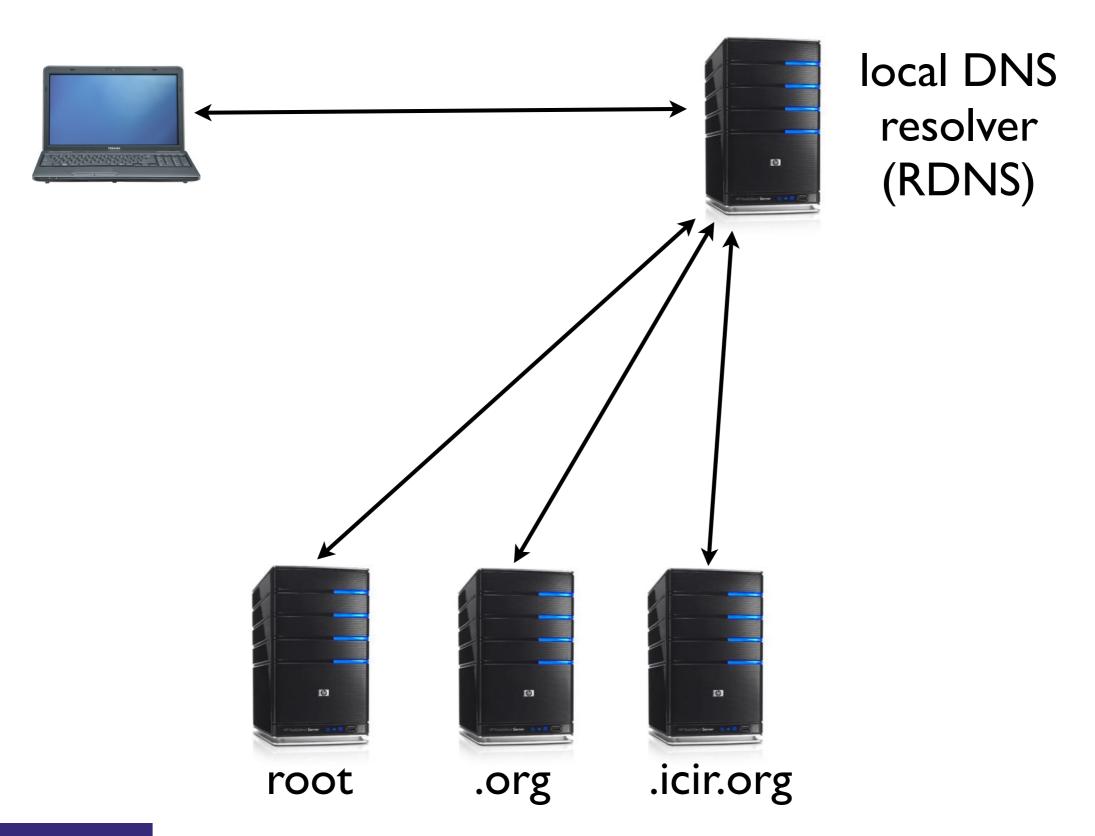




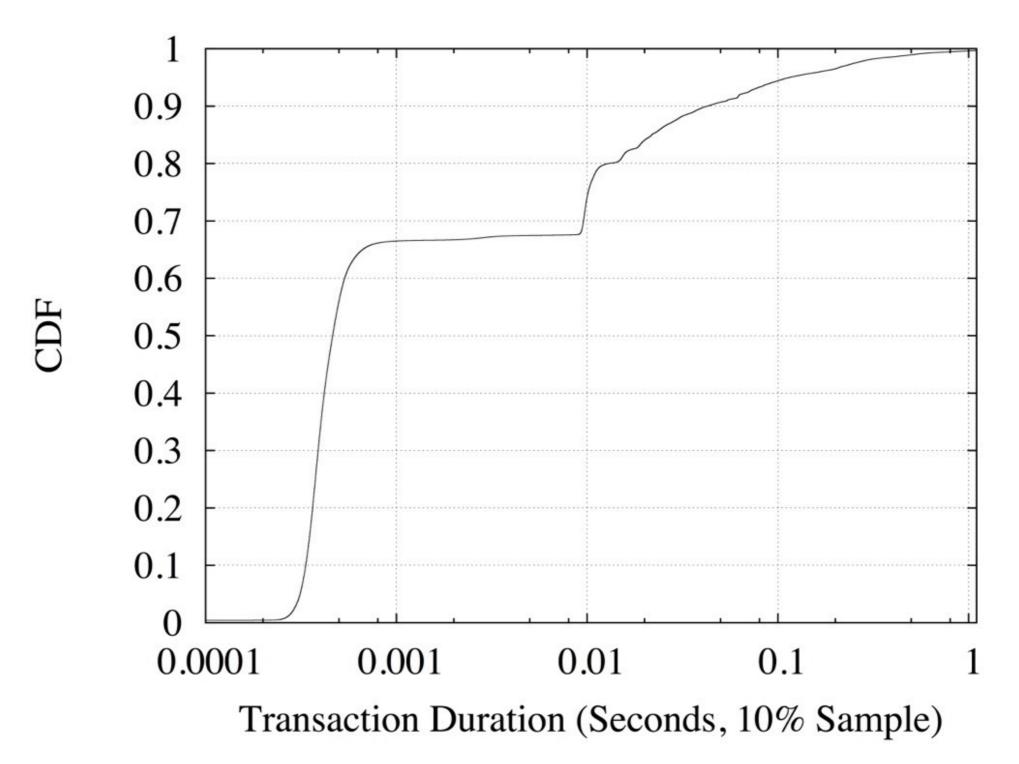
Basic Structure



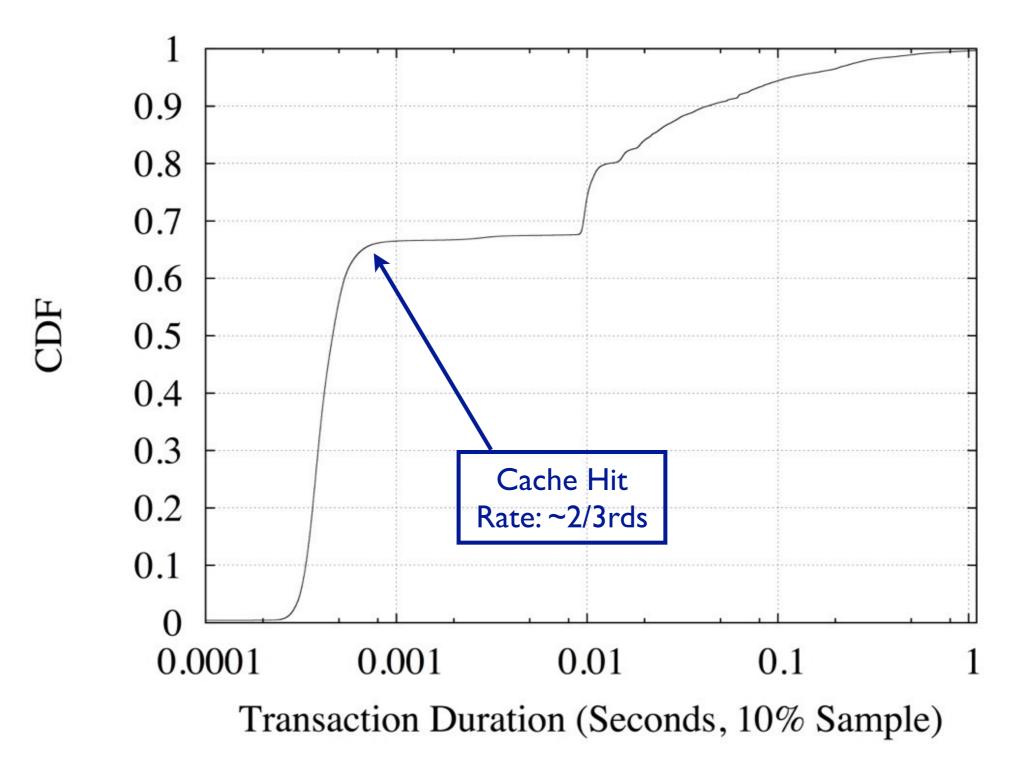
Delegating Resolution



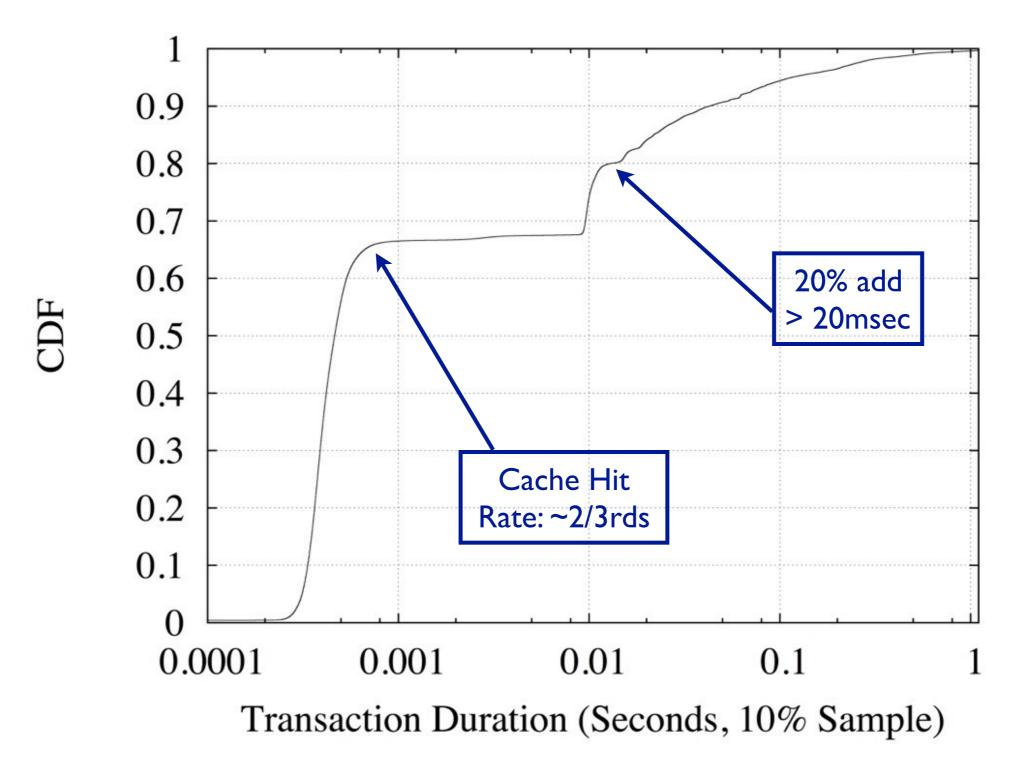
Caching Efficacy



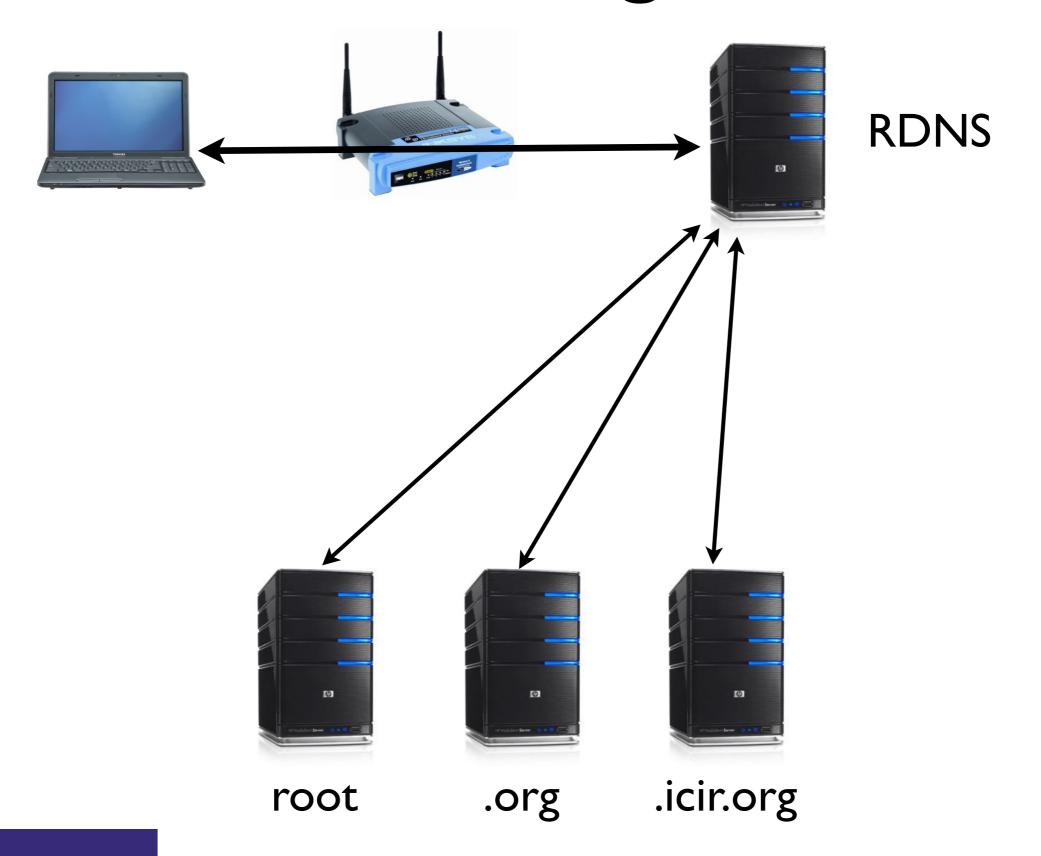
Caching Efficacy



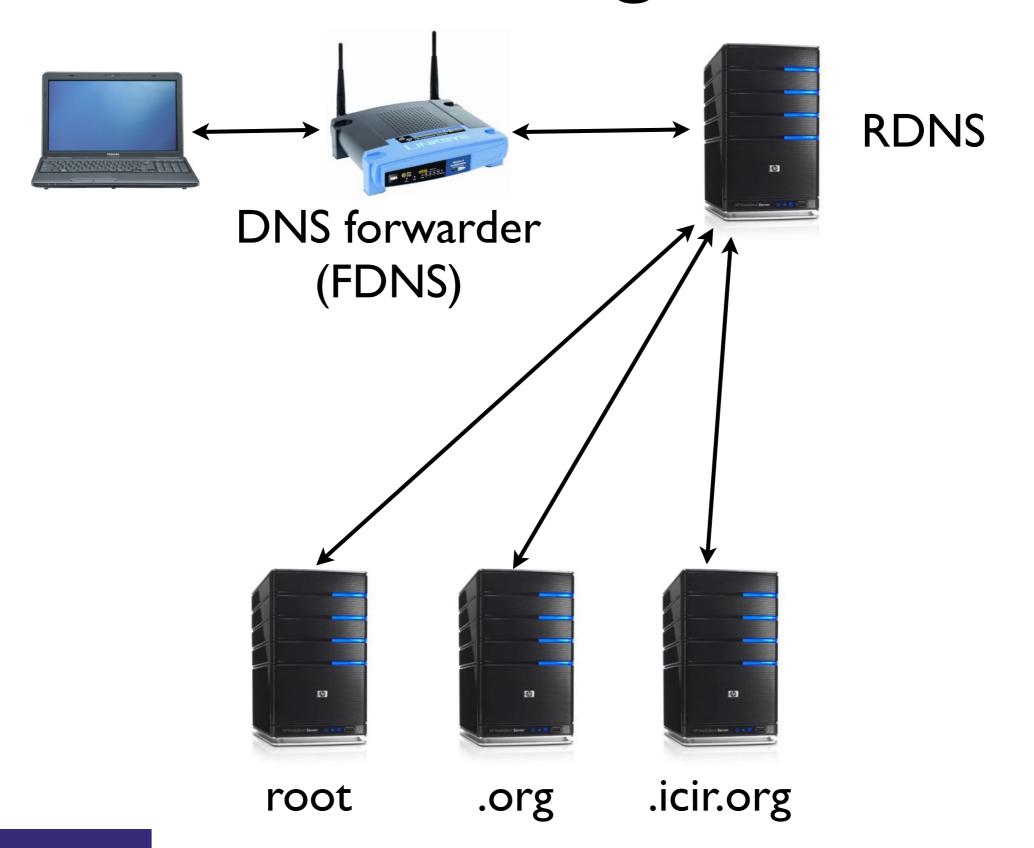
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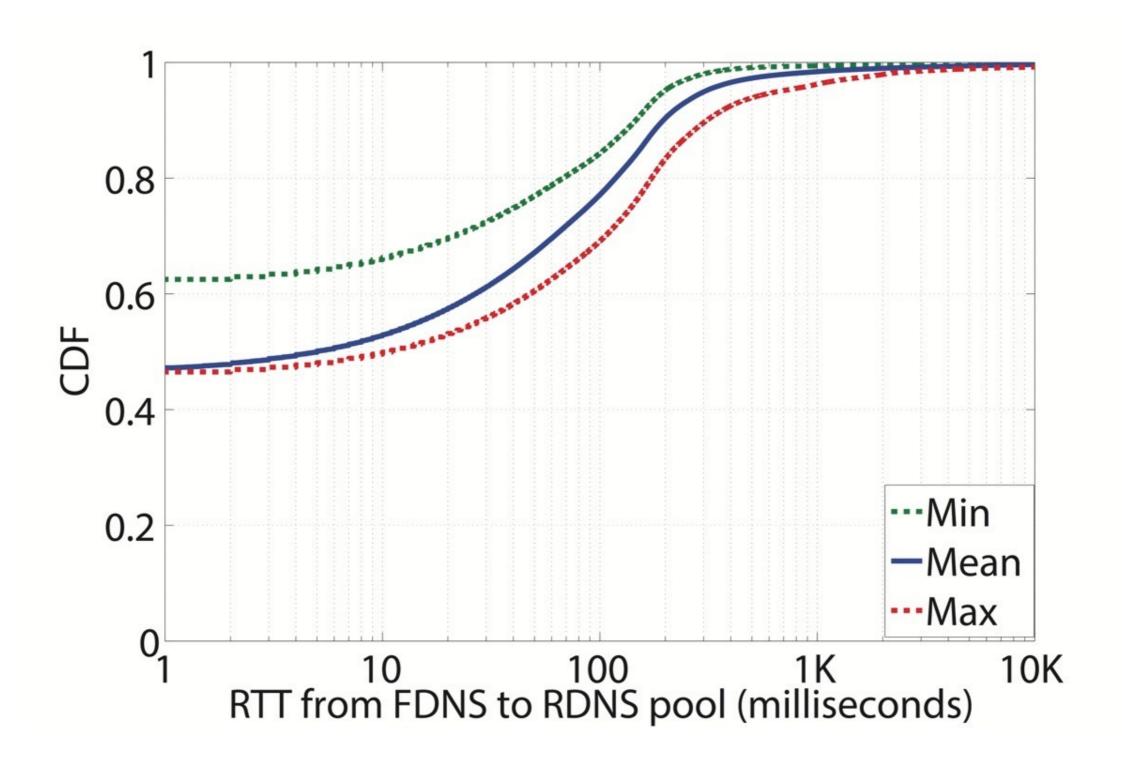
More Delegation!



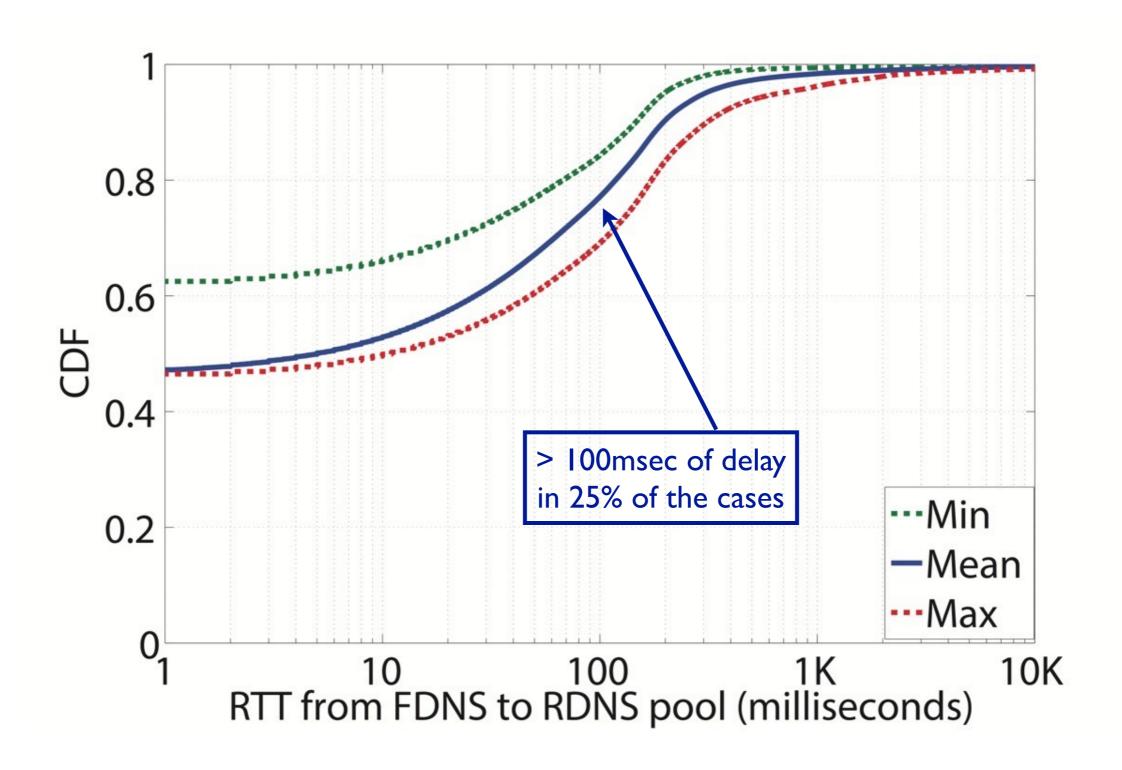
More Delegation!



FDNS-to-RDNS Delay



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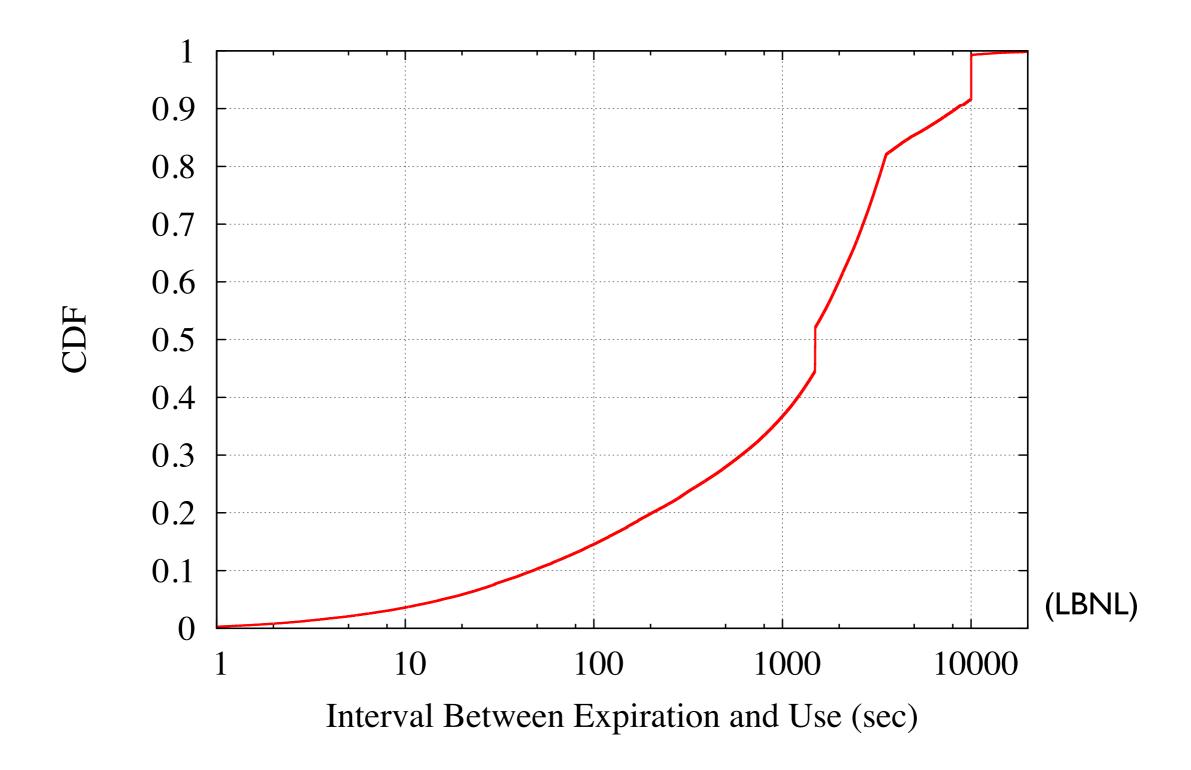
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- E.g., clients do not always adhere to the TTL

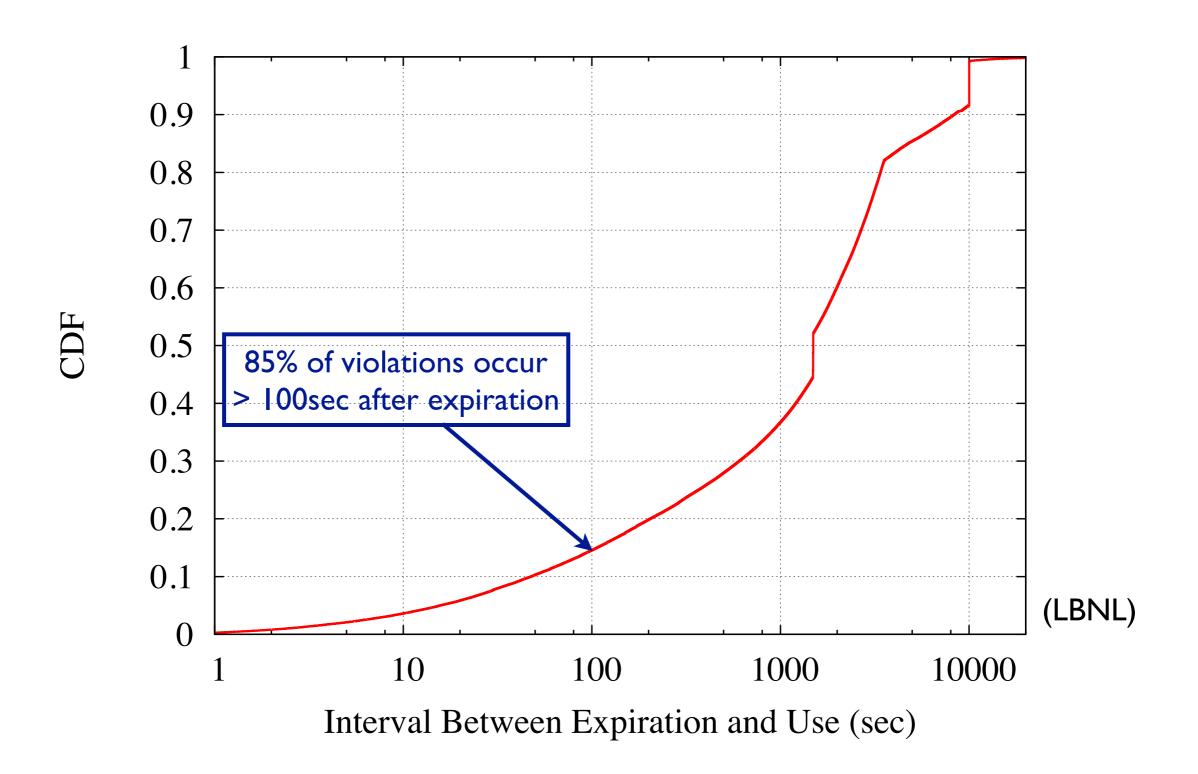
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Network	Cnns.	TTL Exp.
ICSI	443K	30%
CCZ	817K	8%
LBNL	6.IM	14%

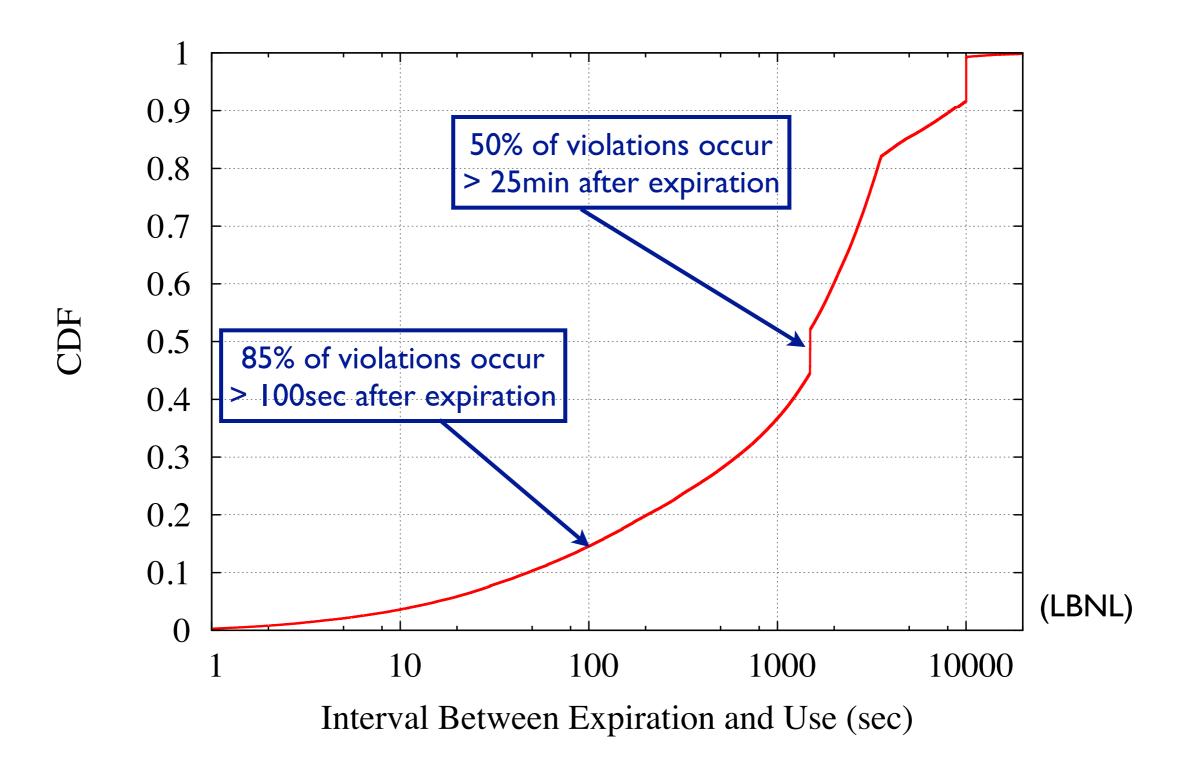
TTL Violations



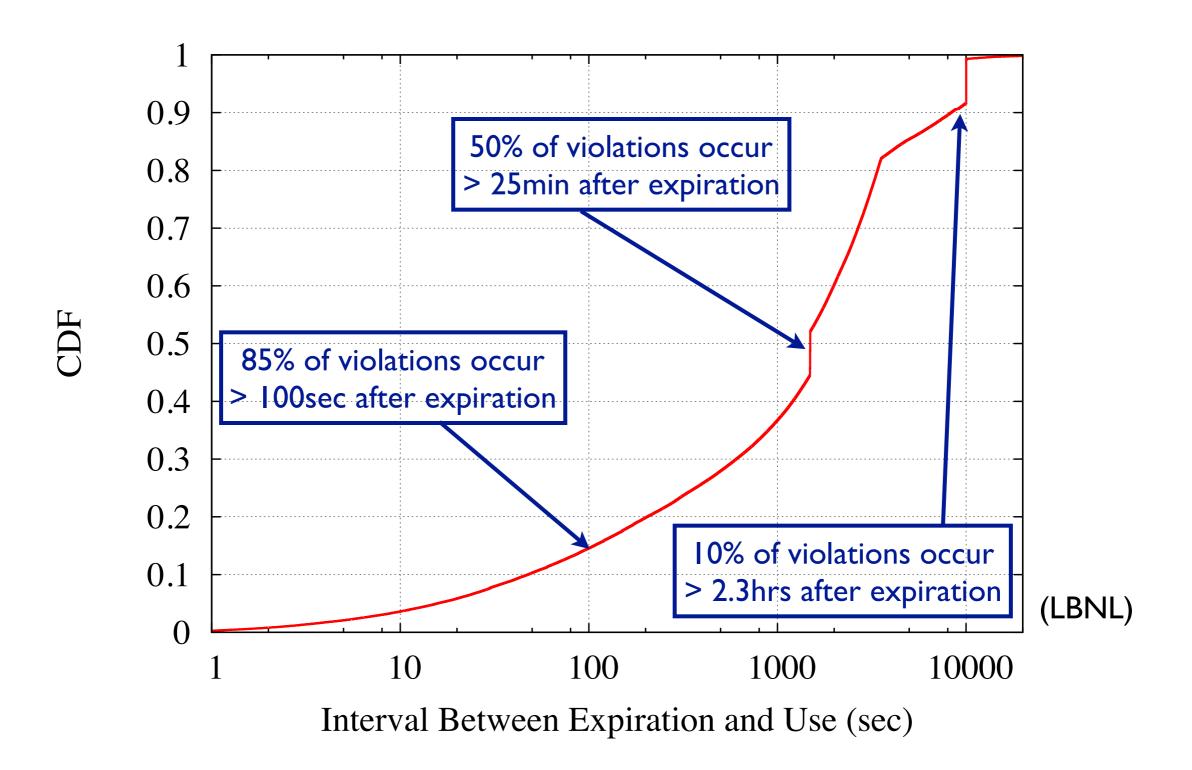
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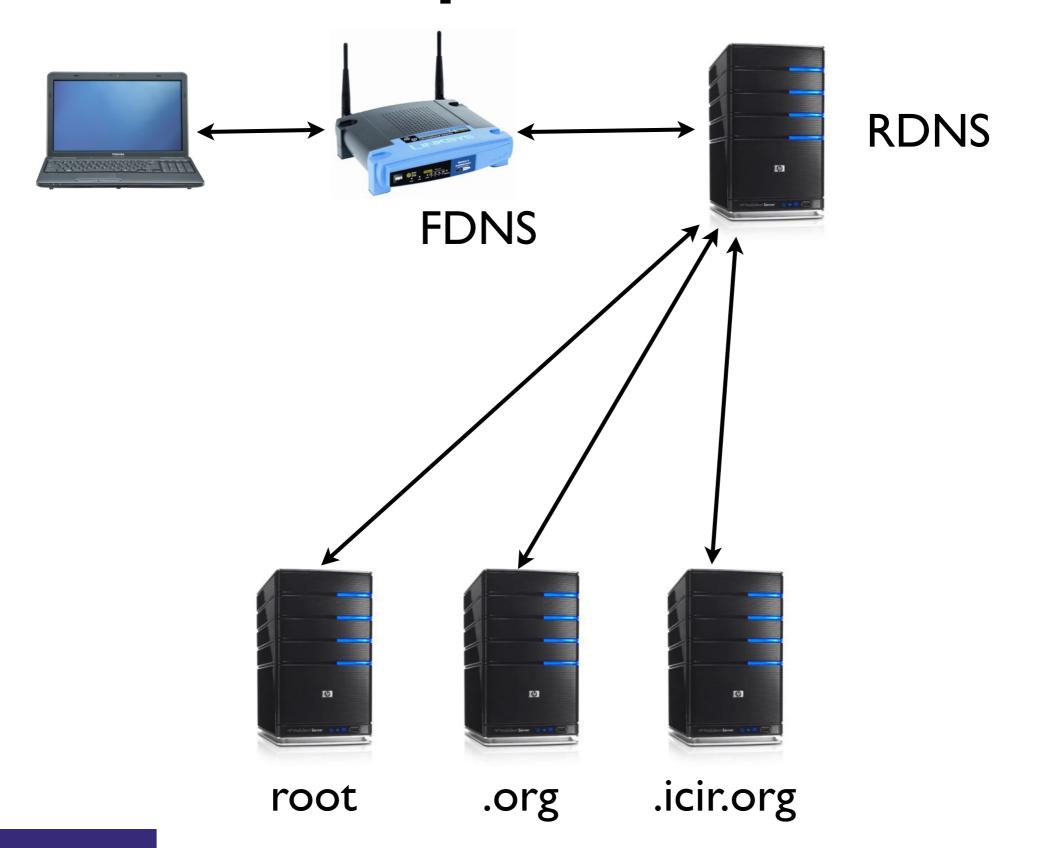
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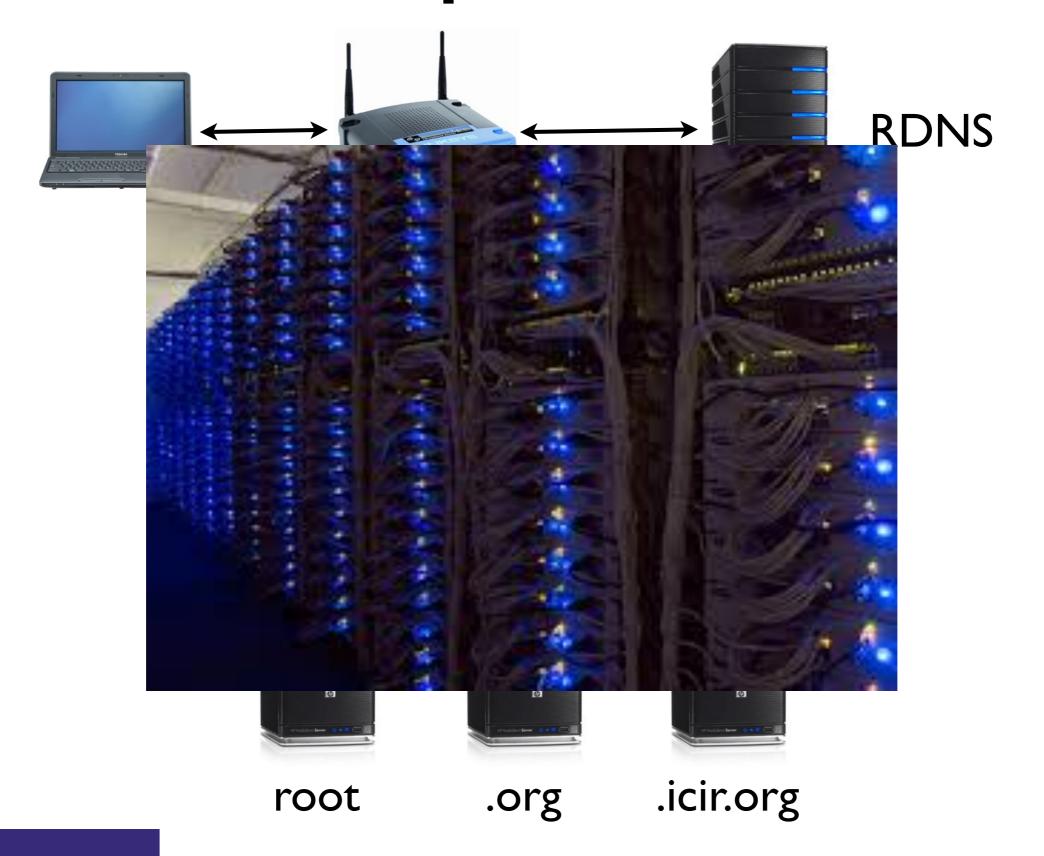
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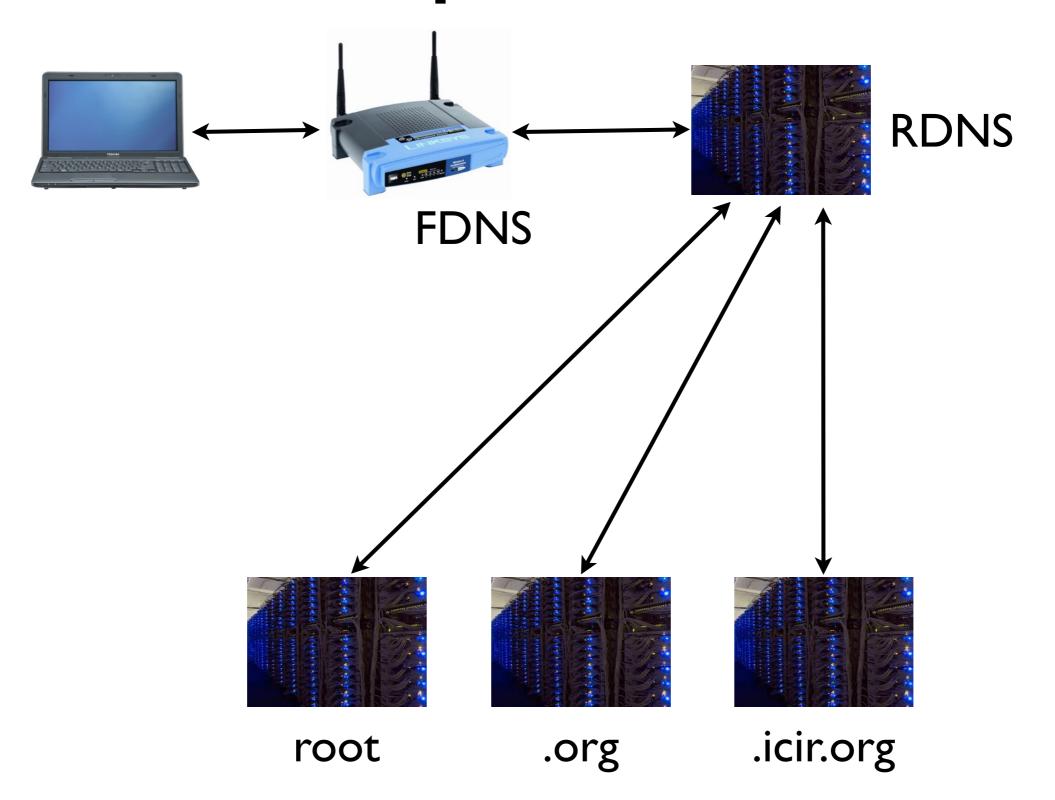
Replication

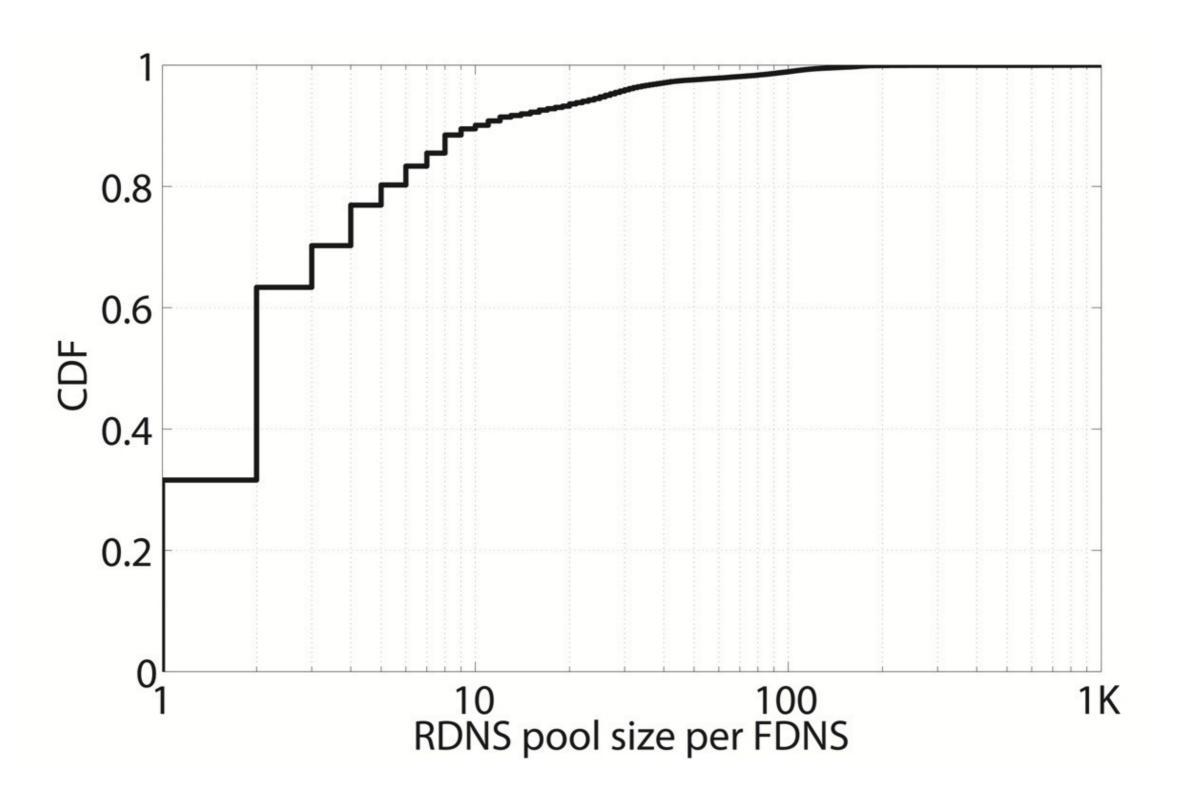


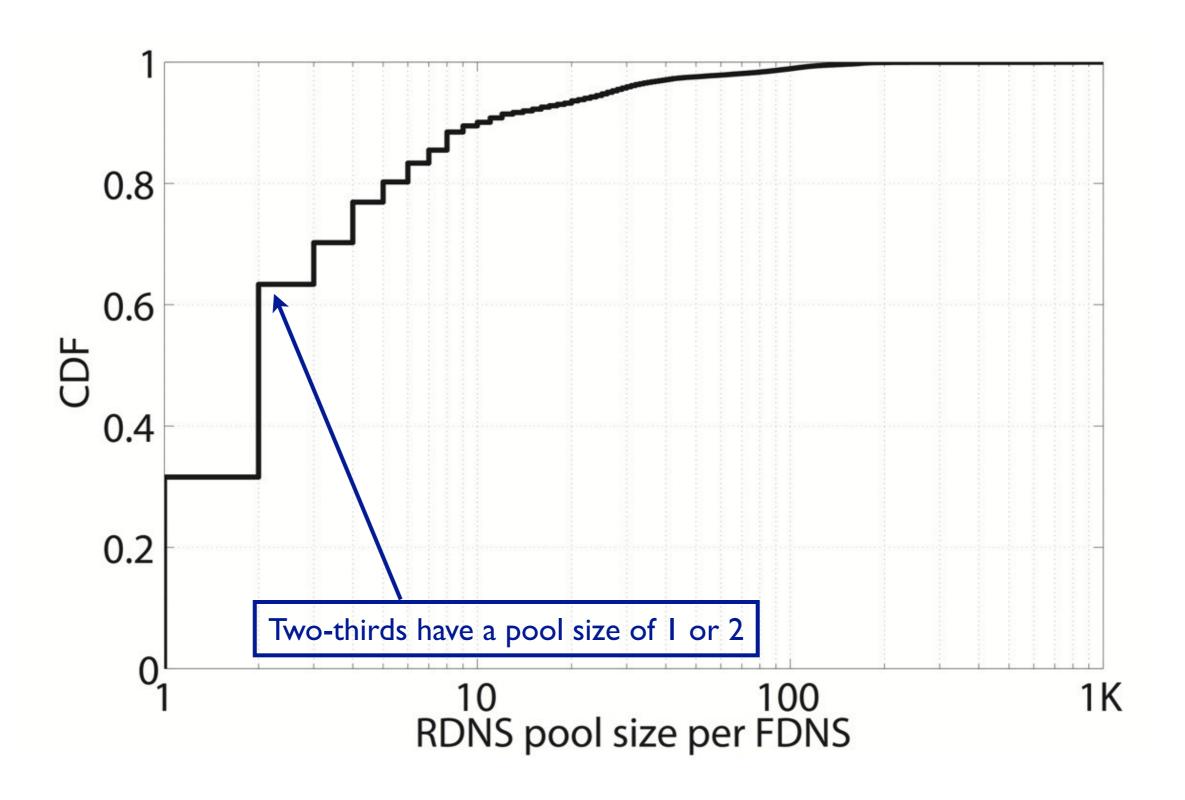
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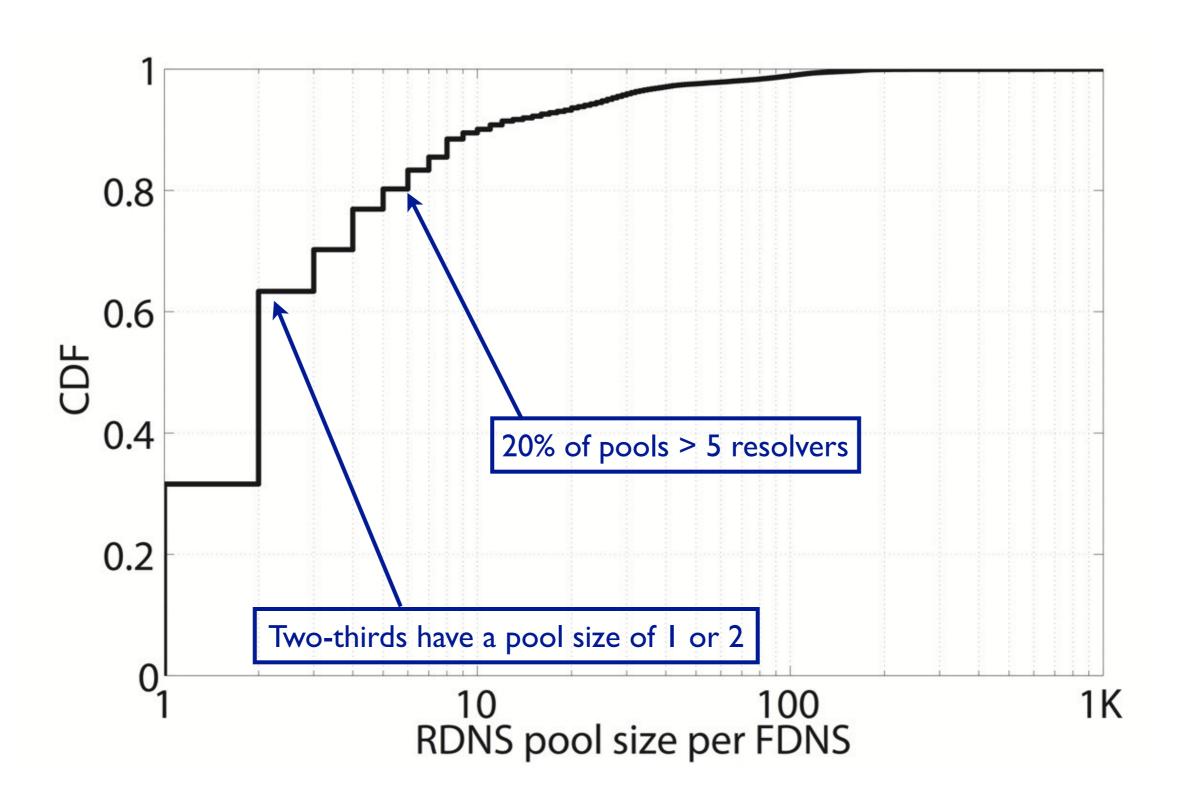


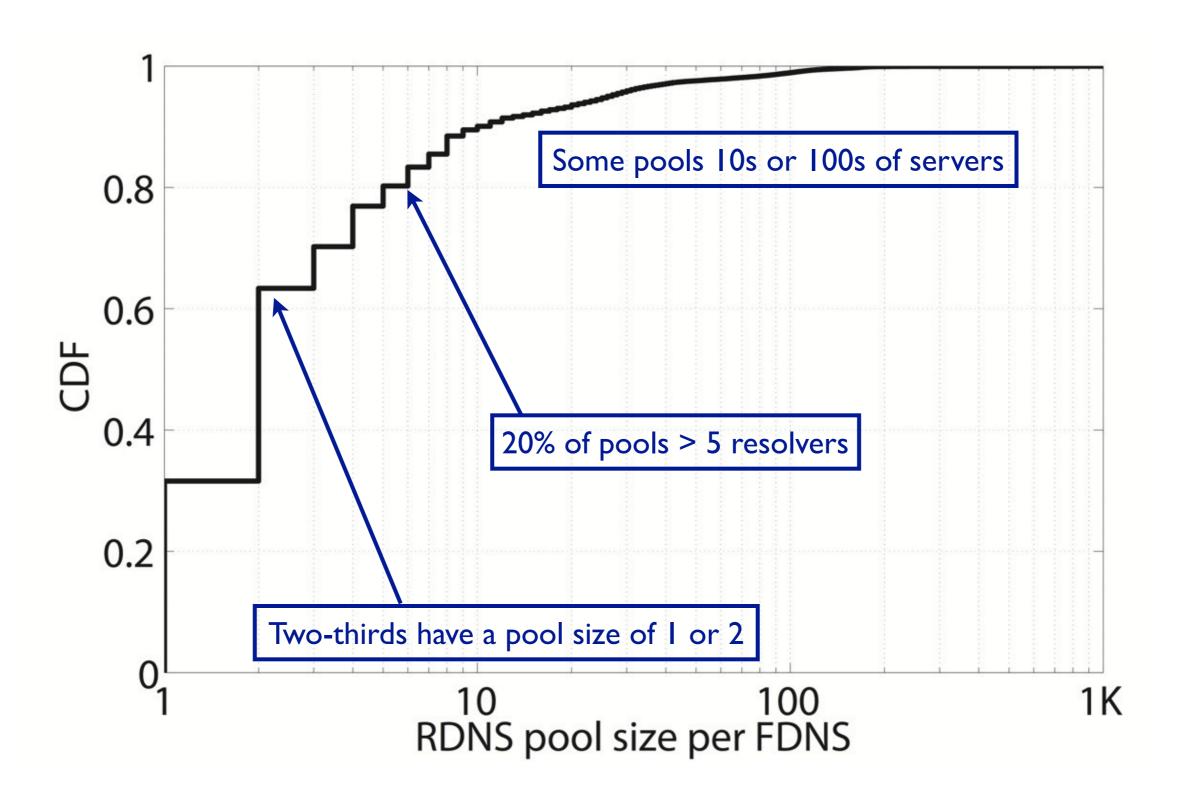
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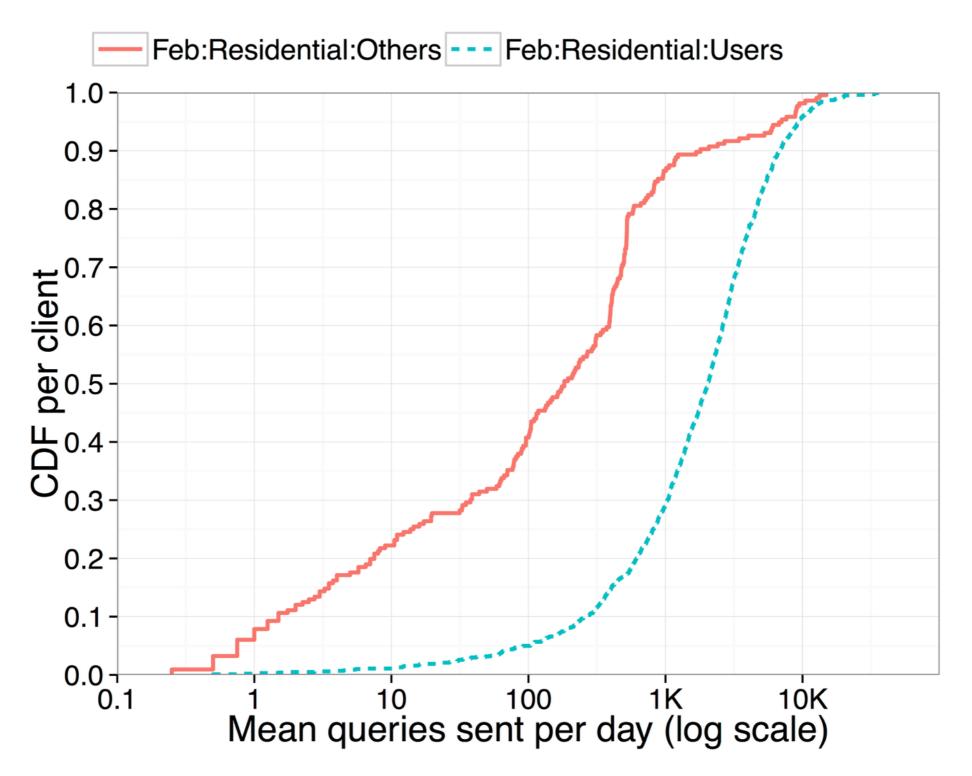




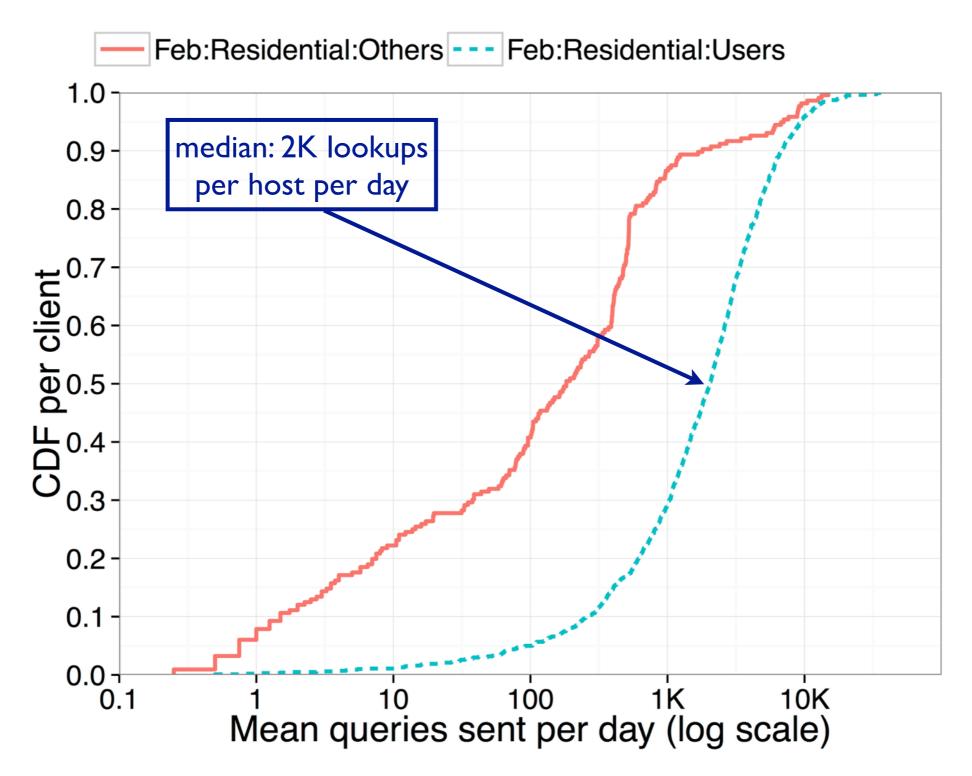
User Behavior

- Moving away from the infrastructure ...
- How about user behavior?

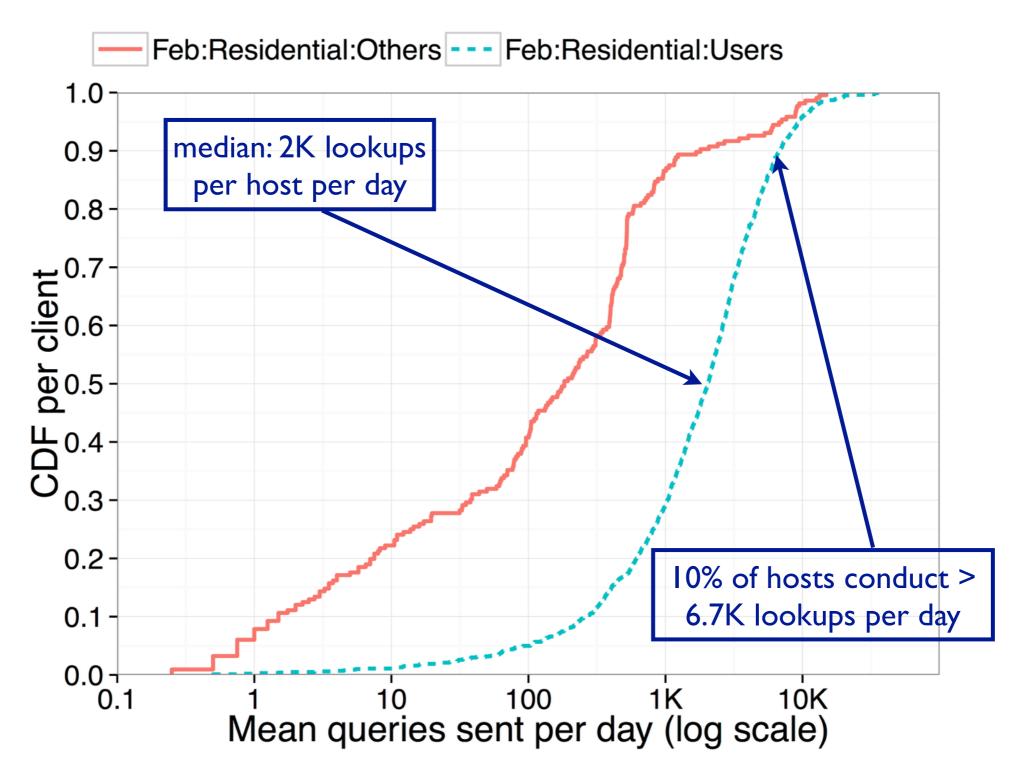
Query Rate



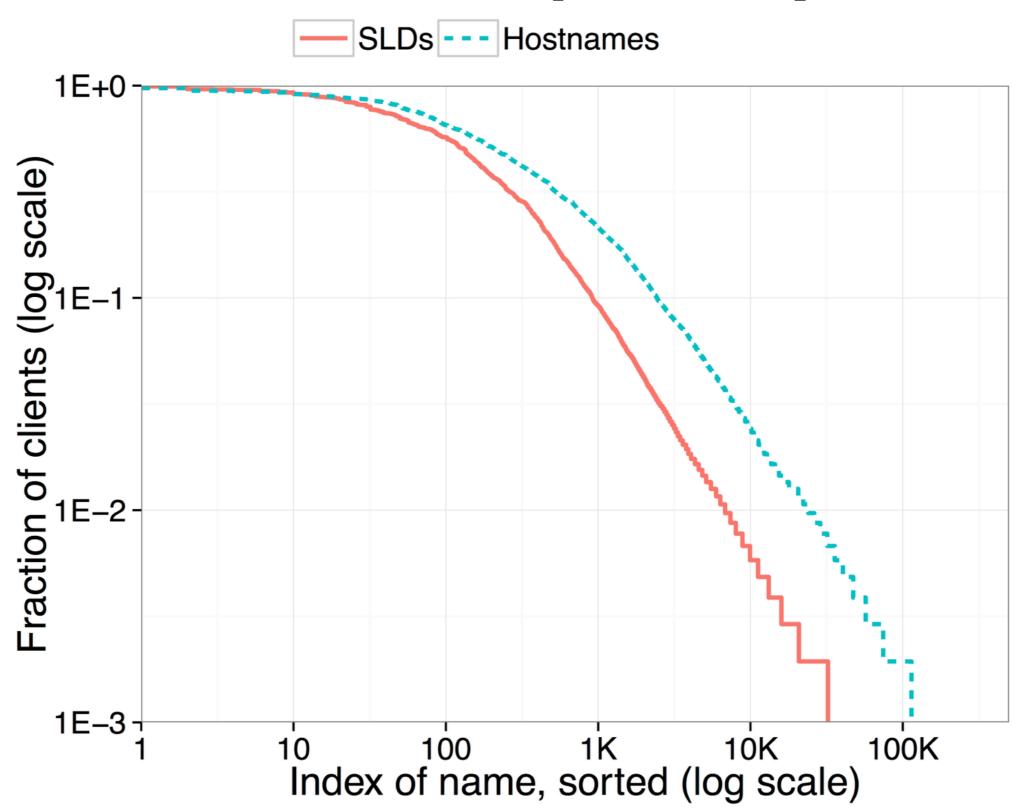
Query Rate



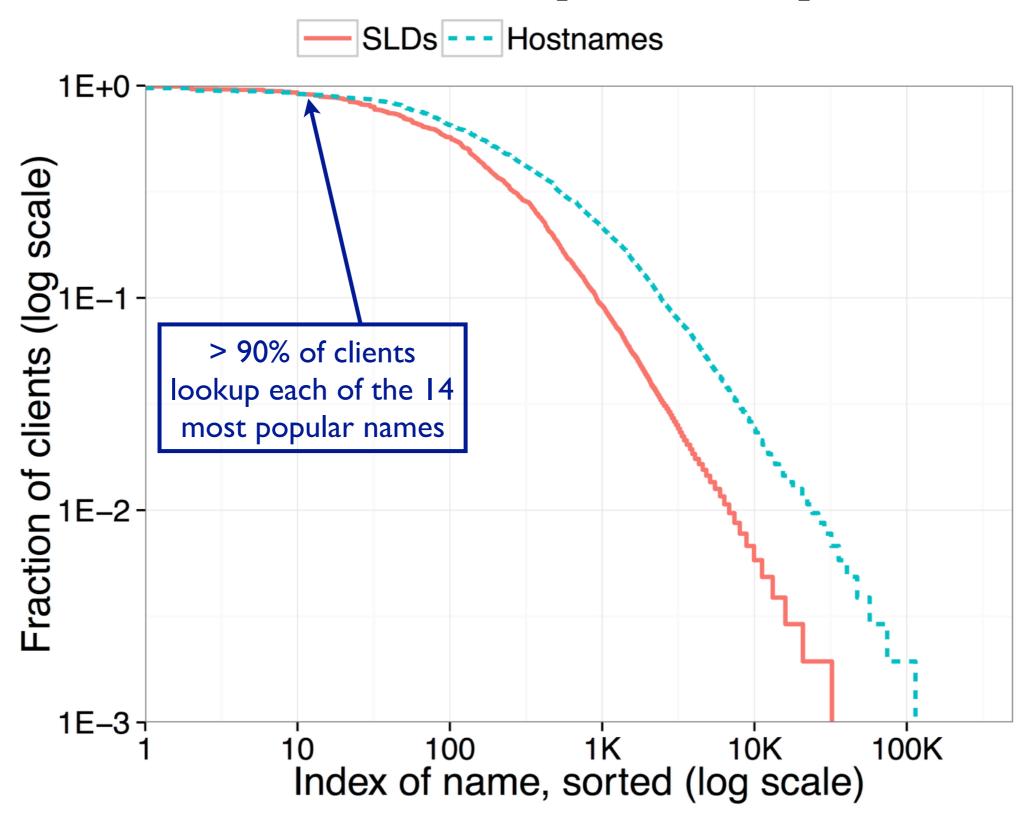
Query Rate



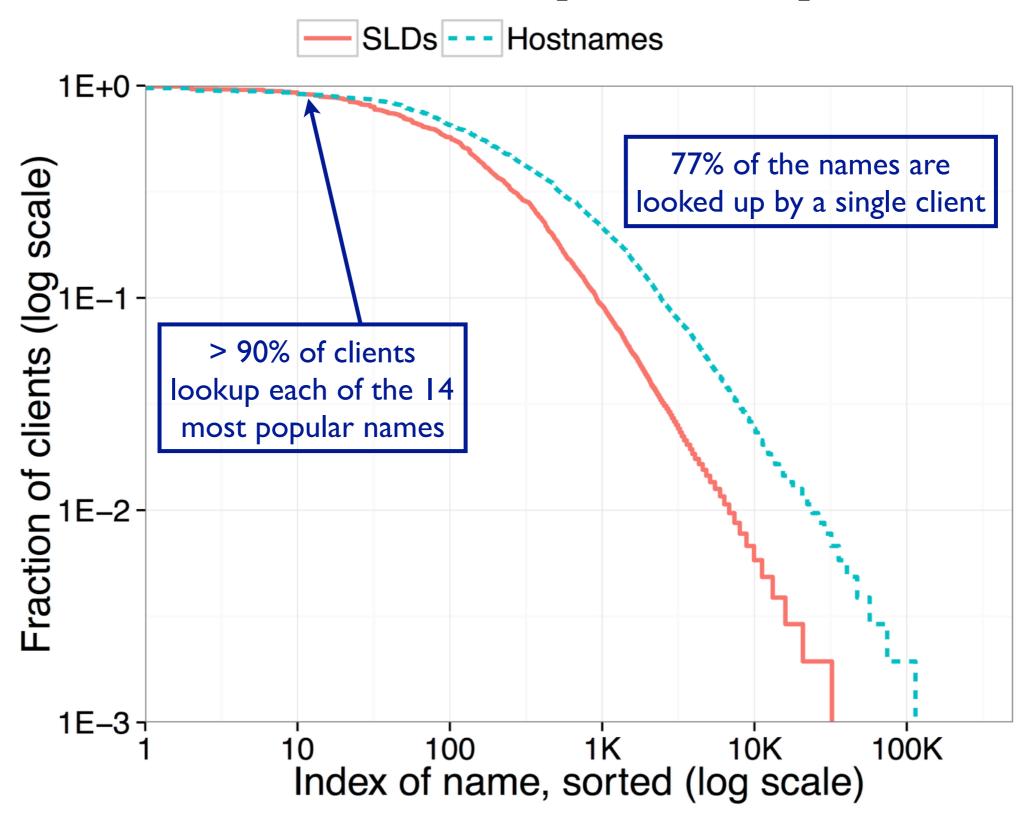
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- Curse: security issues
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 Blessing: provide us with myriad measurement vantage points

 15M open resolvers in 2010 (Leonard & Loguinov, IMC 2010)

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- 15M open resolvers in 2017 (openresolverproject.org)

DNS Security

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• DNS' key flaw:

DNS Security

- DNS' key flaw:
 - all transactions in clear text!
 - no trust, no privacy, no integrity

Security Implications

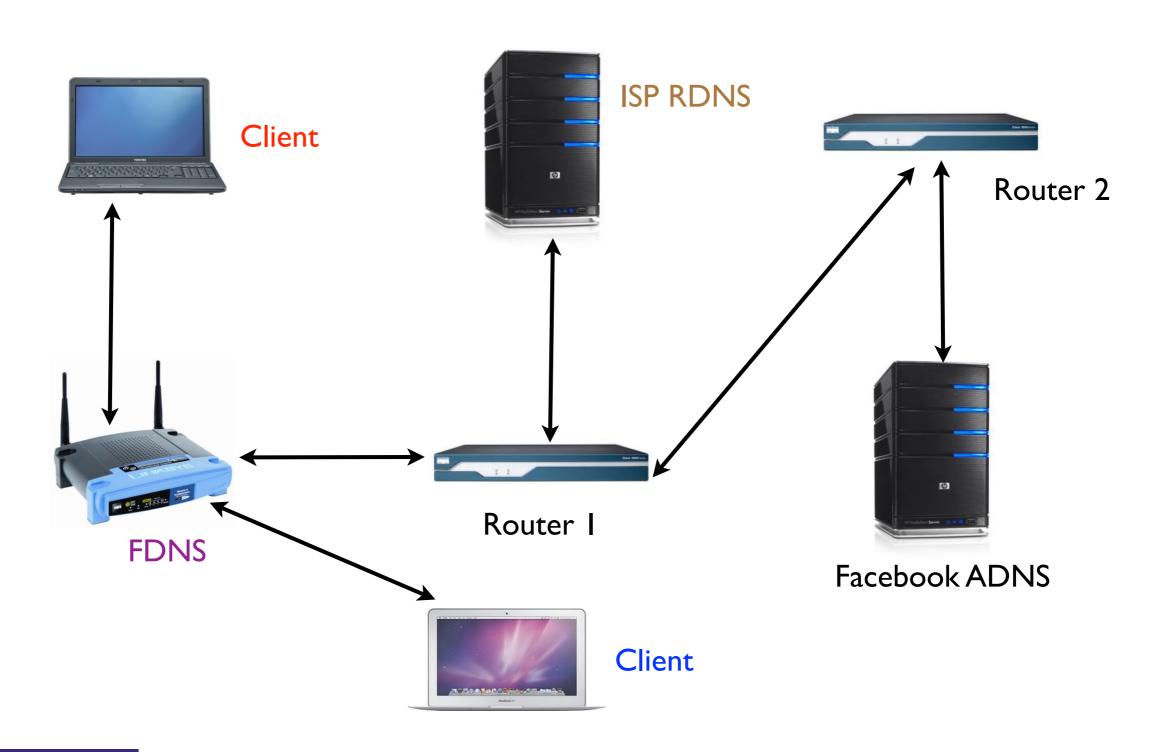
Security Implications

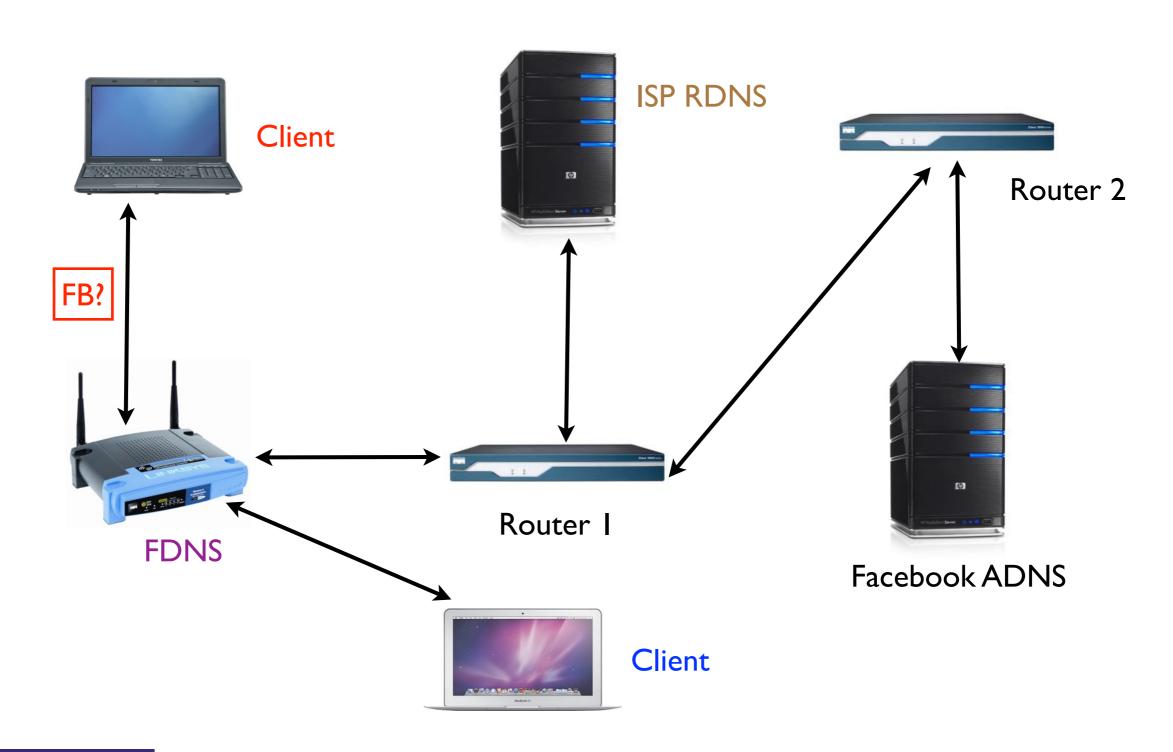
- DNS susceptible to ...
 - ... surveillance
 - ... spoofing
 - ... modification
 - ... injection
 - ... interposition

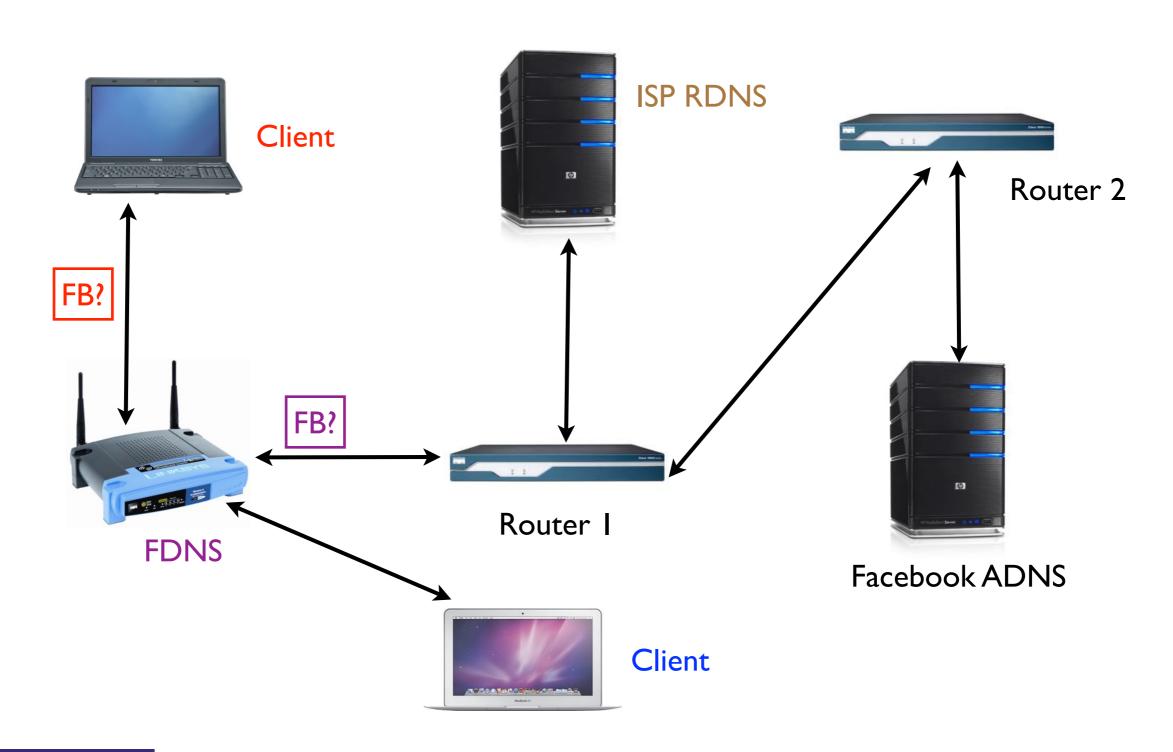
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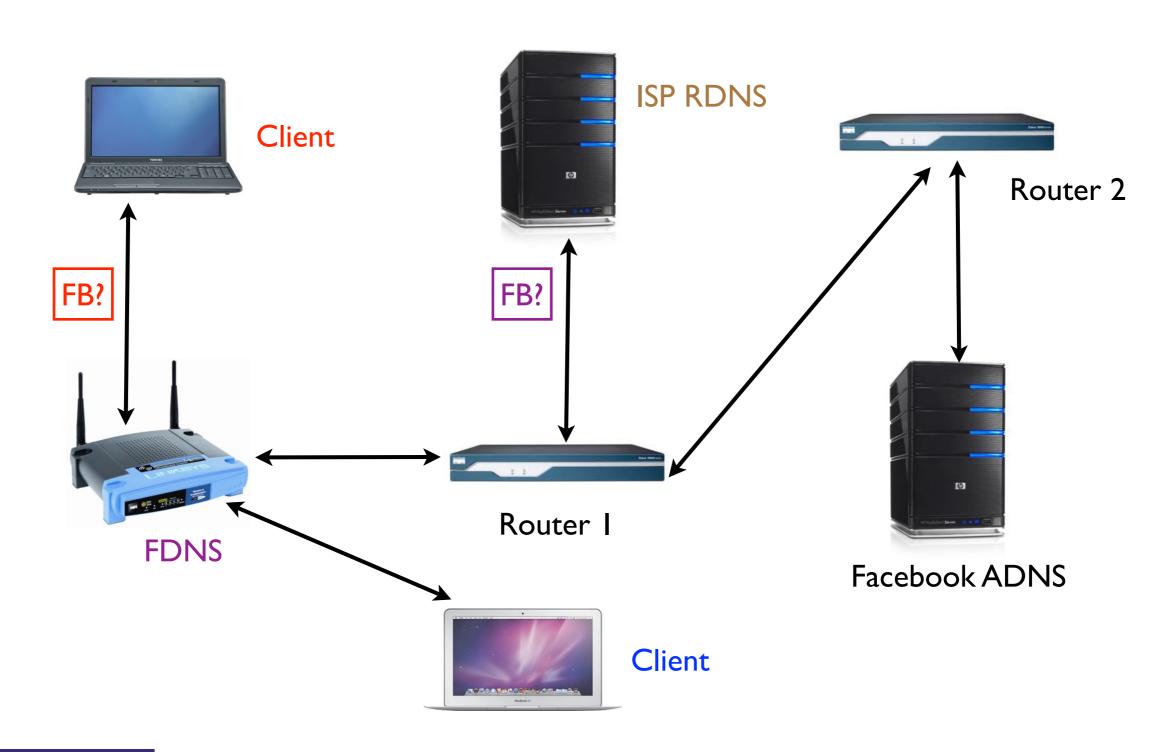
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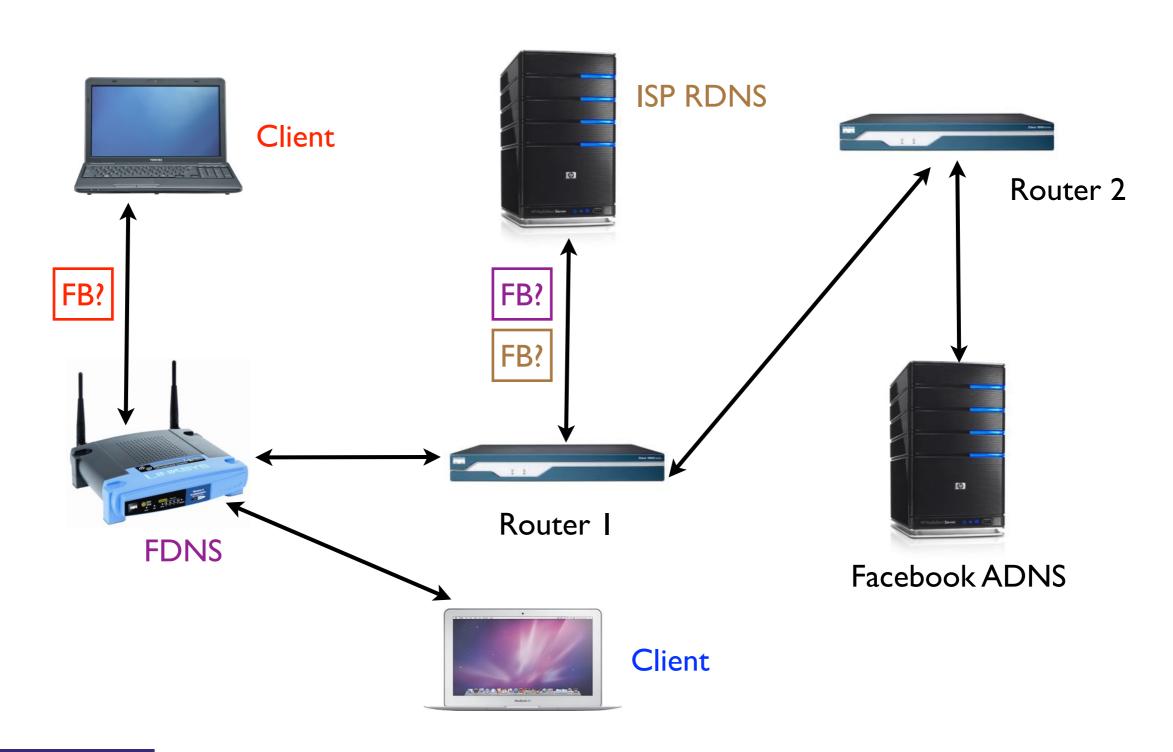
Every component of the DNS ecosystem is a potential adversary

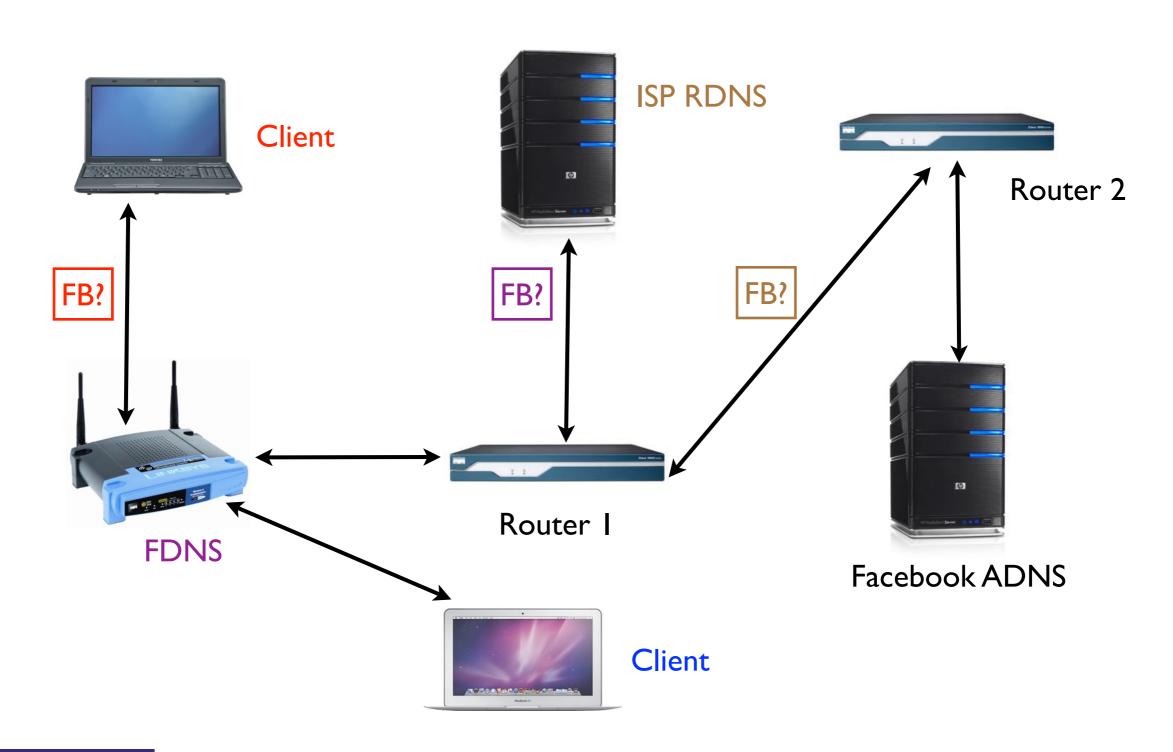


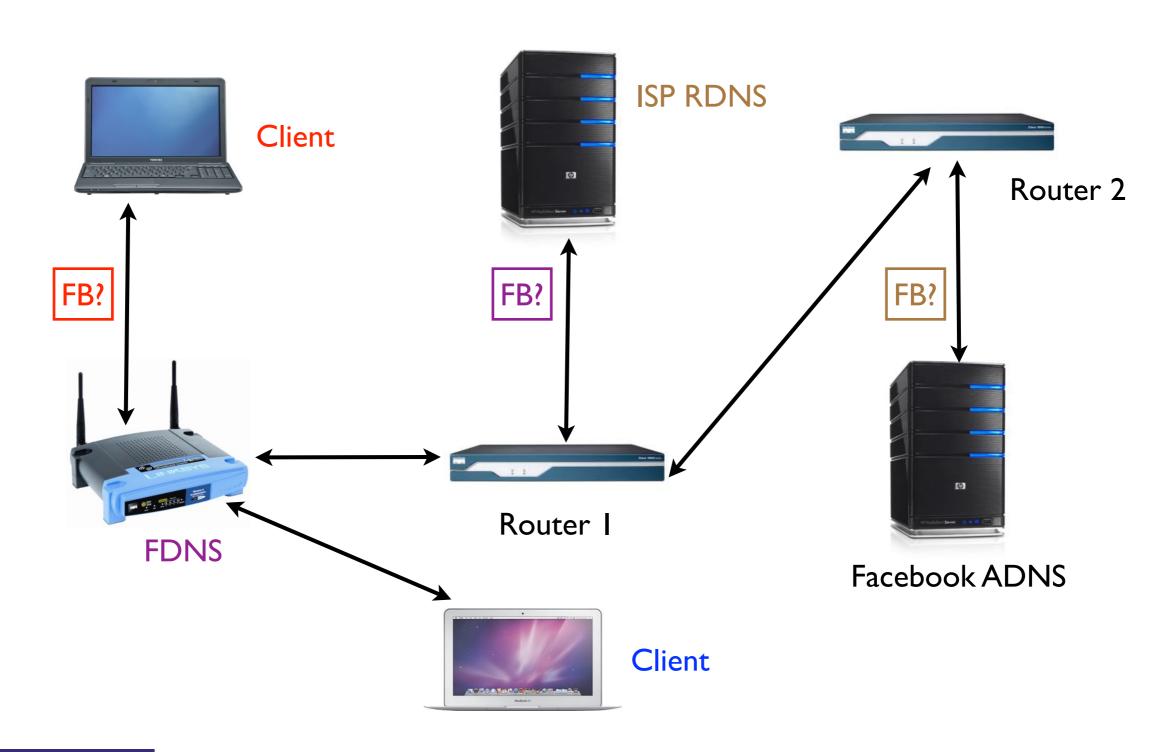


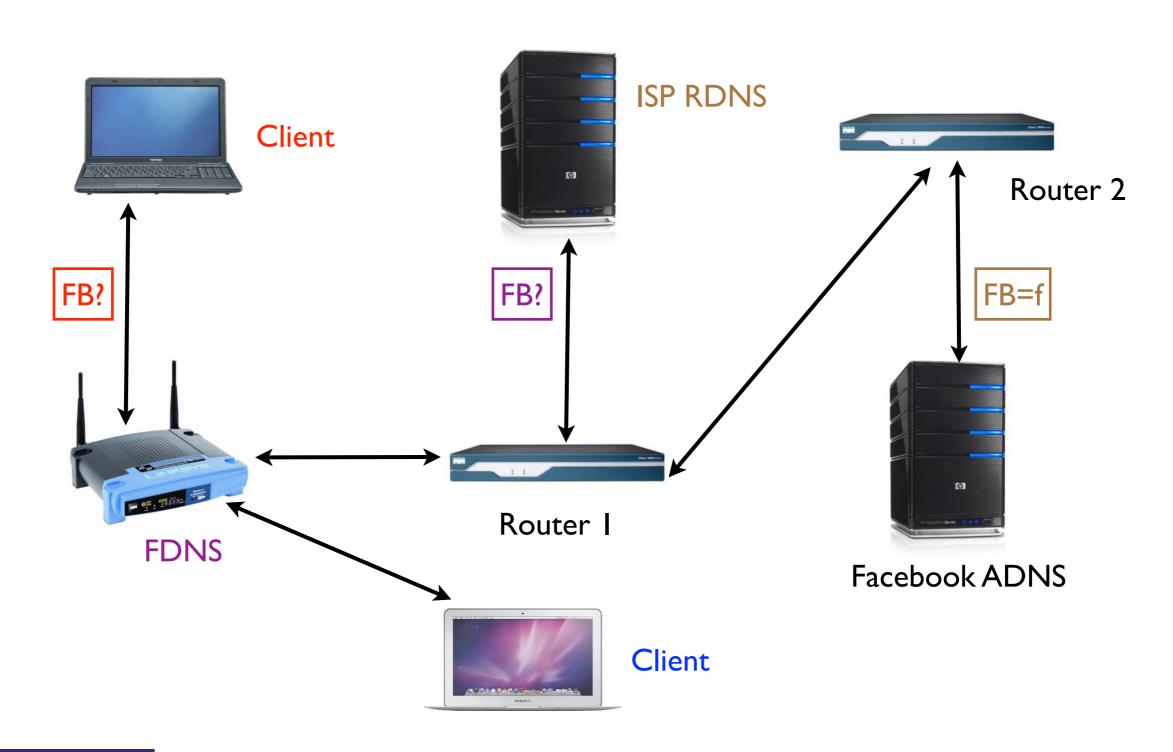


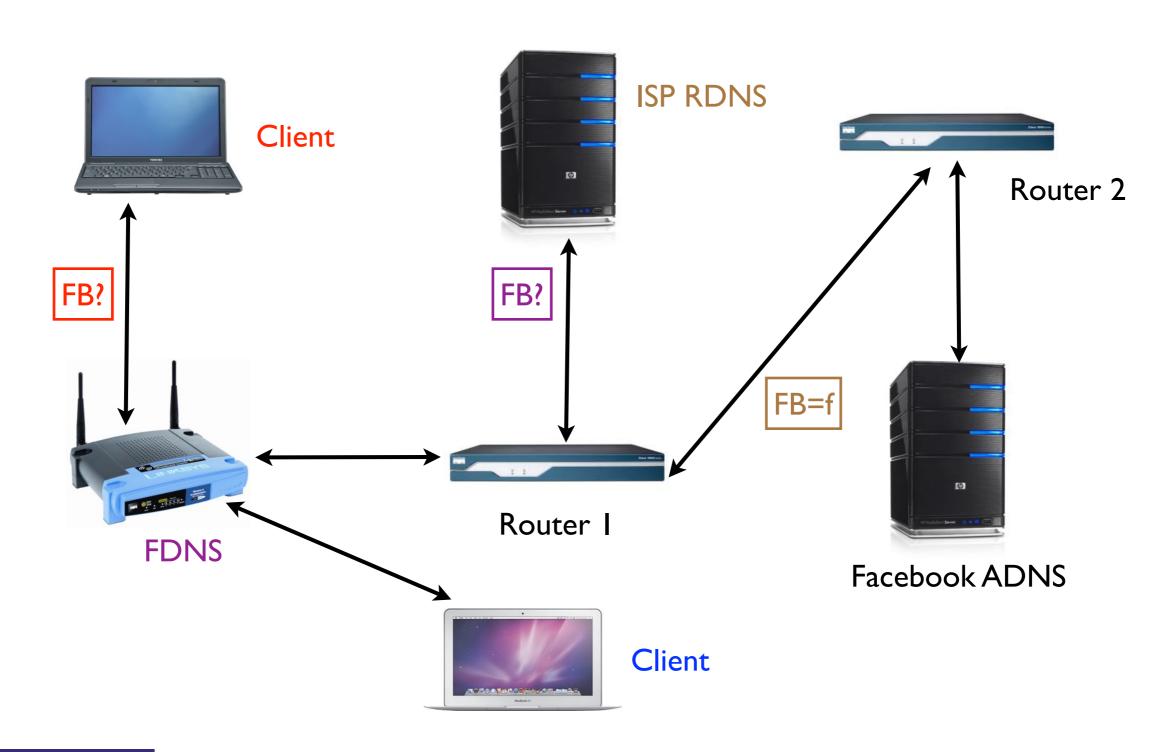


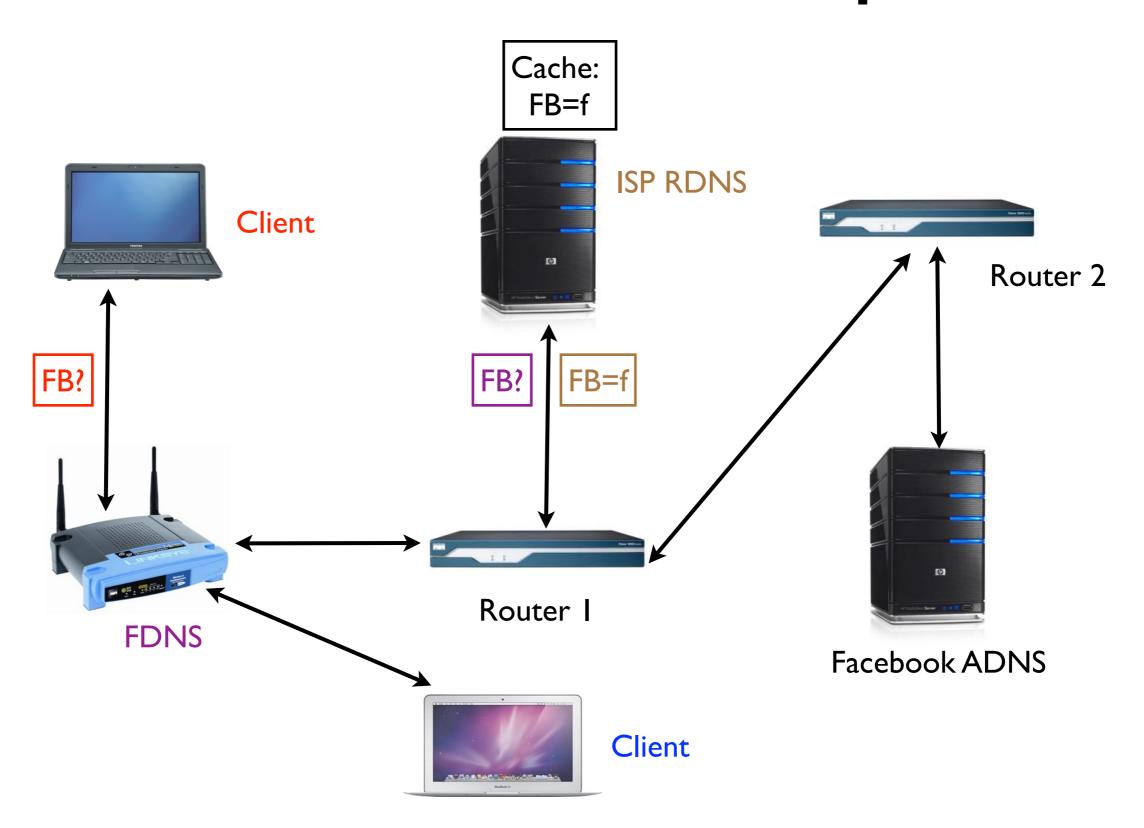


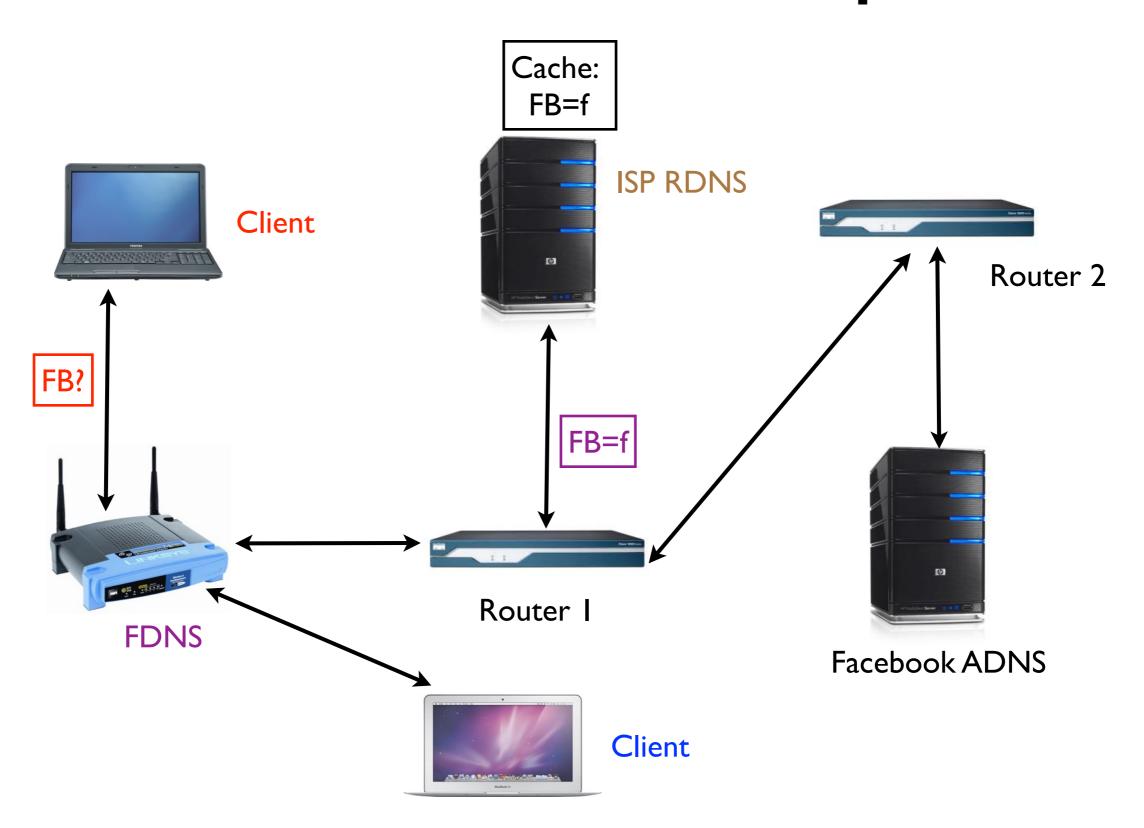


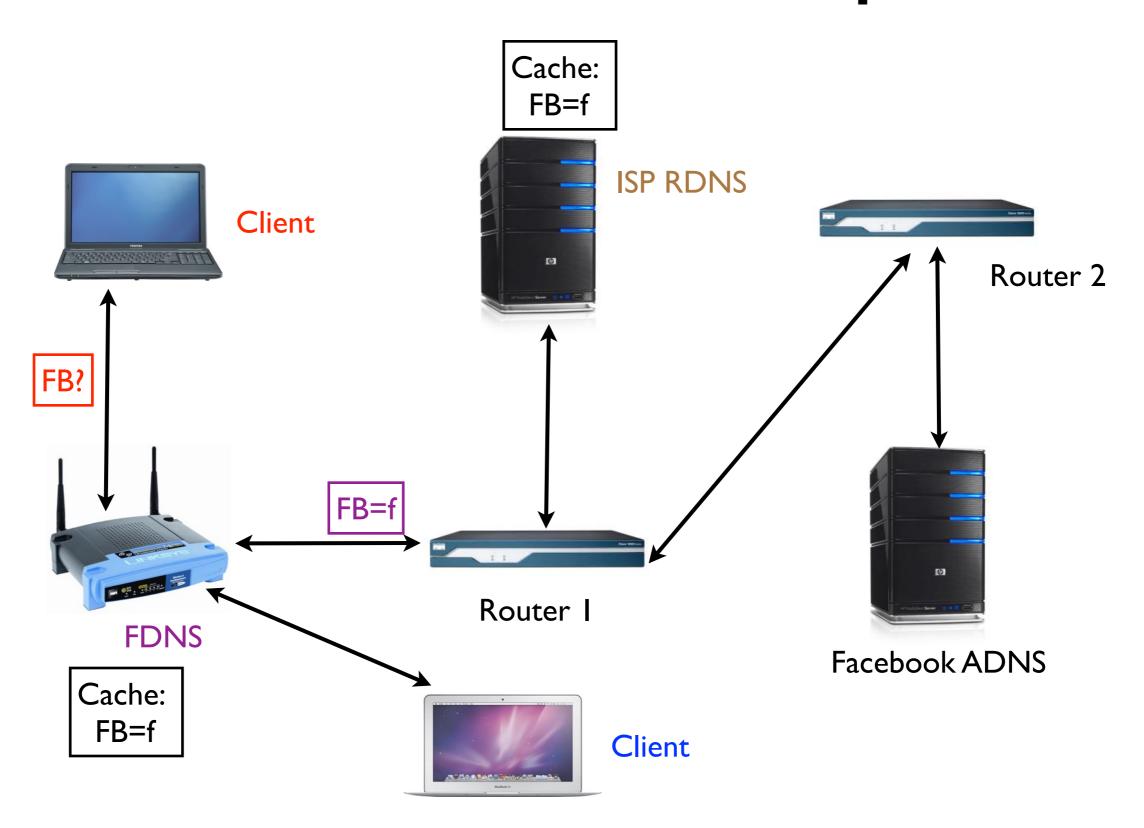


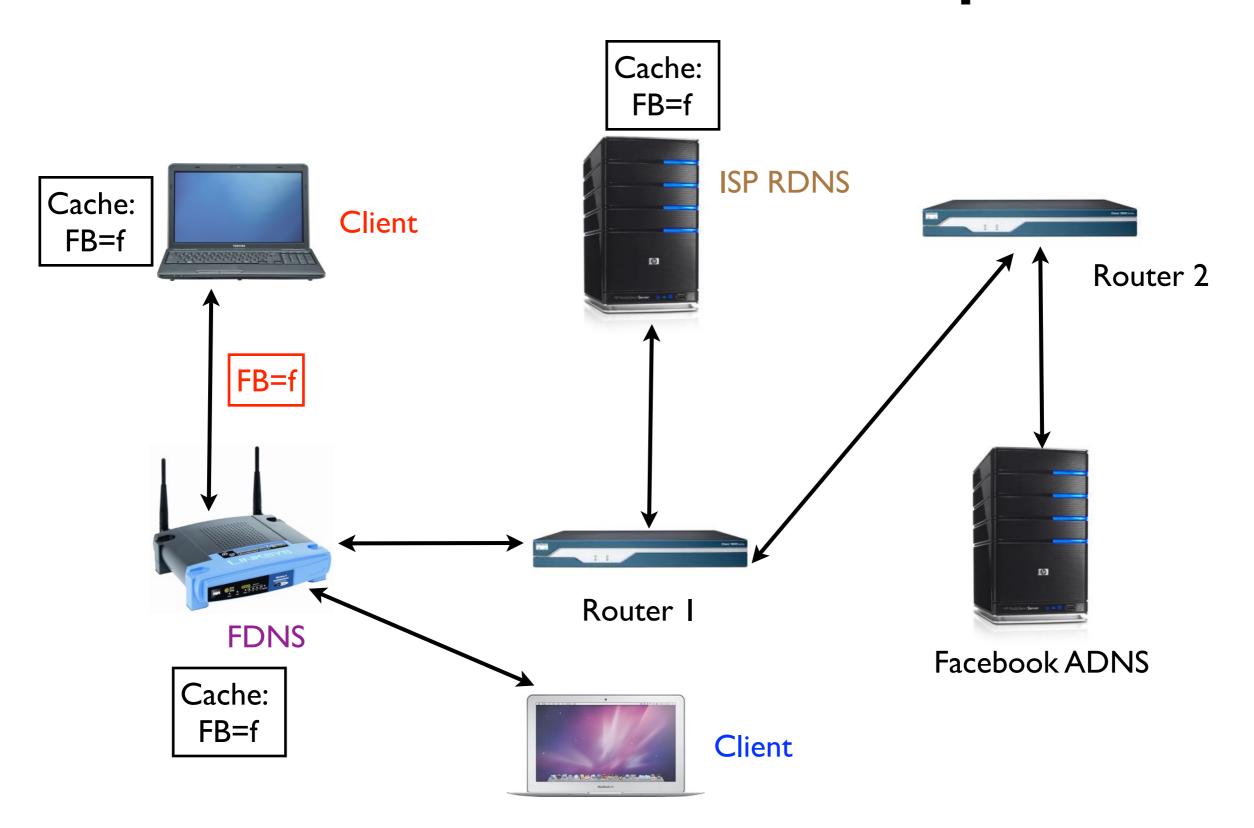


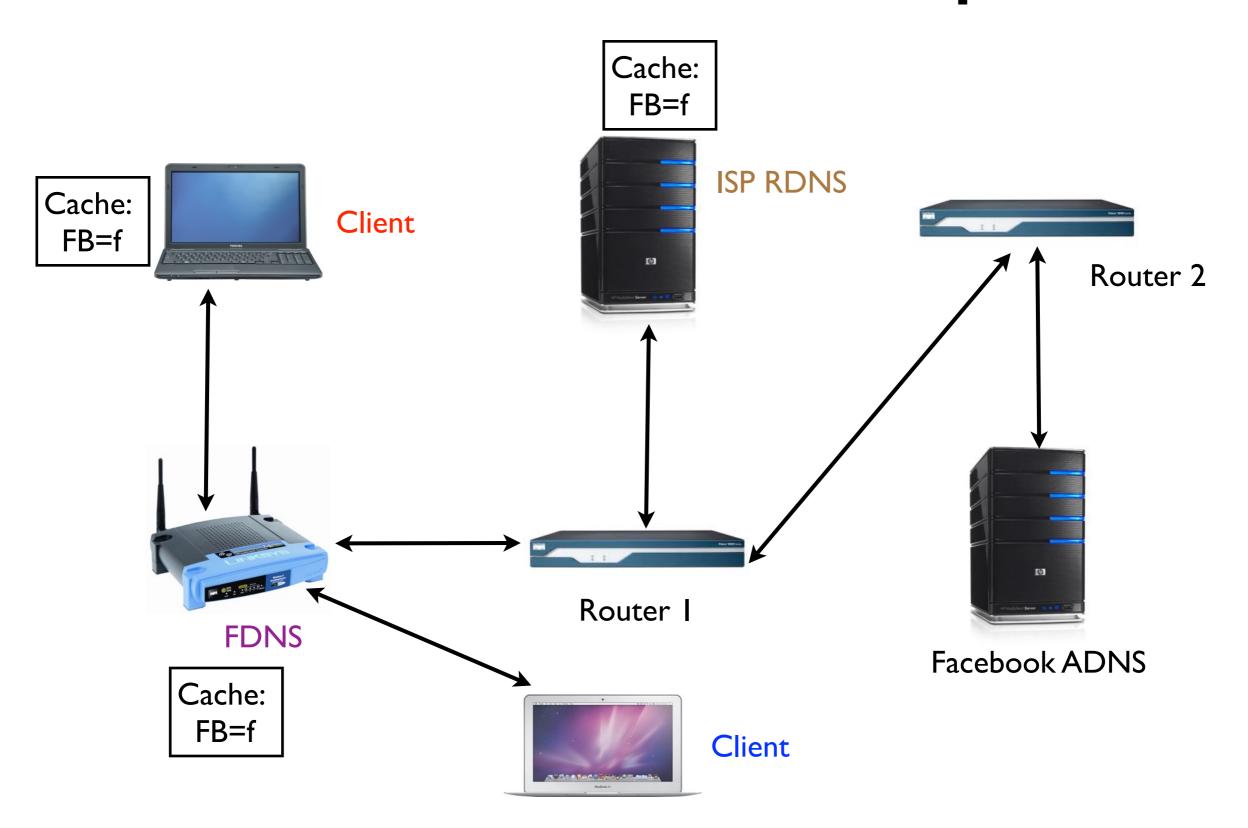


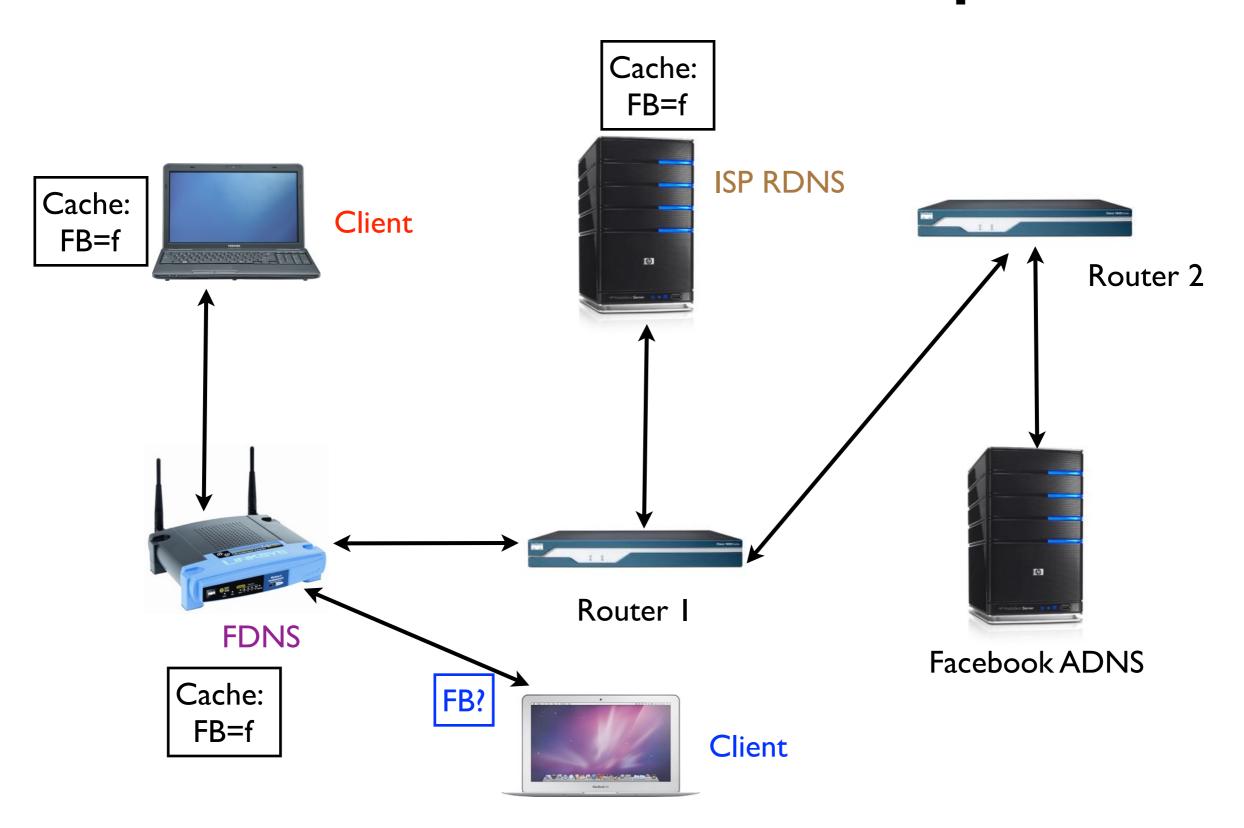


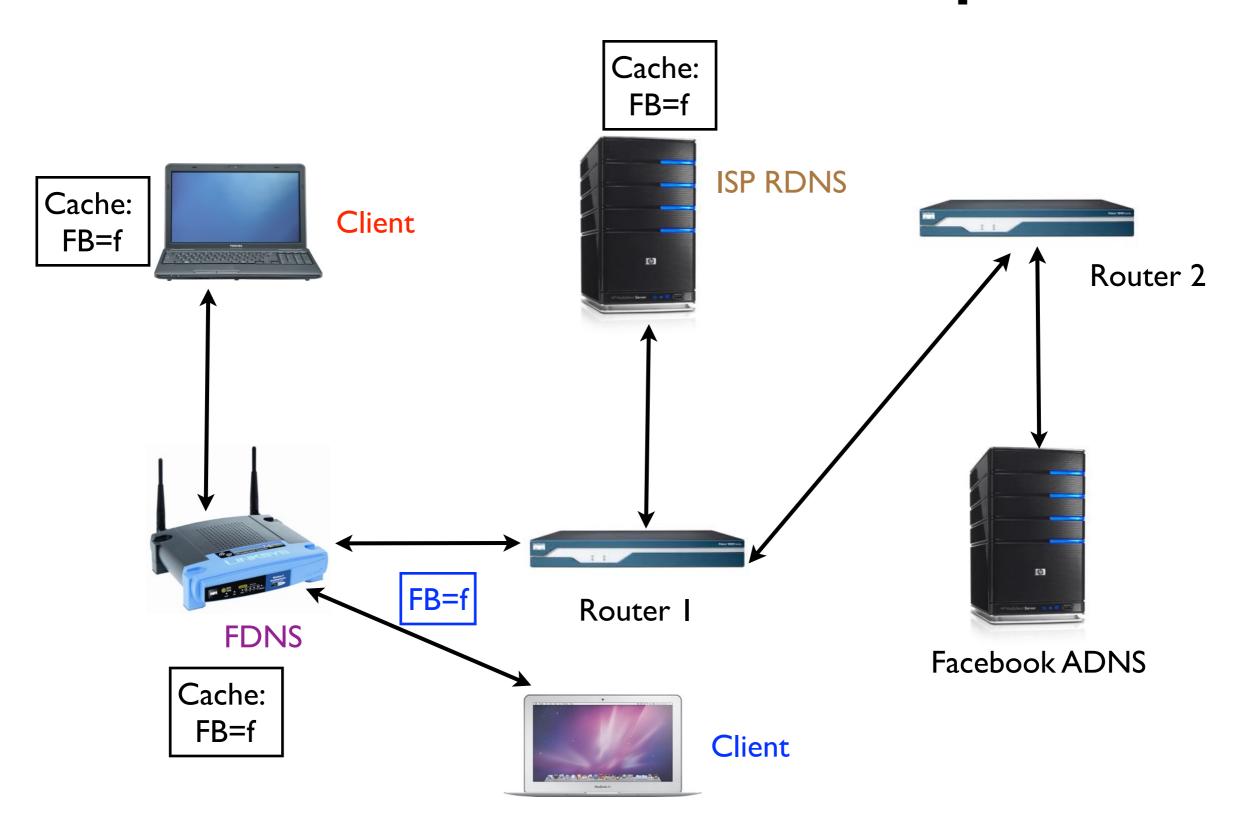


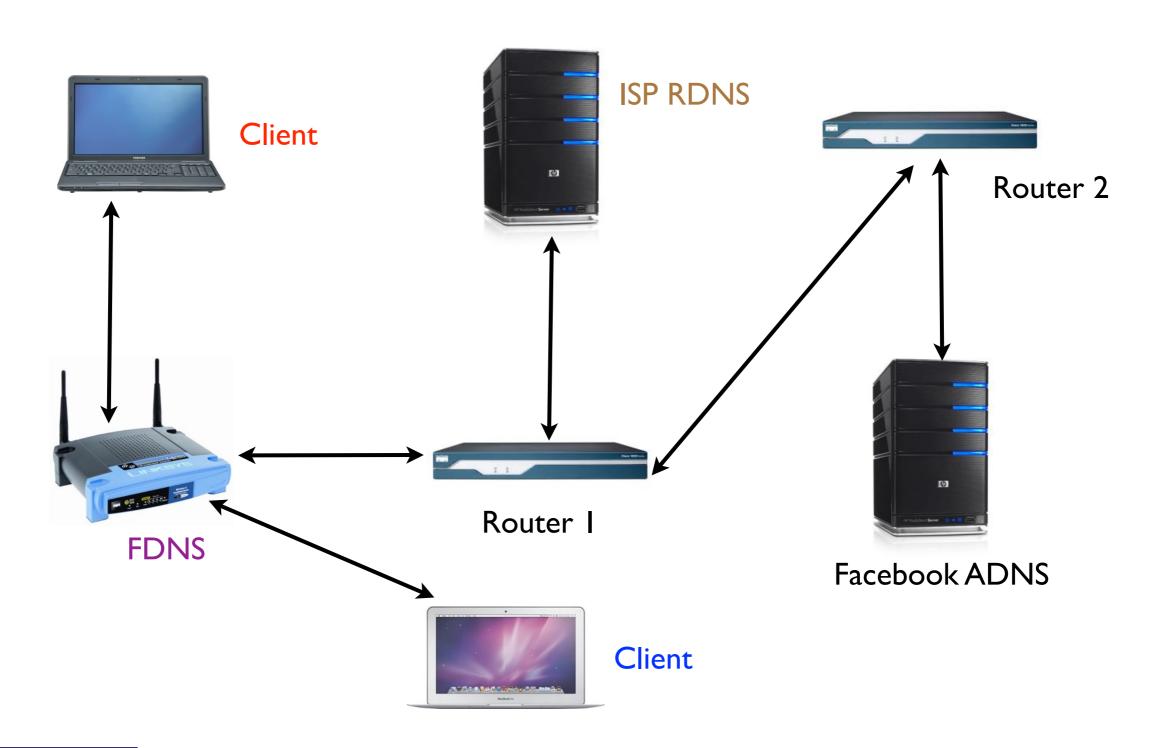


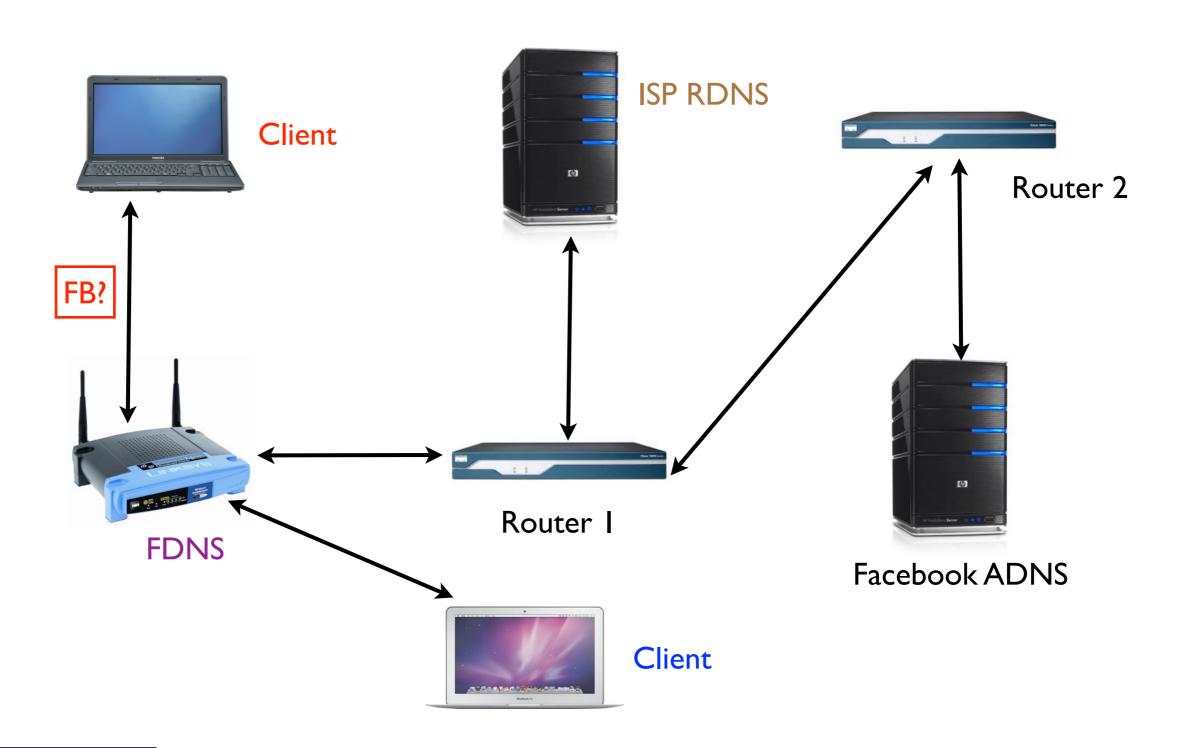


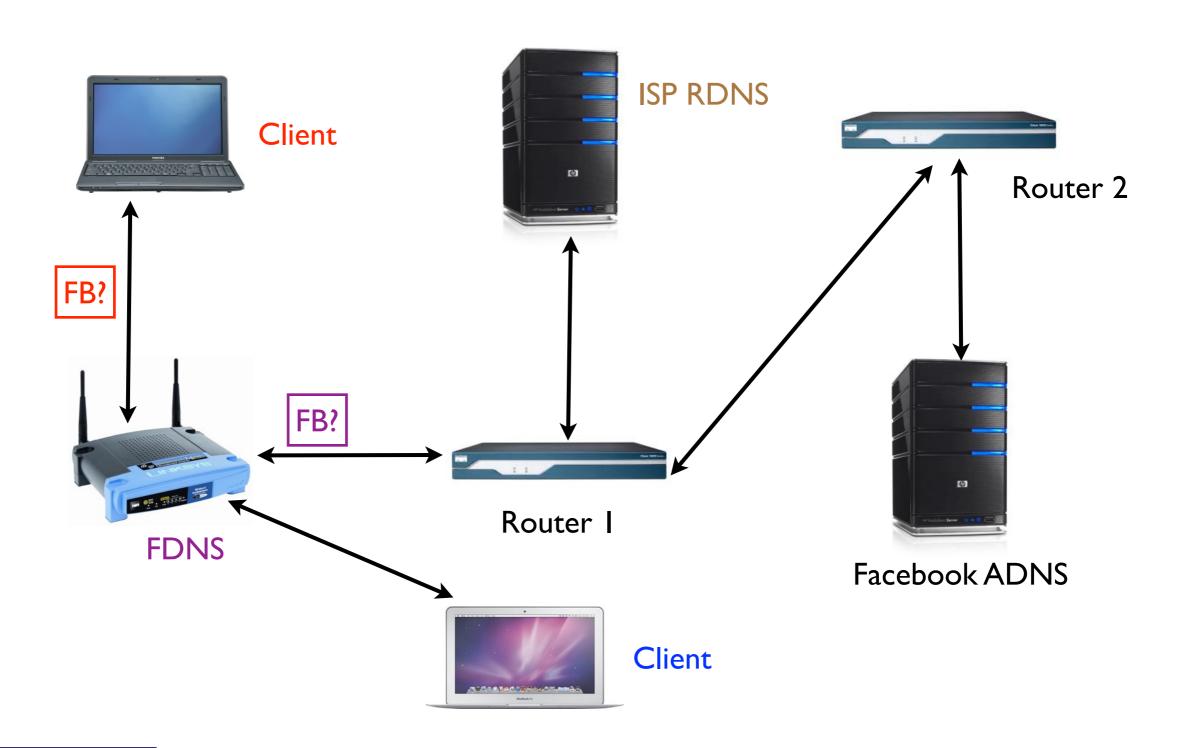


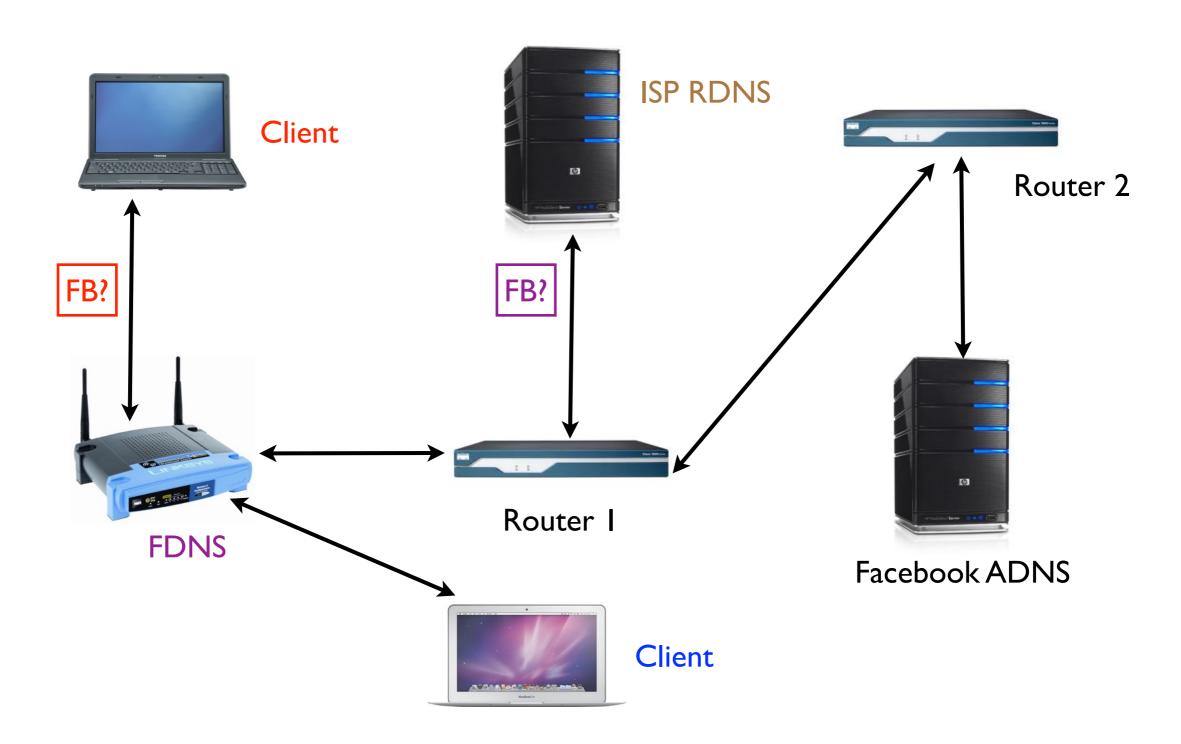


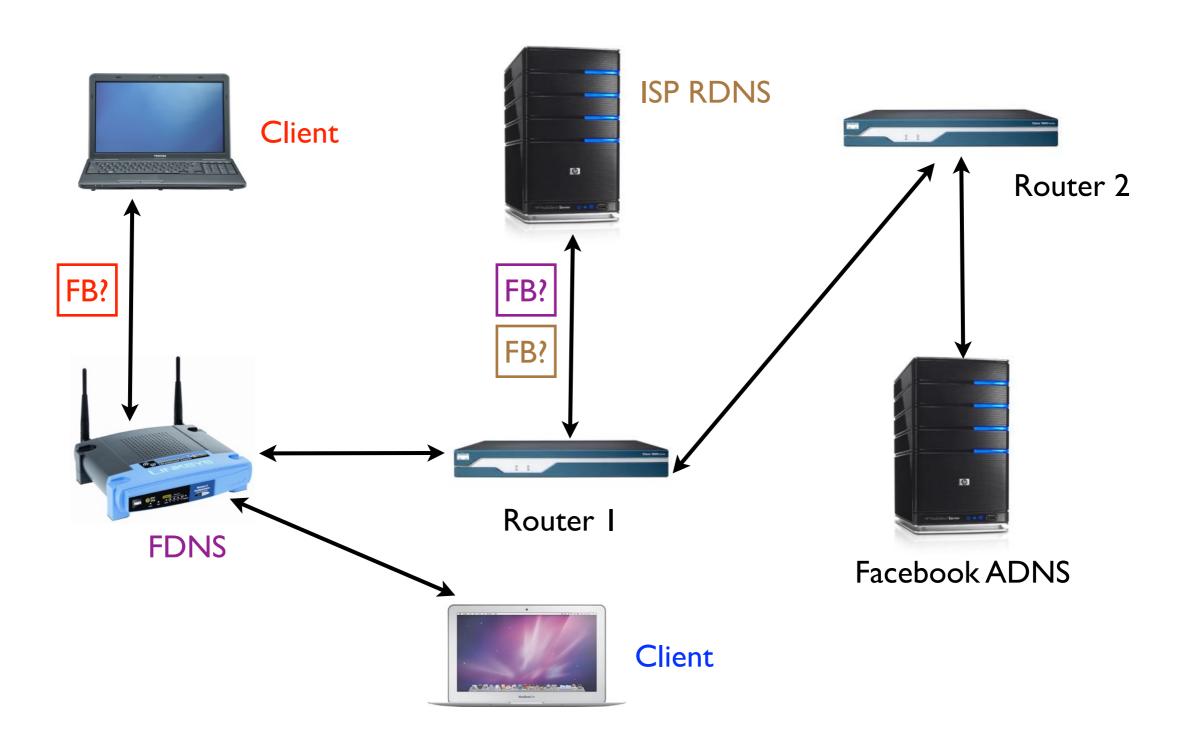


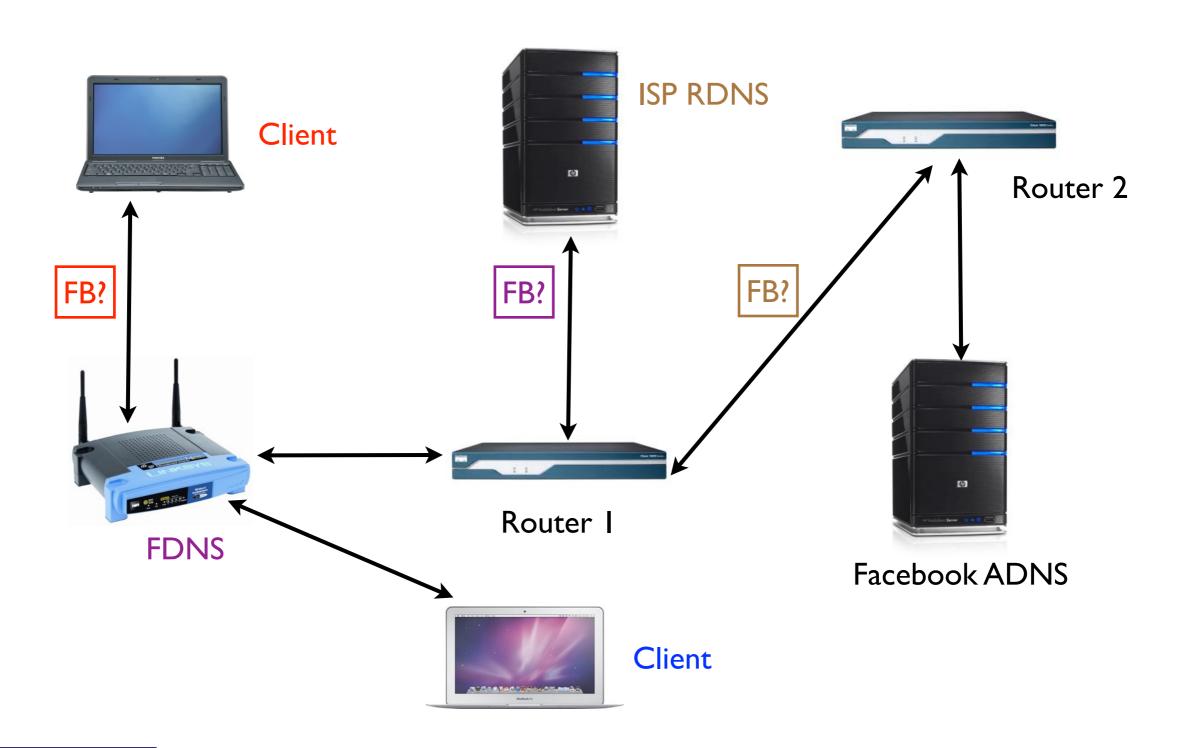


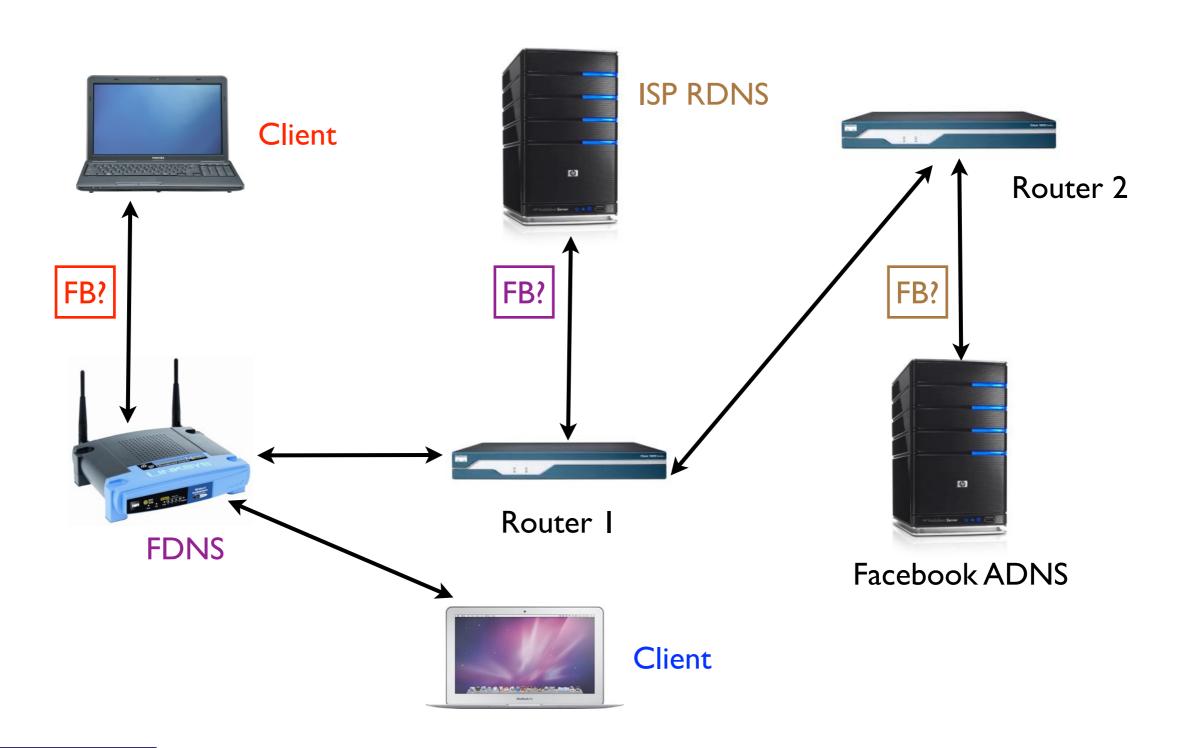


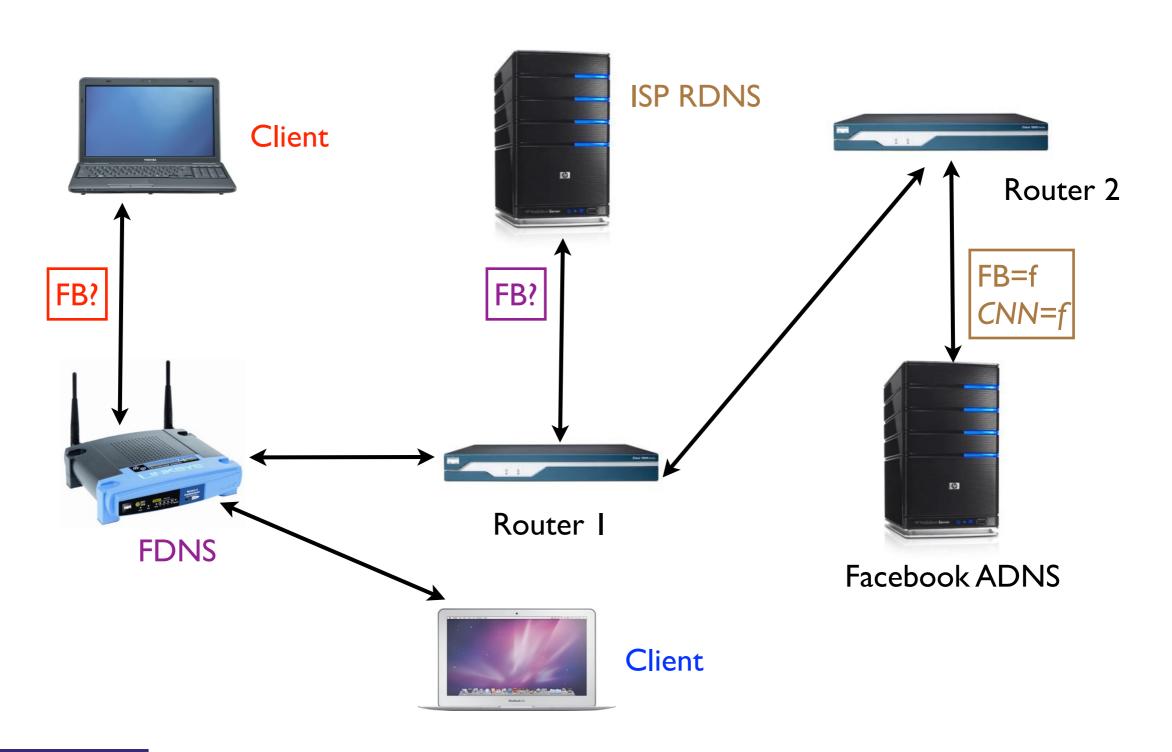


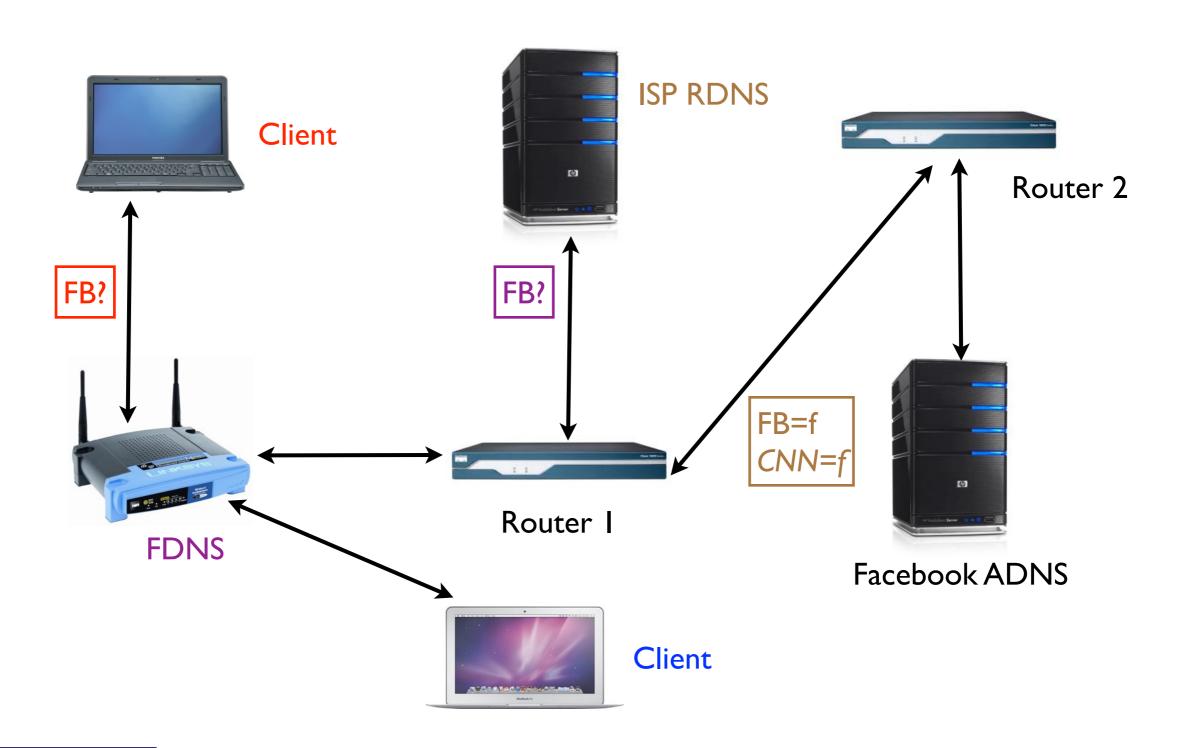


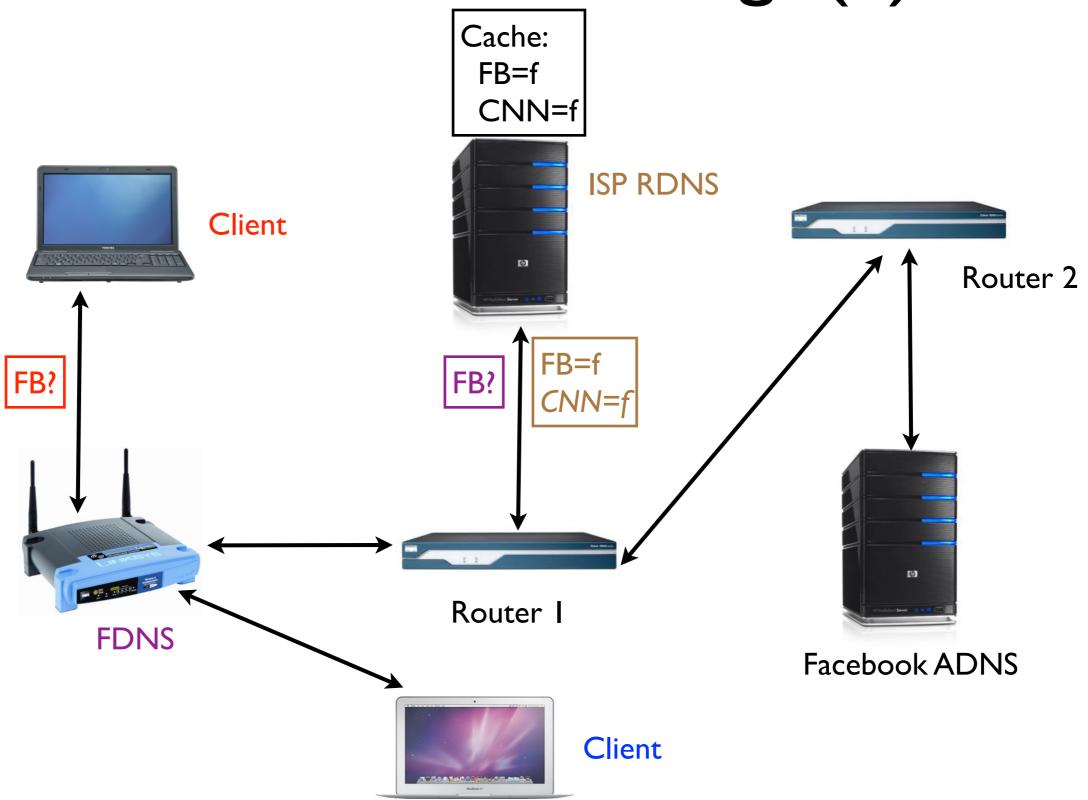


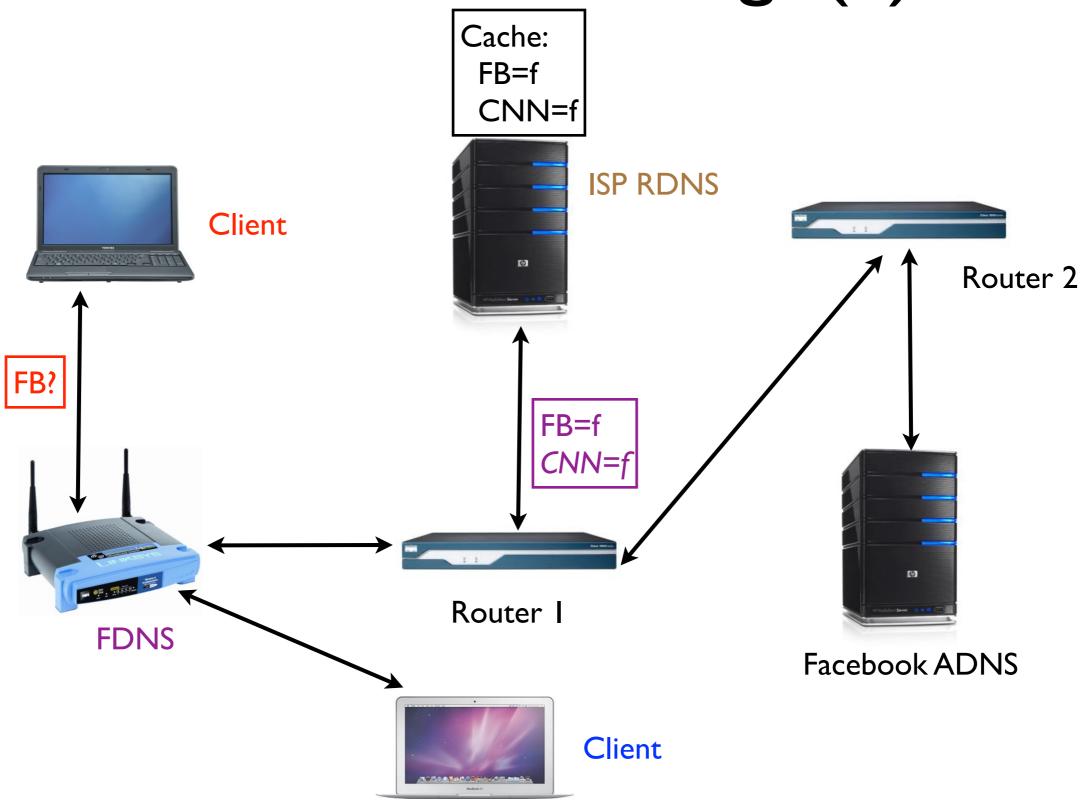


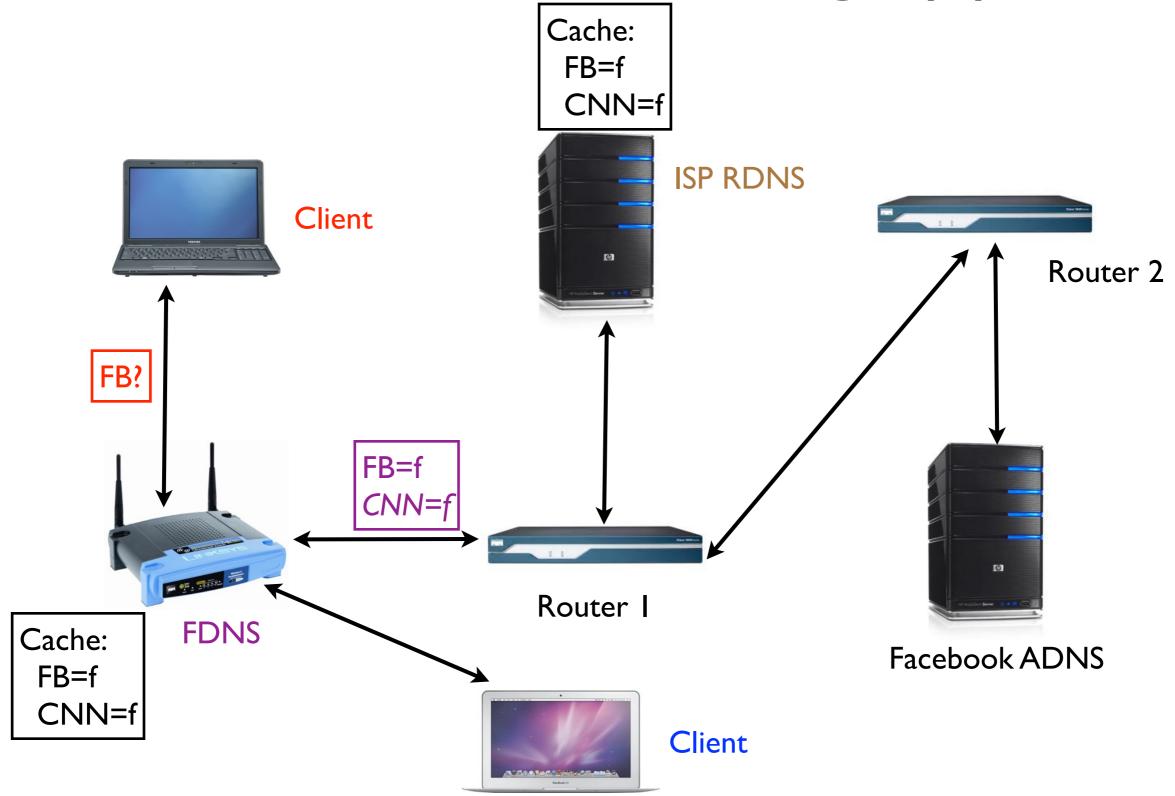


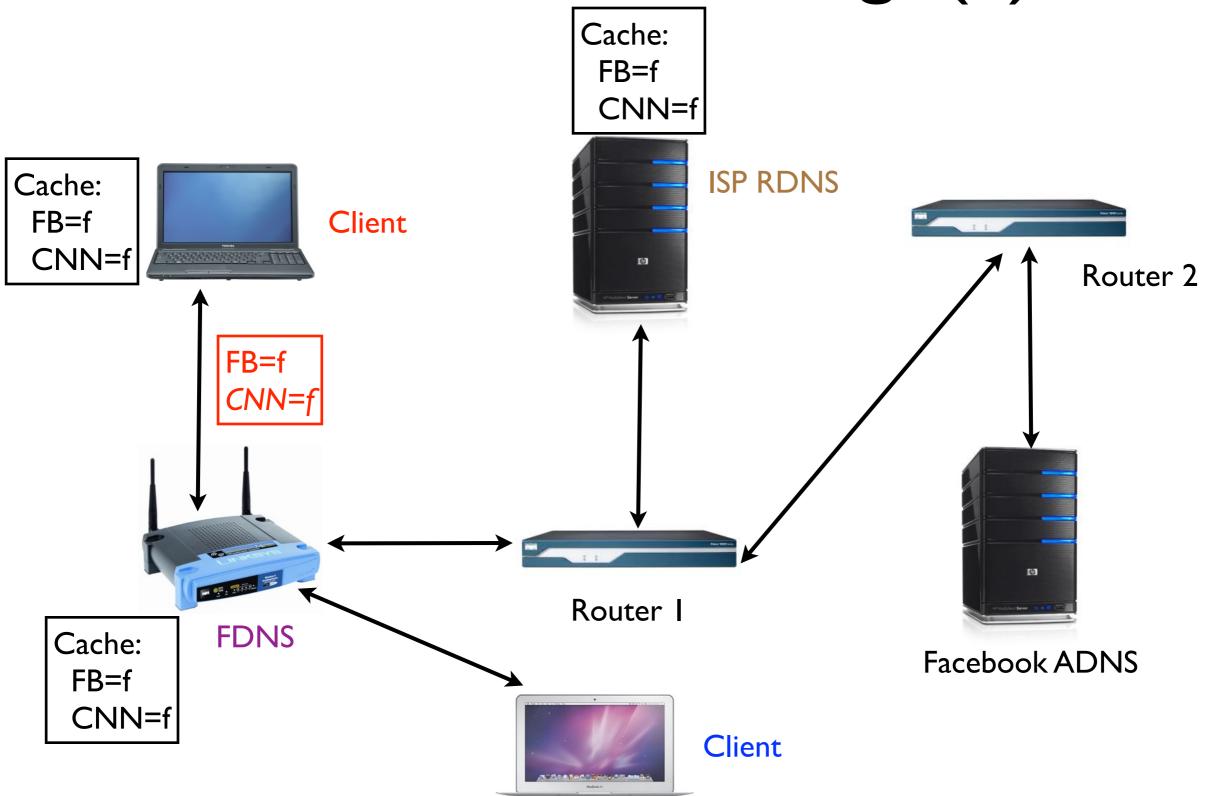


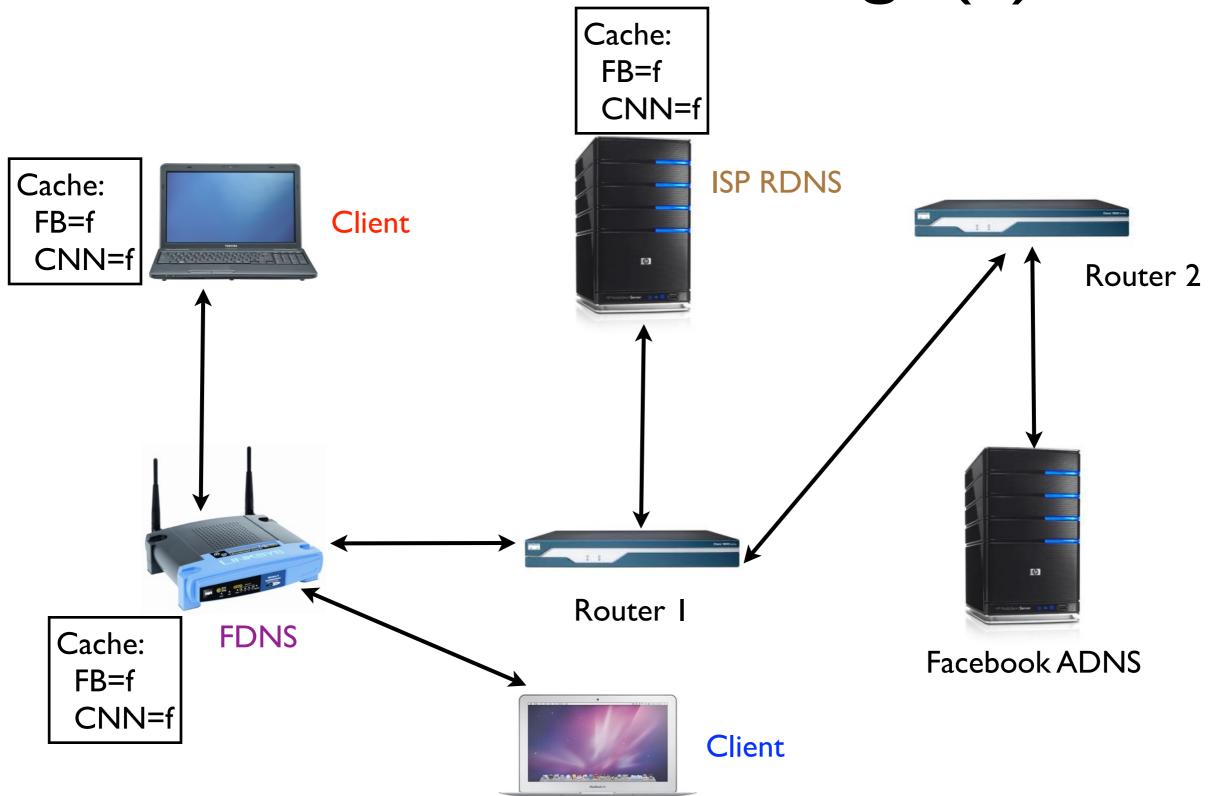


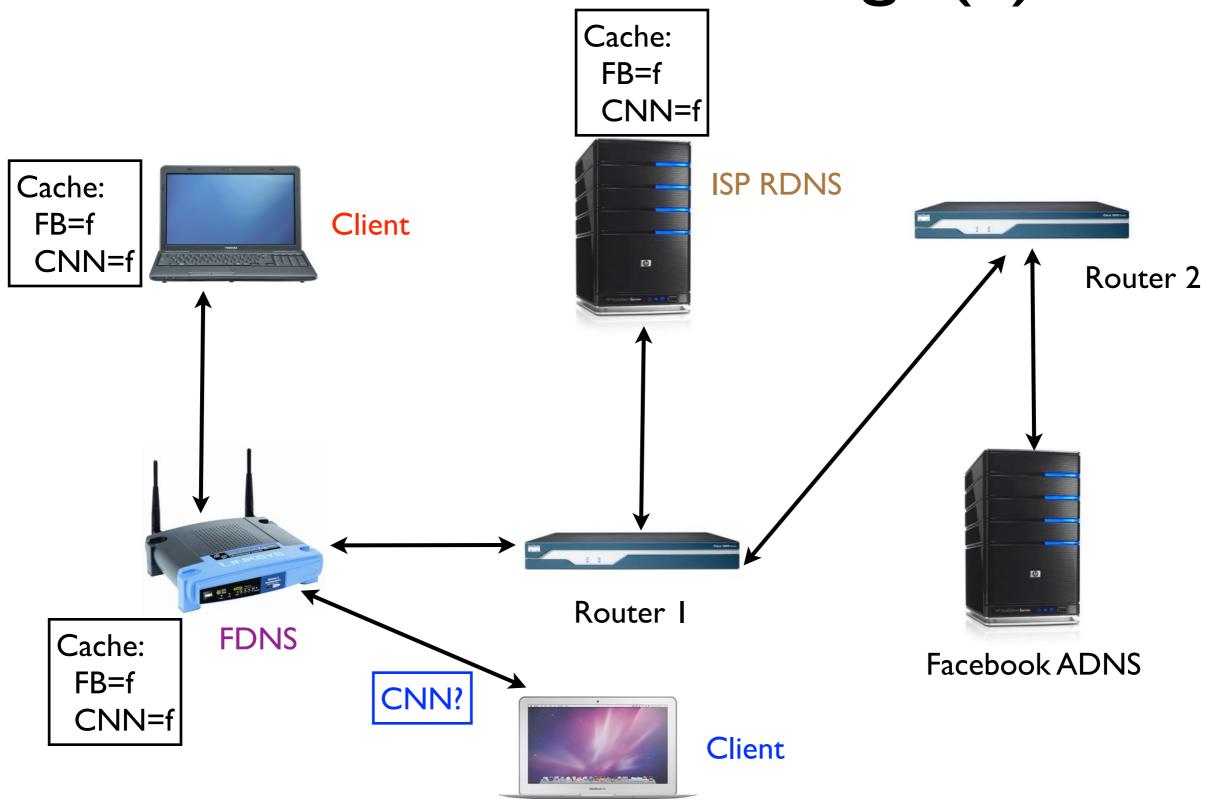


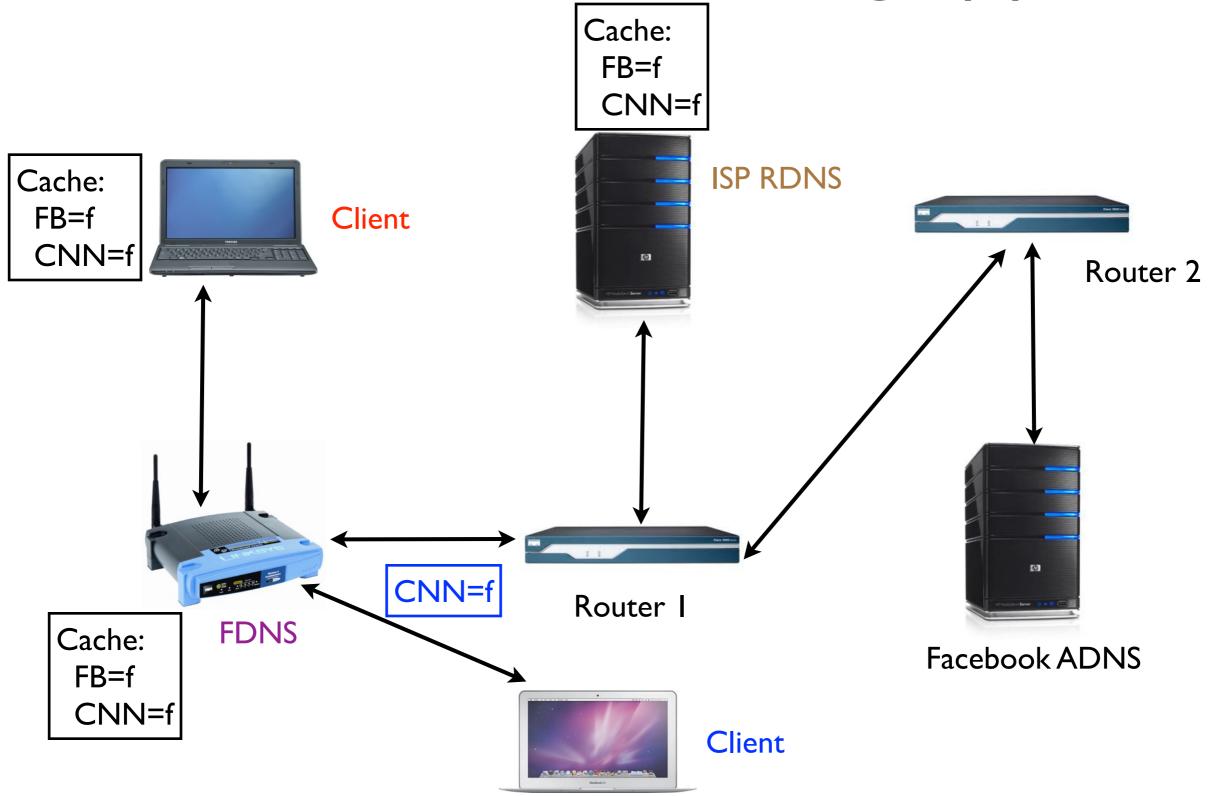












ADNS Lying

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 Out of I.09M open resolvers scanned we find 749 cases where bogus names are cached by part of the DNS infrastructure