

Mark Allman mallman@case.edu

EECS 325/425 Fall 2018

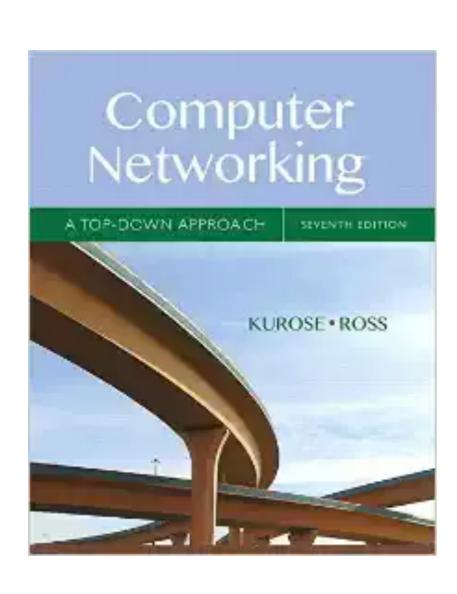
"Well, we busted outa class, had to get away from those fools, We learned more from a three minute record than we ever learned in school."

These slides are more-or-less directly from the slide set developed by Jim Kurose and Keith Ross for their book "Computer Networking: A Top Down Approach, 5th edition".

The slides have been lightly adapted for Mark Allman's EECS 325/425 Computer Networks class at Case Western Reserve University.

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# Reading Along ...

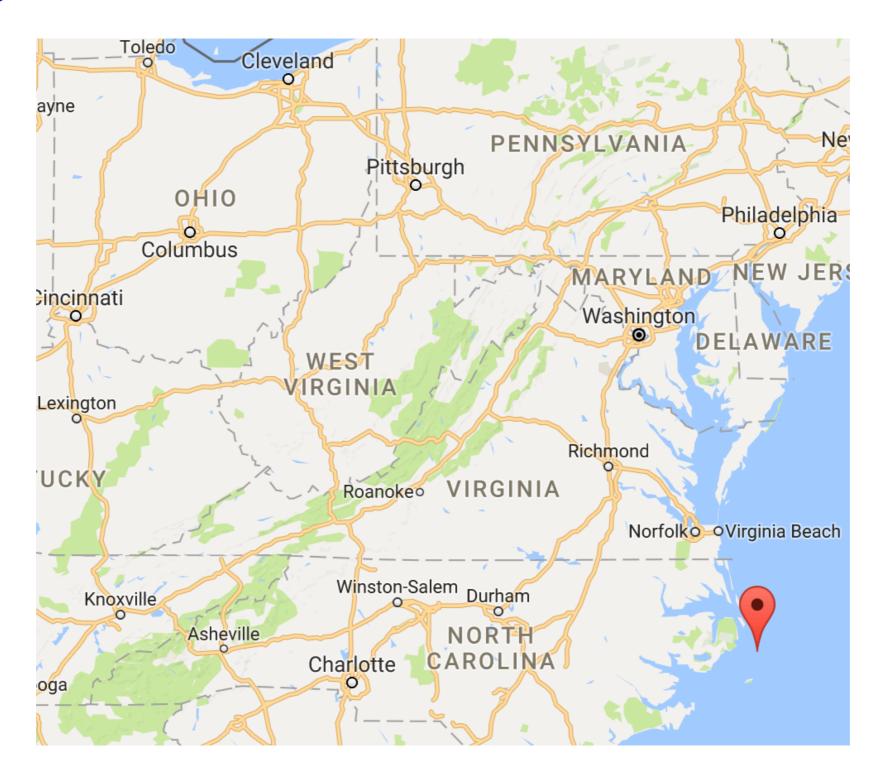


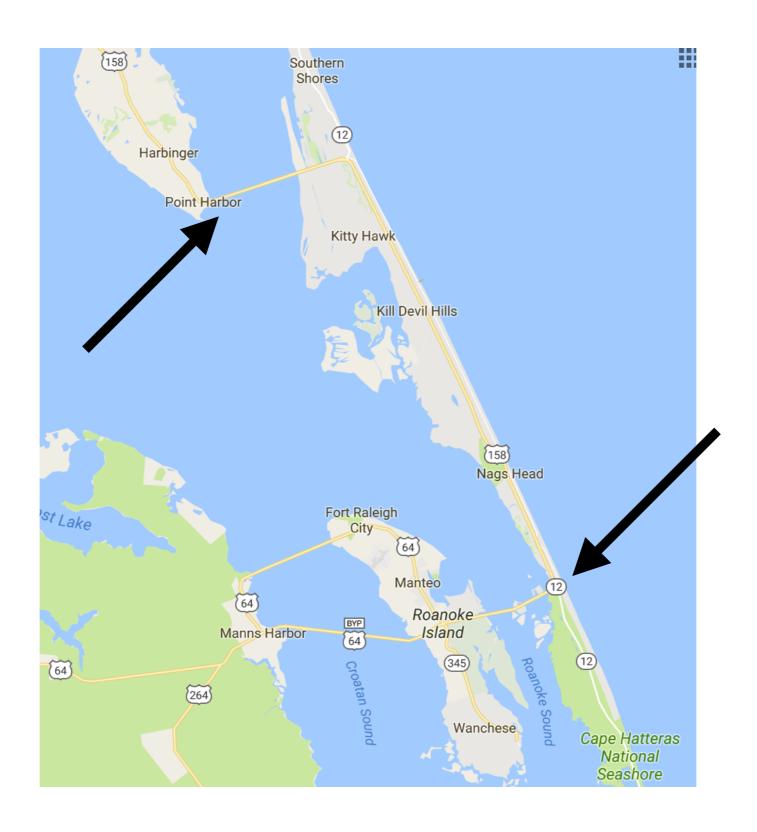
 3.6: Principles of Congestion Control

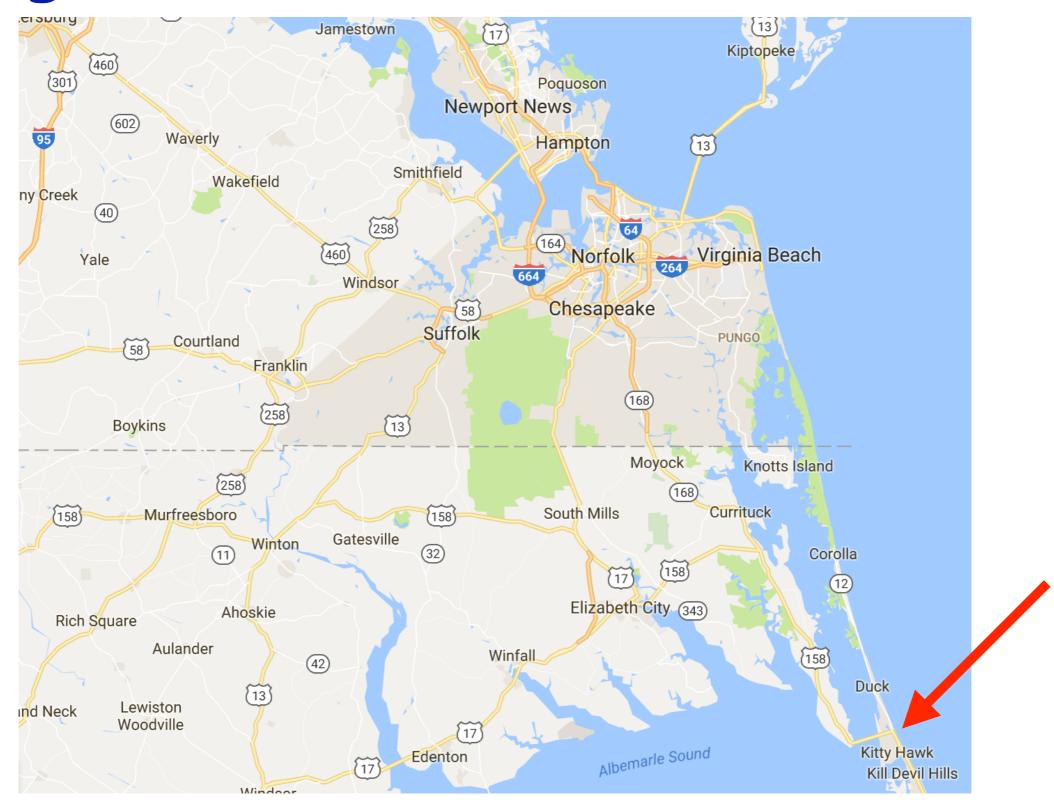
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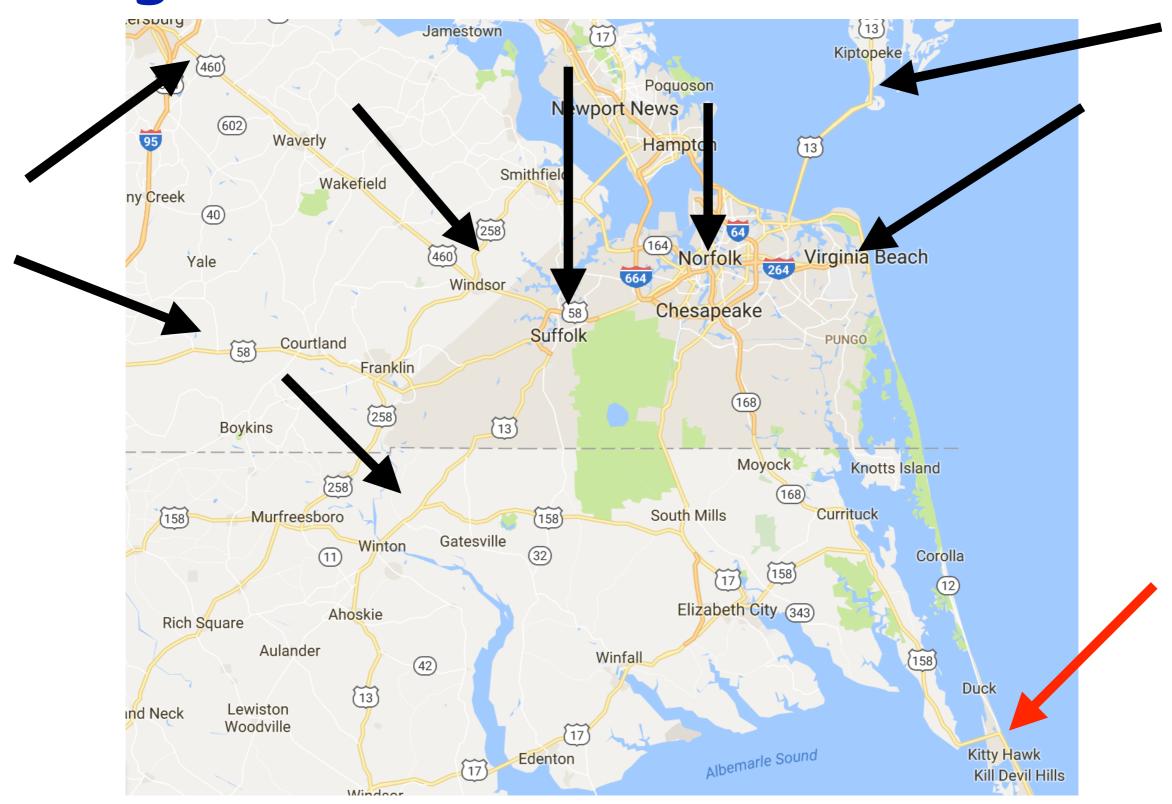
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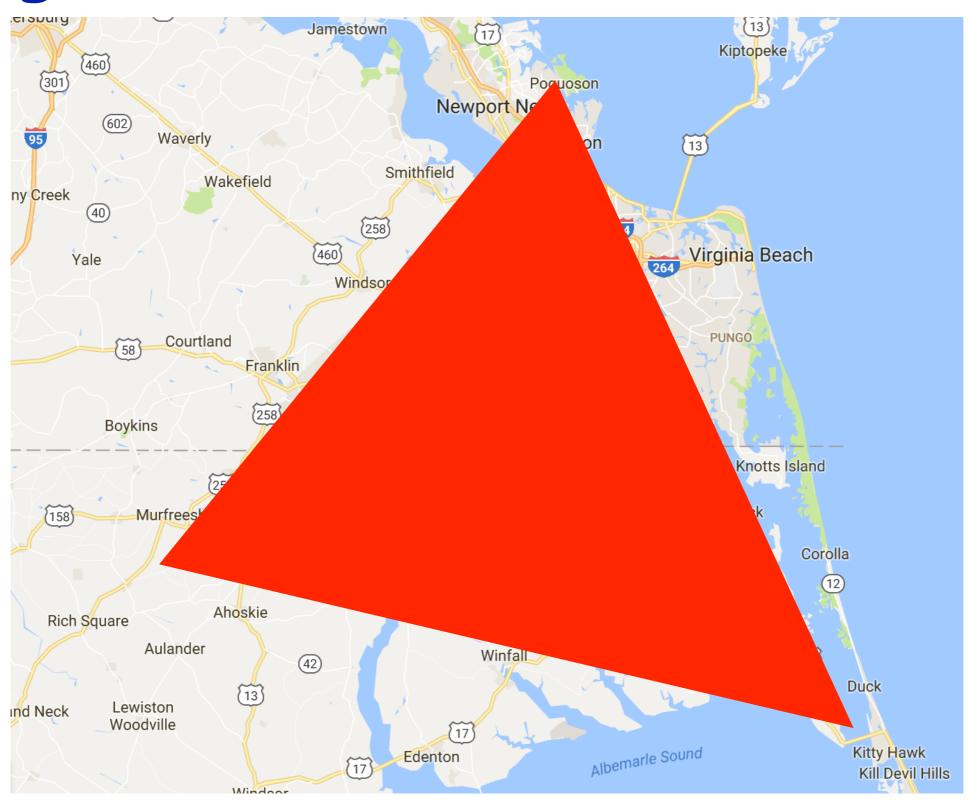
\*Why does it need "controlled"?



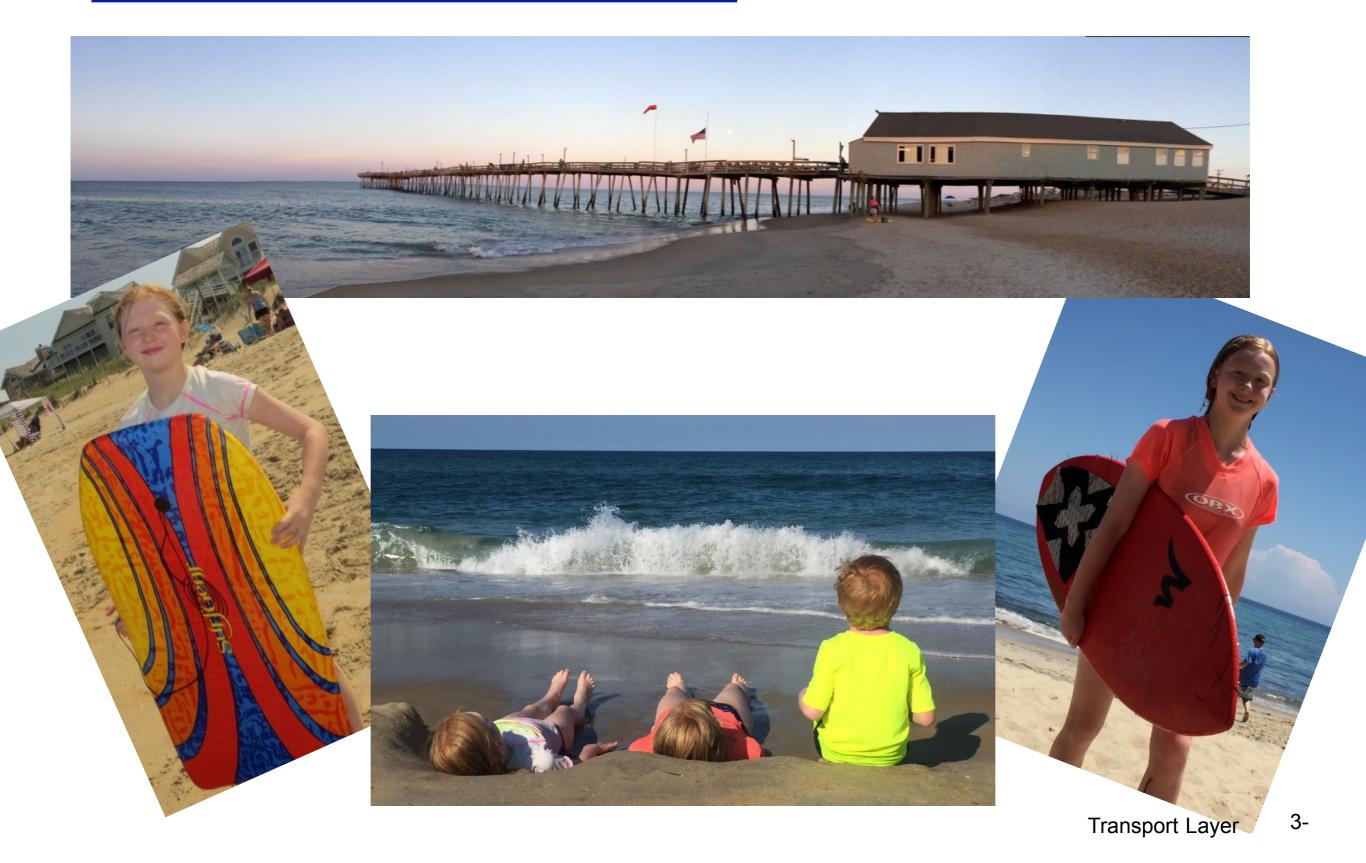






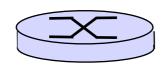


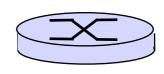
# Worth the Effort!



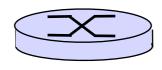






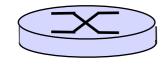


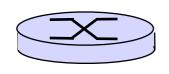










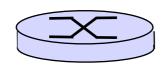


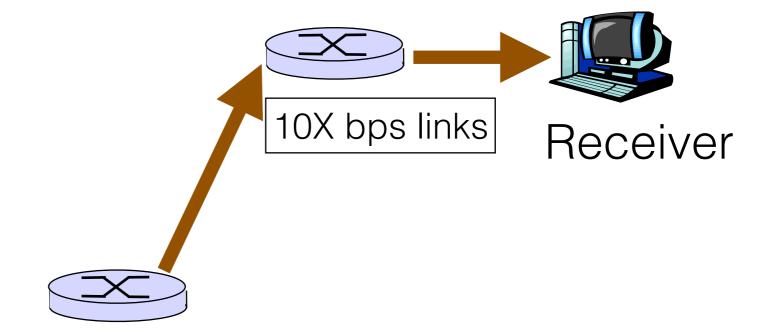




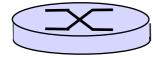


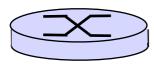








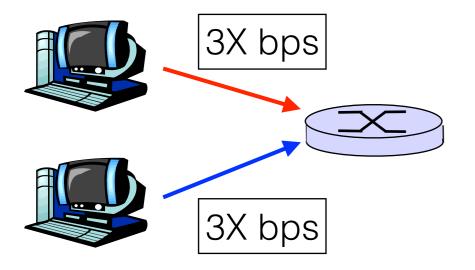


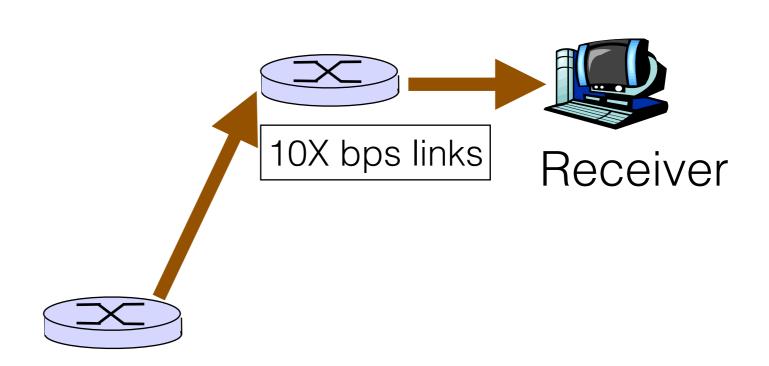




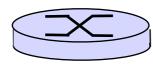


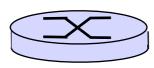








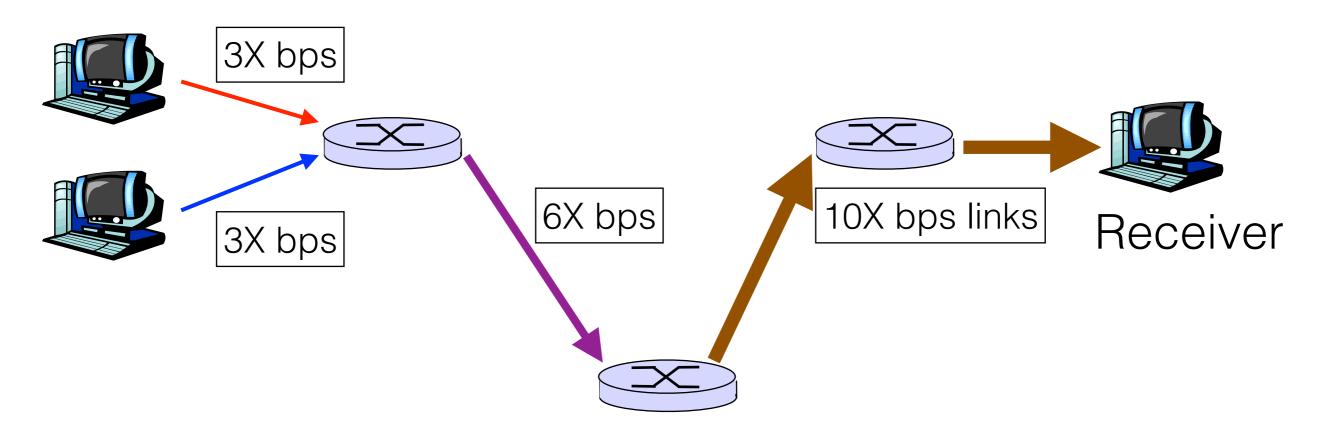




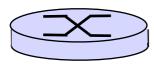








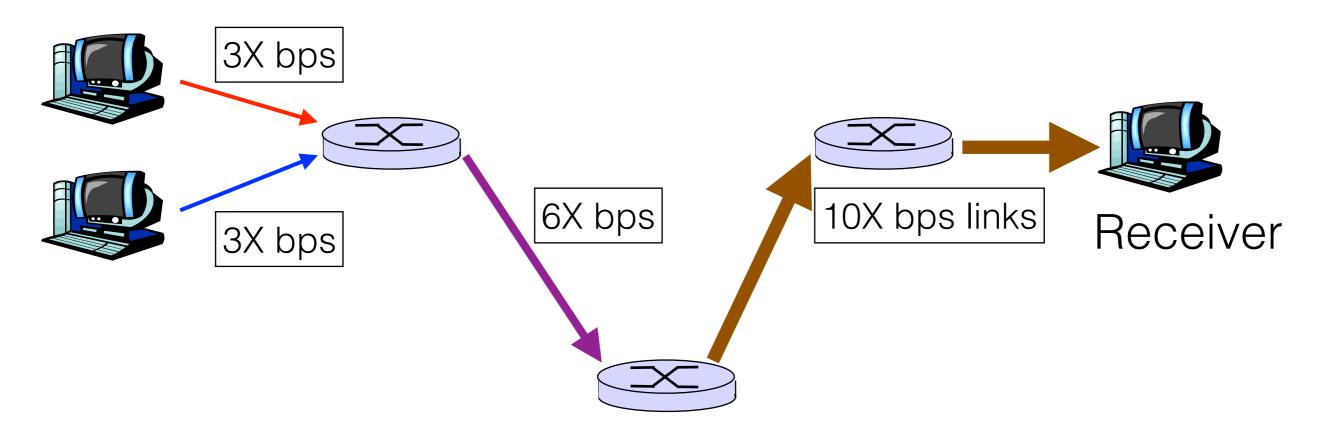


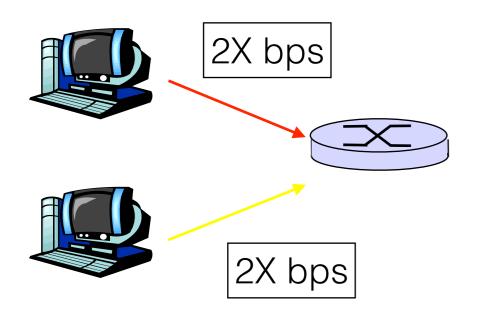


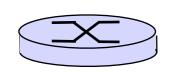






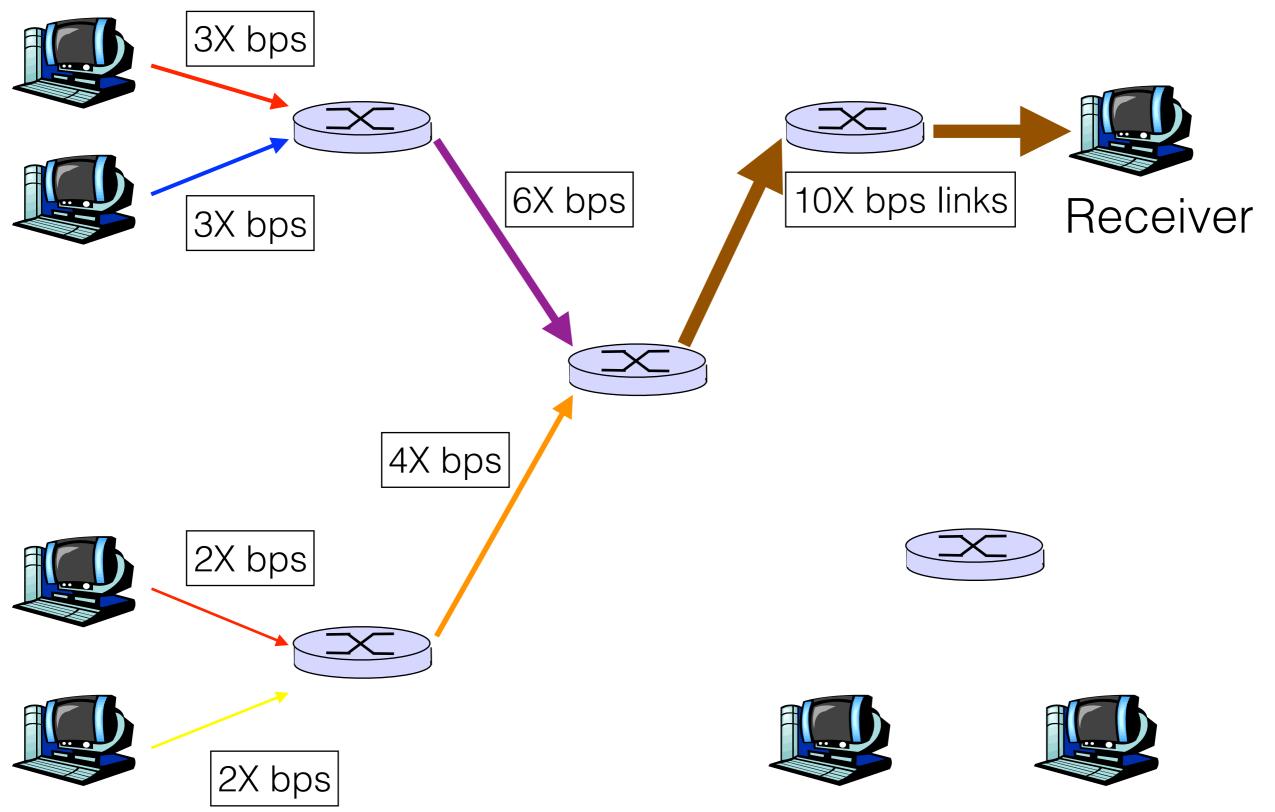


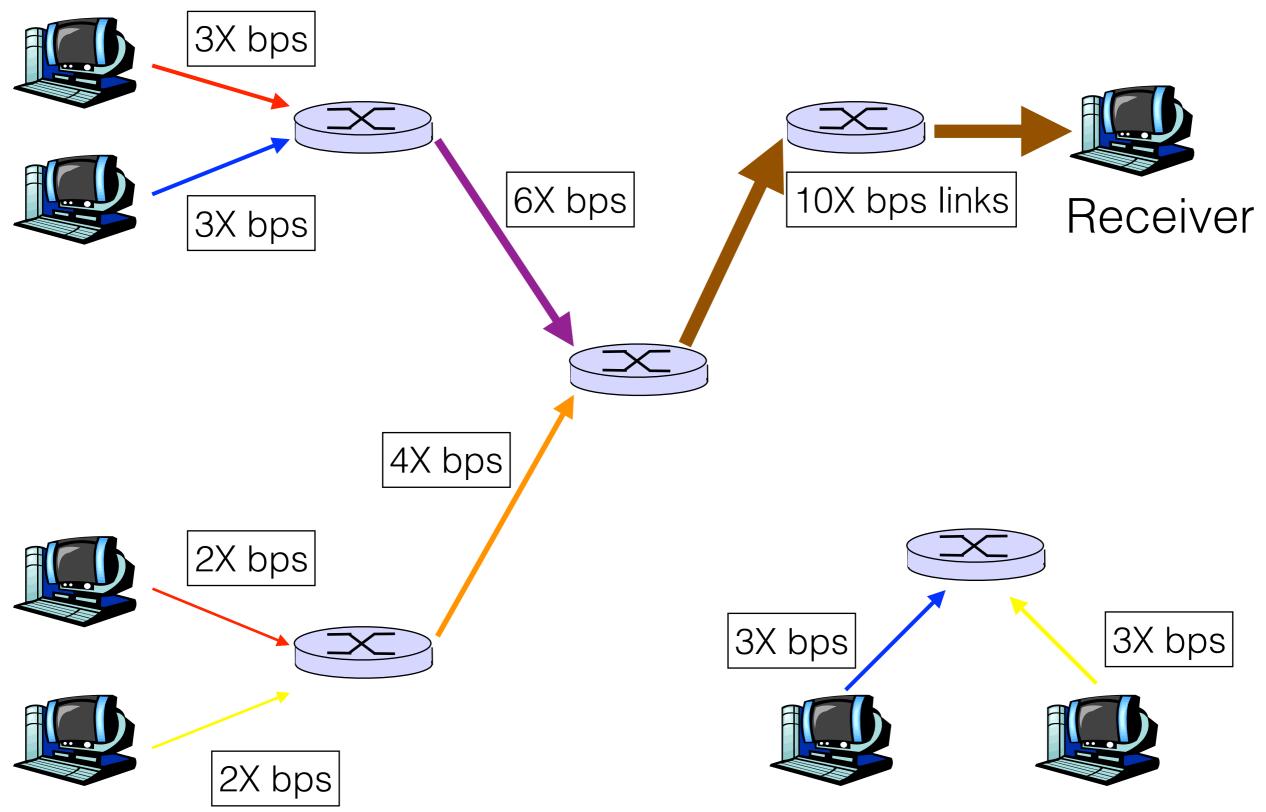


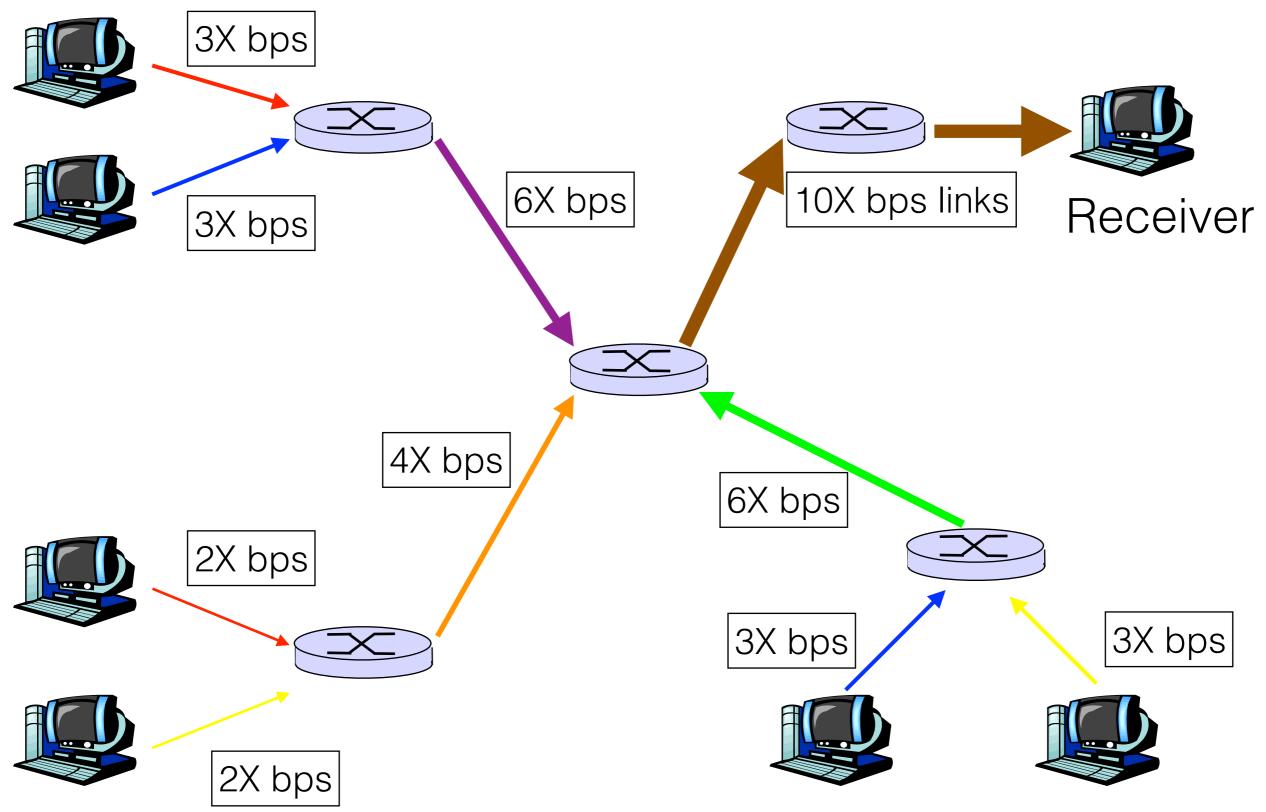


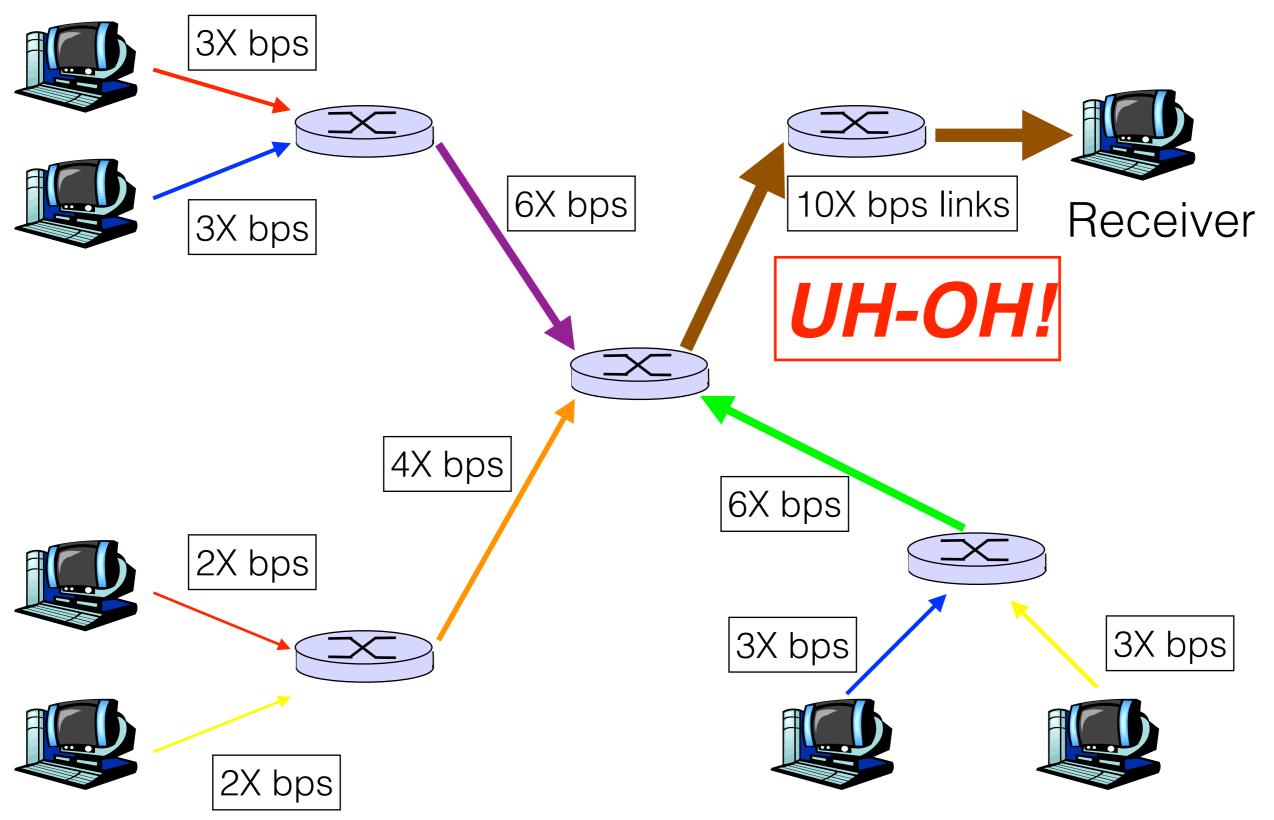












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- \*manifestations:
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- \*a top-10 problem!

- \*more work for a given sending rate
  - i.e., more retransmissions

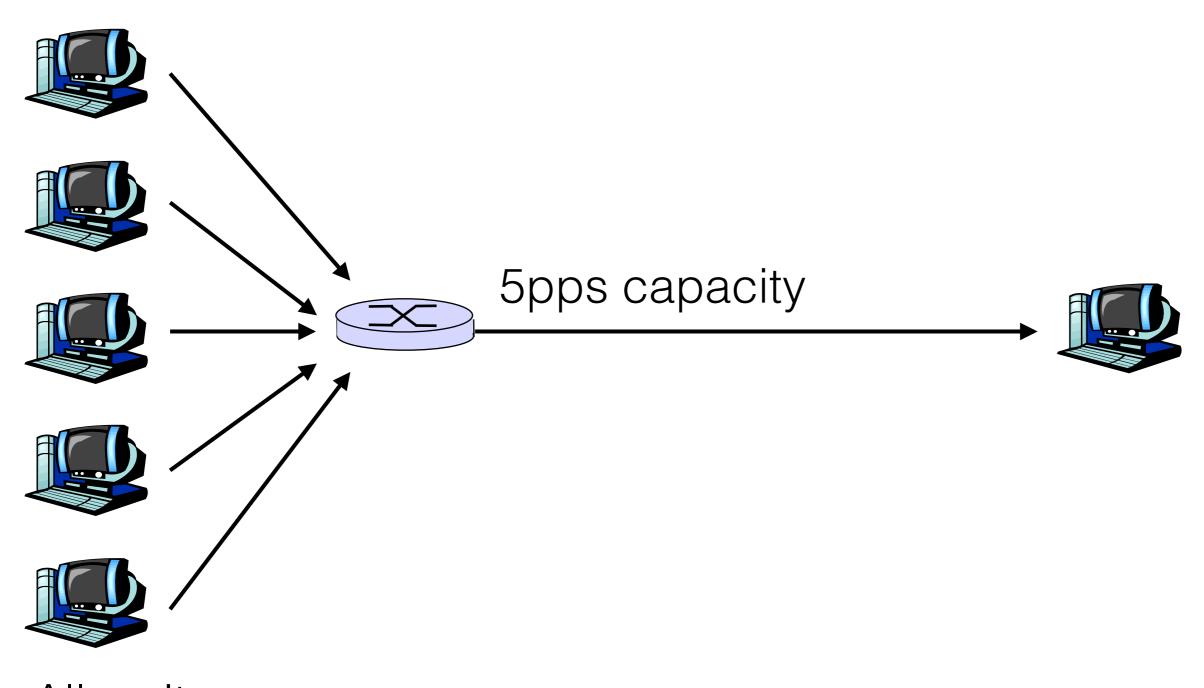
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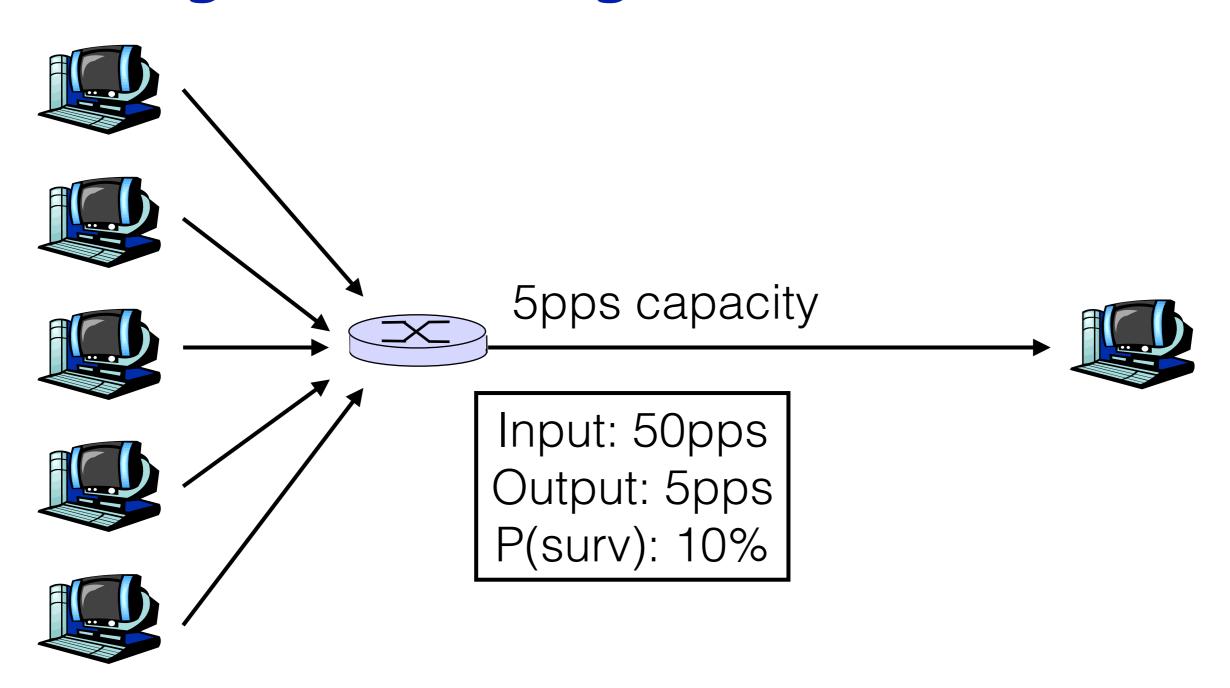
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- \*but, are these annoyance or fundamental issues?!
  - left unchecked congestion can lead to "collapse" whereby a network expends scarce resources on packets that will ultimately be dropped
    - · i.e., accomplish little—if any—meaningful work

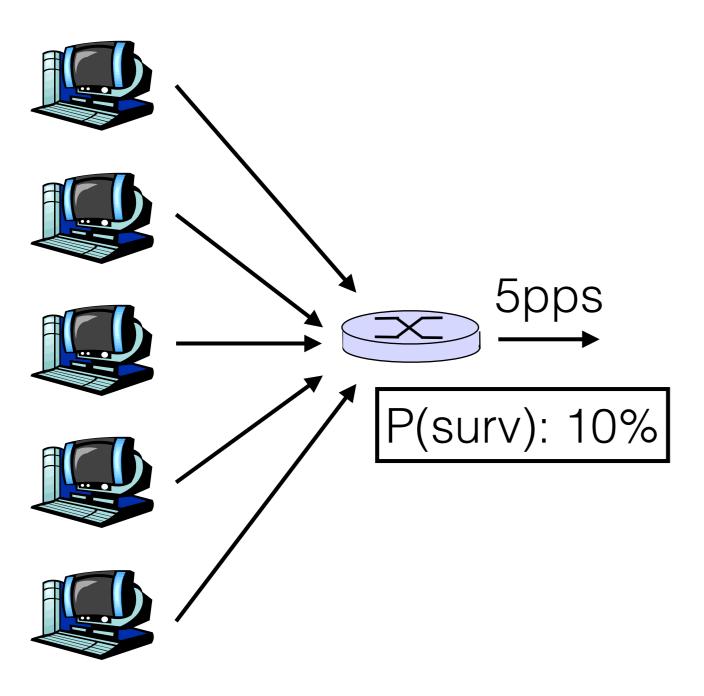
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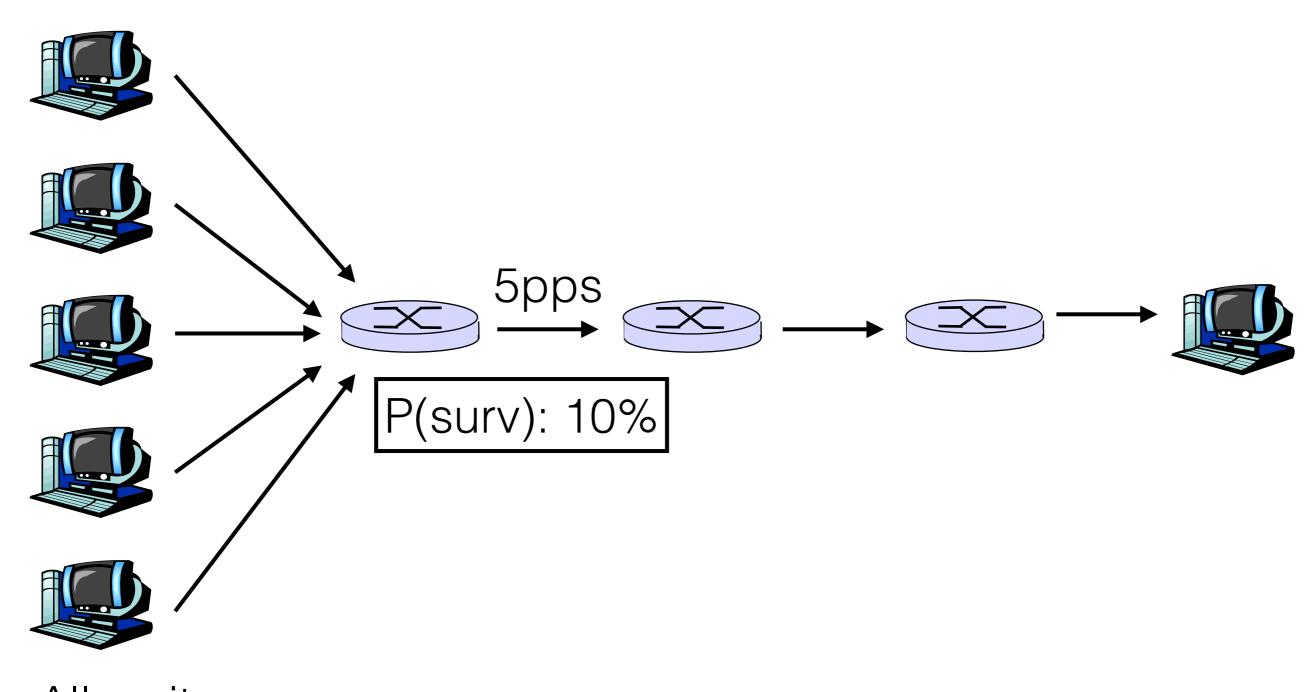


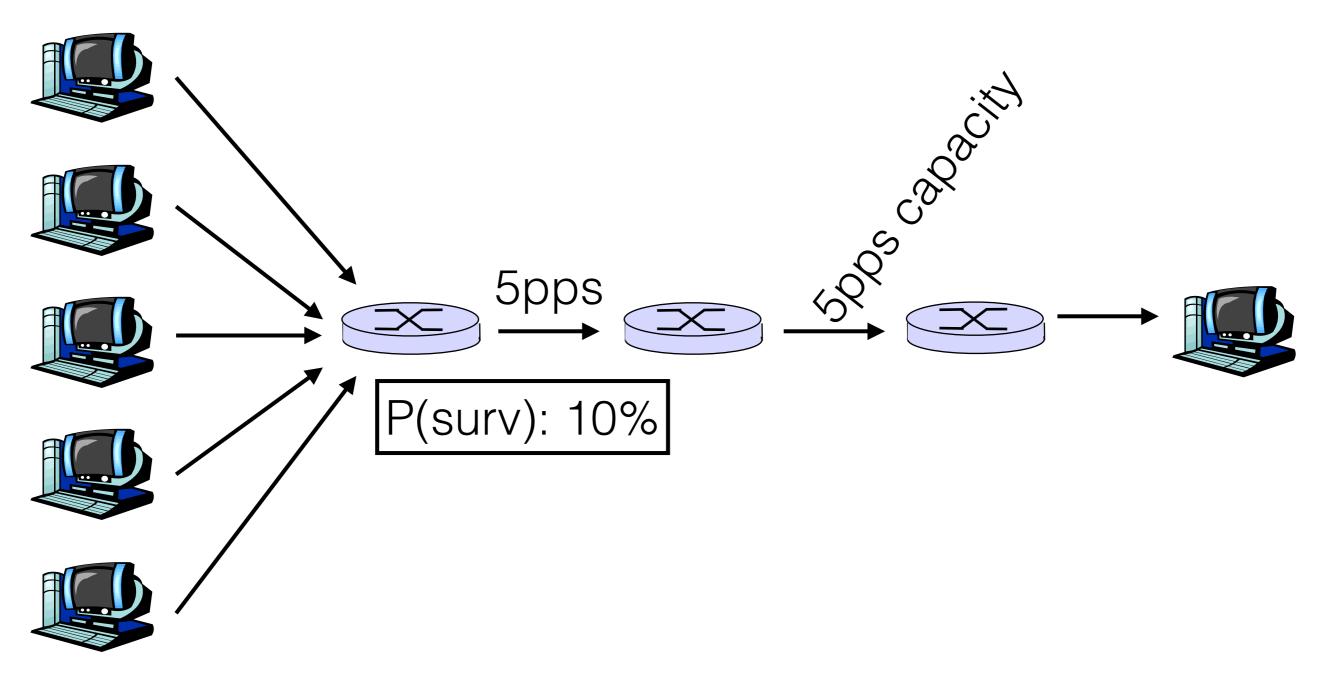
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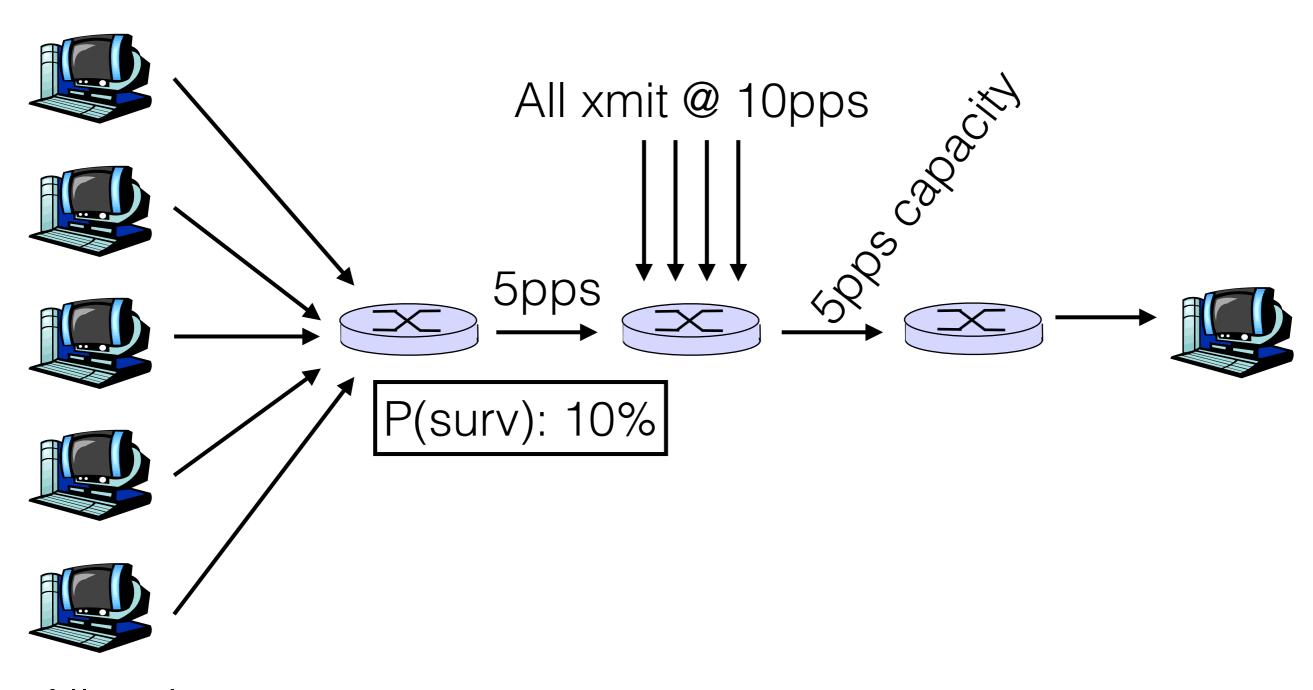
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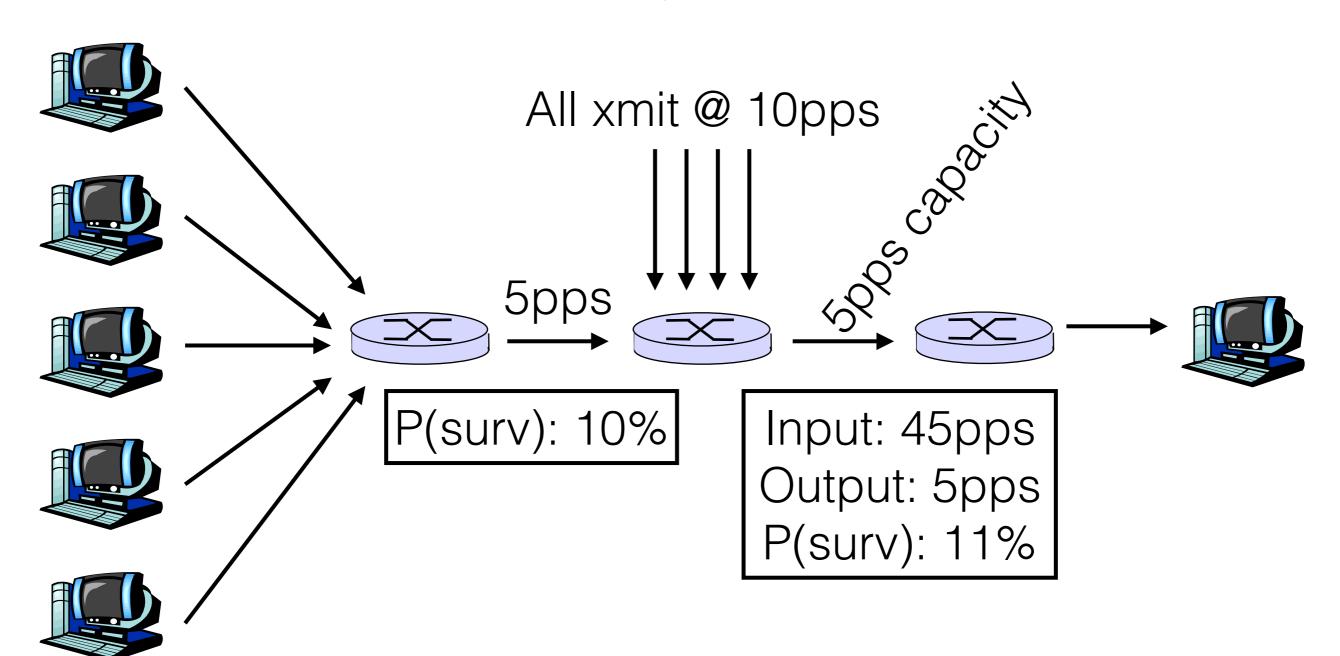


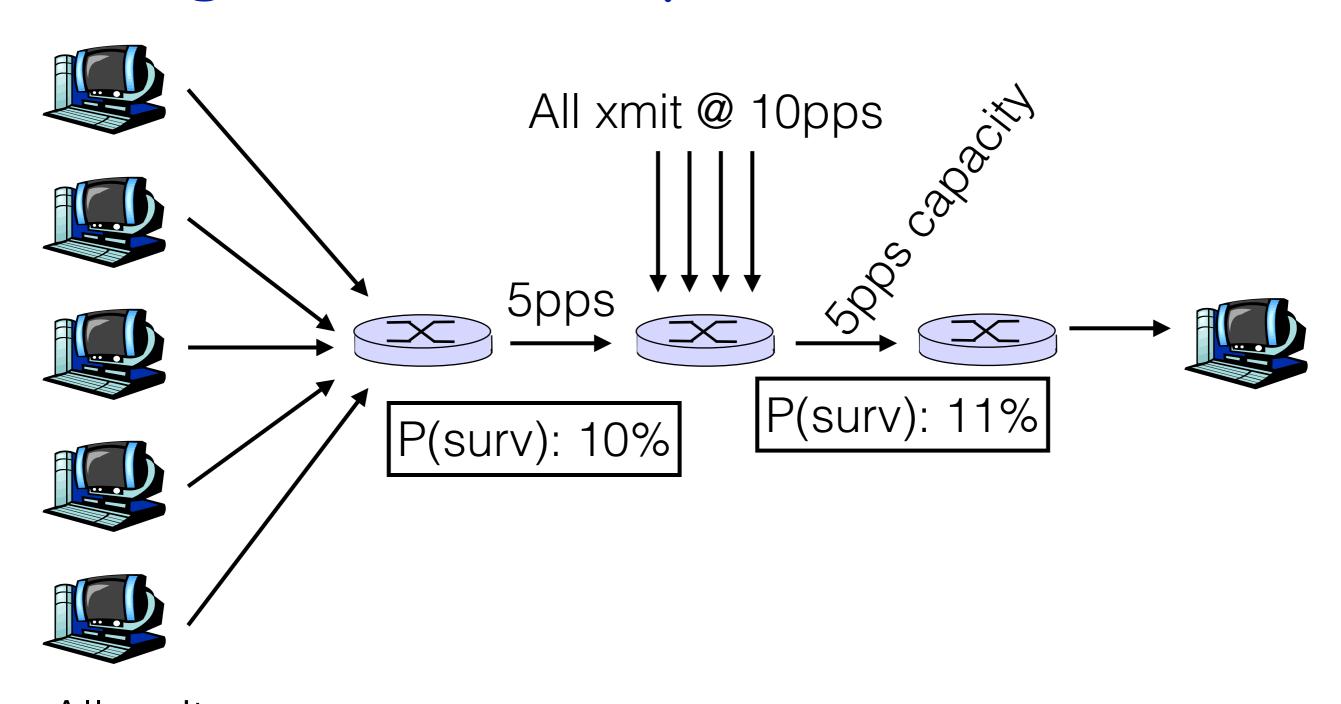


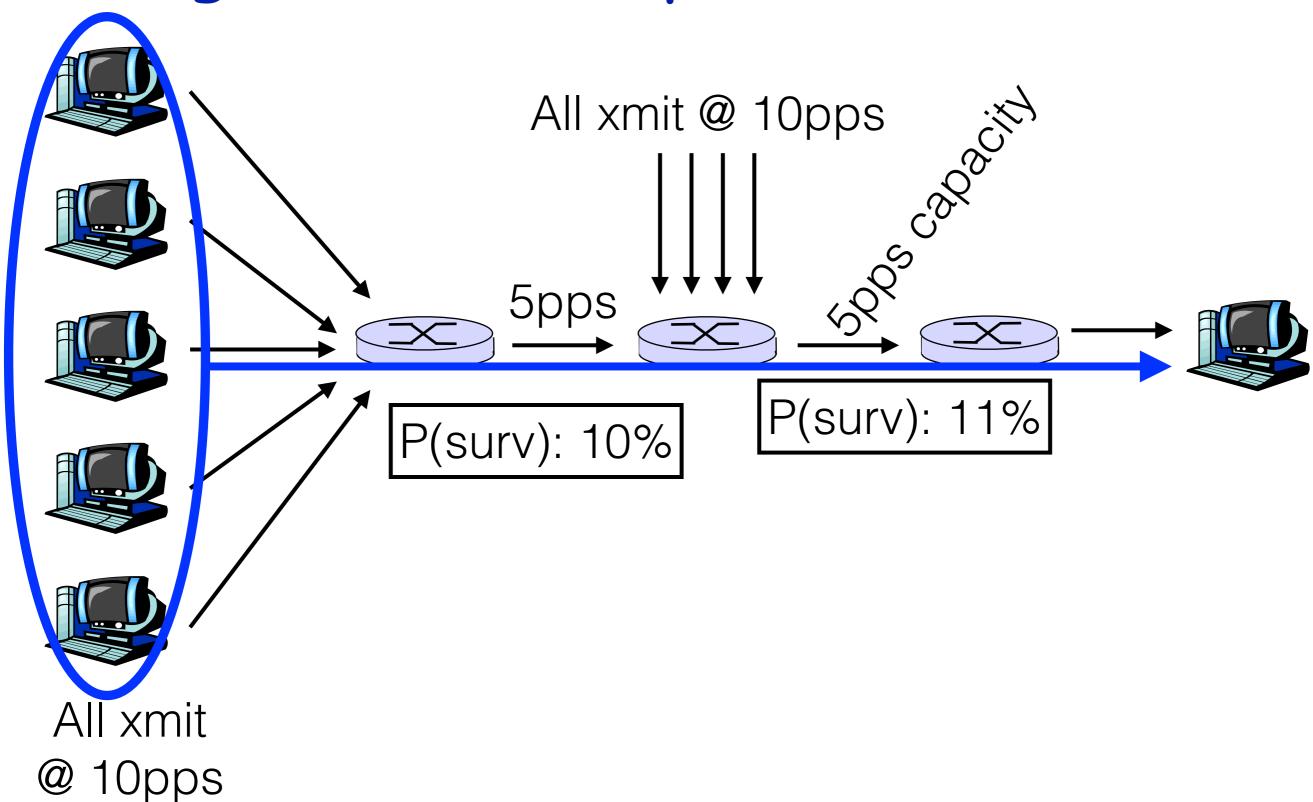
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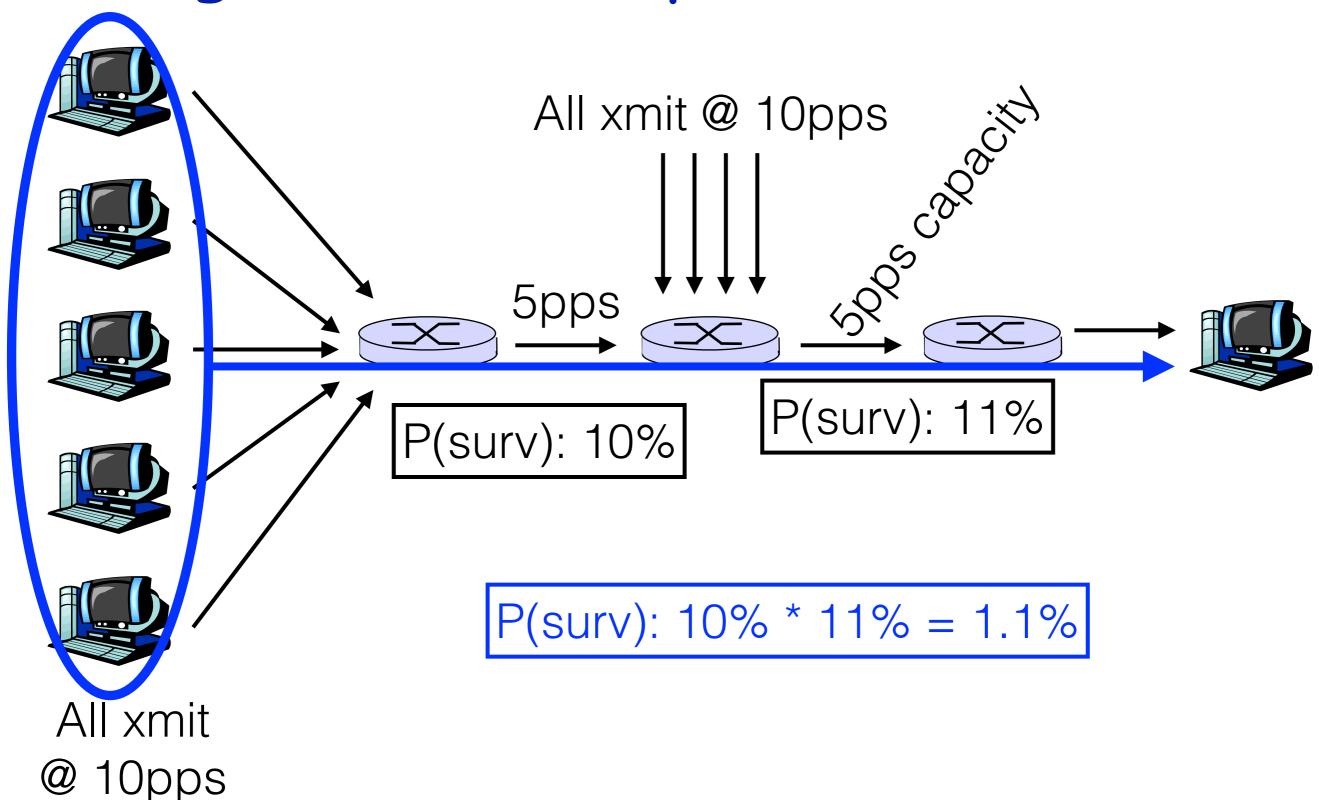
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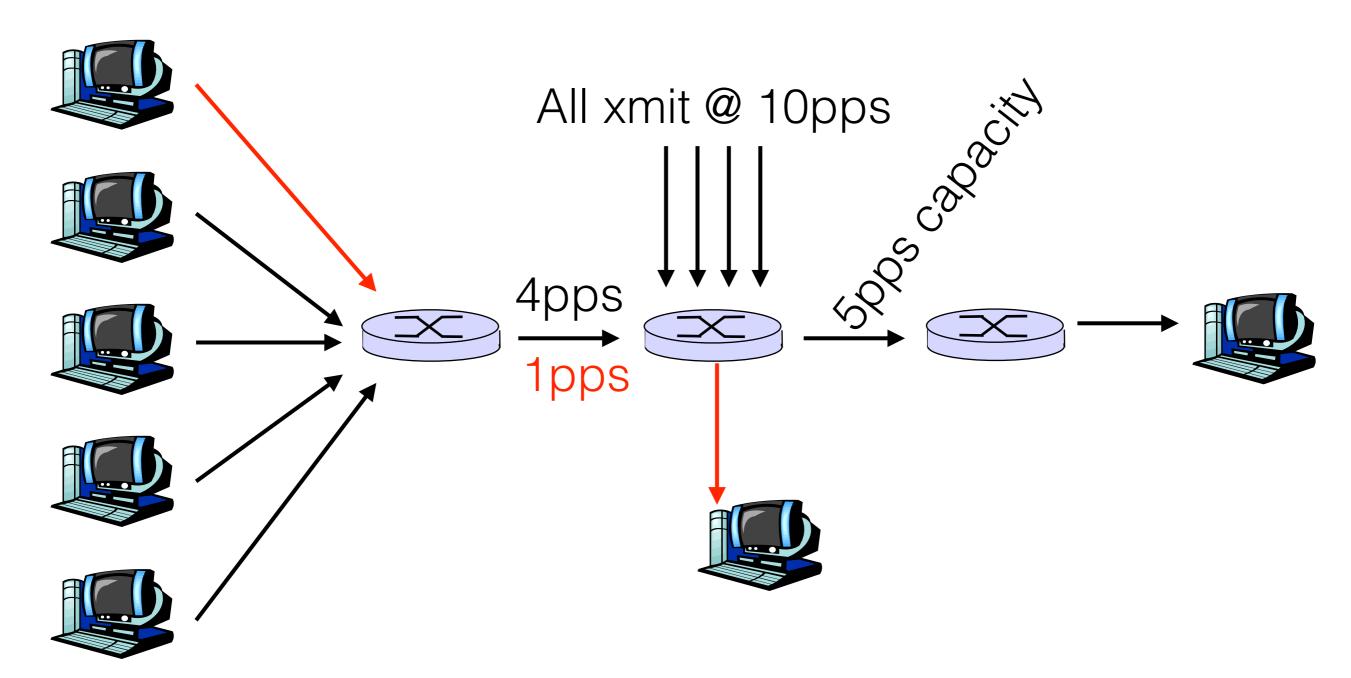






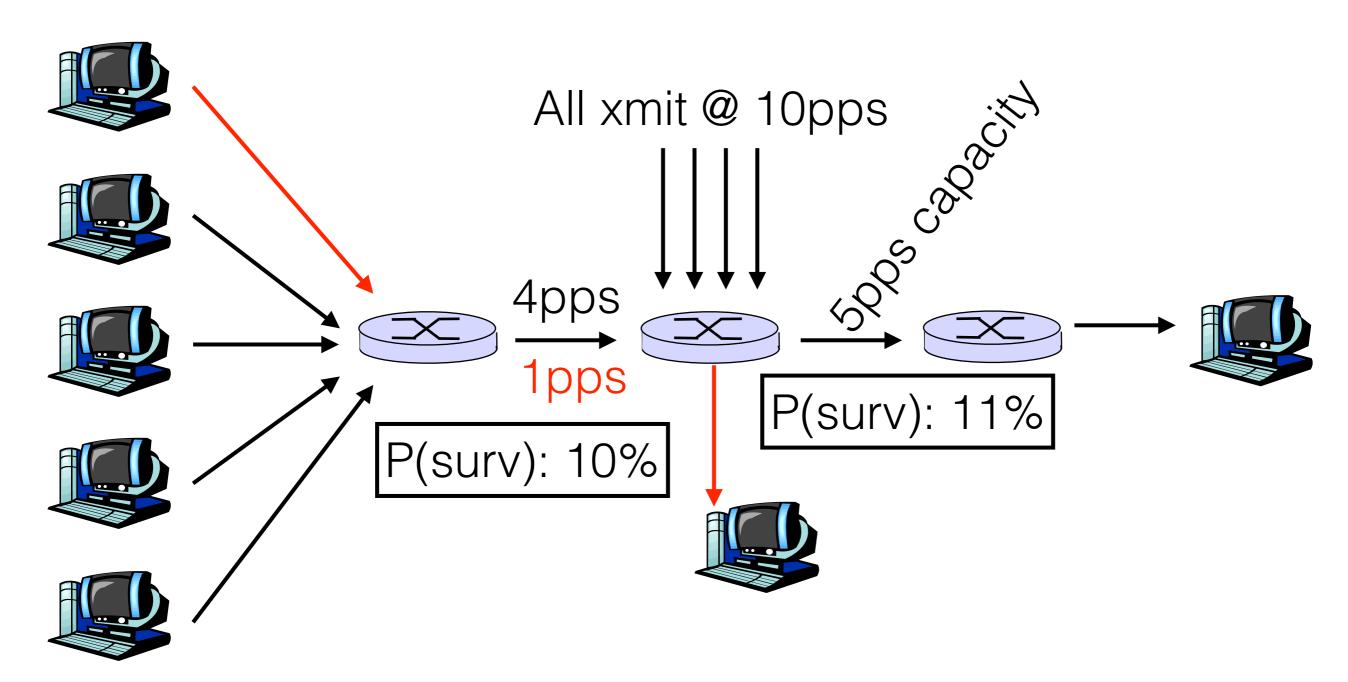


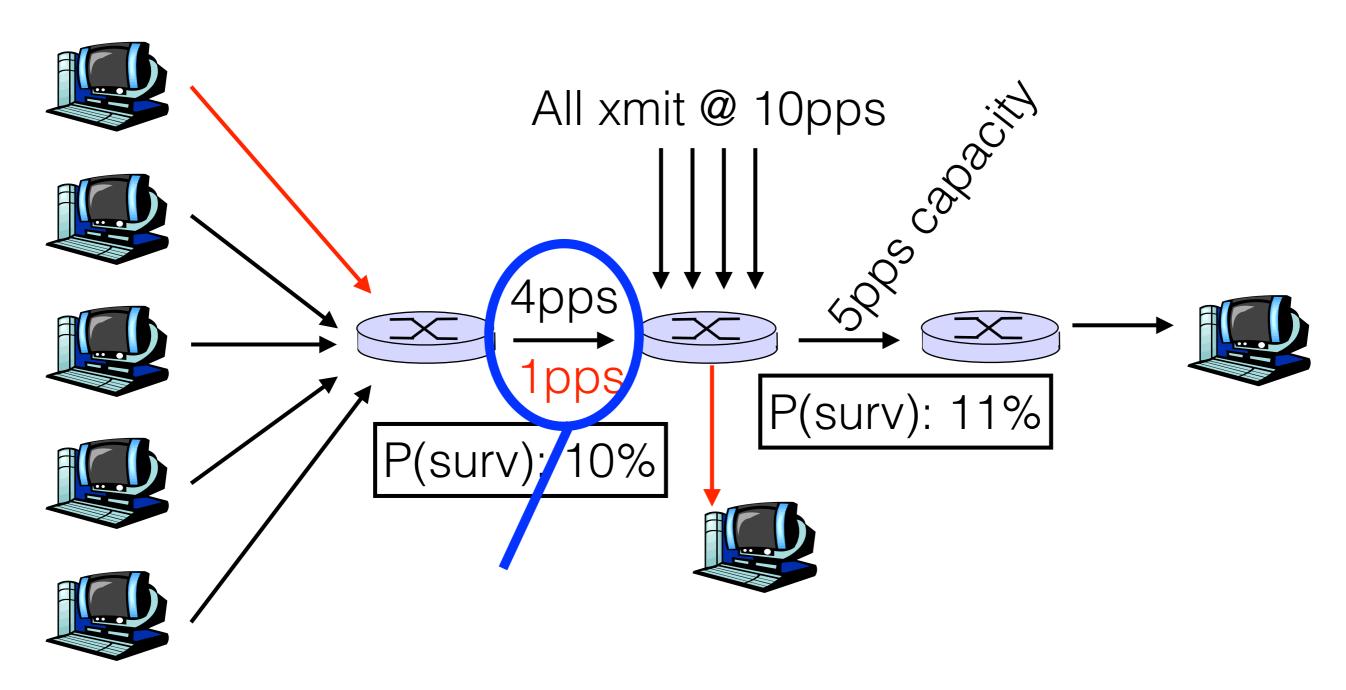


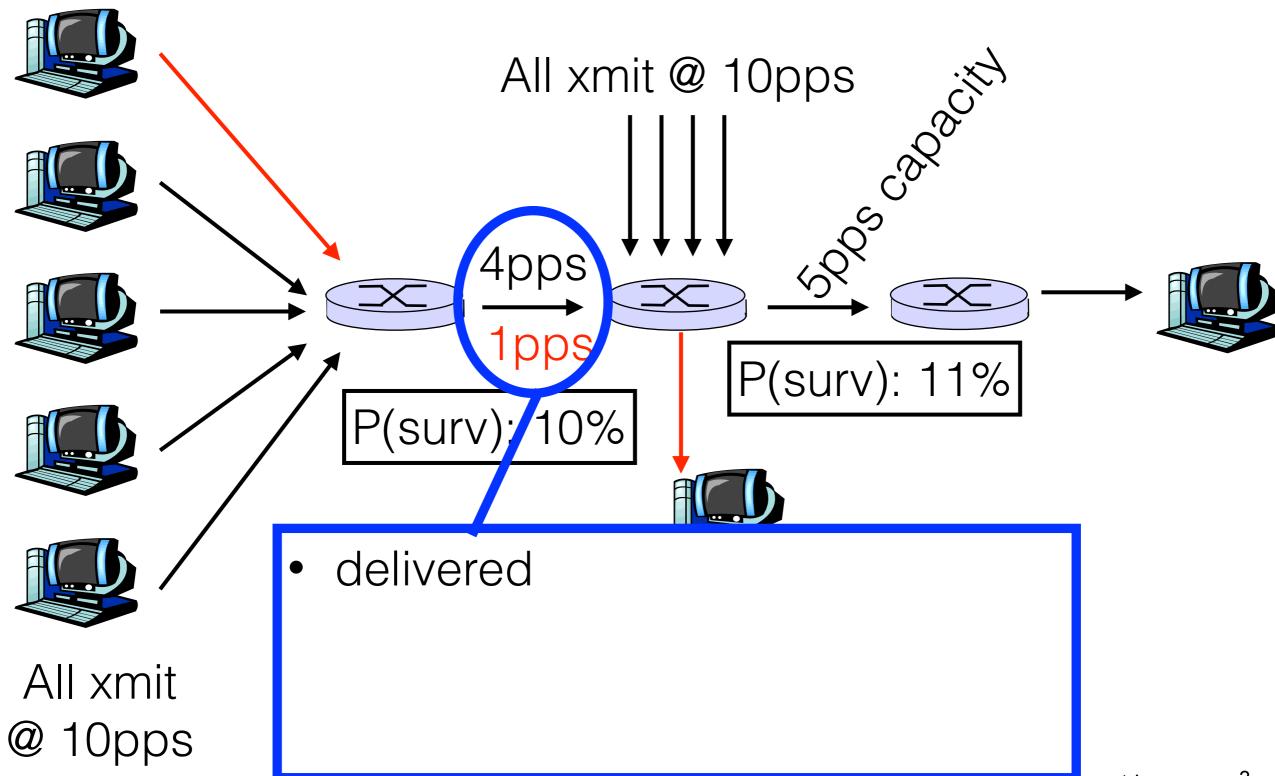


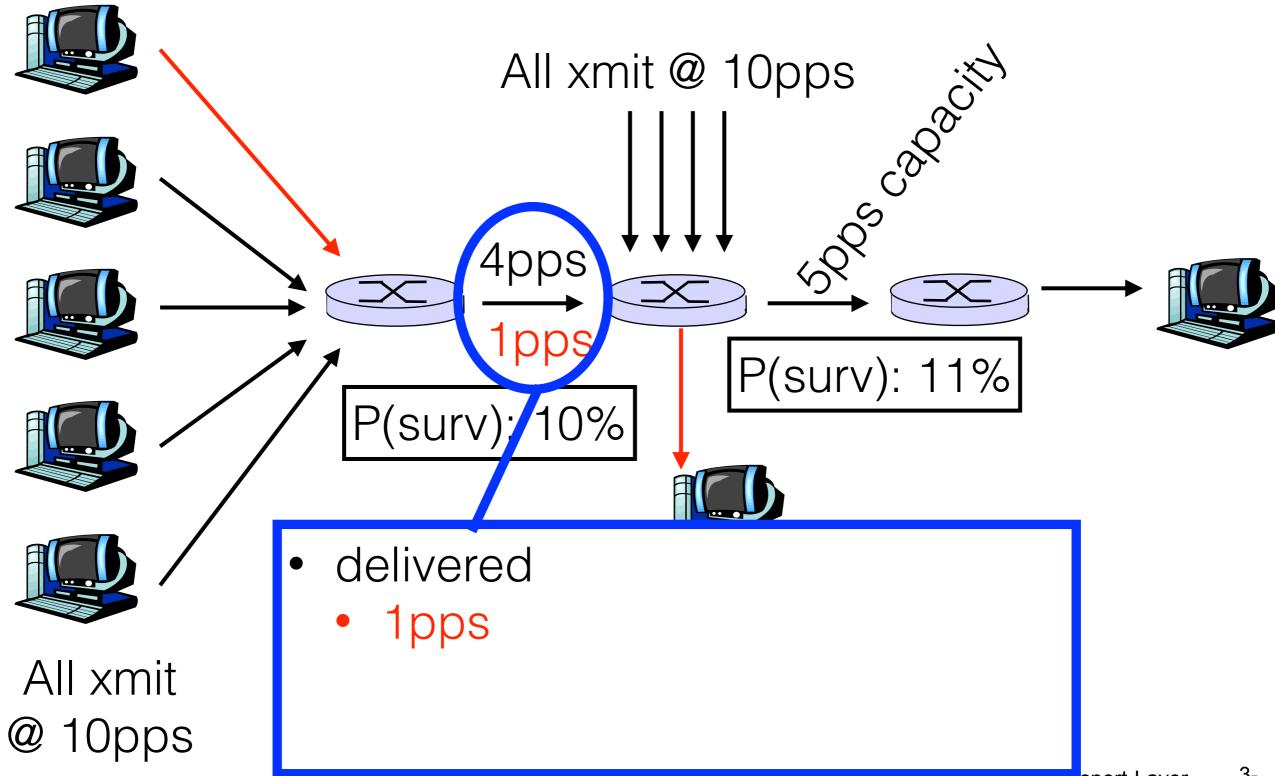
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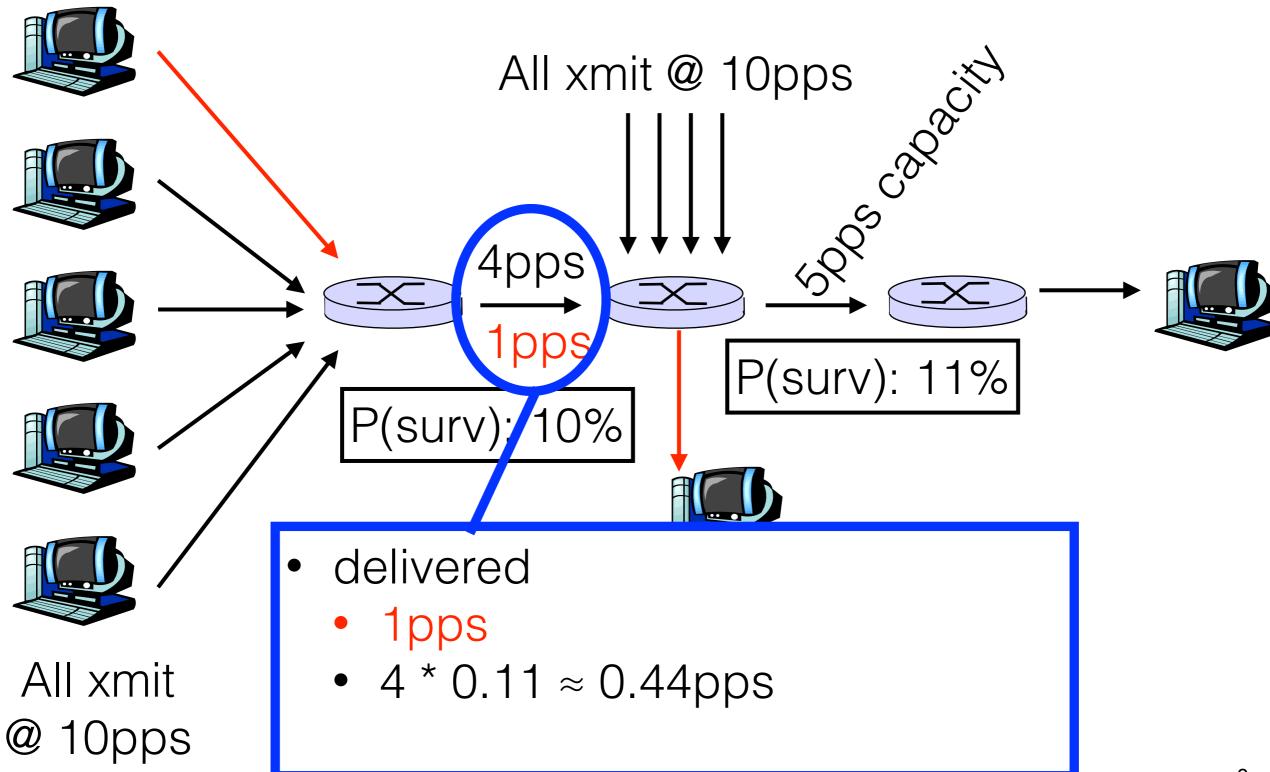
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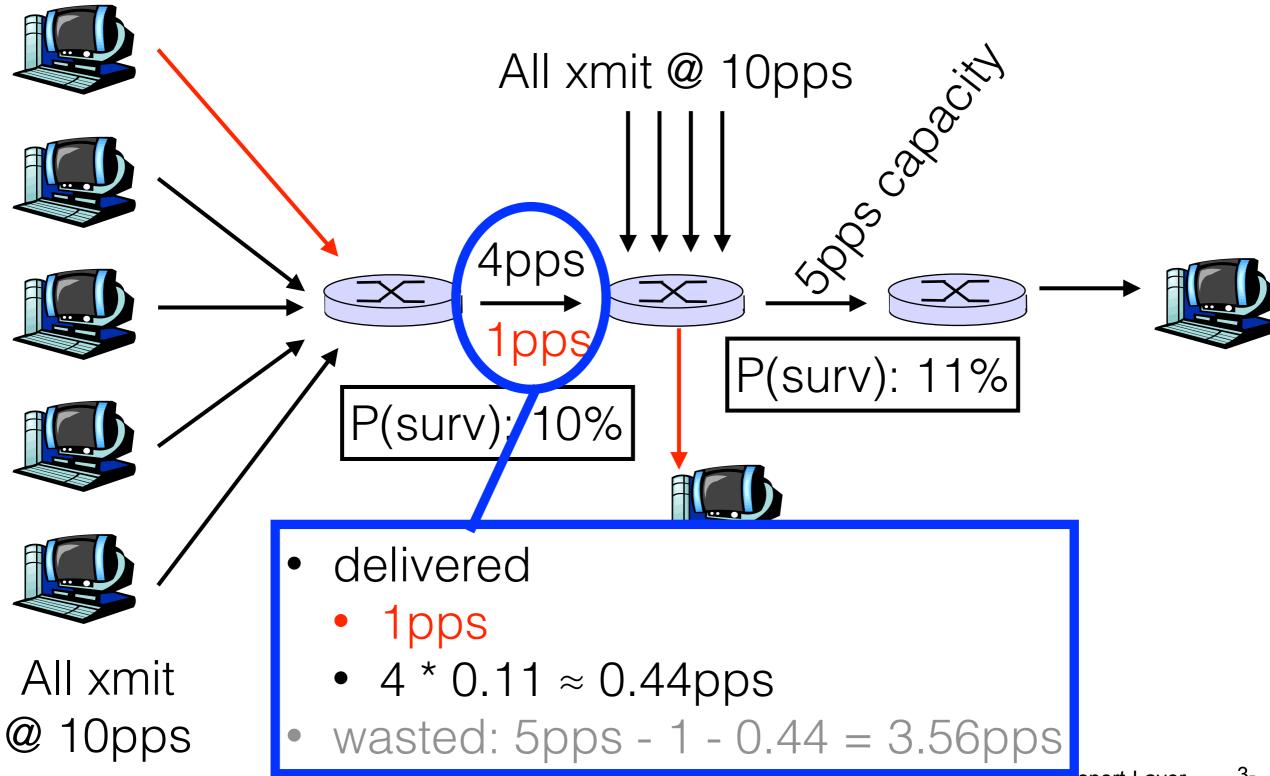










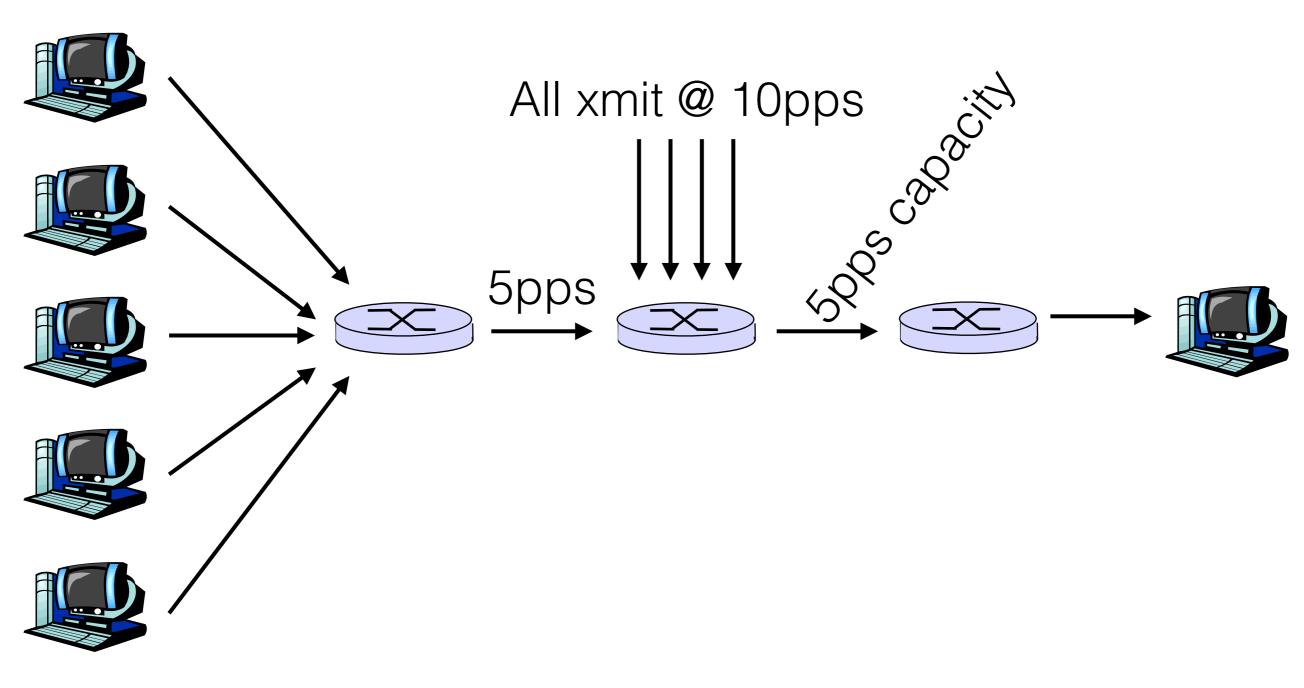


**Application** Transport Network Data Link **Physical** 

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- \*"tragedy of the commons" problem

Two broad approaches towards congestion control:

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- no explicit feedback from network
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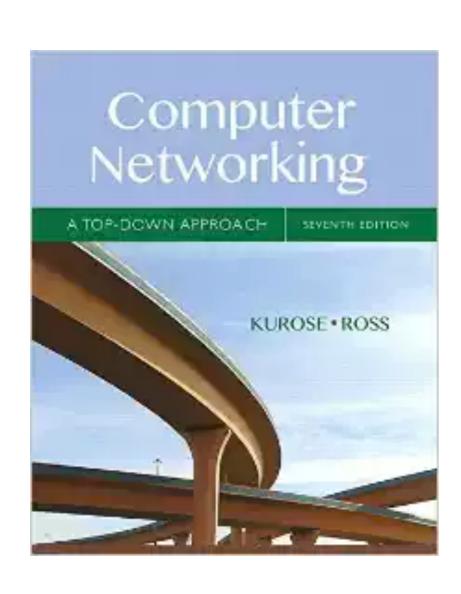
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# network-assisted congestion control:

- routers provide feedback to end systems
  - single bit indicating congestion (SNA, DECbit, TCP/IP ECN, ATM)
  - explicit rate sender should send at

# Reading Along ...



3.7:TCP Congestion
Control

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  - adjusts based on the estimated state of the network
  - cwnd is sender state and not transmitted to the receiver

- \*Assume loss means the network is congested
  - not always a good assumption

#### TCP Congestion Control

- \*Additive increase, multiplicative decrease (AIMD)
  - when no congestion detected, increase the rate (cwnd) linearly
  - when congestion detected, decrease the rate (cwnd) by half

#### TCP Congestion Control: details

\*sender limits transmission:

LastByteSent-LastByteAcked

≤ cwnd

roughly,

rate = 
$$\frac{\text{cwnd}}{\text{RTT}}$$
 Bytes/sec

How does sender perceive congestion?

loss event = timeout or3 duplicate acks

\*cwnd is dynamic, function of perceived network congestion

## Getting Started

\*OK, so, send at most cwnd bytes

\*But, how do we get started?

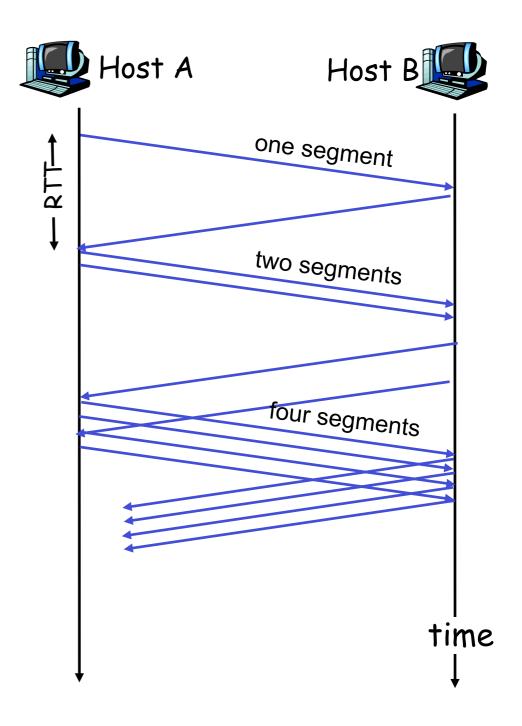
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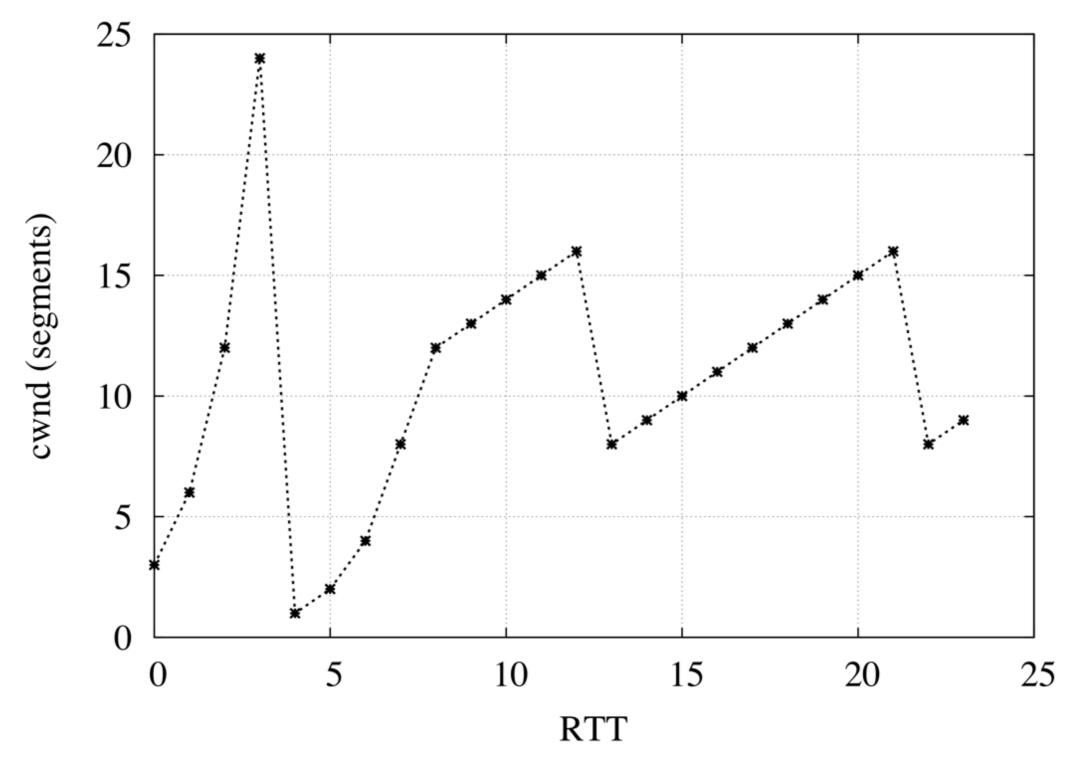
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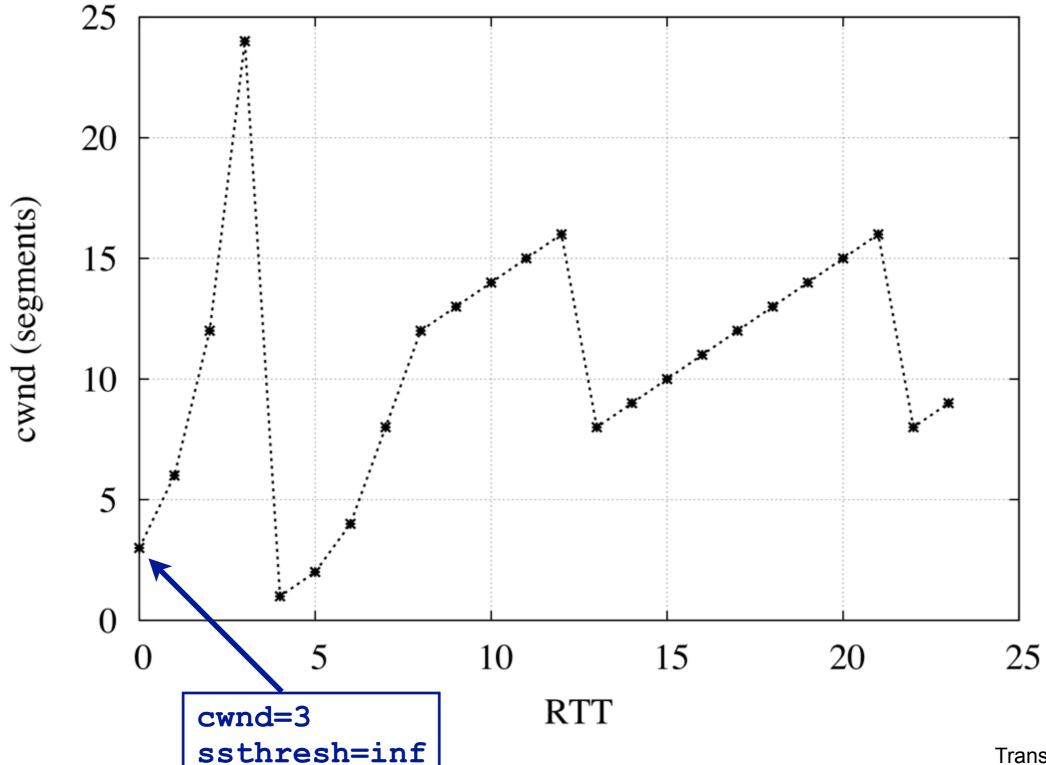
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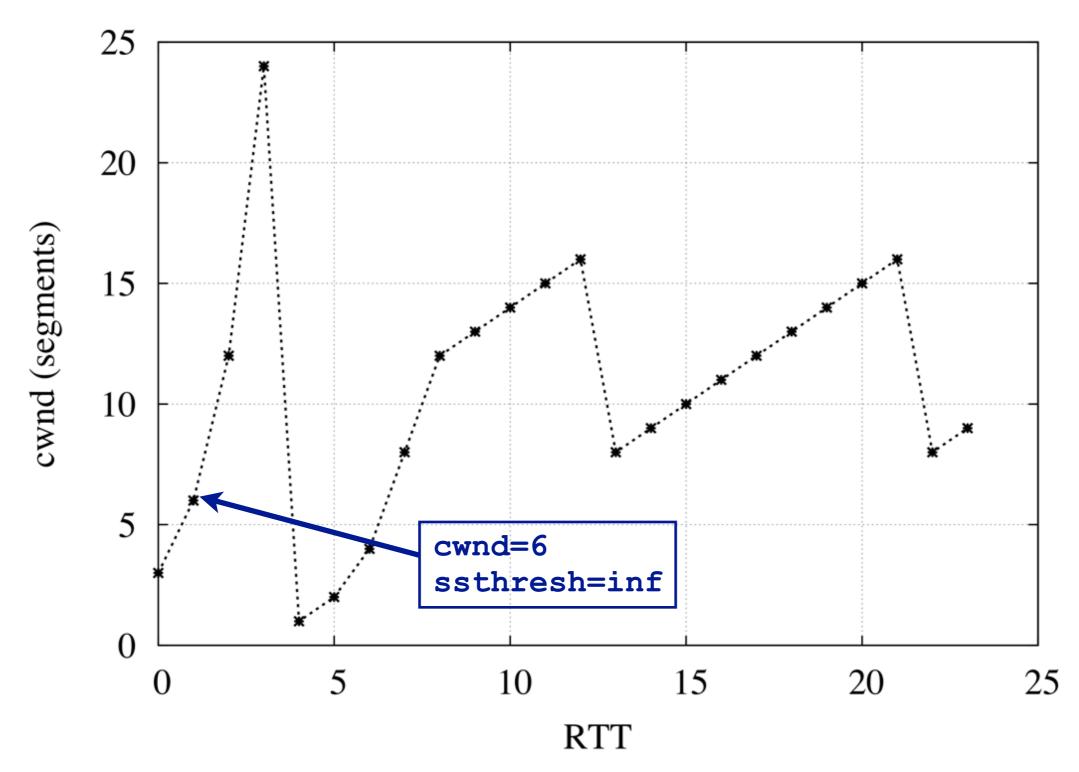
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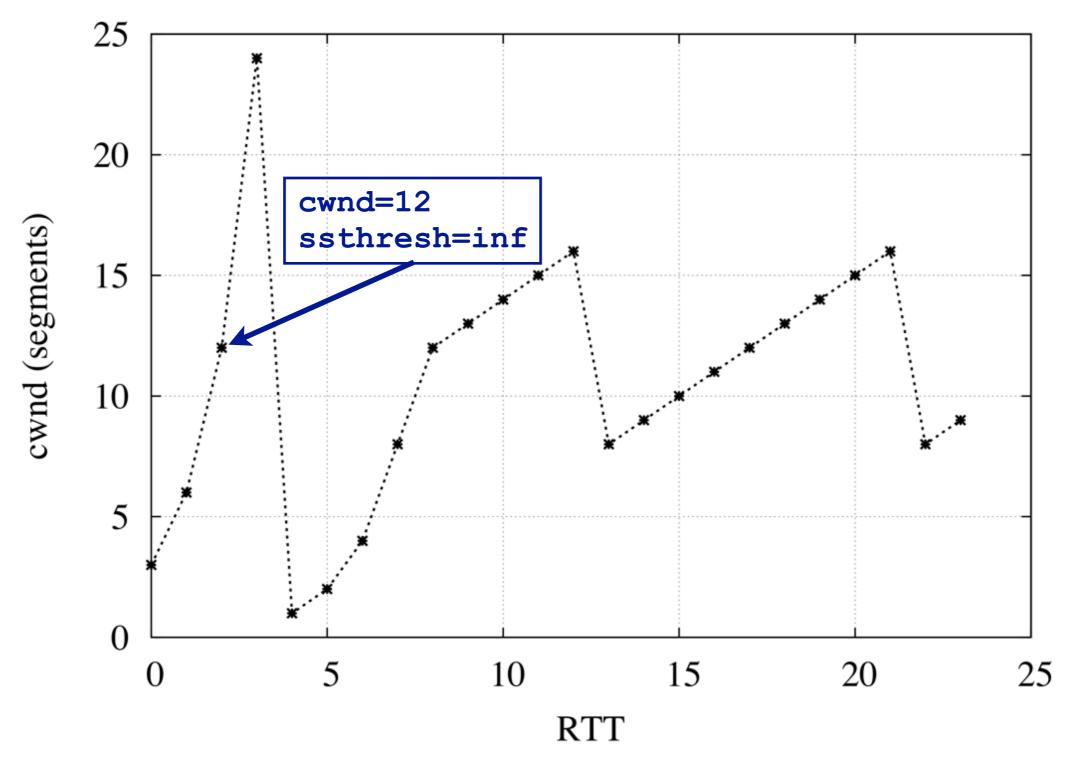
#### Philosophy:

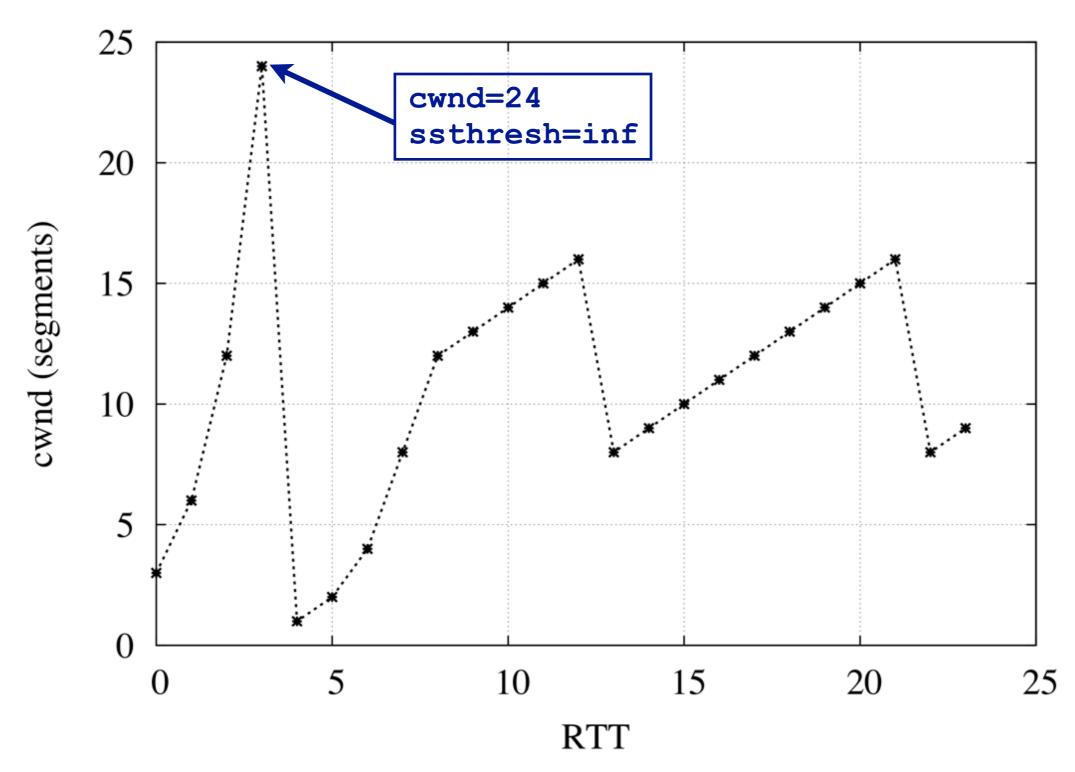
- \* 3 dup ACKs indicates network capable of delivering some segments
- \* timeout indicates a "more alarming" congestion scenario

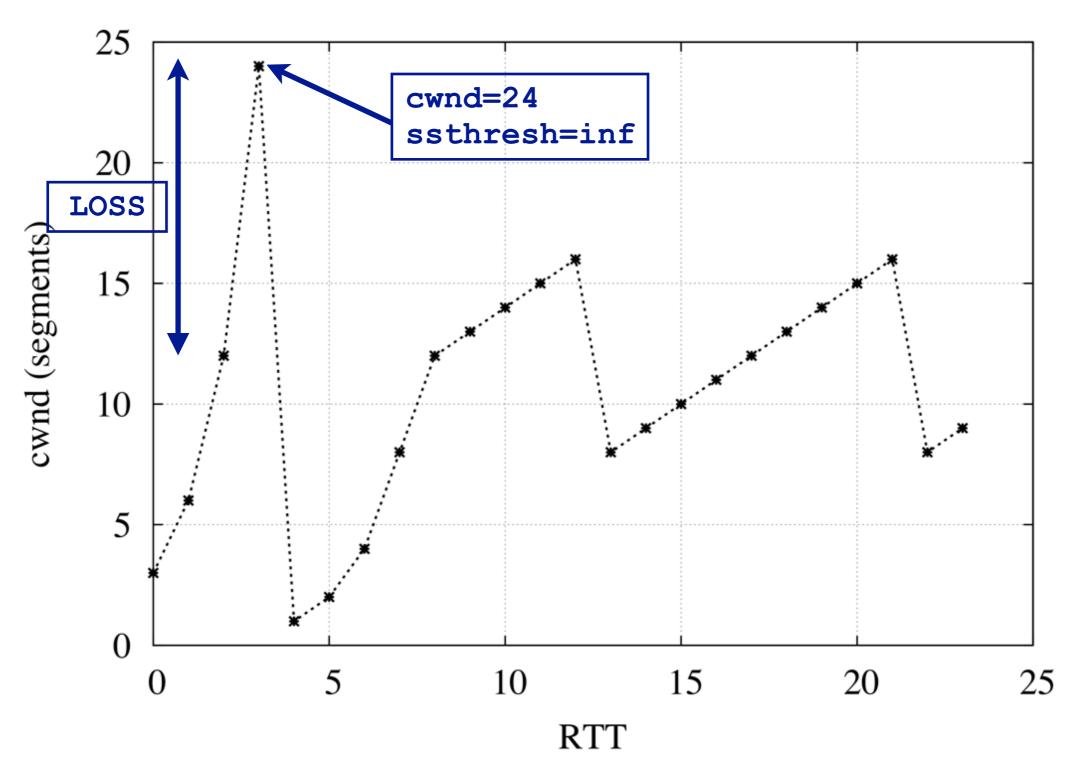


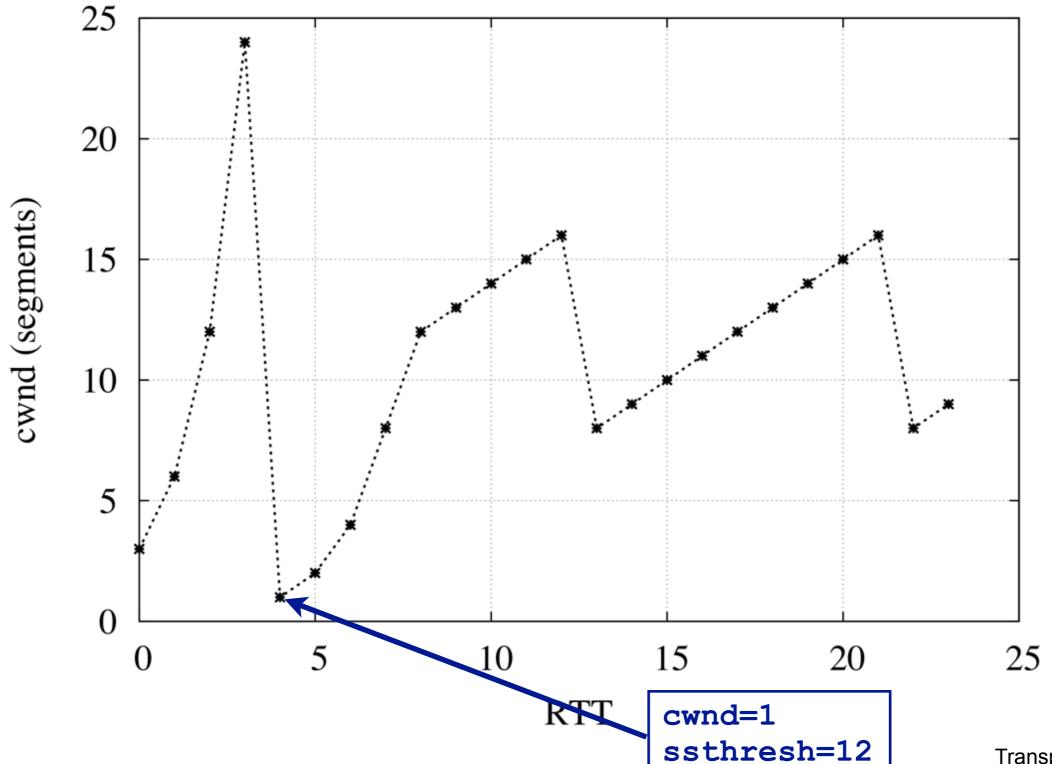


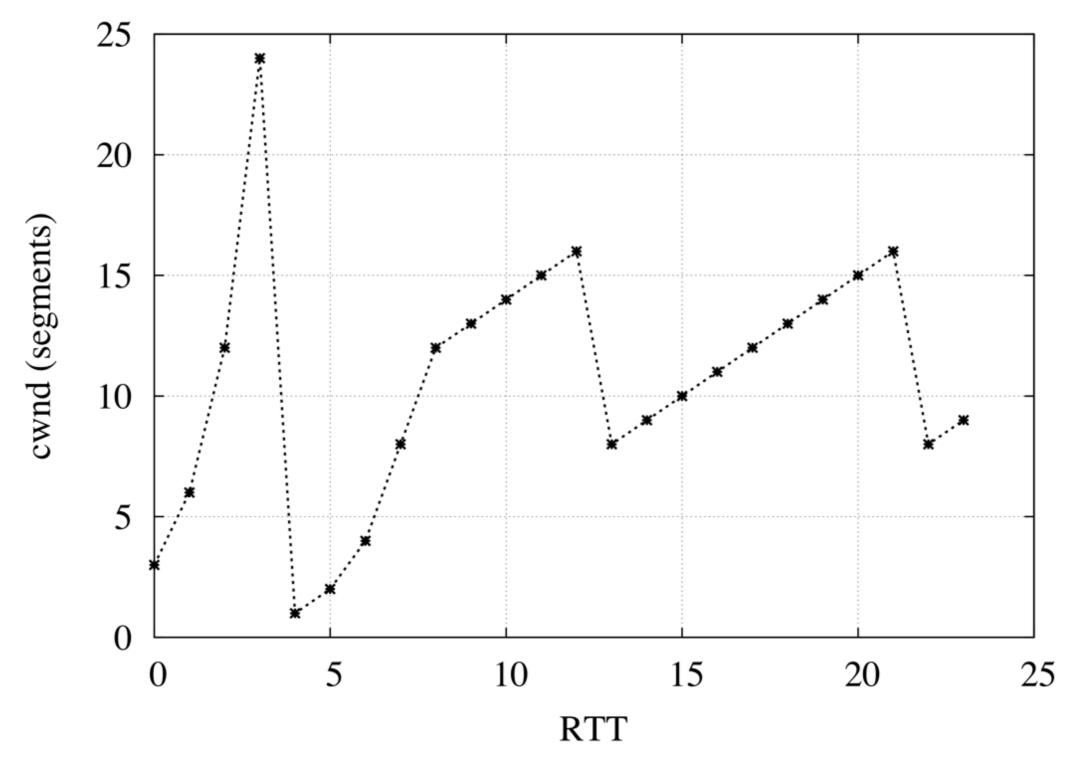


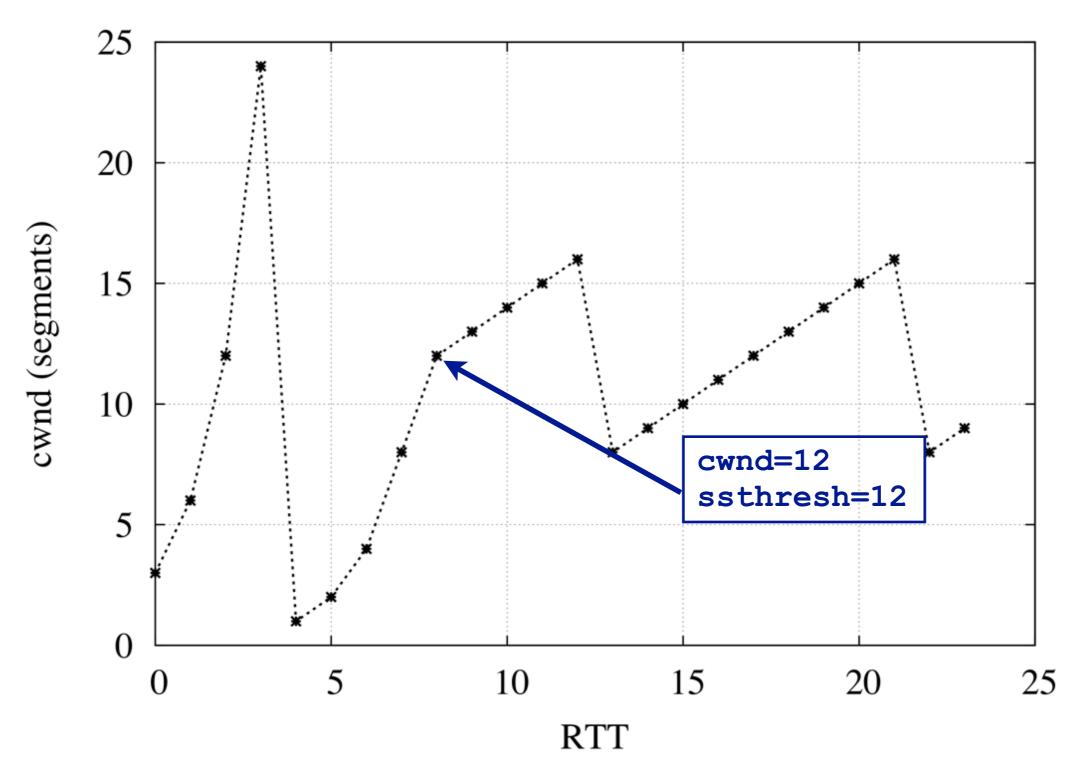


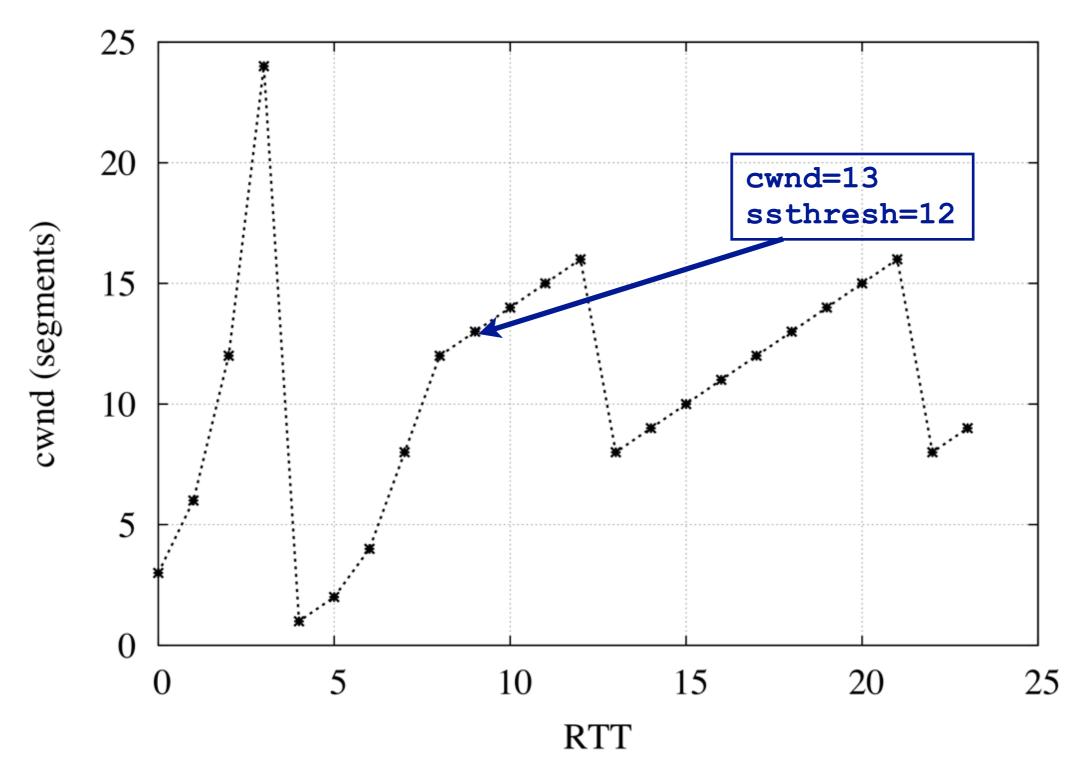


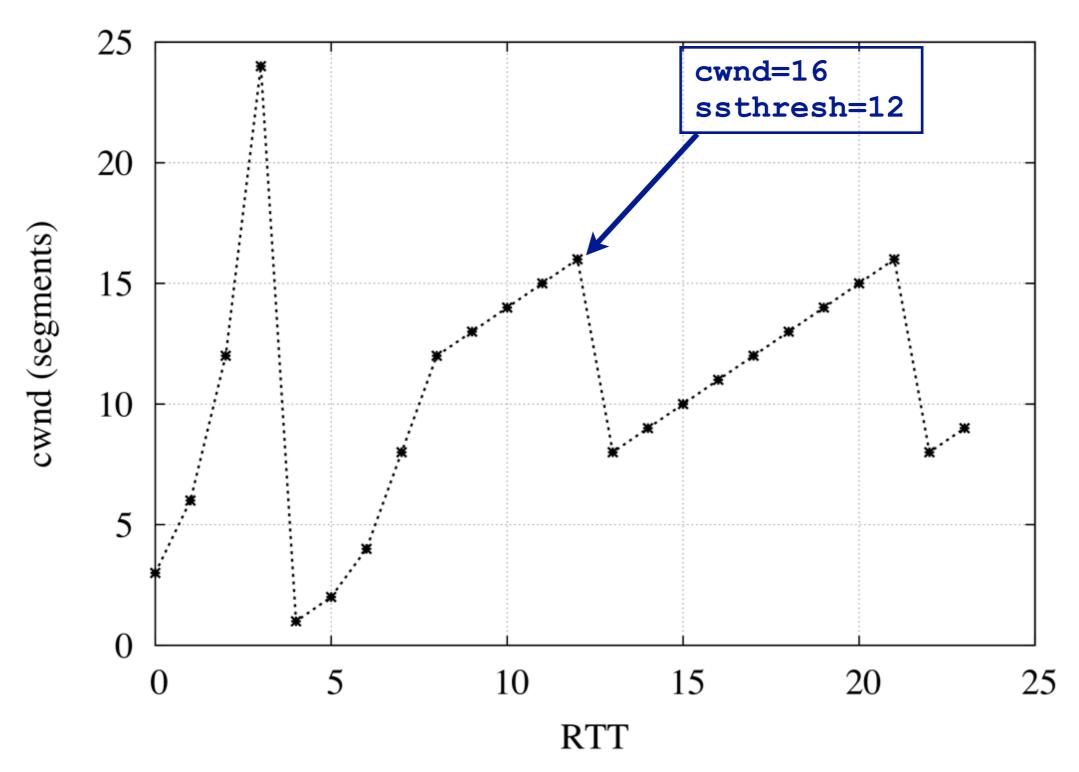


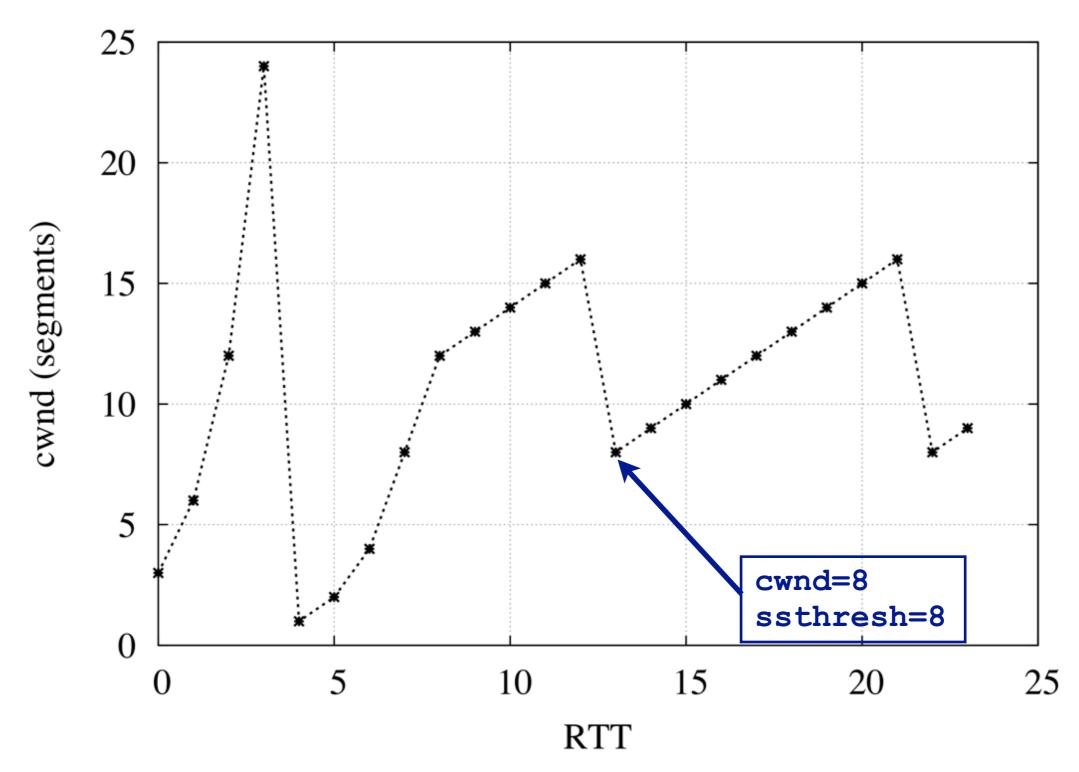




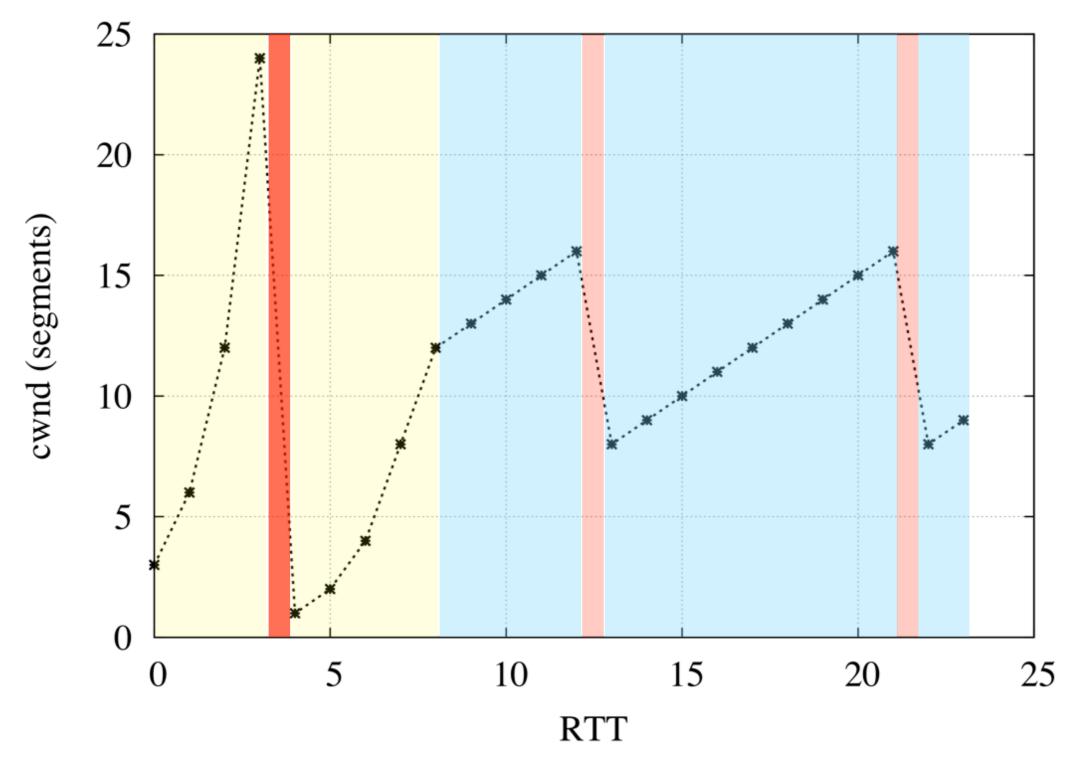




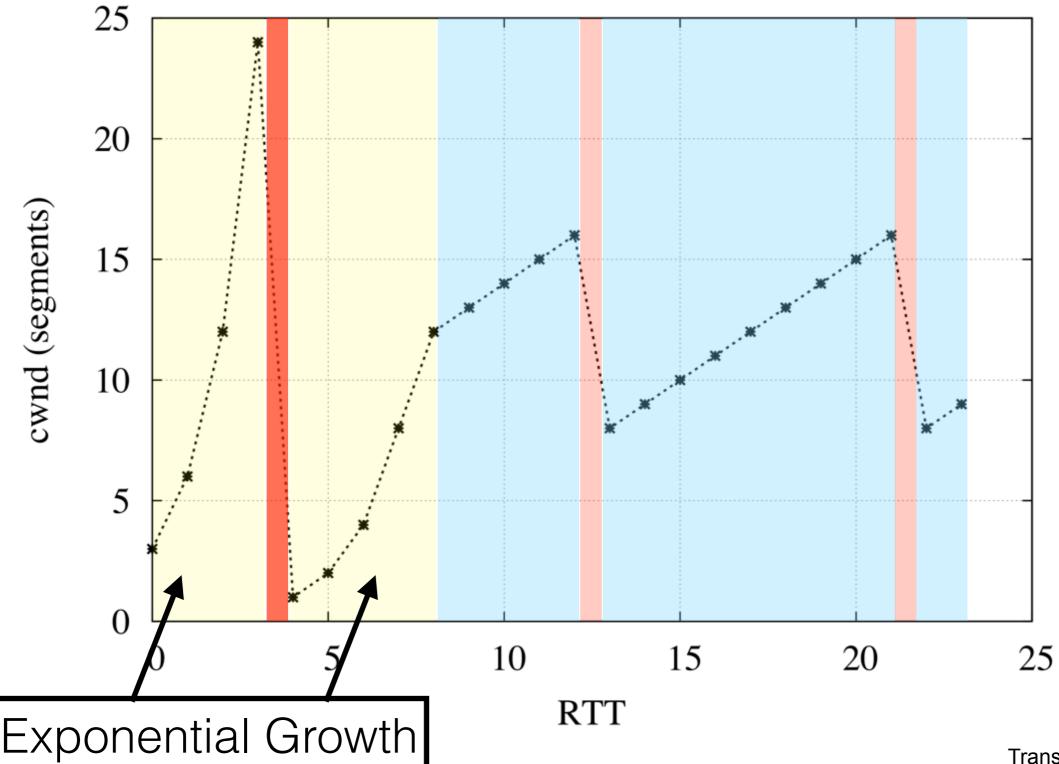




#### cwnd Evolution



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