

# Analysis of the ECMAScript Memory Model : A Program Transformation Perspective

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# Abstract

# Abrégé

# Acknowledgements

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# Chapter 1

## Introduction

## Chapter 2

# Background

## Chapter 3

### Related Work

# Chapter 4

## The Memory Model

This chapter starts with exposing some problems with the existing specifications of the model. The latter part introduces the key components of the model that we find essential to address program transformations. We start by introducing Agents and Event sets, followed by various Binary Relations defined between events. We then introduce certain helper definitions that prove useful in understanding the Axioms of the model. We then use the binary relations and definitions to specify the Axioms of the model. Lastly, we define Races followed by defining what a Consistent Execution is as per the specification of the model.

---

The model we consider is the current draft specification (cite here) of ECMAScript standard. This draft as far as the memory model goes has remain unchanged, so we believe our work will also be of use to those working on this model. The specification is claimed to be axiomatic by definition, which should, in our view remove the complexities of the rest of the standard from the semantics of the model. However, as we noted, there were quite some concerns with it:

**The Model is Quite Algorithmic** Although the standard states that the model is not supposed to be operational, the specifications of the model state otherwise. There are quite a few abstract operations which are not necessary to understand the semantics of the model. As an example, consider one of the Axioms of the model below as stated by the standard.

### 28.7.2 Coherent Reads

A candidate execution *execution* has coherent reads if the following abstract operation returns **true**.

1. For each ReadSharedMemory or ReadModifyWriteSharedMemory event *R* of *SharedDataBlockEventSet(execution)*, do
  - a. Let *Ws* be *execution*.[[ReadsBytesFrom]](*R*).
  - b. Let *byteLocation* be *R*.[[ByteIndex]].
  - c. For each element *W* of *Ws*, do
    - i. If (*R*, *W*) is in *execution*.[[HappensBefore]], then
      1. Return **false**.
    - ii. If there is a WriteSharedMemory or ReadModifyWriteSharedMemory event *V* that has *byteLocation* in its range such that the pairs (*W*, *V*) and (*V*, *R*) are in *execution*.[[HappensBefore]], then
      1. Return **false**.
    - iii. Set *byteLocation* to *byteLocation* + 1.
2. Return **true**.

Figure 4.1: The definition for Coherent Reads

The above axiom is specified as a return value to an abstract operation. While understanding this requires one to know the definitions of *Ws*, *execution*, *SharedDataBlockEventSet* abstract operation, etc. we believe this is not needed as to understand what the axiom is about, which informally can be stated as below in two points:

- A read's value cannot come from a write that has happened after it.
- A read's value cannot come from a write that has been overwritten by some other write.

Axiomatically, we define the above two points using simple binary relations that

we derive from the specification to exist among events. The entire specification is structured in a similar way.

**Certain Unnecessary Definitions** Certain abstract operations are not required to capture the semantics of the model. One such example is in the figure below:

#### 28.5.4 ComposeWriteEventBytes ( *execution*, *byteIndex*, *Ws* )

The abstract operation ComposeWriteEventBytes takes arguments *execution* (a candidate execution), *byteIndex* (a non-negative integer), and *Ws* (a List of WriteSharedMemory or ReadModifyWriteSharedMemory events). It performs the following steps when called:

1. Let *byteLocation* be *byteIndex*.
2. Let *bytesRead* be a new empty List.
3. For each element *W* of *Ws*, do
  - a. Assert: *W* has *byteLocation* in its range.
  - b. Let *payloadIndex* be *byteLocation* - *W*.[[ByteIndex]].
  - c. If *W* is a WriteSharedMemory event, then
    - i. Let *byte* be *W*.[[Payload]][*payloadIndex*].
  - d. Else,
    - i. Assert: *W* is a ReadModifyWriteSharedMemory event.
    - ii. Let *bytes* be ValueOfReadEvent(*execution*, *W*).
    - iii. Let *bytesModified* be *W*.[[ModifyOp]](*bytes*, *W*.[[Payload]]).
    - iv. Let *byte* be *bytesModified*[*payloadIndex*].
  - e. Append *byte* to *bytesRead*.
  - f. Set *byteLocation* to *byteLocation* + 1.
4. Return *bytesRead*.

Figure 4.2: The definition for Compose Write Event Bytes

The above figure is the definition of an abstract operation. To understand what this operation does, one must know the meaning of the terms *ModifyOp*, *Payload*, the list *Ws*, and also know what the argument *ByteIndex* signifies. In its essence, the above operation gives the read-values read by a single write by collecting the values



from their corresponding writes. We realized that one need not know this operation nor understand its function as it does not play a role in the semantics of the model. Other such abstract operations are *ValueOfReadEvent* and *ValidChosenReads*.

**Still a bit verbose** The entire model, is still quite verbose, which makes it difficult to understand the main objective of the model semantics. The following figure is the std specification of another Axiom

#### 28.7.4 Sequentially Consistent Atomics

For a candidate execution *execution*, memory-order is a strict total order of all events in *EventSet(execution)* that satisfies the following.

- For each pair  $(E, D)$  in *execution*.[[HappensBefore]],  $(E, D)$  is in memory-order.
- For each pair  $(R, W)$  in *execution*.[[ReadsFrom]], there is no *WriteSharedMemory* or *ReadModifyWriteSharedMemory* event  $V$  in *SharedDataBlockEventSet(execution)* such that  $V$ .[[Order]] is SeqCst, the pairs  $(W, V)$  and  $(V, R)$  are in memory-order, and any of the following conditions are true.
  - The pair  $(W, R)$  is in *execution*.[[SynchronizesWith]], and  $V$  and  $R$  have equal ranges.
  - The pairs  $(W, R)$  and  $(V, R)$  are in *execution*.[[HappensBefore]],  $W$ .[[Order]] is SeqCst, and  $W$  and  $V$  have equal ranges.
  - The pairs  $(W, R)$  and  $(W, V)$  are in *execution*.[[HappensBefore]],  $R$ .[[Order]] is SeqCst, and  $V$  and  $R$  have equal ranges.

NOTE 1 This clause additionally constrains SeqCst events on equal ranges.

- For each *WriteSharedMemory* or *ReadModifyWriteSharedMemory* event  $W$  in *SharedDataBlockEventSet(execution)*, if  $W$ .[[Order]] is SeqCst, then it is not the case that there is an infinite number of *ReadSharedMemory* or *ReadModifyWriteSharedMemory* events in *SharedDataBlockEventSet(execution)* with equal range that is memory-order before  $W$ .

NOTE 2 This clause together with the forward progress guarantee on agents ensure the liveness condition that SeqCst writes become visible to SeqCst reads with equal range in finite time.

A candidate execution has sequentially consistent atomics if a memory-order exists.

Figure 4.3: The Axiom of Sequentially Consistent Atomics

The above figure, though written concisely, in our view still makes it difficult to understand. We reduced the above entire axiom into three main patterns using

binary relations that should not exist in any execution of a program. While the latter part is less of a semantic specification and more of a programming guideline while using relaxed memory accesses. (one can make countless counter-examples for this.)

Given the above concerns about the model specification in the standard, we formalized it axiomatically in the form of binary relations over events involved and axioms that place restrictions on certain binary relations using the others.

## 4.1 Agents, Events and their Types

**Agents** Agents represent threads in a concurrent program. As per the standard, they have more meaning than what we refer to here. However, with respect to the memory consistency model, we can safely abstract them to just represent threads/processes.

**Agent Cluster** Collection of agents concurrently communicating with each other through means of shared memory form an agent cluster. There can be multiple agent clusters. However, an agent can only belong to one agent cluster. Agents communicating through message passing do not belong in the same agent cluster.

For our purpose, we assume just one agent cluster having one shared memory.

**Agent Event List (*ael*)** Every agent is mapped to a list of events. The list represents the order in which the events are evaluated operationally<sup>1</sup>. We define *ael* as a mapping of each agent to a list of events.

## 4.2 Events

Agent execution is modelled in terms of events. An event is either an operation that involves (shared) memory access or that constrains the order of execution of multiple events.

### 4.2.1 Event Types

Given an agent cluster, an *event set*  $\mathbf{E}$  is a collection of all events from the agent event lists. This set is composed of mainly two distinct subsets as follows:

---

<sup>1</sup>The standard refers this to be an Event List, but we find it a bit misleading as it does not signify a list for each agent. Hence we name it as Agent Event List.

- **Shared Memory (*SM*) Events**

This set is composed of two sets of events; those that write to shared memory called Write events (***W***) and those that read from shared memory called Read events (***R***). Events that belong to both Write and Read events are called Read-Modify-Write.

- **Synchronize (*S*) Events** These events only restrict the ordering of execution of events by agents. For instance *lock* and *unlock* type of events can be categorized under Synchronize events. However, this is not stated in the specification<sup>2</sup>.

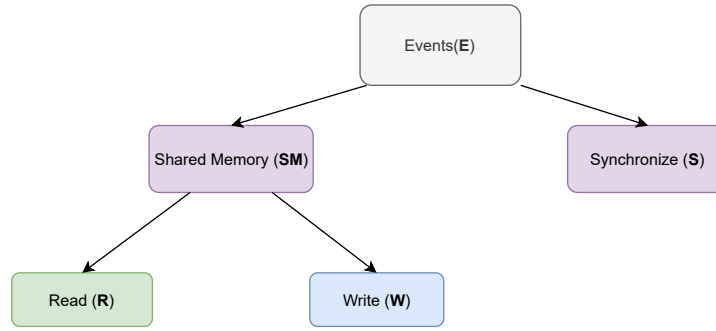


Figure 4.4: The hierarchical categories of different sets of events.

### 4.2.2 Range ( $\mathfrak{R}$ )

Each of the *shared memory events* are associated with a contiguous range of memory on which it operates. Range is a function that maps a shared memory event to the range<sup>3</sup> it operates on. This we represent as a starting index  $i$  and a size. So we could represent the range of a write event  $w$  as

$$\mathfrak{R}(w) = (i, s)$$

---

<sup>2</sup>The features of *Lock* and *Unlock* events is actually not something given to the programmer to use in Javascript. They are used to implement the feature *wait* and *notify* that the programmer can use which adhere to the semantics of *futexes* in Linux. Hence, in the original standard of the model, the distinction between lock and unlock is not made, and it is simply stated as Synchronize Event.

<sup>3</sup>The range as per the ECMAScript standard denotes only the set of contiguous byte indices. The starting byte index is kept separate. We find this to be unnecessary. Hence we define range to have starting index and size.

We define the two binary operators below on ranges:

1. Intersection ( $\cap_{\mathfrak{R}}$ ) - Set of byte indices common to both ranges.
2. Union ( $\cup_{\mathfrak{R}}$ ) - A unique set of byte indices that exist in both the ranges.

Two Ranges can be *disjoint*, *overlapping* or *equal*. We use the binary operators to define these three possibilities between ranges of events  $e$  and  $d$  :

1. Disjoint  $\mathfrak{R}(e) \cap_{\mathfrak{R}} \mathfrak{R}(d) = \phi$
2. Overlapping  $(\mathfrak{R}(e) \cap_{\mathfrak{R}} \mathfrak{R}(d) \neq \phi) \wedge (\mathfrak{R}(e) \cap_{\mathfrak{R}} \mathfrak{R}(d) \neq \mathfrak{R}(e) \cup_{\mathfrak{R}} \mathfrak{R}(d))$  -
3. Equal  $\mathfrak{R}(e) \cap_{\mathfrak{R}} \mathfrak{R}(d) = \mathfrak{R}(e) \cup_{\mathfrak{R}} \mathfrak{R}(d)$  - In simple terms, we define equality as  $\mathfrak{R}(e) = \mathfrak{R}(d)$

### 4.2.3 Event Order / Event Access Mode

Order signifies the sequence in which event actions are visible to different agents as well as the order in which they are executed by the agents themselves. In our context, there are mainly three types (in C11 memory model, they are called access modes) for each shared memory event that tells us the kind of ordering that it enforces.

1. **Sequentially Consistent** (*sc*) - Events of this type are *atomic*<sup>4</sup> in nature. There is a strict global total ordering of such events which is agreed upon by all agents in the agent cluster.
2. **Unordered** (*uo*) - Events of this type are considered *non-atomic* and can occur in different orders for each concurrent process. There is no fixed global order respected by agents for such events.
3. **Initialize** (*init*) - Events of this type are used to initialize the values in memory before they are accessed by agent events.

All events of type *init* are writes and all Read-Modify-Write events are of type *sc*. We represent the type of events in the memory consistency rules in the format “*event : type*”. When representing events in examples, the type would be represented as a subscript: *event<sub>type</sub>*.

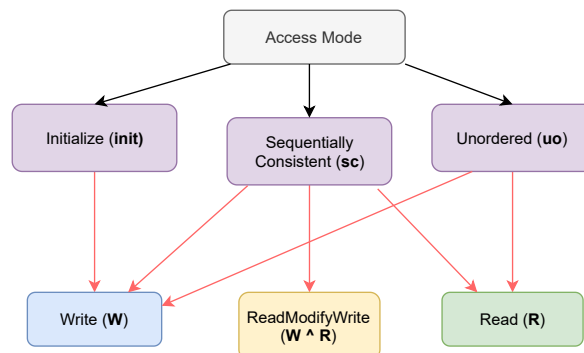


Figure 4.5: Access Modes

#### 4.2.4 Tear Free (*tf*) or Tearing *!tf*)

Additionally, each shared-memory event is also associated with whether they are tear-free or not. OEvents that tear are non-aligned accesses requiring more than one memory access. Events that are tear-free are aligned and should appear to be serviced in one memory fetch<sup>5</sup>.

We represent the tearing of events in the memory consistency rules in the format “*event* : *tf*/*!tf*”. When representing events in examples, the type would be represented as a subscript: *event*<sub>*tf*/*!tf*</sub>.

### 4.3 Relation among events

We now describe a set of binary relations between events. These relations help us describe the consistency rules.

<sup>4</sup>The word *atomic* does not imply the events are evaluated using just one instruction. For example, a 64-bit sequentially consistent write on a 32-bit system has to be done with two subsequent memory actions. But its intermediate state of write must not be seen by any other agent. In an abstract sense, this event must appear ‘*atomic*’. The *atomic* here also refers to implications of whether an event’s consequence is visible to all other agents in the same global total order or not. The compiler must ensure that for each specific target hardware, such guarantees are satisfied.

<sup>5</sup>It is not clear whether the alignment is with respect to specific hardware or not. The notion of one memory fetch may not be possible for all hardware practically, but it is something that must appear so. We will see a rule for ensuring this in the memory consistency rules.

### 4.3.1 Read-Write event relations

There are two basic relations that assist us in reasoning about read and write events.

**Read-Bytes-From** ( $\overrightarrow{rbf}$ ) This relation maps every read event to a list of tuples consisting of write event and their corresponding byte index that is read. For instance, consider a read event  $r[i...(i+3)]$  and corresponding write events  $w_1[i...(i+3)]$ ,  $w_2[i...(i+4)]$ . One possible  $\overrightarrow{rbf}$  relation could be represented as

$$e \xrightarrow{rbf} \{(d1, i), (d2, i+1), (d2, i+2)\}$$

or having individual binary relation with each write-index pair as

$$e \xrightarrow{rbf} (d1, i), e \xrightarrow{rbf} (d2, i+1) \text{ and } e \xrightarrow{rbf} (d2, i+2).$$

**Reads-From** ( $\overrightarrow{rf}$ ) This relation, is similar to the above relation, except that the byte index details are not involved in the composite list. So for the above example, the  $rf$  relation would be represented either as  $e \xrightarrow{rf} (d1, d2)$  or individual binary read-write relation as  $e \xrightarrow{rf} d1$  and  $e \xrightarrow{rf} d2$ . Figure below is an example of a program with its outcome (read values) shown in terms of reads-from relations.

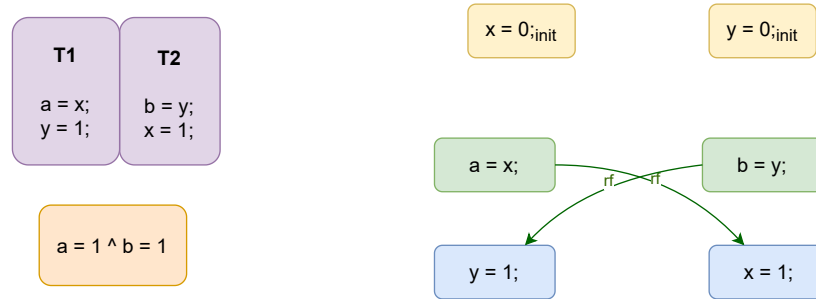


Figure 4.6: An example to show the reads-from relations that are drawn for the example program between read and write events.

### 4.3.2 Agent-Synchronizes With (ASW)

It is a list for each agent that consist of ordered tuples of synchronize events. These tuples specify ordering constraints among agents at different points of execution. So such a list for an agent  $k$  would be represented like:

$$ASW_k = \{\langle s_1, s_2 \rangle, \langle s_3, s_4 \rangle \dots\}$$

For every pair in the list, the second event belongs to the parent agent and the first belongs to another agent it synchronized with<sup>6</sup>.

$$\forall i, j > 0, \langle s_1, s_2 \rangle \in ASW_j \Rightarrow s_2 \in ael(k)$$

The figure below shows an example of this relation among two agents.

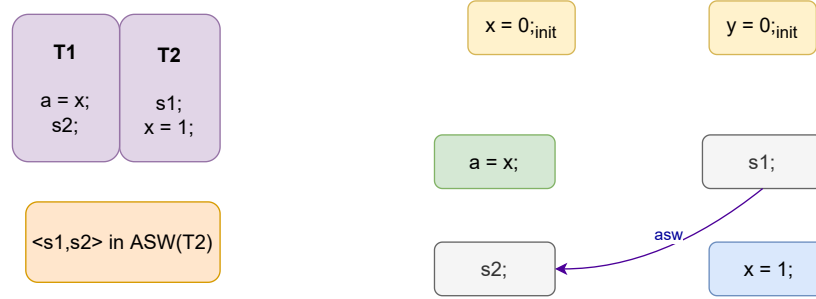


Figure 4.7: An example to show the reads-from relations that are drawn for the example program between read and write events.

## 4.4 Ordering Relations among Events

### 4.4.1 Agent Order ( $\overrightarrow{ao}$ )

It is a union of total order among events belonging to the same agent event list. It is analogous to intra-thread ordering. For example, if two events  $e$  and  $d$  belong to the same agent event list, then either  $e \overrightarrow{ao} d$  or  $d \overrightarrow{ao} e$ .

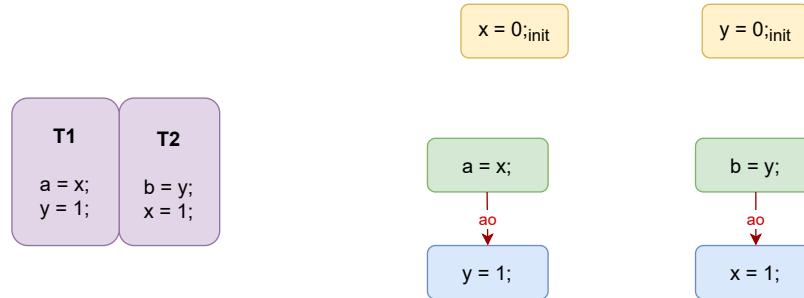


Figure 4.8: An example with agent order among the events.

<sup>6</sup>This is analogous to the property that every unlock must be paired with a subsequent lock, which enforces the condition that a lock can be acquired only when it has been released.

### 4.4.2 Synchronize-With Order ( $\xrightarrow{sw}$ )

Binary relation between two events that establish synchronization between multiple agents. It is a composition of two sets:

1. All pairs belonging to  $ASW$  of every agent belongs to this ordering relation.

$$\forall i, j > 0, \langle e_i, e_j \rangle \in ASW \Rightarrow e_i \xrightarrow{sw} e_j$$

2. Specific reads-from pairs also belong to this ordering relation<sup>7</sup>.

$$(r \xrightarrow{rf} w) \wedge r:sc \wedge w:sc \wedge (\mathfrak{R}(r) = \mathfrak{R}(w)) \Rightarrow (w \xrightarrow{sw} r)$$

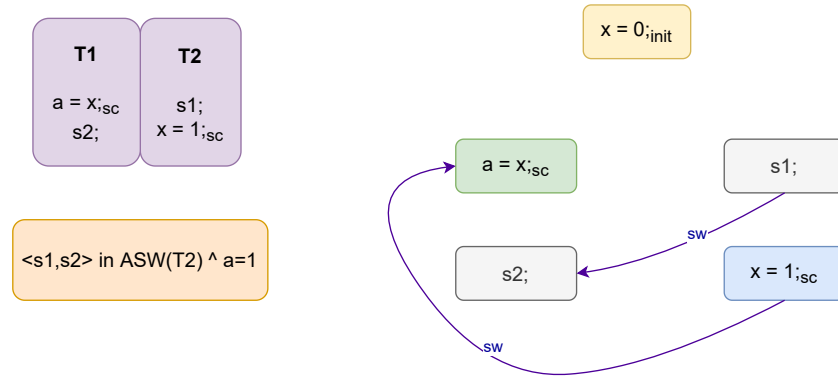


Figure 4.9: An example with synchronize with relations among the events.

### 4.4.3 Happens Before Order ( $\xrightarrow{hb}$ )

A transitive order on events, composed of the following:

1. Every agent-ordered relation is also a happens-before relation

$$(e \xrightarrow{ao} d) \Rightarrow (e \xrightarrow{hb} d)$$

<sup>7</sup>Note that for the second condition, both ranges of events have to be equal. This however, does not mean that the read cannot read from multiple write events. (the read-from relation here is not functional.)



2. Every synchronize-with relation is also a happens-before relation

$$(e \xrightarrow{sw} d) \Rightarrow (e \xrightarrow{hb} d)$$

3. Initialize type of events happen before all shared memory events that have overlapping ranges with them.

$$\forall e, d \in SM \wedge e: init \wedge (\mathcal{R}(e) \cap \mathcal{R}(d) \neq \emptyset) \Rightarrow e \xrightarrow{hb} d$$

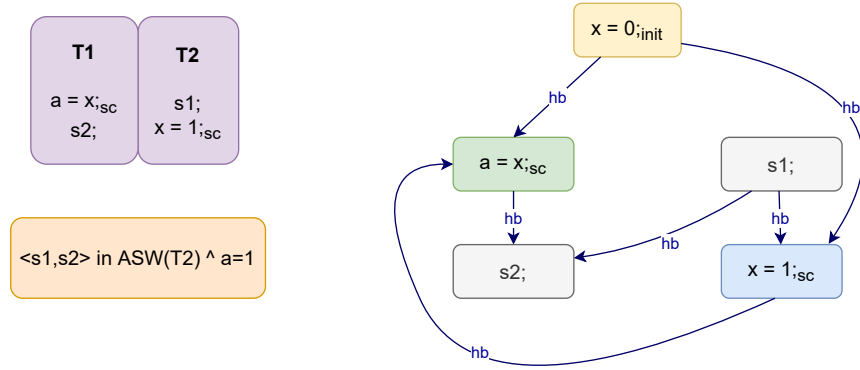


Figure 4.10: An example with all the types of happens-before relations between events.

#### 4.4.4 Memory Order ( $\xrightarrow{mo}$ )

A *total order* on all events that respects happens-before order.

$$e \xrightarrow{hb} d \Rightarrow e \xrightarrow{mo} d$$

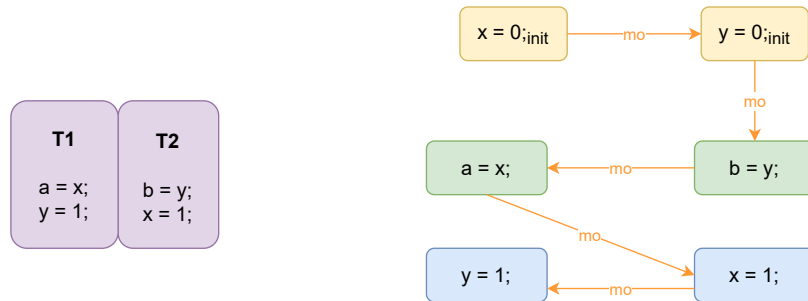


Figure 4.11: An example with a memory order (total) among all events.

An interesting part is that memory order, though total, is a bit undefined as to how it weaves together this total order given different init events. One can certainly make init events weaved among events that occur in each agent, thus making it sort of subjective as to when certain memory fragments are initialized. Discuss with Clark.

## 4.5 Helper Definitions

Before we go into the consistency rules. we define certain preliminary definitions that create a separation based on a program, the axiomatic events and the various ordering relations defined above. This will help us understand where the consistency rules actually apply.

**Definition 4.5.1.** *Program* A program is the source code without abstraction to a set of events and ordering relations. In our context, it is the original Javascript program.

**Definition 4.5.2.** *Candidate* A a collection of abstracted set of shared memory events of a program involved in one possible execution, with the  $\overrightarrow{ao}$  relations. We can think of this as each thread having a set of shared memory events to run in a given intra-thread ordering. An example of a candidate is shown in figure 4.12.

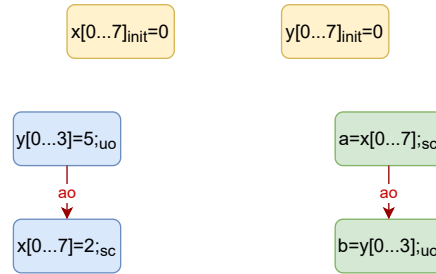


Figure 4.12: An example of a Candidate

**Definition 4.5.3.** *Candidate Execution* A Candidate with the addition of  $\overrightarrow{sw}$ ,  $\overrightarrow{hb}$  and  $\overrightarrow{mo}$  relations. This can be viewed as the witness/justification of an actual execution of a Program. Note that there can be many Candidate Executions for a given Candidate. The following figure shows an example of a candidate execution.

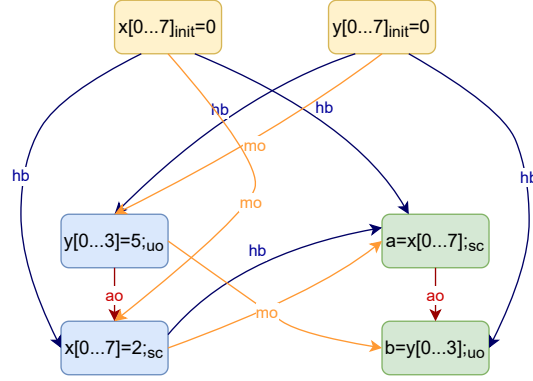


Figure 4.13: An example of an Execution based on Candidate above

**Definition 4.5.4.** *Observable Behavior*

The set of pairwise  $\xrightarrow{rf}/\xrightarrow{rbf}$  relations that result in one execution of the program. Think of this as our outcome of a program execution.

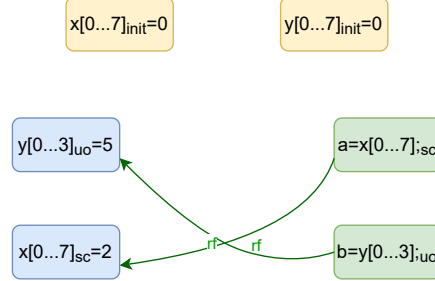


Figure 4.14: Observable Behavior

**Definition 4.5.5.** *Obs* We define  $Obs_P, Obs_C, Obs_{CE}$  as functions that take a program, candidate and candidate execution respectively and give the set of observable behaviors possible by them. We are not concerned with the specific elements in this set, but the relation between the output of each of these functions among each other.

Consider a program  $P$  whose candidates are  $C_1, C_2, \dots, C_n$ . Consider for each candidate  $C_i$ , the candidate executions  $CE_1, CE_2, \dots, CE_m$ <sup>8</sup>. Then, we have the following properties that hold:

$$Obs_P(P) = \cup_{i=1}^n Obs_C(C_i)$$

$$Obs_C(C_i) = \cup_{j=1}^m Obs_{CE}(CE_j)$$

<sup>8</sup>Note that the variables  $n$  and  $m$  need not be finite.

## 4.6 Valid Execution Rules (the Axioms)

We now state the memory consistency rules. The rules are on *Candidate Executions* which will place constraints on the possible *Observable behaviors* that may result from it.

### Axiom 1. Coherent Reads

There are certain restrictions of what a read event cannot see at different points of execution based on  $\overrightarrow{hb}$  relation with write events.

Consider a read event  $e$  and a write event  $d$  having at least overlapping ranges:

$$e \in R \wedge d \in W \wedge (\mathcal{R}(e) \cap_{\mathcal{R}} \mathcal{R}(d) \neq \phi).$$

- A read value cannot come from a write that has happened after it

$$e \overrightarrow{hb} d \Rightarrow \neg e \overrightarrow{rf} d.$$

The figure below pictorially depicts the pattern above where  $e$  cannot read from  $d$ .

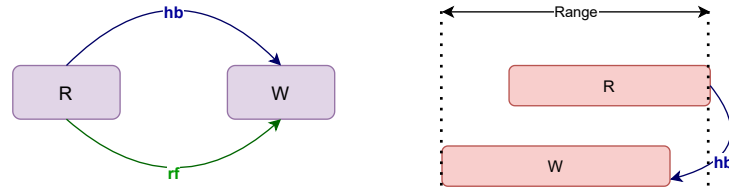


Figure 4.15: A read value cannot come from a write that has happened after it

- A read cannot read a specific byte address value from write if there is a write  $g$  that happens between them which modifies the exact byte address. Note that this rule would be on the  $rbf$  relation among two events.

$$d \overrightarrow{hb} e \wedge d \overrightarrow{hb} g \wedge g \overrightarrow{hb} e \Rightarrow \forall x \in (\mathcal{R}(d) \cap_{\mathcal{R}} \mathcal{R}(g) \cap_{\mathcal{R}} \mathcal{R}(e)), \neg e \overrightarrow{rbf} (d, x).$$

The figure below pictorially depicts the pattern where  $e$  cannot read certain bytes from  $d$ .

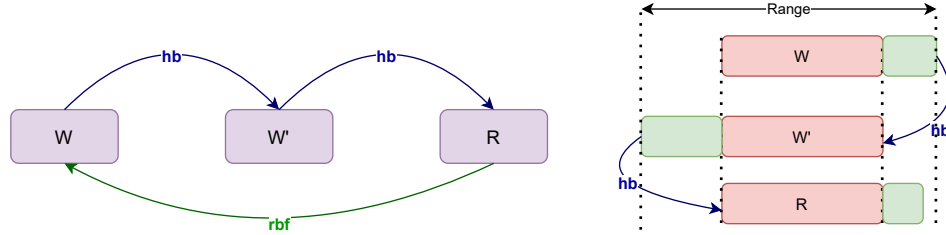


Figure 4.16: A read value cannot come from a write if there is a write that happens between them, writing to the same memory:

**Axiom 2. Tear-Free Reads**

*If two tear free writes  $d$  and  $g$  and a tear free read  $e$  all with equal ranges exist, then  $e$  can read only from one of them<sup>9</sup>.*

$$d:tf \wedge g:tf \wedge e:tf \wedge (\mathfrak{R}(d)=\mathfrak{R}(g)=\mathfrak{R}(e)) \Rightarrow ((e \xrightarrow{rf} d) \wedge (\neg e \xrightarrow{rf} g)) \vee ((e \xrightarrow{rf} g) \wedge (\neg e \xrightarrow{rf} d)).$$

*The following figure shows the pattern that is disallowed among all tear-free events.*

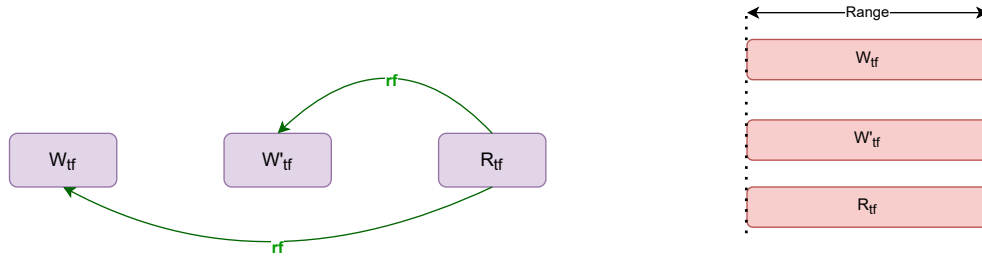


Figure 4.17: Pattern of Tear-free reads

**Axiom 3. Sequentially Consistent Atomics**

*To specifically define how events that are sequentially consistent affects what values a read cannot see, we assume the following memory order among writes  $d$  and  $g$*

<sup>9</sup>To recap a tear-free event cannot be separated into multiple small events that do the same operation. However, considering different hardware architectures, the notion of tear-free need not necessarily mean this. (eg: A 64bit tear-free write to be done in a 32bit system). In a more abstract sense, we need an event to appear 'tear-free'.

and a read  $e$  to be the premise for all the rules:

$$d \xrightarrow{mo} g \xrightarrow{mo} e.$$

- If all three events are of type  $sc$  with equal ranges, then  $e$  cannot read from  $d$

$$d:sc \wedge g:sc \wedge e:sc \wedge (\mathfrak{R}(d)=\mathfrak{R}(g)=\mathfrak{R}(e)) \Rightarrow \neg e \xrightarrow{rf} d.$$

The figure below depicts pictorially the pattern that is not allowed by this rule.

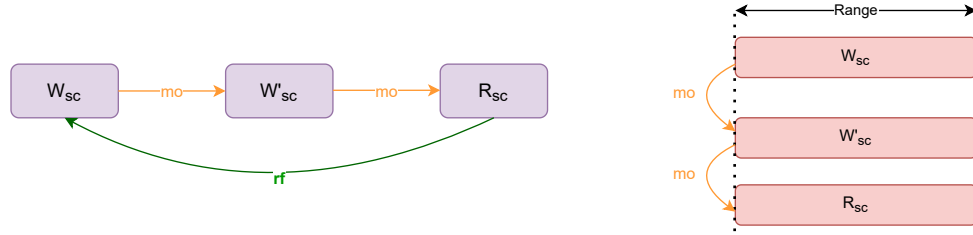


Figure 4.18: A read value cannot come from a write, if there exists a write memory ordered between them and all 3 events are sequentially consistent with equal ranges.

- If both writes are of type  $sc$  having equal ranges and the read is bound to happen after them, then  $e$  cannot read from  $d$

$$d:sc \wedge g:sc \wedge (\mathfrak{R}(d)=\mathfrak{R}(g)) \wedge d \xrightarrow{hb} e \wedge g \xrightarrow{hb} e \Rightarrow \neg e \xrightarrow{rf} d.$$

The figure below depicts pictorially the pattern that is not allowed by this rule.

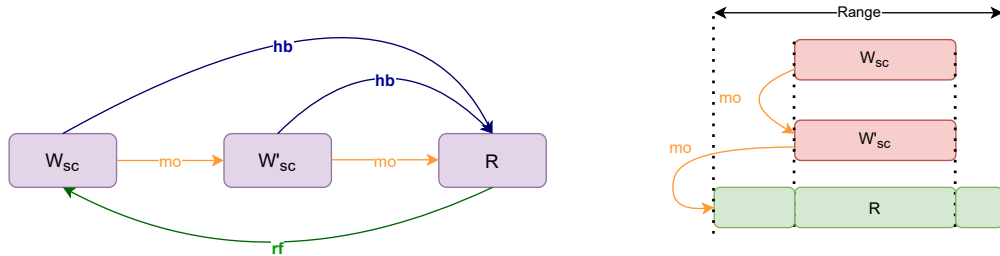


Figure 4.19: A read value cannot come from a write, if there exists a write memory ordered between them and both writes are sequentially consistent with equal ranges.

- If  $g$  and  $e$  are sequentially consistent, having equal ranges, and  $d$  is bound to happen before them, then  $e$  cannot read from  $d$

$$g:sc \wedge e:sc \wedge (\mathfrak{R}(g)=\mathfrak{R}(e)) \wedge d \xrightarrow{hb} g \wedge d \xrightarrow{hb} e \Rightarrow \neg e \xrightarrow{rf} d.$$

The figure below depicts pictorially the pattern that is not allowed by this rule.

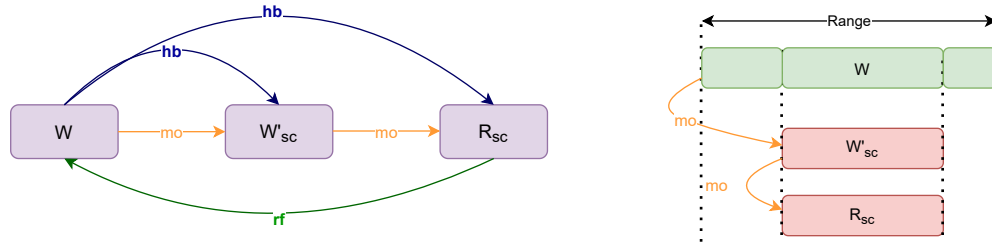


Figure 4.20: A read value cannot come from a write, if there exists a write memory ordered between them and both this write and the read are sequentially consistent with equal ranges.

## 4.7 Race

### 4.7.1 Race Condition $RC$

We define  **$RC$**  as the set of all pairs of events that are in a race. Two events  $e$  and  $d$  are in a race condition when they are shared memory events:

$$(e \in SM) \wedge (d \in SM).$$

having overlapping ranges, not having a  $\xrightarrow{hb}$  relation with each other, and which are either two writes or the two events are involved in a  $\xrightarrow{rf}$  relation with each other. This can be stated concisely as,

$$\neg (e \xrightarrow{hb} d) \wedge \neg (d \xrightarrow{hb} e) \wedge ((e, d \in W \wedge (\mathfrak{R}(d) \cap_{\mathfrak{R}} \mathfrak{R}(e) \neq \emptyset)) \vee (d \xrightarrow{rf} e) \vee (e \xrightarrow{rf} d)).$$

### 4.7.2 Data Race $DR$

We define  **$DR$**  as the set of all pairs of events that are in a data-race. Two events are in a data race when they are already in a race condition and when the two events are not both of type  $sc$ , or they have overlapping ranges. This is concisely stated as:

$$e, d \in RC \wedge ((\neg e:sc \vee \neg d:sc) \vee (\mathfrak{R}(e) \cap_{\mathfrak{R}} \mathfrak{R}(d) \neq \mathfrak{R}(e) \cup_{\mathfrak{R}} \mathfrak{R}(d)))$$

The definition for data race also implies that sequentially consistent events with overlapping ranges are also in a data race. This may be counter-intuitive in the sense that all agents observe the same order in which these events happen.

**Data-Race-Free (DRF) Programs** An execution is considered data-race free if none of the above conditions for data-races occur among events. A program is data-race free if all its executions are data race free. *The memory model guarantees Sequential Consistency for all data-race free programs (SC-DRF).*

## 4.8 Consistent Executions (Valid Observables)

Consistent executions are those which should ideally be possible if the program is actually run on some hardware. For a sequential program, we use the semantics of the programming language to understand what can be the outcome of a program. For a concurrent program, since we can have multiple outcomes of the same program being executed (keeping all inputs constant), we need a semantic model to rely on. The memory model is in essence just this semantic model for programs using shared memory.

In our language, a consistent execution maps to a valid observable behavior, as this is what the user can actually record as an outcome of the program.

As per the standard specification, valid observable behaviour is when<sup>10</sup>:

1. No  $\overrightarrow{rf}$  relation violates the above memory consistency rules.
2.  $\overrightarrow{hb}$  is a strict partial order.

*The memory model guarantees that every program must have at least one valid observable behaviour.*

---

As a summary, this chapter axiomatically defined the ECMAScript memory model. The model is defined using binary relations on events and specifying the constraints of the model in terms of restricting reads-from relations given other binary relations

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<sup>10</sup>There is also some conditions on host-specific events (which we mentioned is not of our main concern) and what is called a chosen read, which is nothing but the reads that the underlying hardware memory model allows. Since we are not concerned with the memory models of different hardware, this restriction on reads is not of our concern.



that can exist between events in a Candidate Execution. In the next chapter, we use this formal model to reason about the validity of instruction reordering under the constraints of the model.

## Chapter 5

# Instruction Reordering

### 5.1 Introduction

Instruction reordering is a common operation in compiler optimization, essential to instruction scheduling of course, but also implicit in loop invariant removal, partial redundancy elimination, and other optimizations that may move instructions. However, whether we can do such reordering freely given a concurrent program using relaxed memory accesses is a bit unclear.

#### **Simple reordering is not straightforward under shared memory semantics**

The main reason is that memory accesses here, do not just perform the desired operation (i.e Read / Write) but also imply certain visibility guarantees across all the other threads. In our observation, we find that, the relaxed memory model of Javascript prescribe semantics for visibility using the  $\overrightarrow{hb}$  relations.

Show an example or multiple examples here that enforces visibility due to having sequentially consistent events involved in a Candidate Execution.

**What can be done?** An example-based analysis exposes to us the problems that might exist when we perform such reordering of events. However, such an analysis, though would work for small programs to identify the possible conditions under which reordering can be done, become infeasible as the programs scale in length and complexity. This is because of the exponential increase in possible executions as the number of threads and program size in general increase. Hence, generalizations by using a small sample size is not something we can afford especially when we want to ensure these program transformations are done by the compiler in contrast to being done manually.

**Our approach** Our solution to this is to construct a proof on Candidate Executions of the original program and the transformed one which exposes the possible observable behaviors it can have. The crux of the proof is to guarantee that reordering does not bring any new  $\xrightarrow{rf}$  (reads-from) relations that did not exist in any Observable Behavior of the original Candidate Execution. It is important to note however, that a proof in this sense would be generalized to any Candidate and is thus conservative. So, it might be the case that for specific programs, reordering can be valid, however, in a general sense may not be valid for others.

**Assumption** We make the following assumptions for every program we consider :

1. All events are tear-free
2. No synchronize events exist
3. No Read-Modify-Write events exist
4. All executions of the candidate before reordering have happens-before as a strict partial order

We first consider when consecutive events in the same agent can be reordered, followed by non-consecutive cases. The crux of the proof is to guarantee that reordering does not bring any new reads-from relations that did not result due to any execution of the original program.

The following definitions and lemmas are not particular to instruction reordering, so I think we can make it a point to put this in a section that introduces our work on optimizations.

### 5.1.1 Preliminaries

Before we go about proving when reordering is valid, we would like to have two additional definitions which would prove useful.

**Definition 5.1.1.** *Consecutive pair of events (cons)* We define cons as a function, which takes two events as input, and gives us a boolean indicating if they are consecutive pairs. Two events  $e$  and  $d$  are consecutive if they have an  $\xrightarrow{ao}$  relation among them and are next to each other, which can be defined formally as

$$(e \xrightarrow{ao} d \wedge \nexists k \text{ s.t. } e \xrightarrow{ao} k \wedge k \xrightarrow{ao} d) \vee (d \xrightarrow{ao} e \wedge \nexists k \text{ s.t. } d \xrightarrow{ao} k \wedge k \xrightarrow{ao} e)$$

**Definition 5.1.2.** *Direct happens-before relation (dir)* We define  $\text{dir}$  to take an ordered pair of events  $(e, d)$  such that  $e \xrightarrow{hb} d$  and gives a boolean value to indicate whether this relation is direct, i.e those relations that are not derived through transitive property of  $\xrightarrow{hb}$ .

We can infer certain things using this function based on some information on events  $e$  and  $d$ .

- If  $e:uo$ , then  $\text{dir}(e, d) \Rightarrow \text{cons}(e, d)$
- If  $d:uo$ , then  $\text{dir}(e, d) \Rightarrow \text{cons}(e, d)$
- If  $e:sc \wedge e \in R$ , then  $\text{dir}(e, d) \Rightarrow \text{cons}(e, d)$
- If  $e:sc \wedge e \in W$ , then  $\text{dir}(e, d) \Rightarrow \text{cons}(e, d) \vee e \xrightarrow{sw} d$
- If  $d:sc \wedge d \in W$ , then  $\text{dir}(e, d) \Rightarrow \text{cons}(e, d)$
- If  $d:sc \wedge e \in R$ , then  $\text{dir}(e, d) \Rightarrow \text{cons}(e, d) \vee e \xrightarrow{sw} d$

### 5.1.2 Lemmas to assist our proof

In order to assist our proof, we define two *lemmas* based on the ordering relations.

**Lemma 1.** *Consider three events  $e, d$  and  $k$ .*

*If*

$$\text{cons}(e, d) \wedge e \xrightarrow{ao} d \wedge ((d:uo) \vee (d:sc \wedge d \in W))$$

*then,*

$$k \xrightarrow{hb} d \implies k \xrightarrow{hb} e$$

When we have two consecutive events  $e$  and  $d$  which are one after the other (i.e.  $e \xrightarrow{ao} d$ ), we can use transitive property of  $\xrightarrow{hb}$  to infer that any event  $k$  that happens before  $e$ , also happens before  $d$ . However, is it possible to derive that the event  $k$  happens before  $e$  using the evidence that  $k$  happens before  $d$ ? This lemma states the condition when this is true.

*Proof.* We will divide the proof for this into two cases, based on what event  $d$  is. For both cases, we have the following to be true :

$$\text{cons}(e, d) \wedge e \xrightarrow{ao} d \tag{0}$$

In the first case, we have  $d:uo$ . From Def 5.1.2 and Def 5.1.1, we have for any event  $k$

$$dir(k, d) \Rightarrow cons(k, d)$$

From (0), we have  $k = e$  satisfying the above property with  $d$ . Because  $\overrightarrow{ao}$  is a total order,  $e$  will be the only event. Thus, for any other  $k \neq e$ , we have

$$k \xrightarrow{hb} d \Rightarrow k \xrightarrow{hb} d$$

The following figure should explain this intuition:

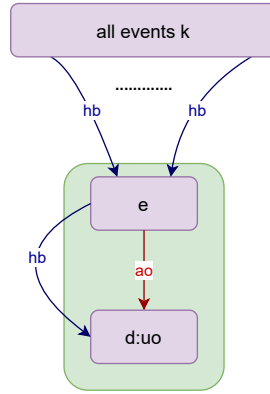


Figure 5.1: For the first case

In the second case, we have

$$d:sc \wedge d \in W \tag{4}$$

Thus, from Def 5.1.2 and Def 5.1.1, for any event  $k$ , we have

$$dir(k, d) \Rightarrow cons(k, d)$$

From (4), event  $e$  satisfies the above. Though there could be direct *happens-before* relation with some event  $k$  from another *agent*, from Def 5.1.2 these are only relations satisfying  $dir(d, k)$ . Thus, we can once again infer that for any  $k \neq e$

$$k \xrightarrow{hb} d \Rightarrow k \xrightarrow{hb} d$$

The following figure explains this intuition:

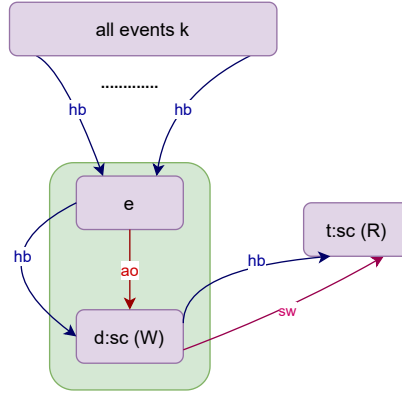


Figure 5.2: For the second case

□

**Lemma 2.** Consider three events  $e$ ,  $d$  and  $k$

If

$$\text{cons}(e, d) \wedge e \xrightarrow{ao} d \wedge ((e:uo) \vee (e:sc \wedge e \in R))$$

then,

$$e \xrightarrow{hb} k \implies d \xrightarrow{hb} k$$

When we have two consecutive events  $e$  and  $d$  which are one after the other (i.e.  $e \xrightarrow{ao} d$ ), we can use transitive property of  $\xrightarrow{hb}$  to infer that any event  $k$  that happens after  $d$ , also happens after  $e$ . However, is it possible to derive that the event  $k$  happens after  $d$  using the evidence that  $k$  happens after  $e$ ? This lemma states the condition when this is true.

*Proof.* Just like the proof for the previous lemma, we will divide the proof for this into two cases, based on what event  $e$  is. Again, for both cases, we have the following to be true:

$$\text{cons}(e, d) \wedge e \xrightarrow{ao} d \tag{0}$$

In the first case, we have  $e:uo$ . From Def 5.1.2 and Def 5.1.1, we have for any event  $k$

$$\text{dir}(e, k) \implies \text{cons}(e, k)$$

From (0), we have  $k = d$  satisfying the above property with  $e$ . Because  $\overrightarrow{ao}$  is a total order,  $d$  would be the only such event. Thus, for any other event  $k \neq d$ , we can infer,

$$e \xrightarrow{hb} k \Rightarrow d \xrightarrow{hb} k$$

The following figure should explain this intuition:

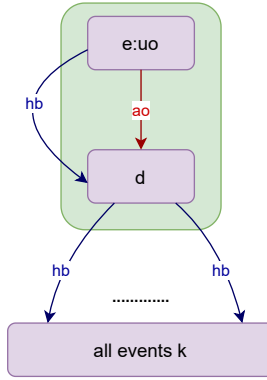


Figure 5.3: Caption

In the second case, we have

$$e : sc \wedge e \in R \tag{4}$$

Thus, from Def 5.1.2 and Def 5.1.1, for any event  $k$ , we have

$$dir(e, k) \Rightarrow cons(e, k)$$

From (4), we have event  $k = d$  satisfying the above property with  $e$ . Though there could be direct *happens-before* relation with some event  $k$  from another *agent*, from Def 5.1.2 these are only relations satisfying  $dir(k, e)$ . Thus, we can infer that for any  $k \neq d$

$$e \xrightarrow{hb} k \Rightarrow d \xrightarrow{hb} k$$

The following figure explains this intuition:

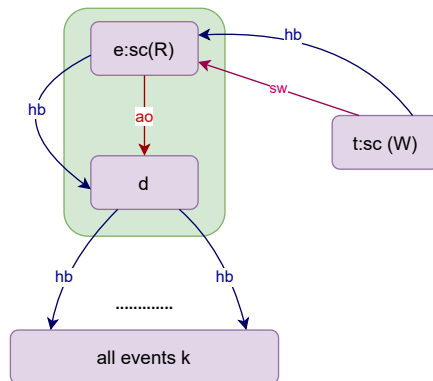


Figure 5.4: Caption

□

## 5.2 Valid reordering

We view reordering as manipulating the agent-order relation. In that sense, reordering two consecutive events  $e$  and  $d$  such that  $e \xrightarrow{ao} d$  becomes:

$$e \xrightarrow{ao} d \mapsto d \xrightarrow{ao} e$$

What implications this change has on the other ordering relations depends on the type of events  $e$  and  $d$  are and would require an analysis on each Candidate Execution. The intuition is that the axioms of the memory model rely on certain ordering relations to restrict observable behaviors in a program. Hence, preserving these ordering relations would help us in turn not introduce new Observable Behaviors. In particular we note that preserving  $\xrightarrow{hb}$  relations (other than the one we eliminate intentionally i.e  $e \xrightarrow{hb} d$ ) would suffice for our purpose. Since  $\xrightarrow{mo}$  respects  $\xrightarrow{hb}$ , we in turn even preserve the memory order which is essential.

Our main objective is to ensure that the set of possible observable behaviors of a program, remain unchanged after reordering. If that is not feasible, then we would want the set of observable behaviors after reordering at the very least to be a subset. This is to ensure that the Candidate of a program does not have some new behaviours that weren't supposed to happen prior to reordering.

We begin by first defining a reorderable pair of events. We then formulate a theorem on the set of observable behaviors of a Candidate before and after reordering a pair of consecutive events which are reorderable. We consider reordering valid if the set of observable behaviours after reordering are a subset of the original.



**Definition 5.2.1.** *Reorderable Pair (Reord)*

We define a boolean function *Reord* that takes two ordered pair of events  $e$  and  $d$  such that  $e \xrightarrow{ao} d$  and gives a boolean value indicating if they are a reorderable pair.

$$\begin{aligned}
\text{Reord}(e, d) = & \\
& (((e:uo \wedge d:uo) \wedge ((e \in R \wedge d \in R) \vee (\mathfrak{R}(e) \cap_{\mathfrak{R}} \mathfrak{R}(d) = \phi))) \\
& \vee \\
& ((e:sc \wedge d:uo) \wedge ((e \in W \wedge (\mathfrak{R}(e) \cap_{\mathfrak{R}} \mathfrak{R}(d) = \phi)))) \\
& \vee \\
& ((e:uo \wedge d:sc) \wedge ((d \in R \wedge (\mathfrak{R}(e) \cap_{\mathfrak{R}} \mathfrak{R}(d) = \phi)))))
\end{aligned}$$

**Theorem 5.2.1.** *Consider a candidate  $C$  of a program and its possible Candidate Executions where  $\xrightarrow{hb}$  is strictly partial order. Consider two events  $e$  and  $d$  such that  $\text{cons}(e, d)$  is true in  $C$  and  $e \xrightarrow{ao} d$ . Consider another candidate  $C'$  resulting after reordering  $e$  and  $d$ . Then if  $\text{Reord}(e, d)$  is true in  $C$ , the set observable behaviors possible due to Candidate Executions of  $C'$  is a subset of that of  $C$ .*

*Proof.* We look at this in terms of performing an instruction reordering on a candidate execution of  $C$ . We would want the resulting candidate execution to preserve all the other  $\xrightarrow{hb}$  relations (except  $e \xrightarrow{hb} d$ ) and that any new  $\xrightarrow{hb}$  relations strictly reduce possible observable behaviors.

The proof is structured as follows. We first show that existing *happens-before* relations in any candidate execution of  $C$  except  $e \xrightarrow{hb} d$  remain intact after reordering. We then identify the cases where new *happens-before* relations could be established. We identify from these cases whether *happens-before* cycles could be introduced. We then show for the remaining cases that new relations do not introduce any new observable behaviors.

The above steps can be summarized as addressing four main questions for any *Candidate Execution* of  $C'$

1. Apart from  $e \xrightarrow{hb} d$ , do other *happens-before* relations remain intact?
2. Apart from  $d \xrightarrow{hb} e$ , are any new *happens-before* relations established?
3. Are any *happens-before* cycles introduced?
4. Do the new relations bring new *observable behaviors*?

**1. Preserving *happens-before* relations** If  $\overrightarrow{hb}$  relations among events are lost after reordering, we may introduce new observable behaviors. The relations that are subject to change can be divided into four parts using events  $e$  and  $d$

a)  $k \xrightarrow{hb} e$

b)  $e \xrightarrow{hb} k$

c)  $d \xrightarrow{hb} k$

d)  $k \xrightarrow{hb} d$

Firstly, note that the relations of the form  $e \xrightarrow{hb} k$  come through either a  $\overrightarrow{sw}$  relation with  $e$  or relations through event  $d$ , i.e. of the form  $d \xrightarrow{hb} k$ . The ones that come due to the latter, may not be preserved after reordering, if we strictly are only able to derive them with relations through  $d$ . Note also that, a similar argument exists for relations of the form  $k \xrightarrow{hb} d$  wherein relations derived through  $e$  ( $k \xrightarrow{hb} e$ ) may be lost after reordering.

Hence, the relations that could be subject to change can be addressed by considering two disjoint sets of events in any *Candidate Execution* of  $C$  as below.

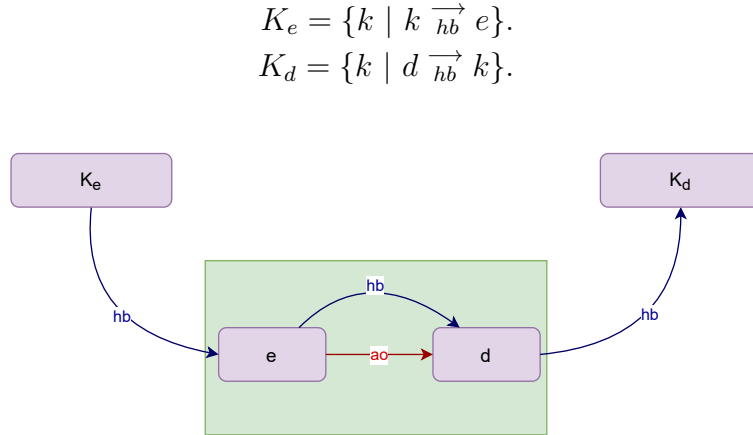


Figure 5.5: For any Candidate Execution of  $C$ , the set  $K_e$  and  $K_d$

Consider two events  $p1 \in K_e$  and  $p2 \in K_d$  (When  $e$  is the first event or  $d$  is the last event, assume dummy events that can act as  $p1$  or  $p2$ .) belonging to the same agent as that of  $e$  and  $d$  such that in  $C$ :

$$dir(p1, e) \wedge dir(d, p2).$$

Note that in terms of direct happens-before relations, on reordering, any *CandidateExecution* of  $C$  will have the following changes

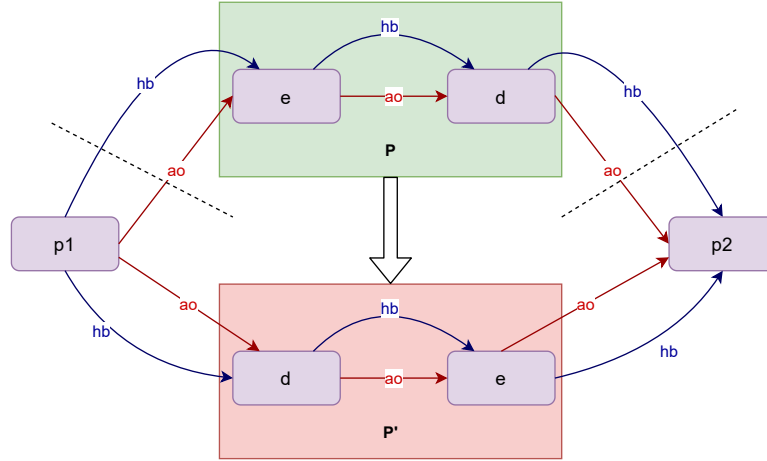


Figure 5.6: The direct relation changes that can be observed while reordering events  $e$  and  $d$

The figure above is to show that, for any *CandidateExecution* of  $C$ , the following is true

$$\text{cons}(p1, e) \wedge \text{dir}(p1, e) \wedge \text{dir}(e, d) \wedge \text{cons}(d, p2) \wedge \text{dir}(d, p2).$$

and for that of  $C'$ ,

$$\text{cons}(p1, d) \wedge \text{dir}(p1, d) \wedge \text{dir}(d, e) \wedge \text{cons}(e, p2) \wedge \text{dir}(e, p2).$$

We need the following key relations to be preserved in Candidate executions of  $C'$

- |                            |                           |
|----------------------------|---------------------------|
| a) $p1 \xrightarrow{hb} e$ | b) $e \xrightarrow{hb} k$ |
| c) $d \xrightarrow{hb} p2$ | d) $k \xrightarrow{hb} d$ |

After reordering, we have (a) and (c) preserved due to transitivity

$$\begin{aligned} p1 \xrightarrow{hb} d \wedge d \xrightarrow{hb} e &\Rightarrow p1 \xrightarrow{hb} e. \\ e \xrightarrow{hb} p2 \wedge d \xrightarrow{hb} e &\Rightarrow d \xrightarrow{hb} p2. \\ p1 \xrightarrow{hb} d \wedge d \xrightarrow{hb} e \wedge e \xrightarrow{hb} p2 &\Rightarrow p1 \xrightarrow{hb} p2. \end{aligned}$$

(b) and (d) may not be preserved due to  $d \xrightarrow{sw} k$  or  $k \xrightarrow{sw} d$ . If we can "pivot" the set  $K_e$  to  $p1$  and  $K_d$  to  $p2$ , it would ensure that our other two intended relations also

remain preserved after reordering by transitivity. To state formally, we have a valid pair of pivots  $\langle p1, p2 \rangle$  when the following two conditions hold

$$\begin{aligned} \forall k \in K_e - \{p1\}, k \xrightarrow{hb} p1. \\ \forall k \in K_d - \{p2\}, p2 \xrightarrow{hb} k. \end{aligned}$$

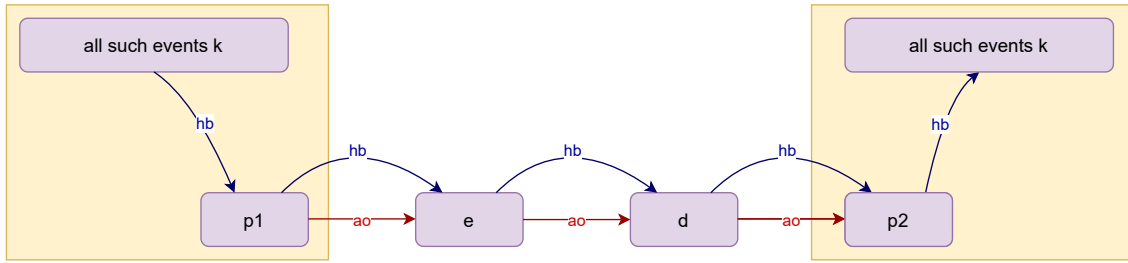


Figure 5.7: For any Candidate execution, the intuition behind valid pivots  $\langle p1, p2 \rangle$

By Lemma 1 and 2 respectively, we have for  $C$ , the following condition where  $\langle p1, p2 \rangle$  is a valid pivot pair

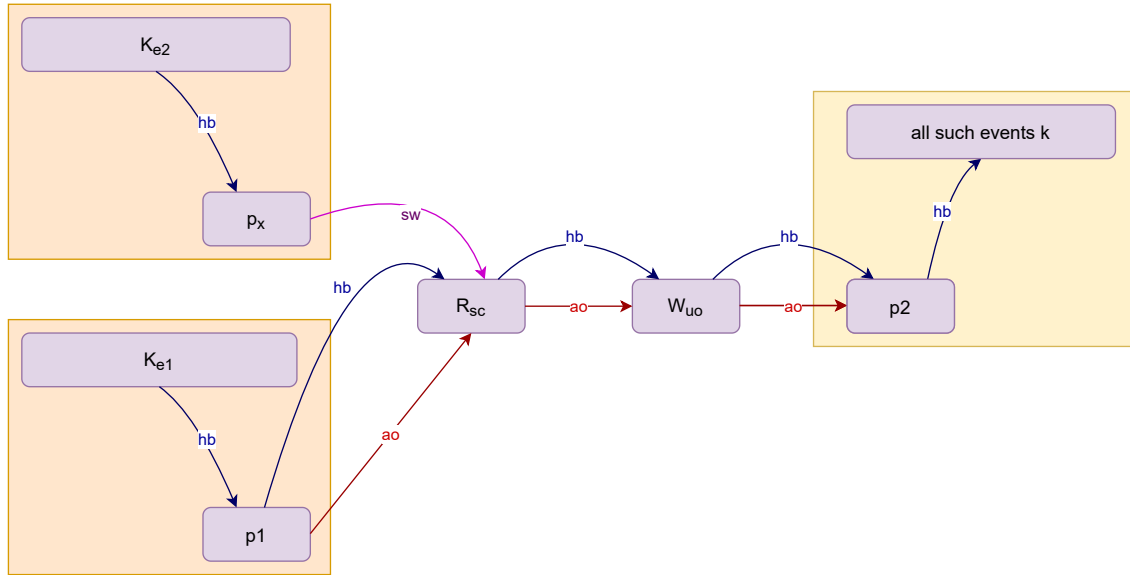
$$\begin{aligned} e : uo \vee (e : sc \wedge e \in W). \\ d : uo \vee (d : sc \wedge d \in R). \end{aligned}$$

The following table summarizes the cases where we have a valid pair of pivots<sup>1</sup>  $\langle p1, p2 \rangle$

$\langle p1, p2 \rangle$	R-R	R-W	W-R	W-W
uo-uo	Y	Y	Y	Y
uo-sc	Y	N	Y	N
sc-uo	N	N	Y	Y
sc-sc	N	N	Y	N

Figure 5.8: Table summarizing whether we have valid pair of pivots based on  $e$  and  $d$

We show a simple example where we do not have a valid pair of pivots, particularly because  $p1$  is not a valid pivot. Note that in this example,  $K_e = K_{e1} + K_{e2} + p1 + p_x$

Figure 5.9: A Candidate Execution where  $p1$  is not a valid pivot

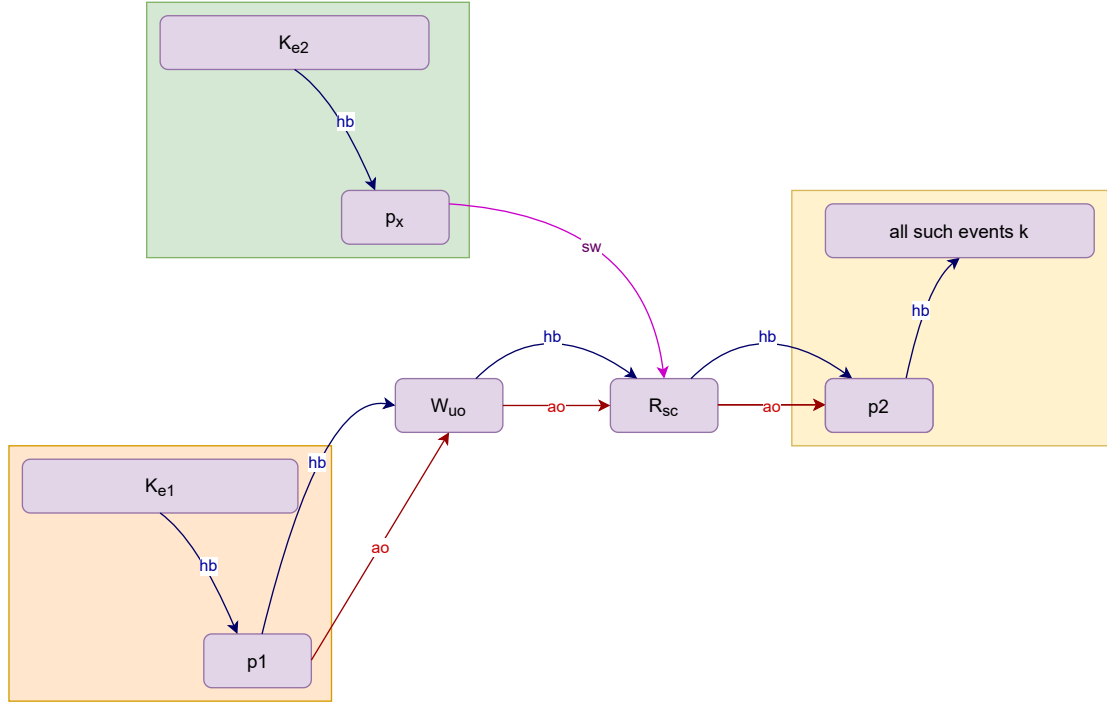


Figure 5.10: The resultant Candidate Execution after reordering, exposing the relations with  $p_x$ ,  $K_{e2}$  and  $d$  that arelost

**2. Additional *happens-before* relations** Although we have identified the cases when *happens-before* relations are preserved, we also get some additional relations in some of them.

As an example, for the case when  $d$  is a sequentially consistent read, by Lemma 1, in any execution of  $C$

$$k \xrightarrow{hb} d \not\Rightarrow k \xrightarrow{hb} e$$

But in *Executions* of candidate  $C'$ , by transitivity, we have

$$k \xrightarrow{hb} d \Rightarrow k \xrightarrow{hb} e$$

This is because, there are *happens-before* relations that come through *synchronize-with* relations with  $d$ . Thus, although we are able to preserve relations that existed

<sup>1</sup>This proof does not go about showing the exact happens-before relations that are preserved; rather it uses the properties between different happens before relations that hold, which would imply that for any possible Candidate Execution after reordering, the set of happens-before relations apart from that between  $e$  and  $d$  remain the same.

in any *CandidateExecution* of  $C$ , we also in the process, introduce new ones in *CandidateExecutions* of  $C'$ . The figure below shows pictorially an example of a *Candidate Execution* of  $C$  for the case above

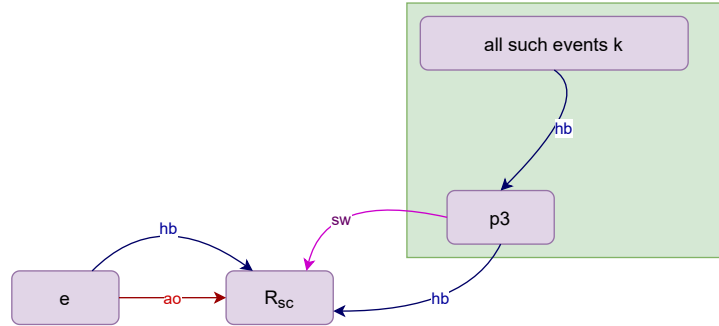


Figure 5.11: A *Candidate Execution* where  $d$  is a sequentially consistent read

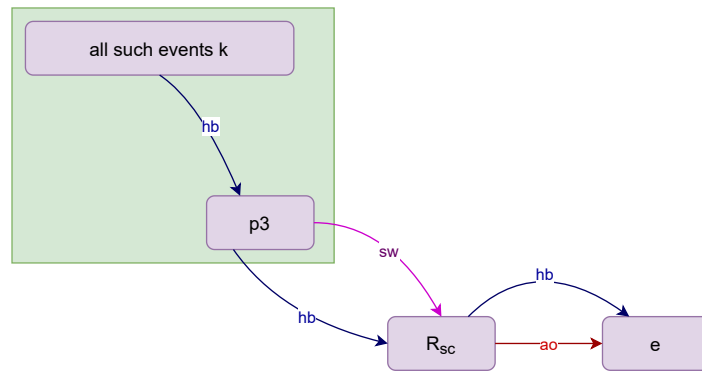


Figure 5.12: The *Candidate Execution* after reordering, exposing the new relations established with  $e$ ,  $p3$  and set  $k$

Explain the above figures or perhaps highlight the new relations that are established.

To summarize, the table below shows the cases where new relations could be introduced.

New ReIn	R-R	R-W	W-R	W-W
uo-uo	N	N	N	N
uo-sc	Y	N	Y	N
sc-uo	N	N	Y	Y
sc-sc	N	N	Y	N

Figure 5.13: Table summarizing when new *happens-before* relations could be introduced based on having valid pair of pivots

For these cases, we must know whether the new relations introduce new observable behaviors.

**3. Presence of cycles?** Before we go into analyzing whether new relations introduce observable behaviours, we first ensure there are no  $\overrightarrow{hb}$  cycles introduced in the process. Consider the example below

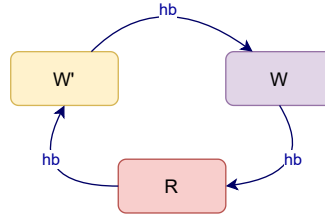


Figure 5.14: Caption

Notice that here, Axiom 1 restricts  $R$  to read from  $W'$ .

$$R \overrightarrow{hb} W' \Rightarrow \neg R \overrightarrow{rf} W'$$

But by transitive property, it is also the case that  $W' \overrightarrow{hb} R$ .

$$W' \overrightarrow{hb} W \wedge W \overrightarrow{hb} R \Rightarrow W' \overrightarrow{hb} R$$

As per this, Axiom 1 will not restrict  $R \overrightarrow{rf} W'$ . To avoid such cases, we will need to ensure that no Candidate Execution of  $C'$  after  $e$  and  $d$  are reordered have  $\overrightarrow{hb}$  cycles.

Note that if a cycle exists after reordering, then

1. The relations preserved do not themselves create a cycle (ref to the theorem)
2. Additional new relations may introduce cycles



The first part is straightforward as we assume we can only do reordering on Candidate Exectuions of  $C$  not having happens-before cycles.

For the second part, we first address the cases where  $d \xrightarrow{hb} e$  may be part of the cycle. The other event  $k$ , may be either from the set  $K_e$ ,  $K_d$  or a new relation that is formed<sup>2</sup>.

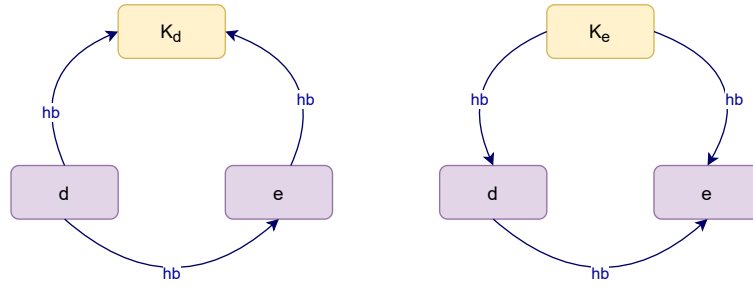


Figure 5.15: If  $k$  belongs to one of the sets  $K_e$  or  $K_d$

The above figure shows that  $k$  cannot belong to either of the sets, as their relations with  $e$  and  $d$  will not result in a cycle.

For cases where  $k \xrightarrow{hb} e$  is the set of new relations, note that by lemma 1

$$k \xrightarrow{hb} e \Rightarrow k \xrightarrow{hb} d$$

For cases where  $d \xrightarrow{hb} k$  is the set of new relations, by lemma 2

$$d \xrightarrow{hb} k \Rightarrow e \xrightarrow{hb} k$$

So for both these cases also, a cycle with  $d \xrightarrow{hb} e$  cannot exist. The following figure shows pictorially this fact.

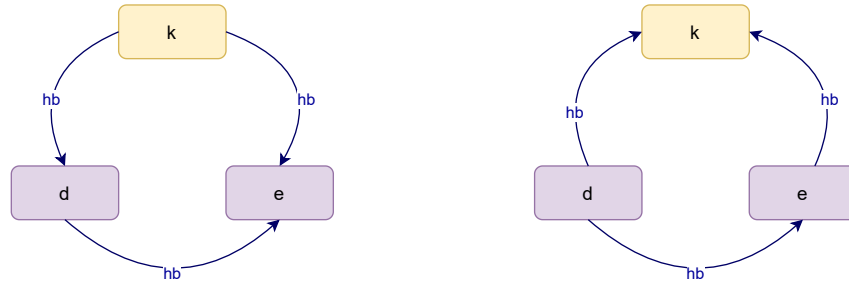


Figure 5.16: If  $k \xrightarrow{hb} e$  or  $d \xrightarrow{hb} k$  are new sets of relations

<sup>2</sup> $K_e$  and  $K_d$  only apart from the new relation because these are the only valid cases where happens-before relations are preserved after reordering. So we need not consider cases such that  $e \xrightarrow{hb} k$  or  $k \xrightarrow{hb} d$  as the old relations they are covered by  $K_e$  and  $K_d$ .

For the one case where we have two new sets of relations formed, i.e  $d \xrightarrow{hb} k$  and  $k \xrightarrow{hb} e$ , we could have a case where  $k$  is a common event for both sets. But, by Lemma 1, we also have  $k \xrightarrow{hb} d$  and by Lemma 2,  $e \xrightarrow{hb} k$ <sup>3</sup>. Thus, we have a cycle. The following figure shows this pictorially

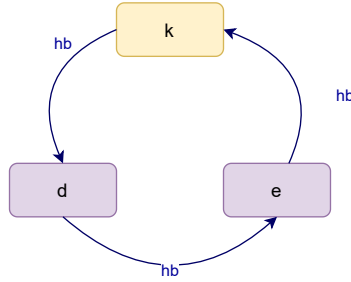


Figure 5.17: A cycle exists in the case where we have two new sets of relations ( $k \xrightarrow{hb} e$  and  $d \xrightarrow{hb} k$ )

Maybe have a better figure, meaning a set of relations where each figure shows clearly which relation is implied due to which lemma

Now for the case when  $d \xrightarrow{hb} e$  may not be part of the cycle, we have only two other new relations,  $k \xrightarrow{hb} e$  or  $d \xrightarrow{hb} k$ .

Considering the first scenario where the new set of relations are of the form  $k \xrightarrow{hb} e$ . Suppose a cycle exists with another event  $k'$ . Then

$$k \xrightarrow{hb} e \wedge e \xrightarrow{hb} k' \wedge k' \xrightarrow{hb} k$$

Note that the latter two relations are not new, since the only new set of relations are of the first form. Now, by Lemma 1 and by transitivity respectively

$$\begin{aligned} k \xrightarrow{hb} e &\Rightarrow k \xrightarrow{hb} d \\ e \xrightarrow{hb} k' &\Rightarrow d \xrightarrow{hb} k' \end{aligned}$$

So, the following is also a cycle

$$k \xrightarrow{hb} d \wedge d \xrightarrow{hb} k' \wedge k' \xrightarrow{hb} k$$

But these relations already exist in the original Candidate Execution, which implies a cycle existed before reordering. This contradicts our assumption that we only

---

<sup>3</sup>It is not actually due to lemmas, but just that the new relations were derived through  $e$  or  $d$ , as these relations existed before reordering.

reorder when the Candidate Executions of  $C$  have no cycles. Thus, by contradiction such a cycle cannot exist.

In similar lines for the cases where the set of new relations are of the form  $d \xrightarrow{hb} k$ , Suppose a cycle exists with another event  $k'$ . Then

$$d \xrightarrow{hb} k \wedge k \xrightarrow{hb} k' \wedge k' \xrightarrow{hb} d$$

Note that the latter two relations are not new, since the only new set of relations are of the first form. Now, by Lemma 2 and by transitivity respectively, we have

$$d \xrightarrow{hb} k \Rightarrow e \xrightarrow{hb} k \xrightarrow{hb} k' \xrightarrow{hb} d \Rightarrow k' \xrightarrow{hb} d$$

Thus, we also have the cycle

$$e \xrightarrow{hb} k \wedge k \xrightarrow{hb} k' \wedge k' \xrightarrow{hb} e$$

But these relations already exist in the original Candidate Execution, which implies a cycle existed before reordering. This contradicts our assumption that we only reorder when the Candidate Executions of  $C$  have no cycles. Thus, by contradiction such a cycle cannot exist.

To summarize, the table below shows the cases where new relations have no happens-before cycles.

New ReIn	R-R	R-W	W-R	W-W
uo-uo	NO	NO	NO	NO
uo-sc	YES	NO	YES	NO
sc-uo	NO	NO	YES	YES
sc-sc	NO	NO	YES	NO

Table 5.1: Insert good caption here.

**4. Do new relations introduce new observable behaviors?** In any candidate execution, reordering events  $e$  and  $d$  eliminates the relation  $e \xrightarrow{hb} d$  and introduces the new relation  $d \xrightarrow{hb} e$ . New behaviours created by the latter directly, if any, are of course intentional (and should normally be avoided by ensuring  $e$  and  $d$  are independent), but we need to ensure that this does not also result in new behaviours indirectly.

Let us first consider the variants of events  $e$  and  $d$  that we need to analyze:

- |   |   |
|---|---|
| a) $e \in R \wedge d \in R \wedge e:uo \wedge d:uo$ | b) $e \in R \wedge d \in R \wedge e:uo \wedge d:sc$ |
| c) $e \in R \wedge d \in W \wedge e:uo \wedge d:uo$ | d) $e \in W \wedge d \in R \wedge e:uo \wedge d:uo$ |
| e) $e \in W \wedge d \in R \wedge e:uo \wedge d:sc$ | f) $e \in W \wedge d \in R \wedge e:sc \wedge d:uo$ |
| g) $e \in W \wedge d \in W \wedge e:uo \wedge d:uo$ | h) $e \in W \wedge d \in W \wedge e:sc \wedge d:uo$ |

We analyze each of the above case one by one by first considering the original relation ( $e \xrightarrow{hb} d$ ) and the reordered one ( $d \xrightarrow{hb} e$ ).

- (a) and (b) do not fit any pattern of our Axioms, hence even after reordering the agent order between them does not match any other axiom. Hence this relation does not introduce any new observable behavior. This is irrespective of the range between the two read events.
- (c) fits in the pattern of Axiom 1, when they have at least overlapping ranges. Before reordering,  $d$  is not allowed to read from  $e$ , but after reordering, it can. Hence this relation can introduce observable behavior if the range between events  $e$  and  $d$  at least overlap.
- (d), (e) and (f) can fit in the pattern of Axiom 1 and 3, if  $e$  and  $d$  at least have overlapping ranges, preventing  $d$  from reading parts of  $e$  or some parts of another write  $k$  due to  $e$  being the intervening write. But after reordering,  $d$  is allowed to read parts of  $k$ , which introduces new observable behaviors.
- (g) and (h) can fit in the pattern of Axiom 1 and 3, if they have at least overlapping ranges. Before reordering, the agent order between  $e$  and  $d$  could prevent some read  $k$  from reading parts of  $e$ . This is not the case after reordering, thus introducing a new observable behavior.

In summary, on observing the role on the axioms on the relation between  $e$  and  $d$ , notice that if both  $e$  and  $d$  are read events then the range does not matter. For all other cases, if events  $e$  and  $d$  have at least overlapping ranges, one could introduce a new observable behavior after reordering them.

We will later show counter examples for each of the above cases that we discard as invalid to reorder. Decide whether to put figures here explaining each pattern that matches or palce counterexamples. The latter will be too much.

Any other new relations that are introduced can be divided into 4 cases, in terms of our events  $e$  and  $d$  and the new relation with some event  $k$ :

- a)  $e:uo \wedge e \in R \wedge k \xrightarrow{hb} e$ .      b)  $e:uo \wedge e \in W \wedge k \xrightarrow{hb} e$ .  
 c)  $d:uo \wedge d \in R \wedge d \xrightarrow{hb} k$ .      d)  $d:uo \wedge d \in W \wedge d \xrightarrow{hb} k$ .

Change the figure above to represent only the first four cases In each of the above cases, note firstly that we need to only consider cases where their ranges are overlapping/equal.

Figure below shows a breakdown of sub-cases for the first case (a), varying based on the nature of event  $k$ .

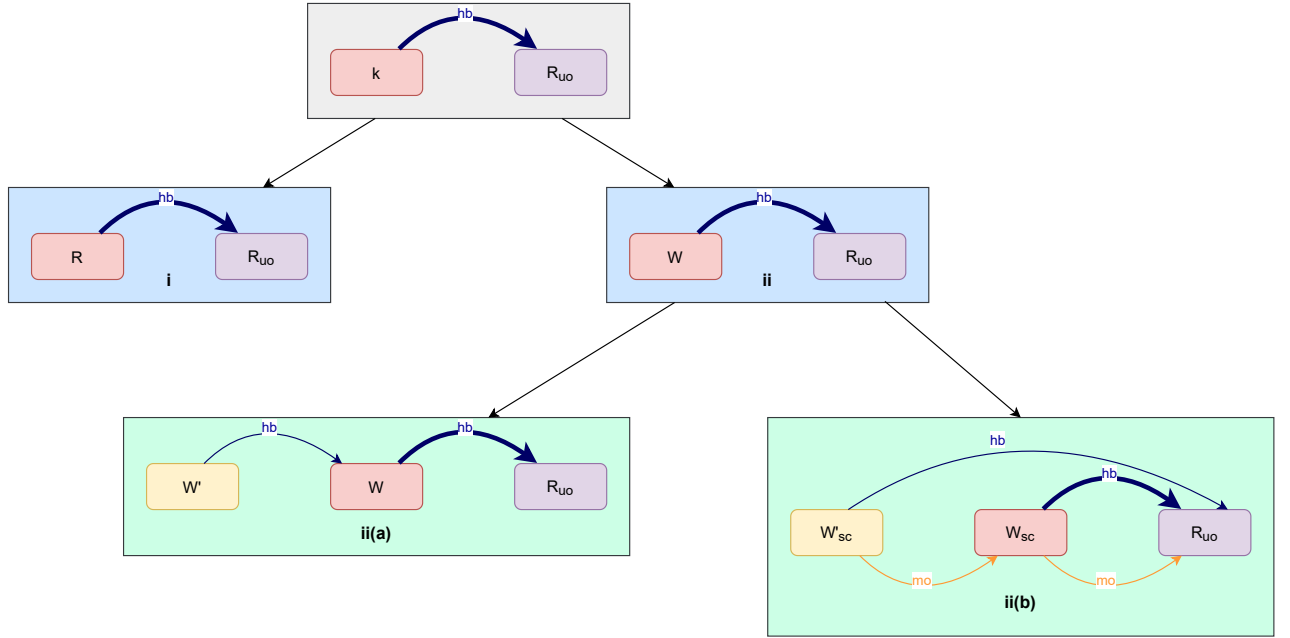


Figure 5.18: The role of the axioms on introducing a new relation between an unordered Read and some event  $k$

1. For (i), when  $k$  is a read, the pattern matches none of the Axioms.
2. For (ii), when  $k$  is a write, Axiom 1 (ii(a)) or Axiom 3 (ii(b)) could restrict the read ( $e$ ) from reading overlapping ranges of  $W'$  with  $W$ .

Figure below shows a breakdown of sub-cases for the case (b), varying based on the nature of event  $k$ .

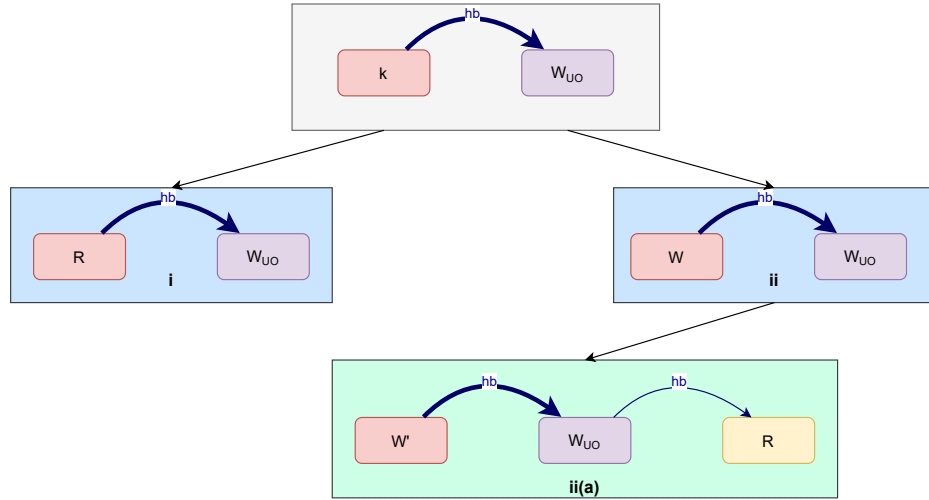


Figure 5.19: (i) and (ii(b)) satisfy the axiom of Coherent Reads

For case (b) we can observe the following from the above figure

1. For (i), when  $k$  is a read, Axiom 1 restricts  $k$  from reading from the write  $e$ .
2. For (ii), when  $k$  is a write, Axiom 1 restricts some read from reading parts of  $k$  due to the write  $e$ .

Figure below shows a breakdown of sub-cases for the first case (c), varying based on the nature of event  $k$ .

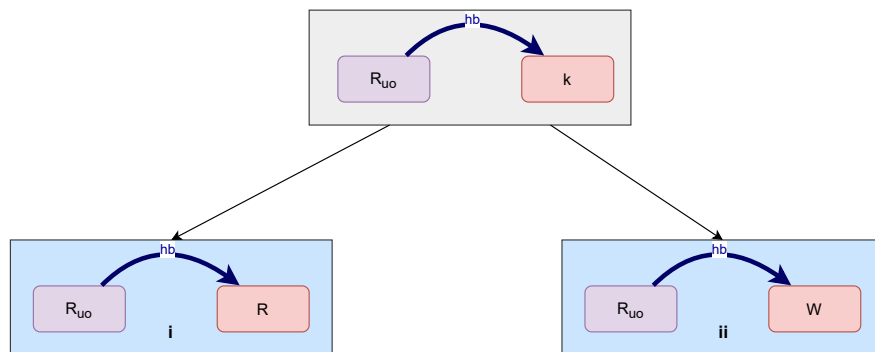


Figure 5.20: (ii) satisfies the axiom of Coherent Reads

For case (c), we can observe the following from the above figure

1. Case (i) does not correspond to any pattern restricted on the model, thus having no impact on the observable behaviors.
2. For (ii), when  $k$  is a write, Axiom 1 restricts the read  $d$  from reading values of write  $k$ .

Figure below shows a breakdown of sub-cases for the first case (d), varying based on the nature of event  $k$ .

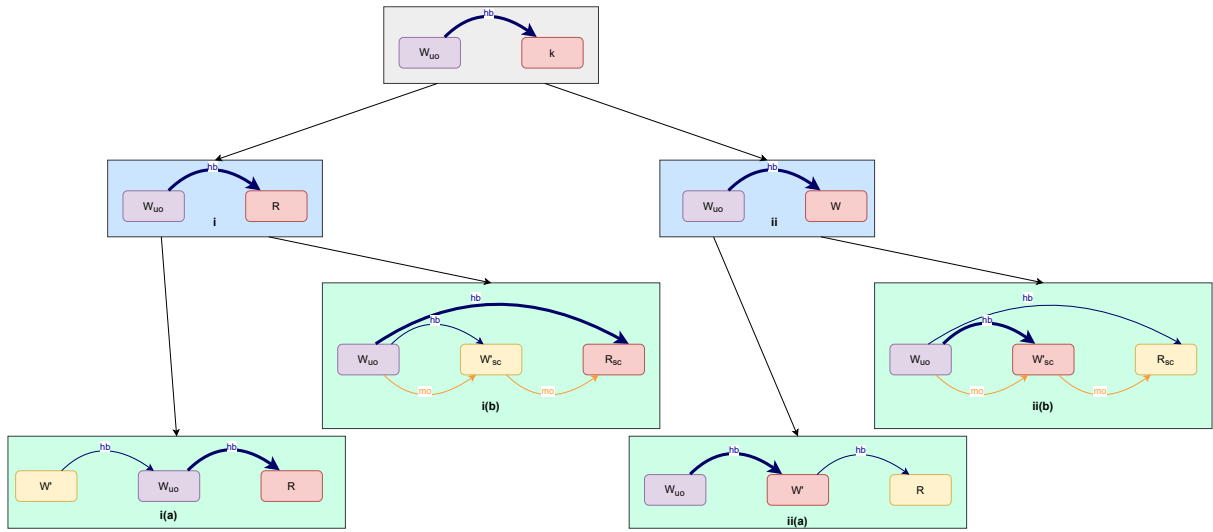


Figure 5.21: (i(a)), (ii(a)) satisfy the axiom of Coherent Reads, whereas (i(b)), (ii(b)) satisfy the axiom of SequentiallyConsistent Atomics

For case (d) we can observe the following

1. For case (i), Axiom 1 (i(a)) or Axiom 3 (i(b)) could restrict a read  $k$  from reading values of write  $d$ ,
2. For case (ii), Axiom 1 (ii(a)) or Axiom 3 (ii(b)) could restrict a read from reading values of write  $d$ ,

The above case wise analysis showed us that any new relation (apart from  $d \xrightarrow{hb} e$ ), matching the patterns of the axioms, only enables in restricting possible observable behaviors, which are  $\xrightarrow{rf}$  relations. Thus, we can infer that no new observable behavior is introduced due to the new set of  $\xrightarrow{hb}$  relations.

In summary, the table below summarizes the valid cases where, we have a pair of valid pivots, where new relations do not introduce new observable behaviors and do not have cycles.

Final	R-R	R-W	W-R	W-W
uo-uo	Y	Y	Y	Y
uo-sc	Y	N	Y	N
sc-uo	N	N	Y	Y
sc-sc	N	N	N	N

Figure 5.22: The final table summarizing the valid cases where observable behaviors will only be a subset after reordering.

Need to label all figures and refer them properly. Also consider elaborating a bit more on each subcase.

Keep in mind that the comparison of ranges is done while addressing question 3 in the proof, so the table above, implicitly also takes into account only the valid cases where ranges are also correct

The table above, precisely is the definition of a reorderable pair. If we write the above table in the form of an expression we have an expanded format of our Reorderable pair function.

$$\begin{aligned}
Reord(e, d) = & \\
& (((e:uo \wedge d:uo) \wedge \\
& \quad ((e \in R \wedge d \in R) \vee \\
& \quad (e \in W \wedge d \in R \wedge (\mathfrak{R}(e) \ \& \ \mathfrak{R}(d) = \phi)) \vee \\
& \quad (e \in R \wedge d \in W \wedge (\mathfrak{R}(e) \ \& \ \mathfrak{R}(d) = \phi)) \vee \\
& \quad (e \in W \wedge d \in W \wedge (\mathfrak{R}(e) \ \& \ \mathfrak{R}(d) = \phi)))) \\
& \vee \\
& ((e:sc \wedge d:uo) \wedge \\
& \quad ((e \in W \wedge d \in R \wedge (\mathfrak{R}(e) \ \& \ \mathfrak{R}(d) = \phi)) \vee \\
& \quad (e \in W \wedge d \in W \wedge (\mathfrak{R}(e) \ \& \ \mathfrak{R}(d) = \phi)))) \\
& \vee \\
& ((e:uo \wedge d:sc) \wedge \\
& \quad ((e \in R \wedge d \in R) \vee \\
& \quad (e \in W \wedge d \in R \wedge (\mathfrak{R}(e) \ \& \ \mathfrak{R}(d) = \phi))))
\end{aligned}$$

□

□

**Corollary 5.2.1.1.** *Consider a Candidate  $C$  of a program and its Candidate Executions which are valid. Consider two events  $e$  and  $d$  such that  $\neg \text{cons}(e, d)$  is true in*



$C$  and  $e \xrightarrow{ao} d$ . Consider another Candidate  $C'$  resulting after reordering  $e$  and  $d$  in  $C$ . If

$$Reord(e, d) \wedge \forall k \text{ s.t. } e \xrightarrow{ao} k \wedge k \xrightarrow{ao} d . Reord(e, k) \wedge Reord(k, d)$$

then, the set of Observable behaviors of  $C'$  is a subset of  $C$ .

*Proof.* We prove this by induction of number of events  $k$  between  $e$  and  $d$ . Let  $n$  denote the number of events.

**Base Case:**  $n = 1$ . This means we have one event  $k$  such that our candidate  $C$  is

$$e \xrightarrow{ao} k \xrightarrow{ao} d \wedge cons(e, k) \wedge cons(k, d).$$

What we want after reordering is  $C'$ , with

$$d \xrightarrow{ao} k \xrightarrow{ao} e.$$

whose observable behaviors is subset of  $C$ .

Without loss of generality, we can choose to first reorder  $k$  and  $d$ . With  $Reord(k, d)$ , we can, from Theorem 5.2.1, reorder them, giving us candidate  $C''$  with

$$e \xrightarrow{ao} d \xrightarrow{ao} k.$$

whose observable behaviors is subset of  $C$ .

Similarly, now with  $Reord(e, d)$ , by Theorem 5.2.1, we get candidate  $C'''$  with

$$d \xrightarrow{ao} e \xrightarrow{ao} k.$$

whose observable behaviors is subset of  $C''$ .

Now lastly, we need to reorder  $e$  and  $k$  for which we need  $Reord(e, k)$  to hold, thus by Theorem 5.2.1, giving us our final candidate  $C'$  with

$$d \xrightarrow{ao} k \xrightarrow{ao} e.$$

whose observable behaviors is subset of  $C'''$ .

By transitive property of subsets, we can conclude that the Observable Behavior of the final candidate  $C'$  after reordering is a subset of  $C$ .

**2. Inductive Case  $n > 1$**  Assume the above corollary holds for  $n = t$ , meaning the observable behaviors of candidate  $C'_t$  is a subset of  $C_t$ .

We need to show that for  $n = t + 1$ , the corollary still holds, for this note firstly that, we have the following ordering relations:

$$e \xrightarrow{ao} k_1 \xrightarrow{ao} k_2 \xrightarrow{ao} k_3 \xrightarrow{ao} \dots \xrightarrow{ao} k_t \xrightarrow{ao} k_{t+1} \xrightarrow{ao} d$$

Without loss of generality, we can first reorder  $k_{t+1}$  and  $d$ . To do this, we need  $Reord(k_{t+1}, d)$  to hold, thus by Theorem 5.2.1, giving us the resultant candidate  $C_t$  with

$$e \xrightarrow{ao} k_1 \xrightarrow{ao} k_2 \xrightarrow{ao} k_3 \xrightarrow{ao} \dots \xrightarrow{ao} k_t \xrightarrow{ao} d \xrightarrow{ao} k_{t+1}$$

whose observable behaviors is a subset of  $C_{t+1}$ .

Now we have  $t$  such events between  $e$  and  $d$ . With our assumption, we can reorder  $e$  and  $d$ , thus giving us candidate  $C'_t$  with

$$d \xrightarrow{ao} k_1 \xrightarrow{ao} k_2 \xrightarrow{ao} k_3 \xrightarrow{ao} \dots \xrightarrow{ao} k_t \xrightarrow{ao} e \xrightarrow{ao} k_{t+1}$$

whose observable behaviors is a subset of  $C_t$ .

Finally, we need to reorder  $e$  and  $k_{t+1}$  to get our final result, for which we need  $Reord(e, k_{t+1})$  to hold, thus by Theorem 5.2.1, giving us finally candidate  $C'_{t+1}$  with

$$d \xrightarrow{ao} k_1 \xrightarrow{ao} k_2 \xrightarrow{ao} k_3 \xrightarrow{ao} \dots \xrightarrow{ao} k_t \xrightarrow{ao} k_{t+1} \xrightarrow{ao} e$$

Whose observable behaviors are a subset of  $C'_t$ .

By transitive property of subsets, we can conclude that the Observable Behavior of the final candidate  $C'_{t+1}$  after reordering is a subset of  $C_{t+1}$ .

Hence, by induction the proof is complete.  $\square$

To investigate the validity of reordering at the program level, we first, in terms of candidates, address code motion. The following two corollaries cover them:

**Corollary 5.2.1.2.** *Consider a Candidate  $C$  of a program and its Candidate Executions which are valid. Consider a set of events  $k_{i \in [1, n]}$  such that  $k_i \xrightarrow{ao} k_{i+1} \wedge cons(k_i, k_{i+1})$ . Consider an event  $e$  such that*

$$cons(e, k_1) \wedge e \xrightarrow{ao} k_1.$$

*Consider another candidate  $C'$  with the only difference from  $C$  being  $cons(e, k_n) \wedge k_n \xrightarrow{ao} e$ . Then the set of observable behaviors of  $C'$  is a subset of that of  $C$  if*

$$\forall i \in [1, n] . Reord(e, k_i). \quad (5.1)$$

*Proof.* Apply theorem of reordering successively, and by transitivity of subset relations, the corollary holds.  $\square$

**Corollary 5.2.1.3.** *Consider a Candidate  $C$  of a program and its Candidate Executions which are valid. Consider a set of events  $k_{i \in [1, n]}$  such that  $k_i \xrightarrow{ao} k_{i+1} \wedge \text{cons}(k_i, k_{i+1})$ . Consider an event  $d$  such that*

$$\text{cons}(d, k_n) \wedge k_n \xrightarrow{ao} d$$

*Consider another candidate  $C'$  with the only difference from  $C$  being  $\text{cons}(d, k_1) \wedge d \xrightarrow{ao} k_1$ . Then the set of observable behaviors of  $C'$  is a subset of that of  $C$  if*

$$\forall i \in [1, n] . \text{Reord}(k_i, d) \quad (5.2)$$

*Proof.* Apply theorem of reordering successively, and by transitivity of subset relations, the corollary holds.  $\square$

Consider whether to write the proof for code motion or not as it is straightforward induction.

### 5.2.1 Counter examples for the invalid cases

For each case where reordering is not safe to do, we also show counter examples of programs where new observable behaviors are introduced. This additionally portrays additional proof of the validity of our approach.

For all the examples we show here, we only show the ordering relations that are important to observe. Putting all the relations among different events in the example will result in confusion, hence we avoid doing so.

**Reads to same memory where  $e$  is of type  $sc$  while  $d$  is of either  $uo/sc$**  The following example involves two reads to the same memory and a write.

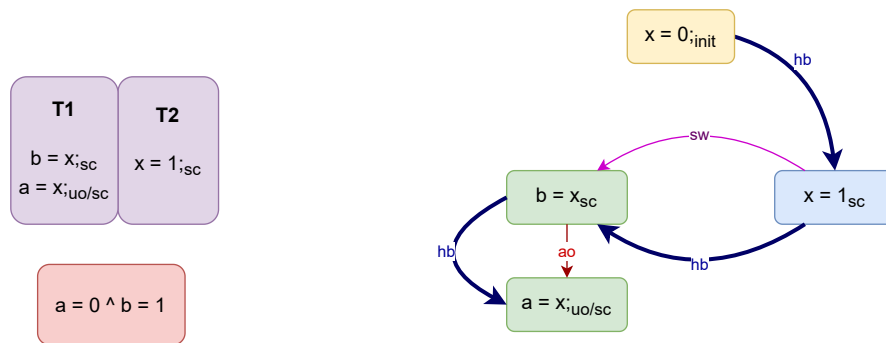


Figure 5.23: Case where  $a = 0$ ,  $b = 1$  is invalid due to Coherent Reads

The figure on the left above shows an example of a candidate where the case of reads in the red box is not possible. The figure on the right shows the Candidate Execution of such a case. Observations:

- We can infer from the Candidate Execution that  $\{x = 0_{init}\} \xrightarrow{hb} \{x = 1_{sc}\} \xrightarrow{hb} \{a = x_{uo/sc}\}$ .
- By the second axiom of coherent reads, it is not possible for  $a$  to read the value of 0 as  $x$  due to the intervening write which changes  $x$  to 1.
- This inference does not rely upon the access mode of the read  $a$ .

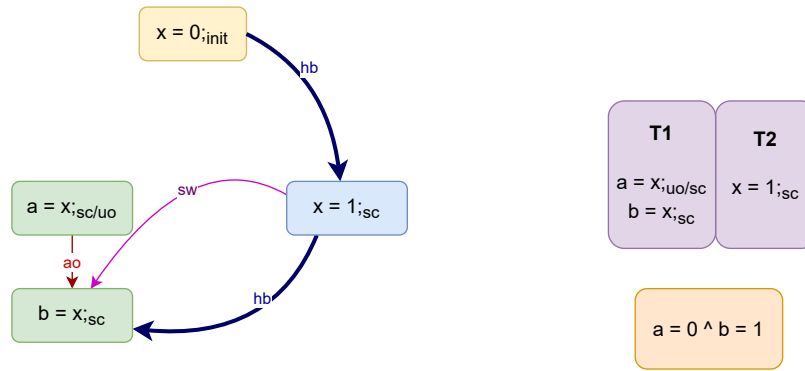


Figure 5.24: Case where the reads are reordered and  $a = 0$ ,  $b = 1$  is valid

The figure on the right shows the program after reordering the two reads in  $T1$ , where the case of reads in the orange box is possible. The figure on the left shows the Candidate Execution of such a case.

Observations:

- From the Candidate Execution, we can infer  $\neg\{x = 0_{init}\} \xrightarrow{hb} \{x = 1_{sc}\} \xrightarrow{hb} \{a = x_{uo/sc}\}$
- We can also infer that  $\{x = 0_{init}\} \xrightarrow{hb} \{a = x_{uo/sc}\}$
- Since none of the axioms disallow the above pattern,  $a$  is allowed to read the value of  $x$  to be 0.
- Hence, the reordering of the two reads is invalid.

**Reads to non-equal range of memory where  $e$  is of type  $sc$  while  $d$  is of either  $uo/sc$**

**A Read  $e$  of type  $sc$  followed by a Write of either  $uo/sc$**  The following is an example of a program with a sequentially consistent read followed by a write of any type.

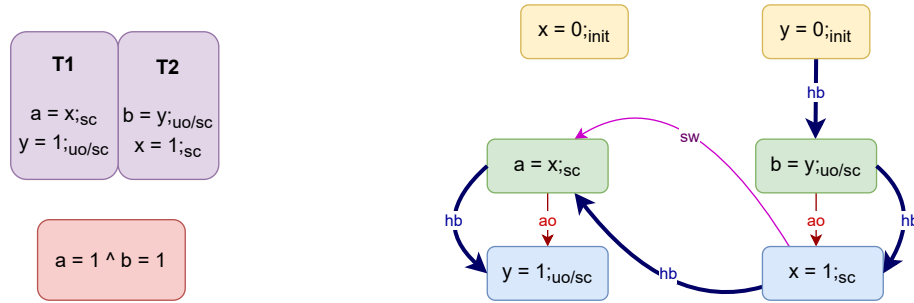


Figure 5.25: Case where  $a = 1$  and  $b = 1$  is invalid due to Coherent Reads.

The figure on the left above shows an example of a candidate where the case of reads in the red box is not possible. The figure on the right shows the Candidate Execution of such a case. Observations:

- From the Candidate Execution, we can infer  $b = y_{uo/sc} \xrightarrow{hb} y = 1_{uo/sc}$
- By the first rule of coherent reads,  $b$  cannot read the value of 1 as  $y$ .
- This inference was due to  $x = 1_{sc} \xrightarrow{hb} a = x_{sc}$

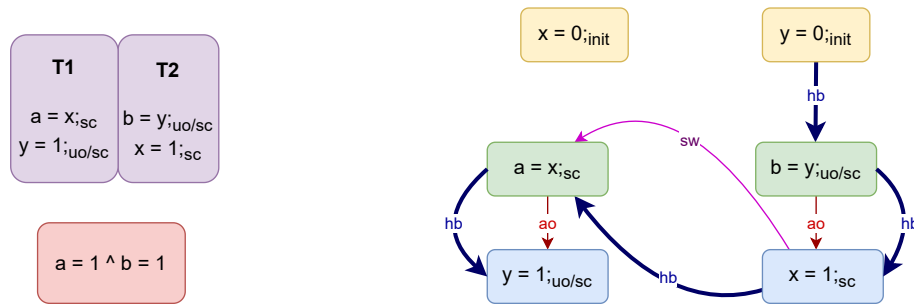


Figure 5.26: Case where events of T1 are reordered, resulting in  $a = 1$  and  $b = 1$  to be valid.

The figure on the right above shows the program after reordering the two events in  $T1$  where case of reads in the orange box is possible. The figure on the left shows the Candidate Execution of such a case. Observations:

- From the Candidate Execution, we can infer  $\neg b = y_{uo/sc} \xrightarrow{hb} y = 1_{uo/sc}$
- Since there is no  $\xrightarrow{hb}$  relation among the above two events,  $b$  can read the value of  $y$  as 1.

**A Read  $e$  of type  $uo$  followed by a write  $d$  of type  $sc$**  For this we can use the same example for the previous part (tag figure of example), where we just reorder  $T2$ 's events.

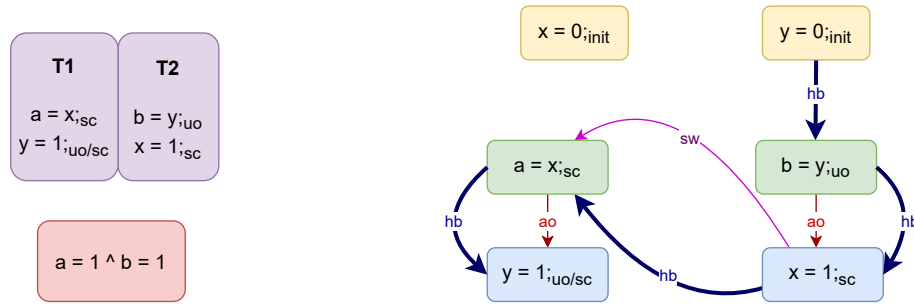


Figure 5.27: Case where  $a = 1$  and  $b = 1$  is invalid due to Coherent Reads.

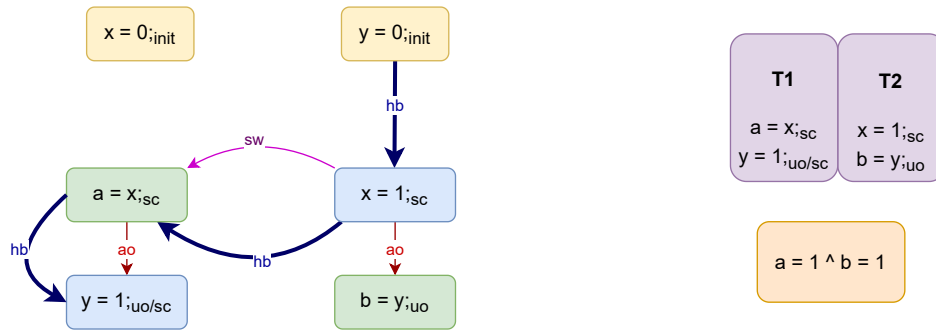
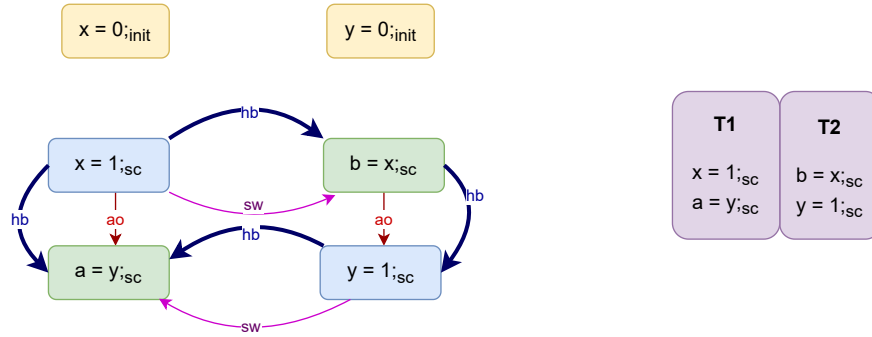
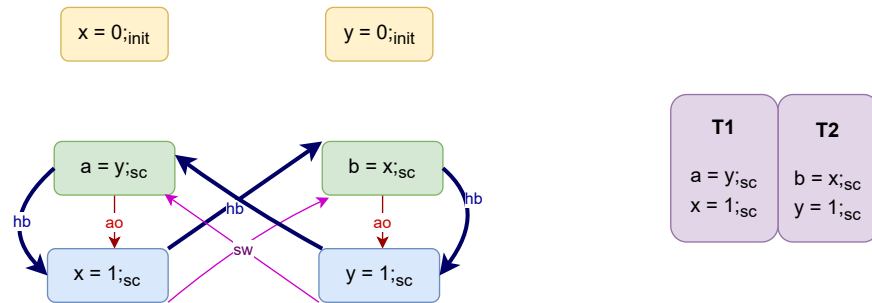


Figure 5.28: Case where events of  $T2$  are reordered, resulting in  $a = 1$  and  $b = 1$  to be valid.

**A Write  $e$  followed by a Read  $d$  both of type  $sc$**  A counter example for this is different. It is not the Observable Behavior we are concerned with that is introduced, but that which is allowed but creates a  $\xrightarrow{hb}$  cycle. The following example is as such:

Figure 5.29: Case where  $a = 1$  and  $b = 1$  is valid and no happens-before cycles

After reordering the two events of  $T1$  in the above example, the same observable behavior holds, but has a cycle introduced. One might think that simply discarding that execution would do. But this would mean discarding  $\xrightarrow{hb}$  relations also, which would require more information to infer which relations are going to create such cycles and which are not. Since we place no assumptions on these relations, but that any happens-before relation other than the one we remove explicitly be reordering are all possible. Hence, the following reordered program outcome is something we do not risk to allow.

Figure 5.30: Case where  $a = 1$  and  $b = 1$  is creates a happens-before cycle

Observation:

- From the read values we can infer that the Candidate Execution should have  $x = 1_{sc} \xrightarrow{hb} a = x_{sc}$  and  $y = 1_{sc} \xrightarrow{hb} a = y_{sc}$ .
- The above relations create the cycle  $a = y_{sc} \xrightarrow{hb} x = 1_{sc} \xrightarrow{hb} a = x_{sc} \xrightarrow{hb} y = 1_{sc} \xrightarrow{hb} a = y_{sc}$ .
- This execution is invalid.

**A Write  $e$  of type  $uo/sc$  followed by a Write  $d$  of type  $sc$**  The following example shows a program with a thread having a write of any access mode( $uo/sc$ ) followed by a write of type  $sc$ .

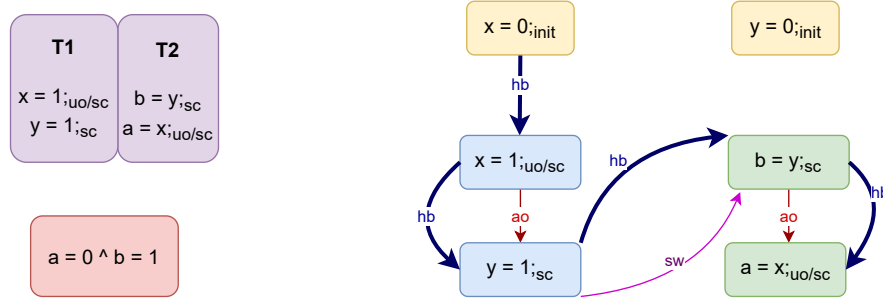


Figure 5.31: Case where  $a = 0$  and  $b = 1$  is invalid due to Coherent Reads.

The figure on the left above shows an example of a candidate where the case of reads in the red box is not possible. The figure on the right shows the Candidate Execution of such a case. Observations:

- From the Candidate Execution, we can infer  $x = 0_{init} \xrightarrow{hb} x = 1_{uo/sc} \xrightarrow{hb} a = x_{uo/sc}$ .
- This is a pattern that matches the second axiom of Coherent reads, thus restricting the read of  $a$  to have the value of 0 as  $x$ .
- This inference was due to  $y = 1_{sc} \xrightarrow{hb} b = y_{sc}$ .

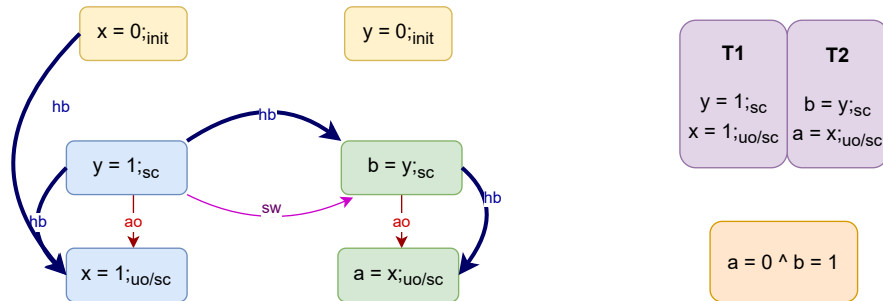


Figure 5.32: Case where events of T1 are reordered, resulting in  $a = 0$  and  $b = 1$  to be valid.



The figure on the right above shows the program after reordering the two events in  $T1$  where case of reads in the orange box is possible. The figure on the left shows the Candidate Execution that explains the orange box case. Observations:

- From the Candidate Execution, we can infer  $\neg x = 1_{uo/sc} \xrightarrow{hb} a = x_{uo/sc}$
- Hence, there is no pattern that the axioms restrict, thus validating  $x$  to be read as 0 by  $a$ .

All the above counter examples have a lot of repetitive text and can have better formal arguments than a list of observations. Make plans to clean up this once filled in with all the counter examples.

### 5.3 From Candidates to Program

We first consider programs with conditionals. The following two properties holds for any candidates of programs having conditional branching.

**Property 1.** *Candidates of Programs with Conditionals* Let  $B1$  be two sets of events based on each branch of a conditional in a program  $P$ . Let  $C$  be any Candidate of  $P$ , Consider  $b1$  to be representative of any event in  $B1$ . Then:

$$\exists C \in P \text{ s.t. } b1 \notin C$$

There exists a candidate of the program such that events from the branch cannot be part of it.

The above property is general for conditionals, being 1-branch or 2-branch. The latter however, has another property which we define below:

**Property 2.** *Candidates of Programs with Conditionals (2-branch)* Let  $B1, B2$  be two sets of events based on each branch of a conditional in a program  $P$ . Let  $C$  be any Candidate of  $P$ , Consider  $b1, b2$  to be representative of any event in  $B1, B2$  respectively. Then:

$$\nexists C \in P \text{ s.t. } b1 \in C \wedge b2 \in C$$

There cannot exist any candidate of the program such that events from both sets can be part of it.

Property 1 may not be true for all programs. Some may always satisfy the conditional in any execution. However, from a general point of view, a program represents all possibilities, hence we can

While the property for 1 branch may not always hold (it can be the case that the branch is always taken in any execution) we are defining it for any program.

*Proof.* Based on an execution of the program, the conditional will either be satisfied or not, but never both. Hence proved both properties. Do we need an elaborate proof of this? As this is direct from existing literature on sequential programs.  $\square$

Perhaps we need a general corollary for program level

**Corollary 5.3.0.1.** *Reordering under Program with Conditionals Consider a program  $P$  and its candidates  $C_1, C_2, \dots, C_n$  in which events  $e$  and  $d$  present in all of them with  $e \xrightarrow{ao} d$ . Consider the set of corresponding candidates  $C'_1, C'_2, \dots, C'_n$  after reordering  $e$  and  $d$  and its corresponding program  $P'$ . Then the set of observable behaviors of  $P'$  is a subset of that of  $P$  if the following two conditions hold:*

$$Reord(e, d) \wedge (\forall C_{i \in [1, n]}, \forall k \in C_i \text{ s.t. } e \xrightarrow{ao} k \wedge k \xrightarrow{ao} d, Reord(e, k) \wedge Reord(k, d))$$

The above condition can be simplified as we already have corollary to show reordering of non-consecutive events No Candidate of  $P$  exists such that only one of  $e$  or  $d$  exists in them.

$$\nexists C \in P \text{ s.t. } (e \in C \wedge d \notin C) \vee (e \notin C \wedge d \in C)$$

*Proof.* Make the above more formal, and use union of sets and subset property.

We prove the second condition first. Suppose the second condition does not hold. Thus we would have

$$\exists C \in P \text{ s.t. } (e \in C \wedge d \notin C) \vee (e \notin C \wedge d \in C)$$

This can only happen, if events  $e$  or  $d$  are part of a conditional branch.

- C1: Both  $e$  and  $d$  are part of conditional branches If  $e$  and  $d$  are in different branches of same conditional, then by Prop 2, we would have

$$\nexists C \in P \text{ s.t. } e \in C \wedge d \in C$$

But our Corollary assumption is that there exists such candidates. Hence, this cannot be the case.

If  $e$  and  $d$  are of the same conditional branch, then our assumption does not hold by Prop 1, 2.

if  $e$  and  $d$  belong to branches of different conditionals, then suppose they are part of conditionals of type Prop 2. Therefore, there exist events  $l1, l2$  in their respective counter branches such that:

$$\begin{aligned} \nexists C \in P \text{ s.t. } e \in C \wedge l1 \in C \\ \nexists C \in P \text{ s.t. } d \in C \wedge l2 \in C \end{aligned}$$

Reordering  $e$  and  $d$  would result in program  $P'$  such that

$$\begin{aligned} \nexists C' \in P' \text{ s.t. } d \in C' \wedge l1 \in C' \\ \nexists C' \in P' \text{ s.t. } e \in C' \wedge l2 \in C' \end{aligned}$$

Thus giving us new Candidates in  $P'$  not in  $P$  such that

$$\begin{aligned} e \in C' \wedge l1 \in C' \\ d \in C' \wedge l2 \in C' \end{aligned}$$

Irrespective of  $e$  and  $d$  being reads or writes, there could be a new  $\xrightarrow{rf}$  relation be formed with some event  $k$ . Thus, we have a new observable behavior.

From the above we can also conclude that even if  $e$  or  $d$  (one of them) are in a conditional branch satisfying Prop 2, a new observable behavior can be introduced due to a new Candidate that violates the original Prop 2 of every candidate of program  $P$ .

Lastly, suppose both  $e$  and  $d$  are part of conditional branches satisfying Prop 1, then we have

$$\begin{aligned} \exists C \in P \text{ s.t. } e \notin C \\ \exists C \in P \text{ s.t. } d \notin C \end{aligned}$$

Let  $B_e$  and  $B_d$  be the respective set of events that belong in the same branch as  $B_e$  and  $B_d$  respectively. Thus, by Prop 1, we also have

$$\begin{aligned} e \notin C \Rightarrow \nexists k \in B_e \text{ s.t. } k \in C \\ d \notin C \Rightarrow \nexists k \in B_d \text{ s.t. } k \in C \end{aligned}$$

Now after reordering  $e$  and  $d$ , we could have a Candidate in  $P'$  such that the above conditions are violated. Irrespective of  $e$  and  $d$  being reads or writes, there could be a new  $\xrightarrow{rf}$  relation be formed with some event  $k$ . Thus, we have a new observable behavior.

Think about this later, discuss with Clark as to how to explain this better. Use diagrams perhaps. This is not well justified. Perhaps place the above condition as a property of each conditional. It would be easier to reason with.

One may think of it as introducing a new event in a Candidate, thus causing new observable behavior.

- C2: Without loss of generality, let us consider  $e$  is part of conditional but  $d$  is not.

By Prop 2, there would exist some event  $l$  in another branch such that

$$\nexists C \in P \text{ s.t. } e \in C \wedge l \in C$$

On reordering  $e$  and  $d$ , we have

$$\nexists C \in P' \text{ s.t. } d \in C \wedge l \in C$$

Thus we have

$$\exists C \in P' \text{ s.t. } e \in C \wedge l \in C$$

giving us a new Candidate in  $P'$  not in  $P$ . Irrespective of  $e$  being a read or a write, there could be a new  $\xrightarrow{rf}$  relation be formed with some event  $k$ . Thus, we have a new observable behavior. By Prop 1, we would have

$$\exists C \in P \text{ s.t. } e \notin C$$

On reordering  $e$  and  $d$ , we have

$$\exists C \in P' \text{ s.t. } d \notin C$$

Thus we have

$$\nexists C \in P' \text{ s.t. } e \notin C$$

giving us a new Candidate in  $P'$  not in  $P$ . Irrespective of  $e$  being a read or a write, there could be a new  $\xrightarrow{rf}$  relation be formed with some event  $k$ . Thus, we have a new observable behavior.

Finding no such event  $k$  would involve gathering more information about other Agent events. Moreover, obtaining such information may be infeasible as the Compiler must ensure that there is No candidate which has such an event  $k$  given that there is event  $e$  as per our condition. Thus, conservatively we can consider doing such reordering to be unsafe.

Now that we have that the second condition must hold, we prove the first condition too must hold. Let  $C_i$  and  $C'_i$  be the candidates before and after reordering  $e$  and  $d$ . From the first condition we have then for  $C_i$

$$\forall k \text{ s.t. } e \xrightarrow{ao} k \wedge k \xrightarrow{ao} d . Reord(e, k) \wedge Reord(k, d).$$

The above is Corollary 1 (tag properly), thus giving us that the observable behaviors of  $C'_i$  is a subset of  $C_i$ . By property of unions of sets, we can conclude that the set of Observable Behaviors of  $P'$  is a subset of that of  $P$ .

Hence proved.

Perhaps we need to define a function *Obs*, that takes a candidate and gives us a set of rf / rbf relations. That would help shaping the argument of observable behavior subset well written.  $\square$

It has come to my notice that in general reordering may not be fine, if we end up removing something outside a conditional. To show this I have a simple counter example. However, the only way to show this, is to say that there is some candidate which suddenly has an agent order relation between two events which were not supposed to have any relation to them. Simply put, we can say that there is a candidate execution of the reordered program, where both the events exist, where as there isn't any candidate of the original program where such a thing can happen.

Next, we consider programs with loops.

Programs with loops might have a little difficulty in defining constraints on Candidates. This is because if we couple them with conditionals then it may be that one iteration would have some events of the loop while the other set of iterations will not have. How then can we define the general case? Perhaps I can consider only programs with loops but no conditionals, then later consider programs having both.

- A :
- B
- C
- D
- E

### 5.3.1

## 6.1 Elimination

- Read Elimination
- Write Elimination

**Theorem 6.1.1.** *Consider a candidate  $C$  of a program and its possible Candidate Executions where  $\overrightarrow{hb}$  is strictly partial order. Consider an event  $e$  which is a read. Consider another Candidate  $C'$  without the event  $e$ . If  $e$  has an unordered access mode, then the set of Observable behaviors of  $C'$  is a subset of  $C$  without the relation  $e \xrightarrow{rf} w$  where  $w$  is some write event in  $C$ .*

*Proof.* We look at this as an elimination of  $e$  that takes place in any candidate execution of  $C$ . We then go about answering the same four questions as we did for reordering. The only major change here being that elimination removes  $\overrightarrow{hb}$  relations. We must check whether the removal of these relations introduce new behaviors, in contrast to that in reordering, where new relations were introduced.

**1. Preserving *happens-before* relations** The relations we want to preserve are those that are derived through relation with  $e$ , meaning the following two relations:

- a)  $k \xrightarrow{hb} e$
- b)  $e \xrightarrow{hb} k$

We can divide the events involved in the above into two sets:

$$K_b = \{k \mid k \xrightarrow{hb} e\}.$$

$$K_a = \{k \mid e \xrightarrow{hb} k\}.$$

Put a figure here for an intuitive understanding of the problem at hand

We need to ensure the following relations hold after elimination.

$$\forall k_a \in K_a \wedge \forall k_b \in K_b . k_b \xrightarrow{hb} k_a \quad (6.1)$$

Slight notational confusion What if the eliminated event is a conditional check? That would mean events in the conditional check are also eliminated. Which would mean one has to check if it is okay to eliminate all events within the conditional.

Similar to reordering, we need to have a valid pivot pair  $\langle p_b, p_a \rangle$  such that

$$\forall k_b \neq p_b \in K_b . k_b \xrightarrow{hb} p_b \quad (6.2)$$

$$\forall k_a \neq p_a \in K_a . p_a \xrightarrow{hb} k_a \quad (6.3)$$

By Lemma 1,  $e : uo$  is the only condition that satisfies our requirement. By Lemma 2,  $e : uo \vee e : sc$  are the options. Considering both the above conditions to be satisfied,  $e : uo$  is the only possibility that holds.

Write an expression which is the conjunction of both lemmas, and show how the conjunction boils down to the result that we come to.

**2. The *happens-before* relations lost** The relations lost are those attached to the event  $e$ , which are:

$$k \xrightarrow{hb} e \vee e \xrightarrow{hb} k \quad (6.4)$$

Do we need to prove that these are the only relations lost? Proof part 1 implicitly shows this.

**3. Presence of Cycles?** Because no new  $\xrightarrow{hb}$  relations are introduced, and because original candidate executions have  $\xrightarrow{hb}$  as a strict partial order, no cycles are introduced after elimination.

Perhaps write this argument a bit better.

**4. Do the lost relations result in New Observable Behaviors?** To answer this, we need to see whether the relations removed had an impact on  $\xrightarrow{rf}$  relations other than those with  $e$ . To prove that it does not have any impact, we divide our argument into two parts, viz. into the two types of relations removed:

$$a) k \xrightarrow{hb} R_u o$$

$$b) R_u o \xrightarrow{hb} k$$

In the first case, we have the following possibilities.

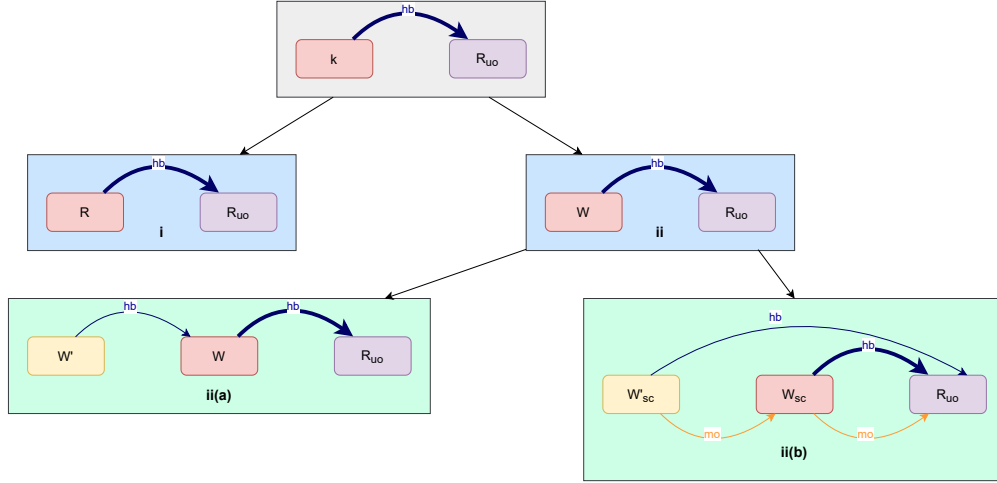


Figure 6.1: The first type of relations removed and the various patterns forbidden by them.

Observations:

- (i) is not a pattern forbidden by the consistency rules
- (ii)(a) is a pattern in Coherent Reads, however, only restricting  $\overrightarrow{rf}$  relation with  $R$  and  $W'$  (which here is our Unordered Read)
- (ii)(b) is a pattern in Sequentially Consistent Atomics, however, once again only restricting  $\overrightarrow{rf}$  relation with  $R$  and  $W'$ .

In the second case, we have the following possibilities.

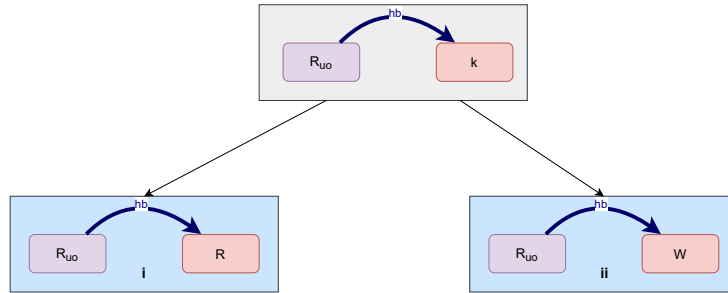


Figure 6.2: The second type of relations removed and the various patterns forbidden by them.



Observations:

- (i) is not a pattern in any Consistency rules
- (ii) is a pattern in Coherent Reads, however, only restricting  $\overrightarrow{rf}$  relation with  $R$  and  $W$

From the above observations, we can see that the relations removed only have restriction on reads-from relations on the event we eliminate. Thus, by case wise analysis we can conclude that no new observable behaviors are introduced due to the removed  $\overrightarrow{hb}$  relations.

□

Explain here why we consider two consecutive write events only. The argument being the Coherent Reads pattern can be triggered anyhow.

**Theorem 6.1.2.** *Consider a candidate  $C$  of a program and its possible Candidate Executions where  $\overrightarrow{hb}$  is strictly partial order. Consider two **write** events  $e$  and  $d$  such that  $\text{cons}(e, d)$  is true in  $C$  and  $e \overrightarrow{ao} d$ . Consider a Candidate  $C'$  without event  $e$ . If  $e$  has an unordered access mode and  $e$  and  $d$  have the same range, then the set of Observable behaviors of  $C'$  is a subset of  $C$ .*

*Proof.* Once again, we look at this as a write elimination done on a Candidate Execution of  $C$ . We start by proving when other happens-before relations remain intact. Followed by identifying relations lost due to elimination and a proof for when these relations do not introduce new observable behaviors.

**Preserving Happens-before relations** The relations we want to preserve are those that are derived through relation with  $e$ , meaning the following two relations:

$$\text{a) } k \overrightarrow{hb} e \qquad \text{b) } e \overrightarrow{hb} k$$

We can divide the events involved in the above into two sets:

$$\begin{aligned} K_b &= \{k \mid k \overrightarrow{hb} e\}. \\ K_a &= \{k \mid e \overrightarrow{hb} k\}. \end{aligned}$$

Put a figure here for an intuitive understanding of the problem at hand

We need to ensure the following relations hold after elimination.

$$\forall k_a \in K_a \wedge \forall k_b \in K_b . k_b \overrightarrow{hb} k_a \tag{6.5}$$

**Slight notational confusion**

Similar to reordering, we need to have a valid pivot pair  $\langle p_b, p_a \rangle$  such that

$$\forall k_b \neq p_b \in K_b . k_b \xrightarrow{hb} p_b \quad (6.6)$$

$$\forall k_a \neq p_a \in K_a . p_a \xrightarrow{hb} k_a \quad (6.7)$$

By Lemma 1,  $e:uo$  is the only condition that satisfies our requirement. It can be our  $p_a$  and by Lemma 2,  $e:uo \vee e:sc$  are the possibilities. Considering both the above conditions to be satisfied,  $e:uo$  is the only possibility that holds.

Again, show the conjunction of both conditions

**2. The *happens-before* relations lost** The relations lost are those attached to the event  $e$ , which are:

$$k \xrightarrow{hb} e \vee e \xrightarrow{hb} k \quad (6.8)$$

Do we need to prove that these are the only relations lost? Proof part 1 implicitly shows this.

**3. Presence of Cycles?** Because no new  $\xrightarrow{hb}$  relations are introduced, and because original candidate executions have  $\xrightarrow{hb}$  as a strict partial order, no cycles are introduced after elimination.

Perhaps write this argument a bit better.

**4. Do the lost relations result in New Observable Behaviors?** To address this, we divide our cases into two parts; one for each type of relation lost:

a)  $k \xrightarrow{hb} e$

b)  $e \xrightarrow{hb} k$

For the first case, we have the following possibilities:

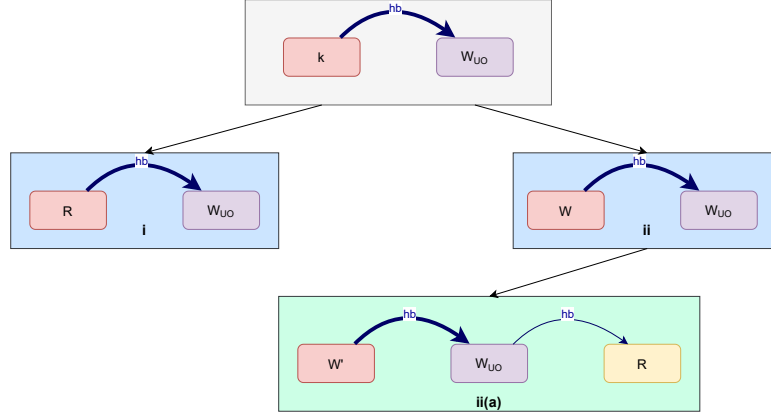


Figure 6.3: First case possibilities (change caption stimiar to that for read elim)

We can observe the following:

- (i) is a pattern from Coherent Reads that restricts the read  $R$  reading from  $W$ . And this will remain the case even after elimination of  $W$ .
- (ii)(a) is a pattern from Coherent reads, forbidding  $R$  to read from some  $W'$ . This will remain the case after elimination of  $W$  if firstly we have  $d \xrightarrow{hb} R$ . By Lemma 2 this is indeed the case. Secondly, we need to ensure that after elimination, the Coherent Reads pattern with  $d$  now restricts the exact set of  $\xrightarrow{rbf}$  relations. Since we have no certain information on the range of  $R$  or  $W'$ , we require the ranges of  $e$  and  $d$  to be same for our requirement to hold in general.
- **PERhaps explain the above argument in more detail**

For the first case, we have the following possibilities:

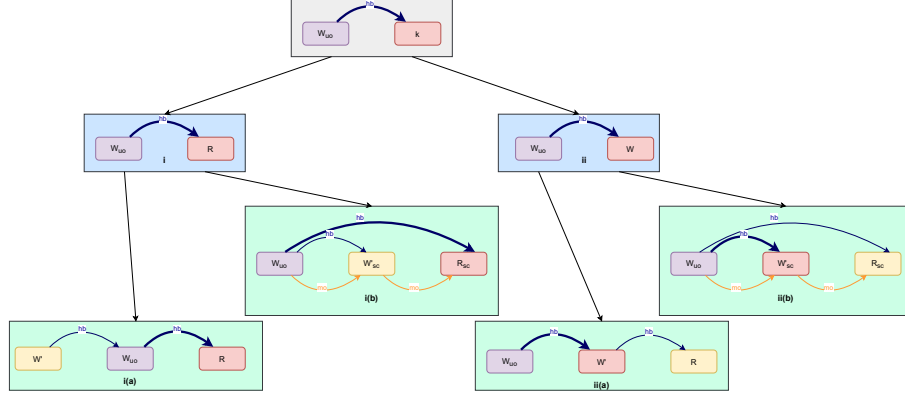


Figure 6.4: Second case possibilities (change caption stimiar to that for read elim)

We make the following observations:

- (i)(a) has the similar argument to the previous case's (ii)(a), requiring  $e$  and  $d$  to have equal ranges.
- (i)(b) is a pattern of Sequentially Consistent Atomics, which restricts  $R$  from reading anything of  $W$ . This will remain the case after  $W$  is eliminated.
- (ii)(a) is a pattern of Coherent Reads, restricting  $R$  from reading  $W$ . This will remain the case after eliminating  $W$ .
- (ii)(b) is the same as (i)(b), hence the argument remains the same.

In all the above cases, observe that on keeping range of  $e$  and  $d$  equal, none of the patterns introduce any new observable behavior. Hence, if we have two consecutive writes of equal ranges, of which the first one has access mode unordered, the set of Observable Behaviors without the write is a subset of that with it present.  $\square$

**Corollary 6.1.2.1.** *Consider a Candidate  $C$  of a program and its Candidate Executions which are valid. Consider two events  $e$  and  $d$  both having equal ranges such that:*

$$e \in W \wedge d \in W \wedge e:uo \wedge e \xrightarrow{ao} d \wedge \neg \text{cons}(e, d)$$

*Consider another Candidate  $C'$  without the event  $e$ . If*

$$\forall k \text{ s.t. } e \xrightarrow{ao} k \wedge k \xrightarrow{ao} d, \text{ Reord}(e, k)$$

*Then, the set of Observable behaviors possible in  $C'$  is a subset of  $C$ .*

*Proof.* We prove by induction on the number of events  $k$  between  $e$  and  $d$ . We verify that if a  $j$  exists that is valid, the Observable behaviors of  $C'$  is a subset of  $C$ .

**Base Case :  $n = 1$**  We have the case when:

$$e \xrightarrow{ao} k_1 \wedge relnk_1 aod$$

By Theorem of Reordering and Def of consecutive events and agent order, we can reorder  $e$  and  $k_1$ , thus giving us a Candidate  $C''$  with :

$$k_1 \xrightarrow{ao} e \wedge e \xrightarrow{ao} d$$

whose observable behaviors are a subset of  $C$ .

By Def of Consecutive instructions and Theorem of Elimination, we can eliminate  $e$ , thus giving us candidate  $C'$  with

$$k_1 \xrightarrow{ao} d$$

whose observable behaviors are a subset of  $C''$ .

By transitive property of subsets, we can conclude that the observable behaviors of  $C'$  is a subset of  $C$ .

**Inductive Case (n)** Let us assume that if the number of events in between are  $n$ , then the corollary holds. Let us consider the Candidate to be  $C_n$  and corresponding candidate after elimination as  $C'_n$ . The observable behavior of  $C'_n$  is a subset of that of  $C_n$ .

If we can show the above holds true for  $n + 1$  events, we are done.

To show this, suppose we have  $C_{n+1}$  as the candidate and  $C'$  as the one after elimination of  $e$ .

Because  $\xrightarrow{ao}$  is a total order, there is a total order among all  $n + 1$  events  $k$  agent ordered between  $e$  and  $d$  such that we can label them  $k_1, k_2, \dots, k_{n+1}$  with the following properties

$$e \xrightarrow{ao} k_1 \xrightarrow{ao} \dots \xrightarrow{ao} k_{n+1} \xrightarrow{ao} d \wedge cons(e, k_1) \wedge cons(k_1, k_2) \wedge \dots \wedge cons(k_{n+1}, d)$$

By Theorem of Reordering and Def of consecutive events and agent order, we can reorder  $e$  and  $k_1$ , thus giving a corresponding candidate  $C_n$  having observable behaviors as a subset of  $C_{n+1}$ .

By our inductive assumption, we have that the observable behaviors of  $C'$  is a subset of  $C_n$ . By transitive property of subsets, we can then conclude that the observable behaviors of  $C'$  are a subset of that of  $C_{n+1}$ .

□

The above proof is clear, but it seems to me that I need to label all definitions and lemmas and theorems and corollary so that I can refer them here.

## 6.2 From Candidates to Program

We first consider the elimination of write in programs with conditional branches. We can summarize the cases that result due to such conditionals into the following parts:

We do not need to know why the compiler is choosing to eliminate a particular write; we only need to assert that whether it is possible given our conservative analysis on observable behaviors. The following corollary states when doing such an elimination is safe:

**Corollary 6.2.0.1.** *Consider a program  $P$  and its candidates  $C_1, C_2, \dots, C_n$  in which events  $e$  and  $d$  present such that*

$$e \in W \wedge d \in W \wedge e : uo \wedge e \xrightarrow{ao} d \wedge \mathfrak{R}(e) = \mathfrak{R}(d)$$

*. Consider the set of corresponding candidates  $C'_1, C'_2, \dots, C'_n$  after eliminating  $e$  and its corresponding program  $P'$ . If*

$$\forall C_{i \in [1, n]}, \forall k \in C_i \text{ s.t. } e \xrightarrow{ao} k \wedge k \xrightarrow{ao} d, \text{ Reord}(e, k)$$

*and*

$$\nexists C \in P \text{ s.t. } e \in C \wedge d \notin C$$

*Then the set of observable behaviors of  $P'$  is a subset of that of  $P$ .*

*Proof.* We first prove that the second condition must hold. We show this by proving that if it does not hold, a new observable behavior can be introduced.

Suppose the second condition does not hold, then we have

$$\exists C \in P \text{ s.t. } e \in C \wedge d \notin C$$

By Prop 1 and Prop 1, we can infer that the above holds if  $e$  or  $d$  are part of a conditional branch.

- Case 1:  $e$  and  $d$  both are part of conditionals By Prop 2 and 1, we have

$$\exists C \in P \text{ s.t. } d \notin C$$

$$\exists C \in P \text{ s.t. } e \notin C$$

After elimination  $e$ , we can have a new observable behavior in a candidate not having  $d$  as above condition states.

Need to refer to part of elimination proof as Coherent Reads would not be triggered anymore for a case and thus we can have a new observable behavior. How to explain this, ask Clark.

- Case 2:  $e$  is part of conditional but  $d$  is not By Prop 2 and 1, we have

$$\exists C \in P \text{ s.t. } e \notin C$$

After elimination  $e$ , we cannot have a new observable behavior in a candidate due to not having  $d$  as above condition states.

- Case 3:  $d$  is part of a conditional but  $e$  is not  
By Prop 2 and 1, we have

$$\exists C \in P \text{ s.t. } d \notin C$$

After elimination  $e$ , we can have a new observable behavior in a candidate not having  $d$  as above condition states.

Need to refer to part of elimination proof as Coherent Reads would not be triggered anymore for a case and thus we can have a new observable behavior. How to explain this, ask Clark.

Add the above property to conditionals with two branches also.

Now that we have that the second condition must hold, we prove the first condition too must hold. Let  $C_i$  and  $C'_i$  be the candidates before and after eliminating  $e$ . From the first condition we have then for  $C_i$

$$\forall k \text{ s.t. } e \xrightarrow{ao} k \wedge k \xrightarrow{ao} d . Reord(e, k).$$

The above is Corollary 1 (tag properly) for elimination, thus giving us that the observable behaviors of  $C'_i$  is a subset of  $C_i$ . Hence this condition must hold for all candidates from which we eliminate  $e$ .

By property of unions of sets, we can conclude that the set of Observable Behaviors of  $P'$  is a subset of that of  $P$ .

Hence proved.

We have not given properly the link between Observable Behaviors, Candidate Executions, Candidates and Programs. Perhaps we need to define a function Obs that gives us the set of Observable Behaviors, where the Domain can be a Program, Candidate, or Candidate Execution.

□

As far as read elimination goes, since we only need the information of read event that is to be eliminated, we do not need to take cases as above for write elimination. Except there can exist one case, in which the read itself is the conditional check. But what is the resultant code after elimination relies on the intention of the compiler, which can be the following:

- It could be plain dead code elimination, wherein both branches of code are eliminated entirely.
- It could also be that the conditional check always returns the same value, which makes the branch taken to be the same.
- It could also be that the choice of branch does not affect the outcome of the program itself.

Since we aren't certain of the reason, it is difficult to identify the target code that is intended after such an elimination. Hence we do not address this case. It is also not within the scope of our analysis to consider the actual mapping between program and candidates. We would need this to prove that the program does not take a particular conditional branch in any execution. This is not easy to do without the mapping in our hands.



## Chapter 7

# Our Critique of The Model

## Chapter 8

# Conclusion, Summary, Future Work