

## 28.7.4 Sequentially Consistent Atomics

For a candidate execution *execution*, memory-order is a strict total order of all events in `EventSet(execution)` that satisfies the following.

- For each pair  $(E, D)$  in *execution*.`[[HappensBefore]]`,  $(E, D)$  is in memory-order.
- For each pair  $(R, W)$  in *execution*.`[[ReadsFrom]]`, there is no `WriteSharedMemory` or `ReadModifyWriteSharedMemory` event  $V$  in `SharedDataBlockEventSet(execution)` such that  $V$ .`[[Order]]` is SeqCst, the pairs  $(W, V)$  and  $(V, R)$  are in memory-order, and any of the following conditions are true.
  - The pair  $(W, R)$  is in *execution*.`[[SynchronizesWith]]`, and  $V$  and  $R$  have equal ranges.
  - The pairs  $(W, R)$  and  $(V, R)$  are in *execution*.`[[HappensBefore]]`,  $W$ .`[[Order]]` is SeqCst, and  $W$  and  $V$  have equal ranges.
  - The pairs  $(W, R)$  and  $(W, V)$  are in *execution*.`[[HappensBefore]]`,  $R$ .`[[Order]]` is SeqCst, and  $V$  and  $R$  have equal ranges.

NOTE 1      This clause additionally constrains SeqCst events on equal ranges.

- For each `WriteSharedMemory` or `ReadModifyWriteSharedMemory` event  $W$  in `SharedDataBlockEventSet(execution)`, if  $W$ .`[[Order]]` is SeqCst, then it is not the case that there is an infinite number of `ReadSharedMemory` or `ReadModifyWriteSharedMemory` events in `SharedDataBlockEventSet(execution)` with equal range that is memory-order before  $W$ .

NOTE 2      This clause together with the forward progress guarantee on agents ensure the liveness condition that SeqCst writes become visible to SeqCst reads with equal range in finite time.

A candidate execution has sequentially consistent atomics if a memory-order exists.