CS-E4500 Problem Set 2

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Two books used as a resources in this report are Cormen et al. (2009), Chapter. 30, Polynomials and FFT and Gathen and Gerhard (2013), Chapter. 8.2, The Discrete Fourier Transform and the Fast Fourier Transform.

Problem 1

Multiplication with the discrete Fourier transform. Let us multiply $f=1+x+x^2\in\mathbb{Z}_{13}[x]$ and $g=2+12x^3\in\mathbb{Z}_{13}[x]$ using the dicrete Fourier transform.

Let ω_n^k be the roots of unity for k = 0, 1, ..., n - 1.

1) Evaluate the polynomials as the roots of unity (DFT)

$$\alpha_k = f(\omega_n^k)$$

$$\beta_k = g(\omega_n^k)$$

2) Multiply the polynomials pointwise

$$\gamma_k = \alpha_k \cdot \beta_k$$

3) Recover the cofficcients of the polynomial using inverse DFT

$$a_j = n^{-1} \sum_{k=0}^{n-1} \gamma_k (\omega_n^{-1})^{kj}$$

The order of the resulting polynomial n=6 and its inverse $n^{-1}=6^{-1}=11\in\mathbb{Z}_{13}.$

The primitive root of unity $\omega_6=4$ and its inverse $w_6^{-1}=4^{-1}=10\in\mathbb{Z}_{13}.$

The power of the root of unity of order 6 are $\omega_6^0,...,\omega_6^5$

The values of polynomial f evaluated at w_6^k are $\alpha_0,...,\alpha_5$

The values of polynomial g evaluated at w_6^k are $\beta_0,...,\beta_5$

The values of the pointwise product $\gamma_k = \alpha_k \cdot \beta_k$ for k = 0, ..., 5 are

The powers of the inverses of the primitive roots of unity are $(\omega_6^{-1})^0,...,(\omega_6^{-1})^5$

Finally, the inverse discrete Fourier tranform gives us the coefficients $a_0,...,a_5$ of the polynomial

We can verify that

$$f \cdot g = 2 + 2x + 2x^2 + 12x^3 + 12x^4 + 12x^5 \in \mathbb{Z}_{13}[x]$$

using for example the naive multiplication algorithm.

Problem 2

The convolution identity. Let $\omega \in R$ be a primitive root of unity of order n in a ring R. Show that for all $f, g \in R[x]/\langle x^n - 1 \rangle$ we have $DFT_{\omega}(fg) = DFT_{\omega}(f) \cdot DFT_{\omega}(g)$.

The polynomial $fg \in R[x]/\langle x^n - 1 \rangle$ is given by the convolution

$$(f*q) \equiv f \cdot q \mod x^n - 1.$$

The by evaluation f * g in the roots of unity ω^j for j = 0, 1, ..., n-1 gives us

$$\begin{split} (f*g)(\omega^j) &= f(\omega^j) \cdot g(\omega^j) + q(\omega^j) \cdot (\omega^j - 1), \quad \omega^j - 1 = 0 \\ &= f(\omega^j) \cdot g(\omega^j). \end{split}$$

This proves the indentity

$$\begin{split} DFT_{\omega}(f*g) &= [(f*g)(\omega^0),...,(f*g)(\omega^{n-1})] \\ &= [f(\omega^0) \cdot g(\omega^0),...,f(\omega^{n-1}) \cdot g(\omega^{n-1})] \\ &= [f(\omega^0),...,f(\omega^{n-1})] \cdot [g(\omega^0),...,g(\omega^{n-1})] \\ &= DFT_{\omega}(f) \cdot DFT_{\omega}(g). \end{split}$$

Problem 3

Input: A polynomial $A(x) = \sum_{j=0}^{n-1} a_j x^j$ in the ring R[x] and roots of unity ω^j for j=0,1,...,n-1 where ω is the primitive root of unity in R.

Output: Vector of values of the polynomial evaluated at the roots of unity $[A(\omega^0), A(\omega^1), ..., A(\omega^{n-1})] \in \mathbb{R}^n$.

Idea: Recursively evaluate polynomial

$$A(x) = A$$

Algorithm

Problem 4

References

Cormen, T.H., Leiserson, C.E., Rivest, R.L. and Stein, C., 2009. *Introduction to Algorithms, Third Edition*. 3rd ed. The MIT Press.

Gathen, J. von zur and Gerhard, J., 2013. *Modern Computer Algebra*. 3rd ed. New York, NY, USA: Cambridge University Press.