

Last time:

think  $k$  perfect field of char  $p > 0$   
 $R$  ess. finite type /  $k$

Then  $D_{R/k} = \bigcup_{e \in \mathbb{N}} \text{Hom}_{R^{\text{op}}}(R, R)$

and this filtration is "linearly cofinal" with the older filtration.

Each  $\text{Hom}_{R^{\text{op}}}(R, R)$  is a subring of  $D_{R/k}$ : often write  $D_R^{(e)}$  "operators of level  $e$ ."

Each  $D_R^{(e)}$  is f.g.  $R$ -module,  
since  $D_R^{(e)} \subseteq D_{R/k}^{(e)} \leftarrow$  f.g.  $R$ -mod.

Easy to see  $D_{R/k}$  is not f.g.  $k$ -algebra in this setting, since any finite subset is contained in some  $D_R^{(e)}$ , so generates a subalgebra of  $D_R^{(e)} \not\subseteq D_{R/k}$ .

For same reason,

if  $k$  has char 0,  $R$  poly ring over  $k$ ,  
then  $D_R^{(e)}$  has no filtration by  
subalgebras that is continual with the older lift.

Ex: Let  $R = k[x_1, \dots, x_n]$  poly ring  
where  $k$  perfect of char  $p > 0$ .

Then  $R^{pe} = k[x_1^{pe}, \dots, x_n^{pe}]$ , and

$$R = \bigoplus_{0 \leq \alpha_1, \dots, \alpha_n < pe} R^{pe} x_1^{\alpha_1} \cdots x_n^{\alpha_n}, \text{ (free } R^{pe}\text{-module).}$$

(collect monomials via congruence classes of exponents  
modulo  $p^e$ ).

Then  $D_R^{(e)}$  is also a free  $R^{pe}$ -module

$$\text{with basis } \{ \ell_{\alpha, \beta} \mid \begin{array}{l} 0 \leq \alpha_1, \dots, \alpha_n < pe \\ 0 \leq \beta_1, \dots, \beta_n < pe \end{array} \}$$

where  $\ell_{\alpha, \beta} (x_1^{\alpha_1} \cdots x_n^{\alpha_n}) = \begin{cases} x^{\beta} & \text{if } \beta = \alpha \\ 0 & \text{otherwise} \end{cases}$

for  $0 \leq \alpha_1, \dots, \alpha_n < p^e$ .

As a ring,  $D_R^{(e)} \cong \text{Mat}_{p^n \times p^n}(R^{pe})$ .

that is, On the other hand,

$$D_{R(K)} = \bigoplus_{\alpha} R D^{(\alpha)}$$

so, can write  $\frac{\partial}{\partial x_i}$  as a matrix,

realizing it in  $\text{Hom}_{R^p}(R, R) \cong \text{Mat}_{p \times p}(R^p)$ ,  
or a different matrix in

$$\text{Hom}_{R^{p^2}}(R, R) \cong \text{Mat}_{p^2 \times p^2}(R^p)$$

Conversely, any  $R^{p^e}$ -linear map from  $R \rightarrow R$   
can be written as an  $R$ -linear combination  
of  $\{D^{(\alpha)}\}$ .

For example, consider

$$\underline{\Phi} := \Phi_{(p^{e-1}, \dots, p^{e-1}), (0, \dots, 0)}, \quad R^{p^e}\text{-linear map}$$

sending  $x^{\frac{p^{e-1}+1}{p-1}} \mapsto 1$   
others  $\mapsto 0$

We have that  $\underline{\Phi} = D^{(p^{e-1}, \dots, p^{e-1})}$ .

To see it, write,  $\alpha = p^e \beta + \gamma$  with  $0 \leq \gamma < p^e$ .

$$\text{Then } \underline{\Phi}(x^\alpha) = x^{p^e \beta} \underline{\Phi}(x^\gamma) = \begin{cases} x^{p^e \beta} & \text{if each } \delta = p^{e-1} \\ 0 & \text{otherwise.} \end{cases}$$

and  $\gamma^{(\varphi_{e-1}, \dots, \varphi_{-1})}(\underline{x}^\alpha) = \binom{\alpha_1}{\varphi_{e-1}} \cdots \binom{\alpha_n}{\varphi_{-1}} x^{\alpha - (\varphi_{e-1}, \dots, \varphi_{-1}) \cdot \underline{1}}$   
 and check that  $\binom{a}{\varphi_{e-1}} = \begin{cases} 1 \bmod p & \text{if } a \equiv \varphi_{e-1} \bmod p \\ 0 \bmod p & \text{otherwise.} \end{cases}$

(Exercise).

# D-modules & D-simplicity

## D-modules

A D-module is just a left  $D_{R,A}$ -module  
 (for some understood  $A \rightarrow R$ ).

We say that  $R$ -module  $M$  is a  
D-module if  $M$  is a left  $D_{R,A}$ -module  
 and the given  $R$ -mod structure agrees with  
 ~~$R$ -module str.~~ on  $M$  via restriction of  
 scalars via  $R \rightarrow D_{R,A}$ ; i.e.,

$$\underset{D_{R,A}\text{-action}}{\overset{R}{\uparrow}} \circ_m = \underset{R\text{-action}}{\overset{R}{\uparrow}} \circ_m \quad \text{for all } r \in R, m \in M.$$

$R = K[x_1, \dots, x_n]$ ,  $K$  char 0,  
 $D_{R/K} \cong R$        $\bar{x}_i$  mult by  $x_i$   
 $\frac{\partial}{\partial x_i}$  differentiate by  $x_i$ .

Any  $K$ -algebra  $R$  (any  $A \rightarrow R$ )  
 $\sim R$  is a  $D$ -module  
 (a ~~left~~  $D_{R/A}$ -module).  
 because  $D_{R/A} \subseteq \underset{\text{subring}}{\text{End}_A(R)}$ .

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Quotient rings?

~~Def~~  $R = S/I$   $S$  poly ring over  $K$   
 $D_{R/K} \cong \frac{(I :_{D_{S/K}} I)}{ID_{S/K}}$

For  $R/I$  to be a  $D_{R/A}$ -module, we  
 need any  $S \in D_{R/A}$  to take  $I$  into  $I$ ;  
 that is:  $(I :_{D_{R/A}} I) = D_{R/A}$ .

We saw that if  $R = S/I$  <sup>spoly ring</sup> Stanley-Reisner ring,  $I = P_1 \cap \dots \cap P_t$   $P_i$  monomial primes then each  $(P_i :_{D_{RIA}} P_i) = D_{RIA}$ .

Thus, in this case, each  $P_i$  or each  $R/P_i$  is a  $D$ -module.

Def:  $A \rightarrow R$  commut. ring

A  $D$ -ideal is a  $D$ -submodule of  $R$ .

Since  $R \hookrightarrow D_{RIA}$ , a  $D$ -ideal is an ideal of  $R$ .

Lem: Let  $I, J$  be  $D$ -ideals, then  $I+J$ ,  $I \cap J$ , and every minimal primary component of  $I$  (if  $R$  is Noeth.) is a  $D$ -ideal.

pf: For "+" and " $\cap$ ", generally fact about submodules.  
If  $Q$  is a primary cpt. of  $I$ , then  $(Q :_{D_{RIA}} Q) = (I :_{D_{RIA}} I) = D_{RIA}$ . □

Rank: If  $I$  is a  $D$ -ideal,  $R/I$  is a  $D$ -module.

The other big source of  $D$ -modules is localization.

$D_{RIA}$  and  $D_{wRIA}$

Prop: Let  $A \rightarrow R$  comm. rings. Let  $M$  be a  $D$ -module. Let  $w \in R$  be mult. closed. Then  $w^{\mathbb{Z}M}$  is a  $D$ -module ( $D_{RA}$ -module) by the rule  $\alpha \cdot \frac{m}{w} = \sum_{i=0}^{\text{ord}(\alpha)} (-1)^i \frac{\alpha^{(i)} \cdot m}{w^{i+1}}$ , where  $\alpha^{(0)} := \alpha$ ,  $\alpha^{(i+1)} := [\alpha^{(i)}, w]$ .

pf: We will use  $\star$  to denote the function  $D_{RA} \times W^{\mathbb{Z}M} \xrightarrow{\star} W^{\mathbb{Z}M}$  defined inductively on order of elt in  $D_{RA}$  by rule

$$\alpha \star \frac{m}{w} := \frac{1}{w} (\alpha \cdot m - [\alpha, \bar{w}] \star \frac{m}{w}), \text{ where}$$

$\star$  denotes the given action  $D_{RA} \times M \xrightarrow{\star} M$ . Then  $\star$  is well-defined, and  $A$ -bilinear. By a straightforward induction,  $\star$  agrees with the action as defined in the statement.

Aside: why  $\star$ ? If  $\star$  worked

$$\alpha \star \left( w \frac{m}{w} \right) = \alpha \star m$$

$$" \qquad \qquad \qquad \alpha \star \frac{m}{w} \qquad \qquad \qquad \alpha \star m$$

$$(w\alpha + [\alpha, \bar{w}]) \star \frac{m}{w}$$

will write  $\alpha' := [\alpha, \bar{w}]$ . Observe first that if  $\bar{r} \in D_{RA}^0$ , then  $\bar{r}$  acts by mult. on  $w^{\mathbb{Z}M}$ .

To check

$$\alpha \star \left( \beta \star \frac{m}{w} \right) = (\alpha \circ \beta) \star \frac{m}{w}.$$

(i) Case  $\alpha \in D_{RA}^0$ :

$$\alpha = \bar{r} \rightsquigarrow$$

$$\bar{r} \star \left( \beta \star \frac{m}{w} \right) = \frac{1}{w} (\beta \cdot m - \beta' \star \frac{m}{w})$$

$$= \frac{1}{w} (\bar{r}(\beta \cdot m) - \bar{r}(\beta' \star \frac{m}{w})) = \frac{1}{w} (\bar{r} \cdot (\beta \cdot m) - \bar{r} \star (\beta' \star \frac{m}{w}))$$

$m$  is a  $D$ -module

$$\alpha \star \frac{m}{w} = \frac{1}{w} (\alpha \cdot m - [\alpha, \bar{w}] \star \frac{m}{w})$$

Induction on order of  $\beta$ .

$$= \frac{1}{w} ((\bar{\alpha} \circ \beta) \cdot m - (\bar{\alpha} \circ \beta') \otimes \frac{m}{w}) = (\bar{\alpha} \circ \beta) \otimes \frac{m}{w}.$$

(ii) Case  $\beta \in D_{\text{PA}}$ : similar.

(iii) General case, by induction on  $\text{ord}(\alpha) + \text{ord}(\beta)$ :

To show:  $w(\alpha \otimes (\beta \otimes \frac{m}{w})) = (\alpha \circ \beta) \otimes \frac{m}{w}$  Definition:  $\alpha \otimes \frac{m}{w} = \frac{1}{w} \alpha \cdot m - [\alpha, w] \otimes \frac{m}{w}$

$$\text{RHS} = (\alpha \circ \beta) \cdot r = (\alpha \circ \beta)' \otimes (r/w)$$

$$= \alpha \cdot (\beta \cdot r) - (\alpha \circ \beta') \otimes (r/w) - (\alpha \circ \beta) \otimes (r/w)$$

$$= \alpha \cdot (\beta \cdot r) - \cancel{\alpha \otimes (\beta \otimes r/w)} - \alpha' \otimes (\beta \otimes r/w)$$

$$\text{LHS} = w(\alpha \otimes (\beta \otimes r/w)) = (\bar{w}\alpha) \otimes (\beta \otimes r/w)$$

$$= (\alpha \bar{w} - \alpha') \otimes (\beta \otimes r/w) = \cancel{\alpha}$$

$$= \alpha \bar{w} \otimes (\beta \otimes (r/w)) - \alpha' \otimes (\beta \otimes r/w)$$

$$= \alpha \otimes (\bar{w}\beta) \otimes r/w - \cancel{\alpha' \otimes \beta} \quad \text{same}$$

$$= \alpha \otimes (\beta \bar{w} - \beta') \otimes r/w - \cancel{\alpha' \otimes \beta} \quad \cancel{w}$$

$$= \alpha \otimes \beta \bar{w} \otimes r/w - \cancel{\alpha \otimes \beta'} \otimes r/w - \cancel{\alpha' \otimes \beta} \quad \cancel{w}$$

$$= \alpha \otimes \beta \otimes \bar{w} \cancel{r/w} - \cancel{\alpha \otimes \beta} \quad \cancel{w} - \cancel{r} \quad \cancel{w}$$

$$= \alpha \otimes \beta \otimes r - \cancel{\alpha \otimes \beta} \quad \cancel{w} - \cancel{r} \quad \cancel{w}$$

$$= \alpha \circ \beta \circ r - \cancel{\alpha \otimes \beta} \quad \cancel{w} - \cancel{r} \quad \cancel{w} \quad \square$$