EE3980 Algorithms

hw06 Trading Stock, II

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Introduction:

In this homework, I will be analyzing, implementing, and observing 2 algorithms, and comparing them with the 2 algorithms in previous homework. The goal of the algorithms is to find the best buying point and selling point for a set of stock price. The input of them will be history of Google stock closing price, and the output of them will be the best buying point, the best selling point and the profit per share.

During the analysis process, I will be using table-counting method to calculate the time complexities of the 2 algorithms. Furthermore, I will try to find the best-case, worst-case, and average-case conditions for the 2 algorithms, respectively. Before implementing on C code, I will try to predict the result based on my analysis. Finally, I will calculate their space complexity for the total spaces used by the algorithm.

The implementation of the 2 algorithms on C code mainly focus on the average time, best-case and worst-case conditions. 9 testing data are given by Professor Chang, and the working environment is on my ubuntu with linux kernel 5.3.0, and gcc version 7.5.0.

Lastly, the observation of them will be focusing on the time complexity of the results from implementation. Moreover, I will compare 4 algorithms with each other and make some rankings. Finally, I will check the experimented results with my analysis.

Analysis:

1. Peak and Valley in an array

Peak and valley are the maximum and minimum values in an array. There are a lot of problems that needs to find peaks and valleys in arrays, for this task, stock trading problem, it can be used to find the best buying and selling point, since the way of acquiring maximum profit is to buy at the lowest price and sell at the highest price.

2. MaxSubArrayBF2 (Brute-force with O(N^2) time complexity)

a. Abstract:

MaxSubArrayBF2 goes through all combinations of buying and selling stocks and finds the 1 way of doing so that has the maximum profit.

The A[i] in the below algorithm indicates the price of each share of stocks. By subtracting the prices of 2 shares, the algorithm gets the profits and find the maximum value among them.

Not clear.

b. Algorithm:

```
1.  // Find low and high to maximize A[low] – A[high], low  high.
2.  // Input: A[1 : n ], int n
3.  // Output: 1 >= low, high <= n and max.
4.  Algorithm MaxSubArrayBF2(A, n, low, high)
5.  {
6.      max := 0 ; // Initialize
7.      low := 1 ;
8.      high := n ;
9.      for j := 1 to n do { // Try all possible ranges: A[j : k ].
10.         for k := j to n do {
11.             sum := A[k] – A[j] ;
12.             if (sum > max) then { // Record the maximum value and range.
13.                 max := sum ;
14.                 low := j ;
15.                 high := k ;
16.             }
17.         }
18.     }
19.     return max ;
20. }
```

c. Proof of correctness:

In this algorithm, it calculates the profits for every possible combination of buying and selling this stock in N*N iterations, where N is the number of stock prices. In line 15 ~ 19, it constantly replaces the stored maximum combinations. Using induction, we can conclude that in every iteration, it either stores the best way so far to buy and sell a stock, or do not store anything. The algorithm terminates when all combination is tested and return the best way.

d. Time complexity:

	s/e	freq	total
1. Algorithm MaxSubArrayBF2(A, n, low, high)	0	0	0
	0	0	0
2. {	1	1	1
3. max := 0 ;	1	1	1
4. low := 1 ;	1	1	1
5. high := n ;	N	1	N
6. for j := 1 to n do {	N/2	N	1/2N^2
7. for k := j to n do {	1	N^2	N^2
8. sum := A[k] – A[j] ;	1	N^2	N^2
9. if (sum > max) then {	1	N^2	N^2
10. max := sum ;	1	N^2	N^2
11. low := j ;	1	N^2	N^2
12. high := k ;	0	N	0
13. }	0	0	0
14. }	0	0	0
15. }	1	1	1
16. return max ;	0	0	0
17. }			
p.s. N/2 sine it goes through 1 ~ N in N iterations	5.5N^2 + N + 4		

The time complexity of *MaxSubArrayBF2* is O(N^2), where N is the number of stock shares.

Best case, Worst case and Average case:

The difference between best case and worst case for this is not obvious. The reason for this is that either way, it goes through all iterations anyway. The only difference is the times of updating the maximum value, which we can neglect since it costs only few steps. And for the above reasons, the average case is quite the same, too.

e. Space Complexity:

The algorithm uses 6 integers and N pairs of elements in array of stocks. **The space complexity would be O(N).**

3. MaxSubArrayN (Search Extreme Values)

a. Abstract:

MaxSubArrayN finds the peak and valley in the array of stock shares. By subtracting the maximum and minimum value of the shares, it gets the maximum profit, the best buying and selling point for this set of stocks. Be aware that the valley should be prior to the peak.

b. Algorithm:

```
1.  // Find low and high to maximize A[low] – A[high], low  high.
2.  // Input: A[1 : n ], int n
3.  // Output: 1 >= low, high <= n and max.
4.  Algorithm MaxSubArrayN(A, n, low, high)
5.  {
6.      minprice := 0 ; // Initialize
7.      sum := 0 ;
8.      for i := 1 to n do { // Try all possible ranges: A[j : k ].
9.          if  A[i] > minprice do {
10.             minprice = A[i];    assignment?
11.         }
12.     }
13.     for i := 1 to n do { // Try all possible ranges: A[j : k ].
14.         if  A[i] < minprice do {
15.             minprice = A[i];
16.             low2 = i;
17.         }  else if (A[i] - minprice) > sum do {
18.             sum = A[i] - minprice;
19.             low1 = low2;
20.             high = i;
```

```
21.         }
22.     }
23.     return sum ;
24. }
```

c. Proof of correctness:

In this algorithm, it first finds the maximum value among the array A in order to set a boundary for the next step, find minimum. It takes N iterations to find it. Using simple induction, we can prove that we can find the maximum value in the array. In each iteration, we store the i-i-th value if it is larger than previous maximum value.

Next, it searches for the valley and peak in an array. In each iteration in N iterations, it updates the valley, which is the best buying point, and updates the peak, which is the best selling point. Using induction, as aforementioned paragraph, we can prove that it is correct. Be aware that the buying point is updated when it finds selling point later than it.

d. Time complexity:

	s/e	freq	total
1. Algorithm MaxSubArrayN(A, n, low, high)	0	0	0
2. {	0	0	0
3. minprice := 0 ;	1	1	1
4. sum := 0 ;	1	1	1
5. for i := 1 to n do {	N	1	N
6. if A[i] > minprice do {	1	N	N
7. minprice = A[i];	1	N	N
8. }	0	N	0

9.	}	0	1	0
10.	for i := 1 to n do {	N	1	N
11.	if A[i] < minprice do {	1	N	N
12.	minprice = A[i];	1	N	N
13.	low2 = i;	1	N	N
14.	} else if (A[i] - minprice) > sum do{	1	N	N
15.	sum = A[i] - minprice;	1	N	N
16.	low1 = low2;	1	N	N
17.	high = i;	1	N	N
18.	}	0	0	0
19.	}	0	0	0
20.	return sum ;	1	1	1
21.	}	0	0	0
		11N + 4		

The time complexity of *MaxSubArrayBF2* is O(N), where N is the number of stock shares.

Best case, Worst case and Average case:

The difference between best case and worst case for this is not obvious. The reason for this is that either way, it finds peak and valley of the array. The only difference is the times of updating the maximum value, which we can neglect since it costs only few steps. And for the above reasons, the average case is quite the same, too.

e. Space Complexity:

The algorithm uses 7 integers and N pairs of elements in array of stocks. **The space complexity would be O(N).**

4. Time & Space Complexity Comparison:

	MaxSubArrayBF	*MaxSubArray*	*MaxSubArrayBF2*	*MaxSubArrayN*
Time	O(N^3)	O(N lg N)	O(N^2)	O(N)
Space	O(N)	O(N)	O(N)	O(N)

Speed (fast>slow):

***MaxSubArrayN > MaxSubArray> MaxSubArrayBF2 > MaxSubArrayBF*.**

Implementation:

1. **Speed Test:**

Speed Test is to find the actual speed and time complexities of the 4 algorithms, *MaxSubArrayBF, MaxSubArray, MaxSubArray* and *MaxSubArrayN*. We use 9 test inputs given by Professor and get the CPU runtimes before and after the algorithms perform their tasks. The implementation is done on my laptop. However, the time recording methods for the 4 algorithms are different. Since *MaxSubArrayBF* runs much slower than the rest of them, I can only run *MaxSubArrayBF* once and record the CPU runtime. However, I will run the others 1000 times and record the average runtime for it.

Workflow for *MaxSubArrayBF* :

```
1.  t_MaxSubArrayBF = GetTime();              // initialize time counter
2.  MaxSubArrayBF();
3.  t_MaxSubArrayBF = GetTime() - t_MaxSubArrayBF;  // calculate CPU time
```

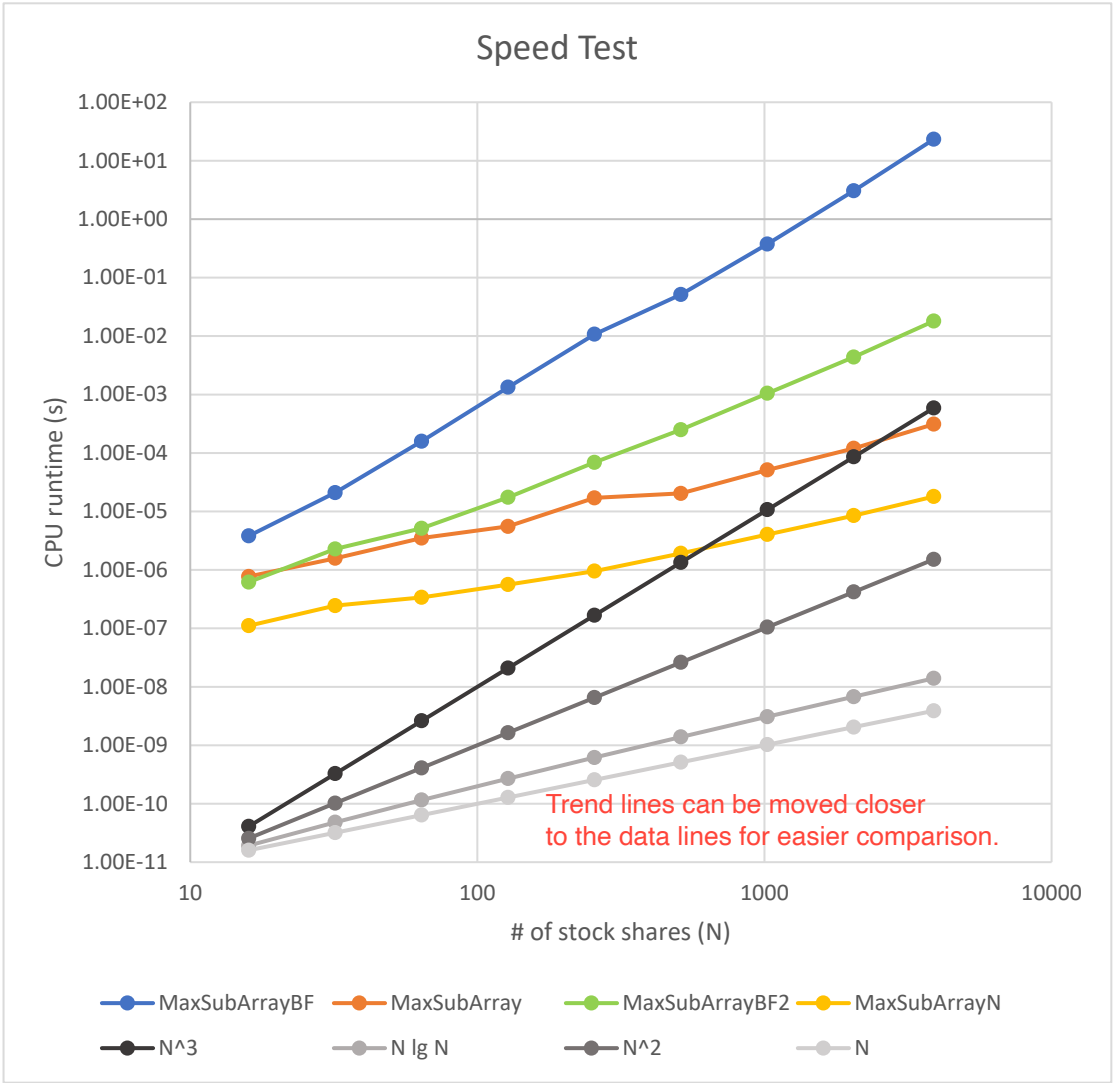
Workflow for the other 3 algorithms :

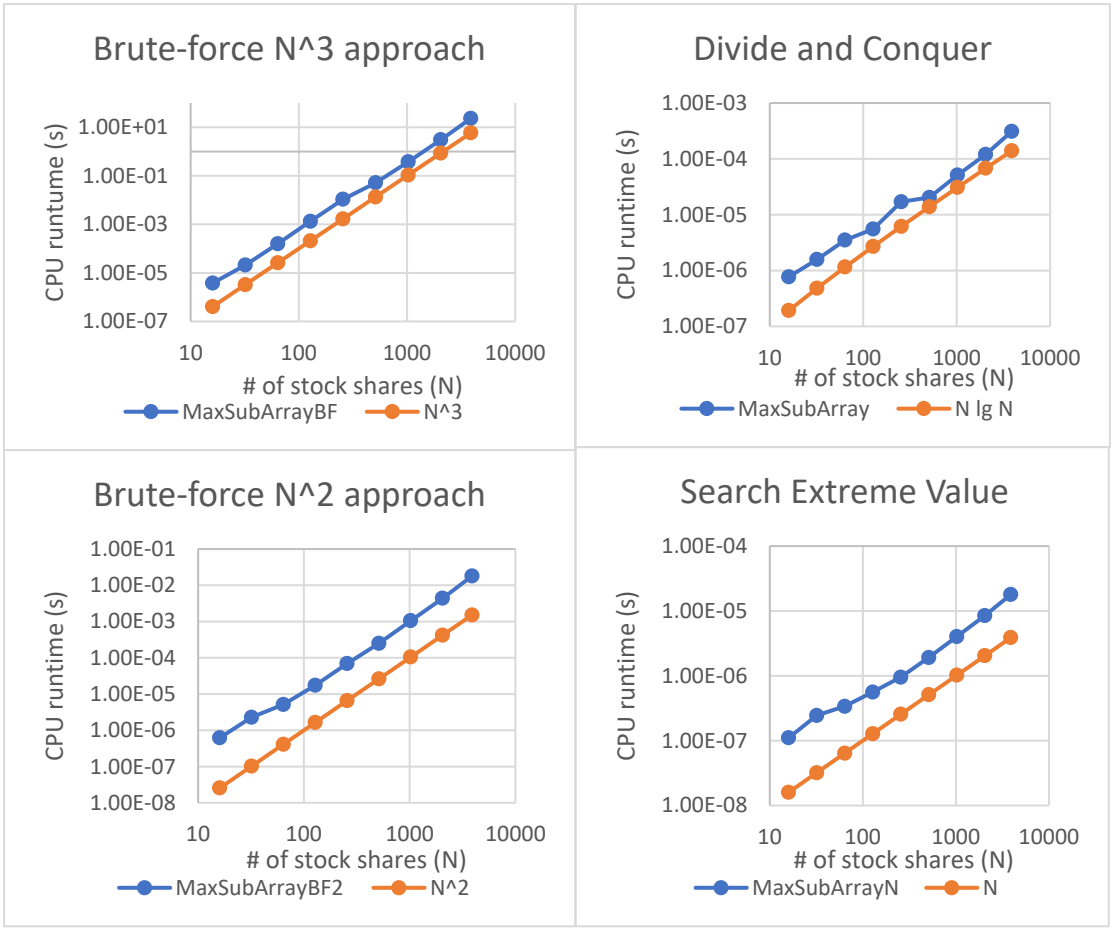
```
4.  t = GetTime();              // initialize time counter
5.  for i := 0 to 1000 do {
```

```
6.     Algorithm();
7.  }
8.  t = (GetTime() - t) / 1000; // calculate CPU time
```

Results:

N (s)	*MaxSubArrayBF*	*MaxSubArray*	*MaxSubArrayBF2*	*MaxSubArrayN*
16	3.81E-06	7.72E-07	6.19E-07	1.11E-07
32	2.10E-05	1.58E-06	2.28E-06	2.45E-07
64	1.59E-04	3.51E-06	5.15E-06	3.39E-07
128	1.34E-03	5.55E-06	1.74E-05	5.62E-07
256	1.07E-02	1.71E-05	6.90E-05	9.52E-07
512	5.15E-02	2.03E-05	2.51E-04	1.92E-06
1024	3.76E-01	5.12E-05	1.05E-03	4.04E-06
2048	3.06E+00	1.20E-04	4.38E-03	8.49E-06
3890	2.33E+01	3.11E-04	1.81E-02	1.80E-05





Observation:

1. Speed, Time complexity:

Actual Speed (> means faster):

***MaxSubArrayN > MaxSubArray> MaxSubArrayBF2 > MaxSubArrayBF*.**

The result matches my analysis precisely. The time complexity of *MaxSubArrayBF2* and *MaxSubArrayN* are **O(N^2)** and **O(N)**. Modify and deviate from maximum subarray approach made the algorithm did made the iteration depth for *MaxSubArrayBF2* lesser by N. By seeking peak and valley in an array did made the time complexity goes by O(N) only.

Overall, the implemented results meet my analysis.

Conclusions:

1. Time and space complexities of the 5 algorithms:

	MaxSubArrayBF	*MaxSubArray*	*MaxSubArrayBF2*	*MaxSubArrayN*
Time	O(N^3)	O(N lg N)	O(N^2)	O(N)
Space	O(N)	O(N)	O(N)	O(N)

2. Actual runtime comparison:

   Speed (> means faster):

   MaxSubArrayN > MaxSubArray> MaxSubArrayBF2 > MaxSubArrayBF.

3. It goes faster when not using the maximum subarray property.

Not clear.

hw06.c

```
1 // EE3980 HW05 Trading Stock, II
2 // 106061146, Jhao-Ting, Chen
3 // 2020/04/18
4
5 #include <stdio.h>
6 #include <stdlib.h>
7 #include <string.h>
8 #include <sys/time.h>
9
10 typedef struct sSTKprice {                  // Storing stock shares
11     int year, month, day;
12     double price, change;
13 } STKprice;
14
15 typedef struct sMaxArray {                  // return found value
16     int low, high;
17     double change;
18 } MaxArray;
19
20 int R = 1000;                               // # of test cycles
21 int N;                                      // # of stock shares
22 STKprice *data;                             // Stock list
23
24 void readInput(void);                       // read all inputs
25 void printInput(void);                      // print Input
26 double GetTime(void);                       // get local time in seconds
27 MaxArray MaxSubArrayBF(STKprice *A, int N);             // Brute-force
28 MaxArray MaxSubArray(STKprice *A, int begin, int end);  // Divide and Conquer
29 MaxArray MaxSubArrayXB(STKprice *A, int begin, int mid, int end);
30 MaxArray MaxSubArrayBF2(STKprice *A, int N);            // Brute-force N^2
31 MaxArray MaxSubArrayN(STKprice *A, int N);          // Search Extreme Value
32
33 int main(void)
34 {
35     int i;                                   // loop index
36     double t;                                // record CPU time
37     MaxArray ans;                            // returned value
38
39     readInput();                             // read input graph
40 //    printInput();
41     printf("N = %d\n", N);
42
43     t = GetTime();                           // initialize timer
44     ans = MaxSubArrayBF(data, N);            // testing Brute-force
45
46     printf("Brute-force approach: time %e s\n", (GetTime() - t));   // result
47     printf("  Buy: %d/%d/%d at %g\n", data[ans.low].year, data[ans.low].month,
48             data[ans.low].day, data[ans.low].price);
```

```
49      printf("  Sell: %d/%d/%d at %g\n", data[ans.high].year,
50              data[ans.high].month, data[ans.high].day, data[ans.high].price);
51      printf("  Earning: %g per share.\n", ans.change);
52
53      t = GetTime();                          // initialize time counter
54      for (i = 0; i < R; i++) {               // testing Divide and Conquer
55          ans = MaxSubArray(data, 0, N - 1);
56      }
57
58      printf("Divide and Conquer: time %e s\n", (GetTime() - t) / R);
59      printf("  Buy: %d/%d/%d at %g\n", data[ans.low].year, data[ans.low].month,
60              data[ans.low].day, data[ans.low].price);
61      printf("  Sell: %d/%d/%d at %g\n", data[ans.high].year,
62              data[ans.high].month, data[ans.high].day, data[ans.high].price);
63      printf("  Earning: %g per share.\n", ans.change);
64
65      t = GetTime();                          // initialize time counter
66      for (i = 0; i < R; i++) {               // testing Brute-force 2
67          ans = MaxSubArrayBF2(data, N);
68      }
69
70      printf("Brute-force #2 approach: time %e s\n", (GetTime() - t) / R);
71      printf("  Buy: %d/%d/%d at %g\n", data[ans.low].year, data[ans.low].month,
72              data[ans.low].day, data[ans.low].price);
73      printf("  Sell: %d/%d/%d at %g\n", data[ans.high].year,
74              data[ans.high].month, data[ans.high].day, data[ans.high].price);
75      printf("  Earning: %g per share.\n", ans.change);
76
77      t = GetTime();                          // initialize time counter
78      for (i = 0; i < R; i++) {               // testing Search Extremes
79          ans = MaxSubArrayN(data, N);
80      }
81
82      printf("Search Extreme Values: time %e s\n", (GetTime() - t) / R);
83      printf("  Buy: %d/%d/%d at %g\n", data[ans.low].year, data[ans.low].month,
84              data[ans.low].day, data[ans.low].price);
85      printf("  Sell: %d/%d/%d at %g\n", data[ans.high].year,
86              data[ans.high].month, data[ans.high].day, data[ans.high].price);
87      printf("  Earning: %g per share.\n", ans.change);
88
89      return 0;
90  }
91
92  void readInput(void)                            // read all inputs
93  {
94      int i;                              // for looping and dynamic store
95
96      scanf("%d\n", &N);                  // read # of Vertices and Edges
97
98      data = (STKprice *)calloc(N, sizeof(STKprice));
```

```
 99
100
101     for (i = 0; i < N; i++) {                   // Store shares in data
102         scanf("%d %d %d %lf\n", &data[i].year, &data[i].month,
103                                 &data[i].day, &data[i].price);  // read shares
104     }
105
106     data[0].change = 0.0;                       // store each share's change
107     for (i = 1; i < N; i++) {
108         data[i].change = data[i].price - data[i - 1].price;
109     }
110 }
111
112 void printInput(void)                           // print stock list
113 {
114     int i;
115     for (i = 0; i < N; i++) {
116         printf("%d %d %d %lf %lf\n", data[i].year, data[i].month, data[i].day,
117                                      data[i].price, data[i].change);
118     }
119 }
120
121 double GetTime(void)                    // demonstration code from 1.1.3
122 {
123     struct timeval tv;
124
125     gettimeofday(&tv, NULL);
126     return tv.tv_sec + 1e-6 * tv.tv_usec;
127 }
128
129 MaxArray MaxSubArrayBF(STKprice *A, int N)
130 {
131     double max = 0;                 // initialize
132     double sum;
133     int low = 0;
134     int high = N - 1;
135     int i, j, k;
136
137     MaxArray ans;
138
139     for (j = 0; j < N; j++) {       // Try all possible ranges: A[j : k].
140         for (k = j; k < N; k++) {
141             sum = 0;
142             for (i = j + 1; i <= k; i++) {  // Summation for A[j : k]
143                 sum = sum + A[i].change;
144             }
145             if (sum > max) {            // Record the maximum value and range
146                 max = sum;
147                 low = j;
148                 high = k;
```

```
149             }
150         }
151     }
152     ans.low = low;
153     ans.high = high;
154     ans.change = max;
155     return ans;
156 }
157
158 MaxArray MaxSubArray(STKprice *A, int begin, int end)
159 {
160     int mid;                    // terminate condition
161     MaxArray ans;
162     MaxArray lsum, rsum, xsum;
163
164     if (begin == end) {
165         ans.low = begin;
166         ans.high = end;
167         ans.change = 0;
168         return ans;
169     }
170     mid = (begin + end) / 2;
171     lsum = MaxSubArray(A, begin, mid);      // left region
172     rsum = MaxSubArray(A, mid + 1, end);    // right region
173     xsum = MaxSubArrayXB(A, begin, mid, end);   // cross bundary
174     if ((lsum.change >= rsum.change) && (lsum.change >= xsum.change)) {
175         ans.low = lsum.low;                     // lsum is the largest
176         ans.high = lsum.high;
177         ans.change = lsum.change;
178         return ans;
179     } else if ((rsum.change >= lsum.change) && (rsum.change >= xsum.change)) {
180         ans.low = rsum.low;                     // rsum is the largest
181         ans.high = rsum.high;
182         ans.change = rsum.change;
183         return ans;
184     }
185     ans.low = xsum.low;                     // cross-boundary is the largest
186     ans.high = xsum.high;
187     ans.change = xsum.change;
188     return ans;
189 }
190
191 MaxArray MaxSubArrayXB(STKprice *A, int begin, int mid, int end)
192 {
193     double lsum, sum, rsum;         // initialize
194     int i;
195     int low, high;
196     MaxArray ans;
197
198     low = mid;                      // initialize for lower half
```

```
199     lsum = 0.0;
200     sum = 0.0;
201
202     for (i = mid; i > begin; i--) {     // find low to maximize A[low : mid]
203         sum = sum + A[i].change;        // continue to add
204         if (sum > lsum) {               // record if larger
205             lsum = sum;
206             low = i;
207         }
208     }
209
210     rsum = 0.0;                         // Initialize for higher half
211     high = mid + 1;
212     sum = 0.0;
213
214     for (i = mid + 1; i <= end; i++) {  // find end
215         sum = sum + A[i].change;        // continue to add
216         if (sum > rsum) {               // record if larger
217             rsum = sum;
218             high = i;
219         }
220     }
221
222     ans.low = low - 1;
223     ans.high = high;
224     ans.change = lsum + rsum;           // overall sum
225     return ans;
226 }
227
228 MaxArray MaxSubArrayBF2(STKprice *A, int N)
229 {
230     double max = 0;                 // initialize
231     double sum;
232     int low = 0;
233     int high = N - 1;
234     int j, k;
235
236     MaxArray ans;
237
238     for (j = 0; j < N; j++) {       // Try all possible ranges: A[j : k].
239         for (k = j; k < N; k++) {
240             sum = A[k].price - A[j].price;      // calculate profit
241             if (sum > max) {            // Record the maximum value and range
242                 max = sum;
243                 low = j;
244                 high = k;
245             }
246         }
247     }
248     ans.low = low;
```

```
249     ans.high = high;
250     ans.change = max;
251     return ans;                         // return the result
252 }
253
254 MaxArray MaxSubArrayN(STKprice *A, int N)
255 {
256     double minprice = 0.0;              // initialize
257     double sum = 0.0;
258     MaxArray ans;
259     int i, low1, low2, high;
260     for (i = 0; i < N; i++){            // get the boundary for finding valley
        for (i = 0; i < N; i++) {            // get the boundary for finding valley
261         if ( A[i].price > minprice) {
            if (A[i].price > minprice) {
262             minprice = A[i].price;
263         }
264     }
265
266     for (i = 0; i < N; i++){
        for (i = 0; i < N; i++) {
267         if (A[i].price < minprice) {    // finding valley
268             minprice = A[i].price;
269             low2 = i;                   // store found valley
270         } else if (A[i].price - minprice > sum) {   // finding peak (profit)
271             sum = A[i].price - minprice;            // calculate profit
272             low1 = low2;
273             high = i;                   // store found peak
274         }
275     }
276     ans.low = low1;
277     ans.high = high;
278     ans.change = sum;
279     return ans;                         // return result
280 }
```

[Program Format] can be improved.
[Coding] hw06.c spelling errors: bundary(1)
[Good] effort in doing the homework, but improvements are still needed.

Score: 91