Algorithm EE398000 hw01

106061146 陳兆廷

1. Introduction:

The analysis and comparisons of the four sorting algorithms: selection sort, insertion sort, bubble sort and shaker sort.

The goal of the algorithms is to sort a list of vocabulary into alphabetical order.

The input of the algorithms is a list of vocabulary, which is in a form of 2d array. The output of the algorithms is a list of sorted, alphabetically ordered vocabulary.

2. Approach:

a. Analysis:

Analysis of Selection Sort:

i. **Intro:** Find minimum value, put it in the front of array.

ii. **Proof:** Each iteration in the ith iterations, it brings the smallest value among the n-1 items to the front of the array. Therefore, the algorithm finds the values in ascending order, undoubtly.

iii. Space complexity:

Fixed part: int: (i, j, k, n), *char: (tmp)

Variable part: *char (A) x n

Total: n (*char) + 4

iv. Time complexity:

| Statement | s/e | Freq. | Total |
|---|--------|--------|-----------------------|
| Algorithm SelectionSort(A, n) | 0 | 0 | 0 |
| 2. { | 0 | 0 | 0 |
| 3. for i := 1 to n do { | n+1 | 1 | n+1 |
| 4. j := i; | 1 | n | n |
| 5. for k := i+1 to n do { | n+1 | n | n(n+1) |
| 6. if (A[k] < A[j]) then j = k | 2 or 1 | n(n+1) | n(n)*2or1 |
| 7. } | 0 | 0 | 0 |
| 8. tmp = list[i]; list[i] = list[j]; list[j] = tmp; | 3 | n | 3n |
| 9. } | 0 | 0 | 0 |
| 10. } | 0 | 0 | 0 |
| Total | | | 3n ² +4n+0 |

*since aveg of i is n/2, I calculate T(P) with i = n/2.

Best case: $2n^2 + 4n + 2 (O(n^2))$.

Worst Case: $3n^2 + 4n + 2 (O(n^2))$.

Average Case: O(n^2).

Therefore, the time complexity of Selection sort is O(n^2).

Analysis of Insertion Sort:

- i. **Intro:** Compare with the minimum value, if true, swap, put it in the front of array.
- ii. **Proof:** Like Selection sort, but Insertion sort gradually inserts the items from 2 to n into array in the front, in the right order.

iii. Space complexity:

Fixed part: int: (i, j, n), *char: (tmp)

Variable part: *char (A) x n

Total: n (*char) + 3

iv. Time complexity:

| Statement | s/e | Freq. | Total |
|---|-------|-----------|-----------|
| Algorithm InsertionSort(A, n) | 0 | 0 | 0 |
| 2. { | 0 | 0 | 0 |
| 3. for j := 2 to n do { | n-1 | 1 | n-2 |
| 4. tmp := A[j]; | 1 | n-2 | n-2 |
| 5. i := j - 1; | 1 | n-2 | n-2 |
| 6. while ((i >= 1) and (tmp < A[i])) do { | n (*) | n-2 | n^2-2n |
| 7. A[i + 1] := A[i]; i := i-1; | 2 | (n-2)(c)n | (2n-4)cn |
| 8. } | 0 | (n-2)(c)n | 0 |
| 9. A[i + 1] = tmp; | 1 | n-2 | n-2 |
| 10. } | 0 | 0 | 0 |
| 11. } | 0 | 0 | 0 |
| Total | | | (1+2c)n^2 |

^{*}Assume >=, <, and are all performed. (I recall that if i<1, computer won't bother do the rest computation.)

Best case: if c = 0 or 1, it is mostly sorted, it won't do anything. (O(1))

Worst case: if c \sim = n, it is almost decreasing, it is \sim (1+2c)n 2 .(O(n 2))

Average case: O(n^2) since most of the case require sorting.

Therefore, the time complexity of Selection sort is $O(n^2)$.

Analysis of **Bubble Sort**:

i. **Intro:** Swap whenever it's in the wrong order.

ii. Proof: It swaps whenever two neighbors are not in the right order.From ith iteration in 1 to n-1 iterations, it does the aforementioned behavior from n to i+1. This way, no single thing is in the wrong order.

iii. Space complexity:

Fixed part: int: (i, j, n), *char: (tmp)

Variable part: *char (A) x n

Total: n (*char) + 3

iv. Time complexity:

| | 1. Algorithm BubbleSort(A, n) | 0 | 0 | 0 |
|----|---|---|--------|---------|
| | 2. { | 0 | 0 | 0 |
| | 3. for i := 1 to n - 1 do { | n | 1 | n |
| | 4. for j := n to i + 1 step -1 do { | n | n-1 | n^2-n |
| | 5. if $(A[j] < A[j-1])$ { | 1 | n-1 | n-1 |
| | 6. $tmp = A[j]; A[j] = A[j-1]; A[j-1] = tmp;$ | 3 | c(n-1) | 3c(n-1) |
| | 7. } | 0 | 0 | 0 |
| | 8. } | 0 | 0 | 0 |
| | 9. } | 0 | 0 | 0 |
| | 10. } | 0 | 0 | 0 |
| То | Total | | | n^2+ |

Best case: if c = 0 or 1, the 2^{nd} iteration does nothing, then it is (O(n)).

Worst case: if c \sim = n, it is almost decreasing, it is \sim n^2+...(O(n^2))

Average case: $O(n^2)$ since most of the case require sorting.

Therefore, the time complexity of Bubble sort is $O(n^2)$.

Analysis of **Shaker Sort**:

- Intro: Swap whenever it's in the wrong order but goes from both ends repeatedly.
- ii. **Proof:** Like Bubble sort, but it reduces the 1 to n cycle to half of it, and repeat the swapping from n to i+1 again, from i to n-1. Same as Bubble sort, no single thing is in the wrong order.

iii. Space complexity:

Fixed part: int: (j, l, r, n), *char: (tmp)

Variable part: *char (A) x n

Total: n (*char) + 4

iv. Time complexity:

| Statement | s/e | Freq. | Total |
|---|-----|-------|------------|
| Algorithm ShakerSort(A, n) | 0 | 0 | 0 |
| 2. { | 0 | 0 | 0 |
| 3. | 2 | 1 | 2 |
| 4. while l <= r do { | n+1 | 1 | n+1 |
| 5. for j := r to l + 1 step -1 do { | n | n | n^2 |
| 6. if $(A[j] < A[j-1])$ { | С | n | cn |
| 7. $tmp = A[j]; A[j] = A[j-1]; A[j-1] = tmp;$ | 3 | n | 3n |
| 8. } | 0 | n | 0 |
| 9. } | 0 | n | 0 |
| 10. | 1 | n | 1n |
| 11. for j := l to r - 1 do { | n | n | n^2 |
| 12. if (A[j] > A[j + 1]) { | С | n | cn |
| 13. $tmp = A[j]; A[j] = A[j+1]; A[j+1] = tmp;$ | 3 | n | 3n |
| 14. } | 0 | n | 0 |
| 15. } | 0 | n | 0 |
| 16. r := r − 1; | 1 | n | 1 n |
| 17. } | 0 | 0 | 0 |
| 18. } | 0 | 0 | 0 |
| Total | | | 2n^2 |

Best case: if c = 0 or 1, 2^{nd} iterations do nothing, then it is (O(n)).

Worst case: if $c \approx 1/2n$, it is almost decreasing, it is(O(n^2)).

Average case: O(n^2) since most of the case require sorting, but it goes 2

times slower than bubble sort.

Therefore, the time complexity of Shaker sort is $O(n^2)$.

Comparing Table:

| | Selection | Insertion | Bubble | Shaker |
|-------|---------------|---------------|---------------|---------------|
| Space | n (*char) + 4 | n (*char) + 3 | n (*char) + 3 | n (*char) + 4 |
| Time | O(n^2) | O(n^2) | O(n^2) | O(n^2) |

Prediction: Shaker > Bubble ? Insertion > Selection

The reason of Insertion > Selection is because the comparing step is lesser in Insertion sort. Shaker sort reduces the swapping steps by comparing from 2 ends.

The speed of the algorithms depends on whether writing data(swapping) or going through more steps is faster. Based on my learning in Computer Architecture, writing data is way slower.

Therefore, my prediction is: Insertion > Selection > Shaker > Bubble.

b. Implementation: (hw01.c on NTHUEE workstation, gcc 4.1.2)

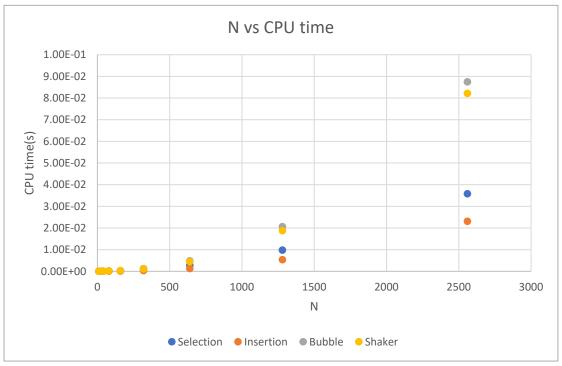
3. Result:

a. Result Table:

| N | Selection | Insertion | Bubble | Shaker |
|------|-----------|-----------|----------|----------|
| 10 | 9.90E-07 | 7.30E-07 | 1.11E-06 | 1.12E-06 |
| 20 | 6.79E-06 | 5.09E-06 | 1.00E-05 | 9.99E-06 |
| 40 | 2.33E-05 | 1.44E-05 | 3.96E-05 | 4.04E-05 |
| 80 | 8.60E-05 | 5.27E-05 | 1.64E-04 | 1.64E-04 |
| 160 | 1.33E-04 | 7.21E-05 | 2.77E-04 | 2.65E-04 |
| 320 | 1.14E-03 | 3.03E-04 | 1.13E-03 | 1.07E-03 |
| 640 | 2.76E-03 | 1.24E-03 | 4.81E-03 | 4.39E-03 |
| 1280 | 9.75E-03 | 5.32E-03 | 2.06E-02 | 1.87E-02 |
| 2560 | 3.58E-02 | 2.31E-02 | 8.74E-02 | 8.21E-02 |

b. Observation:

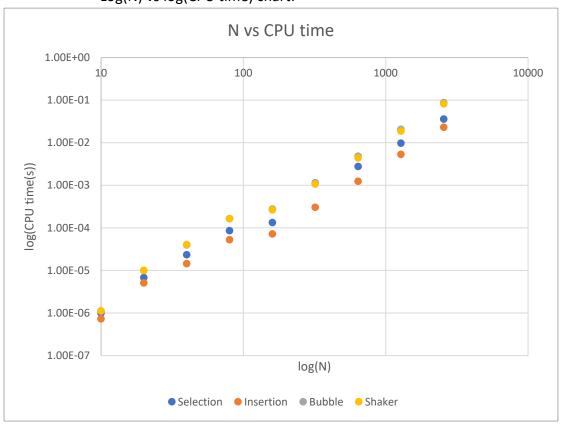
N vs CPU time chart:



Speed: Insertion > Selection > Shaker > Bubble.

It seems that the result meets by prediction. Furthermore, the time complexities of the four algorithms are $O(n^2)$, same as those that I've calculated.

Log(N) vs log(CPU time) chart:



When y-axis is logged, the graph looks linear. It indicates that they are $O(n^2)$.

Logged x-axis is just for good spacing between different Ns.

*The N=10 & N=40 case's bubble sort is faster than shaker sort. They might just be the special case when N is small.

c. Conclusion:

- i. Experimented Speed: Insertion > Selection > Shaker > Bubble.
- ii. Time Complexity: Insertion = Selection = Shaker = Bubble = O(n^2).
- iii. Best Case: Insertion > Shaker = Bubble > Selection.
- iv. Space Complexity: Insertion = Selection = Shaker = Bubble = theta(1)
- v. Swapping (data writing) is much slower than comparing.